EE290 Course Notes

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1.1 Results from random matrix theory

Today we consider random matrices $Z = (Z_{ij}) \in \mathbb{R}^{n\times n}$. IID matrix ensemble is when $Z_{ij} \sim P$ are drawn IID, and the Gaussian Orthogonal Ensemble (GOE) has $Z_{ii} \sim N(0,2)$ and $Z_{ij} = Z_{ji} \sim N(0,1)$ for $i \neq j$.

By convention, normalize and center so $\mathbb{E}Z_{ij} = 0$ and $\mathbb{E}Z_{ij}^2 = 1$.

Intuition: $\|Z\|_{op} \leq C\sqrt{n}$ with high probability.

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Consider Gaussian orthogonal ensemble matrix: $Z_{ij} \sim N(0,1)$ and $Z_{ii} \sim N(0,2)$. View $Z = \left[Z_1, \ldots, Z_n\right]$ with $Z_i \sim N(0, I_n)$. Then

$$\mathbb{E}\|Z_1\|_2^2 = \mathbb{E}[\sum_{i=1}^n Z_{i1}^2] = n \tag{1}$$

$$Z_1^\top Z_2 = \sum_{i=1}^n Z_{i1} Z_{i2} \tag{2}$$

$$\mathbb{E} Z_1^\top Z_2 = 0 \tag{3}$$

$$\mathbb{E}(Z_1^\top Z_2)^2 = n \tag{4}$$

$$|Z_1^\top Z_2| \sim \sqrt{n} \tag{5}$$

$$\frac{Z_1^\top Z_2}{\|Z_1\|\|Z_2\|} \sim \frac{1}{\sqrt{n}} \tag{6}$$

Theorem 1 (*Latała et al. (2006)*)

$$\sup_i \sum_{j=1}^n \mathbb{E}|Z_{ij}|^2 \leq k^2 n \tag{7}$$

$$\sup_j \sum_{i=1}^n \mathbb{E}|Z_{ij}|^2 \leq k^2 n \tag{8}$$

Fourth moment bound

$$\sum_{i=1}^n \sum_{j=1}^n \mathbb{E}|Z_{ij}|^4 \leq k^4 n^2 \tag{9}$$

Then $\mathbb{E}\|Z\|_{op} = O(k\sqrt{n})$

1.2 Gaussian Orthogonal Ensemble

$\|Z\|_{op} = \sigma_{max} = \max_{\|v\|=1} v^\top Z v$

For any fixed $v \in S^{n-1}$, we have a Gaussian tail bound

$$v^\top Z v = \sum_i Z_{ii} v_i + \sum_{i<j} 2 Z_{ij} v_i v_j \tag{10}$$

$$= N(0, \sum_i v_i^4 + \sum_{i<j} 4 v_i^2 v_j^2) \tag{11}$$

$$\Pr(|v^\top Z v| > t) \leq 2e^{-t^2/4} \tag{12}$$

Using an ϵ-net, can find a set of vectors V_ϵ such that

$$\max_{v \in V_\epsilon} |v^\top Z v| \geq (1-2\epsilon) \max_{|v|=1} |z^\top Z v| \geq (1-2\epsilon) t \tag{13}$$

Then by a union bound

$$\Pr[\|Z\|_{op} \geq t] \leq \Pr[\max_{v \in V_\epsilon} |v^\top Z v| \geq (1-2\epsilon)t] \tag{14}$$

$$\leq \sum_{v \in V_\epsilon} \Pr[|v^\top Z v| \geq (1-2\epsilon)t] \tag{15}$$

$$\leq 2|V| e^{-\frac{1}{4}(1-2\epsilon)^2 t^2} \leq \delta \tag{16}$$

If $|V| \leq c^n$, then

$$e^{c(n-ct^2)} \leq e^{\log \delta} \tag{17}$$

$$\log \frac{1}{\delta} \leq ct^2 - n \implies t \geq \sqrt{n + \log \frac{1}{\delta}} \tag{18}$$

Intuition: when dealing with infinite dimensional maximization (Rayleigh quotient for eigenvalue problem), can pass to ϵ-net for cardinality bboud.

Definition 2 (*Covering*)

$V \subset S^{n-1}$ is called an ϵ-net if $\forall u \in S^{n-1}$, $\exists v \in V$ such that $\|u - v\|_2 \leq \epsilon$.

Theorem 3

ϵ-net yields Eq. (13)

Definition 4 (*Packing*)

For $A \subset \mathbb{R}^d$, $V = \{v_i\}_{i=1}^n \subset A$ is an ϵ-packing if $\forall i \neq jJ$, $\|v_i - v_j\|_2 \geq \epsilon$.

Theorem 5

Maximal ϵ-packing is an ϵ-net.

Hence, we can lower bound the packing number (size of largest packing) by the covering number (size of the smallest covering). The following result gives an (obvious?) upper bound:

Lemma 6 (*Volume ratio*)

For any ϵ-packing $V \subset A$,

$$|V| \leq \frac{Vol(A + \frac{\epsilon}{2}B)}{Vol(\frac{\epsilon}{2}B)} \tag{19}$$

where $B = \{x : \|x\|_2 \leq 1\}$.

Why is the diagonal not important? Let $A = \operatorname{diag}(Z)$. Then we have

$$\|Z - A\|_{op} \leq \|Z\|_{op} + \|A\|_{op} \tag{20}$$

$$\max_{x \in S^{n-1}} \|Ax\| = \max_i |Z_{ii}| = O(\sqrt{2 \log n}) \tag{21}$$

So the diagonal term $\|A\|_{op}$ is an order of magnitude smaller that $\|Z\|_{op}$.

Example 7 (*Planted clique*)

Let $G \sim G(1/2, n, k)$. In other words, generate an Erdös-Renyi random graph from $G(n, 1/2)$ and then randomly choose a set $K \subset [n]$ connect together to form a clique.

Goal: find K given G.

Theorem 8 (*Alon et al. (1998)*)

For any c, $k = c\sqrt{n}$, then exists polytime algorithm such that it returns $\hat{K}$ with $P(\hat{K} = K) \to 1$.

Let the adjacency matrix $A_{ij} = \begin{cases} 1 & (i,j) \in K \\ \text{Bern}(1/2) & i \notin K \text{ or } j \notin K, i \neq j \\ 0 & i = j \end{cases}$ and define $W_{ij} = \begin{cases} 2A_{ij} - 1 & i \neq j \\ 0 & i = j \end{cases}$

1. Find top eigenvector u of W

2. Let $\tilde{K}$ index the k largest coordinates $|u_i|$

3. Thresholding

$$\hat{K} = \left\{ v \in [n] : d_{\tilde{K}}(v) \geq \frac{3k}{4} \right\} \tag{22}$$
$$d_{\tilde{K}}(v) = \sum_{j \in \tilde{K}} \mathbb{1}\{(j,v) \text{ connected}\} \tag{23}$$

Goal: show $|\tilde{K} \cap K| \geq (1-\epsilon)k$ whp.
Note that $\mathbb{E}[W] =: 1_k 1_k^\top - \text{diag}(1_k)$ consists of 1s in $K \times K$ and 0 everywhere else. Let

$$W^* = 1_k 1_k^\top \tag{24}$$
$$v = \frac{1}{\sqrt{k}} 1_k \tag{25}$$
$$\tag{26}$$

Notice thresholding over v exactly recovers K, so we want the top eigenvector u of W to be close to v. By Davis-Kahan,

$$\min_{s \in \{\pm 1\}} \|u + sv\|_2 \leq \frac{\|W - W^*\|_{op}}{\lambda_1(W^*) - \lambda_2(W^*)} \tag{27}$$

Note $\lambda_1(W^*) = k$. Suppose extrema attained at $s = -1$, then

$$\|W - W^*\|_{op} \leq \|W - \mathbb{E}W\| + \underbrace{\|\mathbb{E}W - W^*\|}_{=\|\text{diag}\, 1_k\| = 1} \leq c\sqrt{n} + 1 \tag{28}$$

By Weyl's inequality

$$|\lambda_2(W)| = |\lambda_2(W^*) - \lambda_2(W)| \leq \|W^* - W\|_{op} \leq c\sqrt{n} + 1 \tag{29}$$

Finally

$$\|u - v\|_2 \leq \frac{c\sqrt{n} + 1}{c\sqrt{n} - (c\sqrt{n} + 1)} \leq \epsilon \tag{30}$$

NOTE: when you have bounded fourth moments, the rate is always $n^{-1/2}$! Deep result.

