

## SMART CONTRACT AUDIT REPORT

for

STACKER VENTURES

Prepared By: Shuxiao Wang

PeckShield March 11, 2021

### **Document Properties**

Client	Stacker Ventures
Title	Smart Contract Audit Report
Target	Stacker.VC
Version	1.0
Author	Xuxian Jiang
Auditors	Huaguo Shi, Xuxian Jiang
Reviewed by	Shuxiao Wang
Approved by	Xuxian Jiang
Classification	Public

#### **Version Info**

Version	Date	Author(s)	Description
1.0	March 11, 2021	Xuxian Jiang	Final Release
1.0-rc1	January 31, 2021	Xuxian Jiang	Release Candidate #1
0.2	January 23, 2021	Xuxian Jiang	Additional Findings
0.1	January 18, 2021	Xuxian Jiang	Initial Draft

#### Contact

For more information about this document and its contents, please contact PeckShield Inc.

Name	Shuxiao Wang
Phone	+86 173 6454 5338
Email	contact@peckshield.com

### Contents

1	Intr	oduction	4
	1.1	About Stacker Ventures	4
	1.2	About PeckShield	5
	1.3	Methodology	5
	1.4	Disclaimer	7
2	Find	lings	9
	2.1	Summary	9
	2.2	Key Findings	10
3 Detailed Results			11
	3.1	${\sf Safe-Version\ Replacement\ of\ safeTransfer}()/{\sf safeTransferFrom}()\ .\ .\ .\ .\ .\ .\ .\ .\ .\ .$	11
	3.2	Suggested Adherence of Checks-Effects-Interactions	13
	3.3	Incompatibility with Deflationary/Rebasing Tokens	14
	3.4	Non-Functional setEndBlock()/deposit() in LPGauge	16
	3.5	Asset Consistency in VaultGaugeBridge::newBridge()	17
	3.6	Improved Precision By Multiplication And Division Reordering	18
	3.7	Inconsistency Between Document and Implementation	20
4	Con	clusion	21
Re	eferer	nces	22

# 1 Introduction

Given the opportunity to review the design document and related smart contract source code of the Stacker.VC protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

#### 1.1 About Stacker Ventures

Stacker Ventures is a decentralized, community-owned venture capital protocol and accelerator that focuses on early-stage investment with aligned incentives of community investors with founding teams. By creating full functionality for a trust-minimized, decentralized VC investment fund, the protocol provides holders proportional claim over assets in the VC investment fund managed by a decentralized "fund council". These "fund council" members are community-elected to manage subsequent funds and consist of publicly available and auditable individuals. There are policies enforced in the protocol to delineate the responsibilities and limit the power of the fund council members.

The basic information of the Stacker.VC module is as follows:

Table 1.1: Basic Information of The Stacker.VC Module

ltem	Description
lssuer	Stacker Ventures
Website	https://stacker.vc/
Туре	Ethereum Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	March 11, 2021

In the following, we show the Git repository of reviewed files and the commit hash value used in this audit.

https://github.com/jack0x-tech/StackerVC\_VentureFund001/ (3cfb98b)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/jack0x-tech/StackerVC\_VentureFund001/ (87b1581)

#### 1.2 About PeckShield

PeckShield Inc. [13] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

High Critical High Medium

High Medium

Low

Medium Low

High Medium

Low

High Medium

Low

Likelihood

Table 1.2: Vulnerability Severity Classification

#### 1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [12]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [11], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

#### 1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.3: The Full List of Check Items

Category	Check Item		
	Constructor Mismatch		
	Ownership Takeover		
	Redundant Fallback Function		
	Overflows & Underflows		
	Reentrancy		
	Money-Giving Bug		
	Blackhole		
	Unauthorized Self-Destruct		
Basic Coding Bugs	Revert DoS		
Dasic Couling Dugs	Unchecked External Call		
	Gasless Send		
	Send Instead Of Transfer		
	Costly Loop		
	(Unsafe) Use Of Untrusted Libraries		
	(Unsafe) Use Of Predictable Variables		
	Transaction Ordering Dependence		
	Deprecated Uses		
Semantic Consistency Checks	Semantic Consistency Checks		
	Business Logics Review		
	Functionality Checks		
	Authentication Management		
	Access Control & Authorization		
	Oracle Security		
Advanced DeFi Scrutiny	Digital Asset Escrow		
ravancea Ber i Geraemi,	Kill-Switch Mechanism		
	Operation Trails & Event Generation		
	ERC20 Idiosyncrasies Handling		
	Frontend-Contract Integration		
	Deployment Consistency		
	Holistic Risk Management		
	Avoiding Use of Variadic Byte Array		
	Using Fixed Compiler Version		
Additional Recommendations	Making Visibility Level Explicit		
	Making Type Inference Explicit		
	Adhering To Function Declaration Strictly		
	Following Other Best Practices		

