



# SMART CONTRACT AUDIT REPORT

for

## TranchessV2 Protocol



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May 20, 2022

## Document Properties

Client	Trancess Protocol
Title	Smart Contract Audit Report
Target	TrancessV2
Version	1.0
Author	Xuxian Jiang
Auditors	Stephen Bie, Jing Wang, Xuxian Jiang
Reviewed by	Xiaomi Huang
Approved by	Xuxian Jiang
Classification	Public

## Version Info

Version	Date	Author(s)	Description
1.0	May 20, 2022	Xuxian Jiang	Final Release
1.0-rc1	May 10, 2022	Xuxian Jiang	Release Candidate #1

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# 1 | Introduction

Given the opportunity to review the `TranchessV2` design document and related smart contract source code, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

## 1.1 About Tranchess

Elevating from the current tranche-fund model, `TranchessV2` aims to improve the liquidity and accessibility of the current `BISHOP/ROOK/QUEEN` token, through a newly introduced stable swap `AMM` pool. In contrast to `v1`, where users have to cross the bid-ask spread to trade in an orderbook system, or wait up to 24 hours for the creation of `QUEEN` token, the `v2` `AMM` pool allows users to freely convert from `BUSD/BTC/ETH/BNB` into `BISHOP/ROOK/QUEEN`. This `AMM` swap process is instantaneous and more cost efficient. It not only allows `ROOK` holders to timely enter and exit their leverage trading positions, but also enables `BISHOP` holders to earn extra yield through liquidity provision in the `BISHOP-BUSD` pool. The basic information of the audited protocol is as follows:

Table 1.1: Basic Information of TranchessV2

Item	Description
Name	Tranchess Protocol
Website	<a href="https://tranchess.com/">https://tranchess.com/</a>
Type	EVM Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	May 20, 2022

In the following, we show the Git repository of reviewed files and the commit hash value used in

this audit.

- <https://github.com/tranchess/contract-core.git> (6876889)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

- <https://github.com/tranchess/contract-core.git> (bcc1ad0)

## 1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (<https://t.me/peckshield>), Twitter (<http://twitter.com/peckshield>), or Email ([contact@peckshield.com](mailto:contact@peckshield.com)).

Table 1.2: Vulnerability Severity Classification

Impact	High	Medium	Low
	Critical	High	Medium
	High	Medium	Low
Low	Medium	Low	Low
Likelihood			

## 1.3 Methodology

To standardize the evaluation, we define the following terminology based on the OWASP Risk Rating Methodology [10]:

- Likelihood represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Table 1.3: The Full Audit Checklist

Category	Checklist Items
Basic Coding Bugs	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
	Revert DoS
	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
Advanced DeFi Scrutiny	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
	Digital Asset Escrow
	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
Additional Recommendations	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a checklist of items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- Semantic Consistency Checks: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

## 1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
<b>Configuration</b>	Weaknesses in this category are typically introduced during the configuration of the software.
<b>Data Processing Issues</b>	Weaknesses in this category are typically found in functionality that processes data.
<b>Numeric Errors</b>	Weaknesses in this category are related to improper calculation or conversion of numbers.
<b>Security Features</b>	Weaknesses in this category are concerned with topics like authentication, access control, confidentiality, cryptography, and privilege management. (Software security is not security software.)
<b>Time and State</b>	Weaknesses in this category are related to the improper management of time and state in an environment that supports simultaneous or near-simultaneous computation by multiple systems, processes, or threads.
<b>Error Conditions, Return Values, Status Codes</b>	Weaknesses in this category include weaknesses that occur if a function does not generate the correct return/status code, or if the application does not handle all possible return/status codes that could be generated by a function.
<b>Resource Management</b>	Weaknesses in this category are related to improper management of system resources.
<b>Behavioral Issues</b>	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
<b>Business Logic</b>	Weaknesses in this category identify some of the underlying problems that commonly allow attackers to manipulate the business logic of an application. Errors in business logic can be devastating to an entire application.
<b>Initialization and Cleanup</b>	Weaknesses in this category occur in behaviors that are used for initialization and breakdown.
<b>Arguments and Parameters</b>	Weaknesses in this category are related to improper use of arguments or parameters within function calls.
<b>Expression Issues</b>	Weaknesses in this category are related to incorrectly written expressions within code.
<b>Coding Practices</b>	Weaknesses in this category are related to coding practices that are deemed unsafe and increase the chances that an exploitable vulnerability will be present in the application. They may not directly introduce a vulnerability, but indicate the product has not been carefully developed or maintained.






comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.



