

# Network Data Analysis - Assignment 1

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# Contents

<b>1</b>	<b>Pre-processing data</b>	<b>4</b>
<b>2</b>	<b>Analyses of 6 networks</b>	<b>5</b>
2.1	Available networks . . . . .	5
2.1.1	Network 1 . . . . .	5
2.1.2	Network 2 . . . . .	6
2.1.3	Network 3 . . . . .	6
2.1.4	Network 4 . . . . .	6
2.1.5	Network 5 . . . . .	7
2.1.6	Network 7 . . . . .	7
<b>3</b>	<b>Metrics of 6 networks</b>	<b>8</b>
3.1	Number of nodes . . . . .	8
3.2	Number of edges . . . . .	9
3.3	Edge density . . . . .	9
3.4	Degree distribution . . . . .	9
3.5	Average clustering coefficient . . . . .	9
3.6	Number of nodes in strongly connected component (SCC) . . .	11
3.7	Number of nodes in weakly connected component (WCC) . . .	11
3.8	Average path length in SCC . . . . .	11
3.9	Diameter of SCC . . . . .	12
3.10	Community detection . . . . .	12
3.11	Centrality Measures . . . . .	12
3.11.1	Network 1 . . . . .	12
3.11.2	Network 2 . . . . .	13
3.11.3	Network 3 . . . . .	13
3.11.4	Network 4 . . . . .	14
3.11.5	Network 5 . . . . .	14
3.11.6	Network 7 . . . . .	15
<b>4</b>	<b>Insights</b>	<b>17</b>
<b>5</b>	<b>Information Networks</b>	<b>18</b>
5.1	Crawler Implementation . . . . .	18
5.2	Code Explanation . . . . .	24
5.2.1	Imports . . . . .	24

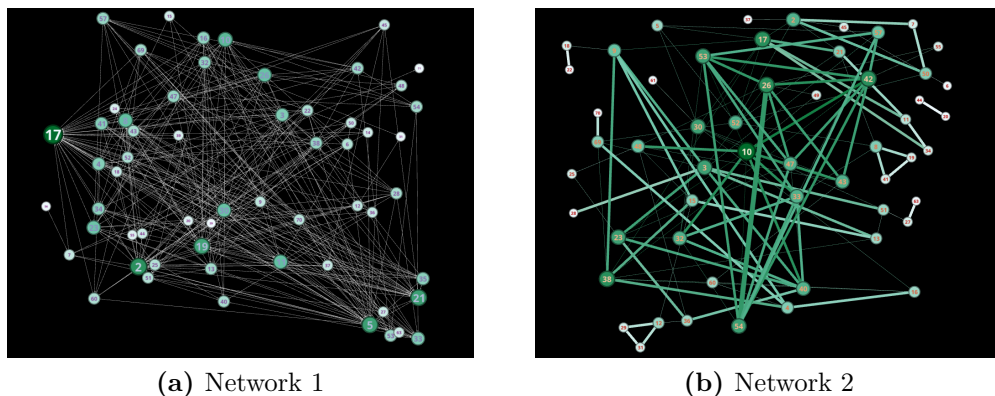
5.2.2	Global Variables . . . . .	24
5.2.3	Functions . . . . .	25
5.2.4	Main . . . . .	26

# 1 Pre-processing data

To complete the pre-processing of the original dataset, I used Pandas and Numpy. The first step for pre-processing the original file, after downloading it, was to remove all rows that had issues. Specifically speaking, this issue is not having either person name or ID. If either of these two information are missing, the row is removed from the original set. After removing those rows, I also removed those columns that did not have a label i.e., a question, and the two last questions that were further apart from the first 7. Furthermore, in order to have a more intelligible data set, I renamed the first column from "*Unnamed: 0*" to "*Name*". Then, I isolated all questions and for each one of them, I created a new data set comprised of three attributes: "*Name*", "*ID*", "*Question*". This resulted in 7 new data sets, one for each question, which were saved as CSV files. The questions are:

1. "Which person you have hear of their voice or seen their faces?"
2. "Which person you have met (in person+online) and exchange conversation?"
3. "Which person you have collaborated with?"
4. "Which person you have eye contact?"
5. "Which peson you have eaten lunch with?"
6. "Which person you have shared a ride?"
7. "Which person you have taken at least two courses with?"