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary		
Configuration	Weaknesses in this category are typically introduced during		
	the configuration of the software.		
Data Processing Issues	Weaknesses in this category are typically found in functional-		
	ity that processes data.		
Numeric Errors	Weaknesses in this category are related to improper calcula-		
	tion or conversion of numbers.		
Security Features	Weaknesses in this category are concerned with topics like		
	authentication, access control, confidentiality, cryptography,		
	and privilege management. (Software security is not security		
	software.)		
Time and State	Weaknesses in this category are related to the improper man-		
	agement of time and state in an environment that supports		
	simultaneous or near-simultaneous computation by multiple		
	systems, processes, or threads.		
Error Conditions,	Weaknesses in this category include weaknesses that occur if		
Return Values,	a function does not generate the correct return/status code,		
Status Codes	or if the application does not handle all possible return/status		
	codes that could be generated by a function.		
Resource Management	Weaknesses in this category are related to improper manage-		
	ment of system resources.		
Behavioral Issues	Weaknesses in this category are related to unexpected behav-		
	iors from code that an application uses.		
Business Logics	Weaknesses in this category identify some of the underlying		
	problems that commonly allow attackers to manipulate the		
	business logic of an application. Errors in business logic can		
	be devastating to an entire application.		
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used		
	for initialization and breakdown.		
Arguments and Parameters	Weaknesses in this category are related to improper use of		
	arguments or parameters within function calls.		
Expression Issues	Weaknesses in this category are related to incorrectly written		
	expressions within code.		
Coding Practices	Weaknesses in this category are related to coding practices		
	that are deemed unsafe and increase the chances that an ex-		
	ploitable vulnerability will be present in the application. They		
	may not directly introduce a vulnerability, but indicate the		
	product has not been carefully developed or maintained.		

# 2 | Findings

#### 2.1 Summary

Here is a summary of our findings after analyzing the Stacker.VC implementation. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings		
Critical	0		
High	0		
Medium	2		
Low	4		
Informational	1		
Total	7		

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

#### 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 2 medium-severity vulnerabilities, 4 low-severity vulnerabilities, and and 1 informational recommendation.

ID	Severity	Title	Category	Status
PVE-001	Medium	Safe-Version Replacement of safeTrans-	Coding Practices	Fixed
		fer()/safeTransferFrom()		
PVE-002	Low	Suggested Adherence of Checks-Effects-	Time and State	Confirmed
		Interactions		
PVE-003	Low	Incompatibility with Deflationary/Rebasing	Business Logics	Confirmed
		Token		
PVE-004	Medium	Non-Functional setEndBlock()/deposit() in	Business Logics	Fixed
		LPGauge		
PVE-005	Low	Asset Consistency in VaultGauge-	Business Logics	Fixed
		Bridge::newBridge()		
PVE-006	Low	Improved Precision By Multiplication And Di-	Numeric Errors	Fixed
		vision Reordering		
PVE-007	Informational	Inconsistency Between Document and Imple-	Coding Practices	Fixed
		mentation		

Table 2.1: Key Stacker.VC Audit Findings

Besides recommending specific countermeasures to mitigate these issues, based on the fact that compiler upgrades might bring unexpected compatibility or inter-version consistencies, it is always suggested to use fixed compiler versions whenever possible. As an example, we highly encourage to explicitly indicate the Solidity compiler version, e.g., pragma solidity 0.6.11 instead of specifying a range, e.g., pragma solidity ^0.6.11.

In addition, we also emphasize that it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms need to kick in at the very moment when the contracts are being deployed in mainnet. Please refer to Section 3 for details.

# 3 Detailed Results

# 3.1 Safe-Version Replacement of safeTransfer()/safeTransferFrom()

• ID: PVE-001

Severity: MediumLikelihood: Medium

• Impact: Medium

• Target: GaugeD1

Category: Coding Practices [7]CWE subcategory: CWE-1126 [2]

#### Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the transfer() routine and possible idiosyncrasies from current widely-used token contracts.