## 2 | Findings

### 2.1 Summary

Here is a summary of our findings after analyzing the implementation of the `TranchessV2` protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings	
Critical	0	
High	0	
Medium	1	
Low	2	
Informational	3	
Total	6	

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in [Section 3](#).

## 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 medium-severity vulnerability, 2 low-severity vulnerabilities, and 3 informational recommendations.

Table 2.1: Key TranchessV2 Audit Findings

ID	Severity	Title	Category	Status
PVE-001	Low	Improved Logic in SwapReward::updateReward()	Business Logic	Resolved
PVE-002	Low	Revisited Logic in InterestRateBallet::syncWithVotingEscrow()	Business Logic	Resolved
PVE-003	Informational	Revisited Reentrancy Protection in VestingEscrow	Time And State	Resolved
PVE-004	Informational	Redundant State/Code Removal	Coding Practices	Resolved
PVE-005	Informational	Consistent Fee Calculation in StableSwap	Business Logic	Resolved
PVE-006	Medium	Trust Issue of Admin Keys	Security Features	Mitigated

Besides the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

## 3 | Detailed Results

### 3.1 Improved Logic in SwapReward::updateReward()

- ID: PVE-001
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: SwapReward
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

#### Description

The TranchessV2 protocol has a helper SwapReward contract that assists the reward dissemination to the associated liquidityGauge. While examining a number of parameters that are updated in SwapReward, we notice the update logic can be improved to accommodate different configurations.

Specifically, we show below the related updateReward() function from the SwapReward contract. It has a rather straightforward logic in assigning the input arguments to internal storage states. However, it comes to our attention that the given start, interval, amount, and ratePerSecond may have internal dependency. As a result, we need to properly compute ratePerSecond within the resulting interval between startTimestamp and endTimestamp where startTimestamp = start.max(block.timestamp) and endTimestamp = start.add(interval). In other words, we can compute ratePerSecond = amount.div(endTimestamp.sub(startTimestamp)).

```
28     function updateReward(  
29         address caller,  
30         uint256 amount,  
31         uint256 start,  
32         uint256 interval  
33     ) external onlyOwner {  
34         require(  
35             endTimestamp < block.timestamp && endTimestamp == lastTimestamp,  
36             "Last reward not yet expired"  
37         );  
38         require(caller != address(0));  
39         ratePerSecond = amount.div(interval);
```

```

40     startTimestamp = start.max(block.timestamp);
41     endTimestamp = start.add(interval);
42     lastTimestamp = startTimestamp;
43     IERC20(rewardToken).safeTransferFrom(
44         msg.sender,
45         address(this),
46         ratePerSecond.mul(interval)
47     );
48 }

```

Listing 3.1: SwapReward::updateReward()

**Recommendation** Properly compute the `ratePerSecond` from the given input arguments in the above `updateReward()` function.

**Status** The issue has been fixed by the following commit: 78c5fbd.

## 3.2 Revisited Logic in InterestRateBallot::syncWithVotingEscrow()

- ID: PVE-002
- Severity: Low
- Likelihood: Low
- Impact: Medium
- Target: InterestRateBallot
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

### Description

The TranchessV2 protocol has primary markets that are designed to allow users to create, redeem, split, or merge `TRANCHE_Q` tokens. In the meantime, it also has a `InterestRateBallot` contract that allows to adjust the interest rate. This `InterestRateBallot` contract will need to timely synchronize with the `VotingEscrowV2` contract to keep track of the voter amount for related users. While examining the synchronization logic, we notice the current implementation can be improved.