After having separated the questions, I processed the inputs of each question in each individual file. The reason for this was to simply adequate each person's input to comma-separated IDs by removing spaces and trailing characters such as commas and hyphens. Finally, for each of the 7 questions, I created 7 other files in the CSV format Gephi expects. For example, a line with the following values "*1,2,3,4,5*" means that node 1 is connected to nodes 2, 3, 4 and 5.



**Figure 1:** Networks 1 & 2

## 2 Analyses of 6 networks

### 2.1 Available networks

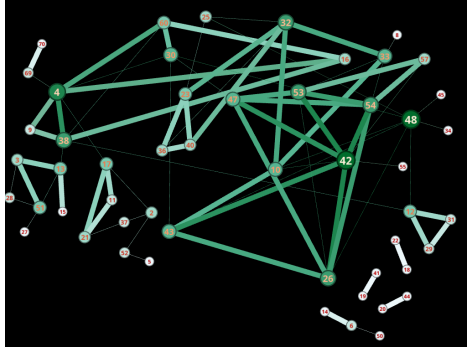
For this assignment, out of the 7 available networks, the ones I chose to work with are networks 1 through 5, and 7.

Before describing each network individually, one characteristic all of them share is that the edges in all 7 graphs are unweighted for the same reason. The questions asked *"which people?"*, meaning that for each person in the list, you only place an edge if the answer is yes. Because it is binary (edge or no edge), all graphs have unweighted edges. However, if we had a question asking *"how many times?"*, then there is a possibility of having weighted edges. Also, Gephi does not render a node if its degree equals zero.

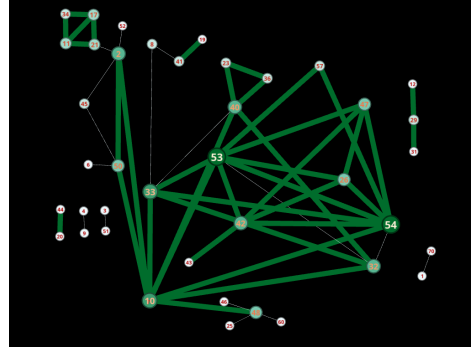
#### 2.1.1 Network 1

The first network asks the following question: *"Which person you have heard of their voice or seen their faces?"*.

**Directed or undirected?** This network is an example of a directed graph. The reason for this is that in online conferences, it is not all participants who speak and/or have their computer cameras on. Therefore, I may see other people's faces and/or hear their voices, but they may not see or hear me.



(a) Network 3



(b) Network 4

**Figure 2:** Networks 3 & 4

### 2.1.2 Network 2

The second network asks the following question: *"Which person you have met (in person+online) and exchange conversation?"*.

**Directed or undirected?** This network is an example of an undirected graph. In order to have a conversation, both parties must engage. If only one of them speak, there is no dialog. Therefore, it is not possible to exchange conversation.

### 2.1.3 Network 3

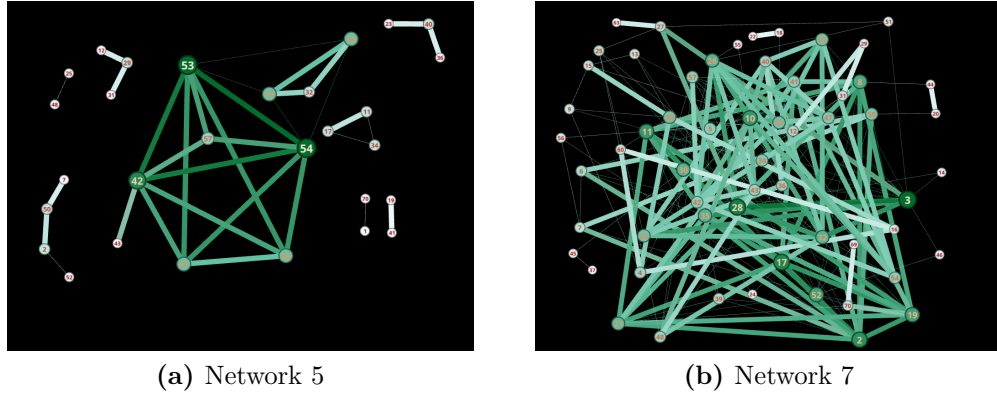
The third network asks the following question: *"Which person you have collaborated with?"*.

**Directed or undirected?** This is an undirected graph because collaboration must be enforced by both parties. It has to be a mutual agreement.

### 2.1.4 Network 4

The fourth network asks the following question: *"Which person you have eye contact?"*.

