

# DIJKSTRA

## Part 1: What problem does Dijkstra solve?

Dijkstra helps you find:

1. The **shortest distance** from a starting node to all other nodes
2. Works only when weights are **positive**
3. Used in **Google Maps, network routing, games**, etc.

Example:

If cities are nodes and road distances are weights, Dijkstra tells you the **minimum travel distance** from your city to all other cities.

## Part 2: Main Idea

Think of the algorithm like this:

1. You start at a node with distance **0**
2. All other distances are  $\infty$  (**infinity**)
3. You always go to the **closest unvisited node**
4. Update (relax) the distances of its neighbors
5. Repeat until all nodes are visited
6. It's like expanding outward from the source step-by-step, always choosing the cheapest next path.

## Part 3: Key Concepts

### 1. Distance Array

Stores the shortest distance found so far:

$\text{dist}[\text{node}] = \text{minimum cost from source to node}$

## 2. Visited Set

Marks nodes whose shortest path is already fixed.

## 3. Priority Queue (Min-Heap)

Always picks the node with the **smallest distance** next.

## Part 4: Step-by-Step Example

Suppose edges:

A --5--> B   A --2--> C  
C --1--> D   B --3--> D

Start from A.

Steps:

1. Initially

$\text{dist}[A] = 0$

$\text{dist}[B] = \infty$

$\text{dist}[C] = \infty$

$\text{dist}[D] = \infty$

2. Pick A (distance 0)

Update:

$B = 5$

$C = 2$

3. Pick C (distance 2)

Update:

$D = 2 + 1 = 3$

4. Pick D (distance 3)

Update:

$B = \min(5, 3+3=6) \rightarrow \text{still}$

5. Pick B (distance 5)

Done.

Shortest distances:

A=0, C=2, D=3, B=5.

### Part 5: Why is this algorithm correct?

Because **when you pick the smallest unvisited distance**, it is always the final shortest path (since weights are positive).

### Part 6: Dijkstra Pseudocode

Dijkstra(Graph, source):

For each node:

$\text{dist}[\text{node}] = \text{infinity}$

$\text{dist}[\text{source}] = 0$

Create a min-priority-queue pq

pq.push( (0, source) )

While pq is not empty:

$(d, u) = \text{pq.pop}()$

If u is already visited:

continue

Mark u as visited

For each neighbor v of u with edge weight w:

If  $\text{dist}[u] + w < \text{dist}[v]$ :

$\text{dist}[v] = \text{dist}[u] + w$

$\text{pq.push}(\text{dist}[v], v)$

### Part 7: What Dijkstra gives you

- 1.Shortest distance from source to all nodes
- 2.You can reconstruct shortest paths too
- 3.Works fast using priority queue

### 8.Dijkstra Diagram

(5)

A ----- B

|       |

(2)    (3)

|       |

C ----- D

(1)

### Edges with weights:

$$A \rightarrow B = 5$$

$$A \rightarrow C = 2$$

$$C \rightarrow D = 1$$

$$B \rightarrow D = 3$$

### How Dijkstra explores (step-by-step visually)

#### Start at A (distance = 0)

$$A(0) \quad B(\infty) \quad C(\infty) \quad D(\infty)$$

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#### Step 1: From A, update neighbors

$$A(0) \rightarrow B = 5$$

$$A(0) \rightarrow C = 2$$

$$A(0) \quad B(5) \quad C(2) \quad D(\infty)$$

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#### Step 2: Pick the smallest distance $\rightarrow C(2)$

Update neighbors of C:

$$C(2) \rightarrow D = 2 + 1 = 3$$

$$A(0) \quad B(5) \quad C(2) \quad D(3)$$

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#### Step 3: Pick next smallest $\rightarrow D(3)$

Update neighbors:

$$D \rightarrow B = 3 + 3 = 6 \text{ But B already has 5} \rightarrow \text{keep 5}$$

No change:

$$A(0) \quad B(5) \quad C(2) \quad D(3)$$

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**Step 4: Pick next smallest → B(5)**

No more updates.

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**★ Final Shortest Distances**

A = 0 C = 2 D = 3 B = 5