

# INTRODUCTION TO MEMORY SYSTEM (II)

---

Lecturer: Hui Annie Guo

[h.guo@unsw.edu.au](mailto:h.guo@unsw.edu.au)

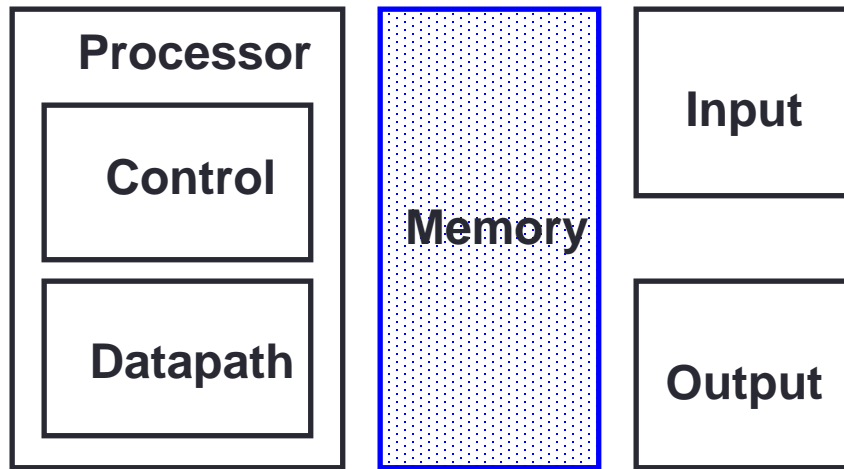
K17-501F

# Lecture overview

- **Topics**
  - **Memory hierarchy**
- **Suggested reading**
  - **H&P Chapter 5.2**

# The big picture:

- **Five classic components of a computer**



# Memory hierarchy

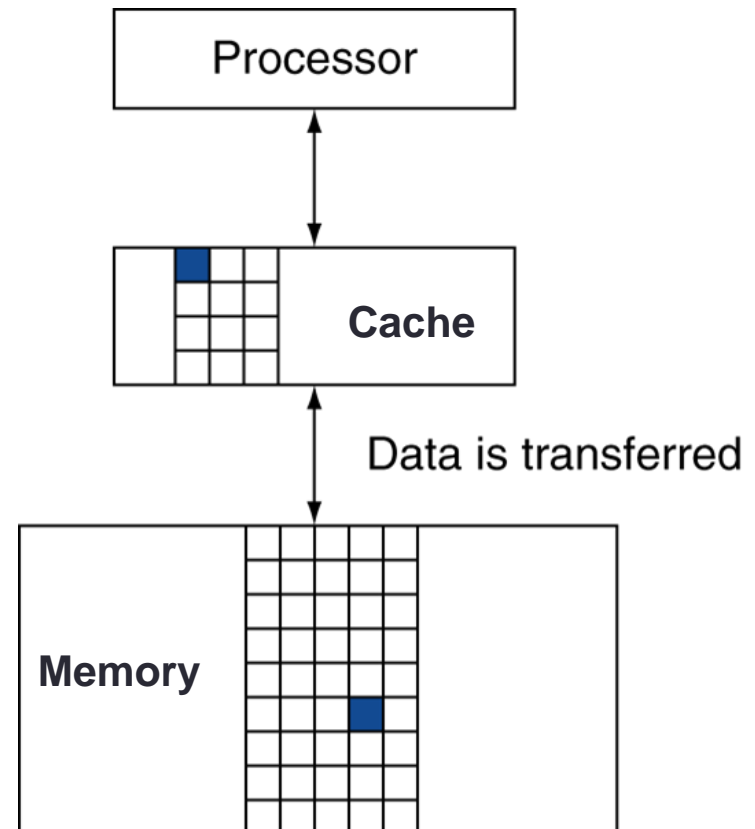
- **Memory hierarchy provides an illusion that a large, fast, cheap memory is available**
  - **How?**
  - **Why?**

# How?

- **Store everything on disk**
- **Copy recently required data from disk to smaller DRAM memory**
  - **main memory**
- **Copy more recently accessed (and nearby) items from DRAM to even smaller SRAM memory**
  - **Cache**
- **See example in the next slide**

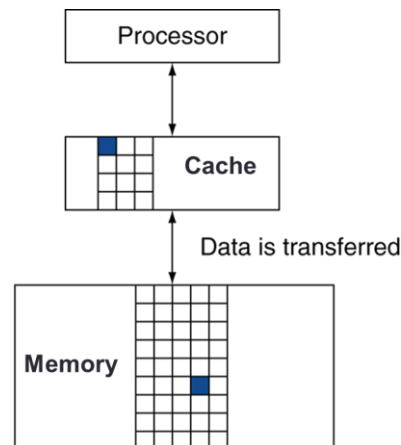
# Memory hierarchy (cont.)

