CS416

Introduction to Computer Science II Spring 2018

Previously in 416

- Data structures
 - Organizing collections of related data objects
 - Linear organizations: arrays, lists, stacks, queues
 - Nonlinear: hash tables, trees

09 Regular Expressions

Preview

- String processing
 - defining string patterns
 - matching patterns
 - pattern substitution
- String processing applications

Preview

- String processing
 - defining string patterns
 - matching patterns
 - pattern substitution
- String processing applications

Regular Expressions

- A regular expression defines a string pattern
 - a pattern can be used to define a valid Java identifier: an unlimited-length sequence of letters, digits, the dollar sign "\$", or the underscore character "_"; the sequence must start with a letter.
- Need a notation (actually it's a simple language) to define patterns such as this.
- Need a notation/interface for using the patterns

Regular Expression Processor

- A regular expression processor:
 - takes as input a string and a pattern
 - decides if/where the pattern is found in the string
 - can also <u>replace</u> the piece of the string that matches the pattern with another string
 - patterns can have multiple parts, which can be used as part of the replacement string
 - Example: convert a list of names from "Last, First" to "First Last"

String: String to be searched

Pattern: to be

String: String to be searched

Replace: being

String: String being searched

Pattern: (Last), (First)

Part 1: Last

Part 2: First

Replace: First Last

The Pattern Class

- Java's *Pattern* class parses regular expressions
- Simplest invocation is a *static* method:
 - boolean Pattern.matches (String regex, String test);

Requires that the pattern match the *entire* string.

```
// the "or" can be in regex: more efficient
if ( Pattern.matches( "Durham|Dover", town )
```

String also has a matches method:

```
if ( town.matches( "Durham" ))
```

"Wildcard" matches

- Suppose we don't have a separate *town* field, but just have a string containing the entire address.
 - We want to see if "Durham" appears <u>anywhere</u> in the address field
 - * in a regex matches 0 or more of <u>any</u> character

```
// check if address is in Durham or Dover
String addr = getNextAddress();

if ( addr.matches( ".*Durham.*" )
{
   System.out.println( "Found: " + addr );
}
```

Pattern matches any number of any characters, followed by "Durham", followed by any number of any characters.

```
// can also write:
if ( addr.matches( ".*Durham.*|.*Dover.*" )
```

How does it work?

- To understand how to use regular expressions well, need to understand the basic matching model, especially for pattern *replications* (wildcards).
- Demonstrate by over-simplified example. There are 2 key concepts:
 - replication matching is greedy; it tries to match as much as possible
 - subsequent match failure causes algorithm to <u>backup</u> and try a less greedy match, if possible, and then go forward again.

Step	Comment	Pattern	Input String	Repl#
1	Match .* greedily	.*Durham.*	1 Main St. Durham, NH	21
2	Match Durham fails; match .* less greedily	.*Durham.*	1 Main St. Durham, NH	20
3	Match Durham fails; match .* less greedily	.*Durham.*	1 Main St. Durham, NH	19
• • •	18, 17, , 12	.*Durham.*	1 Main St. Durham, NH	12
11	Match Durham fails; match .* less greedily	.*Durham.*	1 Main St. Durham, NH	11
12	Match Durham succeeds	.*Durham.*	1 Main St. Durham, NH	1
13	Match .* greedily	.*Durham.*	1 Main St. Durham, NH	4
	Entire pattern matches!	Code is much s	marter, but results must follow t	this mod

String Patterns

- What kinds of patterns might we want to define?
 - We surely want to specify how <u>characters</u> in the string match <u>characters</u> in the pattern
 - exact character match: match the letter "A"
 - subset match: match any letter, or match any digit: [a-z] or [0-9]
 - "anything but" match: any character that is not a space: $[^{\square}]$

 \Box represents a space in the notes.

- wildcard match: match any character:
- We might want to specify how many *replications* of a particular match are part of the pattern
 - 0 or more, 1 or more, 3 or more, between 2 and 5, etc.

