# MODULE 1 / UNIT 2 INSERTION SORT

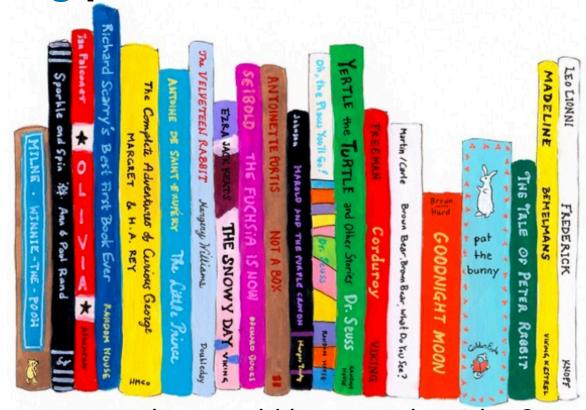


### **Today**

What is an algorithm?

- Understanding time complexities
- Sorting algorithms
- Insertion sort

## A sorting problem



What would be your algorithm?

### **Algorithm**

#### Webster's defintion:

- A procedure for solving a mathematical problem (as of finding the greatest <u>common divisor</u>) in a finite number of steps that frequently involves repetition of an operation;
- broadly: a step-by-step procedure for solving a problem or accomplishing some end especially by a computer

### An algorithm should be described..

- Unequivocally
- Concisely
- With a sufficient level of details for the target reader.
  - e.g. "Pick the smallest book" may or may not be enough:

### Some high level ideas for sorting:

 Seek the smallest book → place in the left → Seek the next smallest (excluding the leftmost) → place in the second left → ....

 Keep the first k elements sorted → Place the k+1-th element in the increasing order → (repeat)

 Keep switching between adjacent elements whenever left is greater than right. → Repeat the procedure n times

What if you have thousands of books?

### Some high level ideas for sorting:

Seek the smallest book → place in the left → Seek the next smallest (excluding the leftmost) → place in the second left → ....
 Selection sort

 Keep the first k elements sorted → Place the k+1-th element in the increasing order → (repeat)

: Insertion sort

 Keep switching between adjacent elements whenever left is greater than right. → Repeat the procedure n times
 : Bubble sort

What if you have thousands of books?

### Visual illustration of insertion sort

https://www.toptal.com/developers/sorting-algorithms/insertion-sort

### **Insertion sort**

• Key idea: Make sure that the first k elements being sorted, and increase k stepwise.

#### Procedure

- When k = 1, the first 1 element is automatically sorted.
- When first k elements are sorted, the (k+1)-th elements can be placed in one of the (k+1) spots.
  - : Keep switching the element with the previous element as long as the previous element is larger.
- Then the first (k+1) elements are sorted.
- By induction, repeating the procedure n times will sort the entire list.

### Computer-friendly description via pseudocodes

```
Data: An unsorted list A[1 \cdots n]
Result: The list A[1 \cdots n] is sorted
for j = 2 to n do
   key = A[j];
   i = j - 1;
   while i > 0 and A[i] > key do
   A[i+1] = A[i];
i = i-1;
    A[i+1] = key;
end
```

# Jupyter notebook:

insertion\_sort.ipynb

### How long does an insertion sort take?

```
for j=2 to n
                                                                   c_1 n
do
    key = A[j];
                                                                  c_2(n-1)
    i = j - 1;
                                                                   c_3(n-1)
                                                                   c_4 \sum_{i=2}^n j
    while i > 0 and A[i] > key
     do
     A[i+1] = A[i];
i = i-1;
                                                                  c_5 \sum_{j=2}^{n} (j-1)
c_6 \sum_{j=2}^{n} (j-1)
     end
     A[i+1] = key;
                                                                   c_7(n-1)
end
```

