



EMBRY-RIDDLE
Aeronautical University™
DAYTONA BEACH, FLORIDA

MemoryRanger Prevents Hijacking **FILE_OBJECT** structures in Windows Kernel

Igor Korkin

2019 ADFSL Conference

WHOAMI

- MEPhI Alumni, PhD in Cyber Security
- Area of interest is Windows Kernel security:
 - Memory Forensics
 - Rootkits Detection
 - Bare-Metal Hypervisors
- Fan of cross-disciplinary research - igorkorkin.blogspot.com
- Love traveling and powerlifting - @igor.korkin

AGENDA



AGENDA

- FILE_OBJECT hijacking: details and demo

-

-

AGENDA

- FILE_OBJECT hijacking: details and demo
- A history of related OS components and memory protection issues
-

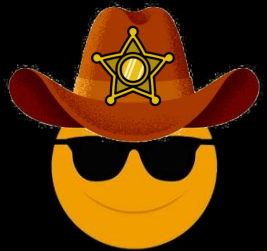
AGENDA

- FILE_OBJECT hijacking: details and demo
- A history of related OS components and memory protection issues
- MemoryRanger hypervisor protects sensitive kernel memory



AGENDA

- FILE_OBJECT hijacking: details and demo
- A history of related OS components and memory protection issues
- MemoryRanger hypervisor protects sensitive kernel memory



File Manager in Kernel Mode

COMPANY BUDGET 2019

	Budget
Team A	\$3,000,000
Team B	\$7,000,000



ENG

9:43 PM





Document1 - Word

File Home Insert Design Layout Referenc... Mailings Review View Tell me what you can do about this problem Share

COMPANY BUDGET 2019

	Budget
Team A	\$3,000,000
Team B	\$7,000,000

Page 1 of 1 10 words English (United States) ENG 9:43 PM

ZWCREATEFILE ROUTINE

```
NTSTATUS ZwCreateFile(..., ShareAccess, ...);
```

ZWCREATEFILE ROUTINE

```
NTSTATUS ZwCreateFile(..., ShareAccess, ...);
```

ShareAccess

- ShareAccess flag determines whether other drivers can access the opened file.
- Calling ZwCreateFile with ShareAccess=0 gives the caller exclusive access to the file.

Document1 - Word

File Home Insert Design Layout Refer Mailin Review View Tell m Igor K... Share

COMPANY BUDGET 2019

	Budget
Team A	\$3,000,000
Team B	00,000

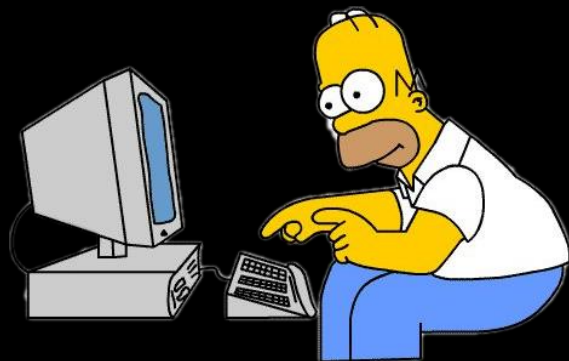
Page 1 of 1 10 words English (United States) 100%

Windows Taskbar: File Explorer, Word, PowerPoint, etc. Time: 9:43 PM





VS.



The Boss's Driver
`ZwCreateFile("budget.txt",
ShareAccess=0)`



The Attacker's Driver



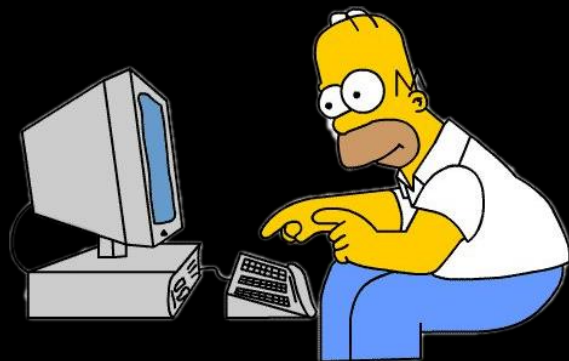
exclusive
mode



budget.txt



VS.



The Boss's Driver
`ZwCreateFile("budget.txt",
ShareAccess=0)`



The Attacker's Driver

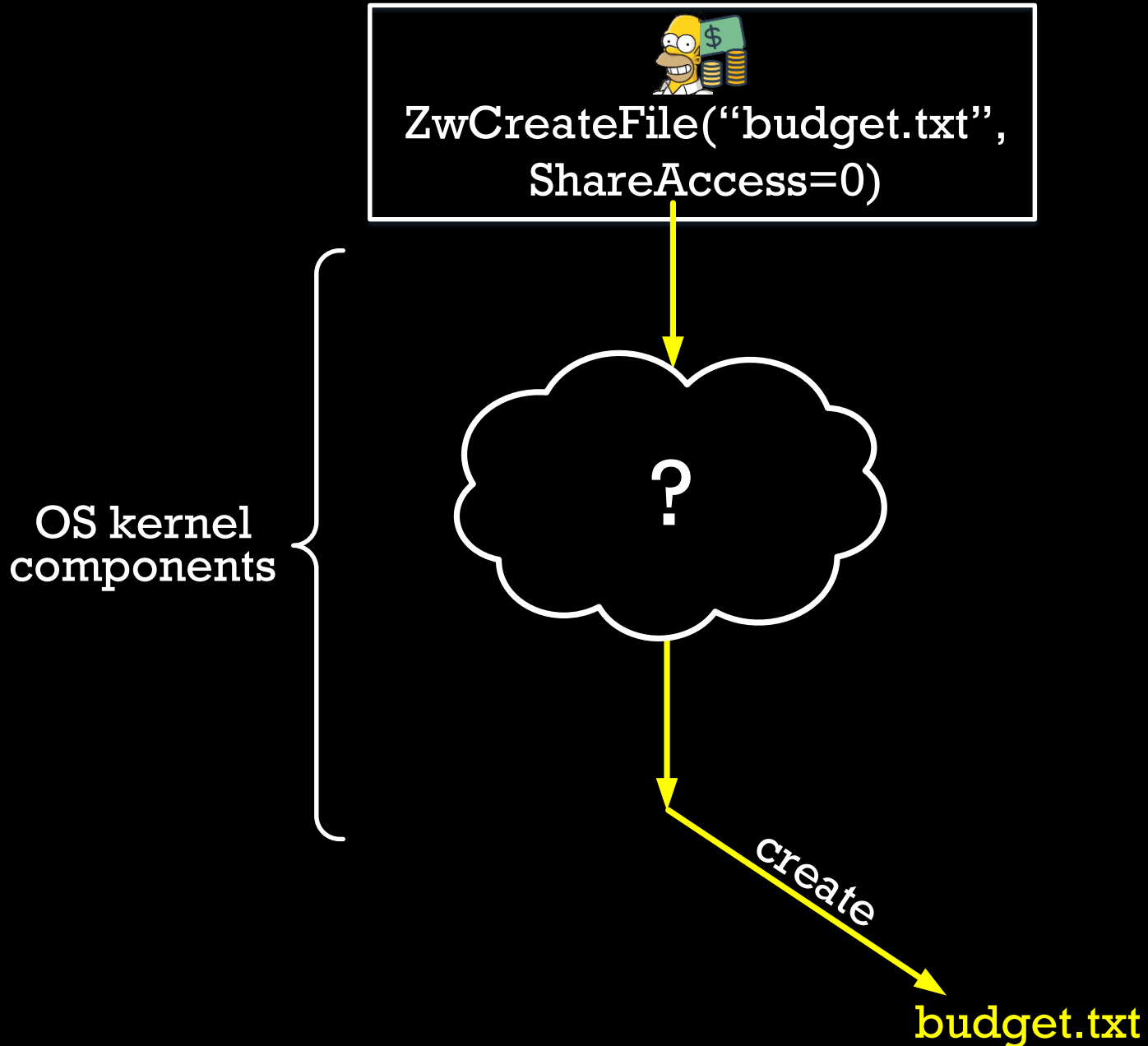


exclusive
mode

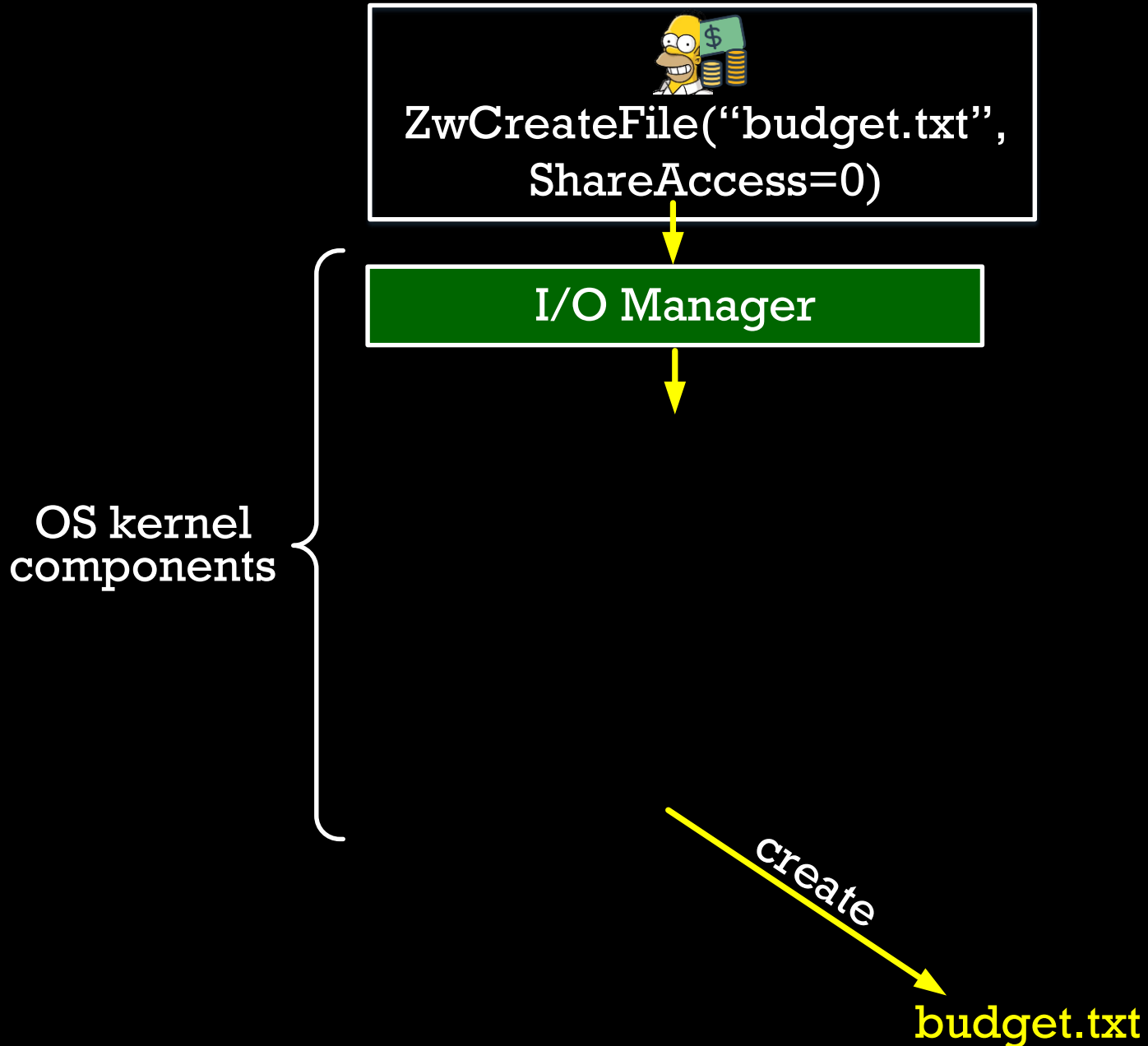


budget.txt

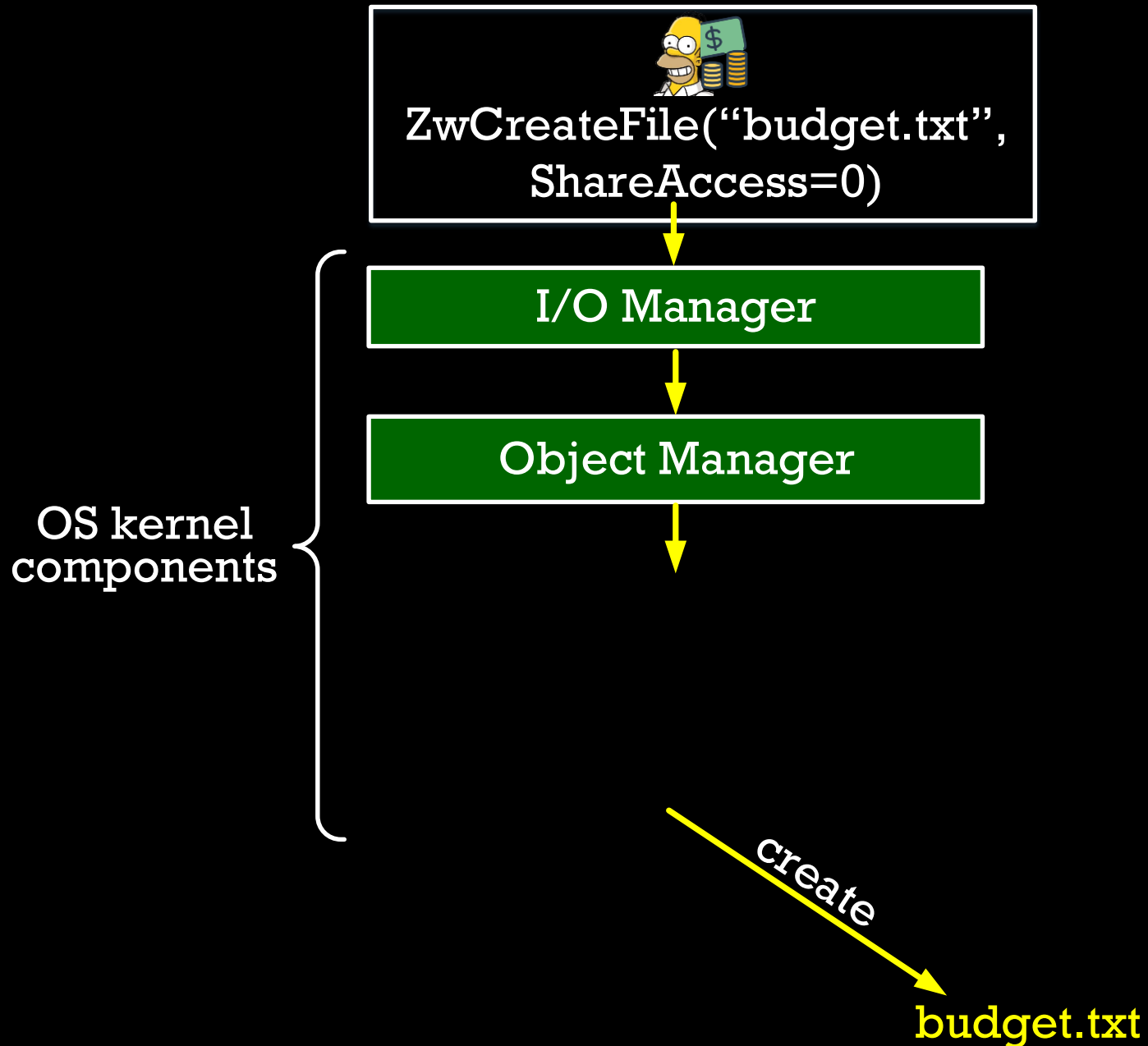
FILE SYSTEM ROUTINES IN WINDOWS KERNEL



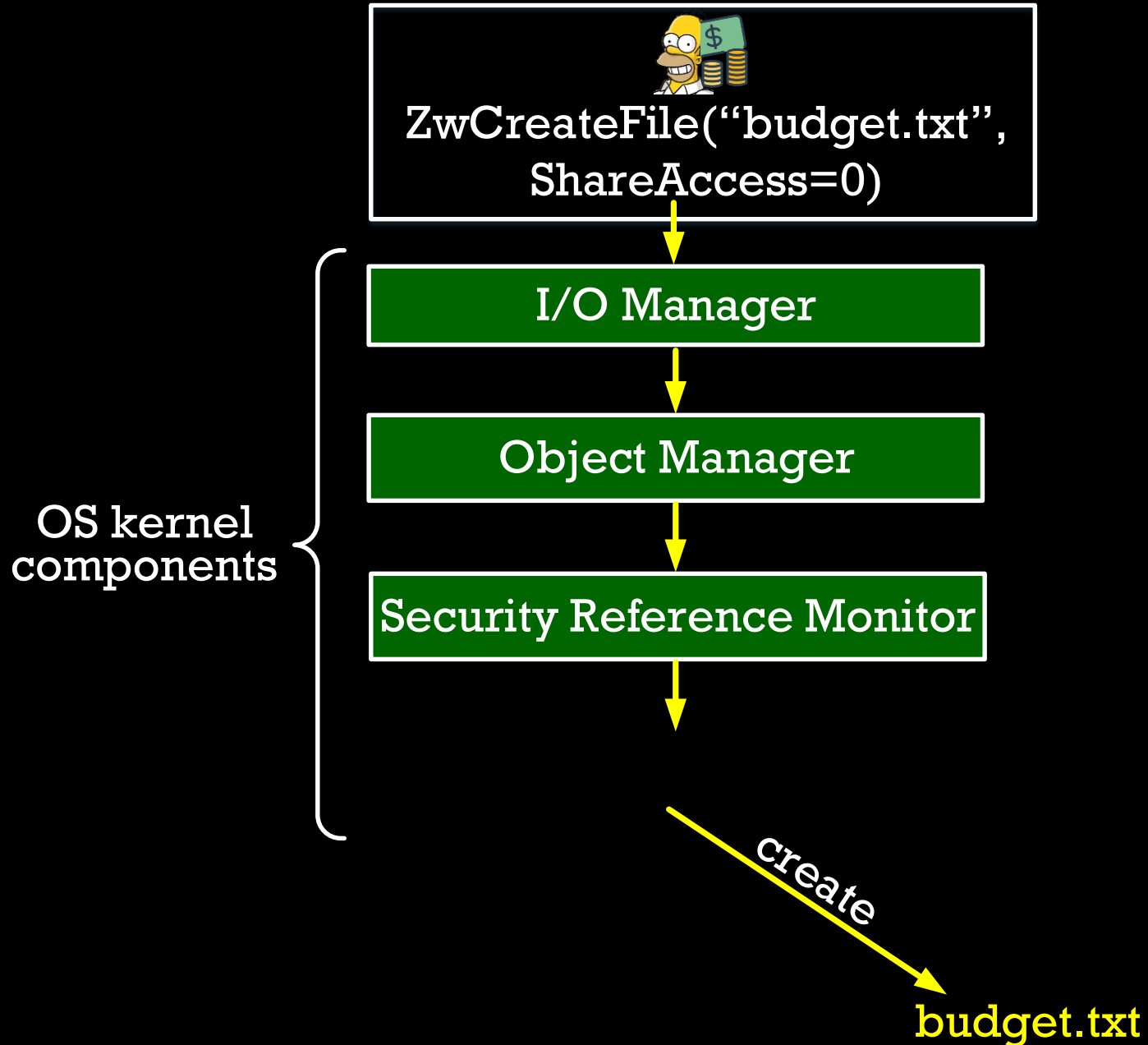
FILE SYSTEM ROUTINES IN WINDOWS KERNEL



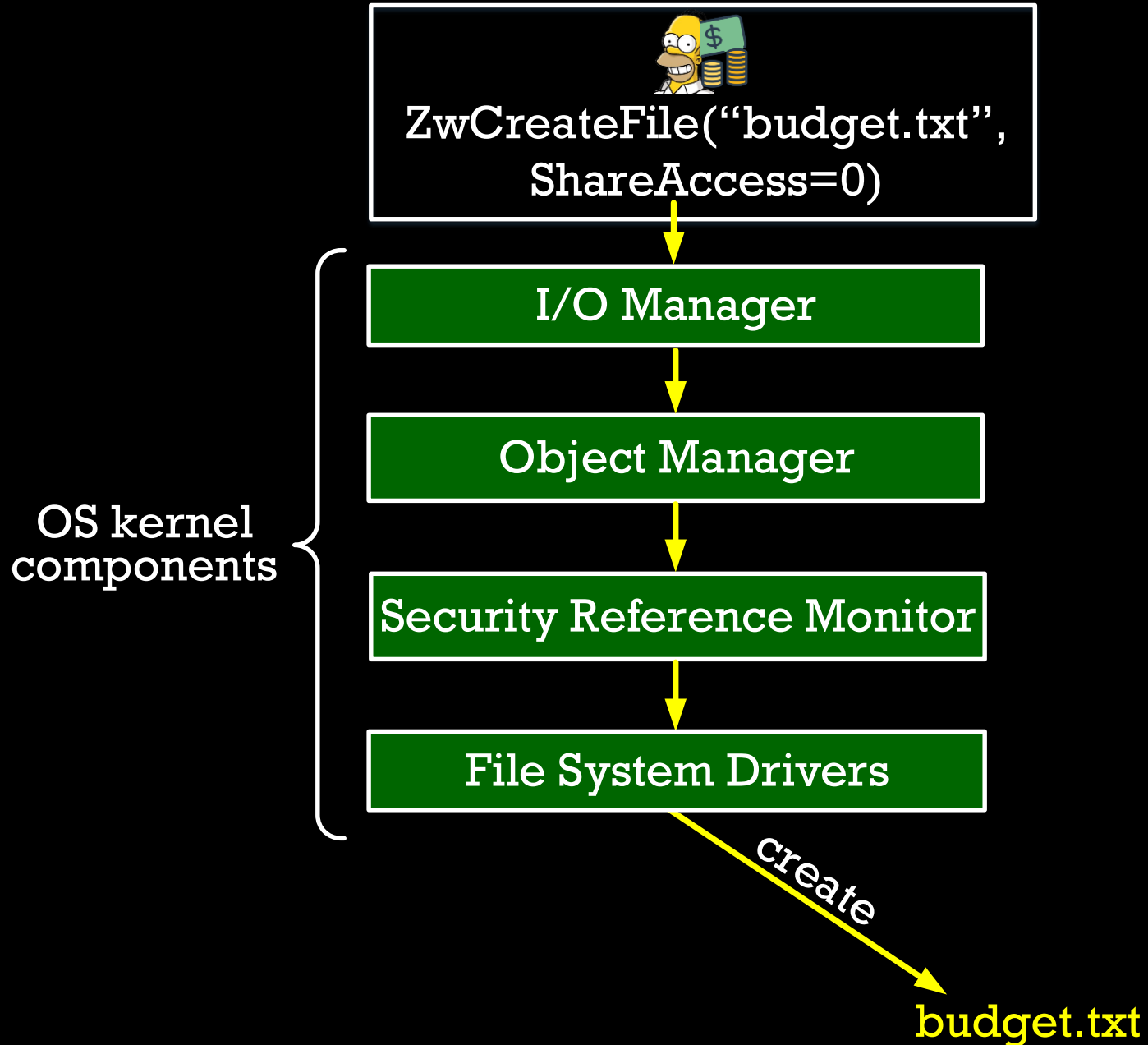
FILE SYSTEM ROUTINES IN WINDOWS KERNEL



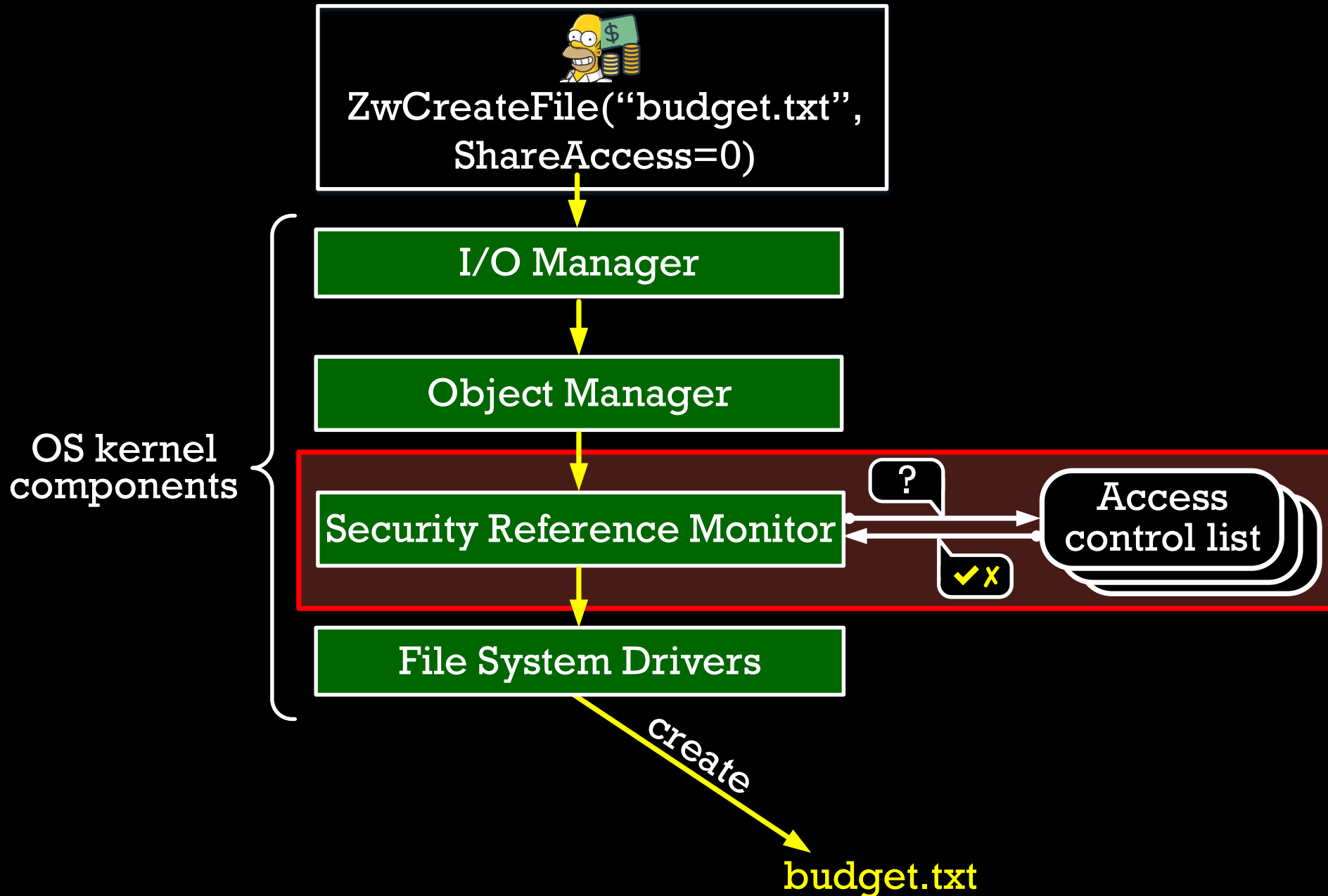
FILE SYSTEM ROUTINES IN WINDOWS KERNEL



FILE SYSTEM ROUTINES IN WINDOWS KERNEL



FILE SYSTEM ROUTINES IN WINDOWS KERNEL



FILE SYSTEM ROUTINES IN WINDOWS KERNEL



ZwCreateFile("budget.txt",
ShareAccess=0)



ZwCreateFile("hijacker.txt")

I/O Manager

Object Manager

Security Reference Monitor

File System Drivers

create

budget.txt

FILE SYSTEM ROUTINES IN WINDOWS KERNEL



ZwCreateFile("budget.txt",
ShareAccess=0)



ZwCreateFile("hijacker.txt")

I/O Manager

Object Manager

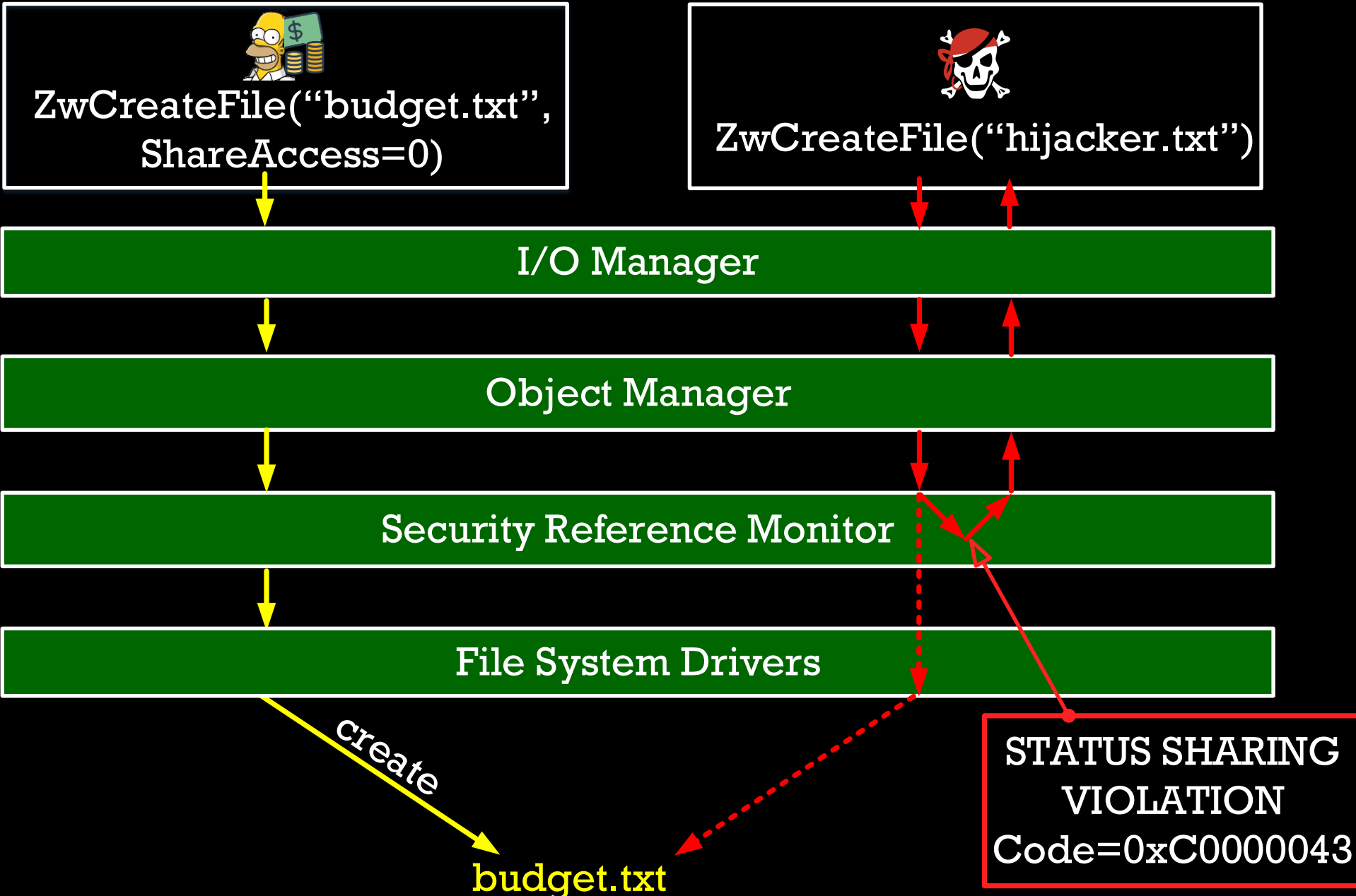
Security Reference Monitor

File System Drivers

create

budget.txt

STATUS SHARING
VIOLATION
Code=0xC0000043



FILE SYSTEM ROUTINES IN WINDOWS KERNEL



ZwCreateFile("budget.txt",
ShareAccess=0)

I/O Manager

Object Manager

Security Reference Monitor

File System Drivers



File Handle



FILE_OBJECT



create

budget.txt

FILE SYSTEM ROUTINES IN WINDOWS KERNEL

 `ZwCreateFile("budget.txt",
ShareAccess=0)` `ZwReadFile( File Handle)`
`ZwWriteFile( File Handle)`

I/O Manager

Object Manager

Security Reference Monitor

 File Handle

 FILE_OBJECT

File System Drivers

create

read-write

budget.txt

FILE SYSTEM ROUTINES IN WINDOWS KERNEL

ZwCreateFile("budget.txt",
ShareAccess=0)



ZwReadFile(File Handle)
ZwWriteFile(File Handle)



Hey! I'm a hacker-attacker!