To elaborate, we show below the related `syncWithVotingEscrow()` function, which updates the vote amount for the given account. It comes to our attention that the current logic simply returns when the user does not have any locked balance (with 0 amount or expired position). Our analysis shows that there is a need to continue the execution to reset the voter amount. Note that the same issue is also applicable to other contracts, including `FeeDistributor`.

```

122     function syncWithVotingEscrow(address account) external override {
123         Voter memory voter = voters[account];
124         if (voter.amount == 0) {
125             return; // The account did not voted before

```

```

126     }
127
128     IVotingEscrow.LockedBalance memory lockedBalance = votingEscrow.getLockedBalance
        (account);
129     if (lockedBalance.amount == 0 lockedBalance.unlockTime <= block.timestamp) {
130         return;
131     }
132
133     // update scheduled unlock
134     scheduledUnlock[voter.unlockTime] = scheduledUnlock[voter.unlockTime].sub(voter.
        amount);
135     scheduledUnlock[lockedBalance.unlockTime] = scheduledUnlock[lockedBalance.
        unlockTime].add(
136         lockedBalance.amount
137     );
138
139     scheduledWeightedUnlock[voter.unlockTime] = scheduledWeightedUnlock[voter.
        unlockTime].sub(
140         voter.amount * voter.weight
141     );
142     scheduledWeightedUnlock[lockedBalance.unlockTime] = scheduledWeightedUnlock[
        lockedBalance.unlockTime
143     ]
144     .add(lockedBalance.amount * voter.weight);
145
146     emit Voted(
147         account,
148         voter.amount,
149         voter.unlockTime,
150         voter.weight,
151         lockedBalance.amount,
152         lockedBalance.unlockTime,
153         voter.weight
154     );
155
156     // update voter amount per account
157     voters[account].amount = lockedBalance.amount;
158     voters[account].unlockTime = lockedBalance.unlockTime;
159
160 }

```

Listing 3.2: InterestRateBallot::syncWithVotingEscrow()

**Recommendation** Properly revise the above logic to update the voter amount, including the cases where the user may have a 0 balance or expired lock position.

**Status** The issue has been fixed by the following commit: 3d0b6c4.

### 3.3 Revisited Reentrancy Protection in VestingEscrow

- ID: PVE-002
- Severity: Informational
- Likelihood: N/A
- Impact: N/A
- Target: VestingEscrow
- Category: Time and State [8]
- CWE subcategory: CWE-663 [3]

#### Description

A common coding best practice in Solidity is the adherence of `checks-effects-interactions` principle. This principle is effective in mitigating a serious attack vector known as `re-entrancy`. Via this particular attack vector, a malicious contract can be reentering a vulnerable contract in a nested manner. Specifically, it first calls a function in the vulnerable contract, but before the first instance of the function call is finished, second call can be arranged to re-enter the vulnerable contract by invoking functions that should only be executed once. This attack was part of several most prominent hacks in Ethereum history, including the DAO [13] exploit, and the recent Uniswap/Lendf.Me hack [12].

We notice the current protocol has taken into consideration the `re-entrancy` protection in current contracts. For example, the `nonReentrant` modifier is widely used, even in certain cases where it is safe to not use the `re-entrancy` protection. For example, we show below the `VestingEscrow::claim()` function, which allows to claim the tokens that were vested before. Our analysis shows that the associated `nonReentrant` modifier is not necessary as it follows strictly the above `checks-effects-interactions` principle with the token-transfer at the end of this function.

```
99      /// @notice Claim tokens which have vested
100     function claim() external nonReentrant {
101         uint256 timestamp = disabledAt;
102         if (timestamp == 0) {
103             timestamp = block.timestamp;
104         }
105         uint256 claimable = _totalVestedOf(timestamp).sub(totalClaimed);
106         totalClaimed = totalClaimed.add(claimable);
107         IERC20(token).safeTransfer(recipient, claimable);
108
109         emit Claim(claimable);
110     }
```

Listing 3.3: VestingEscrow::claim()

**Recommendation** Remove the unnecessary `nonReentrant` modifier in the above `claim()` function.

**Status** The issue has been fixed by the following commit: 9a211be.

### 3.4 Redundant State/Code Removal

- ID: PVE-004
- Severity: Informational
- Likelihood: N/A
- Impact: N/A
- Target: FundV3
- Category: Coding Practices [6]
- CWE subcategory: CWE-563 [2]

#### Description

The TranchessV2 protocol makes good use of a number of reference contracts, such as ERC20, SafeERC20, Ownable, and ReentrancyGuard, to facilitate its code implementation and organization. For example, the FundV3 smart contract has so far imported at least five reference contracts. However, we observe the inclusion of certain unused code or the presence of unnecessary redundancies that can be safely removed.