# How long does an insertion sort take?

$$T(n) = (c_4 + c_5 + c_6)n^2 + \frac{2(c_1 + c_2 + c_3 + c_7) + c_4 - c_5 - c_6}{2}n$$

$$-(c_2 + c_3 + c_4 + c_7)$$

A[i+1] = key;

 $c_7(n-1)$ 

## How long does an insertion sort take?

$$T(n) = (c_4 + c_5 + c_6)n^2$$

$$+ \frac{2(c_1 + c_2 + c_3 + c_7) + c_4 - c_5 - c_6}{2}n$$

$$- (c_2 + c_3 + c_4 + c_7)$$

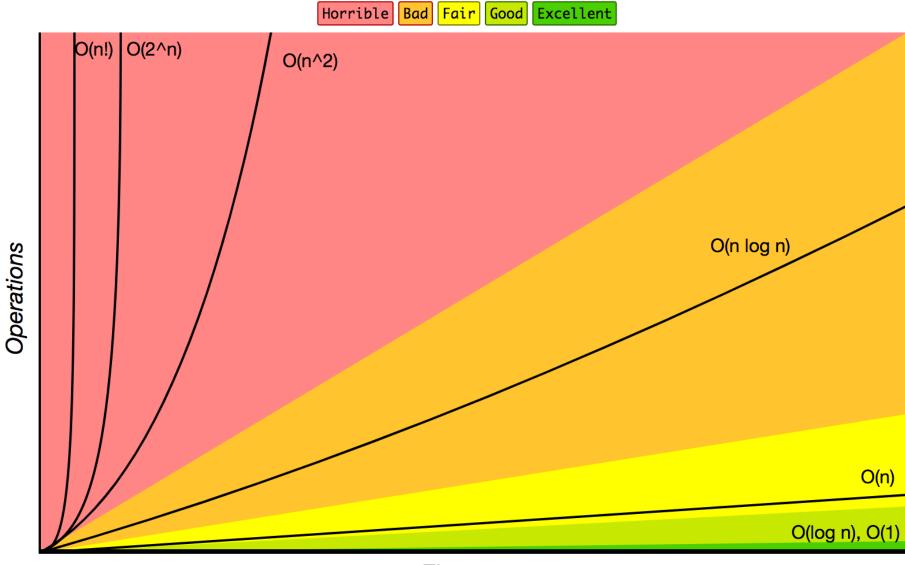
$$= \Theta(n^2)$$

$$+ \frac{1}{2} (n-1)$$

end

### **Time complexity: Notations**

- $O(\cdot)$ : Time complexity is AT MOST  $c(\cdot)$  as n increases
  - Formally, T(n) = O(f(n)) if and only if there exists  $n_0 > 0$  and M > 0 such that  $T(n) \le Mf(n)$  for all  $n > n_0$
- $\Omega(\cdot)$ : Time complexity is AT LEAST  $c(\cdot)$  as n increases.
  - Formally,  $T(n) = \Omega(f(n))$  if and only if there exists  $n_0 > 0$  and M > 0 such that  $T(n) \ge Mf(n)$  for all  $n > n_0$
- $\Theta(\cdot)$ : Time complexity is EXACTLY  $c(\cdot)$  as n increase.
  - $T(n) = \Theta(f(n))$  if and only if T(n) = O(f(n)) and  $T(n) = \Omega(f(n))$



**Elements** Image: http://bigocheatsheet.com

## Q: What would be time complexity of...

- Finding an element from an array?
- Finding an element from a sorted array?
- Inserting an element to an array?
- Inserting an element to a sorted array?
- Finding median value from a sorted array?

## **Summary: Simple Sorting**

- Sorting algorithms
  - Selection sort
  - Insertion sort
  - Bubble sort

- Time complexity
  - Insertion sort is  $\Theta(n^2)$
  - Scalable algorithms should be at most  $\Theta(n \log n)$
- Would it be possible to sort faster than  $\Theta(n^2)$ ?

### **Reading List**

• CLRS I.1 – I.3 (pp. 5-65)