I/O Manager

Object Manager

Security Reference Monitor



File Handle



FILE_OBJECT



File System Drivers

create

read-write

budget.txt

FILE SYSTEM ROUTINES IN WINDOWS KERNEL

ZwCreateFile("budget.txt",
ShareAccess=0)



ZwReadFile(File Handle)
ZwWriteFile(File Handle)



ZwCreateFile("hijacker.txt")

I/O Manager

Object Manager

Security Reference Monitor

File Handle

FILE_OBJECT

File Handle

FILE_OBJECT

File System Drivers

create

read-write

budget.txt

1. Create a file

hijacker.txt

FILE SYSTEM ROUTINES IN WINDOWS KERNEL

ZwCreateFile("budget.txt",
ShareAccess=0)



ZwReadFile(File Handle)
ZwWriteFile(File Handle)



FILE_OBJECT Hijacking

I/O Manager

Object Manager

Security Reference Monitor

File Handle

File Handle

FILE_OBJECT

2.Copy

FILE_OBJECT

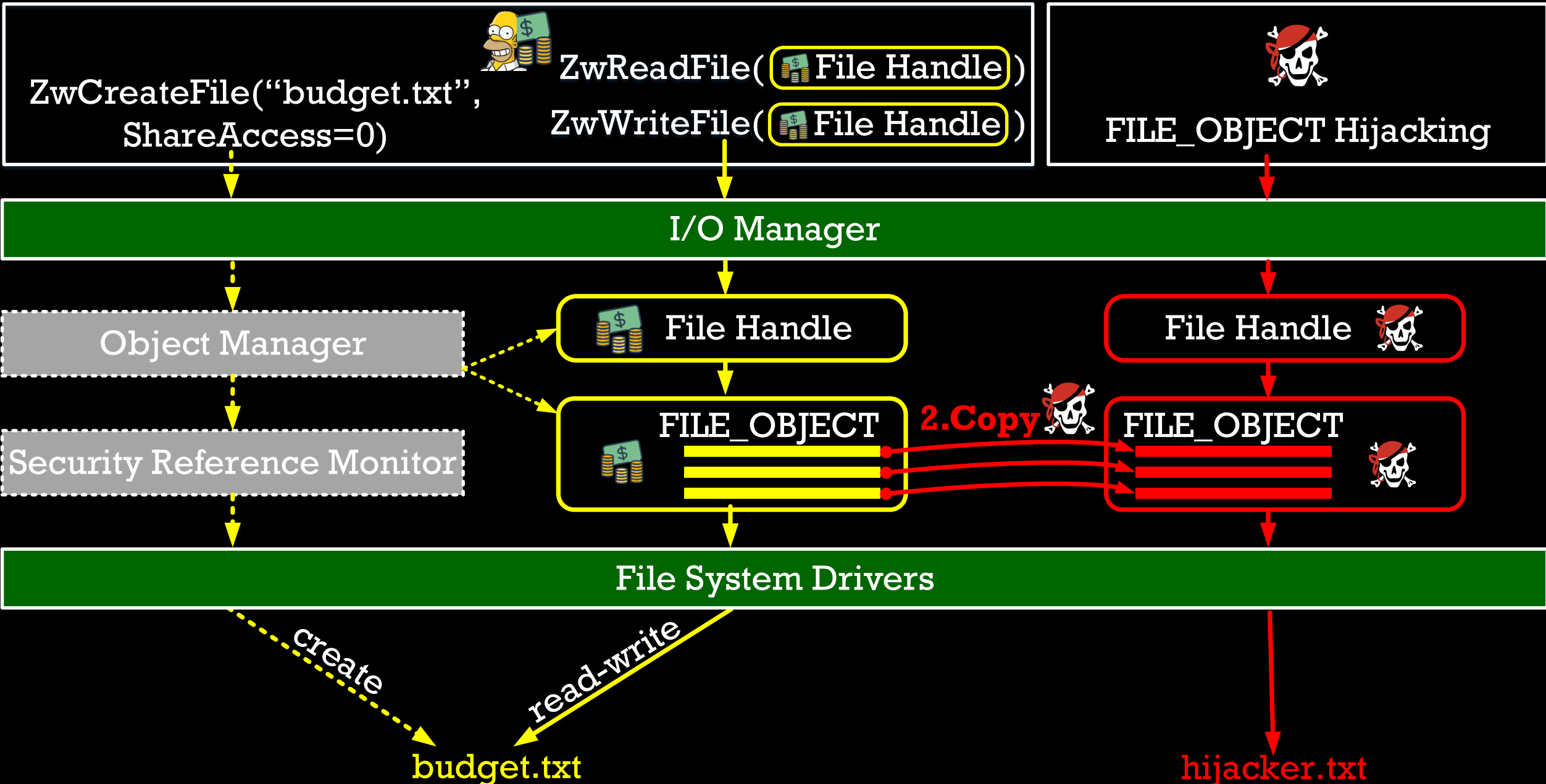
File System Drivers

create

read-write

budget.txt

hijacker.txt



FILE SYSTEM ROUTINES IN WINDOWS KERNEL

ZwCreateFile("budget.txt",
ShareAccess=0)



ZwReadFile(File Handle)
ZwWriteFile(File Handle)



ZwReadFile(File Handle)
ZwWriteFile(File Handle)

I/O Manager

Object Manager

Security Reference Monitor



File Handle



FILE_OBJECT



File Handle



FILE_OBJECT



File System Drivers

create

read-write

3. Read & Write



budget.txt

hijacker.txt

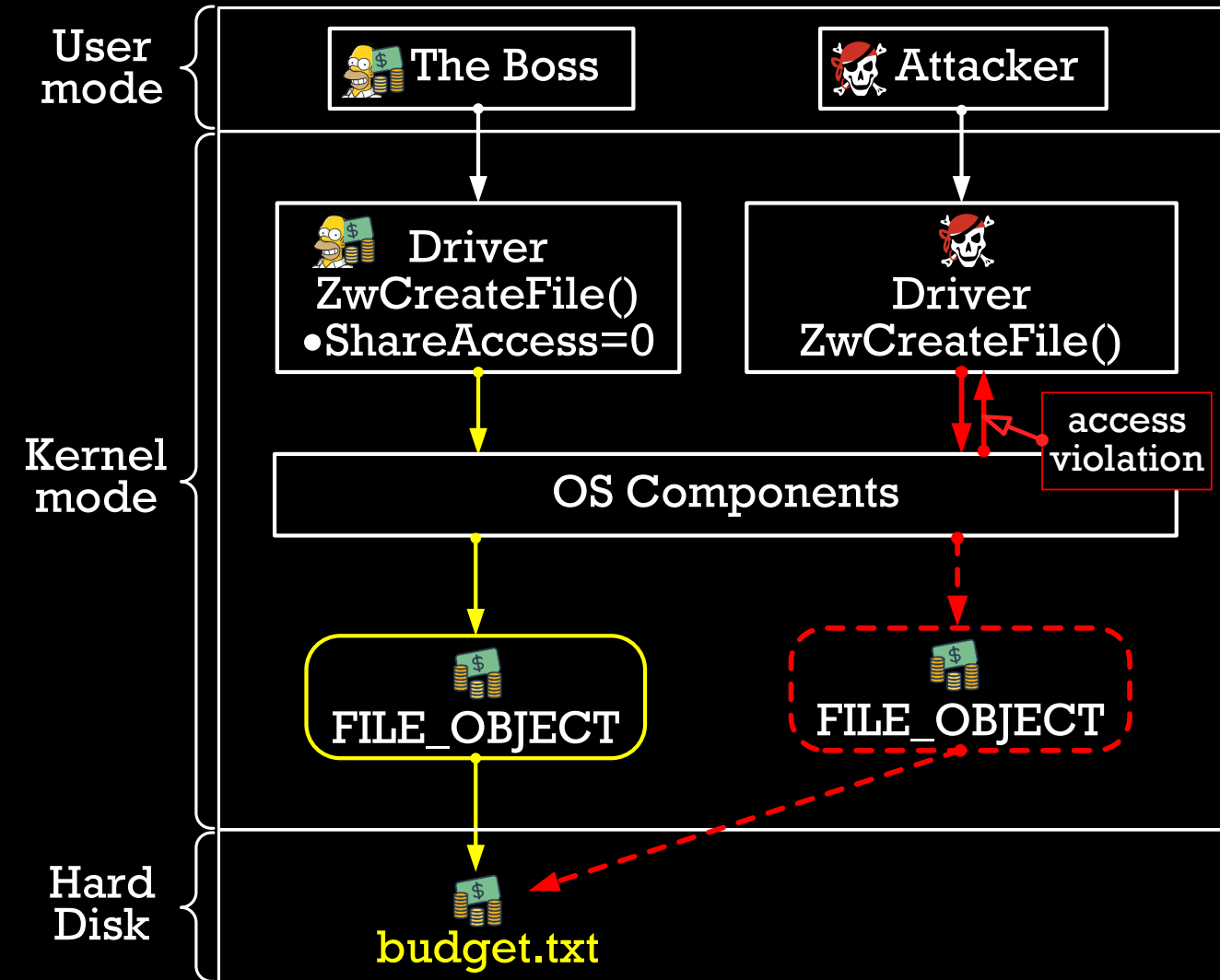
JUST 4 CRUCIAL FIELDS FOR FILES HIJACKING

```
typedef struct _FILE_OBJECT {  
    ...  
    PVPB Vpb;  
    PVOID FsContext;  
    PVOID FsContext2;  
    PSECTION_OBJECT_POINTERS SectionObjectPointer;  
    ...  
} FILE_OBJECT;
```

- The Vpb field points to a mounted Volume Parameter Block (VPB), associated with the target device object.
- FsContext points to the FSRTL_COMMON_FCB_HEADER structure, which has to be allocated by the file driver.
- FsContext2 field refers to the Context Control Block (CBB) associated with the file object
- SectionObjectPointer stores file-mapping and caching-related information for a file stream.

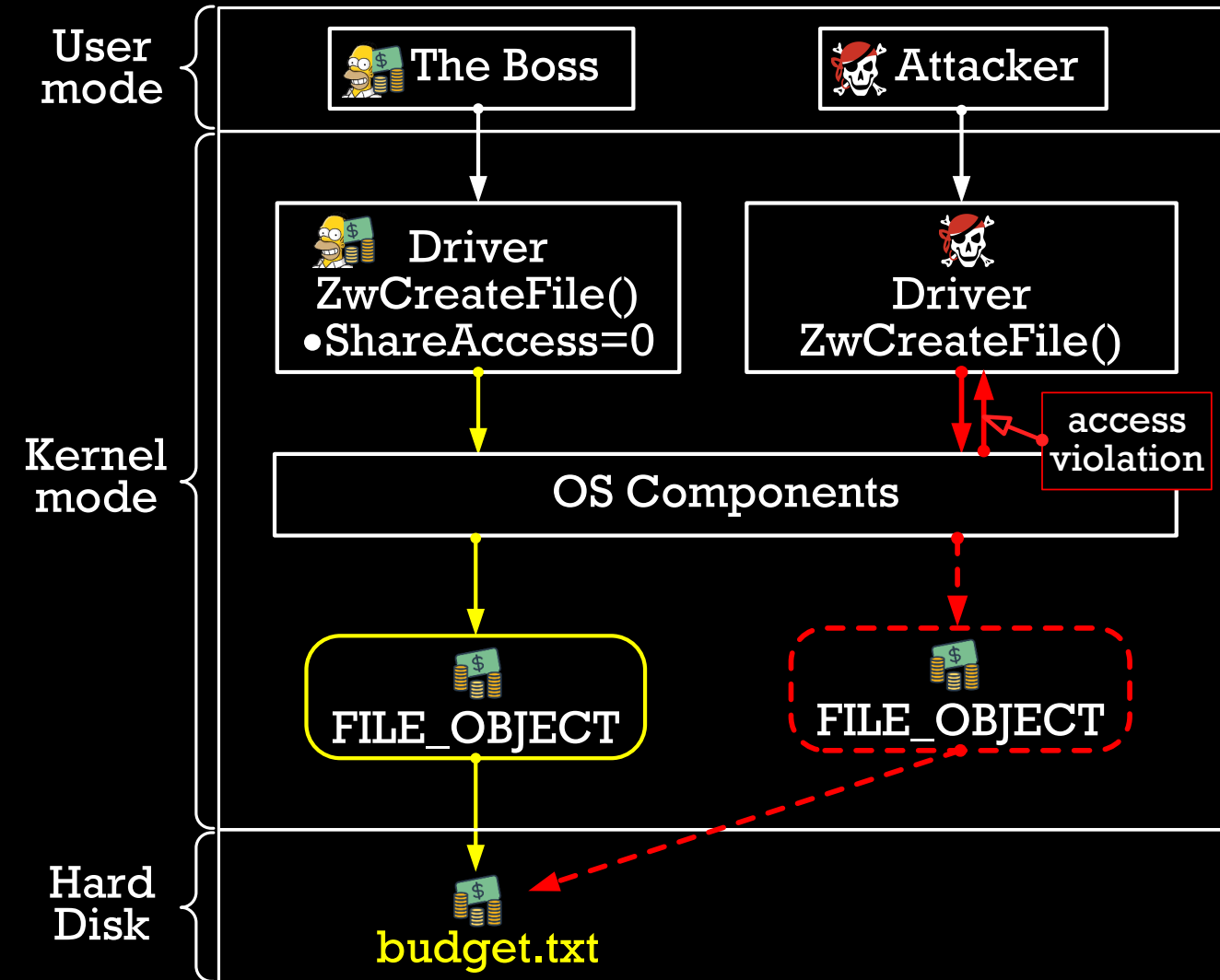
THE ATTACK

Attempt 1: The Legal Access

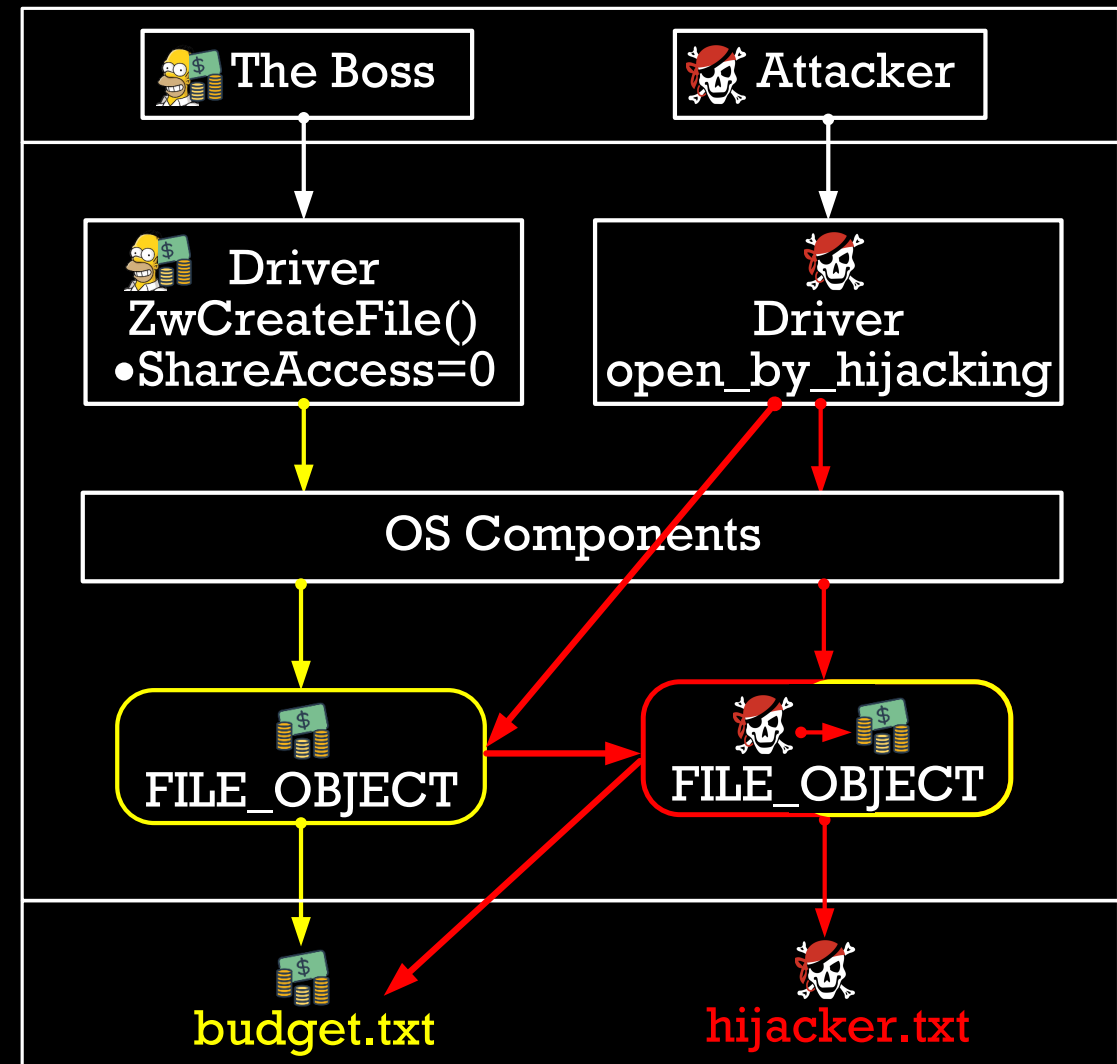


THE ATTACK

Attempt 1: The Legal Access



Attempt 2: The Hijacking Attack



DEMO: THE ATTACK

The online version is here –

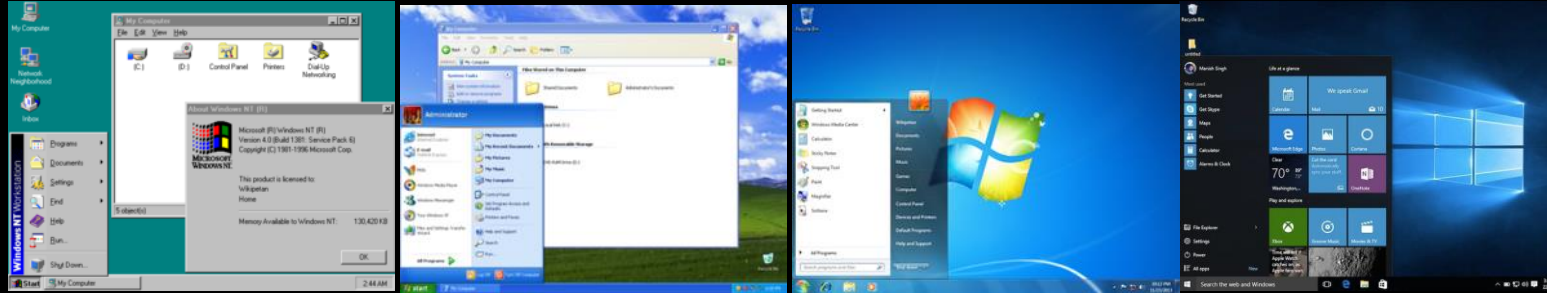
<https://www.youtube.com/watch?v=2mU85RluOSA?vq=hd1080>

- **All Windows OSes since NT 4.0 are vulnerable for FILE_OBJECT hijacking:**



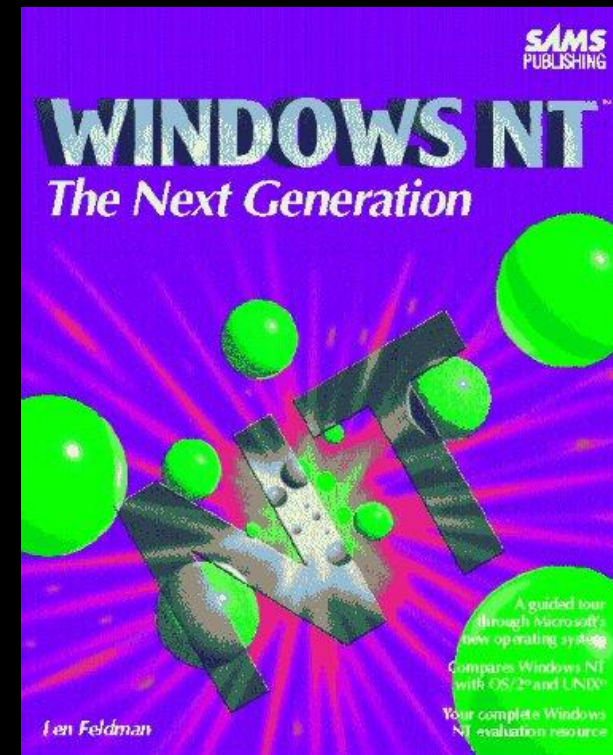
THE ANALYSIS OF THE ATTACK

- All Windows OSes since NT 4.0 are vulnerable for FILE_OBJECT hijacking:



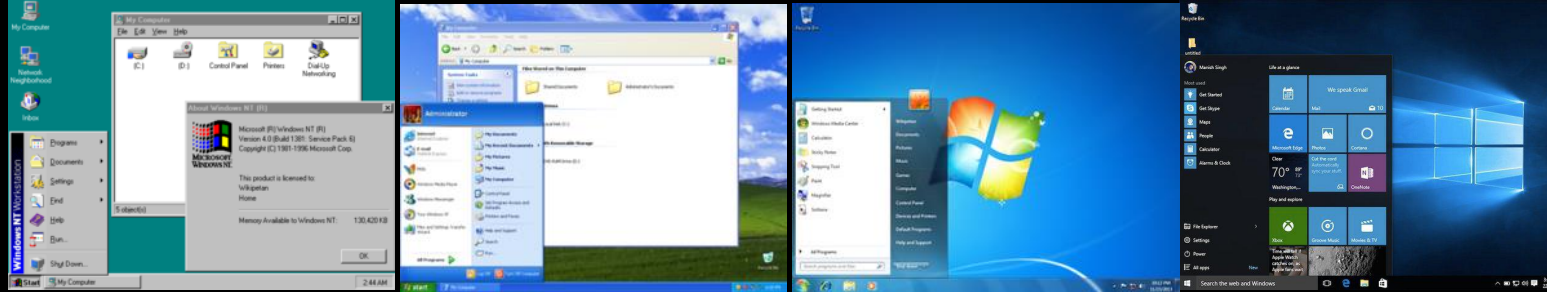
- **1993** - the first mention of Object Manager and Security Reference Monitor

Windows NT: The Next Generation
by Len Feldman, March 1, 1993



THE ANALYSIS OF THE ATTACK

- All Windows OSes since NT 4.0 are vulnerable for FILE_OBJECT hijacking:



- **1993** - the first mention of Object Manager and Security Reference Monitor

- **1965** – the first memory isolation concept Multics* was developed for General Electric 645 mainframe. Multics joined to the ARPANet and gave rise to the Unix.

*DOI: <http://dx.doi.org/10.1145/1463891.1463912>

Two Fathers of Multics

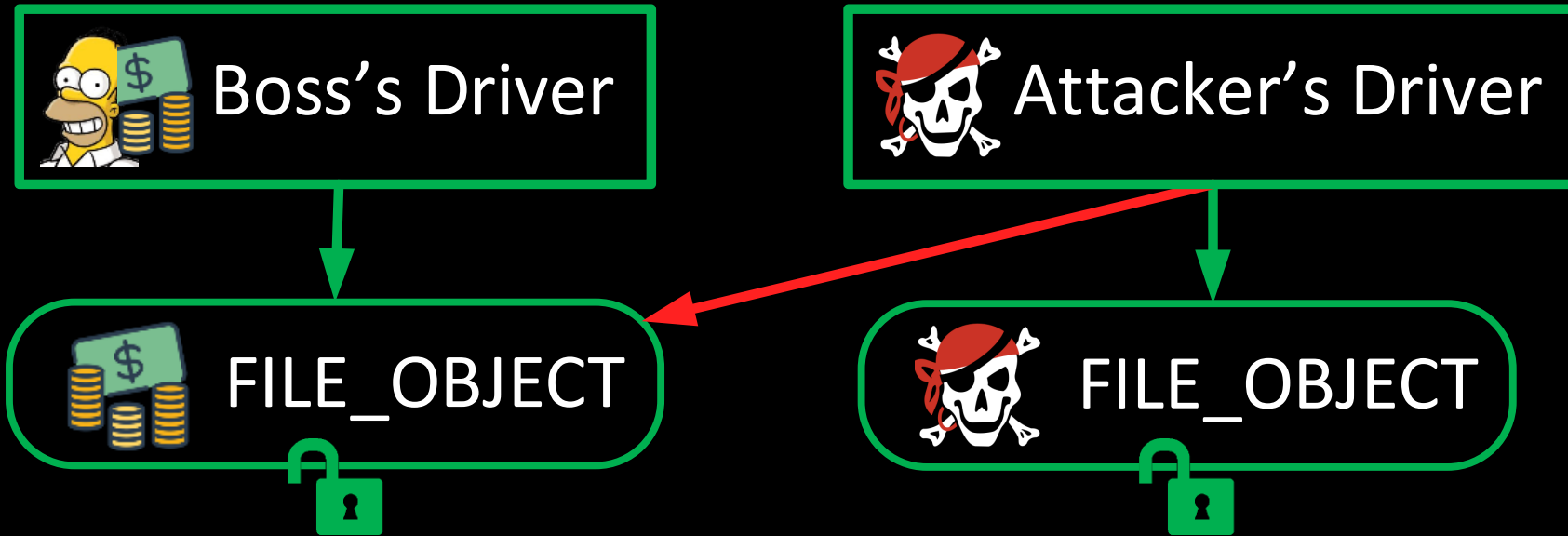


Fernando
Corbato

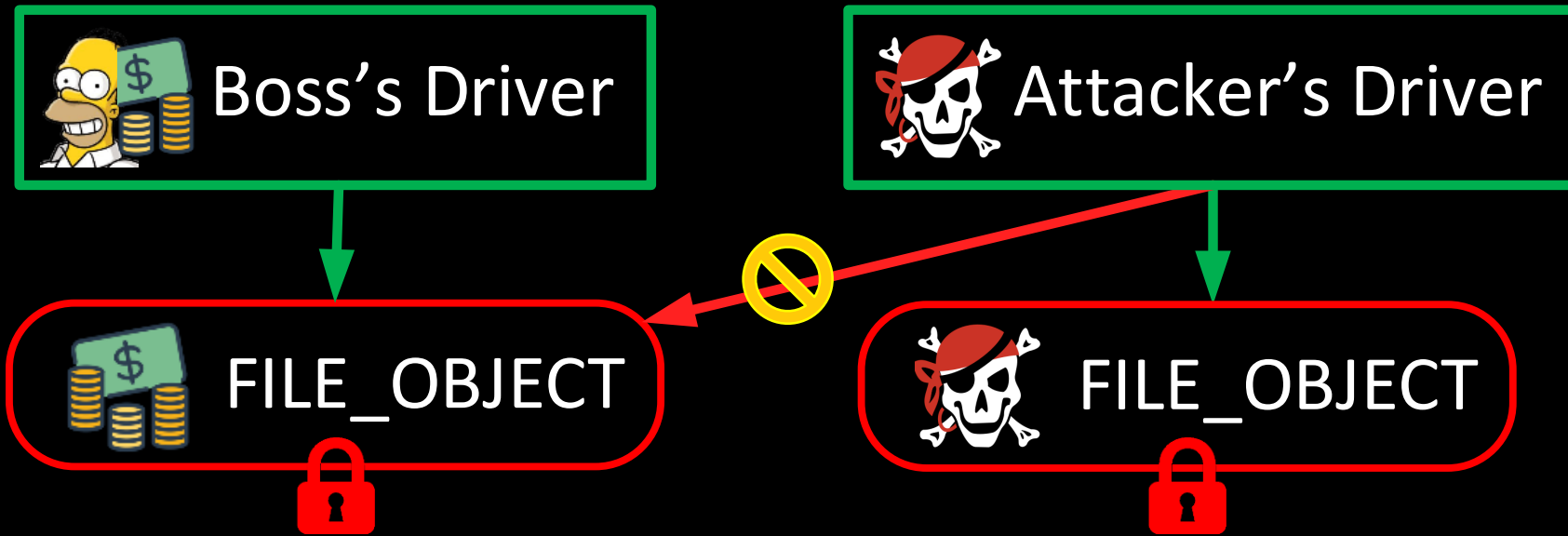


Victor
Vyssotsky

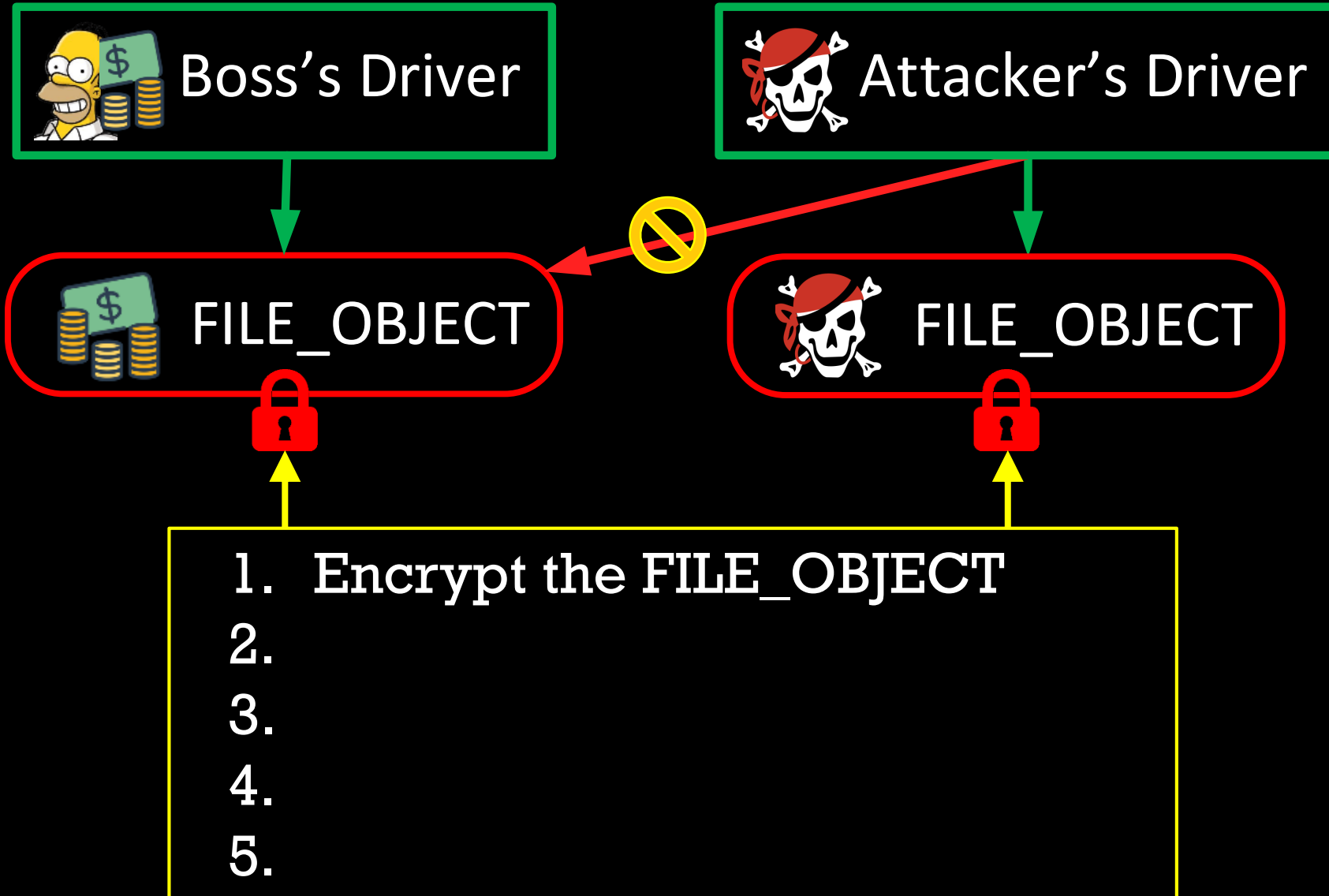
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



