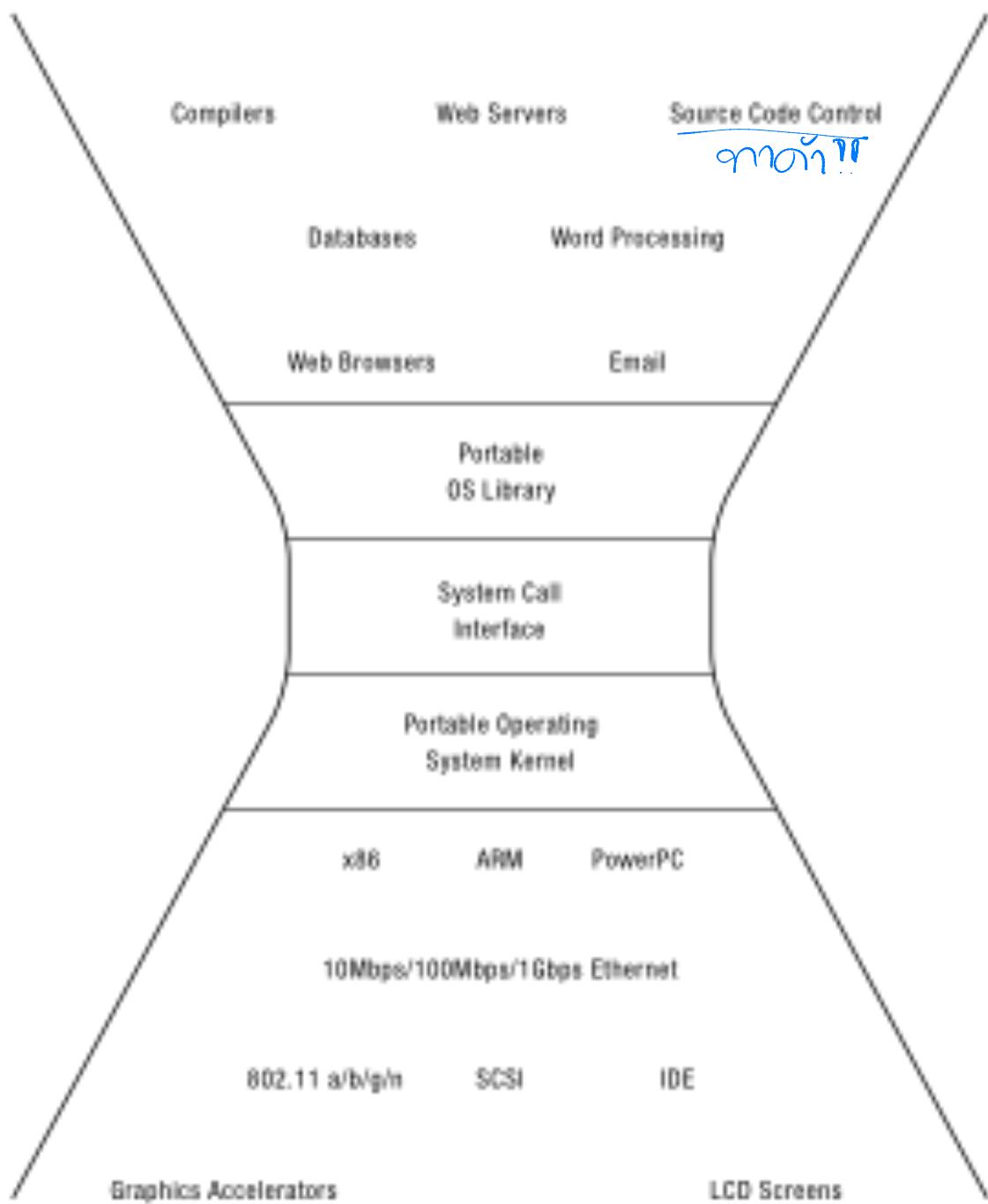


The Process and Programming Interface

ສຳເນົາ
ສຳເນົາ



Main Points

- Creating and managing processes
 - fork, exec, wait
- Performing I/O
 - open, read, write, close
- Communicating between processes
 - pipe, dup, select, connect
- Example: implementing a shell

Shell

→ ສ່ວນ UI

→ ແນວດ
External

→ ໄກສດ

ຄໍາຕັ້ງການ
ໂປຣ ປູມທີ່ກ່ອນນັ້ນ

- A shell is a job control system
 - Allows programmer to create and manage a set of programs to do some task
 - Windows, MacOS, Linux all have shells, Android
luncher
user
Desktop

