HW5: Padding Oracle Attacks and Authenticated-encryption Scheme

Due: April 30, 2018

This homework is divided into two projects:

Project 1: Padding Oracle Attacks

The CBC mode of operation for a block cipher can encrypt messages with length a multiple of the underlying block cipher's block length. What will happen if the length of a message is not multiple of the block length? To extend the functionality of block ciphers to arbitrary length of messages, we use padding. In this exercise we will see how padding gives rise to so-called padding oracle attacks that allow attackers to (very often) recover messages without access to the corresponding secret key.

There are several standard ways to pad a message. For the purpose of this assignment we shall use the padding standard as described in PKCS#7 (https://tools.ietf.org/html/rfc2315#section-10.3). PKCS#7 padding works as follows. If the message requires k bytes to be appended in order for the resulting string to to be a multiple of the block length, then pad the message with k bytes each of value k. If k=0 (that is the message is already a multiple of the block length), then pad it with a new block, each byte of which is set to the block length. Under these rules, note that one never pads by more than the block length in bytes, which for AES is 16 (hex value 0x10). As one example assuming block length of 16 bytes:

Given a message, first, it is padded using the method mentioned above, and then encrypted using CBC mode with some block cipher. Most often these days one uses AES, which has a block length of 16 bytes. At the time of decryption, the ciphertext is first decrypted and then padding is removed to output the plaintext. To remove padding, CBC decryption is performed, and then the last byte of the resulting string is inspected. It is interpreted as the number of bytes of padding, and the decryption routine scans from right to left in the string to remove these values. If at any point the padding value is not the expected one (e.g., the final padding byte was 02, but the second to last most byte is not 02) then a padding error is

raised and decryption is aborted. Often padding errors are observable, either directly because a padding error message is reported or indirectly because of timing or other issues.

Serge Vaudenay in 2002

(https://www.iacr.org/archive/eurocrypt2002/23320530/cbc02_e02d.pdf)

showed how, given access to a public decryption oracle¹ and a ciphertext, it is possible to recover the whole plaintext of some challenge encryption. This is without ever having access to the secret key

What is given?

You are given an implementation of a padding oracle simulator via the PaddingOracle class, which implements the AES-CBC encryption scheme with PKCS#7 padding. This class provides 3 APIs:

What do you have to do?

You have to write a function that takes an oracle instance and an arbitrary length ciphertext, and decrypts the ciphertext by implementing a padding oracle attack. Your attack should work by only interacting with the PaddingOracle object via the decrypt function and should not, for example, read its internal variables. No credit will be given for attacks based on the latter.

Part 1 (30 points): First assume the ciphertext is of length 2 blocks: The first block is $C_0 = iv$, and the second block $C_1 = E_k (msg \oplus C_0)$. Given such a ciphertext write a function that recovers the plaintext by querying the padding oracle instance. Signature of your function should be:

```
po_attack_2blocks(po, ctx)
```

Part 2 (20 points): In the second part you have to extend your attack to arbitrary length messages/ciphertexts.

¹ Oracle in cryptographic context is a black box that implements some API, and anybody can call those APIs to get a result. In our context the padding oracle is a decryption routine.

<u>Hint:</u> You can use the function in the part 1 as a subroutine for this. po_attack(po, ctx)

Turn into CMS your Python source files (.py) with your modifications. Please do not modify paddingoracle.py.

Project 2: Authenticated-encryption Scheme

In this assignment, you will build an implementation of a modified Fernet specification. The new algorithm only relies on AES, as opposed to combining AES-CTR with HMAC. Using just AES can be advantageous in places where one wants small code complexity, or to take advantage of AES-NI hardware support. The specification document can be found at this url.

Part 1 (30 points): Fully implement the Fernet version 0x91 spec, starting from AES. This means you must implement AES CTR and AES CMAC, for example, using only the Hazmat backend to give you access to AES. (More specifically, you can only use AES in ECB mode in the library.) Your code should be clean, well-commented, and easy to read. It should not contain debugging print statements, dead code, or other cruft. We will split this 30% up as follows:

(10 points): Quality and cleanliness of code and comments.

(5 points): Self-correctness (whether it can decrypt something it encrypted).

(15 points): Correctly implements Fernet 0x91 spec. For example, can you decrypt a ciphertext encrypted with our implementation?

Part 1 (20 points): Write a unit testing suite for your implementation. Your tests should cover basic correctness as well as proper handling of error conditions and edge cases. Here you can use the Hazmat backends for CTR mode and CMAC to help you generate test vectors. (That is, you should be able to build an equivalent implementation of Fernet 0x91 using the Hazmat implementations of CTR and CMAC.) Ideally, your tests will cover 100% of the lines of your implementation. Code coverage metrics can be checked using a tool like pytest-cov.

Turn into CMS your Python source files (.py) with your modifications.