

## CSE208: Data Structures and Algorithms II Sessional

### Online week 5: Generalized shortest path (B1/B2)

Time: 35 minutes

*Generalized shortest-paths problem.* In Internet routing, there are delays on lines but also, more significantly, delays at routers. This motivates a generalized shortest-paths problem.

Suppose that in addition to having edge lengths  $\{l_e : e \in E\}$ , a graph also has *vertex costs*  $\{c_v : v \in V\}$ . Now define the cost of a path to be the sum of its edge lengths, *plus* the costs of all vertices on the path (including the endpoints). Give an efficient algorithm for the following problem.

*Input:* A directed graph  $G = (V, E)$ ; positive edge lengths  $l_e$  and positive vertex costs  $c_v$ ; a starting vertex  $s \in V$ .

*Output:* An array  $\text{cost}[\cdot]$  such that for every vertex  $u$ ,  $\text{cost}[u]$  is the least cost of any path from  $s$  to  $u$  (i.e., the cost of the cheapest path), under the definition above.

Notice that  $\text{cost}[s] = c_s$ .

**Input:** The first line of the input file will contain the number of vertices  $n$  ( $\leq 1000$ ) and the number of edges  $m$  ( $\leq 10000$ ) followed by  $n$  lines containing vertex  $u$  and cost  $c$ , and then  $m$  lines each containing origin  $u$ , end  $v$  and weight  $w$  ( $\leq 100000$ ) of an edge of the directed graph. The last line will contain a source vertex  $s$  and a destination vertex  $d$ .

Sample input and output:

5 10 0 5 1 3 2 4 3 10 4 6 0 1 4 0 2 2 1 2 3 1 3 2 1 4 3 2 1 1 2 4 5 2 3 4 3 4 5 4 3 1 0 4	Shortest path cost: 21 0 -> 1 -> 4
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