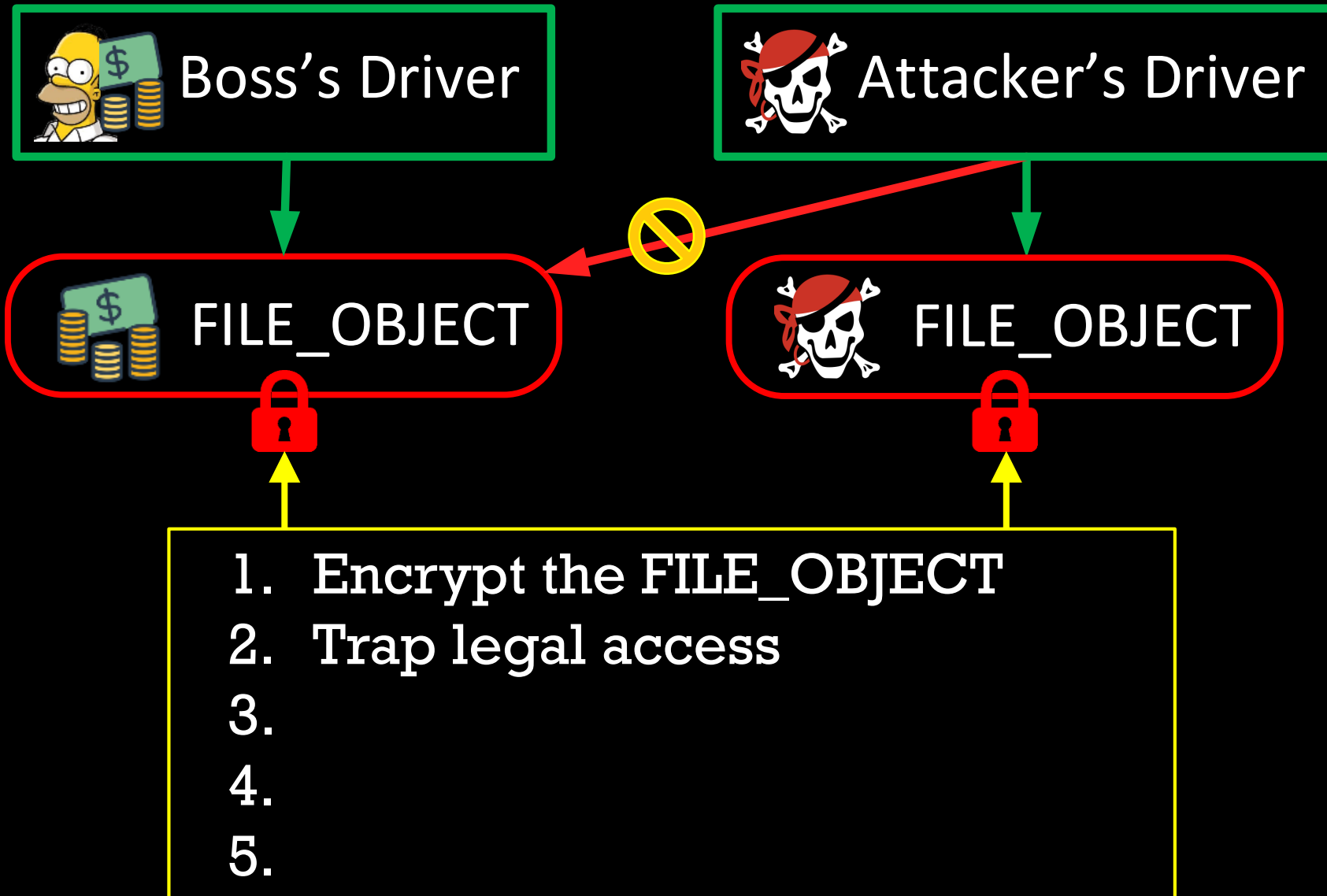
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



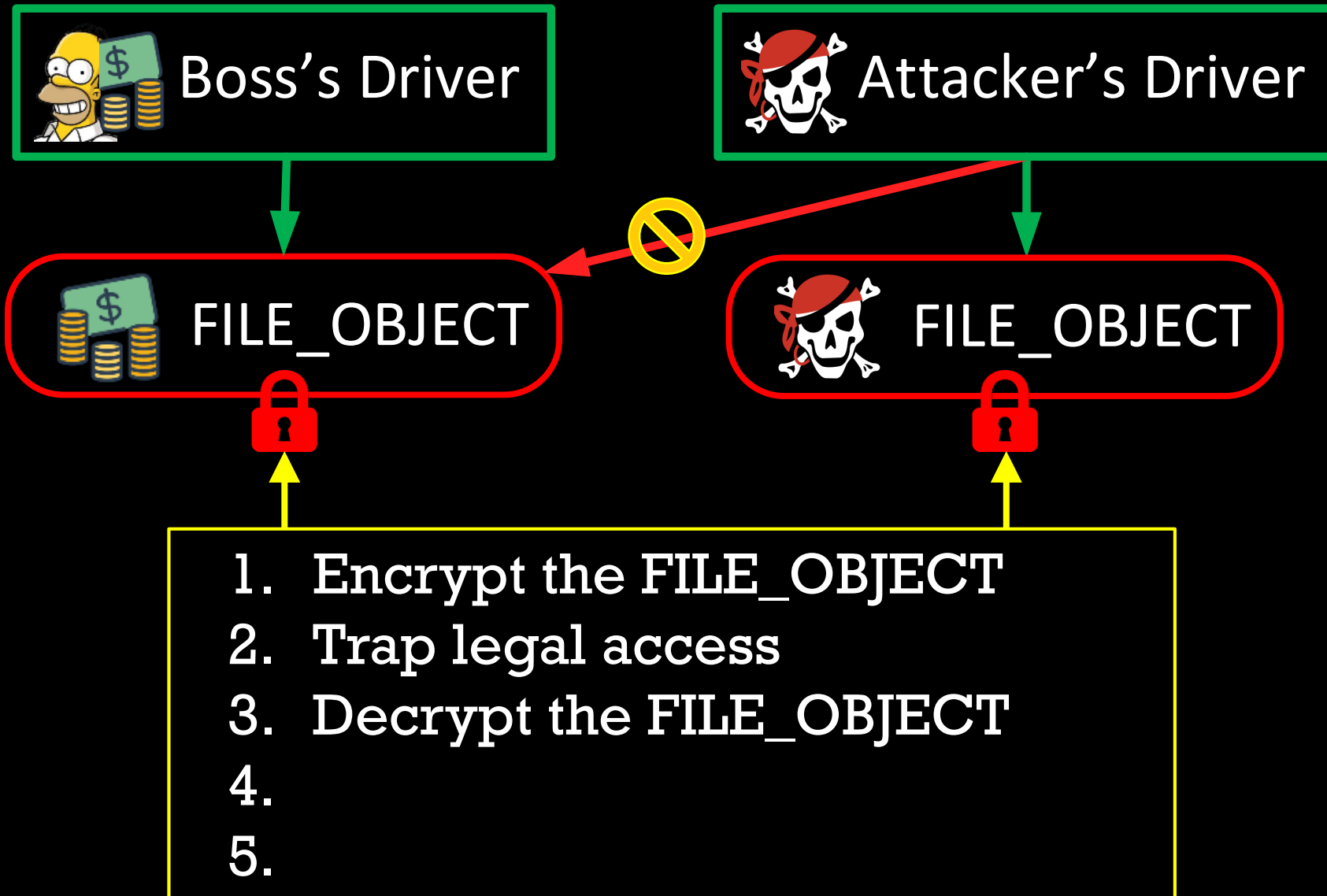
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



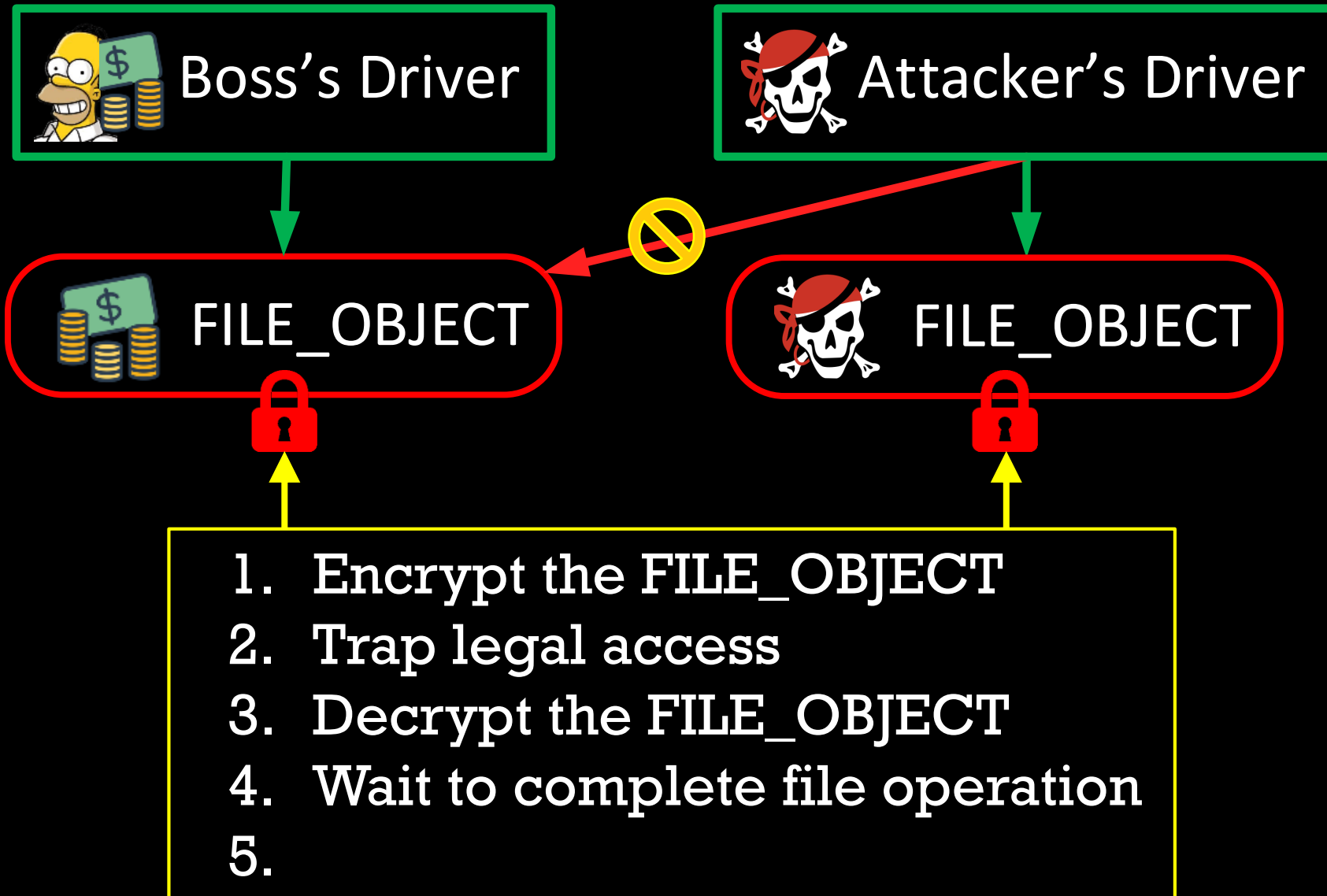
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



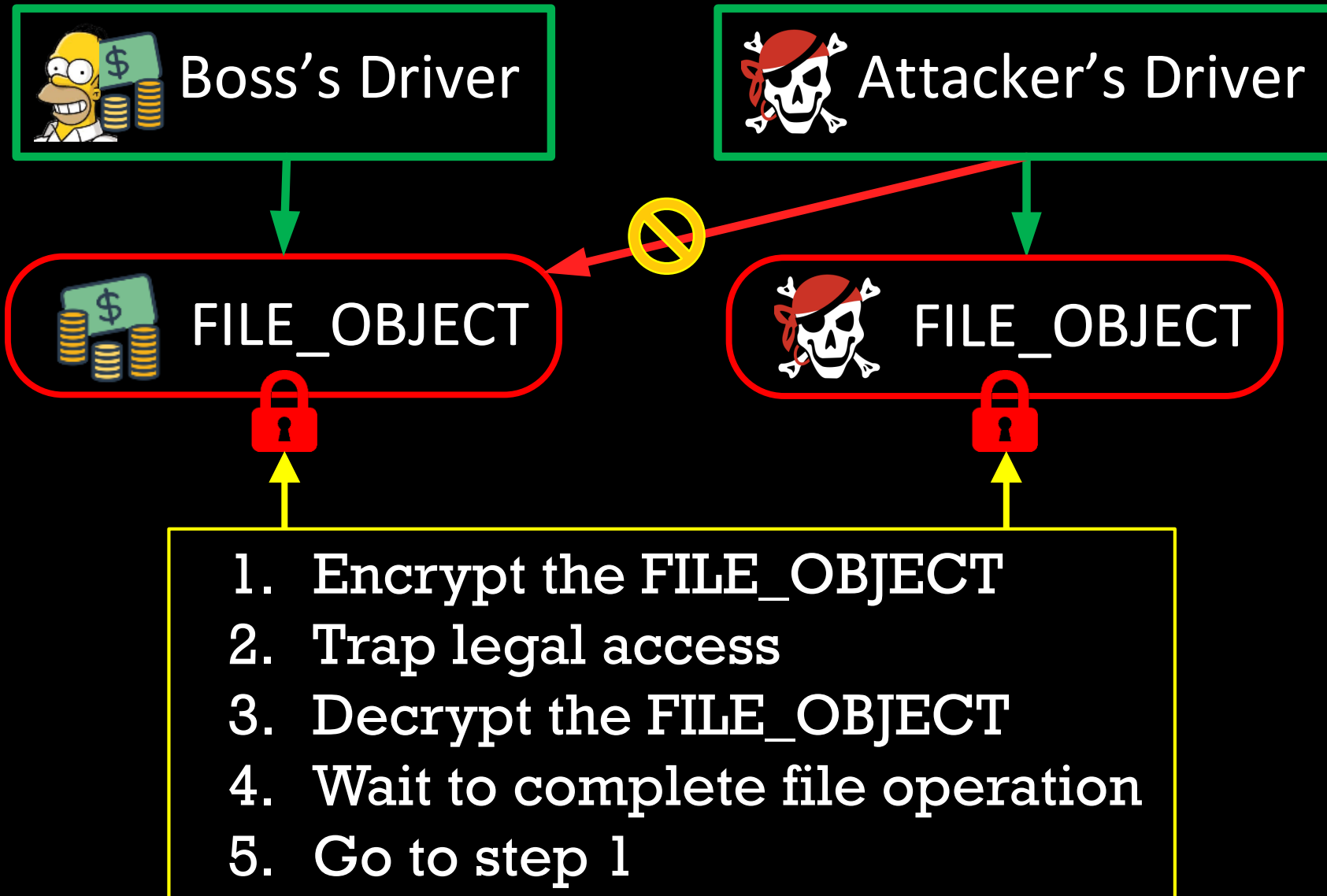
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



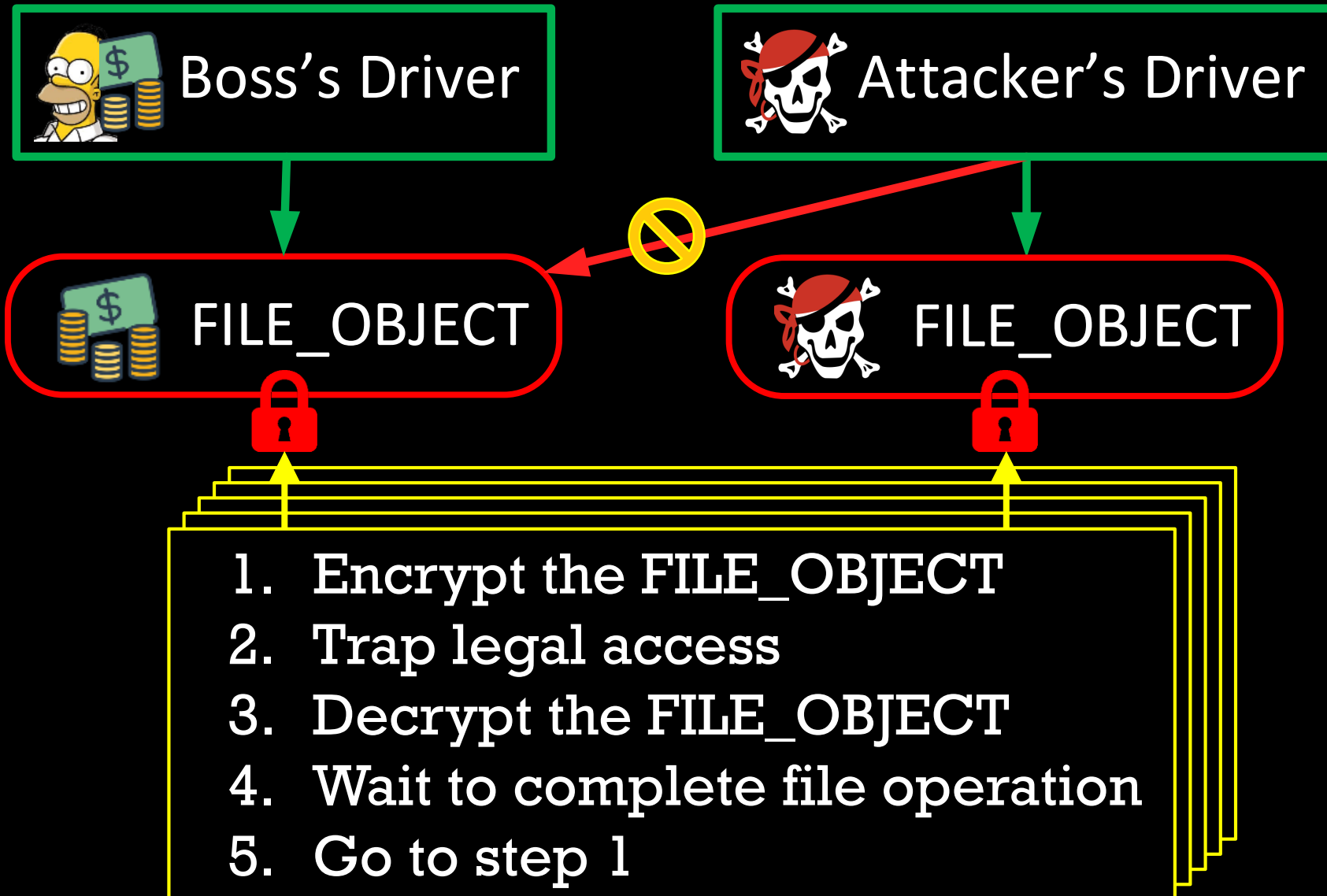
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



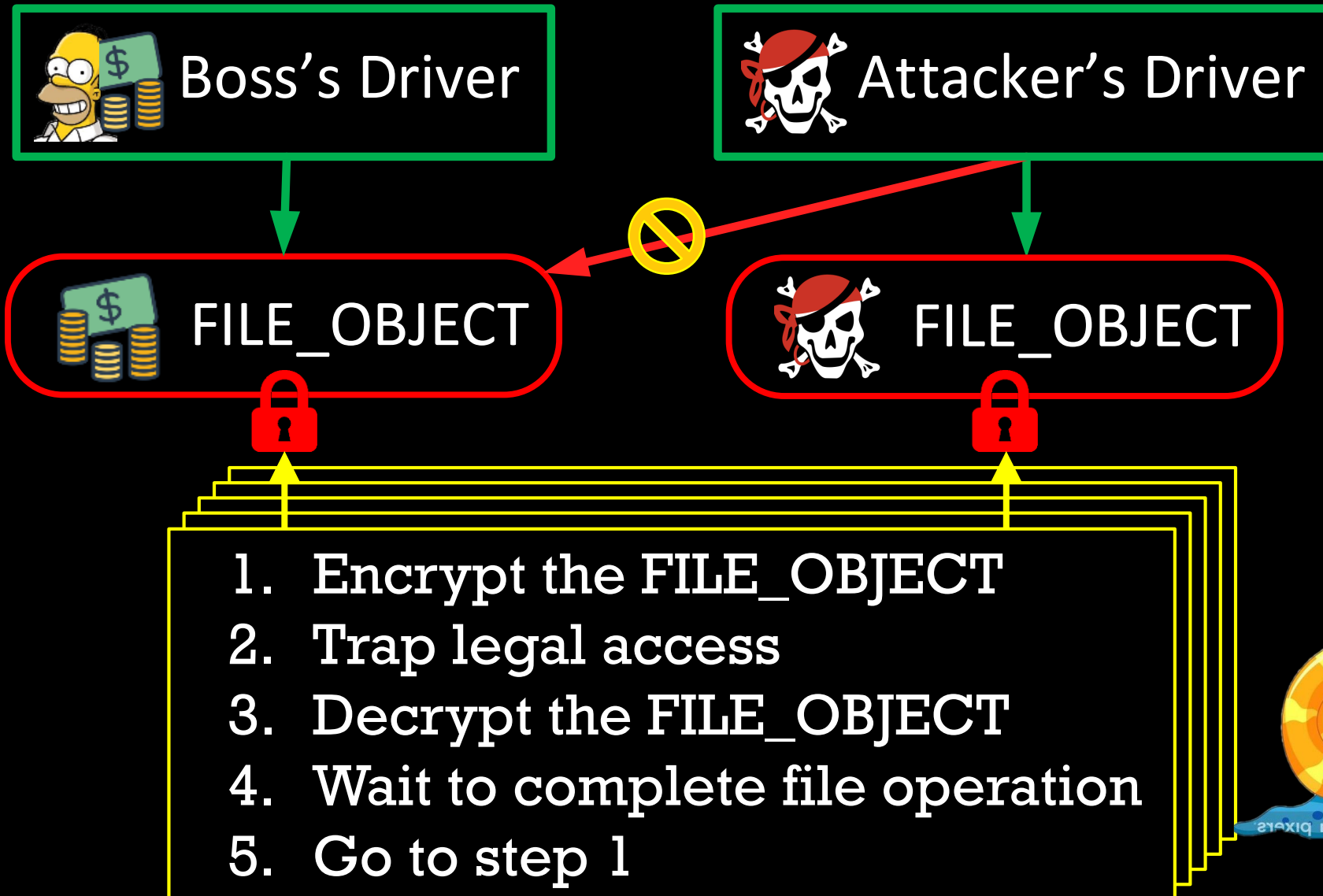
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



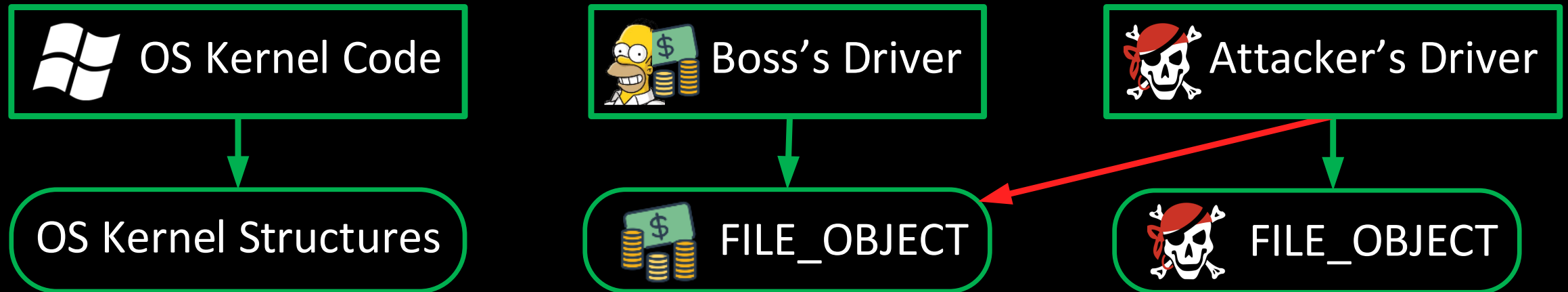
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



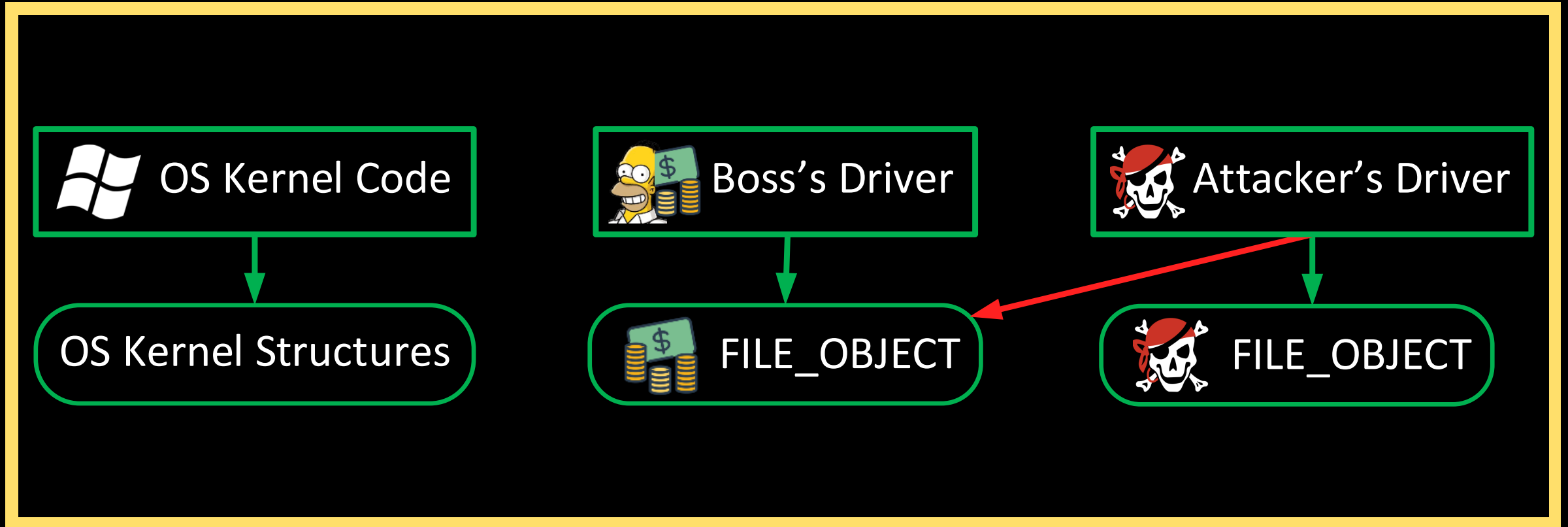
THE FILE_OBJECT PROTECTION VIA ENCRYPTION