Simple Examples

- . means any character
- * means 0or morereplicationsof preceding
- + means 1or morereplicationsof preceding

Pattern	Meaning	Examples	
AAT	AAT	AAT	
A.T	A then any 1 char then T	A <u>A</u> T, A <u>T</u> T, A <u>G</u> T, A <u>C</u> T,	
AAT*	AA then 0 or more T's	AA, AAT, AATT, AATTT,	
AAT+	AA then 1 or more T's	AAT, AATT, AATTT,	
T+A*GG	1 or more T, then 0 or more A, then GG	TTAAAGG, TTTTTAGG, TTGG, TGG,	
AA.+T	AA then 1 or more of anything then T	AAGT, AACGT, AAGAAT, AATT, AATTT, AATTTT,	

These show greedy matching

Grouping

- (...) groups regular expression components
 - can specify a replication on groups

Pattern	Meaning	Examples	
(AT)+ 1 or more AT		AT, ATAT, ATATAT,	
TT(AT)*(CG)*	TT then 0 or more AT then 0 or more CG	TT, TTAT, TTCG, TTATCG, TTATATATCG, TTCGCGCG,	
1 or more of (AT)*(CG)*; ((AT)*(CG)*)+ i.e., 1 or more AT, CG pairs in any order		AT, CG, CGAT, ATCG, ATATATCG, CGATATCGAT,	

Character Classes

Character class version is more efficient

- [...] defines a *character class*
 - a class represents one character that can be any class member
- defines a range of characters: a-z or 0-9

Pattern	Meaning	Examples	
[ACTG] A or C or T or G		A, C, T, G	
(A T G C)	A or C or T or G	A, C, T, G	
[ACTG][ACTG]	(A or C or T or G) then [SEP](A or C or T or G)	AA, AC, AT, AG, CA, CC, CT, CG	
[AC]*[GT]+ 0 or more of (A or C) then 1 or more of (G or T)		ACAAGG, ACTT, AACGGG, AACG	
[a-z]+ 1 or more lower case let		able, ace, abcdefghijklmnopqrstuvwxz	
[a-zA-Z]+ 1 or more letters		Able, aCe, abcdefghIJklmnopqrStuvWxz	

Alternatives:

- | (or) defines alternative patterns
 - the test string matches the (piece of) the pattern if one of the alternatives matches

Pattern	Meaning	Examples	
any order 1 or more AT, CG pairs		AT, CG, CGA, ATCC, CATA	
		AT, CG, CGAT, ATCG, ATATATCG, CGATATCGAT	
((AT)*(CG)*)+	1 or more of (AT)*(CG)*; i.e., 1 or more AT, CG pairs in any order	AT, CG, CGAT, ATCG, ATATATCG, CGATATCGAT	
ATG(A T C G)+(TAG TGA)	start codon, then any dna, then a stop codon ATGAAT		

Character Classes with [^]

- [^ ...] defines a *character class* by what is <u>not</u> in it
 - all characters are in the class except the ones listed

□ symbol represents a space.

Pattern	Meaning	Examples	
[^a-z]	All characters except lower case letters	A, C, T, G, %, 1, #, 7	
[^ACTG][ACTG] anything but A, C, T, then (A or C or T or		aA, rG, 8T, tC, @A	
[□]*[^□]+	0 or more blanks then 1 or more non-blanks	$\Box Abc$, Abc , $\Box \Box \Box 2dfbp/$	
□*[^□]+	Same as above	$\Box Abc$, Abc , $\Box \Box \Box 2dfbp/$	
[^a-z]+	strings with no l.c. letters	ABLE, A1, A+B=123, ^RT^,	
[^0-9a-z]+	strings with no l.c. or digits	ABLE, A#, A+B=EE, ^RT^,	

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Note: ^ is not special except right after [

Metacharacters and Escapes

- . + * [and] have special meaning in regular expressions; they are called <u>metacharacters</u>
- Metacharacters are: . + * [] () { } \$ \ ? | ^
- To use a metacharacter in a pattern, precede it with a \; \ is called the *escape* character

Pattern	Meaning	Examples
$(\Box) +$ string of $\Box $, and *		***
[□ *]+•	string of □ , and *	***
[ATCG\-]+	DNA sequence with missing nucleotides	ATC-GTACGGCAATCG

Escape character use is context dependent.

