

Further Error-Correction Techniques

Or "How to be wrong but not *that* wrong"

Featuring:

- Realm Review
- Decreasing wrongness for less LUT space (ECALE method)
- Further error correction techniques.
- Matlab Guide
- (Potentially) Decreasing wrongness using (QSAM method)

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Absolute error

REALM

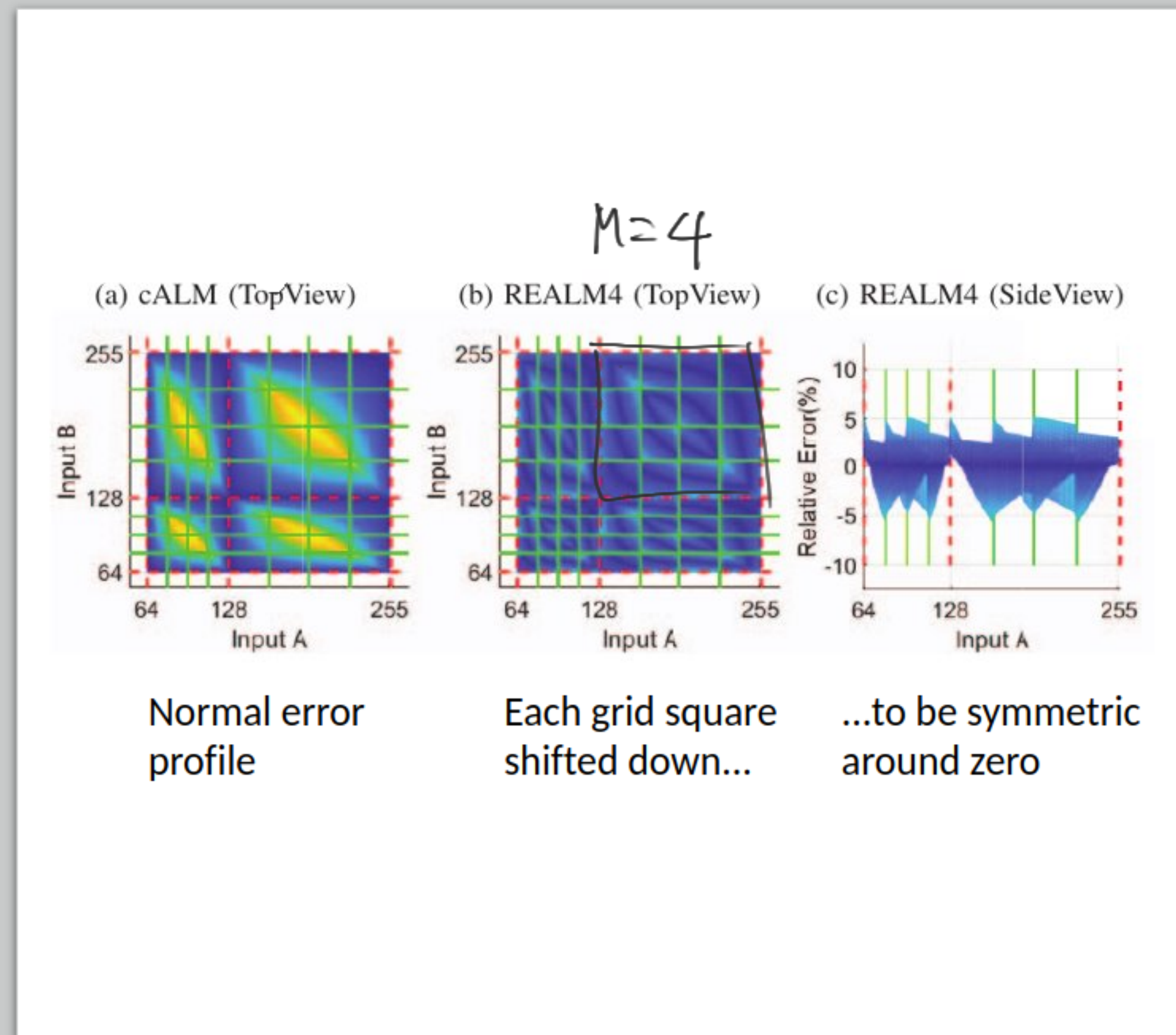
Low relative error & low area
(1% relative error with $M = 16$)

Short crit-path

Arbitrarily configurable Error at the cost of space

Error $\propto 1/M$

Space $\propto M^2$



But at what cost?

For $q = 8191$:

Maximum product is $8191 * 8191 = 67\,092\,481$

=> Need relative error of 0.003%

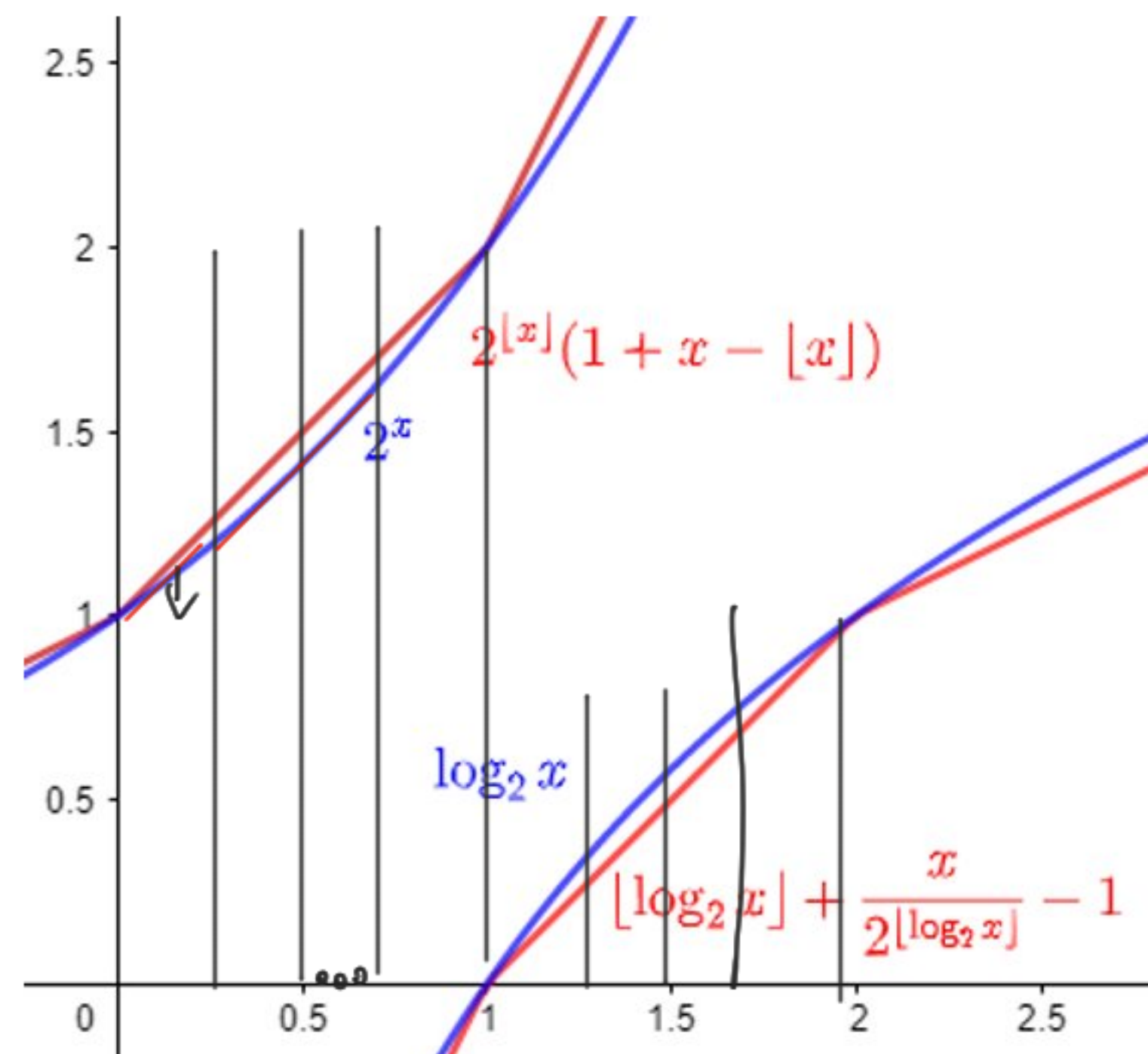
$M=9600$

=> $\sim 9600^2 = 92\,160\,000$ correction terms.

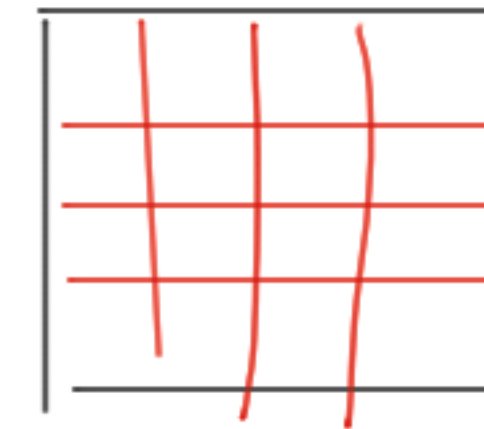
Can't fit in all block ram & distributed ram & LUT's combined

(Derivation Later for those interested)

Looking elsewhere...