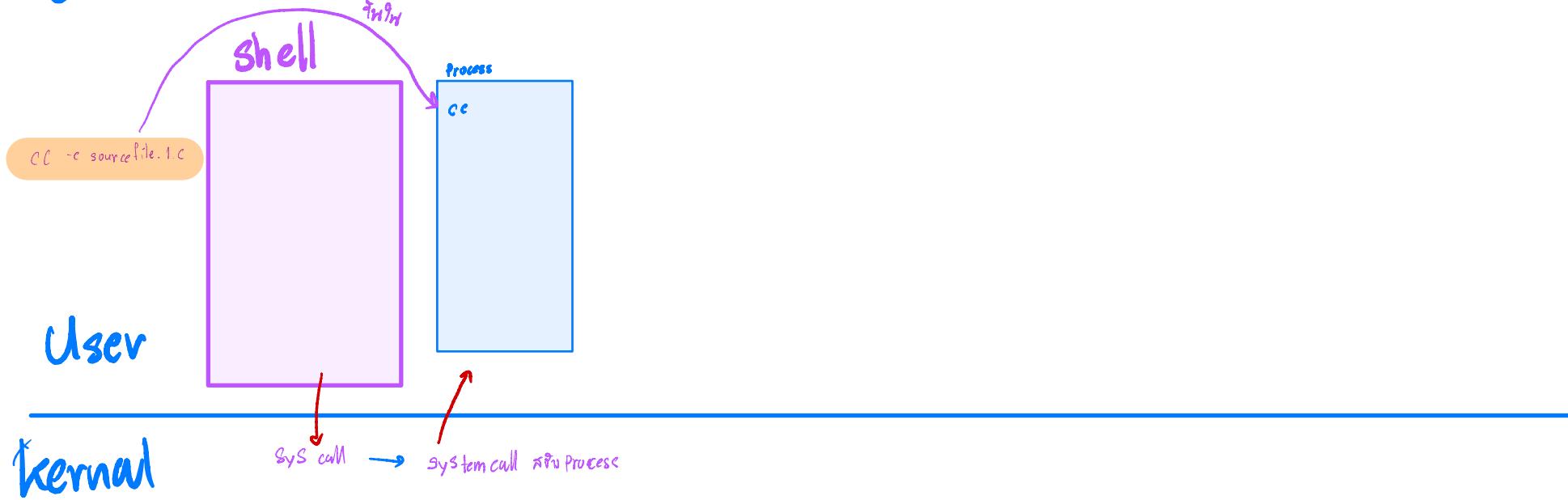
- Example: to compile a C program

cc -c sourcefile1.c

cc -c sourcefile2.c

ln -o program sourcefile1.o sourcefile2.o

Shell និង process



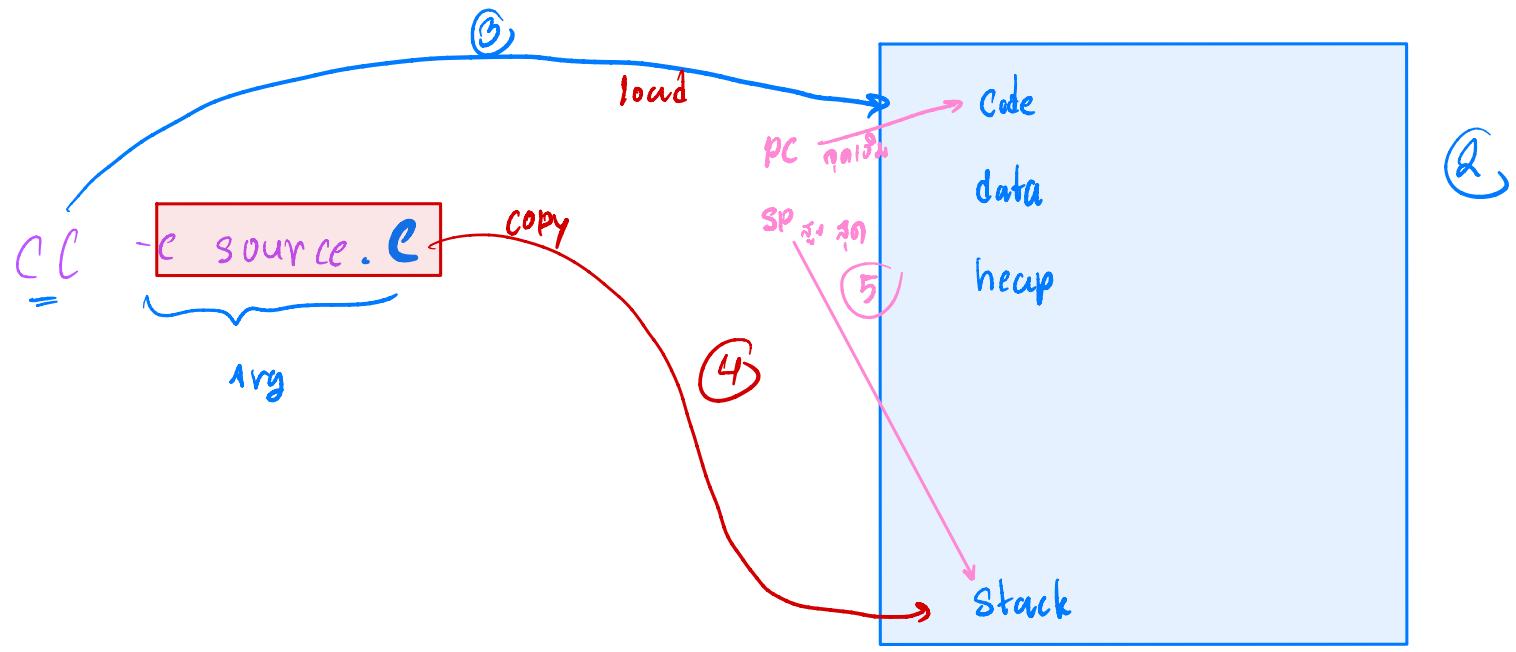
Activity #1

- If the shell runs at user-level, what system calls does it make to run each of the programs?
 - Ex: cc, ln