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Recall the planted clique from Alon et al. (1998): $G \sim G(1/2, n, k)$ is a random graph on $V = [n]$ with some fully connected clique $K \subset [n]$ of cardinality $|K| = k$.

The adjacency matrix

$$A_{ij} = \begin{cases} 1 & \text{if } i, j \in K \\ \text{Bern}(1/2) & i \neq j \text{ ow} \end{cases} \tag{31}$$

Let

$$W_{ij} = \begin{cases} 2A_{ij} & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases} \tag{32}$$

Algorithm 1 of Alon et al. (1998):

1. Find top eigenvector of W, say u

2. Let $\tilde{K}$ index the largest k coordinates $|u_i|$
3. Define $\hat{K} = \{v \in V : d_{\tilde{K}}(v) \geq \frac{3k}{4}\}$

Theorem 9 (*Alon et al. (1998)*)

Algorithm 1 finds $\hat{K}$ such that $\Pr[\hat{K} = K] \to 1$ as $n \to \infty$ if $k \geq c\sqrt{n}$ for sufficiently large c.

Proof. Note that $\mathbb{E}A$ is:

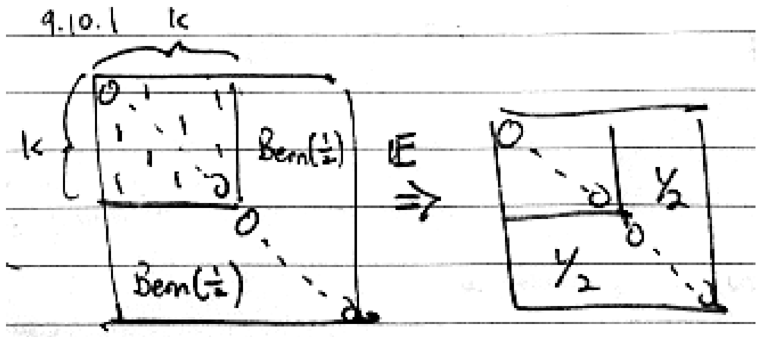

Figure 1: $\mathbb{E}A$ has ones in the upper $k \times k$ block, 0 on the diagonal, and 1/2 everywhere else

From this, we can easily see that the $\mathbb{E}W$ is:

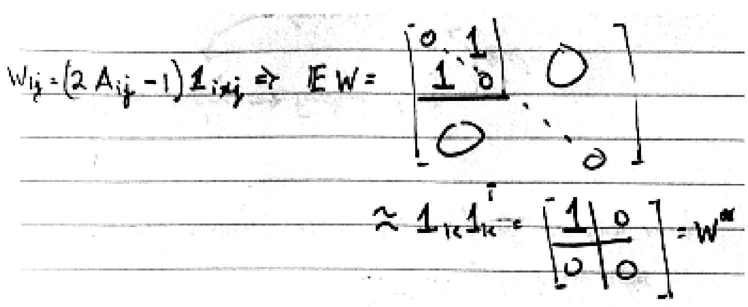

Figure 2: $\mathbb{E}W$ differs from $W^* = 1_k 1_k^\top$ only in the upper k diagonal

Note $\mathbb{E}W = 1_K 1_K^\top - \mathrm{diag}(1_K) \approx 1_K 1_K^\top = W^*$, which is good because we have seen that "differenes in the diagonal are asymptotically negligible."

reference for this? 9-5 lecture

Goal: show $|\tilde{K} \cap K| \geq (1-\varepsilon)k$ whp, $\varepsilon = \varepsilon(c)$.

We first show the top eigenvector of W^* is close to u (the top eigenvector of W). Let $v = \frac{1}{\sqrt{k}} 1_K$ be the top eigenvector of W^*. Note $\lambda_1(W^*) = k$. By Davis-Kahan

$$\min_{s \in \{\pm 1\}} \|u + sv\|_2 \leq \frac{\|W - W^*\|_2}{\lambda_1(w^*) - \lambda_2(w)} \tag{33}$$

Note

$$\|W - W^*\| \leq \|W - \mathbb{E}W\| + \|\mathbb{E}W - W^*\| \leq c\sqrt{n} + 1 \tag{34}$$

Also $\lambda_1(W^*) = k$ and

$$|\lambda_2(W)| \leq |\lambda_2(W^*) - \lambda_2(W)| \leq \|W^* - W\| \tag{35}$$

So by Weyl's inequality

$$\min_{s\in\{\pm 1\}} \|u+sv\|_2 \le \frac{c\sqrt{n}+1}{k-(c\sqrt{n}+1)} \tag{36}$$

$$\le \frac{c\sqrt{n}+1}{c\sqrt{n}-c\sqrt{n}+1} \le \varepsilon \tag{37}$$

Aside: Davis-Kahan to get bound between difference of eigenvectors in 2-norm. Open problem to control others.

Next, if $|K| = k = |\tilde{K}|$ then $|K \setminus \tilde{K}| = |\tilde{K} \setminus K|$.

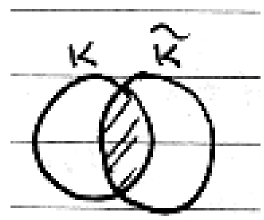

Figure 3: $|K| = |\tilde{K}| \implies |K \setminus \tilde{K}| = |\tilde{K} \setminus K|$ follows from elementary set theory

By definition of v

$$\varepsilon^2 \ge \|u-v\|_2^2 = \sum_{i\in K}(u_i - \frac{1}{\sqrt{k}})^2 + \sum_{i\notin K} u_i^2 \tag{38}$$

Lemma 10

If all $|u_i| \le \frac{1}{2\sqrt{k}}$ for $i \notin \tilde{K}$, then

$$\varepsilon^2 \ge \sum_{i\in K\setminus\tilde{K}} (\frac{1}{\sqrt{k}} - u_i)^2 \ge \sum_{i\in K\setminus\tilde{K}} \frac{1}{4k} \tag{39}$$

This implies $|K \setminus \tilde{K}| \le 4\varepsilon^2 k$.

Lemma 11

If the condition of the previous lemma does not hold, then $\exists i \in \tilde{K}$ with $|u_i| \ge \frac{1}{2\sqrt{k}}$. Then in fact $|u_i| \ge \frac{1}{2\sqrt{k}}$ for all $i \in \tilde{K}$ since

$$\varepsilon^2 \ge \sum_{i\in\tilde{K}\setminus K} u_i^2 \ge \sum_{i\in\tilde{K}\setminus K} (\frac{1}{2\sqrt{k}})^2 = \sum_{i\in\tilde{K}\setminus K} \frac{1}{4k} \tag{40}$$

Hence $|\tilde{K} \setminus K| \le 4\varepsilon^2 k$

So we have achieved our goal.

