#### Finding Partite Hypergraphs Efficiently

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- Hypergraphs
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#### *k*-Graphs

#### Definition

A *k-graph* is a pair G = (V, E) where V is a finite set of *vertices* and  $E \subseteq \binom{V}{k}$  is a set of *edges*.

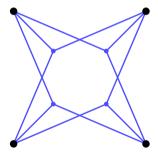


Figure: A complete 3-graph on 4 vertices:  $K_4^{(3)}$ .

#### Partite *k*-Graphs

#### Definition

A k-graph G = (V, E) is r-partite if there exists a partition  $V = V_1 \cup \cdots \cup V_r$  such that every edge of G intersects every part  $V_i$  in at most one vertex. We write  $G = (V_1, \ldots, V_r; E)$ .

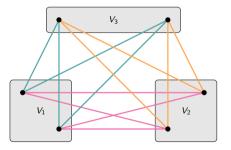


Figure: A complete 3-partite 2-graph:  $K^{(3)}(2,2,2)$ .

#### Partite *k*-Graphs

#### Remark

We may identify E as a subset of  $C = \bigcup_{\{i_1,\dots,i_k\} \in \binom{[r]}{k}\}} V_{i_1} \times \dots \times V_{i_k}$ . If E = C, we say that G is a *complete r*-partite k-graph.

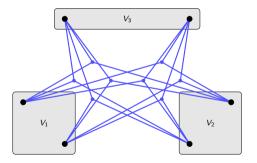


Figure: A complete 3-partite 3-graph:  $K^{(2)}(2,2,2)$ .

#### Turán-Type Problems

#### Definition

Let G = (V, E) be a k-graph and  $n \ge |V|$  an integer. The *Turán number*  $\operatorname{ex}(G, n)$  is the maximum number of edges in a k-graph on n vertices that does not contain a copy of G as a subgraph.

Determining  $\operatorname{ex}(G, n)$  or estimating it as  $n \to \infty$  is known as the *Turán problem* for G.

#### Theorem

For all k-graphs G there exists a constant  $\alpha(G) \in [0,1)$  such that

$$ex(G, n) = (\alpha(G) + o(1)) \cdot \binom{n}{k}$$
 as  $n \to \infty$ .

Furthermore,  $\alpha(G) = 0$  if and only if G is k-partite.

#### The Kővari–Sós–Turán Theorem

The bound  $ex(G, n) = o(n^k)$  can be improved by a lot.

#### Definition

Let  $1 < t_1 \le v_1, \ldots, 1 < t_k \le v_k$  be integers. Then the *generalized Zarankiewicz* number  $z(v_1, \ldots, v_k; t_1, \ldots, t_k)$  is the largest integer z for which there exists a k-partite k-graph  $H = (V_1, \ldots, V_k, F)$  with part sizes  $|V_i| = v_i$  and |F| = z edges such that for all choices of  $W_i \subset V_i$  of sizes  $|W_i| = t_i$ ,  $W_1 \times \cdots \times W_k \not\subset F$ .

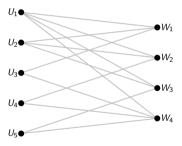
#### Theorem (Kővari–Sós–Turán)

Let  $0 < s \le u$  and  $0 < t \le w$  be integers. Then

$$z(u, w; s, t) \le (s-1)^{1/t}(w-t+1)u^{1-1/t} + (t-1)u$$

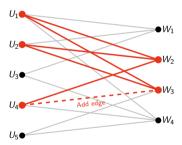
Standard arguments then show that  $ex(n, K(s, t)) = O(n^{2-1/t})$ .

This graph has the maximum number of edges (|E| = 13) to be  $K_{3,2}$ -free.



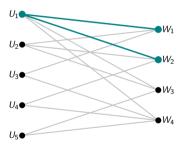
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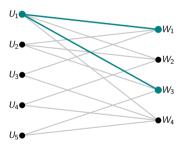
For example, adding the edge  $\{U_4, W_3\}$  creates a  $K_{3,2}$  on vertices  $\{U_1, U_2, U_4\}$  and  $\{W_2, W_3\}$ .

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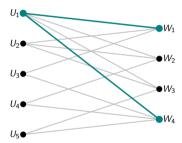
For 
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, we count its  $\binom{4}{2} = 6$  stars.