**Directed or undirected?** This network is undirected because in order to establish eye contact, two people must engage. There is no way for one person to look into someone else's eyes and not be looked back.



**Figure 3:** Networks 5 & 7

### 2.1.5 Network 5

The fifth network asks the following question: *"Which person you have eaten lunch with?"*.

**Directed or undirected?** This is an undirected graph because eating lunch with another person implies that both people had to simultaneously engage in the activity.

### 2.1.6 Network 7

The seventh network asks the following question: *"Which person you have taken at least two courses with?"*.

**Directed or undirected?** This network is an undirected graph because you cannot be simultaneously enrolled in a course another person has not and consider that as taking a course together. Also, the course must be taken in the same year.

<b>Metric</b>	<b>Net. 1</b>	<b>Net. 2</b>	<b>Net. 3</b>	<b>Net. 4</b>	<b>Net. 5</b>	<b>Net. 7</b>
Num. Nodes	60	60	59	60	59	60
Num. Edges	426	120	73	48	30	214
Density	0.120	0.067	0.042	0.027	0.017	0.120
Avg. Clust. Coef.	0.394	0.304	0.296	0.241	0.168	0.488
Num. Nodes SCC	49	-	-	-	-	-
Num. Nodes WCC	58	-	-	-	-	-
Num. Nodes CC	-	52	42	20	10	59
Avg. Path Len. SCC	2.599	-	-	-	-	-
Avg. Path Len. CC	-	3.143	4.234	2.857	1.711	2.654
Diameter of SCC	7	-	-	-	-	-
Diameter of CC	-	7	11	6	4	6

**Table 1:** Network statistics

### 3 Metrics of 6 networks

In this section I will present some statistics about the 6 chosen networks. The results will be presented in tables. All metrics and plots were calculated and generated using the Python package NetworkX.

#### 3.1 Number of nodes

The number of nodes of each graph is the amount of people who have participated in the survey. Whether they are connected to someone else, or not, they will be represented as nodes.



### 3.2 Number of edges

The number of edges of each graph will be the amount of connections that exist between each person for each question.

### 3.3 Edge density

Graph density tells us how connected nodes are between each other. If the density value is high, we say the graph is connected, if the density value is low, we say it is sparse. For undirected graphs, this metric can be calculated as

$$D_{undirected} = \frac{2|E|}{|N|(|N| - 1)} \quad (1)$$

and the density for directed graphs is defined as

$$D_{directed} = \frac{|E|}{|N|(|N| - 1)} \quad (2)$$

where  $E$  is the number of edges and  $V$  is the number of nodes in the graph.

### 3.4 Degree distribution

The degree distribution can superficially tell us what are the preferences for people when connecting to each other. Figure 4 plots the degree distributions for the 6 chosen networks.

### 3.5 Average clustering coefficient

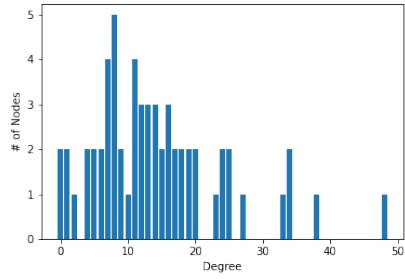
The average clustering coefficient for a graph helps determine how transitive a relationship is. For example, if persons A and B are friends, and persons A and C are friends, there is a high chance that persons B and C are also friends. The clustering coefficient is defined as

$$C_i = \frac{2e_i}{k_i(k_i - 1)} \quad (3)$$

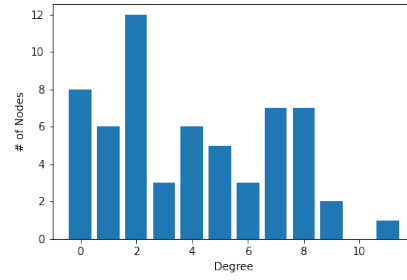
where  $e_i$  is the number of edges between the neighbors of node  $i$ .

