Extending Mubayi and Turán's Algorithm to 3-graphs

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Let G be a 3-graph with n vertices and $m = \epsilon n^3$ edges. A polynomial time algorithm is given to find a $K(q,q,q)$ in G for

$$q = \left\lfloor c_\epsilon^{(3)} \sqrt{\log n} \right\rfloor$$

As long as (insert condition here).

Note that this result is tight up to the constant $c_\epsilon^{(3)}$, as proved in [**?**]. This result is a generalization of the result in 2-graphs by [**?**], and algorithm will be analogous to the one given there. The procedure is as follows:

1. Choose parameters $q < r < n$ depending on n and ϵ.
2. Let R be the set of r vertices with the highest degree in G.
3. find a subset $Q \subset R$ with q vertices such that there is a large $S \subset \binom{[n] \setminus Q}{2}$ satisfying $xyz \in E(G) \, \forall \, \{x,y\} \in S$, $z \in Q$. Say, of size s.
4. Apply the algorithm of [**?**] to find a $K(q,q)$ in the 2-graph induced by S. Say, we find partition $S \supset U \cup V$.

If successful, a $K(q,q,q)$ has been found in G with parts U, V, Q. The problem is now to find parameters q, r such that the above procedure is successful and the algorithm runs in polynomial time.

Lemma 1. *As long as $r \leq \epsilon n$, there are at least $\epsilon r n^2$ edges in G with exactly one vertex in R.*

Proof. The sum of the degrees in G is $3m$. Therefore, by the pigeonhole principle,

$$\sum_{v \in R} d(v) \geq r \cdot \frac{3m}{n} = 3\epsilon r n^2$$

However, here we are overcounting:

- The edges with only one vertex in R are counted exactly once.
- The edges with two vertices in R are counted twice. The contribution of these is at most $r(r-1)(n-r) < r^2 n$

- The edges with all vertices in R are counted three times. The contribution of these is at most $r(r-1)(r-2) < r^3 < r^2 n$

Therefore, the condition will hold as long as $r^2 n \leq \epsilon r n^2 \iff r \leq \epsilon n$. $\square$

Next, a counting argument in the style of [**?**] is used to guarantee the existence of Q and S. The size q of Q will be left as a parameter to be determined later, and the size s of S will be determined by the following lemma:

Lemma 2. *Under the same assumptions as in Lemma 1, and assuming $r \leq n/2$, $r \geq q/\epsilon$, there is a subset $Q \subset R$ of size q and a subset $S \subset \binom{[n]\setminus Q}{2}$ of size*

$$s := \frac{n^2}{8}\left(\frac{\epsilon}{e}\right)^q$$

such that $xyz \in E(G) \,\forall \{x,y\} \in S$, $z \in Q$.

Proof. Let E be the set of edges with exactly one vertex in R. for every $\{x,y\} \in \binom{[n]\setminus R}{2}$, let E_{xy} be the set of edges in E containing x and y. Finally, let

$$T = \left\{ P \subset E_{xy} : x, y \in \binom{[n]\setminus R}{2},\ |P| = q \right\}$$

On the one hand, the number of elements in T is

$$\sum_{\{x,y\}\in\binom{[n]\setminus R}{2}} \binom{|E_{xy}|}{q} \geq \binom{n-r}{2}\binom{\epsilon r n^2/\binom{n-r}{2}}{q} > \binom{n-r}{2}\binom{\epsilon r}{q} \geq \frac{n^2}{8}\binom{\epsilon r}{q}$$

Where the first inequality follows from the convexity of

$$f(x) = \begin{cases} \binom{x}{q} & \text{if } x \geq q-1 \\ 0 & \text{otherwise} \end{cases}$$

and the third from the fact that $r \leq n/2$.

On the other hand, there are only $\binom{r}{q}$ possible $q-$subsets of R. By the pigeonhole principle, one of these (say Q) must be the set of vertices in R associated with P_j for k different $P_j \in T$, where k is

$$\frac{|T|}{\binom{r}{q}} > \frac{n^2\binom{\epsilon r}{q}}{8\binom{r}{q}} \geq \frac{n^2\left(\frac{\epsilon r}{q}\right)^q}{8\left(\frac{er}{q}\right)^q} = \frac{n^2}{8}\left(\frac{\epsilon}{e}\right)^q = s$$

$\square$

Now, the algorithm of [?] is applied to the graph G' with vertex set $[n]$ and edge set S.

This yields a $K(q', q') \subset G'$ with

$$q' = \left\lfloor \frac{\ln(n/2)}{\ln(2en^2/s)} \right\rfloor = \left\lfloor \frac{\ln(n/2)}{\ln(16e^{q+1}/\epsilon^q)} \right\rfloor$$

For the found subgraph to be a $K(q, q, q)$, it is necessary that $q' \geq q$. A sufficient condition is that

$$q \leq \frac{\ln(n/2)}{\ln(16e^{q+1}/\epsilon^q)} - 1 = \frac{\ln(n/2)}{\ln(16e) - q\ln(e/\epsilon)} - 1$$

This is true for

$$0 \leq q \leq \frac{\ln(16\epsilon) + \sqrt{(\ln(16\epsilon))^2 + 4\ln(n/(32e))\ln(e/\epsilon)}}{2\ln(e/\epsilon)}$$

so a valid value for q is

$$q = \left\lfloor \frac{\sqrt{4\ln(n/(32e))\ln(e/\epsilon)}}{2\ln(e/\epsilon)} \right\rfloor = \left\lfloor \frac{\sqrt{\ln(n/(32e))}}{\sqrt{\ln(e/\epsilon)}} \right\rfloor \sim \frac{\sqrt{\ln n}}{\sqrt{\ln(e/\epsilon)}},\ n \to \infty$$

This means that, for small n, we can just find the biggest 3-partite subgraph in G by hand, and for large n, we can use the algorithm described with

$$q = \left\lfloor c_\epsilon^{(3)} \sqrt{\log n} \right\rfloor$$

and

$$c_\epsilon^{(3)} = \frac{1}{\sqrt{\ln(2e/\epsilon)}} = \frac{1}{\sqrt{\ln(2en^3/m)}}$$

Finally, to determine the value of r and the running time of the algorithm, recall that the only conditions we have imposed on r are

$$\begin{aligned} r &\leq \epsilon n \\ r &\leq n/2 \\ r &\geq q/\epsilon \end{aligned}$$

So we can define $r = \lceil q/\epsilon \rceil$ as long as $\lceil q/\epsilon \rceil \leq \min\{n/2,\ \epsilon n\}$. For $\epsilon \geq 1/2$ this clearly holds for n large enough,

References

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[3] Dhruv Mubayi and György Turán. Finding bipartite subgraphs efficiently. *Information Processing Letters*, 110(5):174–177, 2010.