

EE16A - Lecture 23 Notes

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Matching Pursuit Setup

1. Received signal $y = \sum_k \alpha_k S^{N_k} \vec{z}_k$ is a linear combination of delayed versions of the signals sent z_k where S^{N_k} is the circular shift matrix
 2. Sparse Representation: The linear combination (received signal) of the delayed messages contains very few terms relative to the max number of broadcasted signals.
- Scenario: A lot of players are trying to communicate with you, but you can only receive a few at the time.
 - Say you have $L = 2000$ players (trying to communicate), each sends a message $\vec{z}_k \in \mathbb{R}^{N=400}$
 - Usually $L > N$ and typically $L \gg N$
 - Received vector $y = \sum_k \alpha_k S^{N_k} \vec{z}_k$, which is a linear combination of delayed versions of the signals sent z_k where S^{N_k} is the circular shift matrix.
 - Sparse Representation: The linear combination (received signal) of the delayed messages contains very few terms relative to the 2000 possible terms.

Sparsity (Using Dictionary)

1. Use a larger, more specific dictionary to make sparser messages
- The larger the dictionary is, the more specific descriptions can be and the less words you need to use to convey a message.
 - Dictionary contains all the circularly shifted forms of the messages, so it will be of size $L \times N$ (no. of msgs \times no. of possible shifts = msg. vec length)
 - Called a **Redundant Dictionary**
 - $D = \{\vec{z}_1 S_{z_1} \dots \vec{z}_1 S_{z_1}^{N-1} | \vec{z}_2 S_{z_2} \dots \vec{z}_2 S_{z_2}^{N-1} | \dots \vec{z}_n S_{z_n}^{N-1}\}$
 - Assume: D is complete and includes N linearly independent vectors
 - Assume: Every vector is unit length, $\|\vec{z}_1\| = 1$
 - Cannot Assume: Orthogonality
 - Now let $D = \{\phi_1, \dots, \phi_T\}$, $T =$ Dictionary size
 - Need to figure out which subset of ϕ_k 's is represented in the received signal y .
 - Simplification: Assume no delay, so $y = \sum_k \alpha_k \vec{z}_k$ and $D = \{\phi_1 = \vec{z}_1, \dots, \phi_L = \vec{z}_L\}$
 - Residuals will approach 0.

Matching Pursuit Algorithm

1. Take observation vector y as the initial residual vector $r^{[0]}$
2. Project each residual $r^{[m-1]}$ onto the space of each of the vectors in the dictionary D and pick the maximum projection. (Choose a $\phi_k \in D$ s.t. $\vec{v}_m = \text{argmax}_k |\langle \vec{r}^{[m-1]}, \phi_k \rangle|$)
3. Write out equation for residual: $r^{[m-1]} = \langle \vec{r}^{[m-1]}, \vec{v}_m \rangle \vec{v}_m + \vec{r}^{[m]}$ where $r^{[m]} \perp \vec{v}_m$
4. Decompose the next residual ($r^{[m]}$)
5. Resulting signal at iteration M is $\vec{y} = \sum_{m=0}^{M-1} \langle r^{[m]}, \vec{v}_{m+1} \rangle \vec{v}_{m+1} + \vec{r}^{[M]}$