The only "real" metacharacter for character groups is -.

Predefined Character Classes

• There are a bunch of predefined character classes represented by \char where char is a single letter

Predefined class	Equivalent	Meaning	
\d	[0-9]	A digit character	
\D	[^0-9]	A non-digit character	
\s	$[\Box \backslash t \backslash n \backslash r \backslash f]$	"white" space: blank, tab (\t), new line or line feed (\n), carriage return (\r), new page or form feed (\f)	
\S	$[\neg \neg t \ n \ f]$	A non-white space character	
\w	[a-zA-Z_0-9]	A word character	
\W	[^a-zA-Z_0-9]	A non-word character	

Examples with Predefined Character Classes

Pattern	Meaning	Examples	
[\w\$]	word characters plus \$;[sep]valid characters in Java id	a, b, c, z, 0, 1, \$, _	
	valid Java identifier	Nim, _root, temp1, aPlayer	
\s*\S+\s+\S+	optional white space (\s*) then a token (\S+) then required white space (\s+) and another token (\S+)	print□name ++□id%D	
\s*\w+\s+\w+	same as above, except the tokens can only be composed of word characters	print□name add□name2	

Boundary Markers

• Can match a boundary between 2 characters, rather than the character itself

† Not really true in

multi-line mode

- ^ matches the start of the test string†
- \$ matches the end of the test string†
- \b matches a word boundary

Pattern	Meaning	Examples
^\s*\S+\s+\S+\s*\$	optional space then token then space then token then space then end of input	print□name ++□id%D
^\s*\w+\s+\w+\s*\$	same as above, except the tokens can only be composed of word characters	print□name add□name2

Quantifiers

- Replication specifications are called *quantifiers*
 - *Greedy* quantifiers keep trying to match until the match fails: return the match that finds the first occurrence of the pattern and that uses the *maximum* possible amount of the input
 - * 0 or more with a *greedy* match
 - 1 or more with a *greedy* match
 - ? 0 or 1 (an optional piece of the pattern)
 - $\{n\}$ exactly n replications
 - $\{n,m\}$ at least n, no more than m
 - $\{m,m\}$ no more than m (can be 0)
 - $\{n,\}$ at least n

Quantifier Examples

Pattern	Meaning	Examples	
(AT){6,}	6 or more AT: an AT <i>microsatellite</i>	ATATATATATATATAT	
ATG[ATCG]{3}+(TAG TGA)	start codon (ATG), then any number of codons (3 nucleotides), then a stop codon (TAG or TGA) A TGAATCGATCCG A A A A A A A A A A A A A A A A A A A		
(\(\d\d\d\))?\d\d\d\-\d\d\d\d (\(\d{3}\))?\d{3}\-\d{4}	US phone number format with optional area code in ()	862-3780, (603)862-1234	

Two different regular expressions for the same logical pattern

Underlining used to show groups of 3 nucleotides

Reluctant Quantifiers

• *Reluctant* quantifiers stop trying to match as soon as the entire pattern is matched: try to match the first occurrence of the pattern with the *minimum* amount of input

```
• *? 0 or more with a reluctant match
```

- +? 1 or more with a *reluctant* match
- ?? 0 or 1 (an optional piece of the pattern)
- $\{n\}$? exactly *n* replications (same match as greedy)
- $\{n,m\}$? at least n, no more than m
- $\{,m\}$? no more than m (can be 0)
- $\{n,\}$? at least n

Reluctant v. Greedy Quantifiers

Pattern	Input	Match	Notes
TA*	GGGTAAAAC	TAAAA	As many A as can
TA*?	GGGTAAAAC	Т	0 A
(AT)+	GGATATATACTT	ATATAT	As many AT as can
(AT)+?	GGATATATACTT	AT	1 AT
$(GC)\{2,5\}$	AGCTGCGCGCGCAAA	GCGCGCGC	4 GC
$(GC)\{2,5\}$?	AGCTGCGCGCAAA	GCGC	2 GC
(GC){5}	AGCTGCGCGCGCAAA	GCGCGCGCAA	5 GC
(GC){5}?	AGCTGCGCGCGCAAA	GCGCGCGCAA	Exactly 5 is always 5
TATAC?	GGACATATACTT	TATAC	1 C
TATAC??	GGACATATACTT	TATA	0 C