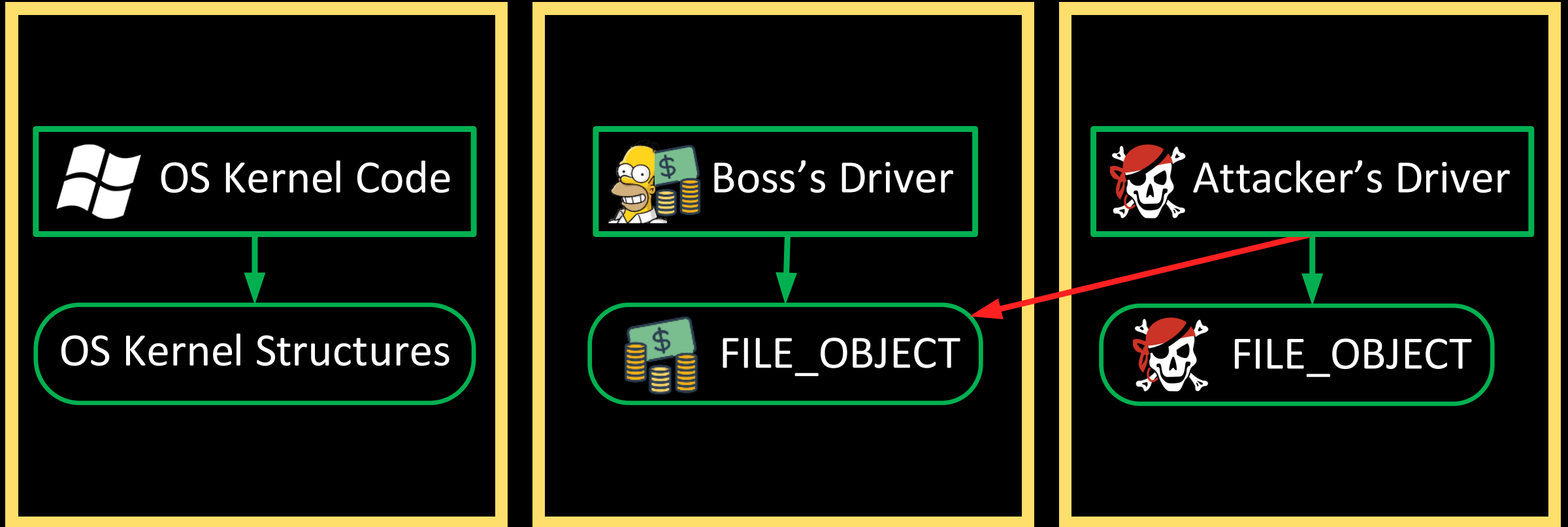
WINDOWS KERNEL MEMORY



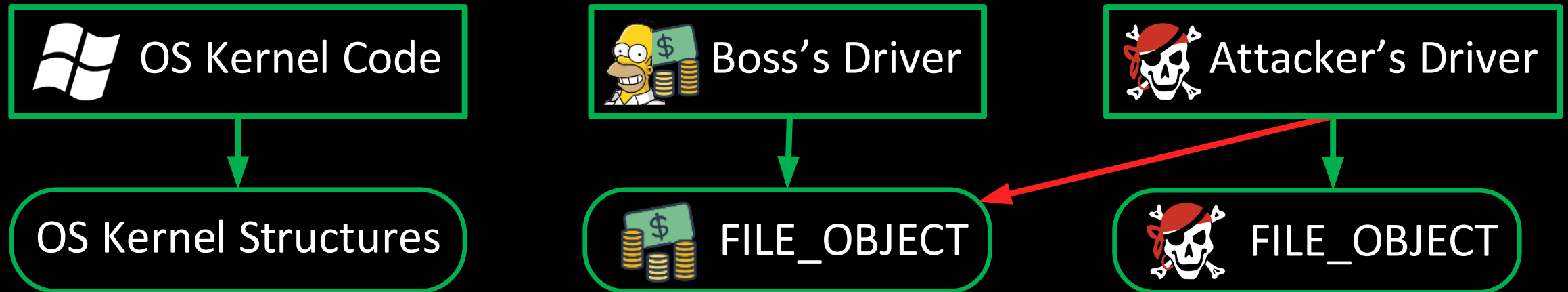
WINDOWS KERNEL MEMORY



WINDOWS KERNEL MEMORY



WINDOWS KERNEL MEMORY

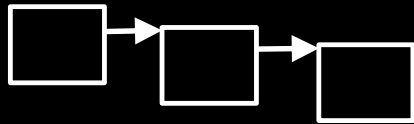


PROCESSING MEMORY ACCESS: EPT FEATURE

VT-x without EPT

Guest OS

Paging
structures

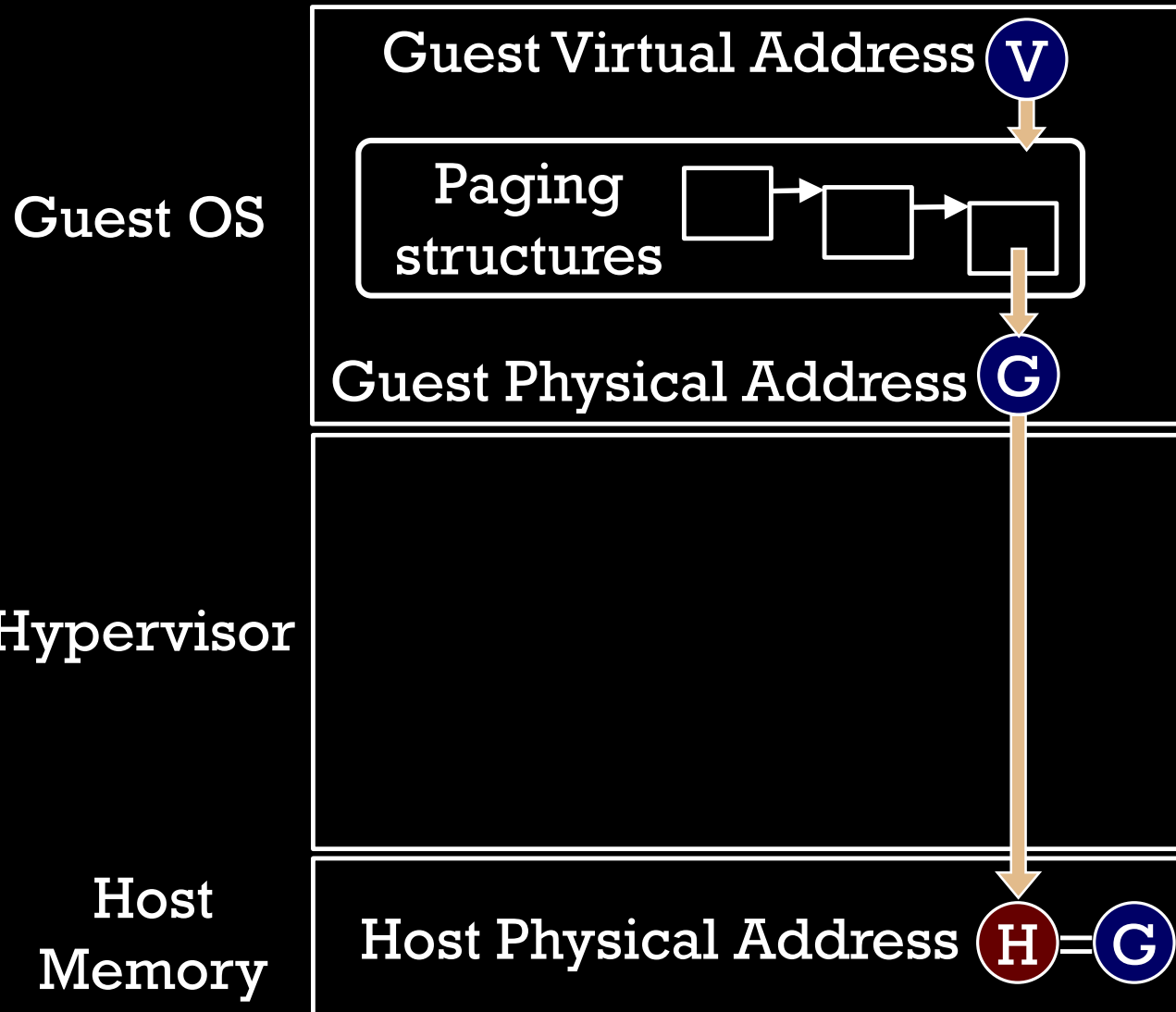


Hypervisor

Host
Memory

PROCESSING MEMORY ACCESS: EPT FEATURE

VT-x without EPT



PROCESSING MEMORY ACCESS: EPT FEATURE

VT-x without EPT

VT-x with EPT

Guest OS

Guest Virtual Address **V**

Paging
structures

Guest Physical Address **G**

Hypervisor

Host
Memory

Host Physical Address **H** = **G**

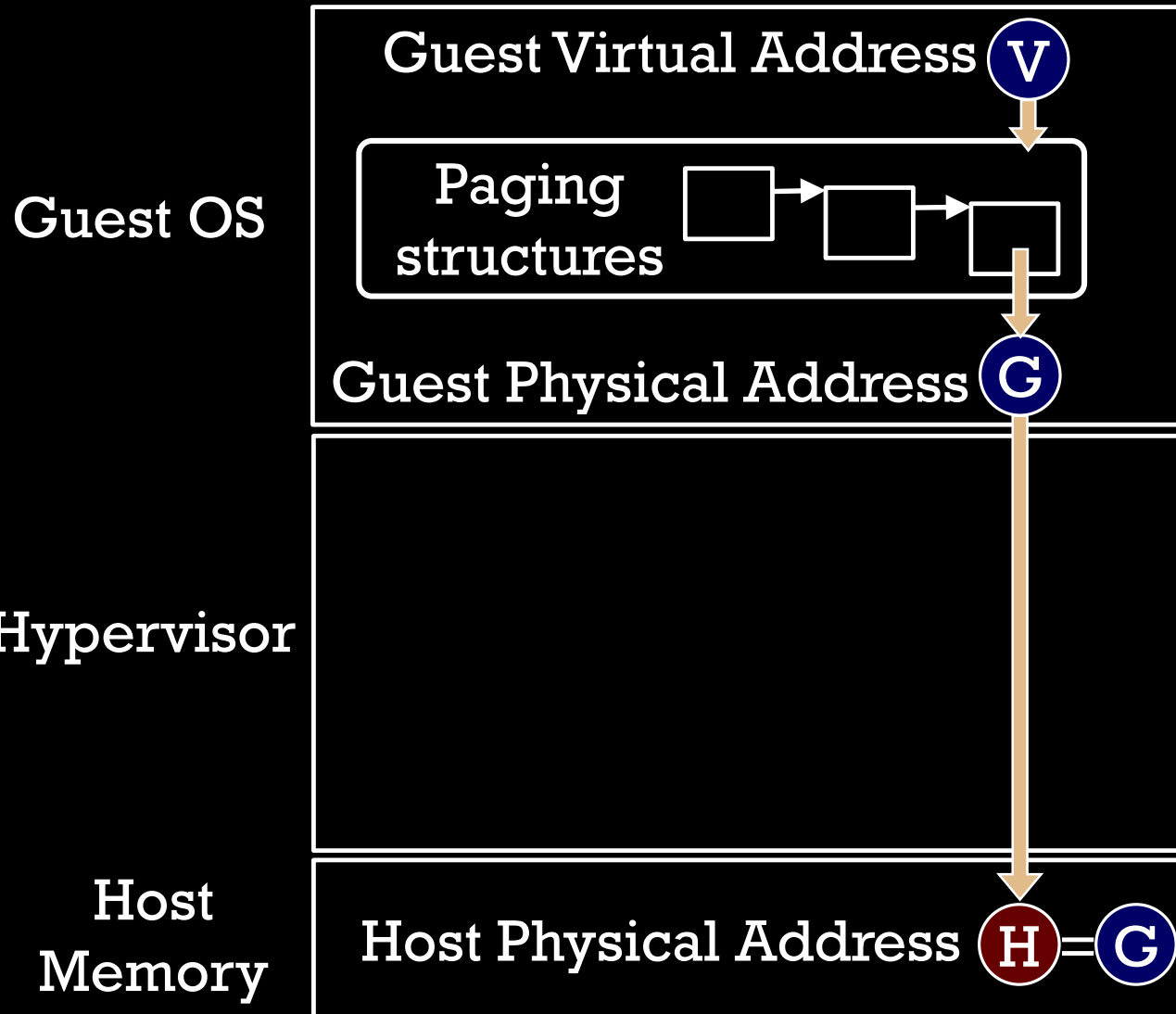
Paging
structures

EPT Paging structures

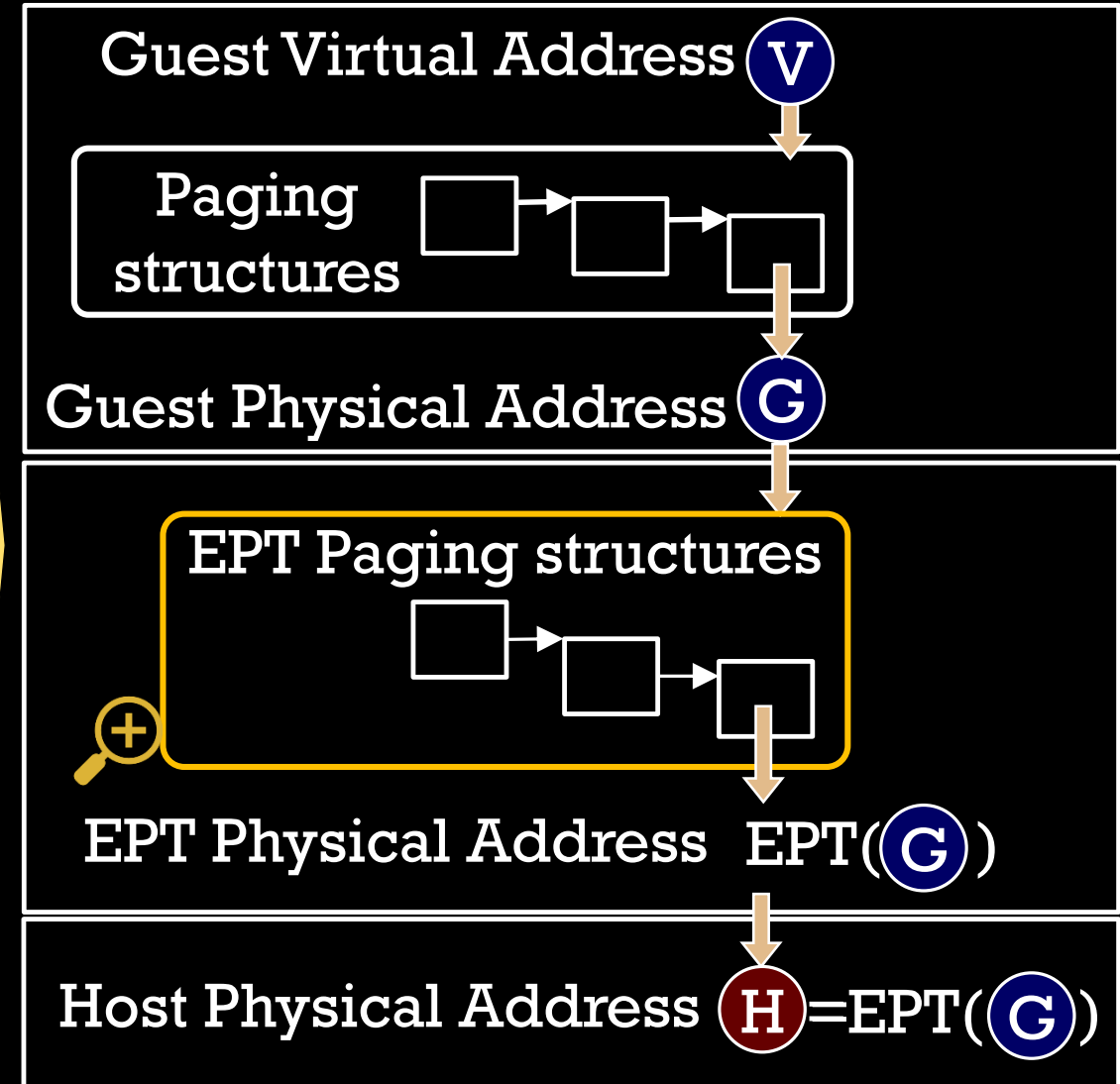


PROCESSING MEMORY ACCESS: EPT FEATURE

VT-x without EPT



VT-x with EPT



EPT PAGING STRUCTURES

EPT Paging structures

EPT Tables



```
graph LR; subgraph EPT_Paging_Structures [EPT Paging structures]; subgraph EPT_Tables [EPT Tables]; direction TB; T1[EPT Table 1]; T2[EPT Table 2]; T3[EPT Table 3]; end; subgraph EPT_Page_Table_Entries [EPT Page Table Entries]; direction TB; E1[EPT Entry for the Page A]; E2[EPT Entry for the Page B]; E3[...]; E4[EPT Entry for the Page Z]; end; T1 --> E1; end
```

The diagram illustrates the EPT paging structures. On the left, a stack of three boxes represents 'EPT Tables'. An arrow points from the top table to a larger box on the right labeled 'EPT Page Table Entries'. This box contains a list of entries: 'EPT Entry for the Page A', 'EPT Entry for the Page B', an ellipsis '...', and 'EPT Entry for the Page Z'.

EPT Page Table Entries

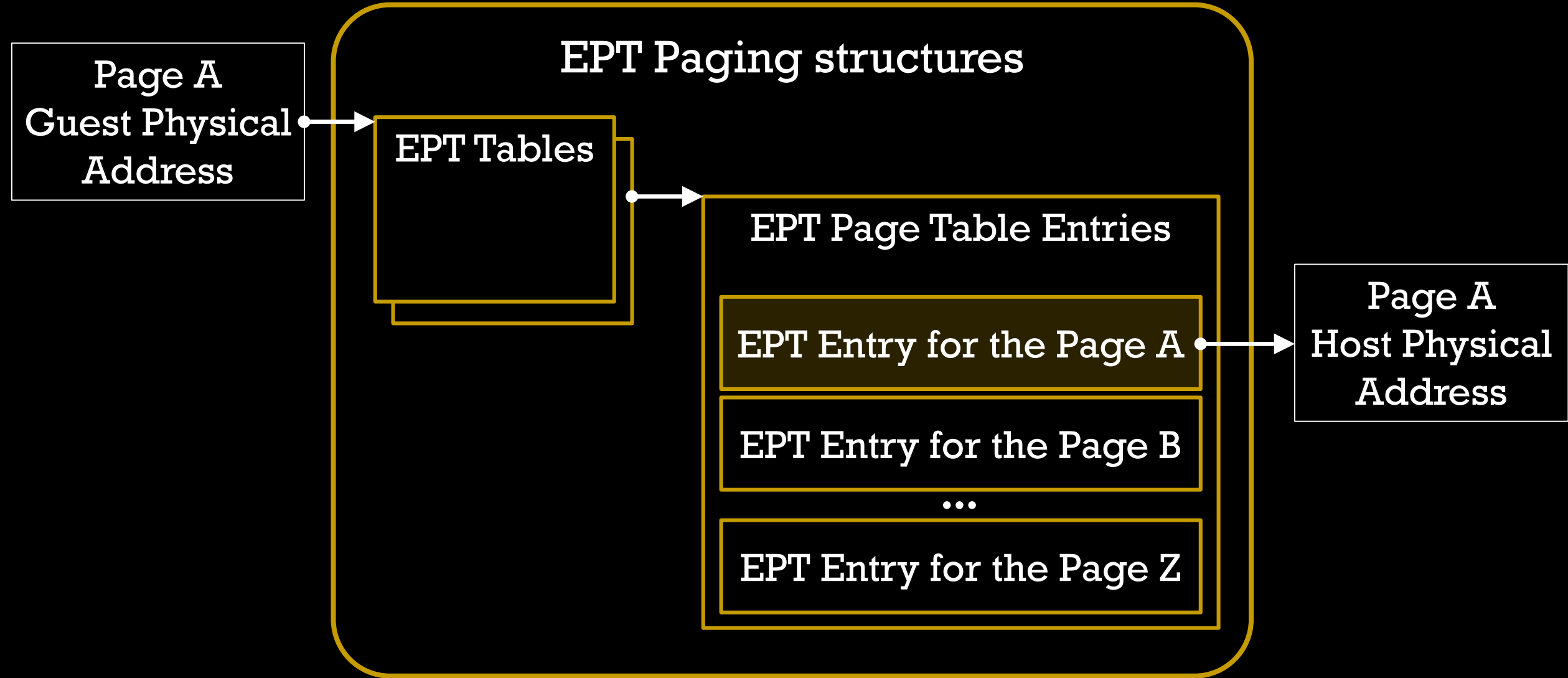
EPT Entry for the Page A

EPT Entry for the Page B

...

EPT Entry for the Page Z

EPT PAGING STRUCTURES



EPT MAIN FEATURES

1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:

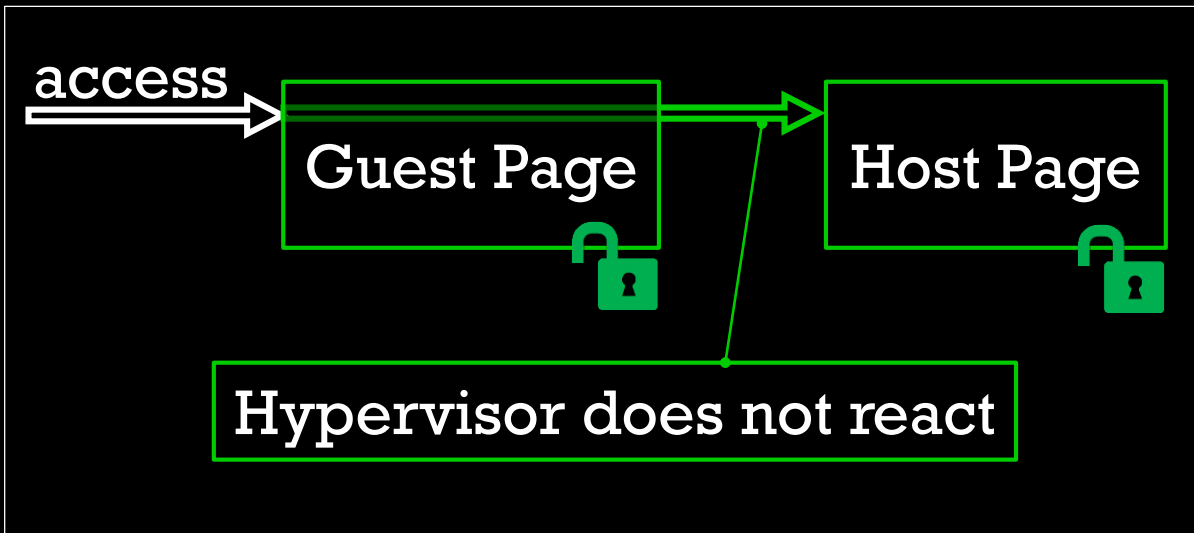


2.

3.

EPT MAIN FEATURES

1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:

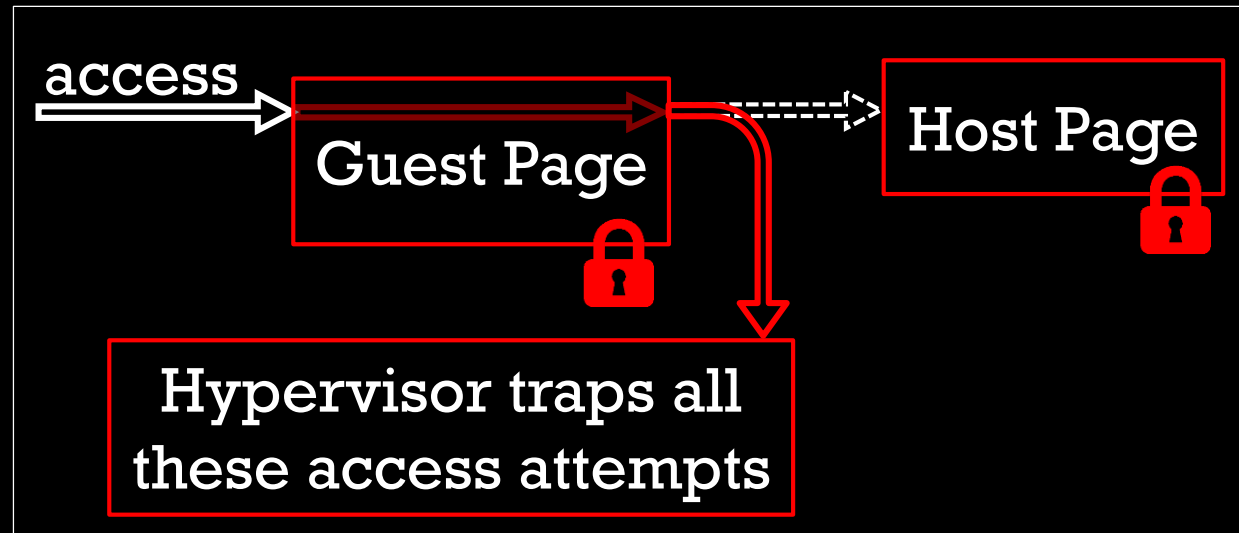
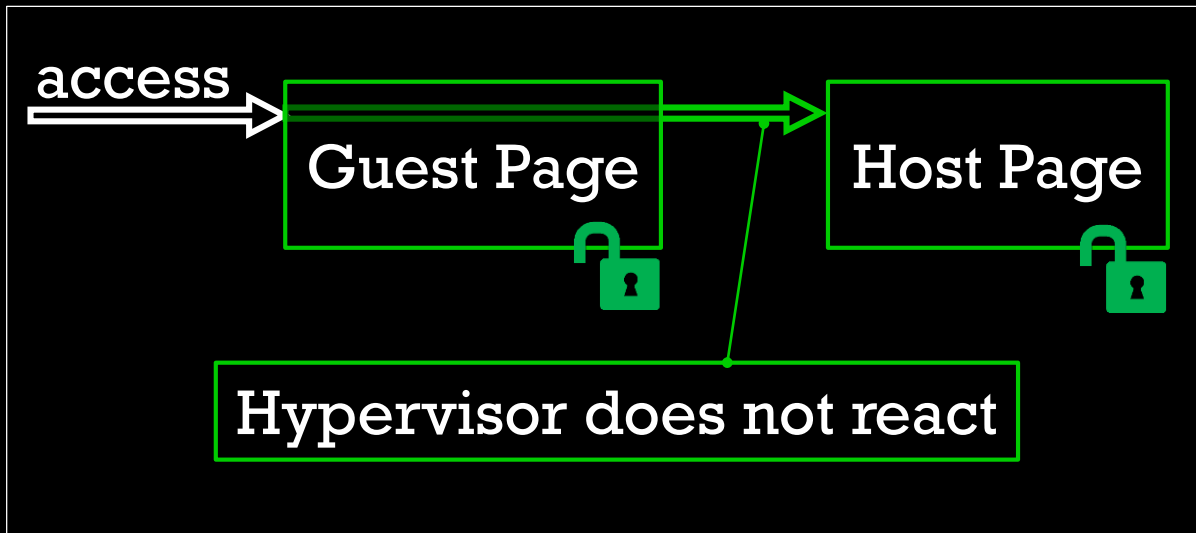


2.

3.

EPT MAIN FEATURES

1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:

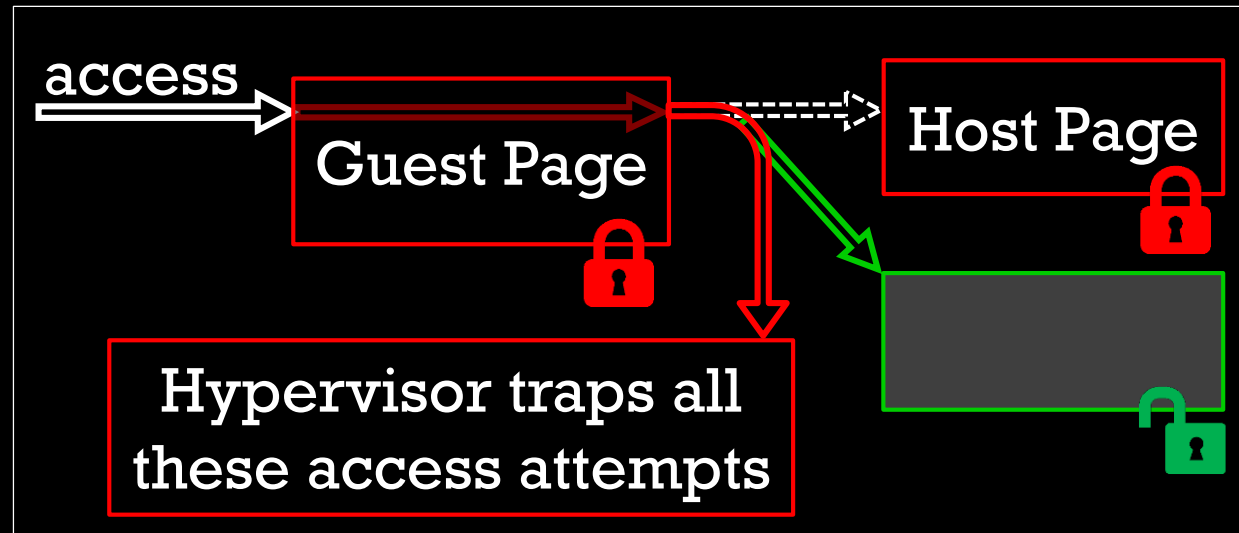
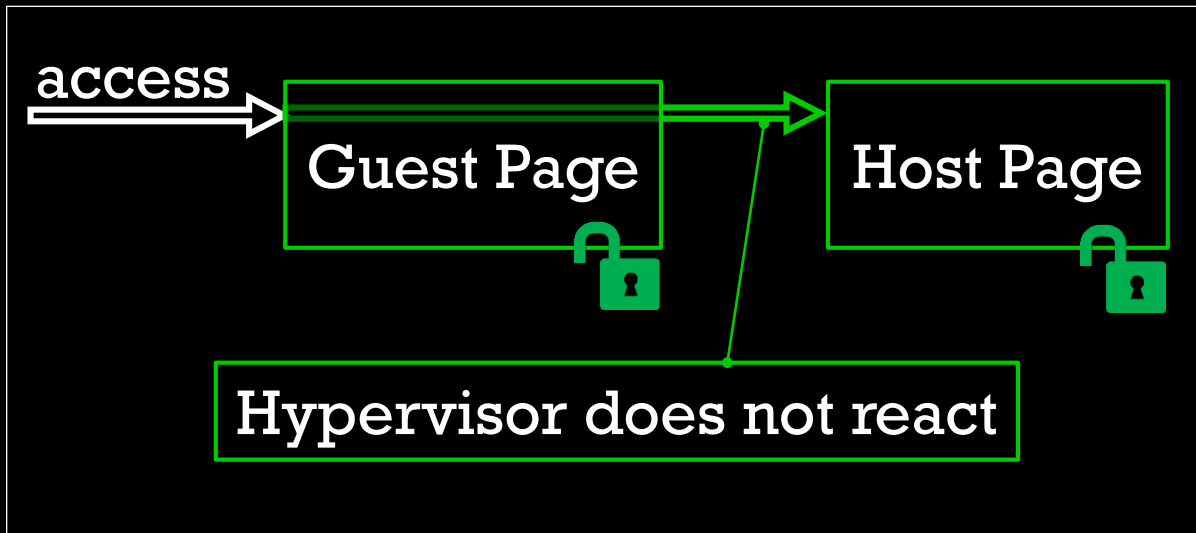


2.