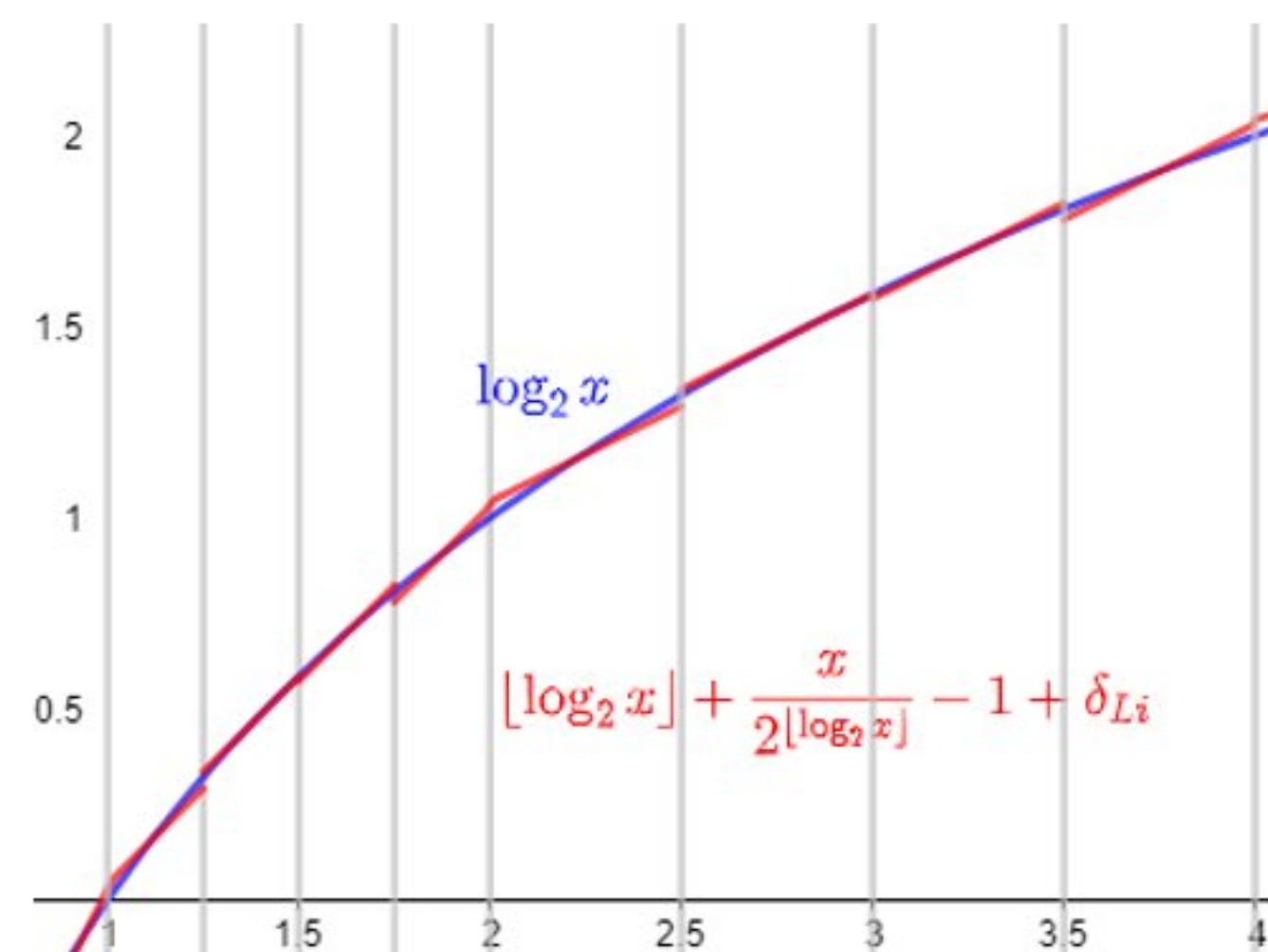
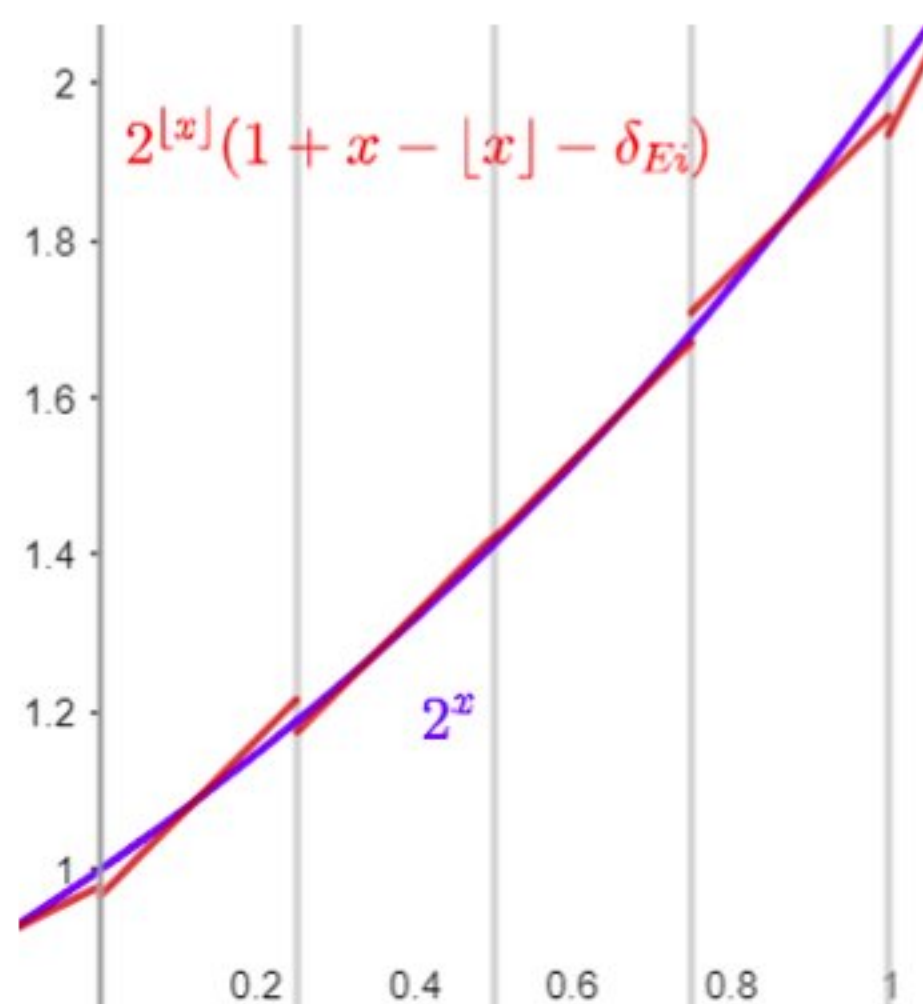
Bel



Approximate
Error corrected Log
and exponential
(ECALE)