Java Regular Expressions

- java.util.regex.Pattern encapsulates the idea of a regular expression and provides some of the functionality associated with matching an input string to the regular expression.
- java.util.regex.Matcher encapsulates the complete matching process and options

Pattern Class

- Create a *Pattern* object with the *static* method, *compile*:
 - Pattern myPat = Pattern.compile(regEx); where regEx can be a *String*, *StringBuffer*, or *CharBuffer*
 - Apply the pattern to an input string as a *separator* for the remaining parts of the input string
 - String[] matches = myPat.split(inputString);
 - if pattern = ":" and
 - inputString = "able: baker: charlie: delta"
 - matches[]: ["able", "baker", "charlie", "delta"]

Matcher

The real power in Java regular expressions comes from the *Matcher* class.

Matcher created by

Matcher created by pattern for a string

```
String dna = readNextSequence();
Pattern orf = Pattern.compile( "atg.{3}*(tag|tga");

Matcher m = orf.matcher( dna );
while ( m.find() )
{
    System.out.println( m.group() );
}
```

First, search from start of input; if match, remember where ended; start next search from there.

The substring matching the pattern

Java Literals

- The Java compiler parses literal strings looking for its own "escape" specifications: \n, \t, \r etc.
 - It has very specific rules about what is allowed to follow \
 - This complicates specification of regular expressions as literal strings
 - Must "double-escape" some stuff.
 - Especially every use of \ has to be doubled to \\
 - So, "([ACTG]{2})\1{5,}" becomes "([ACTG]{2})\\1{5,}"
 - and "(.*),\s*(.*)," becomes "(.*),\\s*(.*),"

Groups

- A group is a portion of the pattern in parentheses ()
- Each group has a unique index (n) based on <u>counting left</u> parentheses from the left of the pattern starting at 1
- As the *regex processor* is trying to match a string to the pattern, it saves the string matched so far by each *group*
- These are called *capturing* groups

Groups

```
String in = "Smith, John, 123 Thayer.";
  Pattern names = Pattern.compile("(.*), \s*(.*),");
  dMatcher m = names.matcher( in );
                                                     Entire match
  while ( m.find() )
     System.out.println( m.group() );
                                                      group 1: Smith
     System.out.println( m.group( 1 ) );
     System.out.println( m.group( 2 ) );
System.out
Smith, John,
                                           group 2: John
Smith
John
```

Pattern Match Substitution

- •Can replace the matched region with something else
 - actually, copy preceding unmatched region to StringBuffer followed by replacement for matched region

```
# replace all occurrences of "hot" with "pot"
String in = "the hot shot had the first shot."
Pattern hot = Pattern.compile( "hot");
StringBuffer out = new StringBuffer();
Matcher m = hot.matcher( in );
while ( m.find() )
{
    m.appendReplacement( out, "pot" );
}
m.appendTail( out );
```

Copies unmatched string from search start to start of match, then copies the replacement ("pot") for the match.

Copies what follows last match

Pattern Match Substitution Example

```
String in =
    "the hot shot had the first shot."
Pattern hot = Pattern.compile( "hot" );
StringBuffer out = new StringBuffer();
Matcher m = hot.matcher( in );
while ( m.find() )
{
    m.appendReplacement( out, "pot" );
}
m.appendTail( out );
```

no match	match	replace
the□	hot	pot
□S	hot	pot
□had the first s	hot	pot
•		

in:

out: |the | pot | spot | had | the | first | spot |.