To finish the proof, first assume $\|u-v\|_2 \le \varepsilon$. For $a \in K$,

$$d_{\tilde{K}}(a) \ge d_{\tilde{K}\cap K}(a) = |\tilde{K}\cap K| - 1 \ge (1-\varepsilon')k \tag{41}$$

so for $a \in K$, we will get $a \in \hat{K}$.

Now if $a \notin K$,

$$d_{\tilde{K}}(a) \le \underbrace{d_K(a)}_{\sim\text{Binom}(k,1/2)} + \underbrace{|\tilde{K}\setminus K|}_{\le\varepsilon' k} \approx \frac{k}{2} \pm c\sqrt{k} \tag{42}$$

where $\approx$ means concentration. To be concrete,

$$\Pr[\hat{K} \neq K] \leq \Pr[\|u-v\|_2 \geq t] + \Pr[\exists a \notin K : d_K(a) \geq (\frac{3}{4} - \varepsilon')k] \tag{43}$$

$$\leq \Pr[\|W - \mathbb{E}W\| \geq c\sqrt{n}] + (n-k)\Pr[B(k,1/2) \geq (\frac{3}{4} - \varepsilon)k] \tag{44}$$

$$\leq ce^{-c'n} + (n-k) \tag{45}$$

Where above we used the multiplicative version of Chernoff bound (useful in combinatorial statistics):

Lemma 12 (***Multiplicative Chernoff Bound***) {lem:mult-che

$$\Pr[X \geq (1+\delta)\mu] \leq \begin{cases} e^{-\delta^2\mu/3} & \delta \in [0,1] \\ e^{-\delta\mu/3} & \delta \geq 1 \end{cases} \tag{46}$$

$$\Pr[X \leq (1-\delta)\mu] \leq e^{-\delta^2\mu/2} \tag{47}$$

As $n \to \infty$, we see that $\Pr[\hat{K} = K] \to 1$. □

Lemma 12 is self-normalizing: let $X = \sum_{i=1}^n X_i$ with X_i independent binary and $\mu = \mathbb{E}X$. Note that after applying, the RHS does not depend on n Verify

AKS Algorithm 2: This algorithm is designed to handle the case when k is not big enough (recall algorithm 1 requires $k \geq c\sqrt{n}$). Search over all S with $|S| = C(c) = 2\log_2 \frac{10}{c} + 2$. For each S:

1. Define $N^*(S) = \{v \in V : v \sim a, \forall a \in S\} \setminus S$
2. Run Algorithm 1 on the induced subgraph (which has distribution $G(1/2, N^*(S), K - S)$), return $Q_S \cup S$
3. Output if $Q_S \cup S$ is a k-clique

Intuition: Suppose $k = 0$ so there's no clique. Then $|N^*(S)| \sim B(n-s, 2^{-s}) \approx \frac{n-s}{2^s}$ so the total number of nodes is much smaller (by order of 2^{-s}). However, the number of clique nodes in $N^*(S)$ is still relatively large, $\geq k - s$. Solving the critical equation (also for algorithm 1) Track htis down

$$k - s \geq C\sqrt{\frac{n}{2^s}} \tag{48}$$

yields the expression for $C(c)$.

Theorem 13

As long as $k \geq (2+\varepsilon)\log_2 n$, then exhaustive search finds k with probability $\to 1$.

Proof. Exhaustive search will always find the clique, but it may return a clique that we didn't plant. So we need to guarantee there is no clique of size $(2+\varepsilon)\log_2 n$ in G whp.

For $S \subset [n]$, $|S| = k$,

$$\Pr[S \text{ is clique}] = \frac{1}{2^{\binom{k}{2}}} \tag{49}$$

$$\Pr[\exists S \subset [n] : S \text{ is clique}] \leq \binom{n}{k}\frac{1}{2^{\binom{k}{2}}} \leq (n2^{-(k-1)/2})^k \to 0 \tag{50}$$

$$\tag{51}$$

as $n \to \infty$ $(k = (2+\varepsilon)\log_2 n)$. □

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3.1 Planted cliques and semidefinite programming

Recall the matrix W from before, which has 1s in the top $k \times k$ block, zero on the diagonal, and Rad(1/2) RVs elsewhere.

Recall the spectral method:

$$\hat{u}_{spec} = \operatorname{argmax}_{\substack{u \in \mathbb{R}^n \\ \|u\|^2 = k}} u^\top W u \tag{52}$$

This needs a cleaning step, which we analyzed previously.

How did they come up with this algorithm? Can we get more insight by analyzing htis method in a more principled framework? Yes, through maximum likelihood!

Consider an alterantive model where within clique we have connection probability p (instead of 1) and other connections with probability q (instead of 1/2), where $p \gg q$.

$$\hat{u}_{MLE} = \operatorname{argmax}_{\substack{u \in \{0,1\}^n \\ \sum_i u_i = k}} u^\top W u \tag{53}$$

From this, we see that the spectral method is a continuous relaxation of the MLE integer program. To make this more precise, consider the SDP

$$\hat{X}_{spec} = \operatorname{argmax}_{\substack{X \succeq 0 \\ \operatorname{Tr} X = k}} \langle W, X \rangle \tag{54}$$

If we let $X = uu^\top$, then we automatically have $X \succeq 0$ and additionally we have $\operatorname{Tr} X = \|u\|_2^2$. Thus, the feasible set of Eq. (52) is the same as Eq. (54).

How do we know the optima of Eq. (54) is attained at a rank 1 matrix $X = uu^\top$? Since $X = \sum_i \lambda_i u_i u_i^\top$ ($\lambda_i \geq 0$) and optima are attained at extremal points, by linearity of $\langle W, X \rangle$ we can put all of the weight on a single λ_i corresponding to the top eigenvector of W.

How can we get Eq. (54) closer to Eq. (53)? Since Eq. (53) is more constrained, we can consider adding more constraints:

$$\begin{aligned}
\tilde{X}_{MLE} &= \operatorname{argmax}_X \langle W, X \rangle && (55)\\
\text{s.t. } & X \succeq 0 && (56)\\
& \operatorname{Tr} X = k && (57)\\
& 0 \leq X \leq J \quad \text{entrywise} && (58)\\
& \langle X, J \rangle = k^2 && (59)\\
& \operatorname{rank}(X) = 1 && (60)
\end{aligned}$$

where $J = 11^\top$.

The solution $X = uu^\top$ where $u \in \{0,1\}^n$, where u indexes the clique.

Conversely, we need to show that the feasible set coincides with Eq. (53). If $X \succeq 0$ and $\operatorname{rank} X = 1$, then we can always write $X = uu^\top$. The trace constraint now reads $k = \operatorname{Tr} X = \sum_i u_i^2$. The third constraint becomes $\langle X, J \rangle = k^2 \implies (\sum_i u_i)^2 = k^2$.

Proposition 14

The optima of Eq. (55) must satisfy: $u_i \in [-1, 1]$, $\sum u_i^2 = k$, $(\sum_i u_i)^2 = k^2$, $\{u_i\} \in \{0,1\}^n$ *or* $\{u_i\} \in \{0, -1\}^n$.

