

**CSCI 3104, Algorithms**  
**Problem Set 1 (50 points)****Due January 22, 2021**  
**Spring 2021, CU-Boulder**

*Advice 1:* For every problem in this class, you must justify your answer: show how you arrived at it and why it is correct. If there are assumptions you need to make along the way, state those clearly.

*Advice 2:* Verbal reasoning is typically insufficient for full credit. Instead, write a logical argument, in the style of a mathematical proof.

**Instructions for submitting your solution:**

- The solutions **should be typed** and we cannot accept hand-written solutions. [Here's a short intro to Latex.](#)
  - You should submit your work through [Gradescope](#) only.
  - The easiest way to access Gradescope is through our Canvas page. There is a Gradescope button in the left menu.
  - Gradescope will only accept **.pdf** files.
  - [It is vital that you match each problem part with your work.](#) Skip to 1:40 to just see the matching info.
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1. (a) Identify and describe the components of a loop invariant proof.  
(b) Identify and describe the components of a mathematical induction proof.

**Solution:**

- (a) .
- **Initialization.** Loop invariant holds true prior to the first iteration of the loop.
  - **Maintenance.** If the loop invariant is true before the  $i$ th iteration of the loop, it remains true before the iteration  $i + 1$ .
  - **Termination.** When the loop terminates, the invariant gives us a useful property that helps show the correctness of the algorithm.
- (b) .
- **Statement of inductive intent.**
  - **Base Case.** Shows that the proposition holds for  $k = 0$  (holds for a simple case).
  - **Inductive Hypotheses.** Assume that the proposition holds true for a particular value  $k$ .
  - **Induction Step.** Prove from hypotheses that if the proposition holds true for  $k$ , it is also true for  $k + 1$ .
  - **Conclusion.** State that based on base case and the inductive step, the proposition holds true for all values.

2. Identify the loop invariant for the following algorithms.

```
(a) function Sum(A)
    answer=0;
    n=length(A);
    for i=1 to n
        answer += A[i]
    end
    return answer
end
```

```
(b) function Reverse(A)
    n=length(A)
    i=ceiling(n/2)
    j=ceiling(n/2) + (n+1) mod 2
    while i>0 and j<=n
        tmp=A[i]
        A[i]=A[j]
        A[j]=tmp
        i=i-1
        j=j+1
    end
end
```

(c) Assume that **A** is sorted such that the largest value is at **A**[n]. Assume **A** contains the value **target**.

```
function Search(A,target) //returns the index of the value target
    left=1
    right=length(A)
    while left <=right
        m=floor((left+right)/2)
        if A[m] < target
            left=m+1
        else if A[m]>target
            right=m-1
        else
            return m
        end
    end
end
```

**Solution:**

- (a) At the start of each iteration of the for loop, the variable *answer* is equal to the sum of all elements in the subarray  $A[1..i-1]$ .
- (b) At the start of each iteration of the while loop, *i* is the  $\text{ceiling}(n/2)$  – the number of iterations.
- (c) At the start of each iteration of the while loop, the variable *target* is in the subarray  $A[\text{left}..\text{right}]$ .

3. Prove the correctness of the following algorithm. (*Hint: You need to prove the correctness of the inner loop before you can prove the correctness of the outer loop.*)

```

function Sort(A)
  n=length(A)
  for i=1 to n
    for j=2 to n
      if A[j]<A[j-1]
        swap(A[j],A[j-1]) //a function that swaps the elements in the array
      end
    end
  end
end

```

**Solution:**

Inner loop:

- **Loop Invariant:** At the start of each iteration of the inner for loop, the subarray  $A[j-1..j]$  is sorted.

Outer loop:

- **Loop Invariant:** At the start of each iteration of the outer for loop, the subarray  $A[1..i]$  is sorted.
- **Initialization:** At the start of the first iteration of the outer for loop, where  $i = 1$ , the subarray  $A[1..i]$  is  $A[1..1]$ , which consists of one single element and is therefore sorted.
- **Maintenance:** At the start of the  $i$ th iteration,
- **Termination:** At th

4. (a) Suppose you have a whole chocolate bar composed of  $n \geq 1$  individual pieces. Prove that the minimum number of breaks to divide the chocolate bar into  $n$  pieces is  $n - 1$ .
- (b) Show that for fibonacci numbers  $\sum_{i=1}^n f_i^2 = f_n f_{n+1}$   
Recall that the fibonacci numbers are defined as  
 $f_0 = 0, f_1 = 1$   
 $\forall n > 1, f_n = f_{n-1} + f_{n-2}$
- (c) For which nonnegative integers  $n$  is  $3n + 2 \leq 2^n$ ? Prove your answer.

**Solution:**