

reminder

$$P_{ij}(m, m+n+r) = \sum_k P_{ik}(m, m+n) P_{kj}(m+n, m+n+r)$$

$$P(X_{n+n+r} = j | X_n = i)$$

re current $\begin{cases} \text{null} & = \mu_i = E_i[T_i] = \infty \\ \text{positive} & = \mu_i = \sum n f_{ii}(u) < \infty \end{cases}$

transient $\mu_i \rightarrow \infty$

- If j is transient $P_{ij}(n) \rightarrow 0$ as $n \rightarrow \infty \quad \forall i$
 $\downarrow \quad \downarrow \quad \downarrow$ null-recurrent $\quad \downarrow \quad \downarrow \quad \downarrow$
- $d(i) = |\{ n : P_{ii}(n) > 0 \}|$

if $d(i) = 1 \Rightarrow$ state i is aperiodic

Classifications of Chairs.

we say that i communicates with j , if the chain can visit j , starting from i , with strictly positive probability: $p_{ij}(n) > 0$ for some $n \geq 0$

And the notation is $i \rightarrow j$

If $i \rightarrow j$ and $j \rightarrow i \Rightarrow i \leftrightarrow j$

Theorem: if $i \leftrightarrow j$

① i & j have the same period

② i is transient iff j is transient

③ i is null-recurrent iff j is null-recurrent

Proof ①: $k \in \mathbb{S}^1$, $D_k = \{m \geq 1, P_{kk}(m) > 0\}$

$$\text{so } d(k) = \gcd(D_k)$$

$$\text{Let } i \leftrightarrow j \Rightarrow \exists m, n \text{ s.t. } \alpha := P_{ij}(m) P_{ji}(n) > 0$$

$$P_{ii}(m+r+n) \geq P_{ij}(m) P_{jj}(r) P_{ji}(n) \geq \alpha P_{jj}(r)$$

$$d(i) \mid m+r+n \quad \text{for } r \in \{0\} \cup D_j$$

$$\text{we know } d(i) \mid m+n \Rightarrow d(i) \mid r \quad \forall r \in D_j$$

$$\Rightarrow d(i) \mid d(j)$$

by the same argument $d(j) \mid d(i)$

$$\text{so, } d(i) = d(j)$$

proof(2): if $i \leftrightarrow j \Rightarrow P_{ii}(m+r+n) \geq \alpha P_{jj}(r) \quad \forall r \geq 0$

$$\sum_r P_{jj}(r) < \infty \quad \text{if} \quad \sum_r P_{ii}(r) < \infty$$

\rightsquigarrow since i is transient

A set C of states is called

(a) Closed if $p_{ij} = 0$ for all $j \notin C$ & $i \in C$

(b) irreducible if $i \leftrightarrow j \quad \forall i, j \in C$

Decomposition theorem: The state space S can be partitioned uniquely as

$$S = T \cup C_1 \cup C_2 \cup \dots$$

transient
irreducible
closed
recurrent

EX: $S = \{1, 2, \dots, 6\}$

$P =$

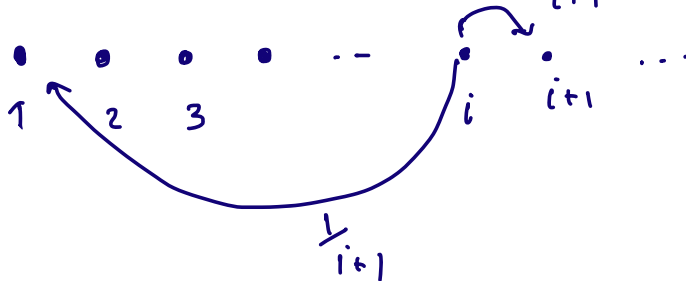
$\frac{1}{2}$	$\frac{1}{2}$	0	0	0	0
$\frac{1}{4}$	$\frac{3}{4}$	0	0	0	0
$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{4}$	0	0
$\frac{1}{4}$	0	$\frac{1}{4}$	$\frac{1}{4}$	0	$\frac{1}{4}$
0	0	0	0	$\frac{1}{2}$	$\frac{1}{2}$
0	0	0	0	$\frac{1}{2}$	$\frac{1}{2}$

irreducible & closed
{3, 4} \Rightarrow transient

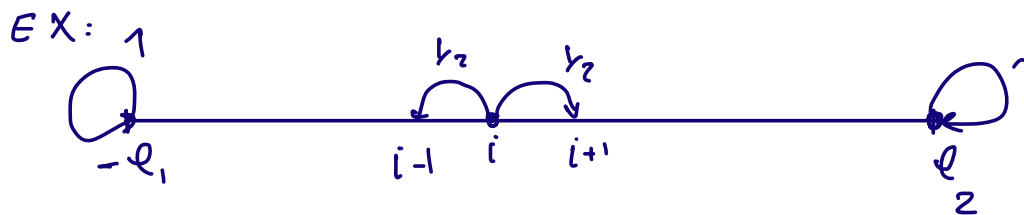
Lemma: If S is finite, one state is recurrent & all recurrent states are positive

$$1 = \lim_{n \rightarrow \infty} \sum_j P_{ij}(n) = \sum_j \lim_{n \rightarrow \infty} P_{ij}(n) = 0$$

EX: All states are null-recurrent $i, i+1$



EX: A symmetric random walk



All states $-l_1 < i < l_2$ are transient

$$P_{ij}(n) \rightarrow 0 \text{ as } n \rightarrow \infty$$

Let q : probability that the player wins l_2 dollars

$\Rightarrow 1-q$: 100% loses l_1

Let $w(n)$ be the gain after playing n rounds

$$0 = E[w(n)] = \sum_{i=-l_1}^{l_2} i \underbrace{P_{0,i}(n)}_{P_i(n)}$$

$$0 = \lim_n E[w(n)] = l_2 q - l_1 (1-q)$$

$$q = \frac{x_1}{l_1 + l_2}$$