unix linux

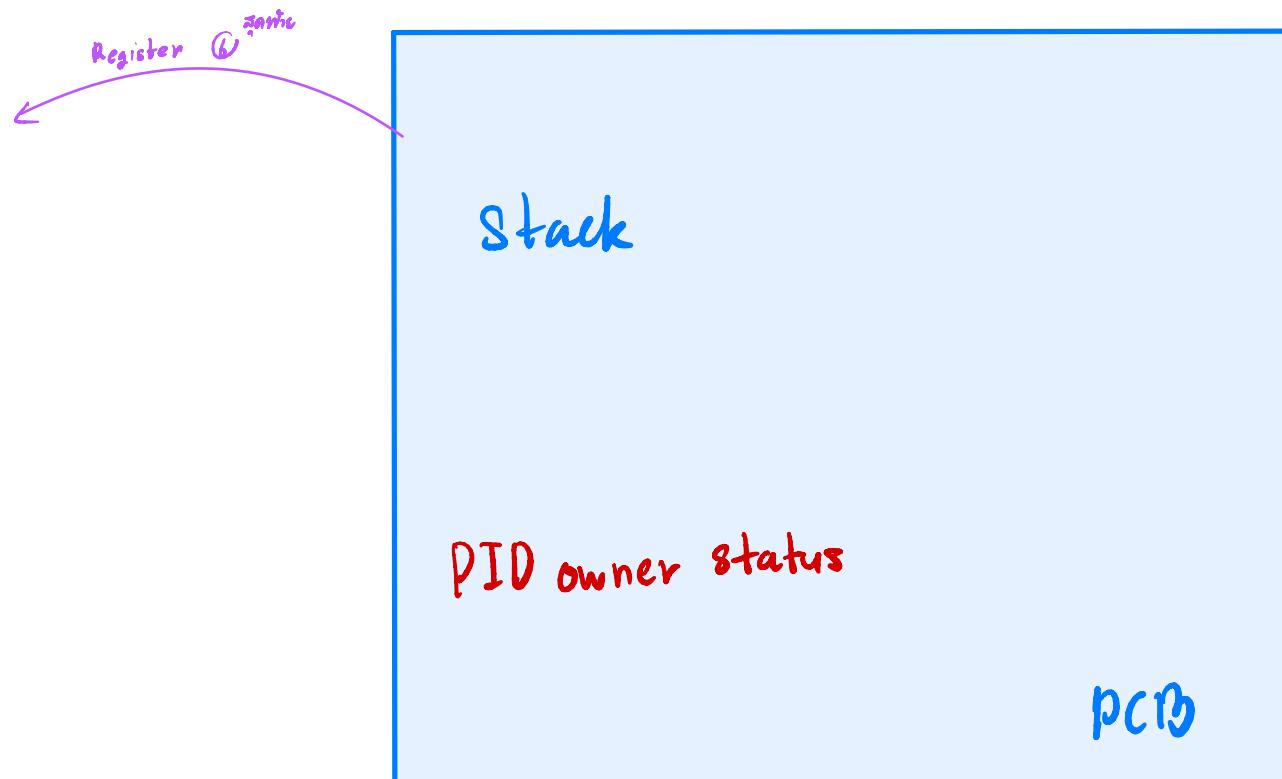
window

Shell minimum



kernel

Scheduler
minimum in process
if qd kernel yarayi emunmawin



Windows CreateProcess

- System call to create a new process to run a program
 - Create and initialize the process control block (PCB) in the kernel
 - Create and initialize a new address space
 - Load the program into the address space
 - Copy arguments into memory in the address space
 - Initialize the hardware context to start execution at “start”¹⁹
 - Inform the scheduler that the new process is ready to run

Windows CreateProcess API

(simplified) ver. 1.0