In fact, the solution is $u = 1_k$ *or* $u = -1_k$.

The linear constraints in Eq. (55) are fine, but the rank constraints are difficult. Here is an easier candidate SDP:

$$\hat{X}_{SDP} = \text{argmax}_X \langle W, X\rangle \tag{61}$$
$$\text{s.t. } X \preceq 0 \tag{62}$$
$$X \geq 0 \tag{63}$$
$$\operatorname{Tr} X = k \tag{64}$$
$$\langle X, J\rangle = k^2 \tag{65}$$

Notice we have dropped the rank constraint as well as the upper entrywise bound.

Theorem 15

$\exists c > 0$ such that for $k \geq c\sqrt{n}$, Eq. (61) has unique maximizer $X^* = 1_k 1_k^\top$ with high probability.

Proof. We first show X^* is a maximizer.

$$\langle W, X^*\rangle = 1_k^\top W 1_k = k^2 - k \tag{66}$$
$$\langle W, X\rangle = \langle W + I, X\rangle - \operatorname{Tr} X \tag{67}$$
$$\operatorname{Tr}(I - X) = \operatorname{Tr} X \leq \langle J, X\rangle - \operatorname{Tr}(X) \tag{68}$$
$$\underbrace{W + I \leq J}_{X \geq 0} \implies \langle J, X\rangle \geq \langle W + I, X\rangle \tag{69}$$
$$\therefore \operatorname{Tr}(I - X) = \operatorname{Tr} X \leq k^2 - k \tag{70}$$

The harder part is uniqueness. We will develop a general technique called dual certificate / KKT condition. Write the Lagrangian for the optimization problem. Introduce dual variables $S \succeq 0$, $B \geq 0$, $\eta \in \mathbb{R}$, $\lambda \in \mathbb{R}$ and

$$\mathcal{L}(X, S, B, \eta, \lambda) = \langle W, X\rangle + \langle S, X\rangle + \langle B, X\rangle + \eta\left(k \operatorname{Tr}(X) + \lambda(k^2 - \langle X, J\rangle\right) \tag{71}$$

Notice

$$\max_{X \text{ feas}} \langle W, X\rangle = \max_X \min_{S,B,\eta,\lambda} \mathcal{L} \tag{72}$$

as desired. Since $\mathcal{L}$ is linear, by Sion's minimax theorem we have

$$\max_X \min_{S,B,\eta,\lambda} \mathcal{L} = \min_{S,B,\eta,\lambda} \max_X \mathcal{L} \tag{73}$$

Note $\langle S, X\rangle = \operatorname{Tr}(S^{1/2} X S^{1/2}) \geq 0$ is non-negative. $\langle B, X\rangle$ is also trivially non-negative.

Lemma 16

The following conditions imply X^* is the unique maximizer: {lem:x-star-u

1. *Stationarity: $W + S + B - \eta I - \lambda J = 0$ (can't improve any more)*
2. *Primal/dual feasibility*
3. *Complementary slackness: $\langle S, X^*\rangle = 0$ and $\langle B, X^*\rangle = 0$.*
4. *Uniqueness: $\lambda_{n-1}(S) > 0$ (second smallest eigenvalue of S)*

The first three conditions are the "KKT conditions." Together, they guarantee X is a maximizer.

Proof of Lemma 16. X^* **is a maximizer**: for feasible variables

$$\langle W, X\rangle \leq \mathcal{L}(X, S, B, \eta, \lambda) \qquad \text{feasible} \tag{74}$$
$$= \mathcal{L}(X^*, S, B, \eta, \lambda) \qquad \text{stationarity} \tag{75}$$
$$= \langle W, X^*\rangle \qquad \text{comp. slackness} \tag{76}$$

Uniqueness: Suppose X' satisfies $\langle W, X'\rangle = \langle W, X^*\rangle$. Then $\langle S, X'\rangle = 0$, and $\langle S, X^*\rangle = 0 \implies 1_k^\top S 1_k = 0 \implies S1_k = 0$. In other words, 1_k is an eignevector with eigenvalue 0 for S. But condition (4) means that 1_k is the only eigenvector with eigenvalue 0, hence $X' = cX^*$ for some $c \in \mathbb{R}$. But by the constrant $\operatorname{Tr} X = k$, we must have $X' = X^*$. □

Hence, if we can find (S, B, η, λ) satisfying Lemma 16, then we have a certificate that X^* is the unique maximizer.

But how can we find this certificate? It's hard in general, but in this case we have an explicit construction.

$$B \geq 0, \quad \eta \in \mathbb{R}, \quad \lambda \in \mathbb{R} \tag{77}$$

$$S = \eta I + \lambda J - B - W \succeq 0 \tag{78}$$

$$S1_k = 0, \quad \langle B, X^*\rangle = 0, \quad \lambda_{n-1}(S) > 0 \tag{79}$$

$$S1_k = 0 \implies \eta I_k + \lambda k 1 = B1_k + W1_k \tag{80}$$

$X^* = 1_k 1_k^\top$. Since we want $\langle B, X^*\rangle = 0$, we want $B_{ij} = 0$ for $(i,j) \in K \times K$. This implies that $(B1_k)i = 0$ for $i \in K$. Let $y = W1_k$.

ith entry, $i \in K$, of Eq. (79) implies $\eta + k\lambda = (B1_k)_i + y_i = k - 1$. Then, choose $\eta = k - 1 - k\lambda$

Now for $i \notin K$, Eq. (79) implies $\lambda k = (B1_k)_i + y_i$. Construct $B = 1_k b^\top + b 1_k^\top$ for some $b \in \mathbb{R}^n$ such that $b_i = 0$ for $i \in K$. Then $B1_k = kb$.

Fig 9.12.1

$b_i = \lambda - \frac{y_i}{k}$ for all $i \notin k$. Check $B \geq 0 \implies b_i \geq 0$. Since $\lambda \geq \frac{y_i}{k}$ for all $i \in K$, $\lambda \geq \max_{i \notin K} \frac{y_i}{k}$. $y_i = W1_k$ which is a sum of Rad(1/2) RVs, so by concentration for some $\lambda \geq c$ this is satisfied whp.

For the last part, we need to show $x^\top S x > 0$ for all x such that $x^\top 1_k = 0$. The exact formula for S is

$$S = \eta + \underbrace{\lambda x^\top J x}_{\geq O(\sqrt{n})} - \underbrace{x^\top B x}_{=0} - \underbrace{x^\top W x}_{\geq O(\sqrt{n})} \tag{81}$$

$$\geq \frac{k}{2} - 1 - x^\top \mathbb{E}[W] x - \|W - \mathbb{E}W\|_{op} \tag{82}$$

$$\geq 0 \qquad \text{for suff large } k \tag{83}$$

□

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4.1 Logistics

HW1 releasted

4.2 Primal method for SDP

Planted Clique model $G(1/2, n, k)$.

$$\hat{X}_{SDP} = \operatorname{argmax}_X \langle W, X\rangle \tag{84}$$

$$st\ X \succeq 0 \tag{85}$$

$$X \geq 0 \tag{86}$$

$$\operatorname{Tr}(X) = k \tag{87}$$

$$\langle X, J\rangle = k^2 \tag{88}$$

where $J = 11^\top$ and $W_{ij} = \mathbb{1}\{i = j\} 2A_{ij} - 1$. Last time we proved (using a dual certificate approach)

Theorem 17

If $k \geq c\sqrt{n}$ for a large enough c, then $X^* = 1_k 1_k^\top$ is the unique maximizer.

