

MODULE *ABProtocol2*

EXTENDS *Integers, Sequences*

CONSTANT *Data*, The set of all possible data objects.
 Bad The bad message (to model broken/lost messages).

Bad must be different from any of the legal messages.

ASSUME $Bad \notin (Data \times \{0, 1\}) \cup \{0, 1\}$

VARIABLES *AVar*, The last $\langle value, bit \rangle$ pair A decided to send.
 BVar, The last $\langle value, bit \rangle$ pair B received.
 AtoB, Sequence of DATA messages in transit from sender to receiver.
 BtoA Sequence of ACK messages in transit from receiver to sender.

$TypeOK \triangleq \wedge AVar \in Data \times \{0, 1\}$
 $\wedge BVar \in Data \times \{0, 1\}$
 $\wedge AtoB \in Seq(Data \times \{0, 1\} \cup \{Bad\})$
 $\wedge BtoA \in Seq(\{0, 1, Bad\})$

$vars \triangleq \langle AVar, BVar, AtoB, BtoA \rangle$ All variables.

$Init \triangleq \wedge AVar \in Data \times \{1\}$
 $\wedge BVar = AVar$
 $\wedge AtoB = \langle \rangle$
 $\wedge BtoA = \langle \rangle$

A sending a data message to B by putting the same message in the channel
 until an ACK is received.

$ASnd \triangleq \wedge AtoB' = Append(AtoB, AVar)$
 $\wedge UNCHANGED \langle AVar, BtoA, BVar \rangle$

B receiving a data message from A.

$BRcv \triangleq \wedge AtoB \neq \langle \rangle$ There is at least one message in the channel.
 $\wedge IF (Head(AtoB) \neq Bad) \wedge (Head(AtoB)[2] \neq BVar[2])$
 THEN $BVar' = Head(AtoB)$ Accept the message if ACK bit is the alternate bit.
 ELSE $BVar' = BVar$ Ignore the message and keep the same local state.
 $\wedge AtoB' = Tail(AtoB)$ Remove the received message from the channel.
 $\wedge UNCHANGED \langle AVar, BtoA \rangle$

B sending an ACK for the last data value received.

$BSnd \triangleq \wedge BtoA' = Append(BtoA, BVar[2])$
 $\wedge UNCHANGED \langle AVar, BVar, AtoB \rangle$

A receiving an ACK from B.

$ARcv \triangleq \wedge BtoA \neq \langle \rangle$ There is at least one message in the channel.
 $\wedge IF Head(BtoA) = AVar[2]$ Check the ACK bit.
 THEN $\exists d \in Data : AVar' = \langle d, 1 - AVar[2] \rangle$ Alternate bit and send another message.
 ELSE $AVar' = AVar$ Keep sending AVar if ACK bit doesn't match.

$$\begin{aligned} & \wedge BtoA' = Tail(BtoA) \quad \text{Remove received message from the channel.} \\ & \wedge \text{UNCHANGED } \langle AtoB, BVar \rangle \end{aligned}$$

$$\begin{aligned} & \text{Corrupt a message in one of the channels.} \\ \text{CorruptMsg} & \triangleq \wedge \vee \begin{aligned} & \text{Corrupt a data message.} \\ & \wedge \exists i \in 1 \dots Len(AtoB) : \\ & \quad AtoB' = [AtoB \text{ EXCEPT } ![i] = Bad] \\ & \wedge BtoA' = BtoA \\ & \vee \text{Corrupt an ACK message.} \\ & \wedge \exists i \in 1 \dots Len(BtoA) : \\ & \quad BtoA' = [BtoA \text{ EXCEPT } ![i] = Bad] \\ & \wedge AtoB' = AtoB \\ & \wedge \text{UNCHANGED } \langle AVar, BVar \rangle \end{aligned} \end{aligned}$$

$$\begin{aligned} \text{Next} & \triangleq \vee ASnd \\ & \vee BRcv \\ & \vee BSnd \\ & \vee ARcv \\ & \vee \text{CorruptMsg} \end{aligned}$$

$$\text{Spec} \triangleq \text{Init} \wedge \Box[\text{Next}]_{vars}$$

$$\text{ABS} \triangleq \text{INSTANCE } \text{ABSpec}$$

$$\text{THEOREM } \text{Live} \triangleq \text{Spec} \Rightarrow \text{ABS!Spec}$$

$$\begin{aligned} \text{FairSpec} & \triangleq \text{Spec} \wedge \text{SF}_{vars}(ARcv) \wedge \text{SF}_{vars}(BRcv) \\ & \wedge \text{WF}_{vars}(ASnd) \wedge \text{WF}_{vars}(BSnd) \end{aligned}$$

NOTE(philiX): This doesn't hold (see the last lectures for the fixes)

$$\text{THEOREM } \text{Live2} \triangleq \text{FairSpec} \Rightarrow \text{ABS!FairSpec}$$