3.

EPT MAIN FEATURES

1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:

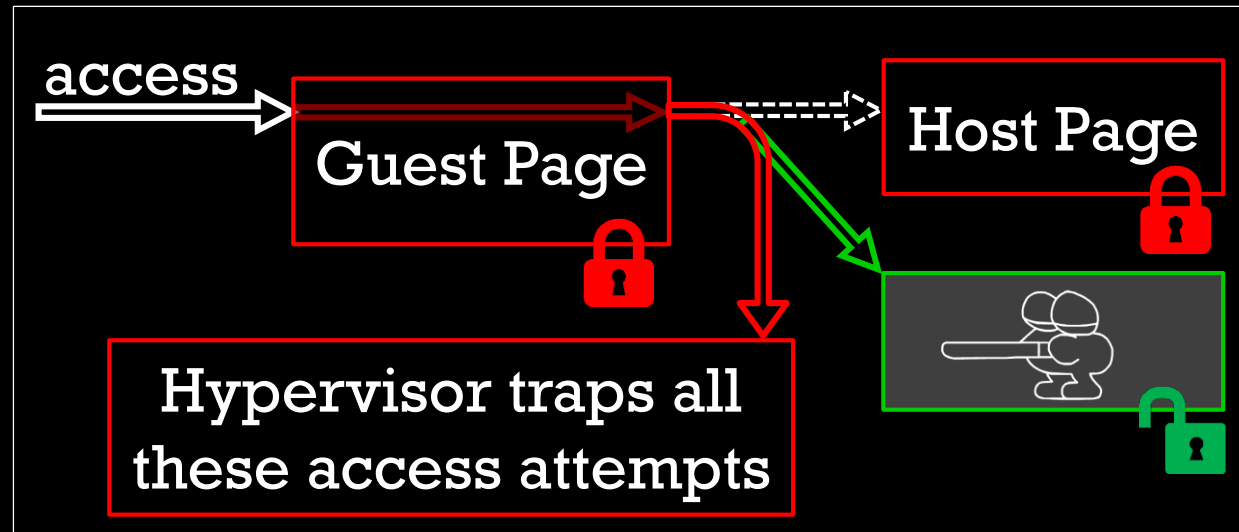
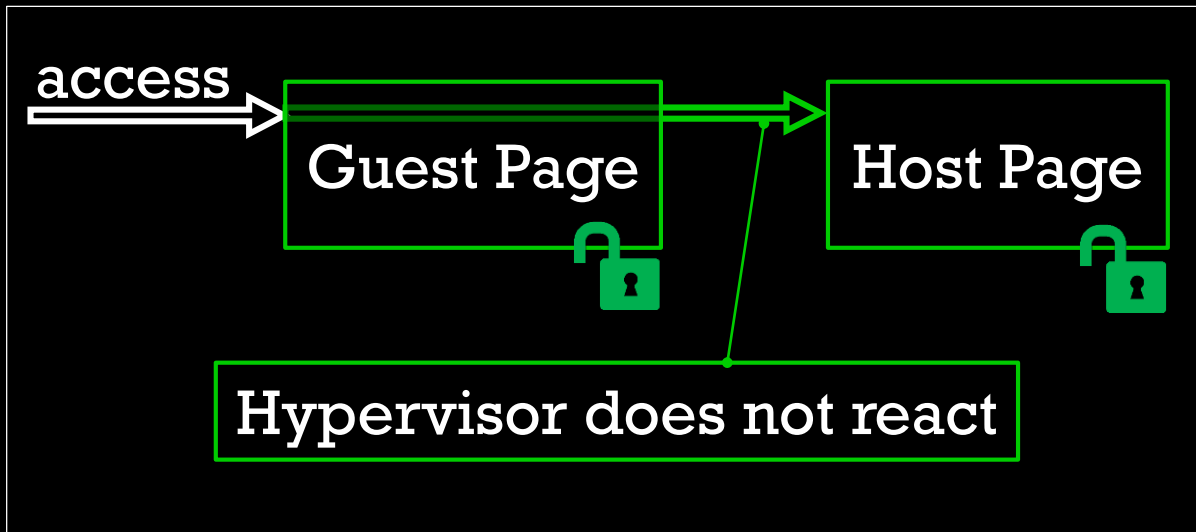


2.

3.

EPT MAIN FEATURES

1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:

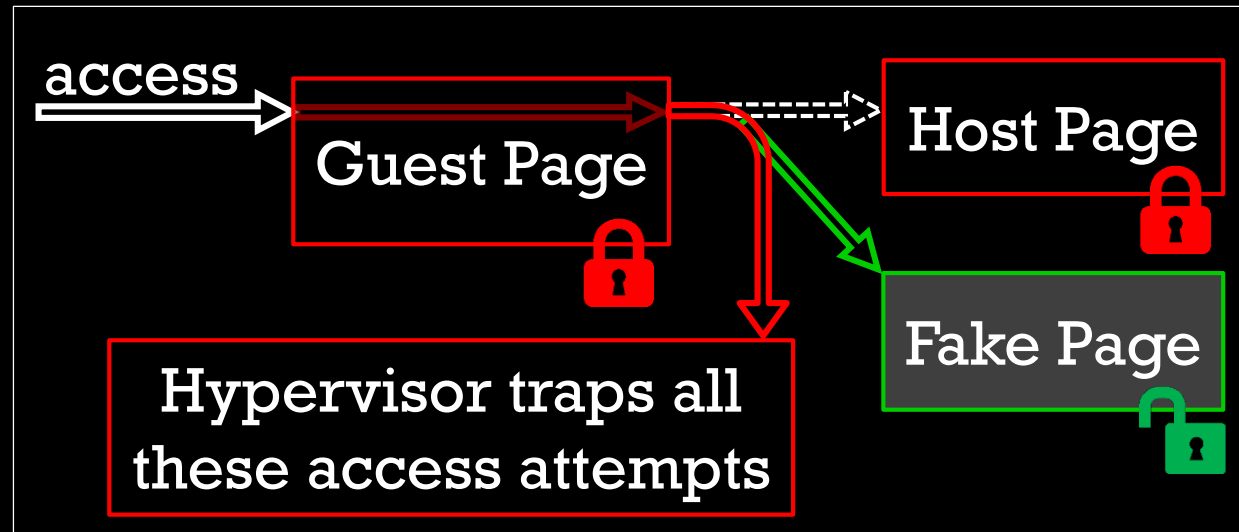
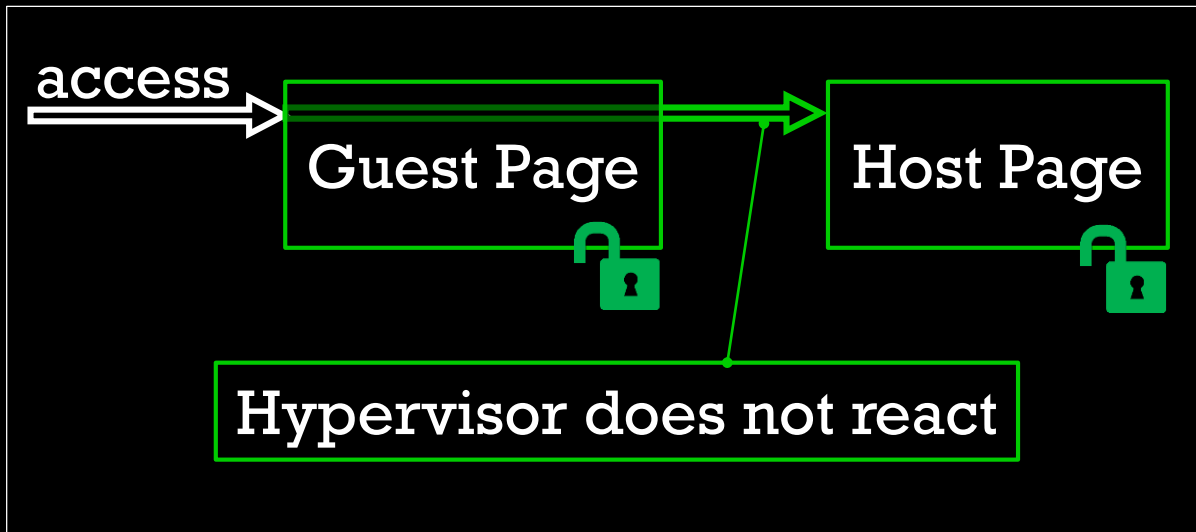


2.

3.

EPT MAIN FEATURES

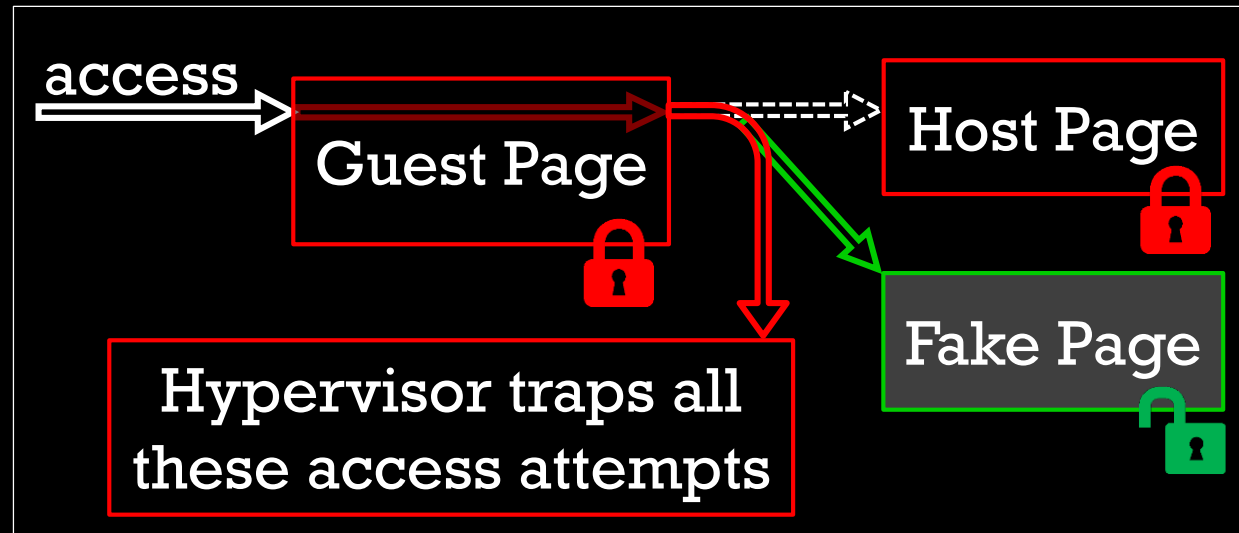
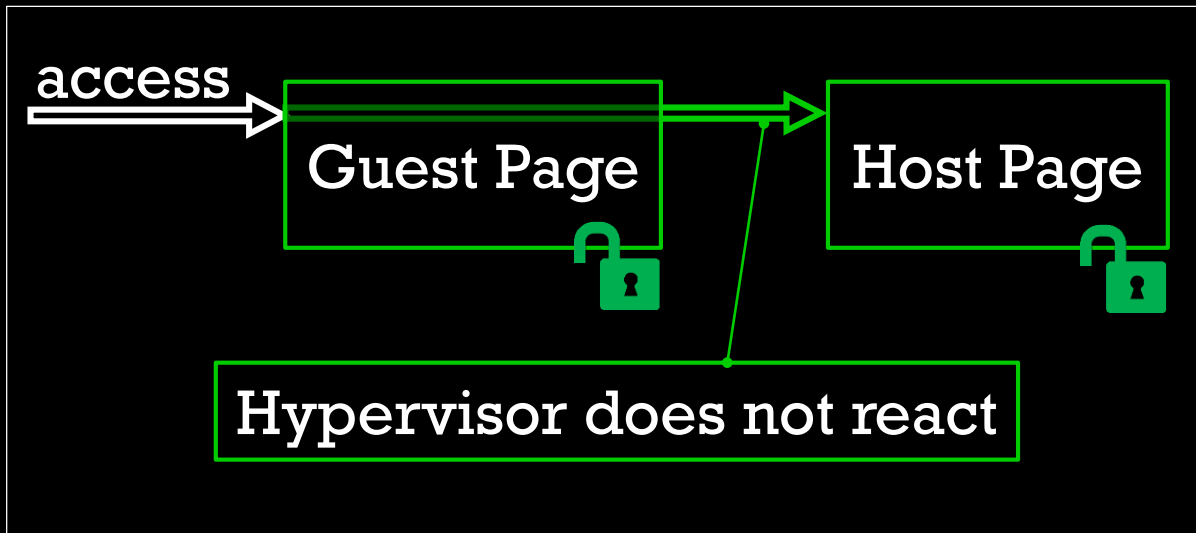
1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:



2. EPT memory settings can be updated in the real time
- 3.

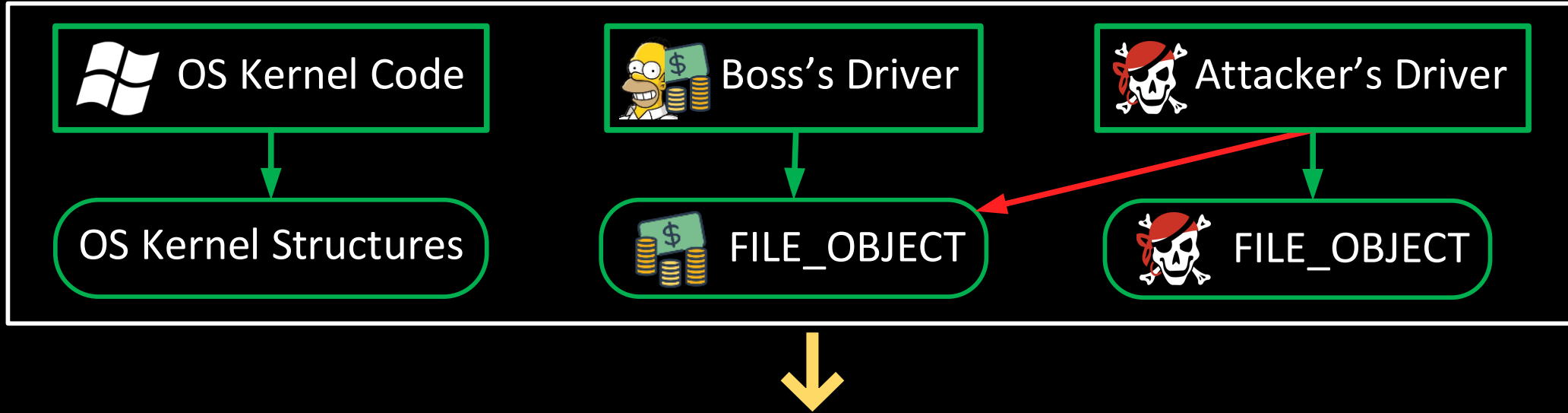
EPT MAIN FEATURES

1. Using EPT we can trap read/write/execute access attempts and redirect them from the secret page to the fake one:

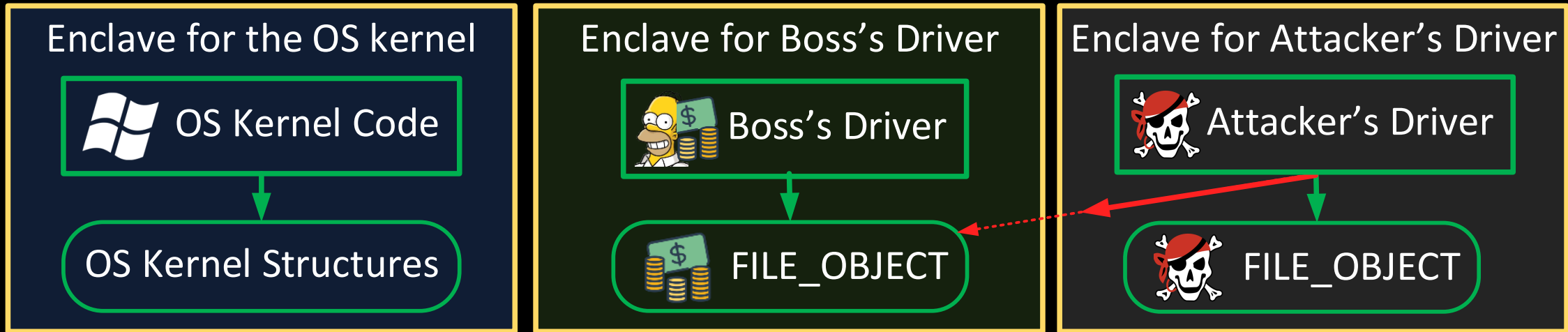
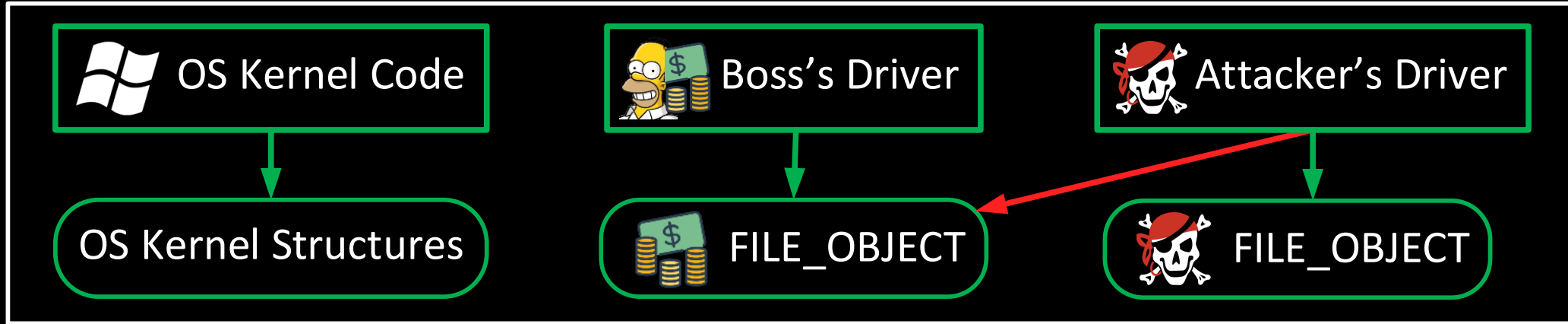


2. EPT memory settings can be updated in the real time
3. We can dynamically allocate several EPTs with different memory setting and switch between them in the real time

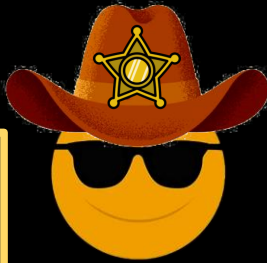
WINDOWS KERNEL MEMORY



WINDOWS KERNEL MEMORY

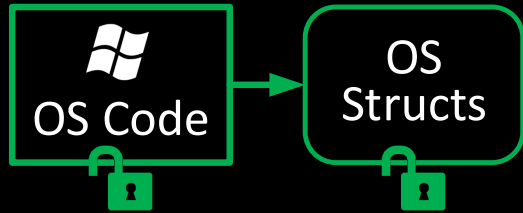


MemoryRanger



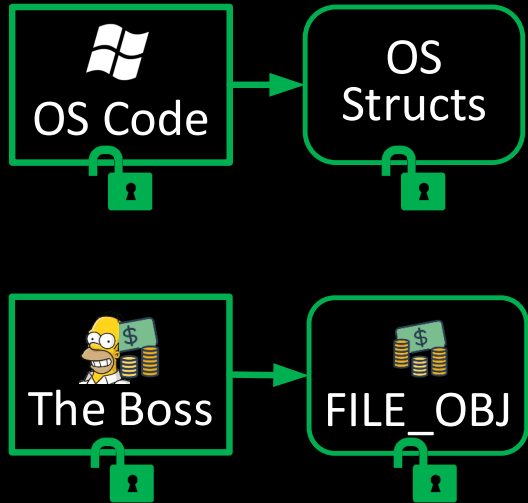
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING

Current Situation



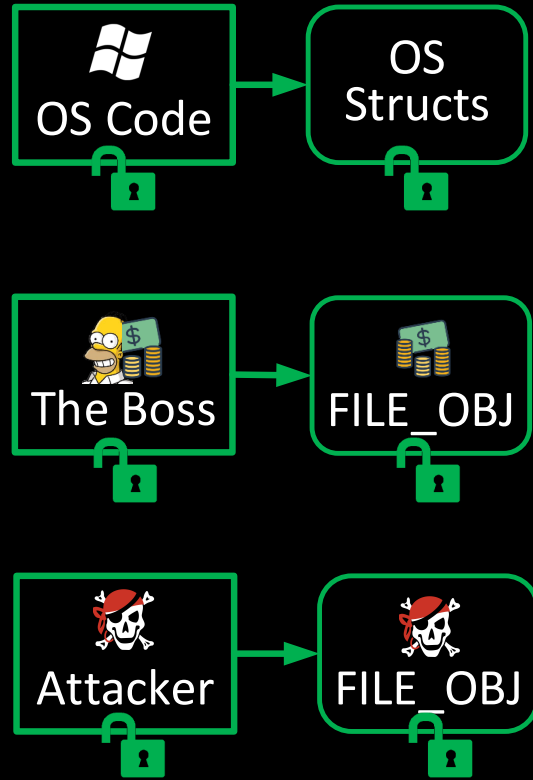
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING

Current Situation



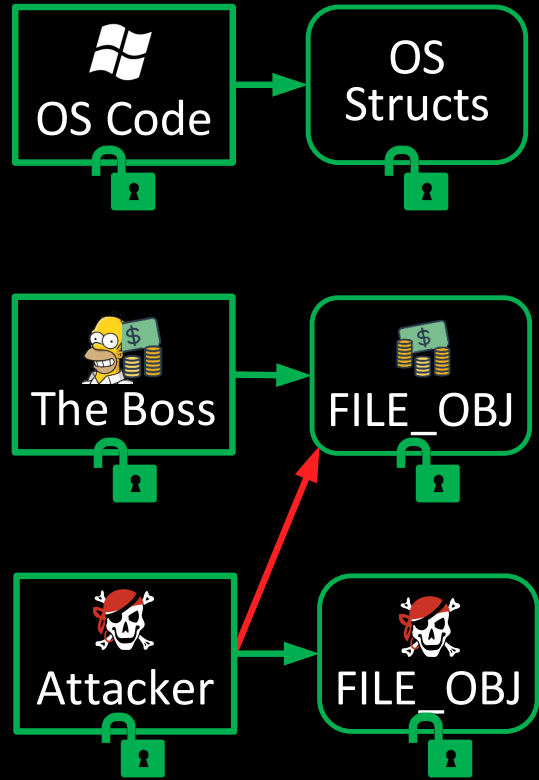
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING

Current Situation



MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING

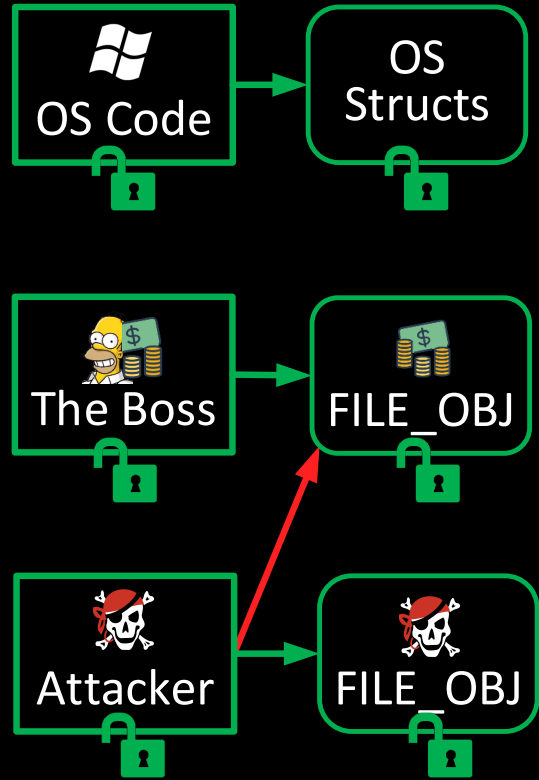
Current Situation



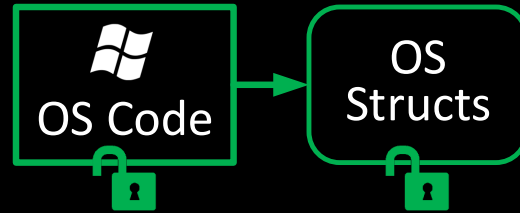
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



Current Situation

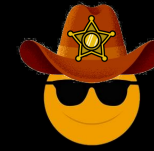


Default enclave for OS

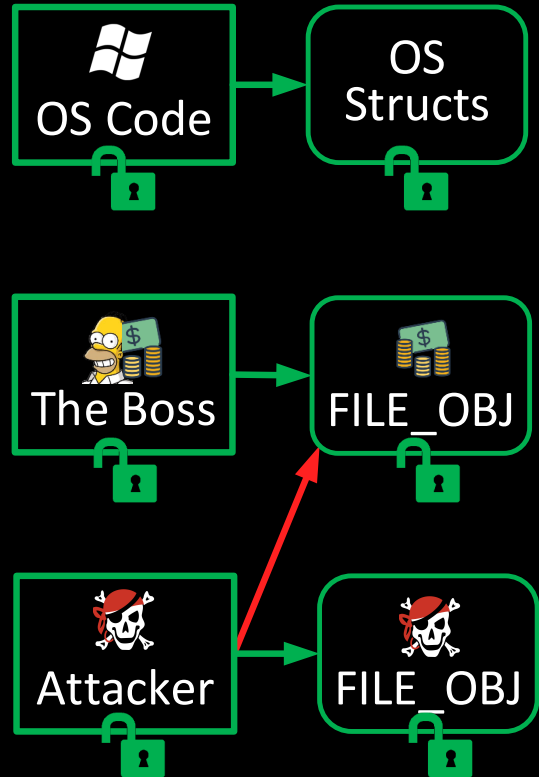


EPT pointer

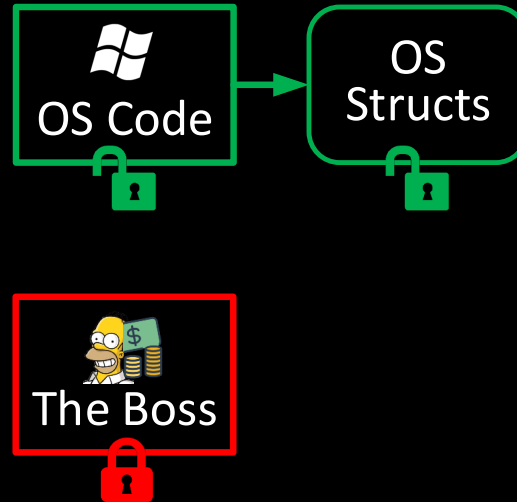
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



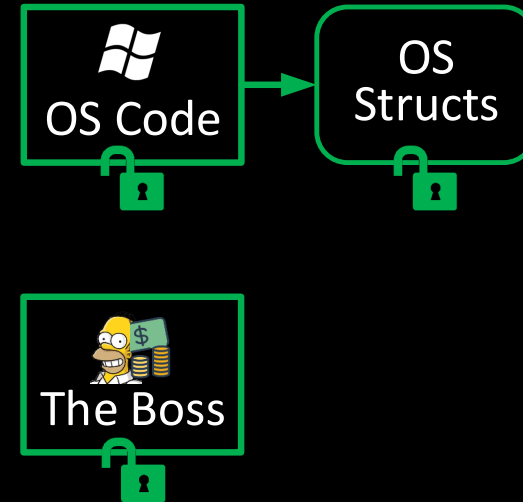
Current Situation



Default enclave for OS

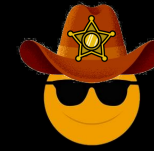


Enclave for Boss's Driver

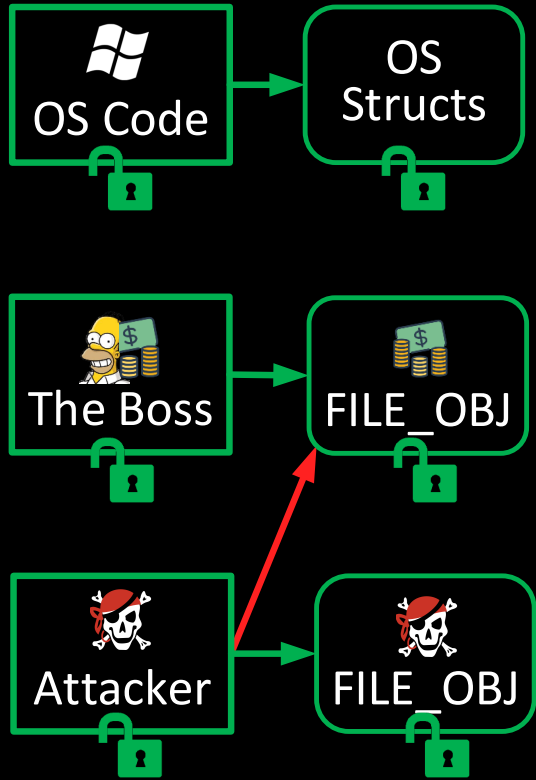


EPT pointer

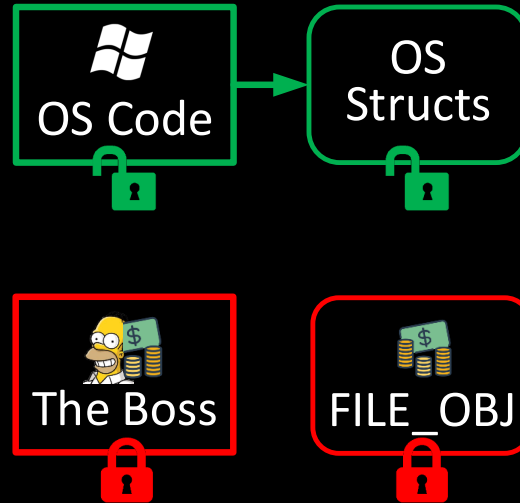
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