In particular, we use the popular token, i.e., ZRX, as our example. We show the related code snippet below. On its entry of transfer(), there is a check, i.e., if (balances[msg.sender] >= \_value && balances[\_to] + \_value >= balances[\_to]). If the check fails, it returns false. However, the transaction still proceeds successfully without being reverted. This is not compliant with the ERC20 standard and may cause issues if not handled properly. Specifically, the ERC20 standard specifies the following: "Transfers \_ value amount of tokens to address \_ to, and MUST fire the Transfer event. The function SHOULD throw if the message caller's account balance does not have enough tokens to spend."

```
64
       function transfer (address to, uint value) returns (bool) {
            //Default assumes totalSupply can't be over max (2^256 - 1).
65
66
            if (balances[msg.sender] >= _value && balances[_to] + _value >= balances[_to]) {
67
                balances [msg.sender] -= value;
68
                balances [ to] += value;
                Transfer(msg.sender, _to, _value);
69
70
                return true;
71
           } else { return false; }
```

```
74
        function transferFrom(address _from, address _to, uint _value) returns (bool) {
75
            if (balances[ from] >= value && allowed[ from][msg.sender] >= value &&
                balances[_to] + _value >= balances[_to]) {
76
                balances [_to] += _value;
77
                balances [ from ] -= value;
78
                allowed [ from ] [msg.sender] -= value;
79
                Transfer ( from, to, value);
80
                return true;
81
            } else { return false; }
82
```

Listing 3.1: ZRX.sol

Because of that, a normal call to transfer() is suggested to use the safe version, i.e., safeTransfer (), In essence, it is a wrapper around ERC20 operations that may either throw on failure or return false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. To use this library you can add a using SafeERC20 for IERC20. Similarly, there is a safe version of transferFrom() as well, i.e., safeTransferFrom().

As an example, we show below the deposit() routine in the GaugeD1 contract. If the USDT token is supported as acceptToken, the unsafe version of IERC20(acceptToken).transfer(vcHolding, \_amountCommitHard) (line 122) may revert as there is no return value in the USDT token contract's transfer() implementation (but the IERC20 interface expects a return value)!

```
99
         function deposit (uint 256 _amount Commit Soft, uint 256 _amount Commit Hard, address
             creditTo) nonReentrant external {
100
             require(block.number <= endBlock, "GAUGE: distribution 1 over");</pre>
101
             require(fundOpen amountCommitHard == 0, "GAUGE: !fundOpen, only soft commit
                 allowed"); // when the fund closes, soft commits are still accepted
102
             require(msg.sender == _creditTo msg.sender == vaultGaugeBridge, "GAUGE: !bridge
                  for creditTo"); // only the bridge contract can use the "creditTo" to
                 credit !msg.sender
104
             claimSTACK( creditTo); // new deposit doesn't get tokens right away
106
             // transfer tokens from sender to account
107
             uint256 _acceptTokenAmount = _amountCommitSoft.add(_amountCommitHard);
108
             if (\_acceptTokenAmount > 0){
109
                 IERC20(acceptToken).safeTransferFrom(msg.sender, address(this),
                     acceptTokenAmount);
             }
110
112
             CommitState memory state = balances[ creditTo];
113
             // no need to update _state.tokensAccrued because that's already done in
                 _claimSTACK
114
             if ( amountCommitSoft > 0){
115
                 \_state.balanceCommitSoft = \_state.balanceCommitSoft.add(\_amountCommitSoft);
                 depositedCommitSoft = depositedCommitSoft.add( amountCommitSoft);
```

```
117
118
                  if ( amountCommitHard > 0){
                         state.balanceCommitHard = state.balanceCommitHard.add( amountCommitHard);
119
120
                       depositedCommitHard = depositedCommitHard.add( amountCommitHard);
122
                       IERC20(acceptToken).transfer(vcHolding, amountCommitHard);
                 }
123
                   \begin{array}{ll} \textbf{emit} & \mathsf{Deposit} \, \big( \, \_\mathsf{creditTo} \, \, , \, \, \, \_\mathsf{amountCommitSoft} \, , \, \, \, \_\mathsf{amountCommitHard} \, \big) \, ; \\ \end{array} 
125
126
                  balances[ creditTo] = state;
127
```

Listing 3.2: GaugeD1::deposit()

**Recommendation** Accommodate the above-mentioned idiosyncrasies about ERC20-related transfer() and transferFrom().

Status The issue has been fixed by this commit: 2c9d64b.