For example, if we examine closely the FundV3 contract, there are a number of events that are defined, but not used. Examples include the `ObsoletePrimaryMarketAdded` and `NewPrimaryMarketAdded` structures.

```

30     event ProfitReported(uint256 profit, uint256 performanceFee);
31     event LossReported(uint256 loss);
32     event ObsoletePrimaryMarketAdded(address obsoletePrimaryMarket);
33     event NewPrimaryMarketAdded(address newPrimaryMarket);
34     event DailyProtocolFeeRateUpdated(uint256 newDailyProtocolFeeRate);
35     event TwapOracleUpdated(address newTwapOracle);
36     event AprOracleUpdated(address newAprOracle);
37     event BallotUpdated(address newBallot);
38     event FeeCollectorUpdated(address newFeeCollector);
39     event ActivityDelayTimeUpdated(uint256 delayTime);
40     event SplitRatioUpdated(uint256 newSplitRatio);
41     event FeeDebtPaid(uint256 amount);

```

Listing 3.4: Various Events Defined in FundV3

**Recommendation** Consider the removal of the non-used events with a simplified, consistent implementation.

**Status** The issue has been fixed by the following commit: `bcc1ad0`.



### 3.5 Consistent Fee Calculation in StableSwap

- ID: PVE-005
- Severity: Informational
- Likelihood: N/A
- Impact: N/A
- Target: StableSwap
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

#### Description

The TranchessV2 protocol has an innovative StableSwap AMM contract, which allows users to freely convert from BUSD/BTC/ETH/BNB into BISHOP/ROOK/QUEEN without the need of crossing the bid-ask spread to trade in an orderbook system, or waiting up to 24 hours for the creation of QUEEN token. The new AMM has the built-in support of collecting swap fee and protocol fee. Our analysis shows the fee collection may be improved.

To elaborate, we show below the related `getQuoteOut()/getQuoteIn()` functions. As their names indicate, these two functions are proposed to calculate the quote amount with the given `baseIn` or the expected `baseOut`. It comes to our attention that the first function computes the fee on the output token, while the second function computes the fee on the input token. The other two functions `getBaseOut()` and `getBaseIn()` share the same issue.

```

181     function getQuoteOut(uint256 baseIn) external view override returns (uint256
        quoteOut) {
182         (uint256 oldBase, uint256 oldQuote, , , , ) =
183             _getRebalanceResult(fund.getRebalanceSize());
184         uint256 newBase = oldBase.add(baseIn);
185         uint256 ampl = getAmpl();
186         uint256 oraclePrice = getOraclePrice();
187         uint256 d = _getD(oldBase, oldQuote, ampl, oraclePrice);
188         uint256 newQuote = _getQuote(ampl, newBase, oraclePrice, d);
189         quoteOut = oldQuote.sub(newQuote).sub(1); // -1 just in case there were some
            rounding errors
190         uint256 fee = quoteOut.multiplyDecimal(feeRate);
191         quoteOut = quoteOut.sub(fee);
192     }
193
194     function getQuoteIn(uint256 baseOut) external view override returns (uint256 quoteIn
        ) {
195         (uint256 oldBase, uint256 oldQuote, , , , ) =
196             _getRebalanceResult(fund.getRebalanceSize());
197         uint256 newBase = oldBase.sub(baseOut);
198         uint256 ampl = getAmpl();
199         uint256 oraclePrice = getOraclePrice();
200         uint256 d = _getD(oldBase, oldQuote, ampl, oraclePrice);
201         uint256 newQuote = _getQuote(ampl, newBase, oraclePrice, d);

```

```

202     quoteIn = newQuote.sub(oldQuote).add(1); // 1 just in case there were some
        rounding errors
203     uint256 fee = quoteIn.mul(feeRate).div(uint256(1e18).sub(feeRate));
204     quoteIn = quoteIn.add(fee);
205 }

```

Listing 3.5: StableSwap::getQuoteOut()/getQuoteIn()

**Recommendation** Properly compute the fee consistently in the above set of functions.

**Status** The issue has been resolved as the team confirms it is one of the design choices, which makes it unique comparing to UNI or Curve. In particular, since the base asset is either BISHOP or QUEEN and both subject to rebalance, the team feels it better to collect all swap fees only in quote asset, making storing and distributing of the fees much easier.