The average clustering coefficient of the graph is calculated as

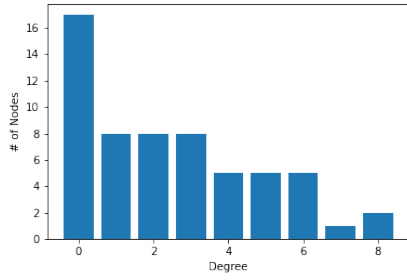
$$\langle C \rangle = \frac{1}{N} \sum_i^N C_i \quad (4)$$



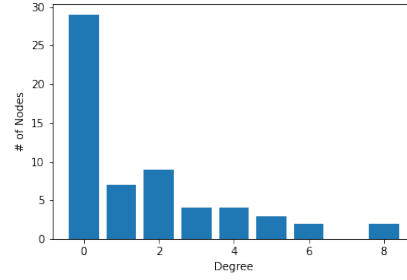
(a) Network 1



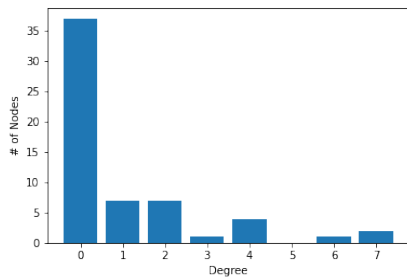
(b) Network 2



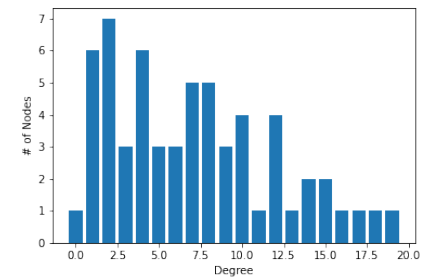
(c) Network 3



(d) Network 4



(e) Network 5



(f) Network 7

**Figure 4:** Degree distributions

where  $N$  is the number of nodes in the graph, and  $C_i$  is the clustering coefficient of node  $i$ .

### 3.6 Number of nodes in strongly connected component (SCC)

The strongly connected component (SCC) metric can only be obtained from directed graphs. Since only the first network is directed, it is the only one that can provide this value. For networks 2 through 5, and 7, the values are from the connected components. Refer to table 1 for the values.

### 3.7 Number of nodes in weakly connected component (WCC)

The weakly connected component (WCC) metric can only be obtained from directed graphs. Since only the first network is directed, it is the only one that can provide this value. For networks 2 through 5, and 7, the values are from the connected components. Refer to table 1 for the values.

### 3.8 Average path length in SCC

The average path length metric indicates how far apart two nodes are from each other in the connected graph. In other words, how many jumps, in average, it takes to reach other nodes. For a directed graph, it is calculated as

$$\langle d \rangle \equiv \frac{1}{2L_{max}} \sum_{i,j \neq i} d_{ij} \quad (5)$$

whereas for undirected graphs, it is calculated as

$$\langle d \rangle \equiv \frac{1}{L_{max}} \sum_{i,j > i} d_{ij} \quad (6)$$

For networks 2 through 5, and 7, I calculated the average path length in the connected component because these networks are undirected graphs, thus not being possible to determine strongly connected components.

### 3.9 Diameter of SCC

This metric represents the maximum shortest distance between two nodes in a connected graph. It is be represented as

$$diameter \equiv \max_{ij} d_{ij} \quad (7)$$

For the first network, I collected the diameter of the strongly connected component. However, since all other networks are undirected graphs, I collected the diameter of the largest connected component.

### 3.10 Community detection

To detect communities in each of the chosen networks, I ran the Girvan–Newman algorithm, implemented in NetworkX. Figure 5 illustrates the communities in each network. The communities the algorithm found make sense, given that it ran by removing the edges with highest betweenness, separating the communities the edge held together. Also, by comparing with figures 1, 2, and 3, we can see that the nodes with most connections between each other form a community the algorithm was able to find. Another interesting feature that the community detection algorithm allows us to perceive is that the more sparse the graph, the more communities we have. Figures 5c, 5d, and 5e depict this behavior.

### 3.11 Centrality Measures

Centrality tries to determine which node is the most central in a graph. The four centrality measures will be used to determine which node is the most important in each network. The measures are in-degree, out-degree, betweenness, and closeness. The choice for each centrality measure for each graph was arbitrary.

#### 3.11.1 Network 1

For this network, in-degree and out-degree will be used to determine the two most central nodes. Because it is a directed graph, we can use these measures. The in-degree centrality of a node is calculated based on how many edges arrive in it. Similarly, the out-degree centrality of a node is calculated based on how many edges leave it. Table 2 summarizes the two most people.