#### 3.2 Suggested Adherence of Checks-Effects-Interactions

• ID: PVE-002

Severity: Low

Likelihood: Low

• Impact: Low

• Target: GaugeD1, LPGuarge

Category: Time and State [9]

• CWE subcategory: CWE-663 [4]

#### Description

A common coding best practice in Solidity is the adherence of checks-effects-interactions principle. This principle is effective in mitigating a serious attack vector known as re-entrancy. Via this particular attack vector, a malicious contract can be reentering a vulnerable contract in a nested manner. Specifically, it first calls a function in the vulnerable contract, but before the first instance of the function call is finished, second call can be arranged to re-enter the vulnerable contract by invoking functions that should only be executed once. This attack was part of several most prominent hacks in Ethereum history, including the DAO [15] exploit, and the recent Uniswap/Lendf.Me hack [14].

We notice there are several occasions the <code>checks-effects-interactions</code> principle is violated. Using the <code>LPGuarge</code> as an example, the <code>deposit()</code> function (see the code snippet below) is provided to externally call a token contract to transfer assets. However, the invocation of an external contract requires extra care in avoiding the above <code>re-entrancy</code>.

Apparently, the interaction with the external contract (line 80) starts before effecting the update on internal states (line 88), hence violating the principle. In this particular case, if the external

contract has some hidden logic that may be capable of launching re-entrancy via the very same deposit() function.

```
75
        function deposit(uint256 _amount) nonReentrant external {
76
            require(block.number <= endBlock, "LPGAUGE: distribution over");</pre>
77
78
            claimSTACK(msg.sender);
79
80
            IERC20(token).safeTransferFrom(msg.sender, address(this), amount);
81
82
            DepositState memory _state = balances[msg.sender];
83
             state.balance = state.balance.add( amount);
84
            deposited = deposited.add( amount);
85
86
            emit Deposit(msg.sender, amount);
87
88
            balances [msg.sender] = state;
89
```

Listing 3.3: LPGuarge::deposit()

Another similar violation can be found in deposit(), withdraw() and upgradeCommit() routines within the GaugeD1 contract.

In the meantime, we should mention that the related routines are properly protected with nonReentrant and the related token contract is not vulnerable or exploitable for re-entrancy.

**Recommendation** Suggest to follow the checks-effects-interactions best practice in addition to current reentrancy protection.

**Status** The issue has been confirmed.

#### 3.3 Incompatibility with Deflationary/Rebasing Tokens

• ID: PVE-003

Severity: Informational

• Likelihood: N/A

Impact: N/A

• Target: Multiple Contracts

Category: Business Logic [8]

CWE subcategory: CWE-841 [6]

#### Description

In the Stacker.VC protocol, the GaugeD1 contract is designed to be the main entry for interaction with investing users. In particular, one entry routine, i.e., deposit(), accepts user deposits of supported assets (e.g., DAI). Naturally, the contract implements a number of low-level helper routines to transfer assets in or out of the GaugeD1 contract. These asset-transferring routines work as expected with

standard ERC20 tokens: namely the vault's internal asset balances are always consistent with actual token balances maintained in individual ERC20 token contract.

```
99
         function deposit (uint 256 amount Commit Soft, uint 256 amount Commit Hard, address
             _creditTo) nonReentrant external {
100
             require(block.number <= endBlock, "GAUGE: distribution 1 over");</pre>
101
             require (fundOpen amountCommitHard == 0, "GAUGE: !fundOpen, only soft commit
                 allowed"); // when the fund closes, soft commits are still accepted
102
             require (msg.sender == creditTo msg.sender == vaultGaugeBridge, "GAUGE: !bridge
                  for creditTo"); // only the bridge contract can use the "creditTo" to
                 credit !msg.sender
104
             claimSTACK( creditTo); // new deposit doesn't get tokens right away
106
             // transfer tokens from sender to account
107
             uint256 acceptTokenAmount = amountCommitSoft.add( amountCommitHard);
108
             if ( acceptTokenAmount > 0){
109
                 IERC20 (acceptToken).safeTransferFrom (msg.sender, address (this),
                     acceptTokenAmount);
110
             }
112
             CommitState memory _state = balances[_creditTo];
113
             // no need to update _state.tokensAccrued because that's already done in
                 _claimSTACK
114
             if ( amountCommitSoft > 0){
                  \_state . balanceCommitSoft = \_state . balanceCommitSoft . add (\_amountCommitSoft);
115
116
                 depositedCommitSoft = depositedCommitSoft.add( amountCommitSoft);
117
             if ( amountCommitHard > 0){
118
119
                 state.balanceCommitHard = state.balanceCommitHard.add( amountCommitHard);
120
                 depositedCommitHard = depositedCommitHard.add( amountCommitHard);
122
                 IERC20(acceptToken).transfer(vcHolding, amountCommitHard);
123
             }
125
             emit Deposit(_creditTo, _amountCommitSoft, _amountCommitHard);
126
             balances[_creditTo] = _state;
127
```