Today we will consider a primal approach.

Round up suffices: Suppose we find X such that $\langle W, X\rangle \geq (1-\varepsilon)\langle W, X^*\rangle$. Let $\hat{X}_{ij} = \mathbb{1}\{X_{ij} > 1/2\}$.

Theorem 18

If $\varepsilon \lesssim \frac{c_0\sqrt{n}}{k^3}$ for sufficiently small $c_0 < 0$, then $\hat{X} = X^*$ whp.

Proof. Suppose $\hat{X} \neq X^*$. Then either:

$\exists (i_0, j_0) \in K \times K$ such that $X^*_{i_0,j_0} = 1$ and $X_{i_0,j_0} \leq \frac{1}{2}$, or

$\exists (i_1, j_1) \notin K \times K$ such that $X^*_{i_1,j_1} = 0$ and $X_{i_1,j_1} > \frac{1}{2}$.

In both acses, $\|X - X^*\|_F \geq \frac{1}{2}$.

Also, we previously showed that the global optimum $\langle W, X^*\rangle = k^2 - k$ because even though W is random, inner product with X^* grabs the upper left $K \times K$ corner where W is deterministic.

Recall the KKT condition: $S \succeq 0$, $S1_K = 0$, $B \geq 0$, $\eta, \lambda \in \mathbb{R}$, $\lambda_{n-1}(S) \geq c_2\sqrt{n}$. ALso

$$\langle W, X^*\rangle - \langle W, X\rangle = \langle S, X\rangle + \langle B, X\rangle =: \delta \tag{89}$$

because last class we had

$$\begin{aligned}
\langle W, X\rangle &\leq L(X, S, B, \eta, \lambda) &(90)\\
&= \langle W, X\rangle + \langle S, X\rangle + \langle B, X\rangle + \eta(k - \operatorname{Tr} X) + \lambda(k^2 - \langle X, J\rangle) &(91)\\
&= \langle W, X^*\rangle &(92)
\end{aligned}$$

We already knew $u = \frac{1}{\sqrt{k}}1_k$ eigenvector of S corresponding to $\lambda_n(S) = 0$ (KKT complementary slackness tells us that $Su = 0$). This gives the matrix inequality

$$S \succeq \lambda_{n-1}(S)(I - UU^\top) \tag{93}$$

Since we previously have a bound on $\langle S, X\rangle$, to look for a sandwich inequality we consider taking an inner product with X

$$\begin{aligned}
\langle S, X\rangle &\geq c_2\sqrt{n}\,\langle X, I - X^*/k\rangle = c_2\sqrt{n}\,\langle X, I\rangle - c_2\frac{\sqrt{n}}{k}\langle X, X^*\rangle &(94)\\
\langle X, X^*\rangle &\geq k^2 - \frac{k\delta}{c_2\sqrt{n}} &(95)
\end{aligned}$$

Where we used the upper bound

$$\delta \geq \langle S, X\rangle \tag{96}$$

This gives a bound on a cross term in the Frobenius norm expansion

$$\begin{aligned}
\|X - X^*\|_F^2 &= \|X\|_F^2 + \|X^*\|_F^2 - 2\langle X, X^*\rangle &(97)\\
\|X^*\|_F^2 &= \|1_k 1_k^\top\|_F^2 = k^2 &(98)\\
\|X\|_F^2 &\leq \|X\|_*^2 = k^2 &(99)\\
\therefore \|X - X^*\|_F^2 &\leq k^2 + k^2 - 2\left(k^2 - \frac{k\delta}{c_2\sqrt{n}}\right) &(100)\\
&= \frac{2k\delta}{c_2\sqrt{n}} \leq \frac{1}{4} &(101)
\end{aligned}$$

□

So we we how to to use approximate KKT conditions. But we need quantitative result of the maximizer (i.e. the second eigenvector $\lambda_{n-1}(S)$) to show the uniqueness of the maximimzer.

4.2.1 SDP Advantage: Robust to monotone adversary

Given adjacency matrix A, allow adversary to delete edges ***not in the clique***.

Failure of spectral methods: they depend too much on edges not in the clique, that by deleting them in a certain way (see Figure) results in their failure.

Figure 9.17.1: spectral methods will fail because there will be two large eigenvalues $\lambda_1 \approx \lambda_2 \approx \frac{n-k}{4}$ corresponding to the ER random blocks and the k-clique will be missed.

In contrast, SDPs enjoy better robust. Consider modification $W \mapsto \tilde{W}$. For any $X \neq X^*$, will show

4.3 Second SDP formulation: primal analysis

This gives another formulation of the same problem, but presents new techniques.

Recall $\operatorname{Tr} X = k = \sum_i \lambda_i(X) = \|X\|_*$ the nuclear norm. We have the SDP formulation

$$\hat{X}_{cvx} = \operatorname{argmax}_X \langle X, W \rangle \tag{102}$$
$$\text{st } \|X\|_* \leq k \tag{103}$$
$$0 \leq X \leq J \tag{104}$$
$$\langle X, J \rangle = k^2 \tag{105}$$

Lemma 19

For any matrix $X \in \mathbb{R}^{m \times n}$, $\|X\|_* \leq 1$ iff $\exists W_1 \in \mathbb{R}^{m \times n}$ and $W_2 \in \mathbb{R}^{n \times n}$ such that $\operatorname{Tr}(W_1) + \operatorname{Tr}(W_2) \leq 2$.

$$\begin{bmatrix} W_1 & X \\ X^\top & W_2 \end{bmatrix} \succeq 0 \tag{106}$$

After this lemmma, we know we can solve the nuclear norm into a PSD constraint and can hence solve this problem with a SDP solver.

Proof. We need the following result:

Lemma 20 (*lSub-differential of nuclear norm*)

$X \neq 0$, $X = U\Sigma V^\top$ and the subgradient for nuclear norm

$$\partial \|\cdot\|_*(X) = \{UV^\top + p^\perp(Y) : \|Y\|_{op} \leq 1\} \tag{107}$$
$$\text{where } p^\perp(Y) = (I - UU^\top)(I - VV^\top) \tag{108}$$

We will show the sufficient condition that for any $X \neq X^*$,

$$\langle W, X^* \rangle - \langle W, X \rangle \gtrsim \|X - X^*\|_{\ell_1} \tag{109}$$

We have $X^* = 1_k 1_k^\top$, with top eigenvector $u = \frac{1}{\sqrt{k}} 1_k$. Analogously, $X^* = kuu^\top$. Letting $E = UU^\top$,

$$p^\perp(Y) = (I - E)Y(I - E) \tag{110}$$
$$p(Y) = Y - P^\perp(Y) = EY + YE - EYE \tag{111}$$

We can decompose

$$\langle W, X^* - X \rangle = \langle X^* - X, X^* \rangle + \langle X^* - X, P^\perp(W - X^*) \rangle + \langle X^* - X, P(W - X^*) \rangle \tag{112}$$

(a)

$$\langle X^* - X \rangle = \sum_{(i,j) \in K \times K} (1 - X_{ij}) = \frac{1}{2}\|X - X^*\|_{\ell_1} \tag{113}$$
$$= \sum_{(i,j) \notin K \times K} (X_{ij} - v) \tag{114}$$