Rank	Person	Score
1	17	0.576
2	5	0.508

(a) In-degree

Rank	Person	Score
1	10	0.271
2	19	0.254

(b) Out-degree

**Table 2:** Network 1 centrality

Rank	Person	Score
1	10	0.186
2	42	0.152

(a) Degree

Rank	Person	Score
1	17	0.169
2	30	0.131

(b) Betweenness

**Table 3:** Network 2 centrality

### 3.11.2 Network 2

The centrality measures analyzed for this network were degree centrality and betweenness centrality. Degree centrality is similar to in-degree and out-degree centrality measures, except for the direction of the edges, which do not exist. The degree centrality measure for a node is defined by the amount of edges connected to it. Betweenness, on the other hand, is defined by the amount of shortest paths that go through a given node, and is calculated as follows

$$C_B(i) = \sum_{j < k} \frac{g_{jk}(i)}{g_{jk}} \quad (8)$$

where  $g_{jk}$  is the amount of shortest paths connecting nodes  $j$  and  $k$ , and  $g_{jk}(i)$  is the node currently being analyzed. Table 3 summarizes the values found.

### 3.11.3 Network 3

The chosen centrality measures for this network were degree and closeness. Degree centrality has been previously defined. However, the closeness centrality measure is defined by the average length of shortest paths between a

Rank	Person	Score
1	48	0.137
2	42	0.137

(a) Degree

Rank	Person	Score
1	48	0.245
2	30	0.243

(b) Closeness

**Table 4:** Network 3 centrality

Rank	Person	Score
1	10	0.039
2	33	0.032

(a) Betweenness

Rank	Person	Score
1	10	0.169
2	33	0.165

(b) Closeness

**Table 5:** Network 4 centrality

node and all other nodes in a graph. It can be calculated as

$$C_C(i) = \left[ \frac{1}{N-1} \sum_{j=1}^N d(i, j) \right]^{-1} \quad (9)$$

where  $N$  is the number of nodes in a graph, and  $d(i, j)$  is the distance between nodes  $i$  and  $j$ . Table 4 summarizes the centrality results for this network.

#### 3.11.4 Network 4

The centrality measures chosen to analyze from this network were betweenness and closeness. These two measures have been previously defined. Table 5 summarizes the values found.

#### 3.11.5 Network 5

The centrality measures analyzed for this network were degree and betweenness centralities. These measures already have been introduced. Table 6 summarizes the two most influential nodes of this network and their centrality score.

Rank	Person	Score
1	53	0.120
2	54	0.120

(a) Degree

Rank	Person	Score
1	42	0.005
2	53	0.004

(b) Betweenness

**Table 6:** Network 5 centrality

Rank	Person	Score
1	3	0.322
2	28	0.305

(a) Degree

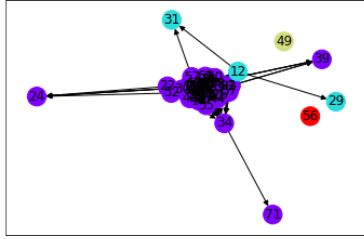
Rank	Person	Score
1	28	0.543
2	17	0.509

(b) Closeness

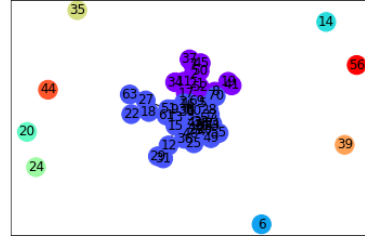
**Table 7:** Network 7 centrality

### 3.11.6 Network 7

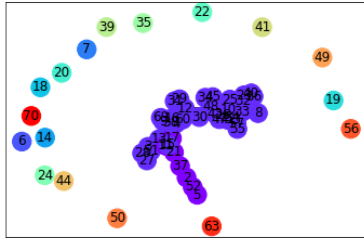
The chosen centrality measures for network 7 were degree and closeness. I will refrain from going deeper into these concepts since they have already been introduced. Table 7 summarizes the centrality findings for this network.