Listing 3.4: GaugeD1::deposit()

However, there exist other ERC20 tokens that may make certain customizations to their ERC20 contracts. One type of these tokens is deflationary tokens that charge a certain fee for every transfer () or transferFrom(). (Another type is rebasing tokens such as YAM.) As a result, this may not meet the assumption behind these low-level asset-transferring routines.

One possible mitigation is to regulate the set of ERC20 tokens that are permitted into the protocol. In our case, it is indeed possible to effectively regulate the set of tokens that can be supported. Keep in mind that there exist certain assets (e.g., USDT) that may have control switches that can be dynamically exercised to suddenly become one.

Recommendation If current codebase needs to support possible deflationary tokens, it is better to check the balance before and after the transfer()/transferFrom() call to ensure the book-keeping amount is accurate. This support may bring additional gas cost. Also, keep in mind that certain tokens may not be deflationary for the time being. However, they could have a control switch that can be exercised to turn them into deflationary tokens. One example is widely-adopted USDT.

**Status** This issue has been confirmed. However, considering the fact that this specific issue does not affect the normal operation, the team decides to address it when the need of supporting deflationary/rebasing tokens arises.

#### 3.4 Non-Functional setEndBlock()/deposit() in LPGauge

ID: PVE-004

Severity: Low

Likelihood: Low

• Impact: Low

• Target: LPGauge

• Category: Coding Practices [7]

• CWE subcategory: CWE-1126 [2]

#### Description

DeFi protocols typically have a number of system-wide parameters that can be dynamically configured on demand. The Stacker.VC protocol is no exception. Specifically, if we examine the LPGauge contract, it has defined a number of system-wide risk parameters: emissionRate, startBlock, and endBlock.

To elaborate, we show below the related configuration routines in the LPGauge contract for the risk parameter — setEndBlock().

```
67
        function setEndBlock(uint256 block) external {
68
            require(msg.sender == governance, "LPGAUGE: !governance");
            require(block.number <= endBlock, "LPGAUGE: distribution already done, must</pre>
69
                start another");
70
            require(block.number <= block, "LPGAUGE: can't set endBlock to past block");</pre>
71
72
            endBlock = block;
73
       }
74
75
        function deposit(uint256 _amount) nonReentrant external {
76
            require(block.number <= endBlock, "LPGAUGE: distribution over");</pre>
77
78
            claimSTACK(msg.sender);
79
80
            IERC20(token).safeTransferFrom(msg.sender, address(this), amount);
81
82
            DepositState memory state = balances [msg.sender];
83
            state.balance = state.balance.add( amount);
```

```
deposited = deposited.add(_amount);

deposi
```

Listing 3.5: Two LPGauge Routines: setEndBlock() and deposit()

These parameters define various aspects of the protocol operation and maintenance and need to exercise extra care when configuring or updating them. Our analysis shows the update logic on these parameters can be improved by applying more rigorous sanity checks. Based on the current implementation, certain corner cases may lead to an undesirable consequence. For example, the above configuration of endBlock is always reverted as endBlock = startBlock + 100 where startBlock is hardcoded as 300! Therefore, every call of setEndBlock() and deposit() fails. In other words, no one is able to deposit into LPGauge, which is clearly not the purpose.

**Recommendation** Revise the hardcoded startBlock to be an appropriate number. In the meantime, validate any changes regarding these system-wide parameters to ensure they fall in an appropriate range. If necessary, also consider emitting relevant events for their changes.

Status The issue has been fixed in the following commit: 87b1581.