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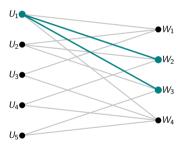
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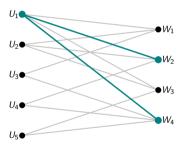
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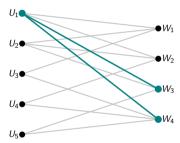
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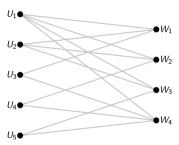
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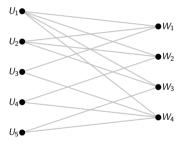
This graph has the maximum number of edges (|E|=13) to be  $K_{3,2}$ -free.



In the example, there are at least  $5\binom{13/5}{2} = 10.4$  stars (there are actually 12)

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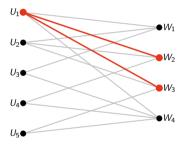
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Each set  $T \subset W$  (in this case,  $T = \{W_1, W_2\}$ ) is in at most s - 1 = 3 - 1 = 2 stars. In total, at most  $2\binom{4}{2} = 12$  stars.

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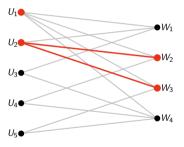
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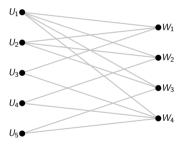
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In the example, we conclude that  $10.4 \le 12$ , which is true. For bigger values of z this would fail, leading to contradiction and therefore upper bounding z.

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- Conclusion:  $u\binom{z/u}{t} \le (s-1)\binom{w}{t}$ , from which the theorem follows.

## Erdős's Bound for Hypergraphs (1964)

#### Theorem (Erdős '64)

For integers 
$$k \geq 2$$
,  $t \geq 2$ ,  $ex(n, K(t, ..., t)) = O(n^{k - \frac{1}{t^{k-1}}})$ .

This generalizes the Kővari–Sós–Turán theorem to k-graphs.

It follows from a similar bound on the corresponding generalized Zarankiewicz number, obtained by induction.

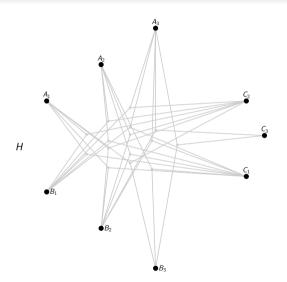
Suppose that  $H = (V_1, ..., V_k; F)$  is a k-graph with  $|W_i| = w$ . Let H have z edges and no copy of K(t, ..., t).

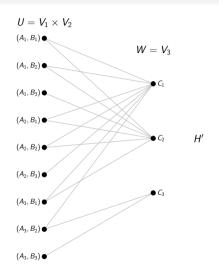
We set up a bipartite k-graph H' = (U, W; F') with

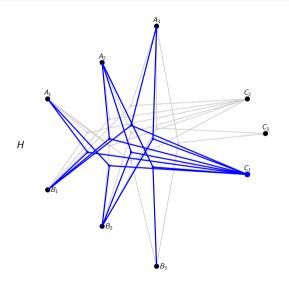
$$U = W_1 \times \cdots \times W_{k-1}$$

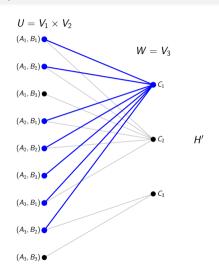
$$W = W_k$$

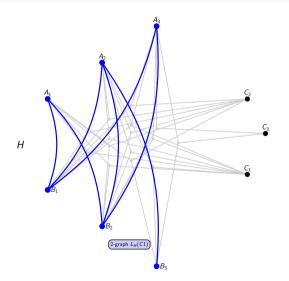
$$F' = \{(X, y) \in U \times W \mid X \cup \{y\} \in F\}.$$