(b)

$$0 \geq \|X\|_* - \|X*_*^\| \tag{115}$$
$$\geq \langle X - X^*, \underbrace{E + p^\perp(Y)}_{\partial\|\cdot\|_*(X^*), \|Y\|_{op}\leq 1} \rangle \tag{116}$$
$$= \langle X - X^*, E\rangle + \langle X - X^*, p^\perp(y)\rangle \tag{117}$$

For the last term, just use Hölder's inequality

$$|\langle X^* - X, P(W - X^*)\rangle| \leq \|P(W - X^*)\|_{\ell_\infty} \|X - X^*\|_{\ell_1} \tag{118}$$

Altogether (remember this, building on this next lecture)

$$\langle X^* - X, W\rangle \geq \left(\frac{1}{2} - \frac{\|W - X^*\|_{op}}{2k} - \|P(W - X^*)\|_{\ell_\infty}\right) \|X - X^*\|_{\ell_1} \tag{119}$$

□

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Recall the SDP relaxation

$$\hat{X}_{cvx} = \text{argmax}_X \langle W, X\rangle \tag{120}$$
$$\text{st } \|X\|_* \leq k \tag{121}$$
$$0 \leq X \leq J = 11^\top \tag{122}$$
$$\langle X, J\rangle = k^2 \tag{123}$$

Theorem 21

If $k \geq c\sqrt{n}$, c sufficiently large, then X^* is the unique maximizer.

Proof. For any feasible X,

$$\langle W, X^*\rangle - \langle W, X\rangle \gtrsim \|X - X^*\|_{\ell_1} \tag{124}$$

□

Last time, defined

$$u = \frac{1}{\sqrt{k}} 1_k \tag{125}$$
$$X^* = 1_k 1_k^\top = k \underbrace{uu^\top}_{=:E} \tag{126}$$
$$P^\perp(Y) = (I - E)Y(I - E) \tag{127}$$
$$P(Y) = Y - P^\perp(Y) = EY + YE - EYE \tag{128}$$

$P^\perp$ is the projection to the orthogonal complement of E, and P is the projection onto E.

We proved last time

$$\langle X - X^*, W\rangle \geq \left(\frac{1}{2} - \frac{\|W - X^*\|_{op}}{2k} - \|P(W - X^*)\|_{\ell_\infty}\right) \|X - X^*\|_{\ell_1} \tag{129}$$

Today, we consider

$$\|W - X^*\|_{op} \leq \underbrace{\|W - EW\|_{op}}_{\lesssim\sqrt{n}} + \underbrace{\|EW - X^*\|_{op}}_{\leq 1} \tag{130}$$

Indeed

$$W - X^* = W - EW - I_k \tag{131}$$
$$\|P(W - X^*)\|_{\ell_\infty} \le \|P(W - EW)\|_{\ell_\infty} + \|P(I_k)\|_{\ell_\infty} \tag{132}$$
$$P(I_k) = EI_k + I_kE - EI_kE = E \tag{133}$$

Also

$$\|P(Y)\|_{\ell_\infty} = \|EY + YE - EYE\|_{\ell_\infty} \tag{134}$$
$$\le \|EY\|_{\ell_\infty} + \|YE\|_\infty + \|EYE\|_\infty \tag{135}$$

The last term is complicated, but notice $\|EYE\|_\infty \le \|EY\|_\infty \|E\|_{\ell_\infty \to \ell_\infty} \le \|EY\|_\infty$ hence

$$\|P(Y)\|_{\ell_\infty} \le 3\|EY\|_{\ell_\infty} \tag{136}$$

Doing the calculation for $\|EY\|_\infty$

$$EY = \frac{1}{k}\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}\begin{pmatrix} 0 & \text{Rad} \\ \text{Rad} & 0 \end{pmatrix} \tag{137}$$

So $\|EY\|_\infty = \frac{1}{k}\max_{j \notin K}\sum_{i \in K} Y_{ij}$.

$n - k$ sub-Gaussian rv with variance $1/k$.

Lemma 22

If X_i satisfies $\mathbb{E}e^{-x_i^2/\sigma^2} \le 2$ for some σ, then

$$\mathbb{E}\max_{i=1}^{n} \lesssim \sigma\sqrt{\log n} \tag{138}$$

5.1 Planted partition model

Let $A_{ij} \sim \begin{cases} P, & \text{if } \sigma_i = \sigma_j \\ Q, & \text{ow} \end{cases}$ with $\sigma = (\sigma_1, \ldots, \sigma_n) \in \{\pm 1\}^n$.

Goal: Recover σ.

Stochastic block model: $P = \text{Bern}(p)$ and $Q = \text{Bern}(q)$. If $p > q$ we call it ***associative*** and $p < q$ is called ***disassociative***.

IID model: $\sigma_i \overset{\text{iid}}{\sim} \text{Rad}$

Bisection: $\sum \mathbb{1}\{\sigma_i = +1\} = \sum \mathbb{1}\{\sigma_i = -1\}$

Some problems we are interested in solving include ***detection***:

$$\mathcal{H}_0 : A_{ij} \overset{\text{iid}}{\sim} \frac{P+Q}{2} \tag{139}$$
$$\mathcal{H}_1 : \text{Planted partition model} \tag{140}$$

Lemma 23

(X, Y) with $Y \in \{\pm 1\}$.

$P_{X|Y=1} = P$ and $P_{X|Y=-1} = Q$.

$P_Y(1) = P_Y(-1) = \frac{1}{2}$.

Observe X, infer Y?

$$\min_{\hat{Y}(X)} \mathbb{E}\mathbb{1}\{\hat{Y} \ne Y\} = \frac{1}{2}(1 - \text{TV}(P, Q)) \tag{141}$$

Another problem is ***correlated recovery***

$$\ell(\sigma, \hat{\sigma}) = \min_{s \in \{\pm 1\}} \|\sigma + s\hat{\sigma}\|_1 \tag{142}$$

If I beat random guess, I win.

Yet another is ***almost exact recovery***

$$\frac{\mathbb{E}\ell(\sigma, \hat{\sigma})}{n} \to 0 \tag{143}$$

Finally in ***exact recovery***

$$\Pr[\sigma \neq \hat{\sigma}] \to 0 \tag{144}$$

Computing TV is not easy usually. ***Ingster-Suslina Trick*** lets us upper bound it with chi squared divergence:

$$\chi^2(P \mid\mid Q) = \left(\int \frac{p^2}{q}\right) - 1 \geq 0 \tag{145}$$

$$\mathrm{TV}(P, Q) \lesssim \sqrt{KL(P \mid\mid Q)} \leq \sqrt{\chi^2(P \mid\mid Q)} \tag{146}$$

Mixture vs single: suppose $\{P_\theta : \theta \in \Theta\}$ family of models, prior Π on Θ,

$$P_\Pi(x) = \int P_\theta(x)\Pi(d\theta) \tag{147}$$

Then sometimes it's easy to write down

$$\chi^2(P_\Pi \mid\mid Q) = \mathbb{E}_{\theta, \hat{\theta}, \Pi} G(\theta, \hat{\theta}) - 1 \tag{148}$$

$$G(\theta, \hat{\theta}) = \int \frac{P_\theta P_{\tilde{\theta}}}{Q} \tag{149}$$

Proof. By Fubini

$$\int \frac{P_\Pi^2}{Q} = \int \frac{\int p_\theta(x)\pi(d\theta) \int p_{\hat{\theta}}(x)\pi(d\hat{\theta})}{Q(x)} dx \tag{150}$$

$$= \int \pi(d\theta)\pi(d\hat{\theta}) \left(\frac{P_\theta(x) P_{\hat{\theta}}(x)}{Q(x)}\right) dx \tag{151}$$

□

5.2 Contiguity between probability measures

Introduced by LeCun in the asymptotic statistics literature.

