Algorithms II Cheat Sheet

Tips

Apply an algorithm you know in a clever way, don't write a new algorithm.

Set notation

$A \in [10] \equiv A \in [1..10]$

Big O

Notation	Intuitive meaning	Analogue
$f(n) \in O(g(n))$	f grows at most as fast as g	$\leq$
$f(n) \in \Omega(g(n))$	f grows at least as fast as g	$\geq$
$f(n) \in \Theta(g(n))$	f at the same rate as g	$=$
$f(n) \in o(g(n))$	f grows strictly less fast than g	$<$
$f(n) \in \omega(g(n))$	f grows strictly faster than g	$>$

Notation	Formal definition
$f(n) \in O(g(n))$	$\exists C,\ n_0 \colon \forall n \geq n_0 \colon f(n) \leq C \cdot g(n)$
$f(n) \in \Omega(g(n))$	$\exists c,\ n_0 \colon \forall n \geq n_0 \colon f(n) \geq c \cdot g(n)$
$f(n) \in \Theta(g(n))$	$\exists c,\ C,\ n_0 \colon \forall n \geq n_0 \colon c \cdot g(n) \leq f(n) \leq C \cdot g(n)$
$f(n) \in o(g(n))$	$\forall C \colon \exists n_0 \colon \forall n \geq n_0 \colon f(n) \leq C \cdot g(n)$
$f(n) \in \omega(g(n))$	$\forall c \colon \exists n_0 \colon \forall n \geq n_0 \colon f(n) \geq c \cdot g(n)$

Interval Scheduling

A **request** is a pair of integers (s, f) with $0 \leq s \leq f$.
We call s the **start time** and f the **finish time**.

A set A of requests is **compatible** if for all distinct $(s, f), (s', f') \in A$, either $s' \geq f$ or $s \geq f'$ — that is, the requests' time intervals don't overlap.

Interval Scheduling Problem
Input: An array $\mathcal{R}$ of n requests $(s_1, f_1), \ldots, (s_n, f_n)$.
Desired Output: A compatible subset of $\mathcal{R}$ of maximum possible size.

Algorithm: GreedySchedule
Input: An array $\mathcal{R}$ of n requests.
Output: A maximum compatible subset of $\mathcal{R}$.

```
1 begin
2     Sort R's entries so that R ← [(s_1, f_1), ..., (s_n, f_n)] where f_1 ≤ ··· ≤ f_n.
3     Initialise A ← [], lastf ← 0.
4     foreach i ∈ {1, ..., n} do
5         if s_i ≥ lastf then
6             Append (s_i, f_i) to A and update lastf ← f_i.
7     Return A.
```

Complexity:

Step 2 takes O(n log n)
Steps 3–6 all take O(1) time and are executed at most n times.
$\therefore$ $totalrunningtime = O(nlogn) + O(n)O(1) = O(nlogn)$.

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interval scheduling proofs

ODEs

1st Order Linear	Use integrating factor, $I = e^{\int P(x)dx}$
Separable:	$\int P(y)dy/dx = \int Q(x)$
HomogEnEous:	$dy/dx = f(x, y) = f(xt, yt)$ sub $y = xV$ solve, then sub $V = y/x$
Exact:	If $M(x, y) + N(x, y)dy/dx = 0$ and $M_y = N_x$ i.e. $\langle M, N\rangle = \nabla F$ then $\int_x M + \int_y N = F$
Order Reduction	Let $v = dy/dx$ then check other types *If purely a function of y,* $\frac{dv}{dx} = v\frac{dv}{dy}$
Variation of Parameters:	When $y'' + a_1 y' + a_2 y = F(x)$ F contains $\ln x$, $\sec x$, $\tan x$, $\div$
Bernoulli	$y' + P(x)y = Q(x)y^n$ $\div y^n$ $y^{-n}y' + P(x)y^{1-n} = Q(x)$ *Let* $U(x) = y^{1-n}(x)$ $\frac{dU}{dx} = (1-n)y^{-n}\frac{dy}{dx}$ $\frac{1}{1-n}\frac{du}{dx} + P(x)U(x) = Q(x)$ *solve as a 1st order*
Cauchy-Euler	$x^n y^n + a_1 x^{n-1} y^{n-1} + \cdots + a_{n-1} y^{n-2} + a_n y = 0$ guess $y = x^r$
3 Cases:	
1) Distinct real roots	$y = ax^{r_1} + bx^{r_2}$
2) Repeated real roots	$y = Ax^r + y_2$ *Guess* $y_2 = x^r u(x)$ *Solve for* $u(x)$ *and choose one* $(A = 1, C = 0)$
3) Distinct complex roots	$y = B_1 x^a \cos(b \ln x) + B_2 x^a \sin(b \ln x)$

Laplace Transforms

$L[f](s) = \int_0^\infty e^{-sx} f(x)dx$

$f(t) = t^n, n \geq 0$	$F(s) = \frac{n!}{s^{n+1}}, s > 0$
$f(t) = e^{at}, a$ *constant*	$F(s) = \frac{1}{s-a}, s > a$
$f(t) = \sin bt, b$ *constant*	$F(s) = \frac{b}{s^2+b^2}, s > 0$
$f(t) = \cos bt, b$ *constant*	$F(s) = \frac{s}{s^2+b^2}, s > 0$
$f(t) = t^{-1/2}$	$F(s) = \frac{\pi}{s^{1/2}}, s > 0$
$f(t) = \delta(t - a)$	$F(s) = e^{-as}$
f'	$L[f'] = sL[f] - f(0)$
f''	$L[f''] = s^2 L[f] - sf(0) - f'(0)$
$L[e^{at} f(t)]$	$L[f](s - a)$

Vector Spaces

$v_1, v_2 \in V$

1. $v_1 + v_2 \in V$
2. $k \in \mathbb{F}, kv_1 \in V$
3. $v_1 + v_2 = v_2 + v_1$
4. $(v_1 + v_2) + v_3 = v_1 + (v_2 + v_3)$
5. $\forall v \in V, 0 \in V \mid 0 + v_1 = v_1 + 0 = v_1$
6. $\forall v \in V, \exists -v \in V \mid v + (-v) = (-v) + v = 0$
7. $\forall v \in V, 1 \in \mathbb{F} \mid 1 * v = v$
8. $\forall v \in V, k, l \in \mathbb{F}, (kl)v = k(lv)$
9. $\forall k \in \mathbb{F}, k(v_1 + v_2) = kv_1 + kv_2$
10. $\forall v \in V, k, l \in \mathbb{F}, (k + l)v = kv + lv$