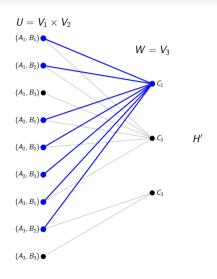


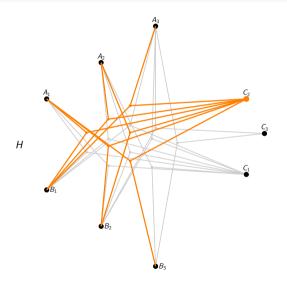


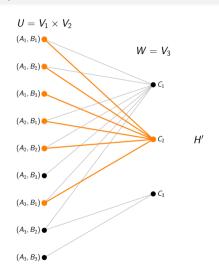


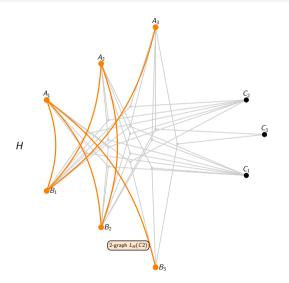


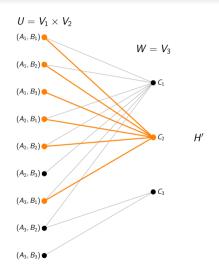


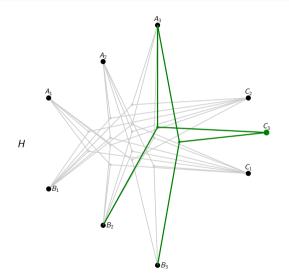


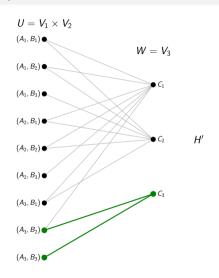


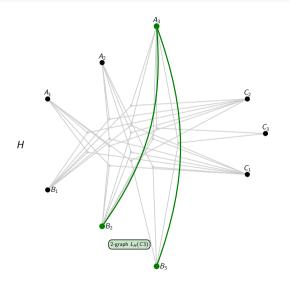


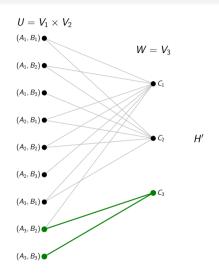


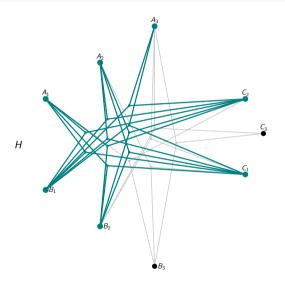


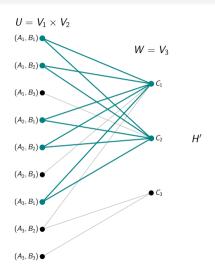


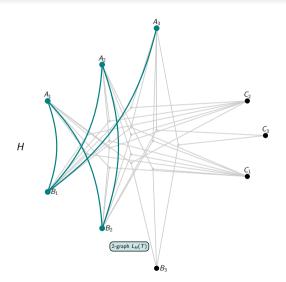


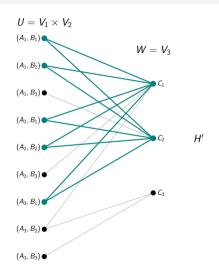


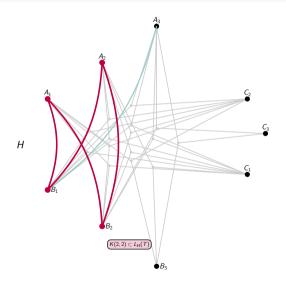


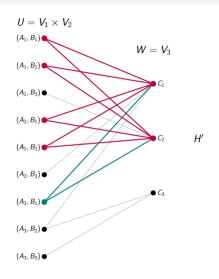


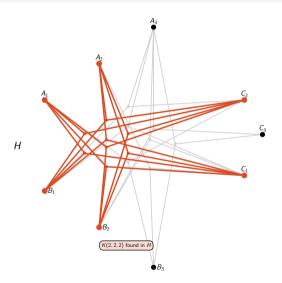


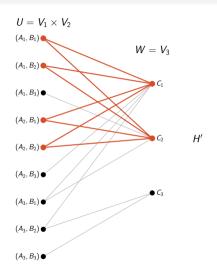












#### Implications of Erdős's Bound

Doing the calculations more carefully, we can show:

#### Theorem

Let  $k \ge 2$  and d > 0. Then there exists a constant  $\delta = \delta(k,d) > 0$  such that every k-graph G with n vertices and density  $d = m/n^k$  contains a complete balanced k-partite k-subgraph K  $(t, .^k, t)$  with

$$t \geq \delta \cdot (\log n)^{1/(k-1)}$$
.

That is, k-graphs of constant density contain complete balanced k-partite k-subgraphs of **growing** part size.

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Can we do it in polynomial time?

Hypergraphs Turán-Type Problems Algorithms Future Work

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