Definition 24

A sequence of probability measures (p_n) is ***contiguous to*** (Q_n) if for any events E_∞,

$$Q_n(E_n) \to 0 \implies P_n(E_n) \to 0 \tag{152}$$

This can be thought of as an asymptotic version of absolute continuity: $P \ll Q$ if for all events E

$$Q(E) = 0 \implies P(E) = 0 \tag{153}$$

To interpret contiguity, let E_n be set X lies in to declare p_n sequence.

$$P_n(E_n) = \mathbb{E}_{Q_n}\left(\frac{P_n}{Q_n}\mathbb{1}(E_n)\right) \tag{154}$$

$$\leq \sqrt{\mathbb{E}_{Q_n}\left(\frac{P_n^2}{Q_n^2}\right)\mathbb{E}_{Q_n}[\mathbb{1}(E_n)]} \tag{155}$$

SBM: Fix label σ.

$$P_\sigma(A) = \prod_{i<j}\left(P\mathbb{1}_{\sigma_i=\sigma_j} + Q\mathbb{1}_{\sigma_i\neq\sigma_j}\right) \tag{156}$$

$$= \prod_{j<j}\left(\frac{P+Q}{2} + \frac{P-Q}{2}\sigma_i\sigma_j\right) \tag{157}$$

$$G(\sigma,\hat{\sigma}) = \int \frac{P_\sigma(A)P_{\hat{\sigma}}(A)}{P_0(A)}dA \tag{158}$$

$$P_0(A) = \prod_{i<j}\frac{P+Q}{2} \tag{159}$$

$$= \prod_{i<j}\left(\int\frac{P+Q}{2} + \int\frac{P-Q}{2}\sigma_i\sigma_j + \int\frac{P-Q}{2}\hat{\sigma}_i\hat{\sigma}_j + \int\underbrace{\frac{(P-Q)^2}{2(P+Q)}}_{=:\rho}\sigma_i\sigma_j\hat{\sigma}_i\hat{\sigma}_j\right) \tag{160}$$

$$= \prod_{i<j}(1+\rho\sigma_i\sigma_j\hat{\sigma}_i\hat{\sigma}_j) \tag{161}$$

$$\leq \exp(\rho\sum_{i<j}\sigma_i\sigma_j\hat{\sigma}_i\hat{\sigma}_j) \tag{162}$$

$$\leq \exp(\frac{\rho}{2}\langle\sigma,\hat{\sigma}\rangle^2) \tag{163}$$

But we know the last term very well. Since $\sigma,\hat{\sigma} \overset{\text{iid}}{\sim} \text{Rad}^n$, we have $\frac{1}{\sqrt{n}}\langle\sigma,\hat{\sigma}\rangle \Rightarrow \mathcal{N}(0,1)$ so

$$\mathbb{E}e^{\frac{\rho}{2}\langle\sigma,\hat{\sigma}\rangle^2} \to \mathbb{E}e^{\frac{\rho}{2}(\sqrt{n}z)^2} = \mathbb{E}e^{\frac{\rho n}{2}z^2} < \infty \tag{164}$$

whenever $\rho_n < 1$. So we have the lower bound

$$\rho = \frac{\tau+o(1)}{n} \quad \tau = \frac{(a-b)^2}{2(a+b)} \tag{165}$$

When $\tau < 1$, then it is impossible to detect.

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6.1 Exact recovery of stochastic block model

Definition 25 (*Symmetric stochastic block model*) {def:9-24-ssbm}

The ***symmetric stochastic block model***, denoted by $SSBM(n, 2, p_{in} = \frac{a\log n}{n}, p_{out} = \frac{b\log n}{n} \mid \sigma)$, is a probability distribution over graphs (V,E) on n vertices where:

- Each vertex $v \in V$ belongs to one of 2 communities, denoted by $\sigma_v \in \{1,2\}$
- **Symmetric**: exactly $n/2$ vertices in each community
- The probability of an edge between two vertices in the same community is $p_{in} = \frac{a\log n}{n}$

- The edge probability between different communities is p_{out}.

Notice that we have chosen to parameterize $p_{in} = \frac{a \log n}{n}$ and $p_{out} = \frac{b \log n}{n}$. Some intuition for the log is to recall that $G(n, c \log n/n)$ is connected whp iff $c > 1$. For SSBM, we have a similar threshold where G is connected whp iff the average of the edge probability coefficients $\frac{a+b}{2} > 1$.

We are interested in ***exact recovery in SSBM***: let $G = (V, E) \sim SSBM(n, 2, p_{in}, p_{out} \mid \sigma^*)$, can we construct an estimator $\hat{\sigma}(G)$ such that as $n \to \infty$

$$\Pr[\sigma^* \neq \hat{\sigma}] \to 0 \tag{166}$$

The goal over the next lectures will be to establish the following phase transition regarding the hardness of exact recovery in SSBM:

Theorem 26

Exact recovery in $SSBM(n, 2, \frac{a \log n}{n}, \frac{b \log n}{n})$ *is efficiently solvable if* $|\sqrt{a} - \sqrt{b}| > \sqrt{2}$ *and unsolvable if* $|\sqrt{a} - \sqrt{b}| < \sqrt{2}$.

Remark 27. We can rewrite $|\sqrt{a} - \sqrt{b}| > \sqrt{2}$ as $\frac{a+b}{2} > 1 + \sqrt{ab}$ and compare against the $\frac{a+b}{2} > 1$ connectivity threshold for SSBM. As expected, exact recovery implies connectivity. Furthermore, exact recovery requires a $\sqrt{ab}$ over-sampling factor.

Remark 28. For $|\sqrt{a} - \sqrt{b}| = \sqrt{2}$, exact recovery is efficiently solvable if $a, b > 0$.

Proof of unsolvable. Consider the one dimensional problem of oracle-aided hypothesis testing problem where the oracle reveals the true communities σ_v of all vertices except for one, say σ_0, and we test $\mathcal{H}_0 = \{\sigma_0 = 1\}$ against $\mathcal{H}_a = \{\sigma_0 = 2\}$.

The probability of error is minimized by the MAP estimator, which picks $\sigma_0 = u$ maximizing the posterior probability

$$\Pr[\sigma_0 = u \mid G = g, \sigma_{\setminus 0} = \sigma_{\setminus 0}] \tag{167}$$

Since $P(\sigma_0 = u) = 1/2$ for $u \in \{1, 2\}$, the posterior probability is

$$\Pr[\sigma_0 = u \mid G = g, \sigma_{\setminus 0} = \sigma_{\setminus 0}] = \frac{\overbrace{\Pr[\sigma_0 = u]}^{=1/2} \Pr[G = g, \sigma_{\setminus 0} = \sigma_{\setminus 0} \mid \sigma_0 = u]}{\Pr[G = g, X_{\setminus 0} = x_{\setminus 0}]} \tag{168}$$

$$\propto \Pr[G = g, \sigma_{\setminus 0} = \sigma_{\setminus 0} \mid \sigma_0 = u] \tag{169}$$

which depends only on the number of edges between vertex 0 and the two communities.