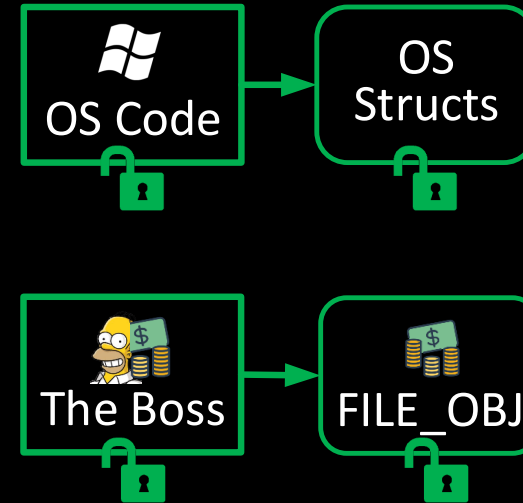
Current Situation



Default enclave for OS



Enclave for Boss's Driver

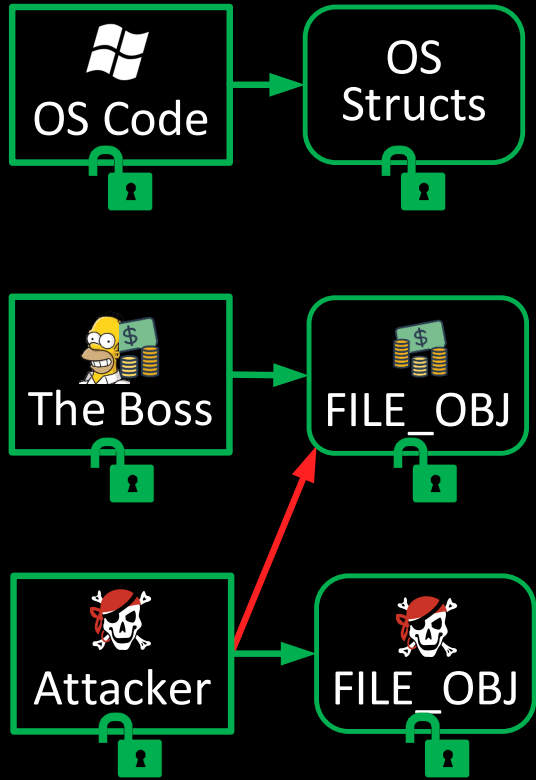


EPT pointer

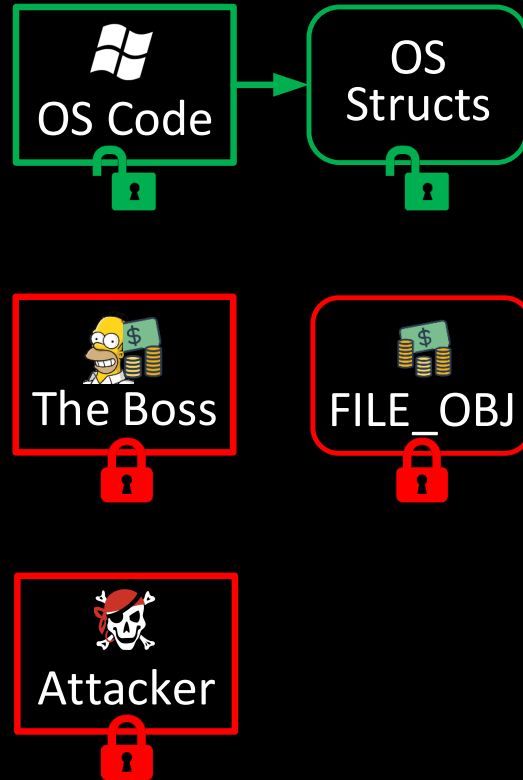
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



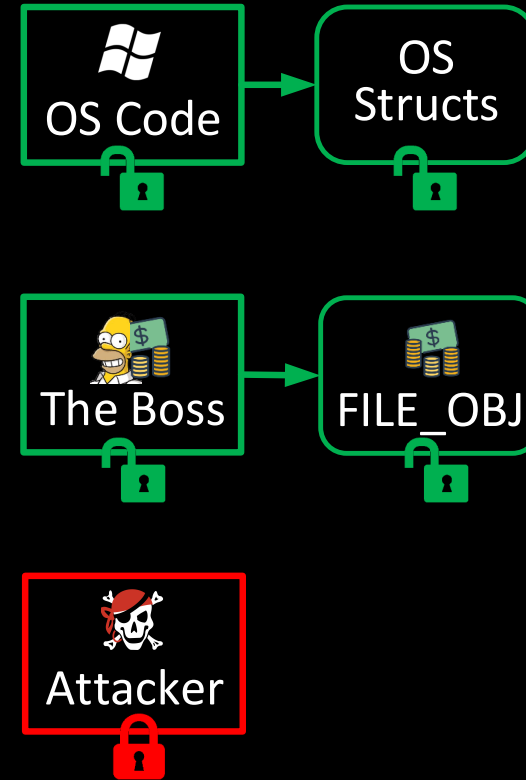
Current Situation



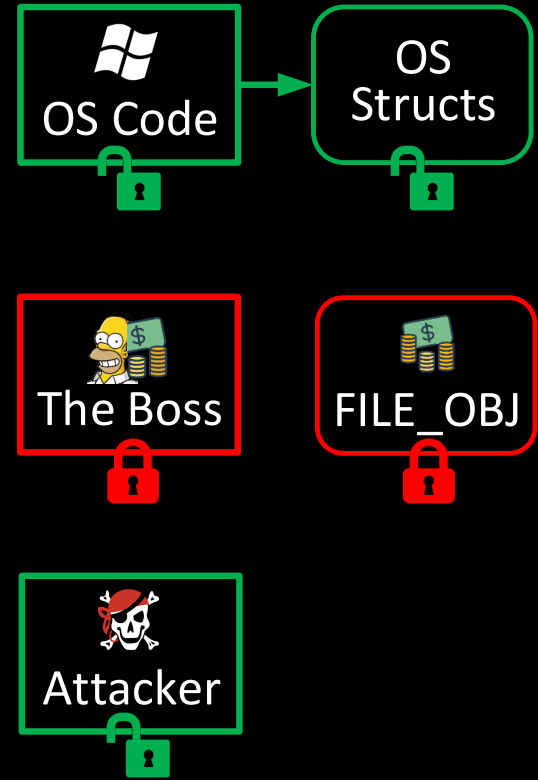
Default enclave for OS



Enclave for Boss's Driver



Enclave for Attacker's Driver

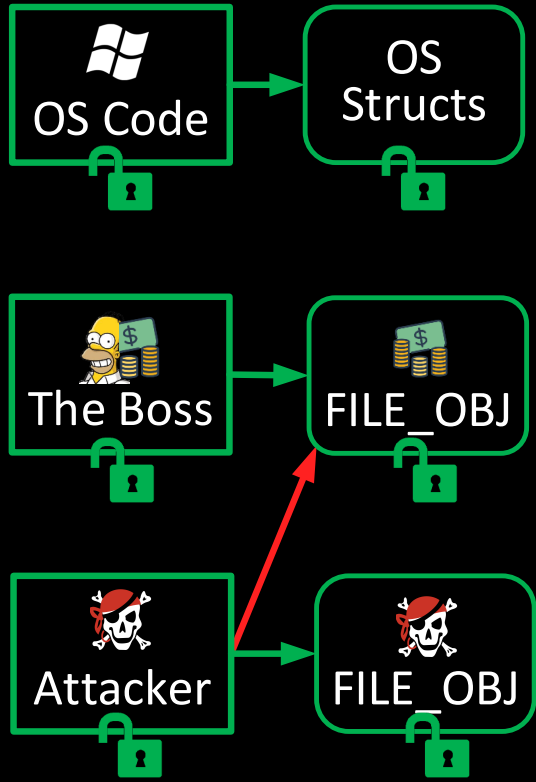


EPT pointer

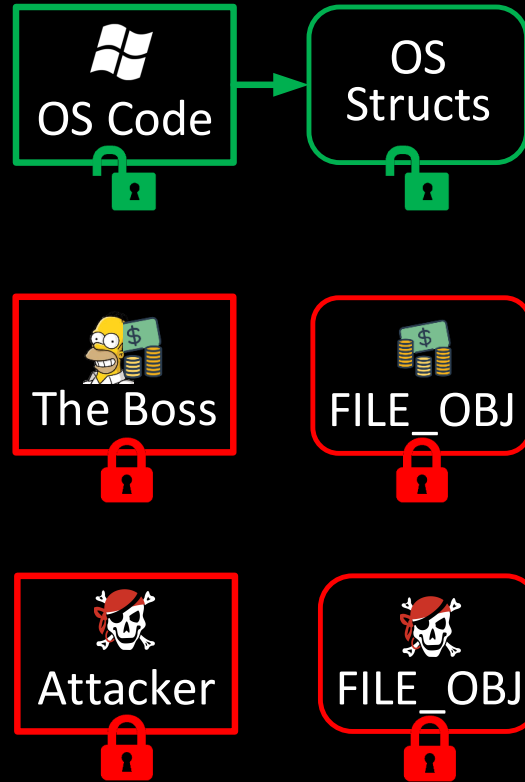
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



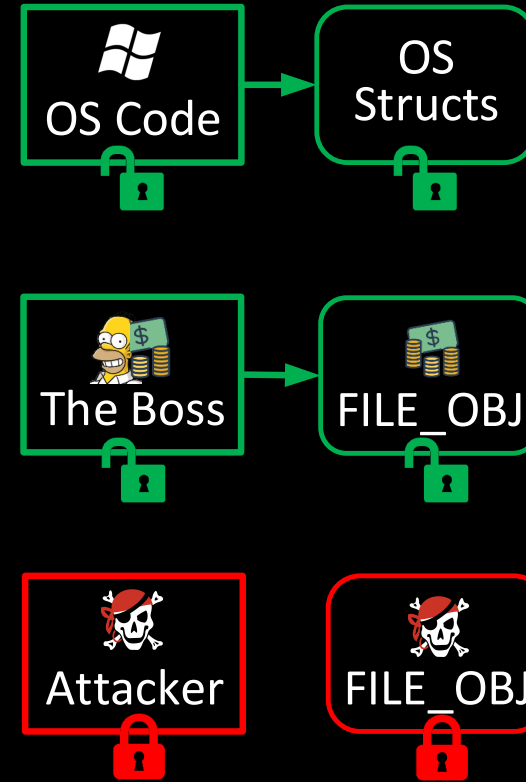
Current Situation



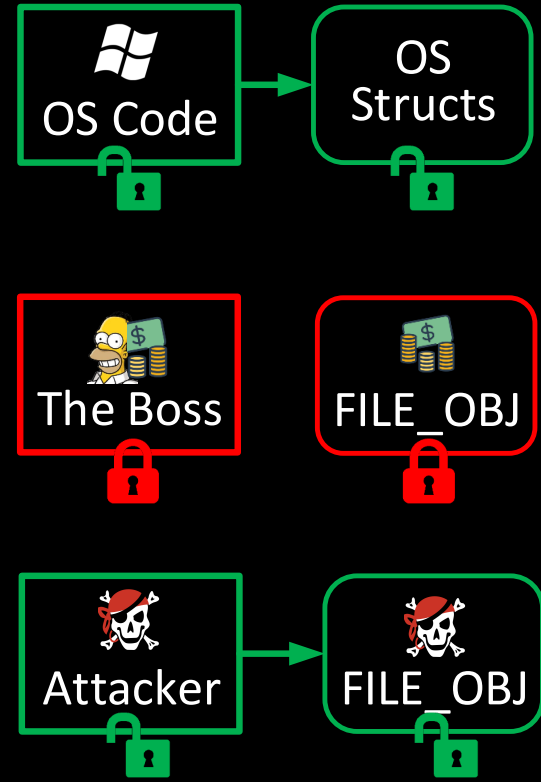
Default enclave for OS



Enclave for Boss's Driver



Enclave for Attacker's Driver

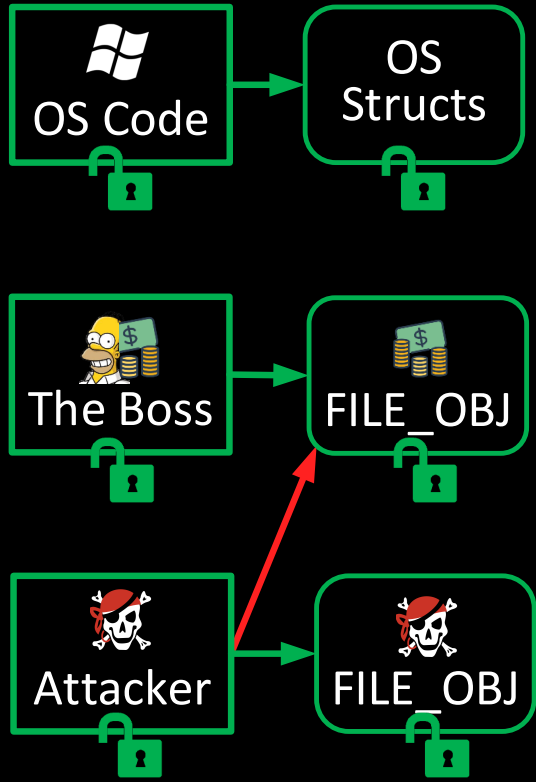


EPT pointer

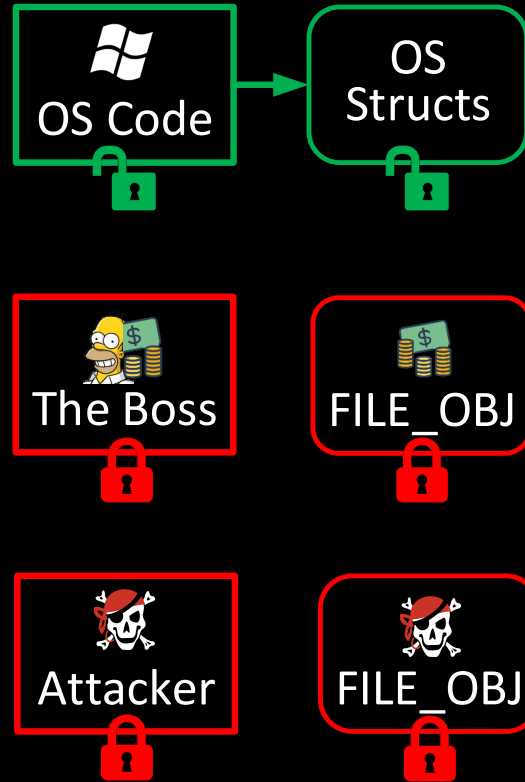
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



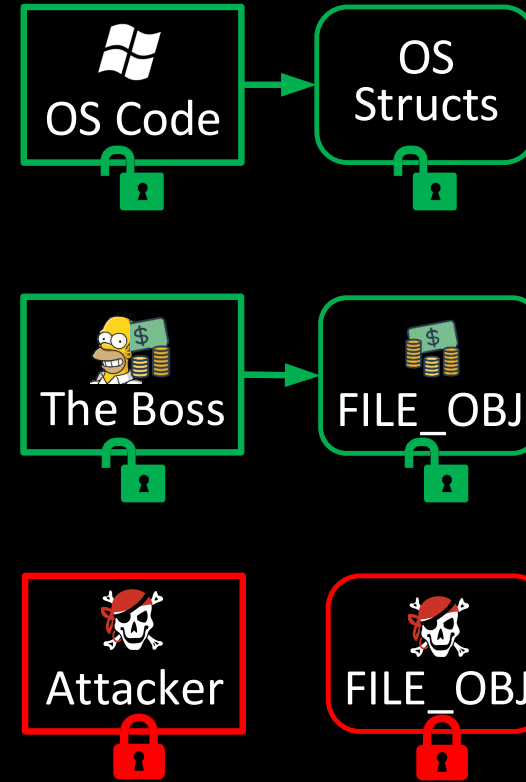
Current Situation



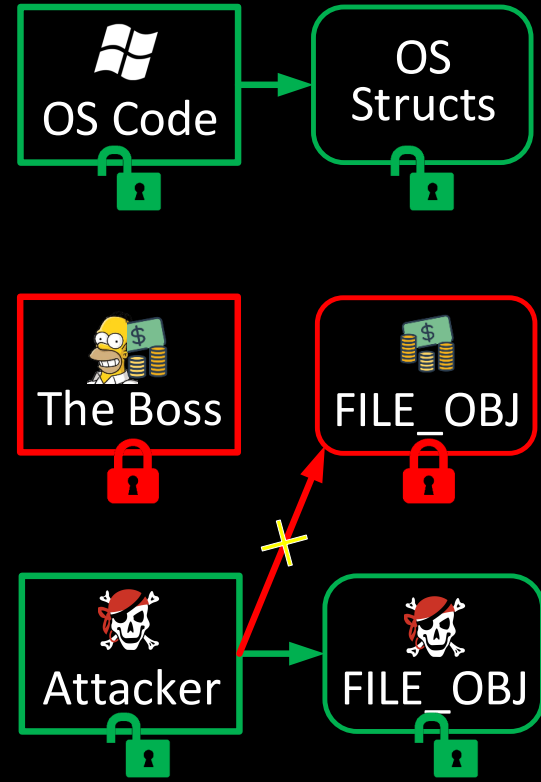
Default enclave for OS



Enclave for Boss's Driver



Enclave for Attacker's Driver

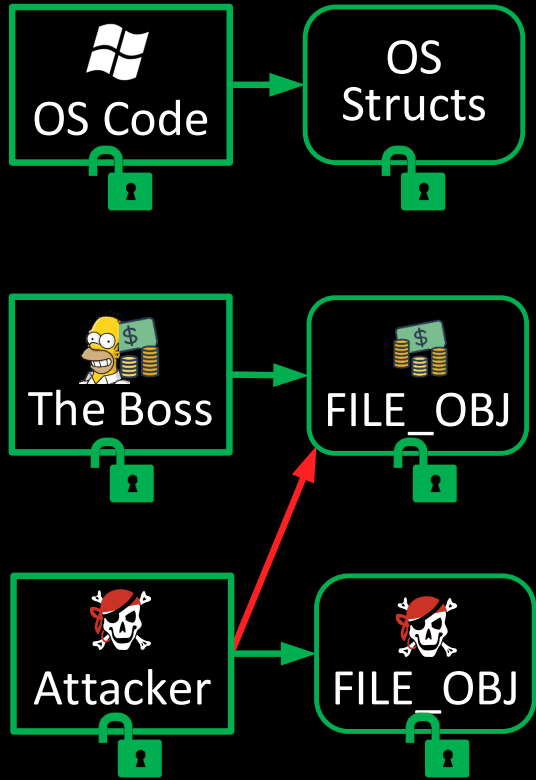


EPT pointer

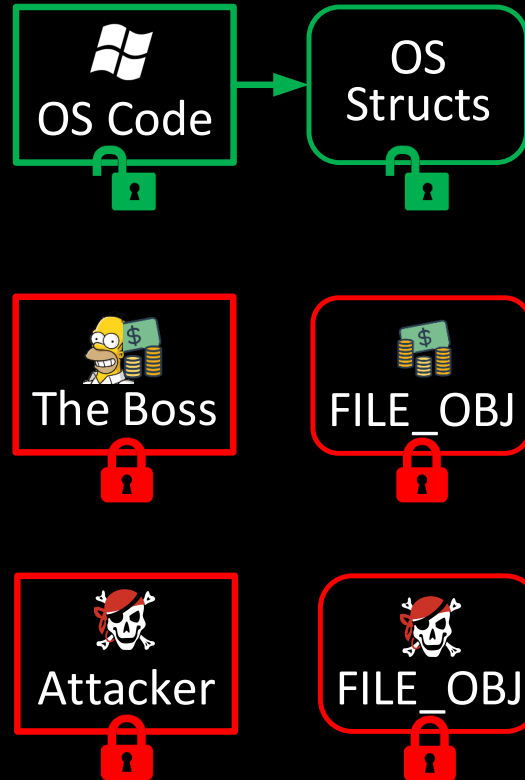
MEMORY RANGER PREVENTS FILE_OBJECT HIJACKING