Groups with Replacement

```
String in = "Smith, John, 123 Thayer.";
 Pattern names = Pattern.compile( "(.*), \\s*(.*)," );
 StringBuffer out = new StringBuffer();
 Matcher m = names.matcher( in );
 while ( m.find() )
                                                       Entire match
    System.out.println( m.group() );
    System.out.println( m.group( 1 ) ←);
                                                        group 1: Smith
    System.out.println( m.group( 2 ) );
    m.appendReplacement(out, "$2 $1,");
 m.appendTail( out );
 System.out.println( out )
System.out
                                           group 2: John
Smith, John,
Smith
John
John Smith, 123 Thayer.
```

Other Matcher Methods

Some other useful Matcher methods:

int start(): index of start of match

int end(): index of 1 character after end of last match

int start(int group): index of start of a group

int end(int group): index of 1 char after end of a group

int groupCount(): number of groups matched

• There are more!

Match Flags

- Flags can control some match characteristics:
 - Pattern.DOTALL normally . does not match line feeds
 - Pattern.MULTILINE treat each line in input as separate input -- (^ matches start of line, \$ matches end of line)
 - Pattern.CASE_INSENSITIVE ignore case differences
 - Pattern.COMMENTS white space is ignored in the pattern (all pattern white space has to be done with \s) and comments are allowed.

Pattern.compile(regex, Pattern.DOTALL | Pattern.MULTILINE)

Comments

- Comments and spaces can be included in patterns
 - Must specify Pattern.COMMENTS "flag"
 - Comments start with "#" and end with "\n".
 - All white space and comments are deleted before compiling

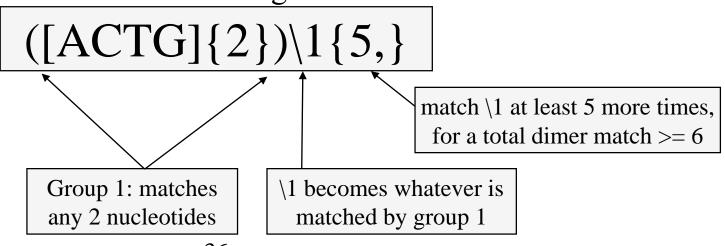
```
Pattern.compile( "ATG([ATCG]{3})+(TAG|TGA)" );
```

Back Referencing

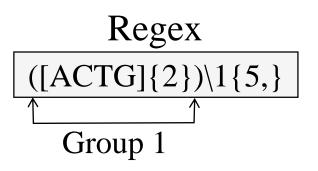
- A regular expression can include components that are based on the input being processed.
- This is done by back referencing
 - As the *regex processor* is trying to match a string to the pattern, it saves the string matched so far by each *group*
 - The pattern can include a reference to a previous group's matching string using the group's id, 1 or 2 or 3, etc.

Back Reference Example

- *Microsatellites* are regions in DNA that contain 6 or more exact copies of some sequence of nucleotides (A, C, T, G)
- a dimer microsatellite is 6 or more replications of 2 nucleotides
 - We could use a brute force r.e.: (AC|AT|AG|CT|CG|TG){6,}
 - this does not look for CA, TA, GA, TC, GC, or GT.
 - the approach would be impossible for longer base patterns.
 - Instead, we'll use back referencing



Dimer Matching Overview



Input

AT**CACACACACACA**AGCTAGC

- Regex processor key steps:
 - match group 1 to AT; this becomes \1
 \1 (AT) does **not** match CA, so <u>fail</u>
 - match group 1 to TC, which becomes \1
 \1 (TC) does not match AC, <u>fail</u>