## 3.6 Trust Issue of Admin Keys

- ID: PVE-006
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: Multiple Contracts
- Category: Security Features [5]
- CWE subcategory: CWE-287 [1]

### Description

In the TranchessV2 protocol, there is a privileged owner account that plays a critical role in governing and regulating the protocol-wide operations (e.g., configuring various system parameters). It also has the privilege to control or govern the flow of assets managed by this protocol. Our analysis shows that the privileged account needs to be scrutinized. In the following, we examine the privileged account and their related privileged accesses in current contracts.

```

569     function _updateFundCap(uint256 newCap) private {
570         fundCap = newCap;
571         emit FundCapUpdated(newCap);
572     }
573
574     function updateFundCap(uint256 newCap) external onlyOwner {
575         _updateFundCap(newCap);
576     }
577
578     function _updateRedemptionFeeRate(uint256 newRedemptionFeeRate) private {
579         require(newRedemptionFeeRate <= MAX_REDEMPTION_FEE_RATE, "Exceed max redemption
            fee rate");
580         redemptionFeeRate = newRedemptionFeeRate;
581         emit RedemptionFeeRateUpdated(newRedemptionFeeRate);

```

```

582     }
583
584     function updateRedemptionFeeRate(uint256 newRedemptionFeeRate) external onlyOwner {
585         _updateRedemptionFeeRate(newRedemptionFeeRate);
586     }
587
588     function _updateMergeFeeRate(uint256 newMergeFeeRate) private {
589         require(newMergeFeeRate <= MAX_MERGE_FEE_RATE, "Exceed max merge fee rate");
590         mergeFeeRate = newMergeFeeRate;
591         emit MergeFeeRateUpdated(newMergeFeeRate);
592     }
593
594     function updateMergeFeeRate(uint256 newMergeFeeRate) external onlyOwner {
595         _updateMergeFeeRate(newMergeFeeRate);
596     }

```

Listing 3.6: Example Privileged Operations in the PrimaryMarketV3 Contract

In addition, we notice the `owner` account that is able to adjust various protocol-wide risk parameters. Apparently, if the privileged `owner` account is a plain EOA account, this may be worrisome and pose counter-party risk to the protocol users. Note that a multi-sig account could greatly alleviate this concern, though it is still far from perfect. Specifically, a better approach is to eliminate the administration key concern by transferring the role to a community-governed DAO. In the meantime, a timelock-based mechanism can also be considered as mitigation.

Moreover, it should be noted that if current contracts need to be deployed behind a proxy, there is a need to properly manage the proxy-admin privileges as they fall in this trust issue as well.

**Recommendation** Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

**Status** This issue has been confirmed and partially mitigated. Especially, for all admin-level operations, the current mitigation is to adopt the standard `TimelockController` with multi-sig `TranchessV2` account as the proposer, and a minimum delay of 1 days. The `TimelockController` address on BSC chain is `0x4BB3AeB5Ba75bC6A44177907B54911b19d1cF8f7`.

## 4 | Conclusion

In this audit, we have analyzed the design and implementation of the `TranchessV2` protocol, which greatly improves the liquidity and accessibility of the current `BISHOP/ROOK/QUEEN` token, through a newly introduced stable swap `AMM` pool. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and fixed.

Moreover, we need to emphasize that `Solidity`-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



## References

- [1] MITRE. CWE-287: Improper Authentication. <https://cwe.mitre.org/data/definitions/287.html>.
- [2] MITRE. CWE-563: Assignment to Variable without Use. <https://cwe.mitre.org/data/definitions/563.html>.
- [3] MITRE. CWE-663: Use of a Non-reentrant Function in a Concurrent Context. <https://cwe.mitre.org/data/definitions/663.html>.
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