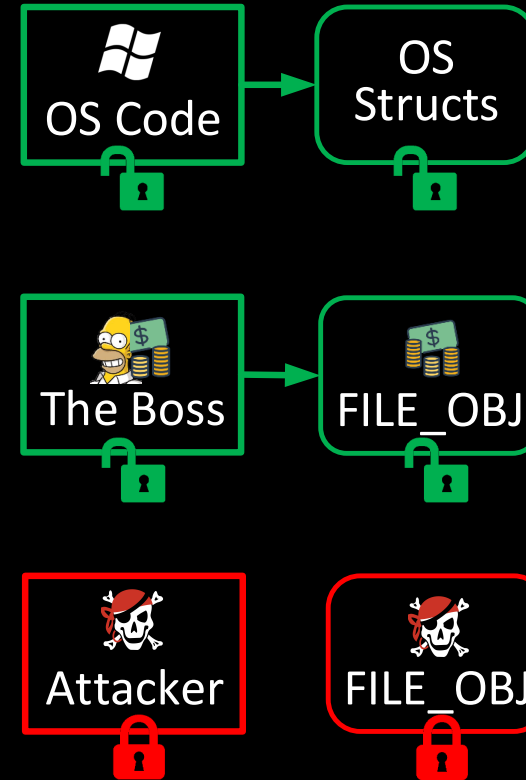
Current Situation



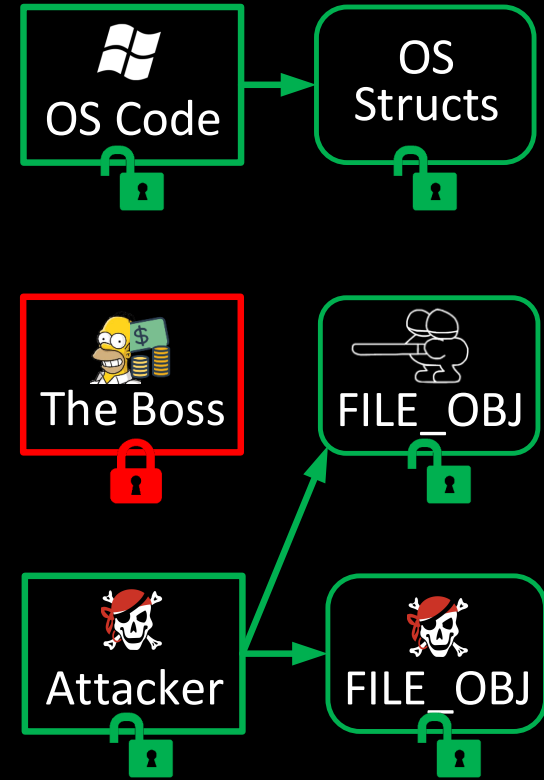
Default enclave for OS



Enclave for Boss's Driver



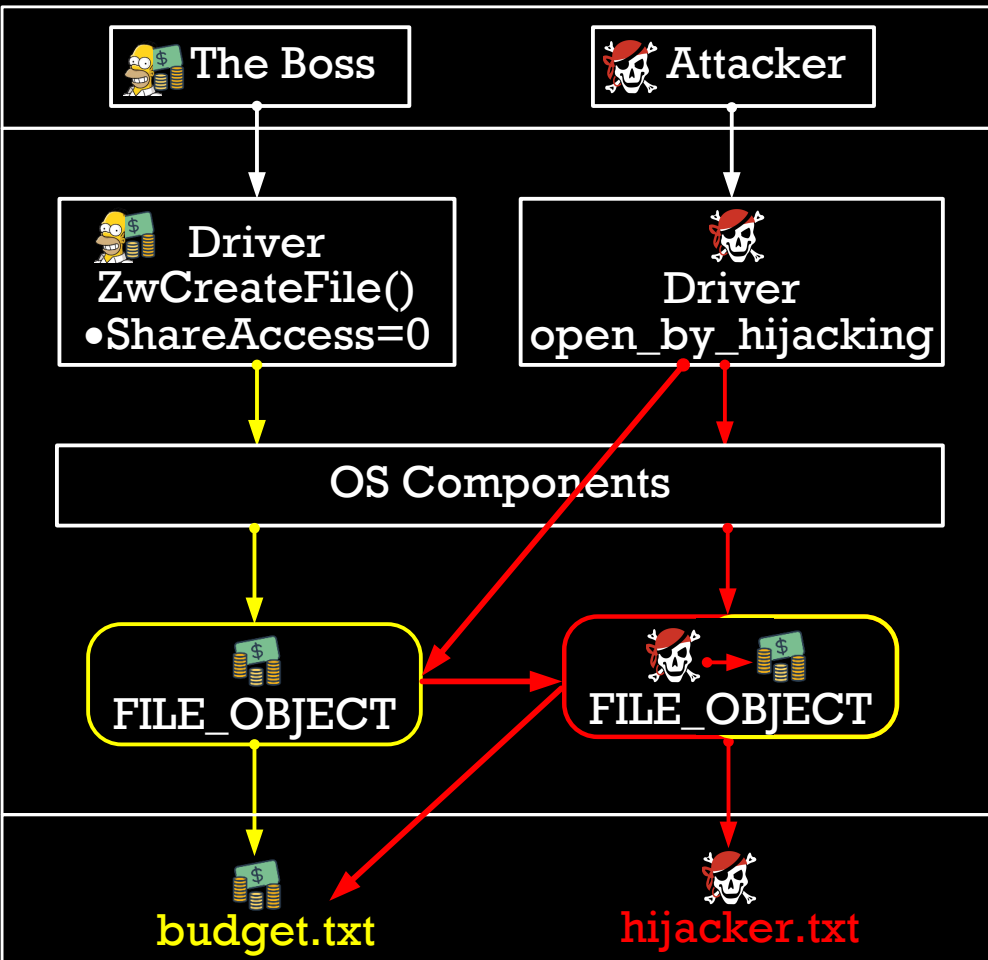
Enclave for Attacker's Driver



EPT pointer

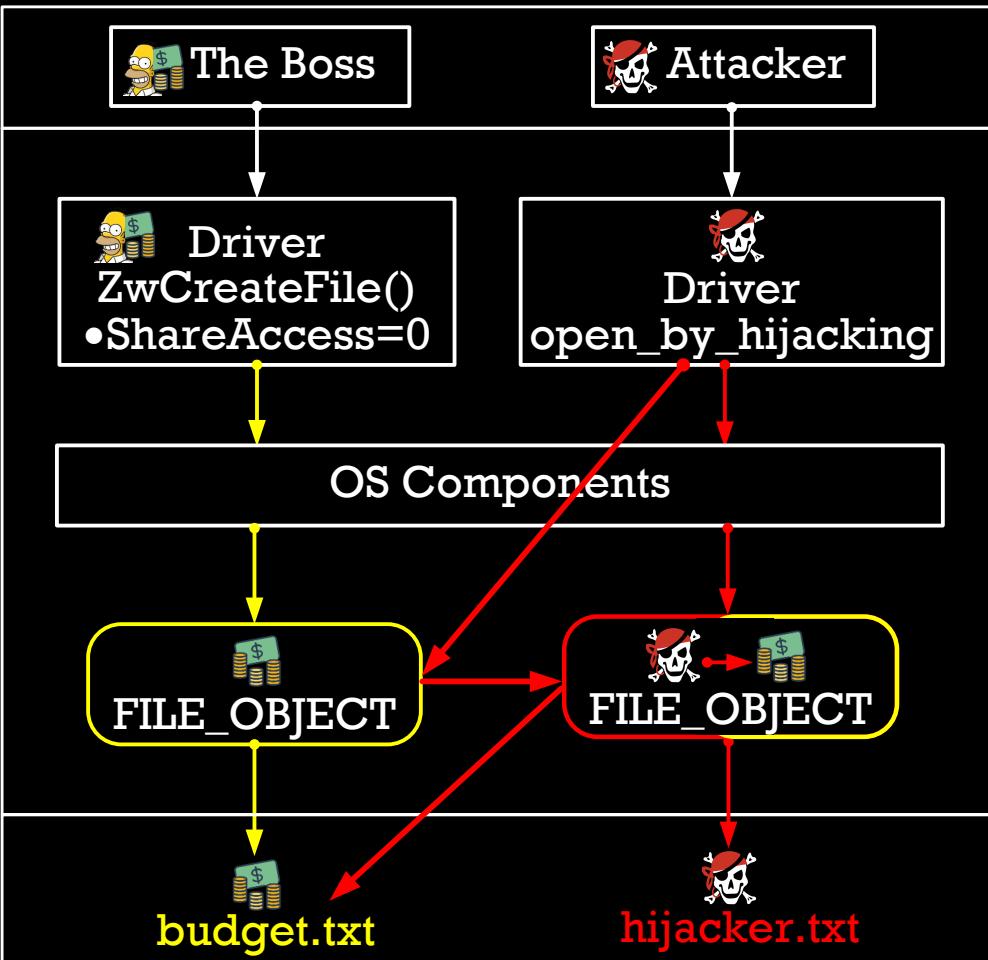
DEMO: PREVENTING THE HIJACKING

Attempt 2: The Hijacking Attack



DEMO: PREVENTING THE HIJACKING

Attempt 2: The Hijacking Attack



MemoryRanger



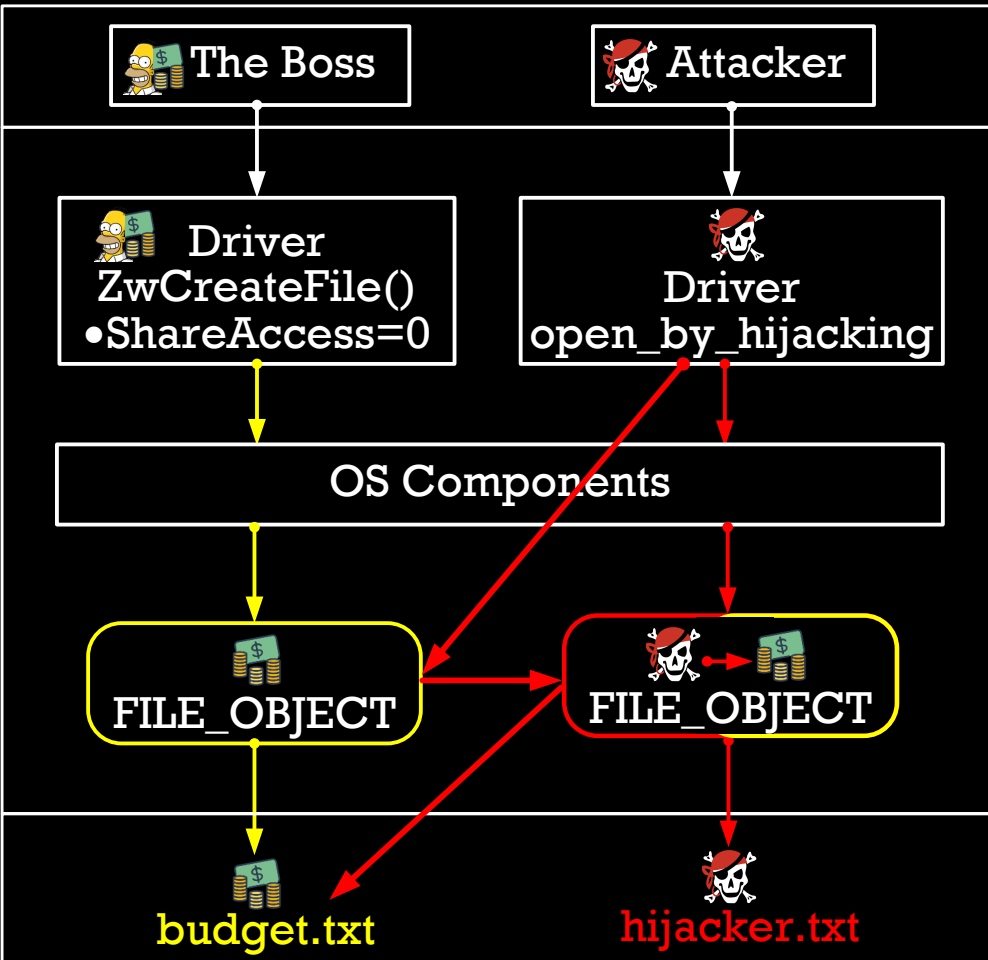
DEMO: THE ATTACK PREVENTION

The online version is here –

<https://www.youtube.com/watch?v=8ONmC5Do4I4?vq=hd1080>

DEMO: PREVENTING THE HIJACKING

Attempt 2: The Hijacking Attack

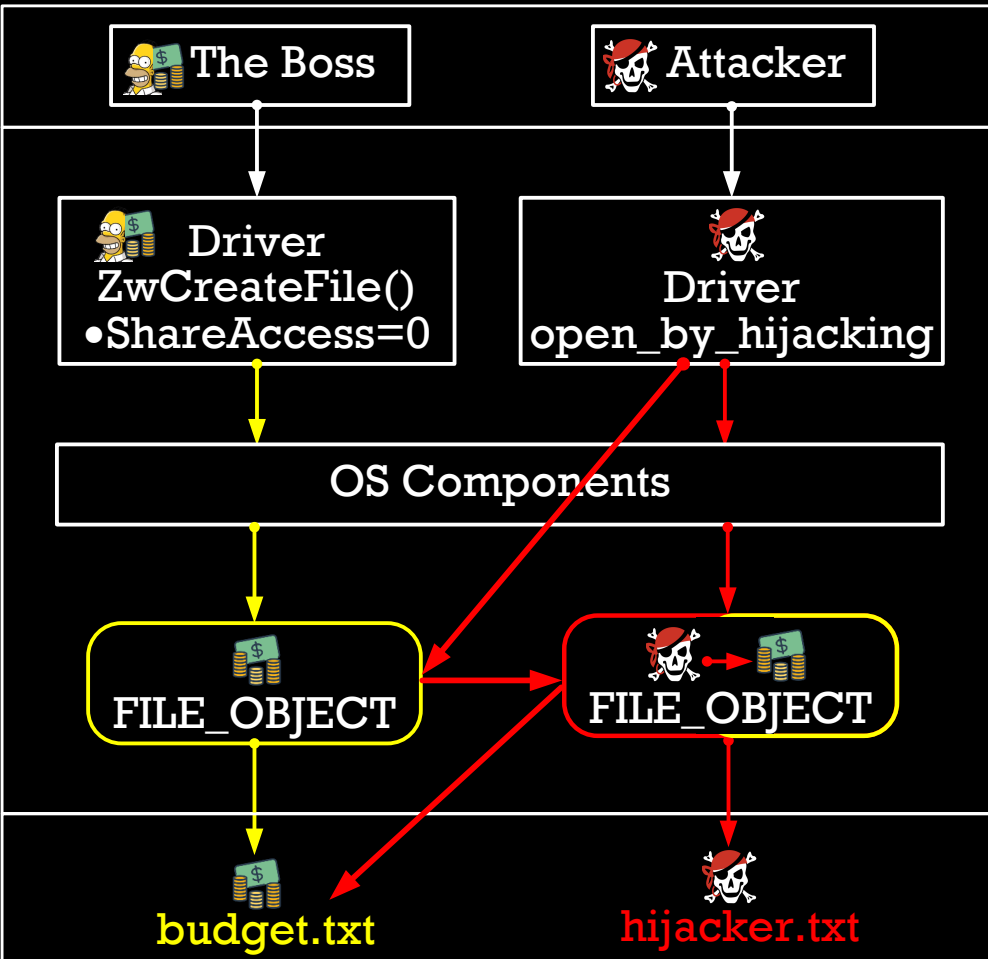


MemoryRanger

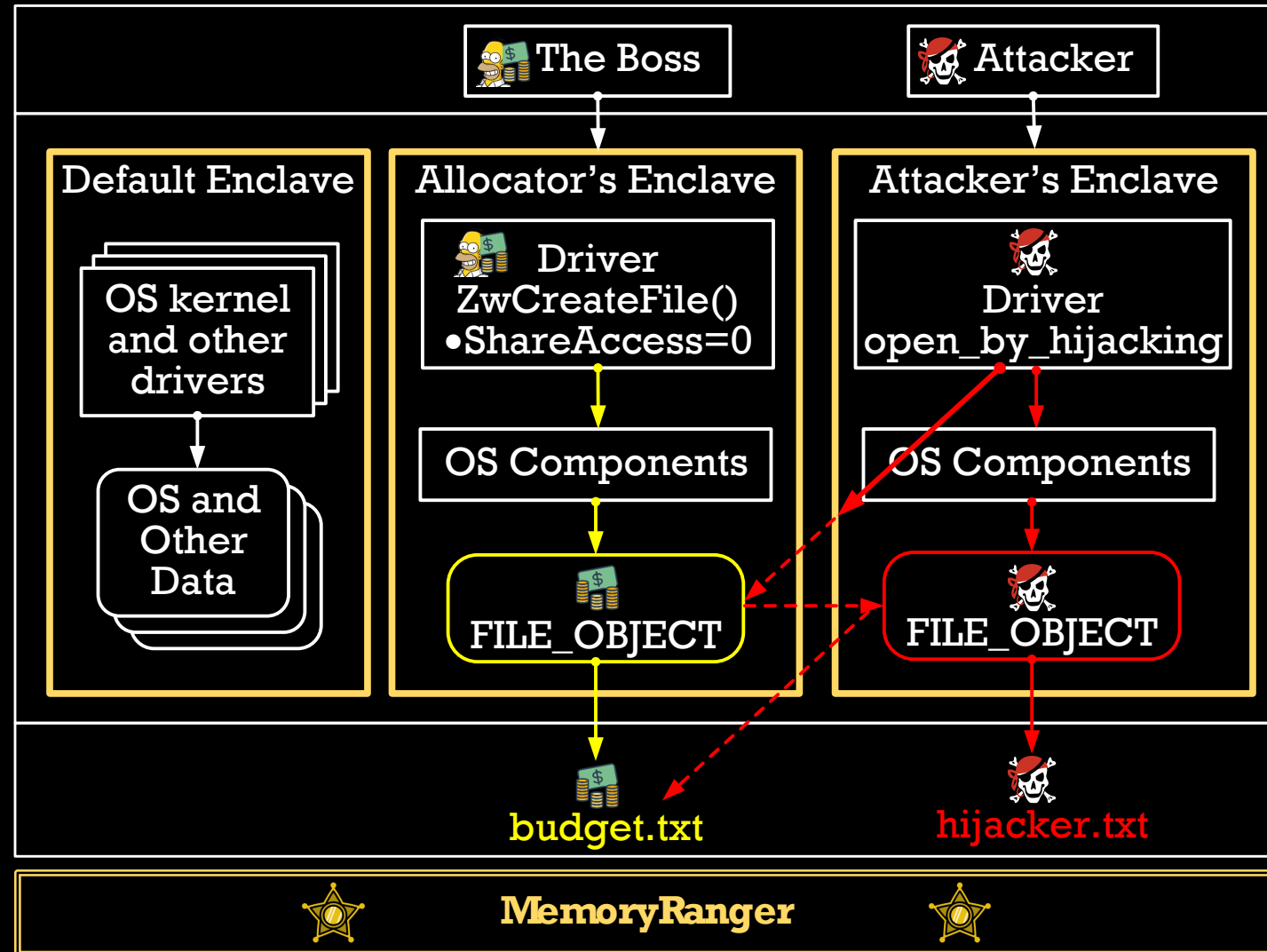


DEMO: PREVENTING THE HIJACKING

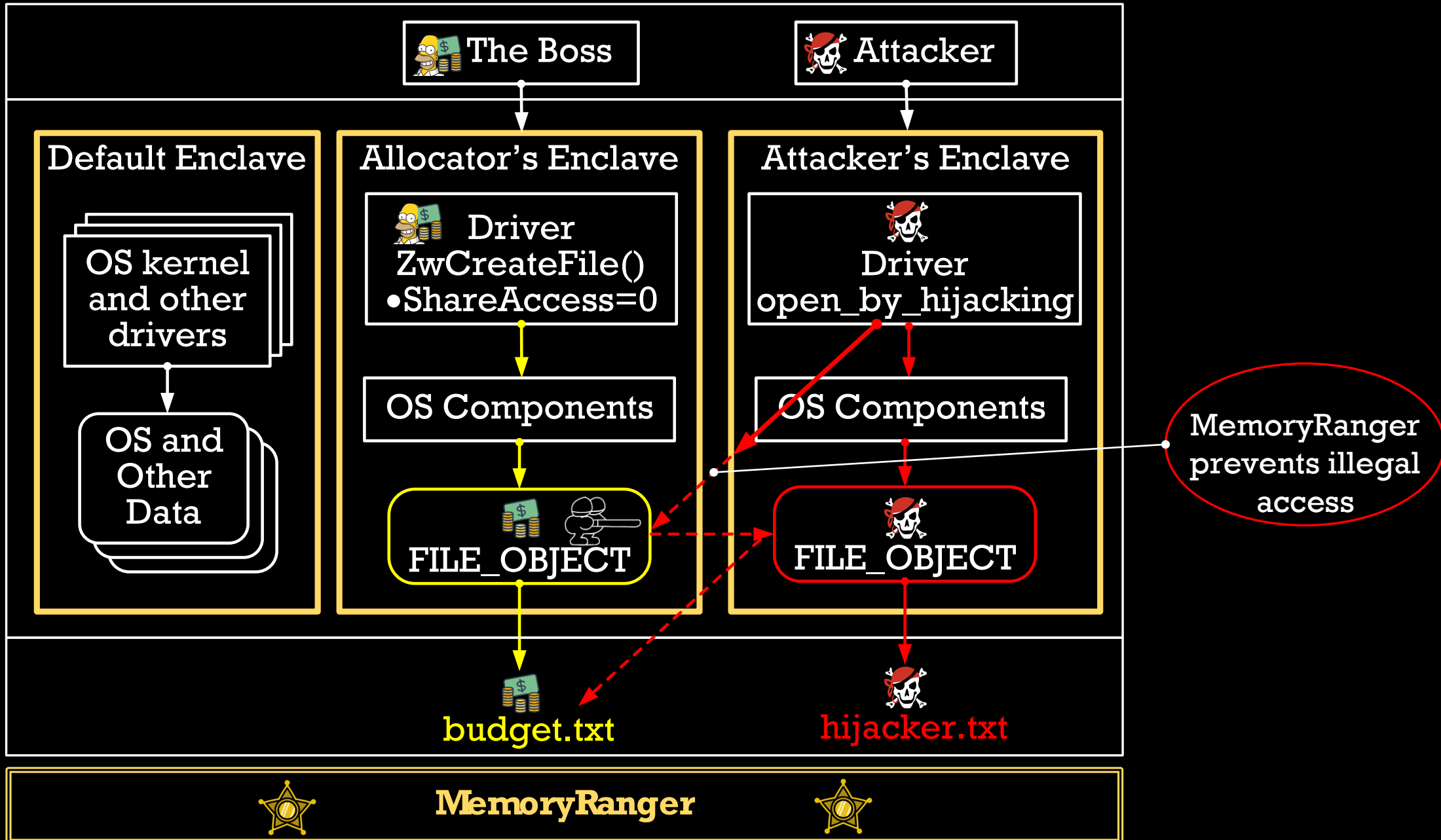
Attempt 2: The Hijacking Attack



Preventing the Hijacking Attack



Preventing the Hijacking Attack



MEMORY RANGER ARCHITECTURE

OS

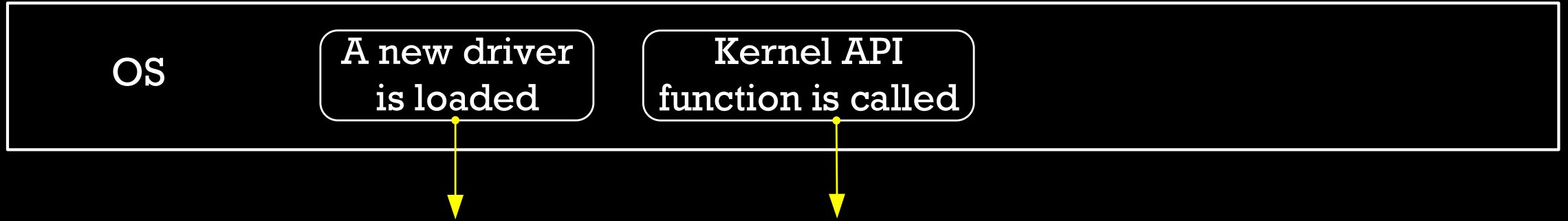
MEMORY RANGER ARCHITECTURE

OS

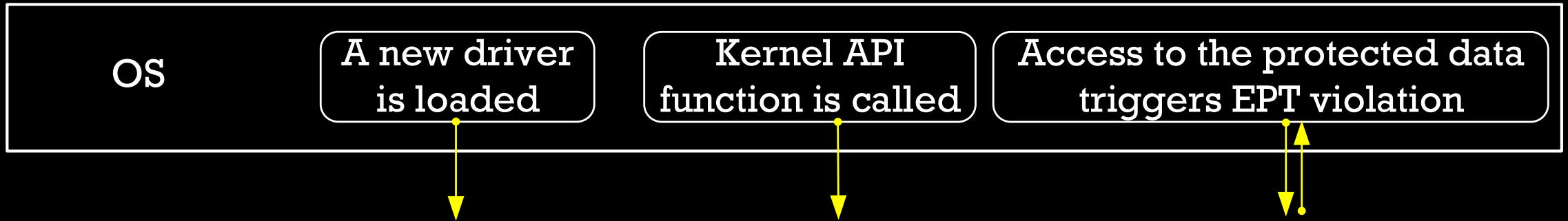
A new driver
is loaded



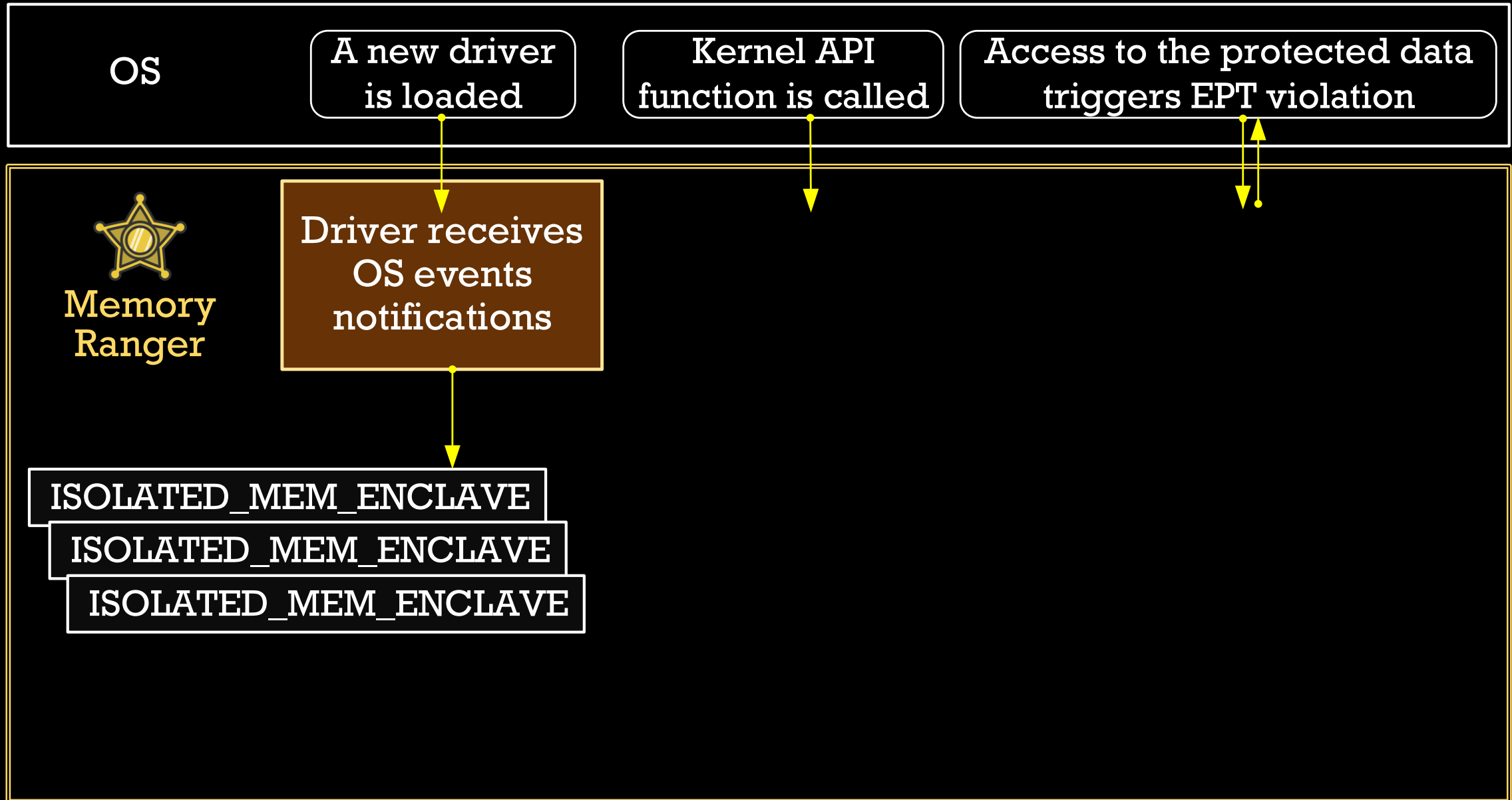
MEMORY RANGER ARCHITECTURE



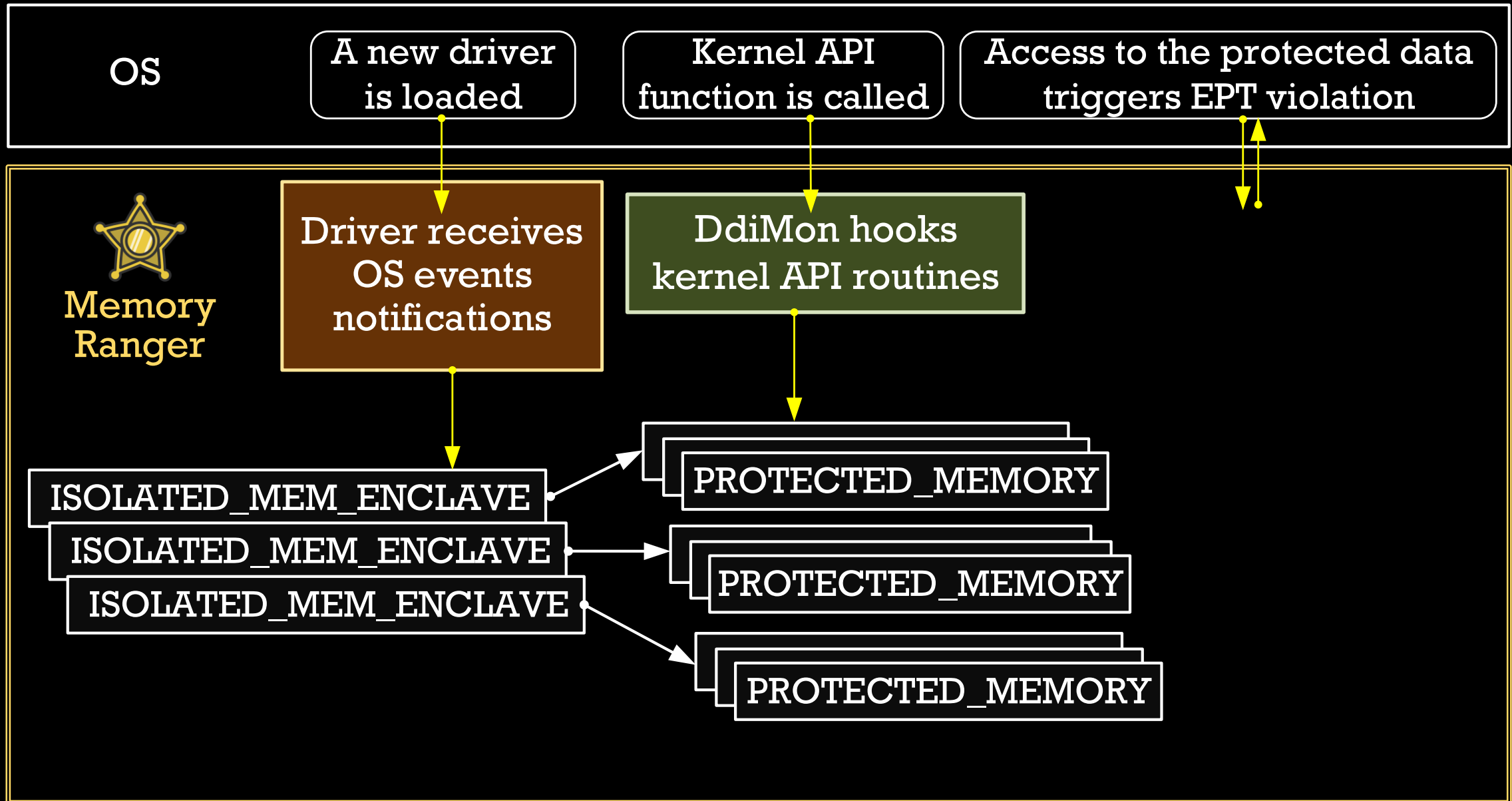
MEMORY RANGER ARCHITECTURE



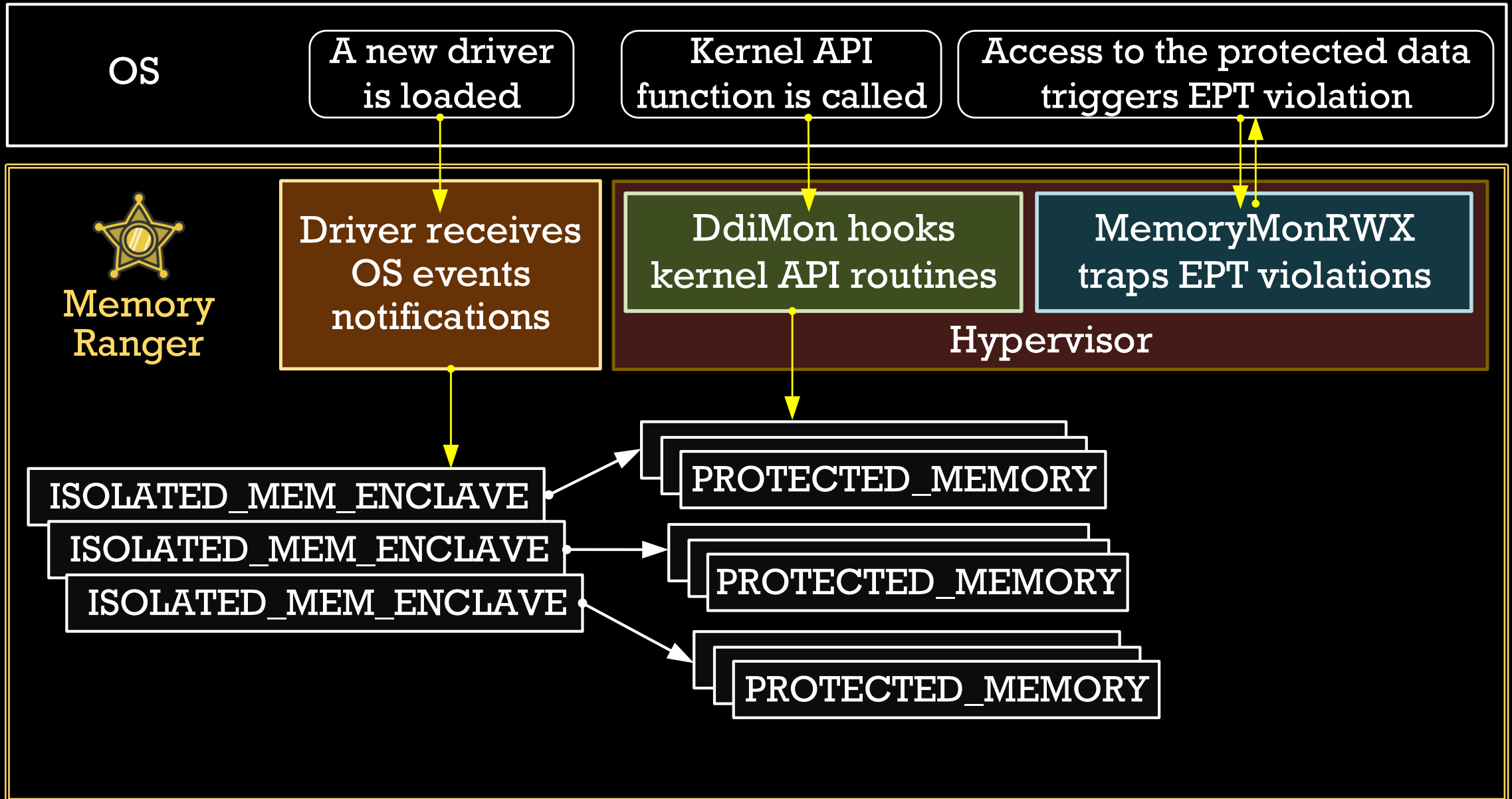
MEMORY RANGER ARCHITECTURE



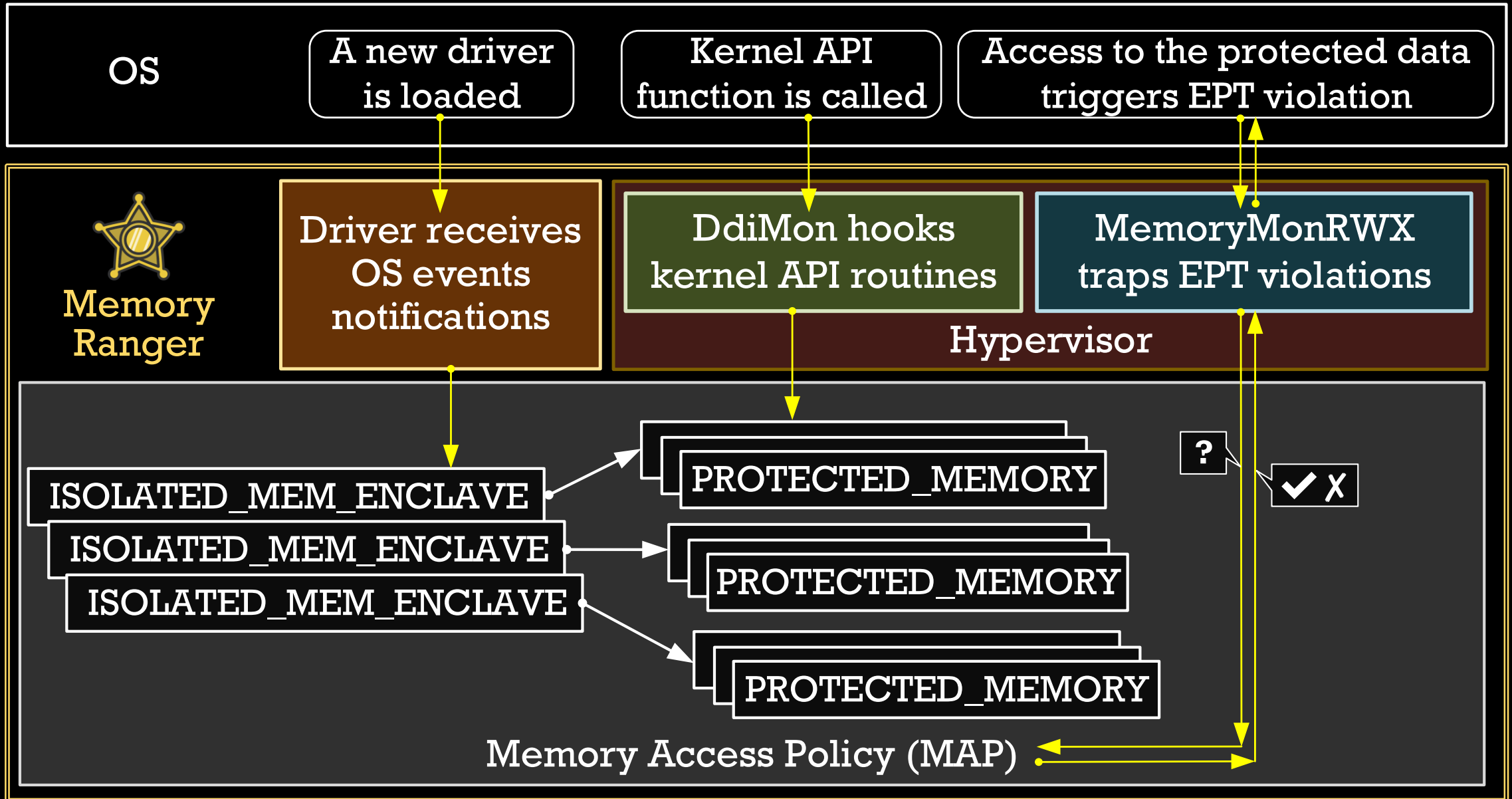
MEMORY RANGER ARCHITECTURE



MEMORY RANGER ARCHITECTURE



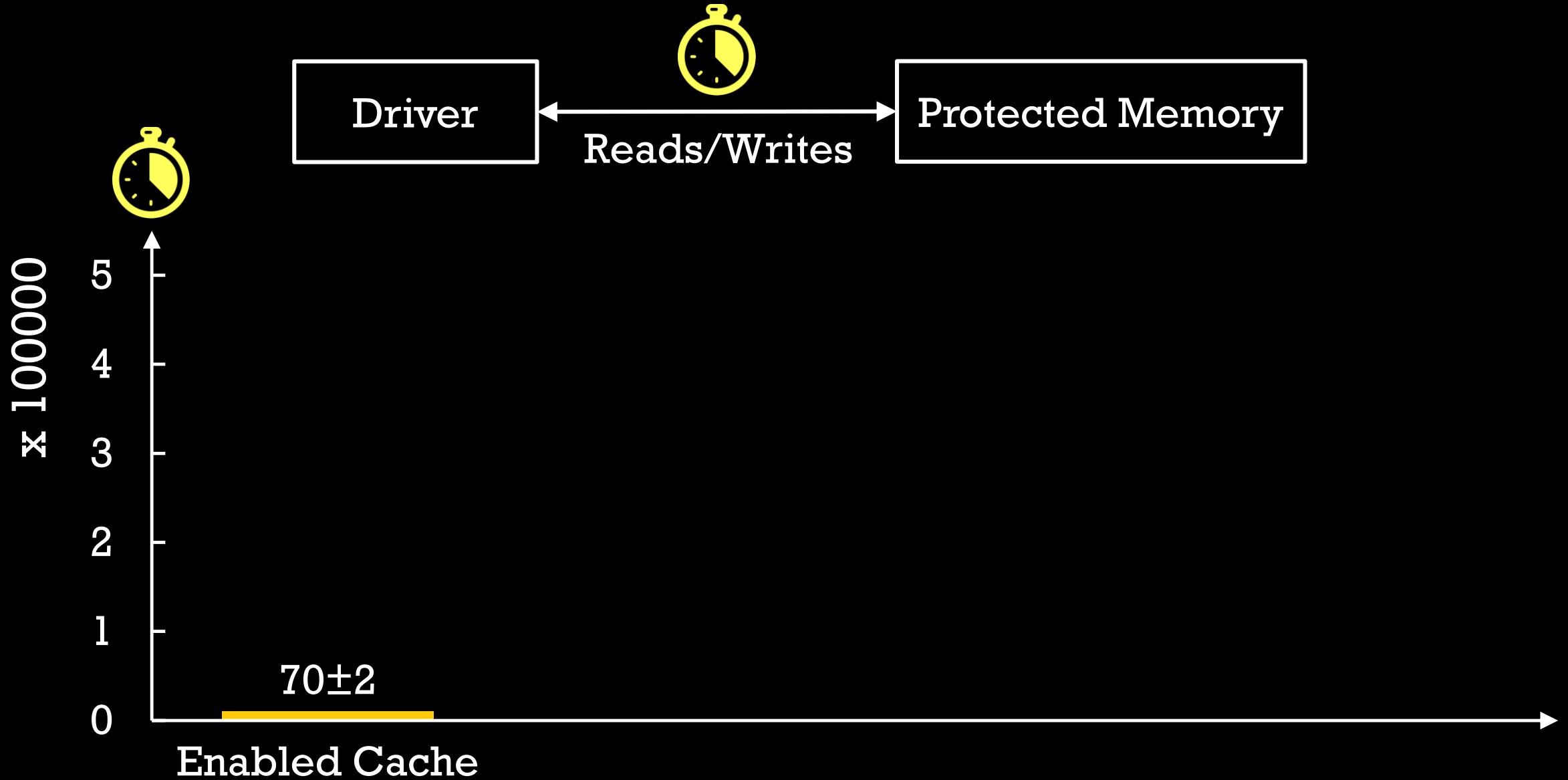
MEMORY RANGER ARCHITECTURE



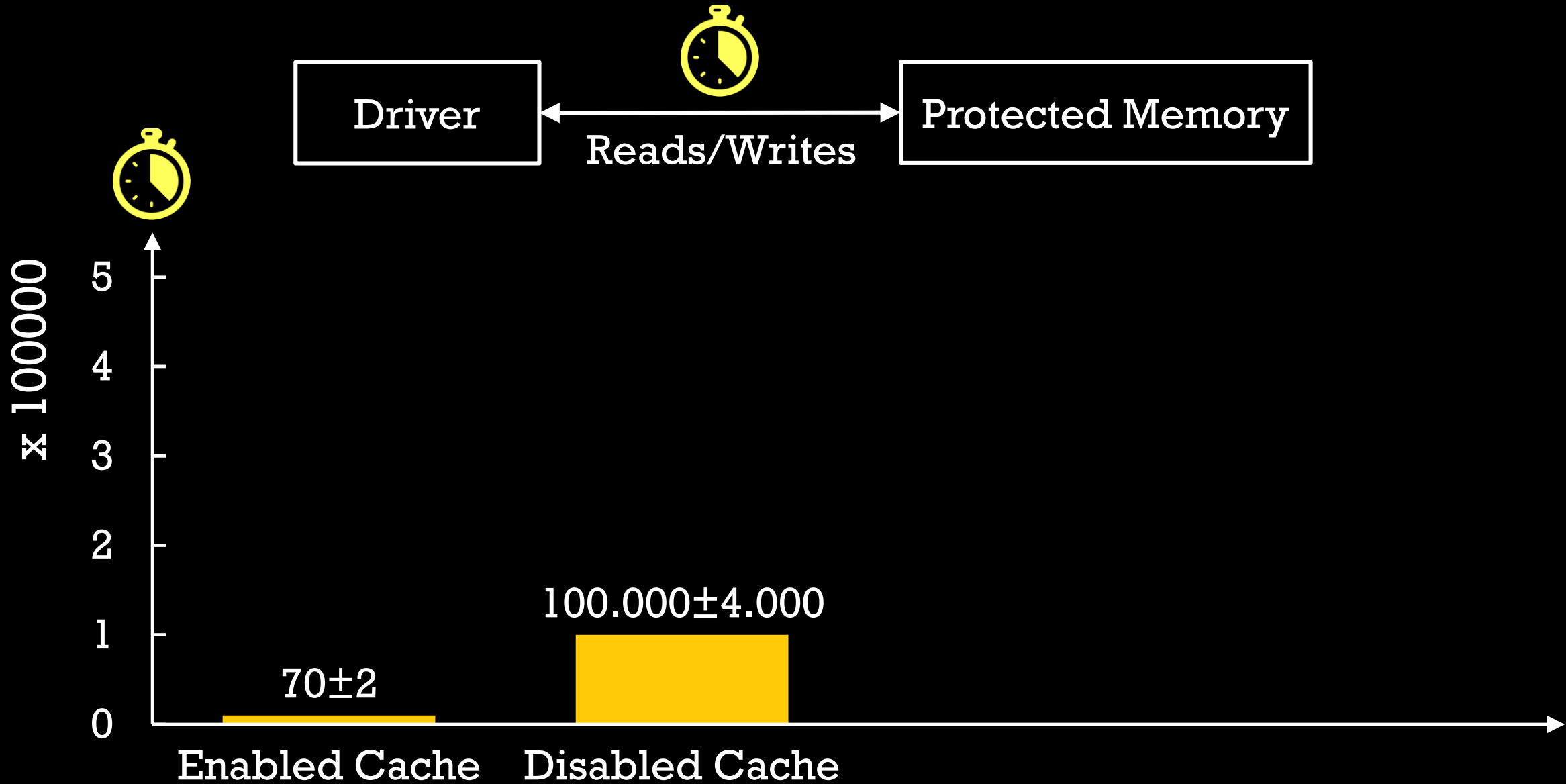
MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



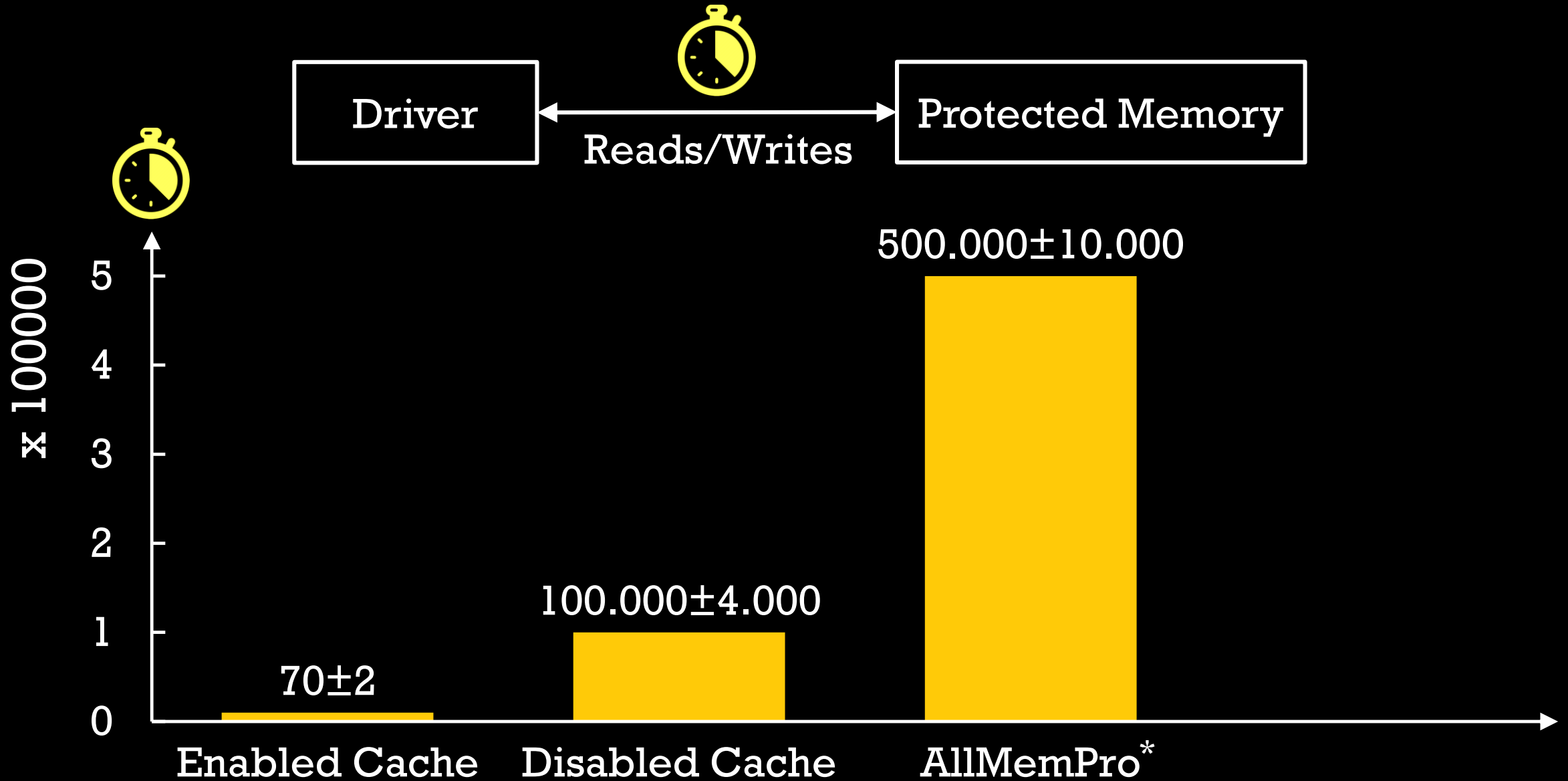
MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME

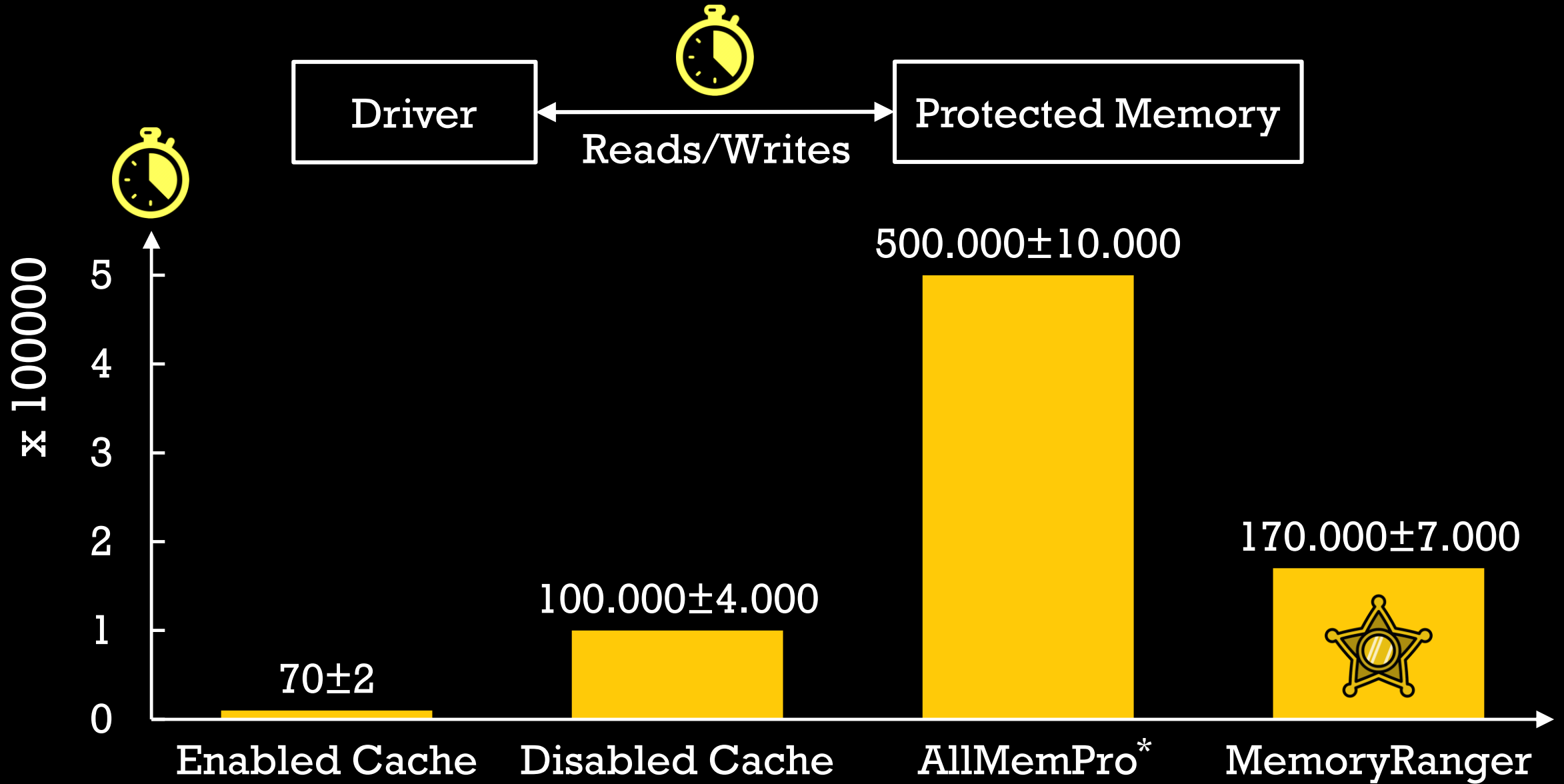


MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



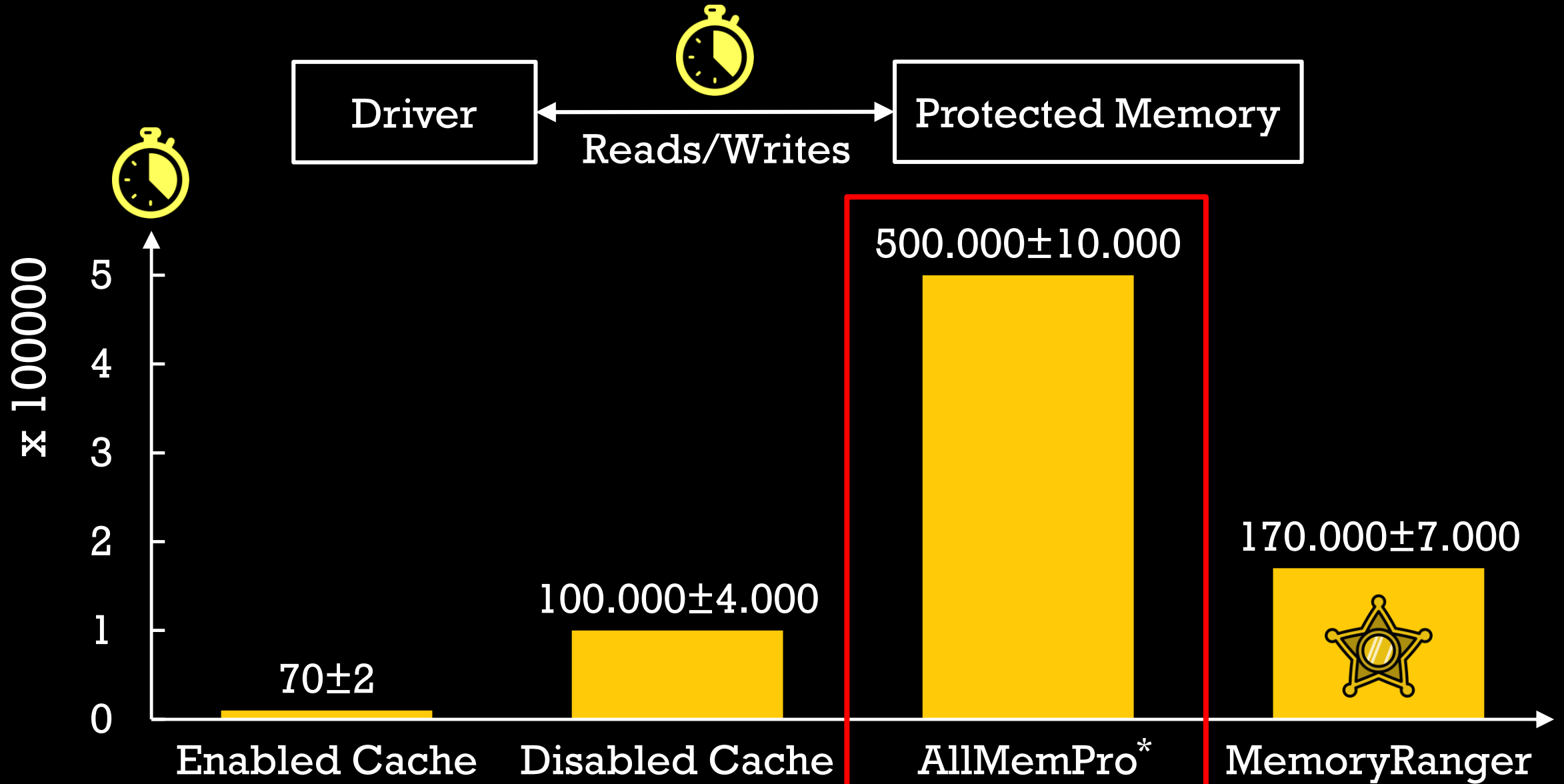
* AllMemPro details - <http://bit.ly/AllMemPro>

MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



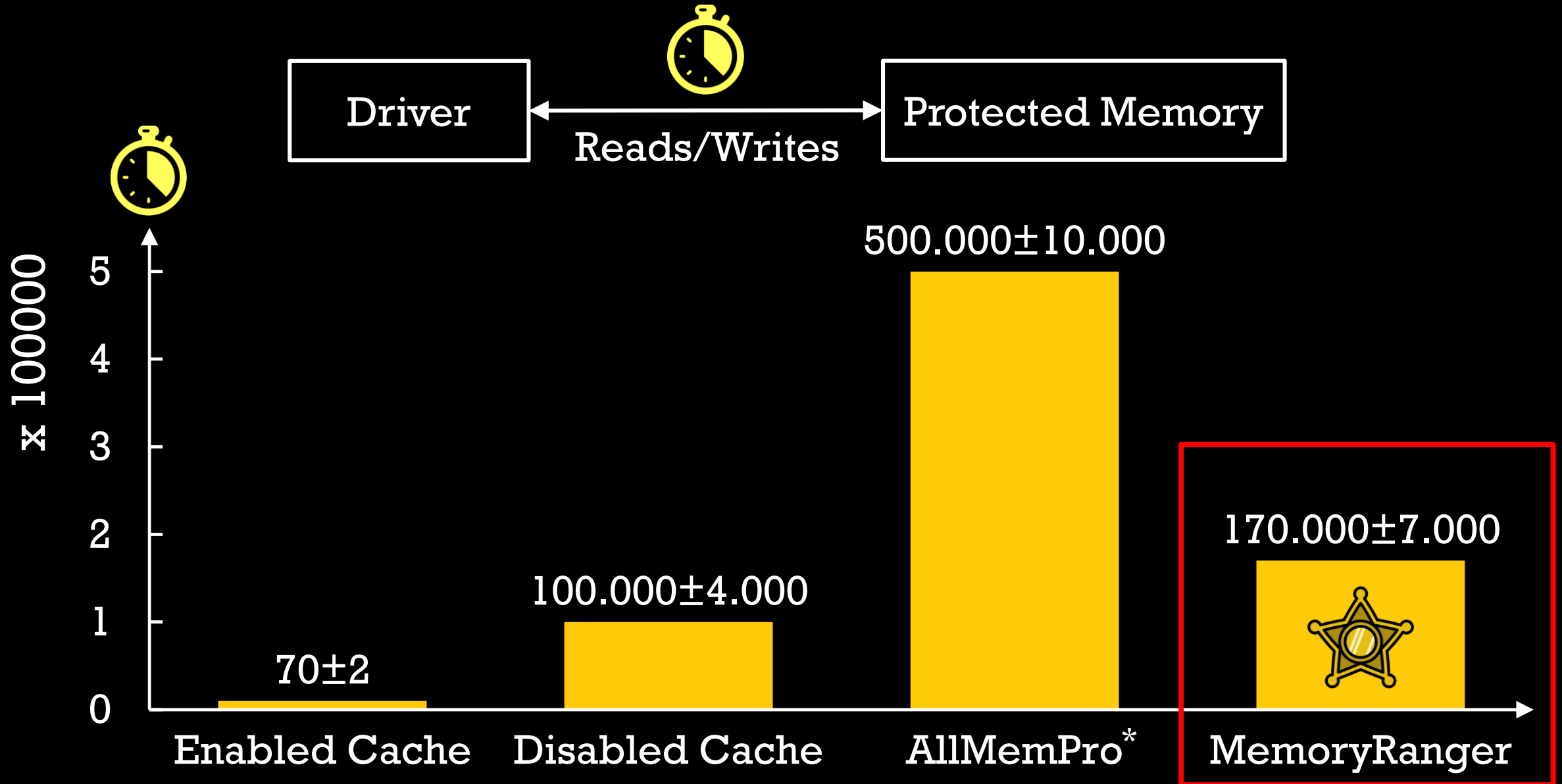
* AllMemPro details - <http://bit.ly/AllMemPro>

MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



* AllMemPro details - <http://bit.ly/AllMemPro>

MEMORY RANGER BENCHMARKS: MEMORY ACCESS TIME



* AllMemPro details - <http://bit.ly/AllMemPro>


THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions					
Integrity					
Confidentiality					



THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code				
Integrity					
Confidentiality					




THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code				
	 OS Code				
Integrity					
Confidentiality					



THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code				
	 OS Code				
Integrity	 Device Guard				
Confidentiality					






THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code				
	 OS Code				
Integrity	 Device Guard				
Confidentiality					





THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div>Drivers code</div> <div>OS Code</div>	<div>Allocated data</div>			
Integrity	 Device Guard				
Confidentiality					






THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div><div>Drivers code</div><div> OS Code</div></div>	<div>Allocated data</div>			
Integrity	 Device Guard				
Confidentiality					






THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div>Drivers code</div> <div>→</div> <div>Allocated data</div>				
	<div>OS Code</div>		<div>EPROCESS structures</div> <div>PsActiveProcessLinks</div>		
			<div>LDR_DATA_TABLE_ENTRY structures</div> <div>PsLoadedModuleList</div>		
			<div>DRIVER_OBJECT structures</div> <div>MajorFunction[]</div>		
Integrity	 Device Guard				
Confidentiality					









THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div>Drivers code</div>	<div>Allocated data</div>			
	<div> OS Code</div>		EPROCESS structures		
			PsActiveProcessLinks		
			LDR_DATA_TABLE_ENTRY structures		
			PsLoadedModuleList		
			DRIVER_OBJECT structures		
			MajorFunction[]		
Integrity	 Device Guard				
Confidentiality			(skipped)		









THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code	Allocated data			