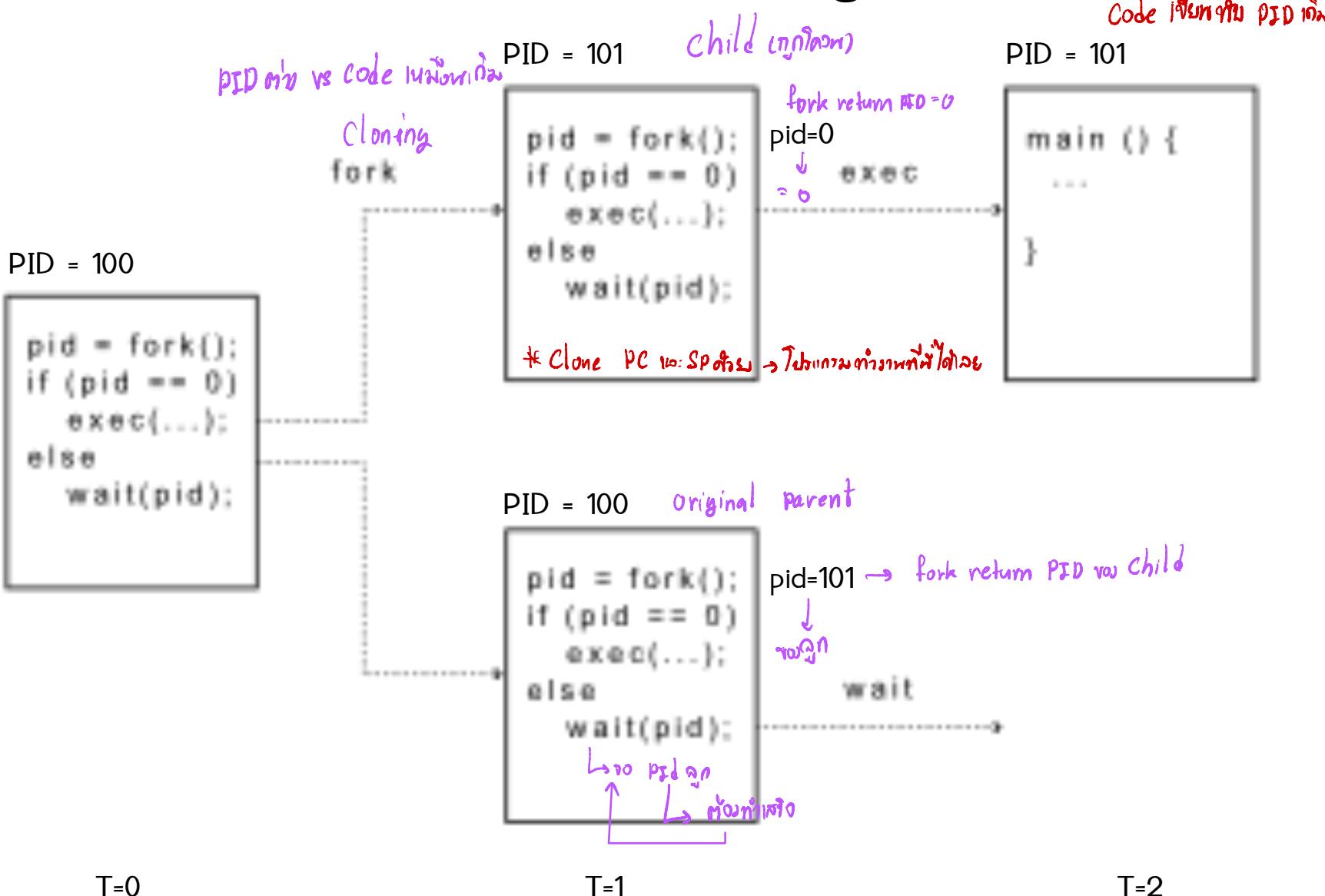
```
if (!CreateProcess(  
    NULL,          // No module name (use command line)  
    argv[1],       // Command line  
    NULL,          // Process handle not inheritable  
    NULL,          // Thread handle not inheritable  
    FALSE,         // Set handle inheritance to FALSE  
    0,             // No creation flags  
    NULL,          // Use parent's environment block  
    NULL,          // Use parent's starting directory  
    &si,           // Pointer to STARTUPINFO structure  
    &pi )          // Pointer to PROCESS_INFORMATION structure  
)
```