$$M_E = 4$$

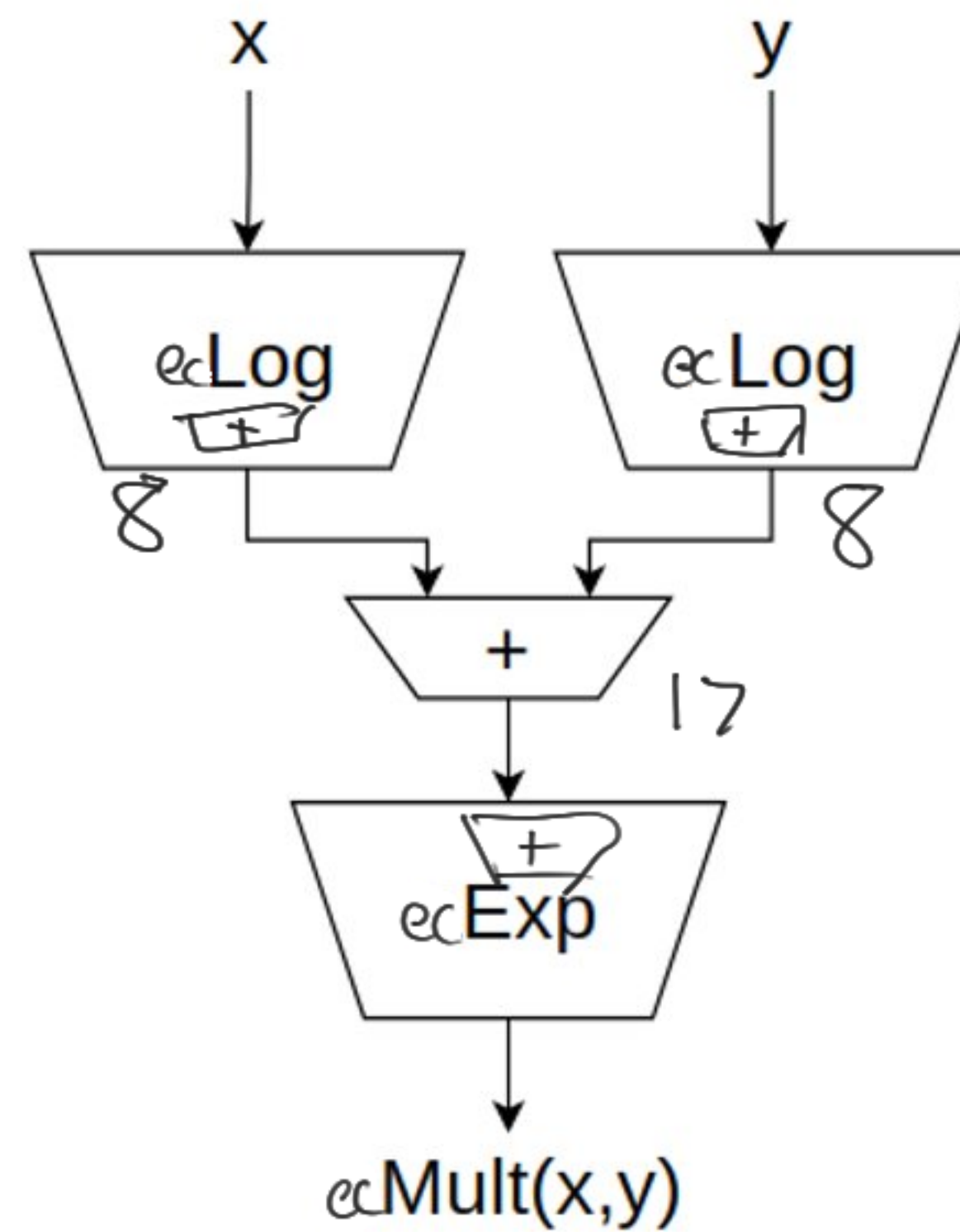
$$M_L = 4$$



Multiplication

As easy as $2^{\log(x) + \log(y)} = xy$

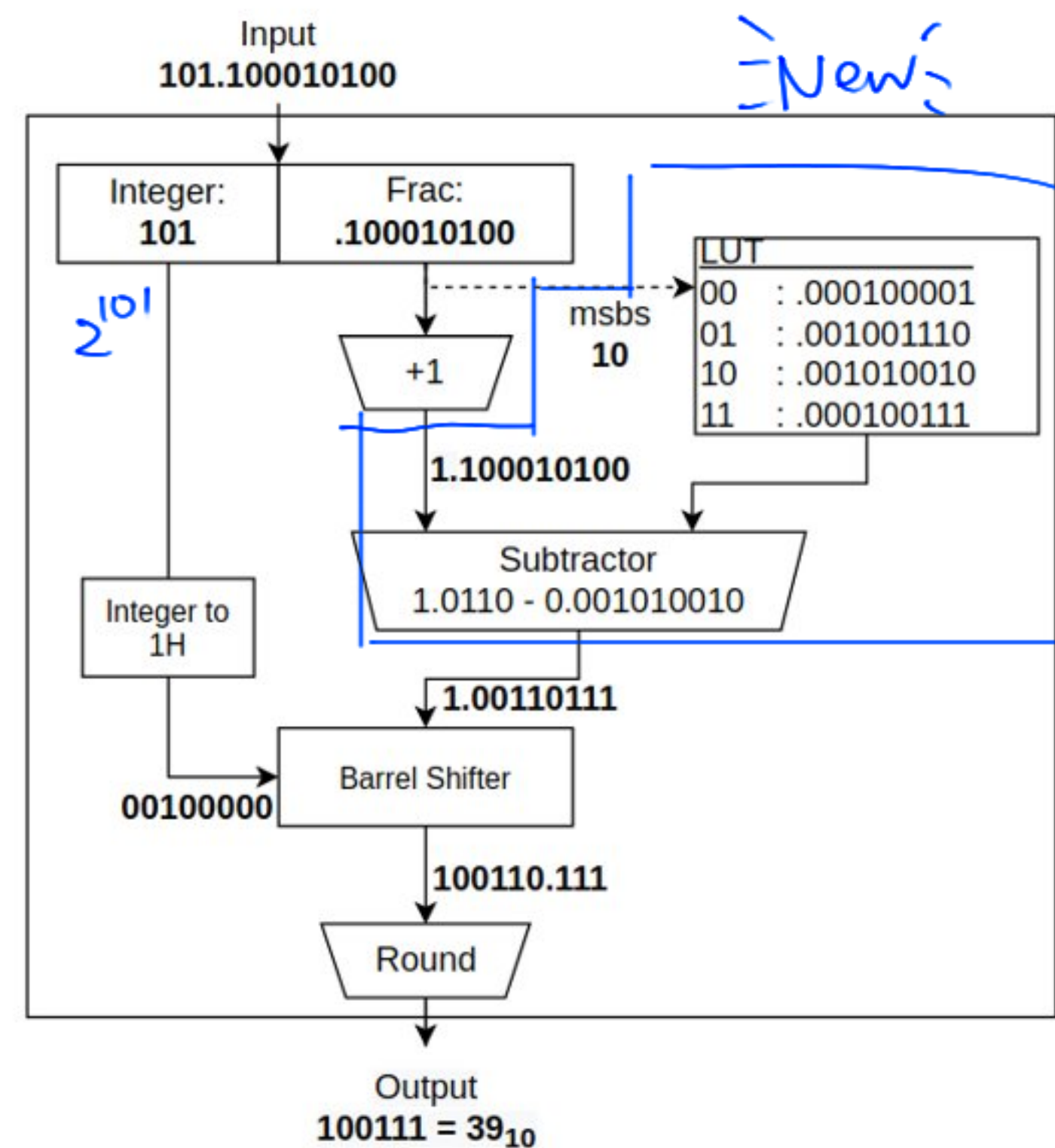
ecLog: "error corrected log"



ECALE 2^x (ecExp)