- A greedy algorithm: Step by step. At each step, search for a local optimum to reach a global optimum.
- Goal: Estimate $y \in \mathbb{R}^N$ by finding the main signal contributors
 - Let $M =$ no. of terms in the estimation (signal contributors)
 - Let $D = \{\phi_l\}, L = 1, \dots, L$ where $L \geq M$ and $L > N$ ($\phi_L \in \mathbb{R}^N$)
 - D is complete and $\|\phi_l\| = 1$
- At each step you have **Residue of** $y = \vec{r}^{[0]}$ (unaccounted portion)
- Step 0: $\vec{r}^{[0]} = y$
- Step 1: Choose a $\phi_k \in D$ s.t. $\vec{v}_1 = \text{argmax}_k |\langle y, \phi_k \rangle|$
 - Choose ϕ_k that gives the largest projection of y onto ϕ_k
 - Write $y = \langle \vec{y}_1, \vec{v}_1 \rangle \vec{v}_1 + \vec{r}^{[1]}$ where $r^{[1]} \perp \vec{v}_1$
 - $\|\vec{y}\|^2 = \|\langle \vec{y}_1, \vec{v}_1 \rangle \vec{v}_1\|^2 + \|\vec{r}^{[1]}\|^2 = |\langle \vec{y}_1, \vec{v}_1 \rangle|^2 \|\vec{v}_1\|^2 + \|\vec{r}^{[1]}\|^2$
 - $\|\vec{y}\|^2 = |\langle \vec{y}_1, \vec{v}_1 \rangle|^2 + \|\vec{r}^{[1]}\|^2$ (becuase \vec{v}_1 has unit length)
- Step 2: Decompose $\vec{r}^{[1]}$: Find the vector in D that best represents $\vec{r}^{[1]}$
 - Choose a $\phi_k \in D$ s.t. $\vec{v}_2 = \text{argmax}_k |\langle \vec{r}^{[1]}, \phi_k \rangle|$
 - $\vec{r}^{[1]} = \langle \vec{r}^{[1]}, \vec{v}_2 \rangle \vec{v}_2 + \vec{r}^{[2]}$ where $r^{[2]} \perp \vec{v}_2$
 - $r^{[2]}$ may not be (probably not) $\perp \vec{v}_1$
- Step m : Decompose $\vec{r}^{[m-1]}$
 - $\vec{r}^{[m-1]} = \langle \vec{r}^{[m-1]}, \vec{v}_m \rangle \vec{v}_m + \vec{r}^{[m]}$
 - To stop at iteration M , express $\vec{y} = \sum_{m=0}^{M-1} \langle r^{[m]}, \vec{v}_{m+1} \rangle \vec{v}_{m+1} + \vec{r}^{[M]}$
- Residuals approach 0
 - $\|\vec{y}\|^2 = \sum_{m=0}^{M-1} |\langle r_y^{[M]}, \vec{v}_{m+1} \rangle|^2 + \|r_y^{[M]}\|^2$
 - Fixed term on left
 - Sum of positive terms on right plus residual
 - As sum increases last residual must decrease for each step

Another View of the Matching Pursuit Algorithm

1. If A_m has Orthonormal Columns: $A_m^T A_m = I$
- Start with $y = r^{[m]}$ (observed measurement)
 - Create a matrix A_m at each step that has the columns of the vectors in the dictionary you have obtained up to that point. (Matrix of principal directions up to this step)
 - Start from $m = 1$: while $r^{[m]} \neq 0$, look for principal direction: $\vec{v}_m = \operatorname{argmax}_k |\langle r^{[m]}, \vec{z}_k \rangle|$
 - argmax : Scan all indices k , on k for which the inner product is the largest, use that inner product and assign $\vec{v}_m = \vec{z}_k$ as the m th principal direction
 - Augment matrix of principal directions: $A_m = [A_{m-1} | \vec{v}_m]$
 - Project y onto sol. space of A_m : $\vec{y} = A_m \alpha_m + \vec{\varepsilon}$
 - Approximate y in Least Squares sense
 - $\vec{y}_m = A_m \hat{\alpha}_m$ where $\hat{\alpha}_m$ is the LS soln to $A_m \alpha_m = \vec{y}$
 - $\hat{\alpha}_m = (A_m^T A_m)^{-1} A_m^T \vec{y}$
 - $\vec{y}_m = A_m (A_m^T A_m)^{-1} A_m^T \vec{y}$
 - $\vec{r}_m = \vec{y} - \vec{y}_m$
 - Problem: When columns of A_m , the principal directions \vec{v}_k are not orthogonal, the LS soln, is very costly.
 - Keep iterating forward: $m = m + 1$ until end.