以下的報告是根據助教提供的任務描述而寫的。引述助教為

Hw2範例.docx檔是該作業的繳交格式範例

### 要做的事情有:

- 1.用自己的話敘述Binary Trie / Fourbit Trie比較
- 2.使用我們提供的程式碼去跑binary\_trie.c 4\_bit\_trie.c 並分析自己實作的結果
- 3.用自己的話敘述leaf pushing algorithm
- 4.使用我們提供的程式碼去跑leaf\_push.c (執行的部分可以參考範例檔案的指令)

請注意:第二部提到的'自己的實作'我解讀為我們需要實作binary trie和 fourbit trie. 我的 github repository為https://github.com/felixweiss1999/trie\_performance。我希望這個報告滿足您的要求!

### 1. Binary Trie / Fourbit Trie比較

Binary Trie / Fourbit Trie都在最長前綴匹配中是兩種常用的資料結構。

#### Binary Trie:

- 1. 結構:二元搜尋樹, 每個節點最多有兩個子節點:左子節點和右子節點。
- 2. 內存佔用:每個節點需要存儲一個二進制位,因此在內存方面佔用較少。

#### Fourbit Trie:

- 1. 結構: 是一種特殊的前綴樹, 每個節點代表一個四位二進制數(0-15), 並且可以具有多個子節點。
- 2. 內存佔用:每個節點需要存儲四個二進制位. 因此在內存方面佔用相對較多。

最長前綴匹配是通過遍歷樹結構來查尋與輸入IP地址的最長前綴匹配項。選擇哪種資料結構取決 於具體的應用需求和性能要求。

## **2.** 跑binary\_trie.c 4\_bit\_trie.c 並分析自己實作的結果 跑binary\_trie.c的輸出:

```
# of prefix = 997952
# of prefix is len < 16 = 4409
# of prefix is len = 16 = 14910
max_segment_size =1852, index = 48771
empty_segment =38053 , non_empty_segment =27482
number of layer0: 32
number of sum: 0</pre>
```

# 跑4\_bit\_trie.c的輸出:

```
0
0
0
memory access time table:

1 times = 0
2 times = 0
3 times = 1
4 times = 189
5 times = 1895
6 times = 89968
7 times = 84645
9 times = 34645
9 times = 0
10 times = 0
11 times = 0
12 times = 0
13 times = 0
14 times = 0
15 times = 0
16 times = 0
17 times = 0
18 times = 0
19 times = 0
20 times = 0
21 times = 0
22 times = 0
23 times = 0
24 times = 0
25 times = 0
26 times = 0
27 times = 0
28 times = 0
29 times = 0
30 times = 0
31 times = 0
31 times = 0
32 times = 0
33 times = 0
34 times = 0
35 times = 0
36 times = 0
37 times = 0
38 times = 0
39 times = 0
39 times = 0
39 times = 0
30 times = 0
31 times = 0
31 times = 0
32 times = 0
33 times = 0
34 times = 0
35 times = 0
36 times = 0
37 times = 0
38 times = 0
39 times = 0
39 times = 0
30 times = 0
31 times = 0
31 times = 0
32 times = 0
33 times = 0
34 times = 0
35 times = 0
36 times = 0
37 times = 0
38 times = 0
39 times = 0
39 times = 0
31 times = 0
31 times = 0
31 times = 0
32 times = 0
33 times = 0
34 times = 0
35 times = 0
36 times = 0
37 times = 0
38 times = 0
39 times = 0
39 times = 0
31 times = 0
```