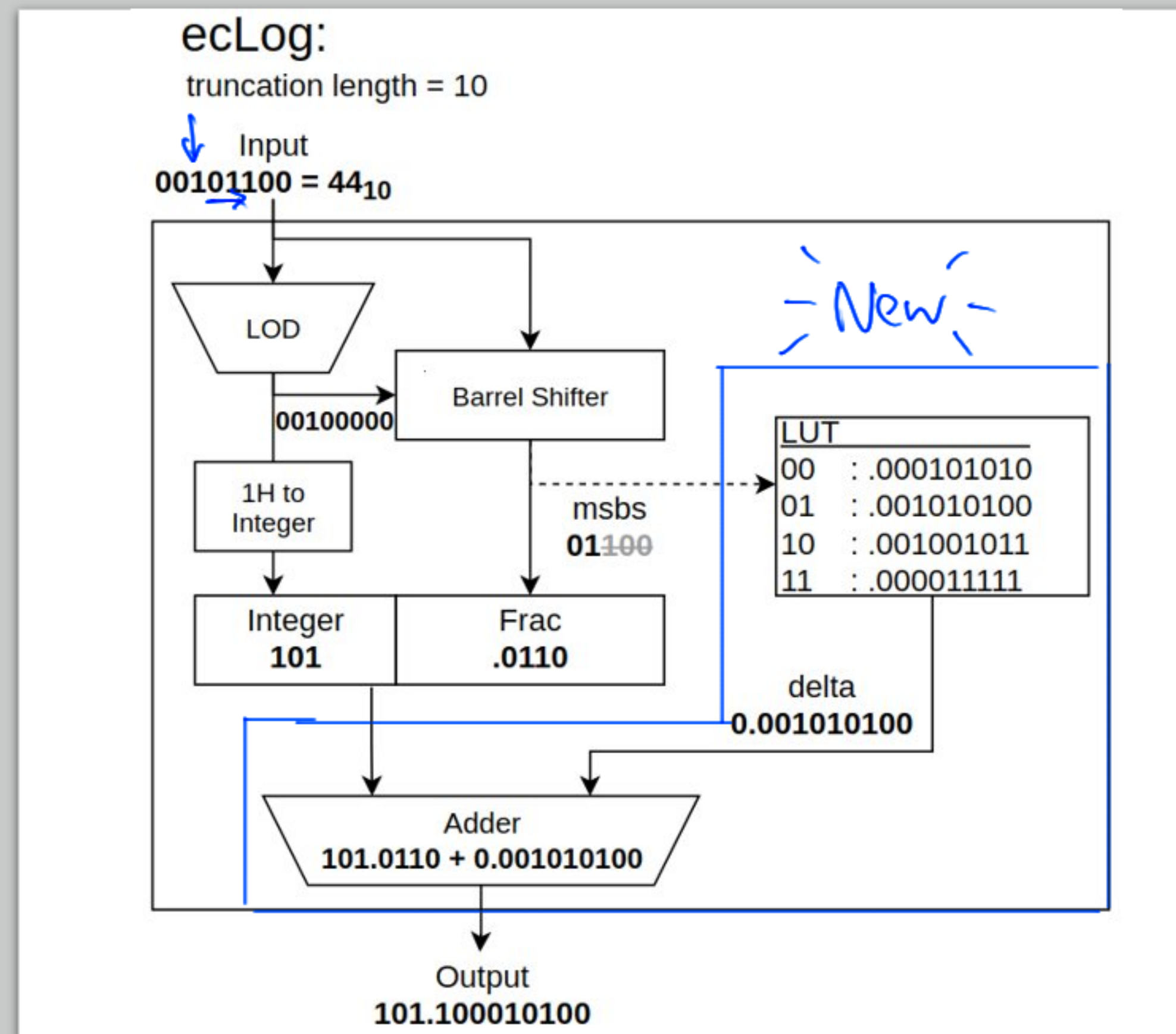
- Input is fixed point binary.
- Truncation length is
 - = Length of fractional part.
 - = Length of correction values
 - > Maximum integer part number. For max precision.
 - $> bitlen \times 2 + 1$

ecExp:
truncation length = 10



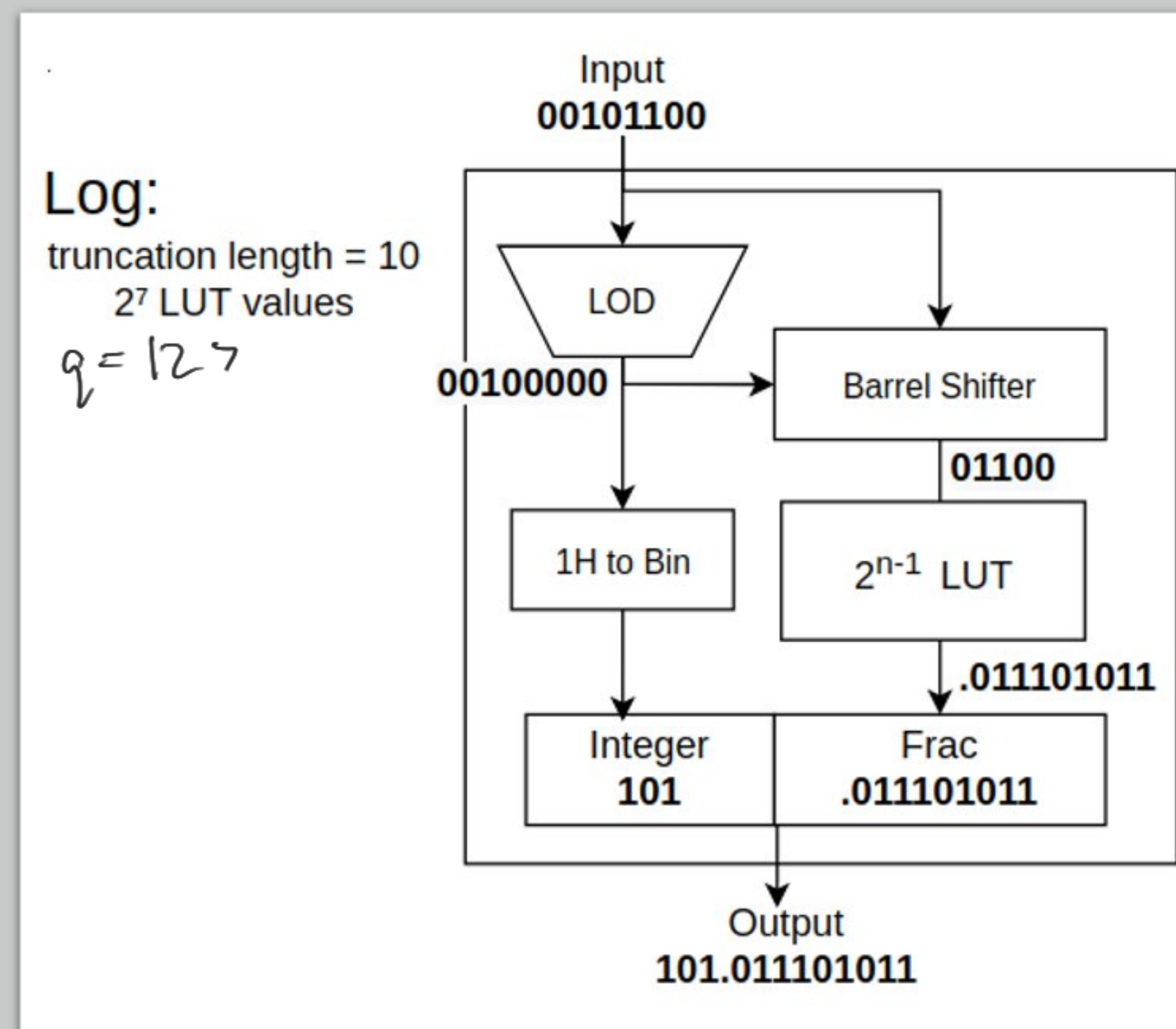
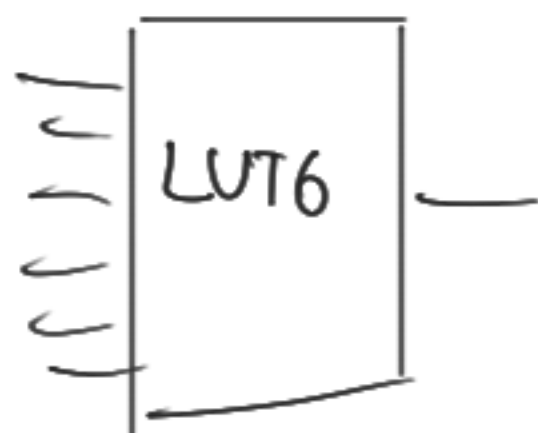
ECALE $\log_2 x$ (ecLog)