#### 3.5 Asset Consistency in VaultGaugeBridge::newBridge()

• ID: PVE-005

Severity: Low

• Likelihood: Low

Impact: Medium

• Target: VaultGaugeBridge

• Category: Business Logics [8]

• CWE subcategory: CWE-837 [5]

#### Description

In Stacker.VC, the VaultGaugeBridge contract facilitates users in allowing them to deposit into yEarn vault and then into the GaugeD1 in a single transaction. Naturally, there is a one-to-one mapping between a yEarn vault and its gauge. To properly link a vault with its gauge, it is a need for the two to operate on the same underlying asset. For example, the yEarn vault allows for DAI-based deposits and withdraws. The associated gauge naturally has DAI as the underlying asset. If these two have different underlying assets, the link should not be successful.

If we examine the newBridge() routine in the VaultGaugeBridge contract, this routine allows for dynamic binding of the vault with a new gauge (line 55). A successful binding needs to satisfy a number of requirements. One specific example is shown as follows: \_gauge.acceptToken == \_vault.

Apparently, this requirement guarantees the consistency of the underlying asset between the vault and its associated gauge.

```
// create a new bridge, warning this allows an overwrite
function newBridge(address _vault, address _gauge) external {
   require(msg.sender == governance, "BRIDGE: !governance");

bridges[_vault] = _gauge;
}
```

Listing 3.6: VaultGaugeBridge::newBridge()

However, if we examine the above logic, the requirement of having the same underlying asset is not enforced. An unmatched deployment or configuration between vault and gauge may cause unintended consequences, including possible asset loss. With that, we suggest to maintain an invariant by ensuring their consistency when the mapping is being linked.

**Recommendation** Ensure the consistency of the underlying asset between the vault and its associated gauge. An example revision is shown below.

```
// create a new bridge, warning this allows an overwrite
function newBridge(address _vault, address _gauge) external {
    require(msg.sender == governance, "BRIDGE: !governance");
    require{_gauge.acceptToken == _vault},
    bridges[_vault] = _gauge;
}
```

Listing 3.7: VaultGaugeBridge::newBridge()

Status This issue has been resolved as the team decides to only support ibETHv2 deposits into the GaugeD1 contract, instead of many different yEarn tokens.

# 3.6 Improved Precision By Multiplication And Division Reordering

• ID: PVE-006

Severity: Low

• Likelihood: Low

Impact:Low

• Target: Multiple Contracts

• Category: Numeric Errors [10]

• CWE subcategory: CWE-190 [3]

#### Description

In Stacker.VC, there is a VCTreasuryV1 contract that creates full functionality for a trust-minimized, decentralized VC investment fund. Within this contract, there is a routine, i.e., startFund(). This

routine is called once all tokens have been issued. By calling this routine, the "fund council" actually seeds the fund with ETH, and the seeded fund will go into an ACTIVE state. At this point, new investments can then be made.

To elaborate, we show below this routine's implementation. It comes to our attention that this routine computes the maxInvestment threshold as maxInvestment = msg.value.div(max).mul(investmentCap) (line 180).

```
170
        // seed the fund with ETH and start it up. 3 years until the fund is dissolved
171
         function startFund() payable external {
172
             require(currentState == FundStates.setup, "TREASURYV1: !FundStates.setup");
173
             require(msg.sender == councilMultisig, "TREASURYV1: !councilMultisig");
174
             require(totalSupply() > 0, "TREASURYV1: invalid setup"); // means fund tokens
                 were not issued
176
             fundStartTime = block.timestamp;
177
             fundCloseTime = block.timestamp.add(ONE YEAR);
179
             initETH = msg.value;
180
             maxInvestment = msg.value.div(max).mul(investmentCap);
182
             changeFundState(FundStates.active); // set fund active!
183
```

Listing 3.8: VCTreasuryV1::startFund()

It is important to note that the lack of float support in Solidity may introduce subtle, but troublesome issue: precision loss. One possible precision loss stems from the computation when both multiplication (mul) and division (div) are involved. Specifically, the computation at line 180 is better performed as follows: maxInvestment = msg.value.mul(investmentCap).div(max).

A better approach is to avoid any unnecessary division operation that might lead to precision loss. In other words, the computation of the form A / B \* C can be converted into A \* C / B under the condition that A \* C does not introduce any overflow.