A**TC**ACACACACACAAGCTAGC

ATCACACACACACAAGCTAGC

AT**CA**CACACACACAAGCTAGC

a complete match to the pattern

ATCACACACACACAAGCTAGC

Dimer Matching Detail

Regex

 $([ACTG]{2})\1{5,}$

Input

ATCACACACACACACACA
AGCAGCAGCAAGCAACA
ACAACAACAGCACGTCGATG
AAGGAAGTCATAGCAGTTTC
AGCGACGCAGCGCCCTTAA
TTTAGCAGAACGGTCGCCTC

Pattern [SEP] Position	Pattern Char	Input [SEP] Position	Action
([ACTG]{2})\1{5,}	[ACTG]	<u>A</u> TCACACA	match A (1 of 2)
([ACTG]{2})\1{5,}	[ACTG]	A <u>T</u> CACACA	match T(2 of 2); group 1 matches AT; 1 = AT
([ACTG]{2})\ <u>1</u> {5,}	<u>A</u> T	<i>ат</i> <u>C</u> ACACA	C fails to match A, backup
([ACTG]{2})\1{5,}	[ACTG]	A <u>T</u> CACACA	match T (1)
([ACTG]{2})\1{5,}	[ACTG]	ATCACACA	match C(2); group 1 matches; \1=TC
([ACTG]{2})\ <u>1</u> {5,}	<u>T</u> C	ATCACACA	A fails to match T, backup
([ACTG]{2})\1{5,}	[ACTG]	AT <u>C</u> ACACA	match C (1)
([ACTG]{2})\1{5,}	[ACTG]	ATc <u>A</u> CACA	match A(2); group 1 matches; \1=AC

italic smaller font: characters matched so far

Dimer Matching (2)

Regex

 $([ACTG]{2})\1{5,}$

Input

ATCACACACACACACACA
AGCAGCAGCAAGCAACA
ACAACAACAGCACGTCGATG
AAGGAAGTCATAGCAGTTTC
AGCGACGCAGCGCCCTTAA
TTTAGCAGAACGGTCGCCTC

Pattern [SEP] Position	Pattern Char	Input Position	Action
([ACTG]{2})\1{5,}	[ACTG]	ATc <u>A</u> CACA	match A(1); group 1 matches; \1=CA
([ACTG]{2})\ <u>1</u> {5,}	<u>C</u> A	AT <i>ca</i> <u>C</u> ACA	match C
([ACTG]{2})\ <u>1</u> {5,}	C <u>A</u>	AT <i>cac<u>A</u>CA</i>	match A, done with \1 match(#1)
([ACTG]{2})\ <u>1</u> {5,}	<u>C</u> A	AT <i>caca<u>C</u>A</i>	match C
([ACTG]{2})\ <u>1</u> {5,}	C <u>A</u>	AT <i>cacac<u>A</u></i>	match A (\1 #2)
([ACTG]{2})\ <u>1</u> {5,}	<u>C</u> A	AT <i>cacaca</i> CAC	match C
([ACTG]{2}) <u>\1</u> {5,}	C <u>A</u>	AT <i>cacacac<u>A</u>C</i>	match A (\1 #3)
([ACTG]{2})\ <u>1</u> {5,}	<u>C</u> A	AT <i>cacacaca</i> <u>C</u>	match C
([ACTG]{2})\ <u>1</u> {5,}	C <u>A</u>	CACACACACACAC	match A (\1 #4)
([ACTG]{2})\ <u>1</u> {5,}	<u>C</u> A	CACACACACACAC	match C
([ACTG]{2})\ <u>1</u> {5,}	C <u>A</u>	CACACACACACAC	match A (\1 #5)

Dimer Matching (3)

Regex

 $([ACTG]{2})\1{5,}$

Input

ATCACACACACACACACA
AGCAGCAGCAAGCAACA
ACAACAACAGCACGTCGATG
AAGGAAGTCATAGCAGTTTC
AGCGACGCAGCGCCCTTAA
TTTAGCAGAACGGTCGCCTC

Pattern [SEP] Position	Pattern Char	Input Position	Action
([ACTG]{2})\ <u>1</u> {5,}	C <u>A</u>	CACACACACACACACA	match A (\1 #5)
([ACTG]{2})\ <u>1</u> {5,}	<u>C</u> A	CACACACACACACACA	match C
([ACTG]{2})\ <u>1</u> {5,}	C <u>A</u>	CACACACACACAC <u>A</u> A	match A (\1 #6)
([ACTG]{2}) <u>\1</u> {5,}	<u>C</u> A	CACACACACACAA	C fails to match A but pattern is matched, no backup

Note that 7 copies of the dimer were matched, even though 6 was the minimum needed. This was a *greedy* match; it kept trying to match the pattern until it failed. You can specify a *reluctant quantifier* that quits matching as soon as pattern is satisfied.