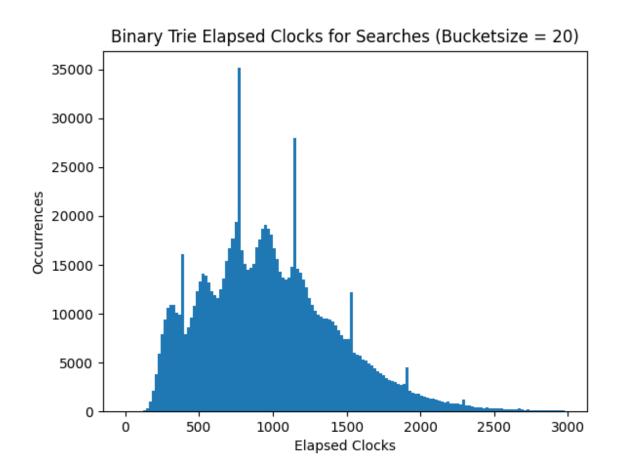
自己的實作結果(Build: ipv4\_rrc\_all\_90build.txt, Insert: ipv4\_rrc\_all10insert.txt, Search: ipv4\_rrc\_all 90build.txt) Executed on Windows. MaxClock, MinClock measure search times.

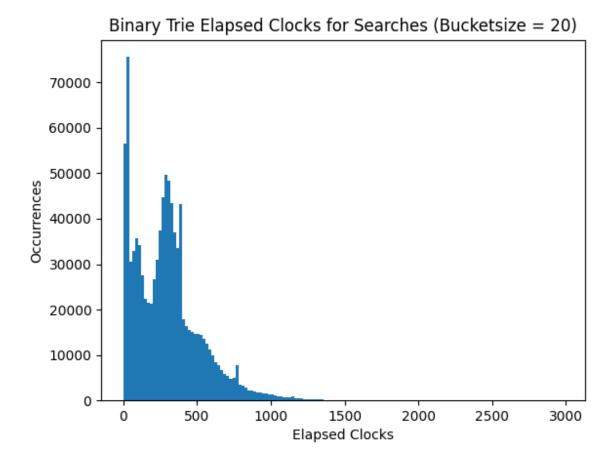
	Fourbit Trie	Binary Trie
Average Build Time	541.483055	1091.195962
Average Insert Time	292.669077	1018.485173
Average Search Time	470.270782	1026.658968
Number Of Nodes (before insert)	1096379	2282296
Total memory requirement	(8 * 16 + 4) * 1096379	(8+8+4) * 2282296
	= 141330 KB	= 44576 KB
(MaxClock, MinClock)	(90269, 18)	(57380, 57)

#### 我的評論:

正如預期的那樣,與我實作的binary trie相比,我fourbit trie的實作在建構、插入和搜索的方便掌握明顯的優勢。這是可以預料的,因為fourbit trie的深度應該比binary trie小得多,從而減少了緩慢的記憶體查詢次數。它也具有相對少的節點,雖然其個別節點佔用更多空間,導致其佔用的総空間超過對應的binary trie的三倍以上。我發現測量的MaxClock值在不同運行中變化很大,因此應該對這些極端值持保留態度。我不確定是什麼原因導致了這些極端情況的出現。

### 個別搜尋時間分佈如下:





# 3. 敘述Leaf pushing algorithm

葉推算法將binary trie中每個內部節點的前綴搬移到其對應的葉節點,確保所有前綴僅存儲在葉節點中。舉個例子,考慮一個擁有三個前綴的樹:10\*、110\*和1110\*。應用葉推算法後,結果的樹將包含四個前綴:110\*、1110\*、1111\*和01\*。前綴01\*和1111\*從前綴10\*繼承了路由資訊。儘管前綴數量增加了,葉推算法下的每個前綴仍然是獨立且互不相交的。

# 4. 跑1level\_push.c的輸出:

```
Before one-level push have prefix nodes: 997952
After one-level push have prefix nodes: 959645
add space: 16206
Avg. Build Time: 1667
Total memory requirement: 53871 KB
There are 3258148 nodes in binary trie
Avg. Search: 1
(MaxClock, MinClock) = ( 1,
                                     1)
997952
0
0
0
0
0
0
0
0000
0
0
0
0000
0
0
00000
0
0
0
00000
0
0
0000
0
0
0
0
0
0
Avg. insert Time:1402
```