Similar precision improvements can also be observed in the arithmetic operations in other routines, e.g, getUtilization(), killQuorumRequirement(), pauseQuorumRequirement(), and \_assessFee().

**Recommendation** Avoid unnecessary precision loss due to the lack of floating support in Solidity. An example revision to the above computation is shown as: maxInvestment = msg.value. mul(investmentCap).div(max).

Status The issue has been fixed by this commit: 2c65e20.

#### 3.7 Inconsistency Between Document and Implementation

• ID: PVE-007

• Severity: Informational

• Likelihood: N/A

• Impact: N/A

Target: VCTreasuryV1

• Category: Coding Practices [7]

• CWE subcategory: CWE-1041 [1]

#### Description

There exist a misleading comment embedded among lines of solidity code, which brings unnecessary hurdles to understand and/or maintain the software.

The comment can be found in line 170 of VCTreasuryV1::startFund(). the preceding function header indicates that the fund will be dissolved after 3 years. However, the fundCloseTime state (line 177) indicates it takes one year to dissolve.

```
170
        // seed the fund with ETH and start it up. 3 years until the fund is dissolved
171
        function startFund() payable external {
172
            require(currentState == FundStates.setup, "TREASURYV1: !FundStates.setup");
173
            require(msg.sender == councilMultisig, "TREASURYV1: !councilMultisig");
174
            require(totalSupply() > 0, "TREASURYV1: invalid setup"); // means fund tokens
                were not issued
            fundStartTime = block.timestamp;
176
177
            fundCloseTime = block.timestamp.add(ONE YEAR);
179
            initETH = msg.value;
180
            maxInvestment = msg.value.div(max).mul(investmentCap);
182
             changeFundState(FundStates.active); // set fund active!
183
```

Listing 3.9: VCTreasuryV1::startFund()

**Recommendation** Ensure the consistency between documents (including embedded comments) and implementation.

**Status** This issue has been fixed by removing the aforementioned inconsistency.

# 4 Conclusion

In this audit, we have analyzed the Stacker.VC design and implementation. The protocol aims to develop a decentralized, community-owned venture capital protocol and accelerator with a focus on early-stage investment that aligns incentives of community investors with founding teams. During the audit, we notice that the current implementation is well organized and those identified issues are promptly confirmed and fixed.

Meanwhile, we need to emphasize that smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



# References

- [1] MITRE. CWE-1041: Use of Redundant Code. https://cwe.mitre.org/data/definitions/1041. html.
- [2] MITRE. CWE-1126: Declaration of Variable with Unnecessarily Wide Scope. https://cwe.mitre.org/data/definitions/1126.html.
- [3] MITRE. CWE-190: Integer Overflow or Wraparound. https://cwe.mitre.org/data/definitions/190.html.
- [4] MITRE. CWE-663: Use of a Non-reentrant Function in a Concurrent Context. https://cwe.mitre.org/data/definitions/663.html.
- [5] MITRE. CWE-837: Improper Enforcement of a Single, Unique Action. https://cwe.mitre.org/data/definitions/837.html.
- [6] MITRE. CWE-841: Improper Enforcement of Behavioral Workflow. https://cwe.mitre.org/data/definitions/841.html.
- [7] MITRE. CWE CATEGORY: Bad Coding Practices. https://cwe.mitre.org/data/definitions/1006.html.
- [8] MITRE. CWE CATEGORY: Business Logic Errors. https://cwe.mitre.org/data/definitions/840.html.
- [9] MITRE. CWE CATEGORY: Concurrency. https://cwe.mitre.org/data/definitions/557.html.

- [10] MITRE. CWE CATEGORY: Numeric Errors. https://cwe.mitre.org/data/definitions/189.html.
- [11] MITRE. CWE VIEW: Development Concepts. https://cwe.mitre.org/data/definitions/699. html.
- [12] OWASP. Risk Rating Methodology. https://www.owasp.org/index.php/OWASP\_Risk\_Rating\_Methodology.
- [13] PeckShield. PeckShield Inc. https://www.peckshield.com.
- [14] PeckShield. Uniswap/Lendf.Me Hacks: Root Cause and Loss Analysis. https://medium.com/ @peckshield/uniswap-lendf-me-hacks-root-cause-and-loss-analysis-50f3263dcc09.
- [15] David Siegel. Understanding The DAO Attack. https://www.coindesk.com/understanding-dao-hack-journalists.