Other Microsatellites

Regex for dimer microsatellite

 $([ACTG]{2})\1{5,}$

• What regular expression would we use to find any *trimer* (3 nucleotide) microsatellite? ([ACTG]{3})\1{5,}

- What about a *tetramer* (4 nucleotides)? ([ACTG]{4})\1{5,}
- One regular expression to find <u>all</u> microsatellites: ([ACTG]{2,})\1{5,}

{2,}: all microsatellites of length greater than 2

Huntington's Disease

- Huntington's disease is a severe neurological disorder
 - symptom onset usually occurs between 30-50 years old
 - deterioration is progressive
 - although the cause is not fully understood, the condition can be identified by an excessive replication count of a *trimer* (3 nucleotides) microsatellite in a gene on chromosome 4
 - the "normal" gene has between 6 and 35 repeats of the trimer, CAG
 - Huntington's sufferers have 36 (or more) repeats
 - What is a regex to detect this?
 - $(CAG){36,}$

String vs. StringBuffer

- Pattern replacement methods use *StringBuffers*, rather than *Strings* because they are <u>much</u> more efficient
 - *String* objects are *immutable*; they cannot be changed

```
String results = "";
String[] parts;
parts = pattern.split( inputString );
for ( int i = 0; i < parts.size; i++ )
  results += parts[ i ] + "\n";</pre>
```

Java object creation is expensive!

Every execution of the loop body means the creation of a **new** *String* object whose value is the concatenation of 3 strings.

StringBuffer objects can be modified, so can do the same operation with **no** object creation, except for 1

StringBuffer.

StringBuffer

- StringBuffers are dynamically sized
 - they have a current length and a current capacity where length <= capacity
- StringBuffer has many variations of the following mutation operations
 - append new characters to the end
 - insert new characters inside the existing string
 - delete characters from the string

StringBuffer append/insert

The *append* and *insert* methods are overloaded so that any kind of primitive or *Object* can be converted to a string representation and added to the end of the current string; if necessary, the *capacity* is increased.

Add to the end of the current string.

```
append( String s );
append( float f );
append( double d );
append( char c );
append( char[] c );
append( byte b );
append( byte b );
append( int i );
append( long l );
append( boolean b );
append( Object o );
// and a couple more
```

Insert converted object to start at position *p*

```
insert( int p, String s );
insert( int p, float f );
insert( int p, double d );
insert( int p, char c );
insert( int p, char[] c );
insert( int p, byte b );
insert( int p, short s );
insert( int p, int i );
insert( int p, long l );
insert( int p, boolean b );
insert( int p, Object o );
// and a couple more
```

Replacement and deletion

- There are several mechanisms for modifying or deleting portions of the current string:
 - deleteCharAt(int position);
 - delete(int start, int end);
 - setCharAt(int position, char newChar);
 - replace(int start, int end, String s);
 - the length of the s does **not** need to match the length of the string being replaced.

StringBuffer search

- *StringBuffer* has search methods:
 - int indexOf(String s)
 - int indexOf(String s, int fromIndex)
 - int lastIndexOf(String s)
 - int lastIndexOf(String s, int fromIndex)
- It does **not** have a *matches*(*String regex*) method

StringBuilder

MacBookPro (2GB) performance:

- StringBuffer is a thread-safe class; you can use it in a multi-threaded environment
 - that makes it more robust, but slower
- If you are using a singlethread implementation, you can use *StringBuilder*; it has the same interface as *StringBuffer*, but is faster, but not thread-safe.

14M appends, 7M deletions, and 7M
indexOf with small strings.

	String	StringBuffer	StringBuilder	
sec	5.25	3.23	2.06	
Without indexOf:				
	Ctuins	String Duffor	Ctain a Duildon	

	String	StringBuffer	StringBuilder
sec	4.74	2.37	1.16

18818 cumulative appends of 60 chars

	String	StringBuffer	StringBuilder
sec	168.43	0.24	0.24
ratio	702	1	1

Review

- Regular expressions are a powerful representation tool
- Java supports full regular expressions
 - character classes, both inclusion and exclusion
 - alternatives (|)
 - replications: *, +, { ... }, ?
 - greedy and non-greedy matching
 - back-referencing patterns
 - substitution

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Sorting and searching algorithms