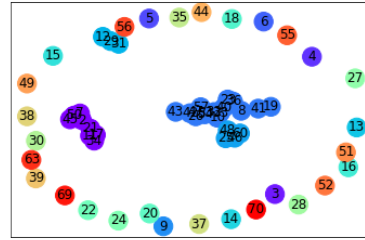
(a) Network 1



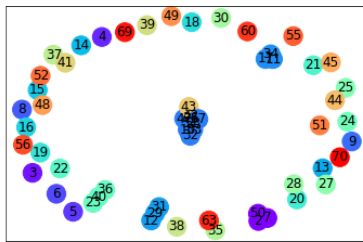
(b) Network 2



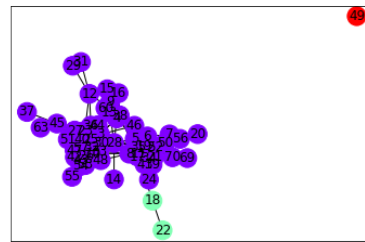
(c) Network 3



(d) Network 4



(e) Network 5



(f) Network 7

**Figure 5:** Detected Communities



## 4 Insights

All the metrics previously calculated and visualized allow us to extract insightful knowledge about the networks that exist in our graduate program. Tackling the aspects of low degree distribution, low density and high number of communities, we can see that the questions "*Which person you have collaborated with?*", "*Which person you have eye contact?*", and "*Which person you have eaten lunch with?*" create the most sparse graphs. That may be due to the fact that most classes are still online and people haven't gotten the chance to interact with each other in a more meaningful way.

It is also interesting to notice unexpected behaviors on these sparse networks. Let us look at network 3 for example. It was one of the three networks that were low in density and very sparse. However, the existing connected component was well connected. This means that despite existing many people who did not collaborate with their peers, the ones who did were thorough in doing it. However, because there were few edges in the whole graph, the average path length in the connected component was high. This is an indication that people collaborate together but they have their usual partners.

For network 5, which asks the question "*Which person you have eaten lunch with?*", we see that not many people have participated in this activity. Those who did, usually do frequently and there seems to be a closed group. Visualizing the network we are able to see that the connected component is almost a complete graph. This means that even though there are few connections in the graph, the ones that exist are meaningful.

Finally, discovering which nodes in each graph are more central is important. The reason for this is that knowing which people connect two or more components may contribute to how the people in these networks interact and collaborate with each other. Removing some of the key people from can result in disrupting an intricate collaboration network.

## 5 Information Networks

### 5.1 Crawler Implementation

```
1  ''' IMPORTS '''
2  # import HTML parser
3  from bs4 import BeautifulSoup
4
5  # import HTTP library
6  import requests
7
8  # import URL utility library
9  from urllib.parse import urlsplit
10
11 # import another URL utility library
12 import tldextract
13
14 # import sleep to prevent IP lockout
15 #     or intentional delays
16 from time import sleep
17
18 # import datetime to improve logging
19 from datetime import datetime
20
21 # python native queue data structure
22 from collections import deque
23
24 # native python json module, allows to serialize/
25 #     deserialize json
26 import json
27
28 '''GLOBAL VARIABLES'''
29 # list of accepted domains
30 university_domains = ['uoit', 'ontariotechu']
31
32 # file extensions to be ignored
33 bad_file_extensions = ["mp4", "mkv", "pdf", "docx", "doc", "
```

```

        mp3", "wav", "webp", "jpg", "png"]
34
35 # web page to begin traversing
36 INITIAL_URL = "https://ontariotechu.ca"
37
38 # this is a dictionary. each key is a URL and each
   value is a set of URLs that are referenced by the
   key
39 graph = {}
40
41 # this will be the queue of unvisited URLs
42 queue = deque()
43 '''END GLOBALS VARIABLES'''
44
45 ''' BEGIN FUNCTIONS '''
46 # helper function to get formatted current datetime
47 def get_now():
48     return datetime.now().strftime("%Y-%m-%d_%H:%M:%S")
49
50 # lambda function to disregard all those hrefs that are
   not http or https pages
51 def filter_non_http(url):
52     splitted_url = urlsplit(url)
53     scheme = splitted_url.scheme
54     if scheme not in ['http', 'https']:
55         return False
56     return True
57
58 # lambda function to filter off URLs with bad file
   extensions
59 def detect_bad_file_extensions(url):
60     splitted_url = urlsplit(url)
61     path = splitted_url.path
62
63     for ext in bad_file_extensions:
64         if ext in path:
65             return False
66     return True