- Should have same truncation length as ecExp
- Can have a different number of LUT values to ecExp.



Precise LUT $\log_2 x$

- "Let's memorize every possible value of log"
- Use if need $M_L > q/2$ (Quite probable)
- LUT's on the Kintex-7 Are 6-bit \Rightarrow 64 possible values



ECALE Summary

Pros:

- Significantly decreased error vs space complexity

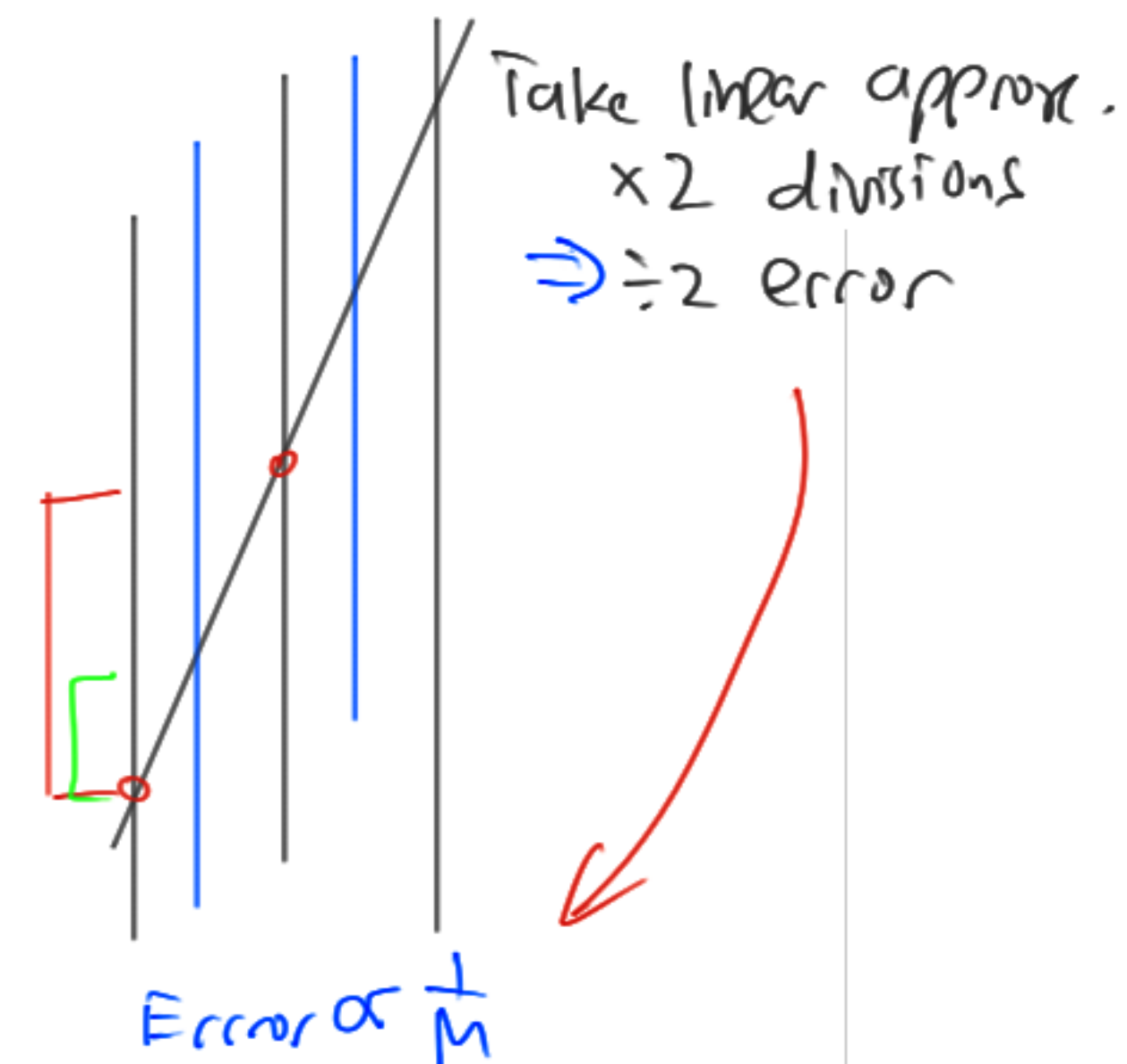
Cons:

- Triples number of LUT's needed compared to realm
- Increases critical path by one LUT, and one adder
- Possible to implement, but still not space efficient for LWE

Real
space $\propto M^2$

ECALE
space $\propto M_L + M_E$

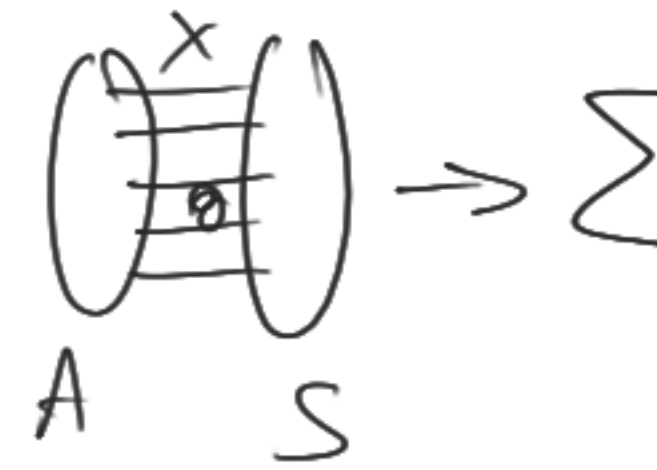
Excels when low range relative error rates are needed.



Question/Break time!

Further considerations:

$$A S = B$$



• LUT Sharing

- One look up table to rule them all?
- How can multiple resources share the same table?

• LUT Delay

- Block ram \rightarrow clocking
- Pipelining

$$2^{15}$$

• Truncation length

- $2 * \text{bitlen} + 1$ for max precision

Bit splitting

$$q+1 = 2^{13} \Rightarrow 2^{13} \equiv 1 \pmod{8191}$$

$$q = 2^{13} - 1$$

$$m=6, n=7$$

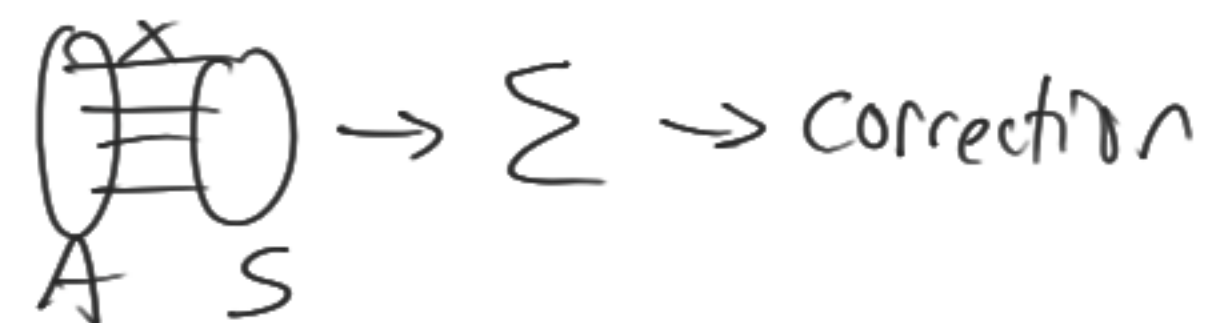
$$010101101001 \xrightarrow{12} 010101 \mid 101001$$

- $(a2^n + b)(c2^m + d) = ac2^{n+m} + ad2^n + bc2^m + bd$
- Precise logarithms achieved earlier
- Take advantage of modulus operation to reduce 2^{n+m} term
- Mixing in precise / approximate multiplication

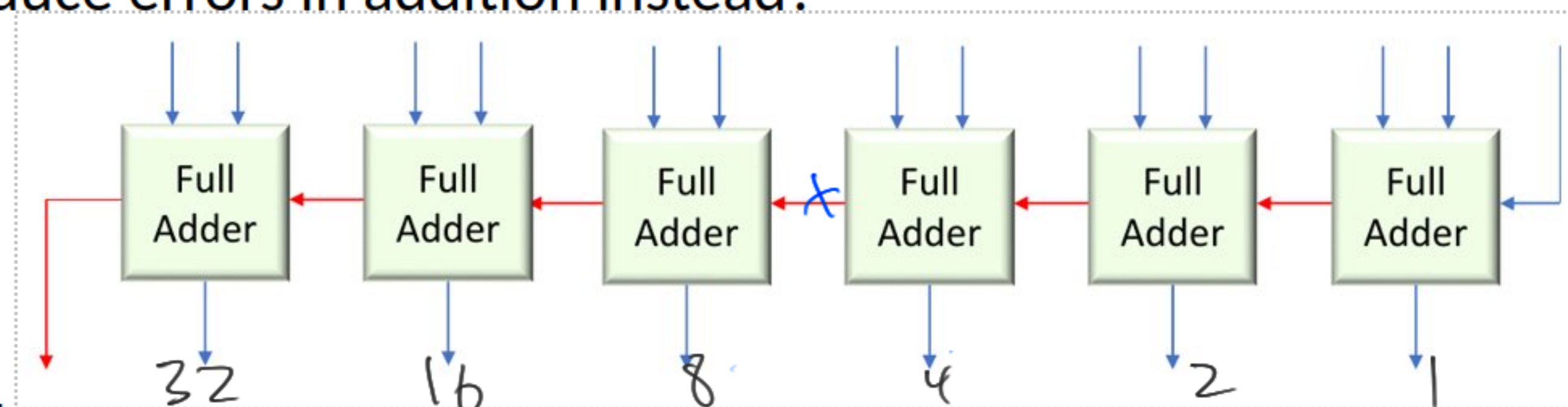
$$ad2^n + bc2^m + bd + ac$$

↑
Precise

Approximate addition



Multipliers on the Kintex-7 are fast $\rightarrow 6.5ns$
Introduce errors in addition instead? dsp units



To achieve full range of errors, want to break at different locations.

