Yale University

CPSC 516, Spring 2023

Assignment 5

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P.1.

+2 (a

From our work in class, we know that $\nabla^2 f \succeq mI$ implies that f is m-strongly convex. This yields the inequality

$$f(y) - [f(x) + \langle \nabla f(x), y - x \rangle] \ge \frac{m}{2} ||y - x||_2^2$$

by the definition of strong convexity.

On the other hand, we know by the second order Taylor expansion about x that

$$f(y) = f(x) + \langle \nabla f(x), y - x \rangle + \frac{1}{2} (y - x)^T \nabla^2 f(\xi)(y - x) \qquad \xi \in [x, y]$$

$$f(y) - f(x) + \langle \nabla f(x), x - y \rangle \le \frac{M}{2} \|y - x\|_2^2.$$

$$MI \succeq \nabla^2 f(\xi)(y - x)$$

This shows both inequalities.

+2 (b)

Suppose $f(z^*) = y^*$ and consider the inequality from P.1.(a)

$$y^* \le f(z) \le f(x) + \langle \nabla f(x), z - x \rangle + \frac{M}{2} ||z - x||_2^2.$$

Let us minimize the RHS with respect to z by taking the derivative and setting it to 0. We must have

$$\nabla f(x) + M[z - x] = 0$$
$$z = x - \frac{1}{M} \nabla f(x).$$

Substituting this particular value of z to the RHS above yields

$$f(x) + \left\langle \nabla f(x), -\frac{1}{M} \nabla f(x) \right\rangle + \frac{1}{2M} \|\nabla f(x)\|_{2}^{2}$$
$$= f(x) - \frac{1}{2M} \|\nabla f(x)\|_{2}^{2}$$

which is one of the desired inequalities.

To see the other inequality, We note that we wish to prove

$$\frac{1}{2m} \|\nabla f(x)\|_2^2 \ge [f(x) - f(z^*)]$$

which is known as the *Polyak-Lojasiewicz* (*PL*) condition.

By the definition of strong convexity,

$$\begin{split} f(x) - f(z^*) &\leq \langle \nabla f(x), x - z^* \rangle - \frac{m}{2} \|x - z^*\|_2^2 \\ &= \langle \nabla f(x), x - z^* \rangle - \frac{m}{2} \|x - z^*\|_2^2 - \frac{1}{2m} \|\nabla f(x)\|_2^2 + \frac{1}{2m} \|\nabla f(x)\|_2^2 \\ &= -\frac{1}{2} \left\| \sqrt{m} (x - z^*) - \frac{1}{\sqrt{m}} \nabla f(x) \right\|_2^2 + \frac{1}{2m} \|\nabla f(x)\|_2^2 \\ &\leq \frac{1}{2m} \|\nabla f(x)\|_2^2. \end{split}$$

Having shown both inequalities, we conclude the proof.

(c)

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In P.1.(b), we have shown that

$$f(z) \le f(x) - \frac{1}{2M} \|\nabla f(x)\|_2^2$$

for $z := x - \frac{1}{M} \nabla f(x)$. and all $\mathbf{q} \in \mathbb{R}^n$

But since we chose the step size α to minimize $f(x_{t+1})$, it must be at least as good as $\alpha = \frac{1}{M}$. Thus

$$f(x_{t+1}) \le f(x_t) - \frac{1}{2M} \|\nabla f(x_t)\|_2^2$$

as desired.

–) (d)

We argue by induction, the base case

$$f(x_0) - y^* \le 1 \cdot [f(x_0) - y^*]$$

holds trivially. Suppose it holds up to some t and consider $f(x_{t+1}) - y^*$.

We have

$$f(x_{t+1}) - y^* = f(x_{t+1}) - f(x_t) + f(x_t) - y^*$$

$$\leq f(x_t) - y^* - \frac{1}{2M} \|\nabla f(x_t)\|_2^2 \qquad \qquad \text{P.1.(c)}$$

$$\leq f(x_t) - y^* - \frac{m}{M} [f(x_t) - y^*] \qquad \qquad \text{P.1.(b) LHS}$$

$$= \left(1 - \frac{m}{M}\right) [f(x_t) - y^*]$$

$$= \left(1 - \frac{m}{M}\right)^{t+1} [f(x_0) - y^*]. \qquad \text{induction hypothesis}$$

By induction, we conclude the proof.