```

```

67
68 # lambda function to
69 def rstrip_url(url):
70     splitted_url = urlsplit(url)
71     scheme = splitted_url.scheme
72     netloc = splitted_url.netloc
73     path = splitted_url.path\
74         .rstrip('/')\
75         .replace('index.php', '')\
76         .replace('index.html', '')\
77         .replace('//', '/')\
78         .rstrip('/')
79     new_url = f'{scheme}://{netloc}{path}'.rstrip('/')
80     return new_url
81
82 # lambda function to convert uoit domain to
ontariotechu
83 def uoit_to_ontariotechu(url):
84     splitted_url = urlsplit(url)
85     scheme = splitted_url.scheme
86     netloc = splitted_url.netloc.replace('uoit', '
    ontariotechu')
87     path = splitted_url.path
88     new_url = f'{scheme}://{netloc}{path}'.rstrip('/')
89     return new_url
90
91 # key function for BFS. This expands the current node (
URL)
92 def generate_children(url):
93     try:
94         # 1) sleep for 2 seconds to avoid ip blocking
or slow down
95         sleep(2)
96
97         # 2) perform HTTP request to given URL
98         print(get_now(), "GET_HTTP_request", url)
99         r = requests.get(url)
100        html_content = r.content

```

```

101
102     # 3) parse HTML response to Python object
103     soup = BeautifulSoup(html_content, "html.parser
        ")
104
105     # 4) get all anchor tags
106     all_anchors = soup.find_all("a")
107
108     # 5) get all hrefs from all anchors
109     all_hrefs = []
110     for anchor in all_anchors:
111         try:
112             all_hrefs.append(anchor['href'])
113         except:
114             pass
115
116     # 6) filter off hrefs that are not in the UOIT
        or OTU domains
117     university_hrefs = list(filter(lambda d:
        tldextract.extract(d).domain in
        university_domains, all_hrefs))
118
119     # 7) filter off all URLs that are not http or
        https (e.g. mailto)
120     university_hrefs = list(filter(filter_non_http,
        university_hrefs))
121
122     # 8) filter off URLs that point to files
123     university_hrefs = list(filter(
        detect_bad_file_extensions, university_hrefs
        ))
124
125     # 9) remove index.php or index.html, trailing
        slashes and double slashes
126     university_hrefs = list(map(rstrip_url,
        university_hrefs))
127
128     # 10) convert uoit domains to ontariotechu

```

```

129         since uoit redirects to ontariotech
        university_hrefs = list(map(
            uoit_to_ontariotechu, university_hrefs))
130
131         # 11) return all found children
132         print(get_now(), 'Returning_children...')
133         return list(set(university_hrefs))
134     except:
135         return []
136
137 # bfs driver function
138 def bfs(url):
139     print(get_now(), "Started_BFS...")
140     # 1) start by populating the queue with the initial
        node
141     queue.append(url)
142
143     # 2) while the queue isnt empty
144     while len(queue) != 0:
145         # 3) dequeue a node
146         u = queue.popleft()
147
148         # 4) get current nodes of the graph
149         graph_keys = list(graph.keys())
150
151         # 5) if the popped item is not a node in the
            graph
152         if u not in graph_keys:
153             print(get_now(), f'Generating_children_of_{
                u}...')
154             # 6) expand all nodes from current node
155             children = generate_children(u)
156             print(get_now(), 'Generated', len(children)
                , 'children!')
157
158             # 7) append to queue all chosen links in
                current web page
159             for child in children:

```

```

160         queue.append(child)
161
162         # 8) add current node's adjacency list
163         graph[u]=children
164         # 9) if web page has been visited before, skip
165         else:
166             print(get_now(), f'Skipped_{u}!')
167
168         # 10) after traversing all web pages, dump graph to
169         a json file for posterior processing and
170         analytics
171         print(get_now(), 'Writing_adjacency_list_to_file ...')
172         with open('bfs_adj_list.json', 'w') as f:
173             f.write(json.dumps(graph))
174         print(get_now(), 'Wrote_adjacency_list_to_file!')
175
176     ''' END FUNCTIONS '''
177
178 # Python's main function
179 ''' BEGIN MAIN '''
180 if __name__ == '__main__':
181     program_begin = get_now()
182     print(program_begin, "Starting_BFS...")
183     # start BFS from initial URL
184     bfs(INITIAL_URL)
185     program_end = get_now()
186     print(program_end, "Program_finished!")
187     delta = program_end - program_begin
188     print("Program_took", delta, "to_complete!")
189 ''' END MAIN '''

```

## 5.2 Code Explanation

### 5.2.1 Imports

BeautifulSoup is responsible for converting the HTML string retrieved from the website into a serialized Python object. This allows for easy querying of all HTML tags and their attributes.

The *requests* library is used to perform HTTP GET requests. The actual web page information is stored in the response's *content* property.

From *urllib* I imported the *urlsplit* function to easily split the URL into scheme, domain, and path. Query string parameters were ignored. In the same line of thought, *tlextract* was also used to easily extract the domain of the URL and change it in a lambda function from *uoit* to *ontariotechu* if it was the case.