- **A typical example – two-level memory hierarchy**
  - **Processor accesses cache for data**
  - **Cache is fast but small**
  - **Memory is larger but slow**



# Memory hierarchy (cont.)

- **Cache hit**
  - If accessed data is present in cache
- **Cache miss**
  - If accessed data is not in cache
  - **Solution: copy the data block from memory to the cache**
  - **The accessed data is then available in cache**



# Memory hierarchy (cont.)

- **Cache hit rate**
  - hits/accesses
- **Hit time**
  - Time taken to access cache
- **Cache miss rate**
  - misses/accesses  
=  $1 - \text{hit rate}$
- **Miss penalty**
  - Time taken to copy data block from memory to cache



# Why?

- **Principle of locality**
  - **Programs tend to access relatively small portions of the address space over a small period of time.**

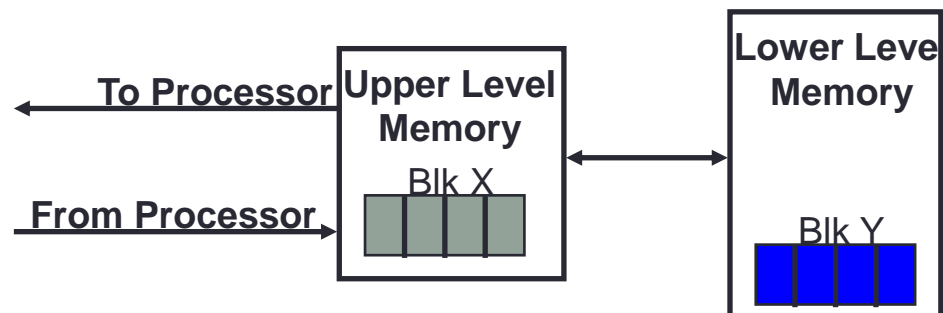


# Principle of locality

- **Two types of localities:**
  - **Temporal Locality** (locality in time)
    - If an item is referenced, it will tend to be referenced again soon.
  - **Spatial Locality** (locality in space)
    - If an item is referenced, items whose addresses close by tend to be referenced soon thereafter.

# Principle of locality (cont.)

- **Apply principle of locality on memory hierarchy, we can achieve “a large, fast, cheap memory”**
  - **Temporal locality**
    - **Keep most recently accessed data items closer to the processor**
  - **Spatial locality**
    - **Move multiple neighbourhood data items together (a **data block**) to cache**

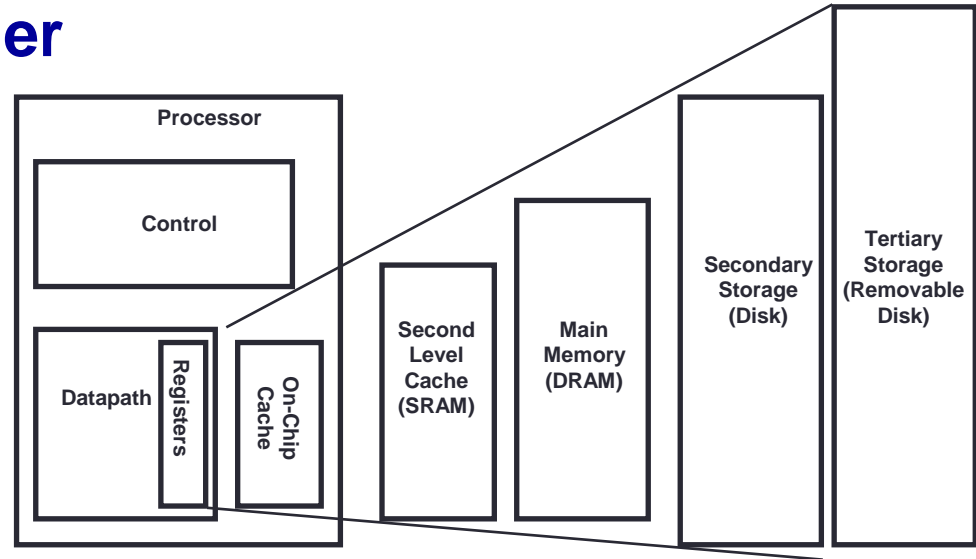


# Average access time

- **Average memory access time (AMAT)**
  - $\text{AMAT} = \text{Hit time} + \text{Miss rate} \times \text{Miss penalty}$
- **Example**
  - CPU with 0.2 ns clock, hit time = 1 cycle, miss penalty = 20 cycles, l-cache miss rate = 5%
  - What is AMAT of instruction memory?
    - $\text{AMAT} = (1 + 0.05 \times 20) \times 0.2 = 0.4\text{ns}$
    - 2 cycles per instruction

# Overview of memory system hierarchy

- **registers  $\leftrightarrow$  memory**
  - by compiler/programmer
- **cache  $\leftrightarrow$  memory**
  - by the hardware
- **memory  $\leftrightarrow$  disks**
  - by hardware and operating system (virtual memory)
  - by programmer



**We want AMAT as small as possible!**

# About Tutorial Solutions

# About Group Project