The number of iterations to reach ε error can be computed as follows

$$\left(1 - \frac{m}{M}\right)^{t} \left[f(x_0) - y^*\right] \le \exp(-mt/M) \left[f(x_0) - y^*\right]
\le \varepsilon
- \frac{mt}{M} + \log[f(x_0) - y^*] \le \log \varepsilon
t \ge \frac{M}{m} \log \frac{f(x_0) - y^*}{\varepsilon}.$$

+4 (e)

Suppose we are given A, b as input.

Consider the minimization problem

$$\min \|Ax - b\|_2^2$$
$$x \in \mathbb{R}^n$$

The objective is the composition of an affine function and a convex, separable, and non-decreasing (in each coordinate) function, which is therefore convex.

We explicitly compute its first and second derivatives

The objective is certainly twice differentiable, and since the eigenvalues of A^2 are just the eigenvalues of Asquared,

$$\lambda_1(A)^2 I \leq \nabla f^2 \leq \lambda_n(A)^2 I.$$

By our work above, if we start with an initial solution $x_0 := 0$ and run gradient descent with step size $\alpha = \frac{1}{\lambda_n(A)^2}$, this yields a solution x such that $||Ax - b||_2^2 \le \varepsilon$ after

$$T = O\left(\frac{\lambda_n(A)^2}{\lambda_1(A)^2} \log \frac{\|b\|_2^2}{\varepsilon}\right)$$

iterations. In each iteration, we need to compute the gradient and subtract it from the current iterate. The number of arithmetic operations is dominated by the gradient computation A(Ax), which requires $O(n^2)$ operations if we compute Ax and then A(Ax).

Thus the algorithm terminates after performing

 $O(n^2T) = O\left(n^2\kappa^2\log\frac{\|b\|_2^2}{\varepsilon}\right)$ Definding to K canbe Aron proved to linear by considery a duppoint convex longram.

arithmetic operations.

P.2.



Lemma 1:

 $(BB^T)^{\dagger}Bg$ is a minimizer of the optimization problem

$$\min_{y \in \mathbb{R}^n} \|B^T y - g\|_2^2.$$

Thus $B^T(BB^T)^{\dagger}Bg$ is the Euclidean projection of g onto the row space of B.

Proof: Lemma 1

The objective function is a composition of an affine (convex) function with a convex, separable, and coordinate-wise non-decreasing function, which is therefore convex. We can thus solve this problem by taking the gradient and setting it to zero.

The objective is equivalent to

$$(B^Ty - g)^T(B^Ty - g) = y^TBB^Ty - 2g^TB^Ty - g^Tg.$$

Taking the derivative and setting it to 0 yields

$$2y^TBB^T - 2g^TB^T = 0$$

$$y = (BB^T)^\dagger Bg.$$

Recall from elmentary linear algebra that the kernel is the orthogonal complement of the rowspace so that we can write

$$\mathbb{R}^m = \ker B \oplus \operatorname{row}(B).$$

In particular, if $r := \operatorname{rank} B$, we can find an orthonormal basis of \mathbb{R}^m , say $v_1, \ldots, v_r, w_1, \ldots, w_{m-r}$, where $v_i \in \operatorname{row}(B), w_j \in \ker B$. Thus we can write

$$g = \sum_{i=1}^{r} \langle g, v_i \rangle v_i + \sum_{j=1}^{m-r} \langle g, w_j \rangle w_j.$$

By elementary linear algebra,

$$x_g = \sum_{i=1}^r \langle g, v_i \rangle v_i$$
$$\Pi g = \sum_{j=1}^{m-r} \langle g, w_j \rangle w_j.$$

Thus

$$x_q + \Pi g = g$$

as desired.