The *sleep* function from module *time* was imported to force a 2-second wait before issuing the next HTTP request. This was done to avoid throttling, which would slow down the crawler's activity, and to prevent the domain from locking out my IP.

The *datetime* function from module *datetime* was imported to simply print out current date and time. The reason behind this was to improve logging quality.

The *deque* class was imported from module *collections* to ease the development of the BFS algorithm. It is a queue data structure and allows for queueing and dequeuing elements in it. It works as the queue of web pages that have not been visited yet.

Finally, the *json* module was imported to dump the assembled graph into a JSON file. Because the crawling was a time-consuming task, having the graph readily available is important.

### 5.2.2 Global Variables

A *university\_domains* list containing *uoit* and *ontariotechu* was used to hold the two University's known domains. It is employed on the filtering of URLs that point to outside of the University's network. This way I can keep crawl only the pages that belong and refer to the Ontario Tech University.

The *bad\_file\_extensions* list holds a few of the common file extensions that may be present on a URL. It, too, is used to filter off web pages that would host files. Since they are not actual web pages, it makes no sense to go through them.



Simply, the *INITIAL\_URL* is just a variable that is used to trigger the Breadth First Search Algorithm. It is passed in the main method as the starting node.

The *graph* variable is a dictionary that will be the assembled graph. The keys will be the URLs and the adjacency list will be a set of URLs the node has hyperlinks to.

Finally, the *queue* variable starts as an empty queue. This queue is used to hold the web pages that have yet to be visited.

### 5.2.3 Functions

The *get\_now* function is a wrapper for the *datetime* function. It's only purpose is to return a formatted string containing the current date and time, for logging purposes.

Next, the *filter\_non\_http* function is a lambda function. It is used in a *filter* list function to remove all URLs that are not HTTP or HTTPS schemes. An example of a URL that was removed was a *mailto*.

The *detect\_bad\_file\_extensions* function is another lambda filter function. It works by splitting the URL and grabbing its path - the path is everything that comes after the country (e.g. *.ca*) - and searching for any of the aforementioned bad file extensions.

Following the list of helper functions, the *rstrip\_url* is another lambda function. This one, on the other hand, is used in a list mapping. It works by splitting a URL and altering the path to remove trailing slashes, any files that are index files (HTML or PHP). This function then returns the stripped URL.

The next lambda function is called *uoit\_ontariotechu*. The sole purpose of this function is to map all URLs that are in the UOIT domain to Ontario Tech University's domain. This avoids adding nodes to the graph that are different in name, but that point to the same content.

Next, the core of the BFS algorithm, is the *generate\_children* function. It works by trying to perform an HTTP GET request. When the request fails by any reason, I say that no children were found, since no web page was collected. If the request does not fail, I collect the content (namely, the HTML page as a string), parse it with *BeautifulSoup*, query all anchor HTML tags, then try to collect every href attribute from the anchors. Next, I process all hrefs using a series of mappings and filters. I first filter off those URLs that are not in the UOIT/Ontario Tech University domain. Then the

remaining URLs are filtered to remove every scheme that is not HTTP not HTTPS. The remaining URLs are again filtered to remove those that host files. Next, I map the URLs to remove trailing slashes and index pages. Finally, the remaining set of URLs is mapped again, but this time to change all UOIT domains into Ontario Tech University domains. After this series of mappings and filters, the remaining URLs are put into a set, which in Python removes duplicates from any list, and returned. This finishes the execution of the *generate\_children* function.

Lastly, the *bfs* function is the actual implementation of the Breadth First Search algorithm. It starts by pushing to the queue the URL that was passed as a parameter. Then, until the queue is not empty, it dequeues a node that has not yet been visited and gets the current list of nodes in the graph. If the dequeued node is not in the list of nodes, expand the graph by generating the current node's children. Then, for each child of the current node, add it to the queue then add all children as the adjacency list of the current node in the graph. if the dequeued node is in the list of nodes in the graph, skip it to avoid performing unnecessary HTTP requests. When the queue is empty, the function dumps the assembled graph to a JSON file for persistent storage and avoid running the crawler every time the graph is needed.

#### 5.2.4 Main

The main function in Python is used to call the *bfs* function. It passes the *INITIAL\_URL* as parameter. This function is the starting point of a Python program. It also calculates how long it took for the crawler to finish traversing all web pages in the Ontario Tech University's domain.