	 OS Code		EPROCESS structures		
			PsActiveProcessLinks		
			LDR_DATA_TABLE_ENTRY structures		
Integrity	 OS Code		PsLoadedModuleList		
			DRIVER_OBJECT structures		
			MajorFunction[]		
Confidentiality	Device Guard		Patch Guard		
			(skipped)		









THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div>Drivers code</div>	<div>Allocated data</div>			
	<div> OS Code</div>		EPROCESS structures		
			<div>PsActiveProcessLinks</div>	<div>Token</div>	
			LDR_DATA_TABLE_ENTRY structures		
			<div>PsLoadedModuleList</div>		
			DRIVER_OBJECT structures		
			<div>MajorFunction[]</div>		
Integrity	 Device Guard		 Patch Guard		
Confidentiality			(skipped)		










THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code	Allocated data			
	 OS Code		EPROCESS structures		Token
			LDR_DATA_TABLE_ENTRY structures		
			DRIVER_OBJECT structures		
					FILE_OBJECT structures
Integrity	 Device Guard		 Patch Guard		
Confidentiality			(skipped)		












THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code	Allocated data			
	 OS Code		EPROCESS structures	Token	? ? ?
			LDR_DATA_TABLE_ENTRY structures		
			DRIVER_OBJECT structures		
				FILE_OBJECT structures	
Integrity	 Device Guard		 Patch Guard		
Confidentiality			(skipped)		













THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	Drivers code	Allocated data 	EPROCESS structures		
	 OS Code		PsActiveProcessLinks	Token	? ? ?
			LDR_DATA_TABLE_ENTRY structures PsLoadedModuleList		
			DRIVER_OBJECT structures MajorFunction[]		
				FILE_OBJECT structures	
Integrity	 Device Guard		 Patch Guard		
Confidentiality			(skipped)		

THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div>Drivers code</div> <div></div>	<div>Allocated data</div> <div></div>	<div>EPROCESS structures</div> <div>PsActiveProcessLinks</div>	<div>Token</div> <div></div>	<div>???</div>
	<div>OS Code</div> <div></div>		<div>LDR_DATA_TABLE_ENTRY structures</div> <div>PsLoadedModuleList</div>		
			<div>DRIVER_OBJECT structures</div> <div>MajorFunction[]</div>		
				<div>FILE_OBJECT structures</div>	
Integrity	<div> Device Guard</div>	<div></div>	<div> Patch Guard</div>	<div></div>	
Confidentiality	<div></div>	<div></div>	<div>(skipped)</div>	<div></div>	

THE CURRENT SITUATION WITH ATTACKS ON WINDOWS MEMORY

	Code	Drivers allocations	Dynamically Allocated Data by the OS		
Memory Regions	<div>Drivers code</div> <div></div>	<div>Allocated data</div> <div></div>	<div>EPROCESS structures</div> <div>PsActiveProcessLinks</div>	<div>Token</div> <div></div>	<div>?</div> <div>?</div> <div>?</div>
	<div>OS Code</div> <div></div>		<div>LDR_DATA_TABLE_ENTRY structures</div> <div>PsLoadedModuleList</div>		
			<div>DRIVER_OBJECT structures</div> <div>MajorFunction[]</div>		
				<div></div> <div>FILE_OBJECT structures</div>	
Integrity	 Device Guard		 Patch Guard		
Confidentiality			(skipped)		

CONCLUSION

- All modern Windows OSes are vulnerable to FILE_OBJECT hijacking
- MemoryRanger prevents the hijacking attack by running drivers into isolated memory enclaves
- Research is ongoing

Thank you!

Igor Korkin igor.korkin@gmail.com

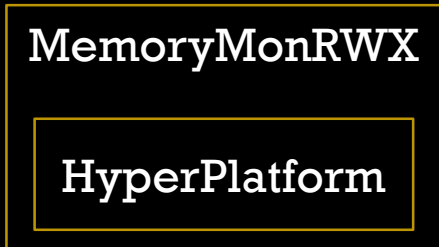
All the details & my CV are here igorkorkin.blogspot.com



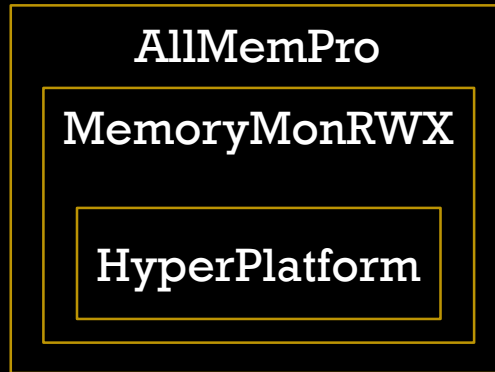
MEMORY RANGER HISTORY



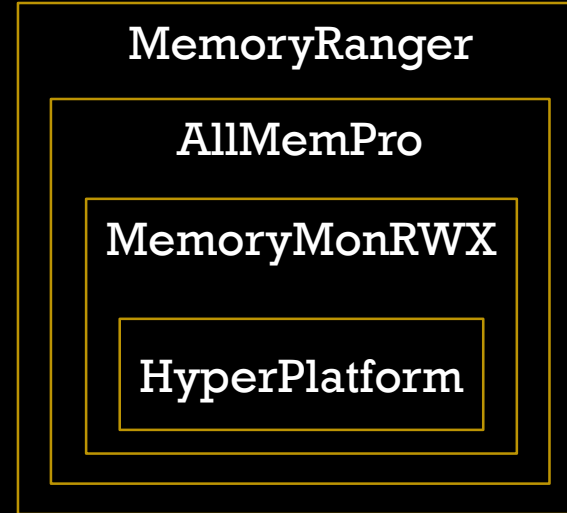
Step 1



Step 2



Step 3



Step 4



Step 5



1. Korkin, I., & Tanda, S. (2016). Monitoring & controlling kernel-mode events by HyperPlatform. Recon, Canada.
2. Korkin, I., & Tanda, S. (2017). Detect Kernel-Mode Rootkits via Real Time Logging & Controlling Memory Access. ADFS, USA.
3. Korkin, I. (2018). Hypervisor-Based Active Data Protection for Integrity and Confidentiality of Dynamically Allocated Memory in Windows Kernel. ADFS, USA.
4. Korkin, I. (2018). Divide et Impera: MemoryRanger Runs Drivers in Isolated Kernel Spaces. BlackHat, UK
5. Korkin, I. (2019). MemoryRanger Prevents Hijacking FILE_OBJECT structures in Windows Kernel. ADFS, USA.