



zero-knowledge proof of disjunctive statements (OR proofs)

Asked 4 years, 9 months ago Modified 11 months ago Viewed 1k times



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I know there are standard ways to prove disjunctive statements about discrete logs, e.g. [OR proof](#). But are there similar approaches for other class of language? For example, how can one go about proving either G_1 or G_2 has a Hamiltonian path (without leaking which one)?

zero-knowledge-proofs

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asked May 26, 2017 at 11:53



qweruiop

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1 Answer

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You can use the [CDS94](#)-technique for that. Suppose you have two zero-knowledge proof systems σ_1 and σ_2 , both of which consist of three messages: a commitment com , a (public-coin) challenge ch , and a response r . And suppose you want to prove the OR of both claims, meaning that you can generate a valid proof for (wlog.) σ_1 for any challenge even after generating the commitment, but in order to generate a valid proof for σ_2 you must get the challenge first and compute the commitment from there. So you have essentially one degree of freedom in the choice of (ch_1, ch_2) , i.e., one is fixed beforehand but you want to hide which one. The clue is to get a new challenge ch' from the verifier and guarantee that (ch_1, ch_2, ch') satisfies a suitable relation, for instance sum-to-zero. So the proof system for σ_1 -OR- σ_2 looks like this:

- preprocessing by the prover: compute com_2, ch_2, rsp_2 for a random ch_2
- prover sends commitment: com_1, com_2
- verifier sends challenge: ch'
- prover computes $ch_1 \leftarrow ch' - ch_2$ and uses this to complete σ_1
- prover sends response: $(ch_1, rsp_1), (ch_2, rsp_2)$
- verifier verifies that (com_1, ch_1, rsp_1) is valid, that (com_2, ch_2, rsp_2) is valid, and that $ch_1 + ch_2 + ch' = 0$

You can use $\sum_i ch_i = ch'$ for OR-proofs consisting of any number of claims. However, in some cases you want to prove more specific facts such as " t out of these n claims are true". In this case, you want to use Shamir's secret sharing and exchange the sum-to-zero relation for a polynomial of degree $n + 1 - t$: this guarantees that all $n - t$ degrees of

freedom for choosing ch_i must be used up by the false claims.

The CDS-94 technique applies to any zero-knowledge proof system that follows this three-pass public coin structure. Schnorr's protocol for proving discrete logarithm knowledge is just one particular case of that structure.

[CDS94](#): Cramer, Ronald, Ivan Damgård, and Berry Schoenmakers. "Proofs of partial knowledge and simplified design of witness hiding protocols." CRYPTO 1994.

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edited Apr 11, 2021 at 9:39



Community Bot

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answered May 26, 2017 at 12:55



Alan

1,350

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Can you give an example for the "t out of n claims" or a reference? – [Jus12](#) Jan 8, 2019 at 18:23

The CDS'94 paper referenced above treats this case extensively. – [Alan](#) Jan 9, 2019 at 19:15