Other points

- Sacrifice Security
 - Decrease size of #samples or q to reduce error
- Do more research!
 - Tons more designs exist in the wild.

Q. Sam!

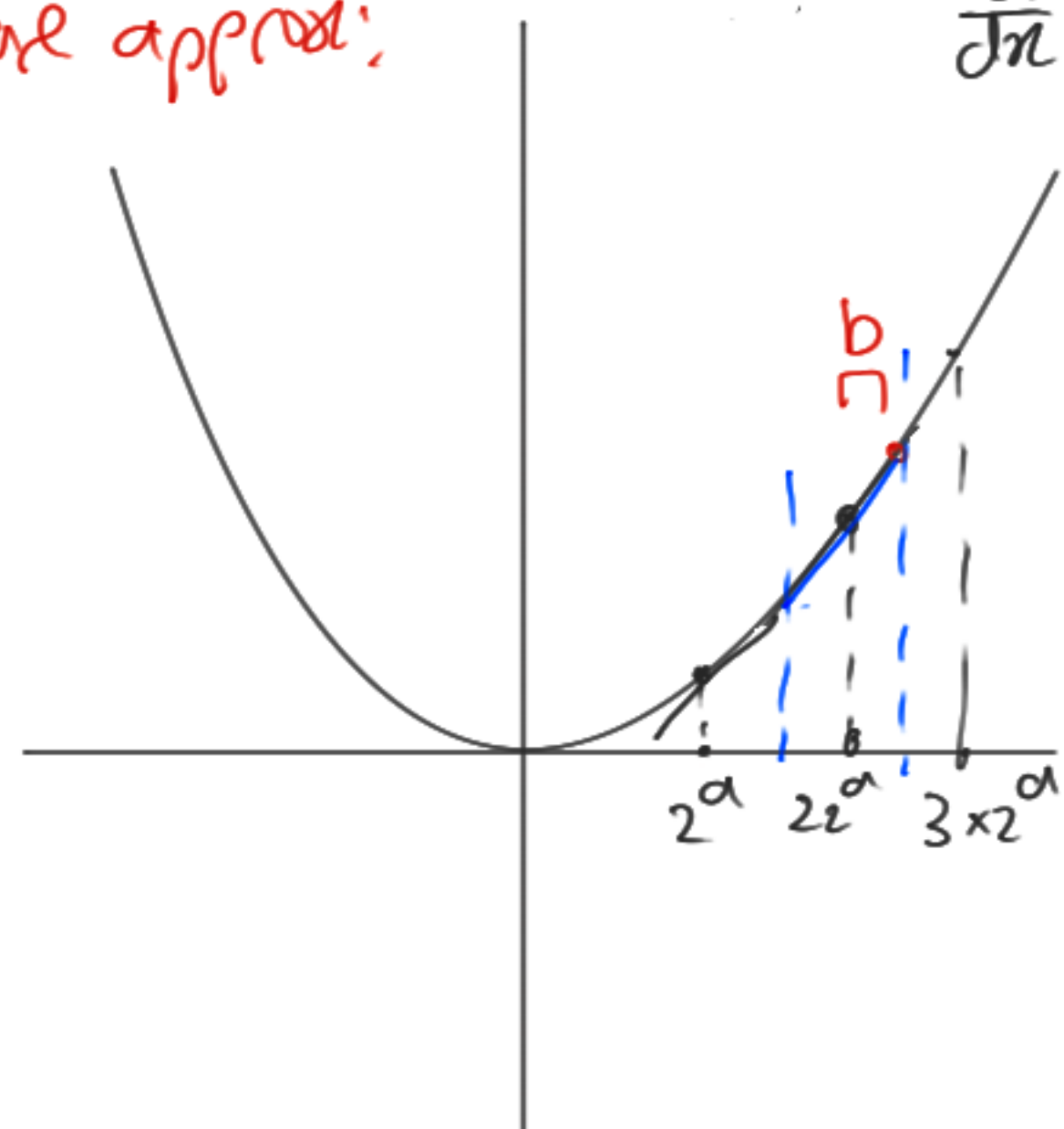
$$\frac{g(x+y) - g(x-y)}{4} = f(x, y)$$

$$\frac{(x+y)^2 - (x-y)^2}{4} = xy$$

More \rightarrow abs error is $\frac{2b_{max}^2}{4}$

\Rightarrow Want to approximate squares.

Square approx:



$\frac{dx^2}{dx} = 2x \rightarrow$ take linear approx.

$$(n \times 2^a \pm b)^2 = \underbrace{n^2 2^{2a} + b^2}_{\text{approx}} + \underbrace{2bn 2^a}_{\text{error}}$$

Annotations: LUT $\rightarrow n^2 2^{2a}$, Mult $\rightarrow b^2$, $b < 2^{a-1} \rightarrow b^2 < 2^{2a-2}$

Absolute error!! \therefore

$a=3 \rightarrow b < 4 \rightarrow$ easy
 $a=4 \rightarrow b < 8 \rightarrow$ okay
 $\rightarrow a=5 \rightarrow b < 16 \rightarrow$ Too hard! needs full multiplier.

More viable reduction is 2^4 . \therefore

Round x to nearest 2^a . eg!
 $\rightarrow x = 2^a \pm b$

$a=3$
 $x = 4025$
 $x = 4024 + 1$

$y = 4021$
 $= 4024 - 3$

$z = 10110$
 $z = 11000 - (10 + 1)$

Derivations of REALM ballpark

For $q = 8191$:

Maximum product is $8191 \times 8191 = 67\,092\,481$

Want absolute error to be much less than $q/4$

\therefore Relative error $< (q/4)/q^2 = 0.003\%$.

Compare against Realm's 1.8% max error @ $M = 16$. Recall Error $\times M \propto 1$

\therefore Need $M \times 0.003\% > 1.8\% \times 16$

$\therefore M > 9600$

Recall realm #Terms = M^2

$\therefore 9600^2$ correction terms

Can't fit in all block ram & distributed ram & LUT's combined