UNIX Process Management

Clone process ทำใหม่

- UNIX fork - system call to create a copy of the current process, and start it running
 - No arguments!
- UNIX exec - system call to change the program being run by the current process
- UNIX wait - system call to wait for a process to finish ดำเนิน shell
- UNIX signal - system call to send a notification to another process

UNIX Process Management



Activity #2: What does this code print?

ស្ថាប់នេះ ទាំងពីរ

[នៅ 101]
[== 0 false
else នៅណា else]

```
int child_pid = fork(); [ នៅ 100

if (child_pid == 0) {           // I'm the child process

    printf("I am process #%d\n", getpid());

    return 0;

} else {                      // I'm the parent process

    printf("I am parent of process #%d\n", child_pid);

    return 0;

}
```

Activity #3

- Can UNIX fork() return an error? Why?

ເລັກຕິບ / ດົນເກີມຂອງ

ໃຫຍດ ແລ້ວ Error ອັບ

- Can UNIX exec() return an error? Why?

ສົ່ງ -> ໄປນູຍຄອງນິ້ນໄປດີເລີຍ

- Can UNIX wait() ever return immediately? Why?

ຢູ່ນຳນົມ → ຖືກາຕົວ → return ຕັກຕົວ ບໍລິຫານ

ວິທະຍາ ວິທະຍາ

Implementing UNIX fork

Steps to implement UNIX fork

กี่ใน mem ที่ต้อง

- Create and initialize the process control block (PCB) in the kernel ||| กากตัน = ก็อปปี้ไปด้วย
- Create a new address space ||| ใหม่ต้องห้าม user space
- Initialize the address space with a **copy** of the entire **contents** of the address space of the parent No ผิดพลาด
- Inherit the execution context of the parent (e.g., any open files) No ผิดพลาด (นี่)
- Inform the scheduler that the new process is ready to run เริ่มต้น process + user space ให้

Implementing UNIX exec

วิธีการ
ตอนนี้จะสอน
วิธี execute ให้ดู
ในรูปแบบ pseudocode
ก่อน

- Steps to implement UNIX exec

Linux Pro ใจกลาง

- Load the program into the current address space
- Copy arguments into memory in the address space
- Initialize the hardware context to start execution at "start"

SP / process

No ข้อมูล

PC
as main
(start)

SP จุดเริ่มต้น stack

Linux code จะ รีเซ็ตตัวแปร

UNIX I/O

1975

- Uniformity ទូទាត់សំគេងកំណើនែងតាំង **system calls**
 - All operations on all files, devices use the same set of **system calls**: open, close, read, write
 - Open before use ការឱកាសីវេលា → ដោយក្នុងបច្ចេកទេស 1 program អាជីវកម្មរបស់លើកការណ៍ open នៅក្នុង
 - Open returns a handle (file descriptor) for use in later calls on the file
 - Byte-oriented **byte array** ឈរ
 - Kernel-buffered **read/write** អង់ត់ហ៊ី **Buffer**
 - Buffer នៃការសេដ្ឋកិច្ច និងការបញ្ចូល នៅក្នុងកុងហ៊ី
 - Explicit close ចាប់ OS ដោយកិច្ច kernel
 - To garbage collect the open file descriptor
- ការបញ្ចូល Buffer ត្រូវបានបញ្ចូល
កុងហ៊ី Buffer

UNIX File System Interface

- UNIX file open is a Swiss Army knife:
 - Open the file, return file descriptor
 - Options:
 - if file doesn't exist, return an error
 - If file doesn't exist, create file and open it
 - If file does exist, return an error
 - If file does exist, open file
 - If file exists but isn't empty, nix it then open
 - If file exists but isn't empty, return an error
 - ...

Activity #4

Interface Design Question

- Why not separate syscalls for open/create/exists?

ການມ່ວຍເປົນ 3 System call

ନୀମି - କର୍ମ

if (!exists(name))

check won't
in pro check won't

ក្នុងការអនុវត្ត → fail ក្នុង / ក្នុងសេចក្តី / fail នៅលាន Pro
និងនាមពេលរាយការណ៍រាយការណ៍

```
create(name); // can create fail?
```

```
fd = open(name); // does the file exist?
```

→ Sys 'call' តើដំឡើងនៅក្នុងការបញ្ជូនការណា

→ ជាកំណត់ កំណើនការ ដែលរាយ

↳ គណន៍ការងារខ្លួន

Implementing a Shell

ឧប្បម្ពការ + Arg

```
char *prog, **args;  
int child_pid;
```

```
// Read and parse the input a line at a time  
while (readAndParseCmdLine(&prog, &args)) {  
    child_pid = fork();      // create a child process  
    if (child_pid == 0) {  
        exec(prog, args);    // I'm the child process. Run program  
        // NOT REACHED  
    } else {  
        wait(child_pid);     // I'm the parent, wait for child  
        return 0;  
    }  
}
```

shell គាំទ្រ (រូប)

រាយការណាស់
shell ឱ្យការ

↓ ჩამოვარა გარეთ მისვლის

UNIX Process Management

PID ის vs კოდის მნიშვნელება

Cloning
fork

PID = 100

```
pid = fork();
if (pid == 0)
    exec(...);
else
    wait(pid);
```

PID = 101

child (უკიდო)

```
pid = fork();
if (pid == 0)
    exec(...);
else
    wait(pid);
```

fork return PID=0
pid=0
↓
exec
= 0

PID = 101

```
main () {
    ...
}
```

Code მუნიშვნელი PID ის

PID = 100 original parent

```
pid = fork();
if (pid == 0)
    exec(...);
else
    wait(pid);
```

↳ ის პიდ იქნა
↳ მიუნიშვნელი

pid=101 → fork return PID ვა child
↓
უკიდო

wait

T=0

T=1

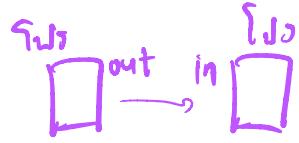
T=2

In Unix

- A program can be a file of commands
สุ เพาค์ฟิล์ม → แฟ้มคำสั่ง
dir > ชื่อพิมพ์.Txt เมื่อพิมพ์: ဂາເພັ່ນທີ່ນ ທີ່ນີ້ແມ່.Txt
- A program can send its output to a file
ອ່ານ ຕາກໄຟຟົນ
xxx < ທີ່ນີ້ແມ່.Txt
- A program can read its input from a file
ລົງຈາກນິກົມກັບລົງຈາກວຸນ
Input
- The output of one program can be the input to another program
ດົກ ໂປ່ງ
out → in



Interprocess Communication



↑
ກົດຕາມເກີບ
↓
ສຳຜົນໄຟຟ້າ

- Producer-consumer (2 program ຮ່ວມກຳທາງຫຼືອນກິນ ນົ່ວຍພົນລະ ພົບອົນຮັບ)

— Output of one program is accepted as input of another program

- One-way communication
- Pipe

- Client-server ສູ່ສະໜອບທາງ socket

- Two-way communication

client ↔ server

- Server implements specialize task **qັດ services** ຖຸນຂາຍ

— Print serve

- File system

ຮູບພົບ ເນື້ອຫຼາກຄາງໃນການ-ສູ່
ຈົບັດ ຢູ່

- Write data to a file then read file as an input

ອຸນດ້ານ

- Reader and writer are not need to running at the same time

ຕາມມີການນັກງານ ພົບອົນຮັບ-ສູ່ນ້າຍ

Operating system structure

- Monolithic kernel
- Microkernel