Let $T = \#\{v \in V \setminus \{0\} : \sigma_v = 1 \text{ and } (0, v) \in E\}$ count the number of edges between vertex 0 and all the vertices in community 1 (provided by the oracle through $\sigma_{\setminus 0}$). Notice $T \mid \sigma_0 = 1 \sim B(n/2, p_{in})$ and $T \mid \sigma_0 = 2 \sim B(n/2, q_{out})$, so the error probability for a hypothesis test using T is bounded as

$$p_e \leq P(B(n/2, p_{in}) \leq B(n/2, p_{out})) \tag{170}$$

$$= n^{-\left(\frac{\sqrt{a} - \sqrt{b}}{\sqrt{2}}\right)^2 + o(1)} \tag{171}$$

We will spend the remainder of this lecture showing that exact recovery is not solvable if $np_e \to \infty$. □

Important intuition: Let $X = (X_1, \ldots, X_n) \overset{\text{iid}}{\sim} P$ or Q, $\mathcal{H}_0$ be the hypothesis that the samples are from P, and $\mathcal{H}_1$ that they are from Q. The minimum probability of error (under an equally probable prior) is

$$\frac{1}{2}\left(1 - \text{TV}(p^{\otimes n}, q^{\otimes n})\right) \tag{172}$$

To bound this quantity, there is a (not commonly used) Chernoff bound of

$$\text{TV}(p^{\otimes n}, q^{\otimes n}) = 1 - e^{-nc(P,Q) + o(n)} \tag{173}$$

where $c(P, Q) = -\log \inf_{\alpha \in [0,1]} \int p^\alpha q^{1-\alpha}$.

We will instead be concerned with bounds involving a different discrepancy metric.

Definition 29 (***Squared hellinger distance***)

The ***squared Hellinger distance***

$$H^2(P,Q) = \mathbb{E}_Q\left[\left(1-\sqrt{\frac{P}{Q}}\right)^2\right] \geq 0 \tag{174}$$

$$= \mathbb{E}_Q\left[1+\frac{P}{Q}-2\sqrt{\frac{P}{Q}}\right] \tag{175}$$

$$= 1+1-2\int\sqrt{PQ} = 2\left(1-\int\sqrt{PQ}\right) \tag{176}$$

It sandwiches total variation distance in the following sense:

$$0 \leq \frac{1}{2}H^2(P,Q) \leq \mathrm{TV}(P,Q) \leq H(P,Q)\sqrt{1-\frac{H^2}{4}} \leq 1 \tag{177}$$

Lemma 30

For any sequence $\{p_n\}$, $\{q_n\}$, *as* $n \to \infty$

$$\mathrm{TV}(p_n^{\otimes n}, q_n^{\otimes n}) \to 0 \iff H^2(p_n,q_n) = o(1/n) \tag{178}$$

$$\mathrm{TV}(p_n^{\otimes n}, q_n^{\otimes n}) \to 1 \iff H^2(p_n,q_n) = \omega(1/n) \tag{179}$$

So H^2 *provides us with*

Without loss of generality, let $C_1 = [1 : n/2] = \{v : (\sigma_0)_v = 1\}$ and $C_2 = [n/2+1 : n] = \{v : (\sigma_0)_v = 2\}$ where σ_0 are the true labels. Let $G \sim P_{G|\sigma}(\cdot \mid \sigma_0)$ be the SSBM graph generated from this community assignment.

Definition 31 (***Bad pairs***)

For a community assignment $\sigma \in \{0,1\}^n$, let $\sigma[u \leftrightarrow v]$ denote σ except with the community assignments for u and v swapped.

The ***bad pairs*** of vertices are

$$\mathcal{B}(G) = \{(u,v) : u \in C_1, v \in C_2, \Pr_{G|\sigma}[G \mid \sigma_0] \leq \Pr_{G|\sigma}[G \mid \sigma_0[u \leftrightarrow v]] \tag{180}$$

The reason why these pairs are bad is because if $(u,v) \in \mathcal{B}(G)$ then the MAP estimator would assign greater probability to the incorrectly swapped $\sigma_0[u \leftrightarrow v]$ labels than the true σ_0 labels, therefore:

Corollary 32

If $\mathcal{B}(G)$ *is non-empty with non-vanishing probability, then exact recovery is not possible.*

To characterize the bad vertices involved in bad pairs, notice that swapping vertices u and v flips the edge probabilities $p_{out} \leftrightarrow p_{in}$ for all the edges containing u and v ***except*** for the (u,v) edge (if it exists). When $p_{in} > p_{out}$, we have

$$\Pr_{G|\sigma}[G \mid \sigma_0] \leq \Pr_{G|\sigma}[G \mid \sigma_0[u \leftrightarrow v]] \iff d_+(u) + d_+(v) \leq d_-(u \setminus v) + d_-(v \setminus u) \tag{181}$$

This motivates the following definition:

Definition 33 (***Bad vertices for each community***)

For $i \in \{1,2\}$, the ***bad vertices within community*** i are

$$\mathcal{B}_i(G) = \{u \in C_i : d_+(u) \leq d_-(u) - 1\} \tag{182}$$

where $d_+(u) = \#\{\text{edges } u \text{ has in its own comunity}\}$ and $d_-(u)$ similarly but with the other community.

Notice if $u \in \mathcal{B}_1(G)$ and $v \in \mathcal{B}_2(G)$, then

$$d_+(u) + d_+(v) \leq d_-(u) + d_-(v) - 2 \leq d_-(u \setminus v) + d_-(v \setminus u) \tag{183}$$

and therefore $(u, v) \in \mathcal{B}(G)$ and exact recovery fails.

Lemma 34

$\sqrt{a} - \sqrt{b} < \sqrt{2} \implies \Pr[\exists u \in \mathcal{B}_1(G)] = 1 - o(1)$

Let $\mathcal{B}_u = \mathbb{1}(d_+(u) \leq d_-(u) - 1)$.

$$\Pr[\forall u \in c_I, u \notin \mathcal{B}_1(G)] = \Pr[\sum_{u=1}^{n/2} \mathcal{B}_u = 0] \leq ? \tag{184}$$

Theorem 35 (*Paley-Zygmund Inequality*)

Let $X \geq 0$, $0 < \mathbb{E}X^2 < \infty$. For any $c \in [0,1]$

$$\Pr[X > c\mathbb{E}[X] \geq (1-c)^2 \frac{(\mathbb{E}X)^2}{\mathbb{E}X^2} \tag{185}$$

Some intuition for Paley-Zgymund: Figure 9.24.1
Applying Paley-Zygmund on the complement event with $c = 0$.

$$\Pr[\forall u \in c_I, u \notin \mathcal{B}_1(G)] = \Pr[\sum_{u=1}^{n/2} \mathcal{B}_u = 0] \leq \frac{\mathrm{Var}(\sum \mathcal{B}_u)}{\mathbb{E}(\sum \mathcal{B}_u)^2} \tag{186}$$

$$nP(B_1 = 1) + \frac{n(n-1)}{2} P(B_1 = 1, B_2 = 1) + \frac{n^2}{2} P(B_1 = 1, B_{n/2+1} = 1) \tag{187}$$

$$\begin{aligned} P(B_1 = 1 \mid B_2 = 1) &= P(d_+(1) \leq d_-(1) - 1 \mid d_+(2) \leq d_-(2) - 1) && (188)\\ &= P(B(n/2-2, q_{in}) + B_{1,2} \leq B(n/2, q_{out}) - 1 && (189)\\ &\qquad \mid B'(n/2-2, q_{in}) + B_{12} \leq B'(n/2, q_{out}) - 1) && (190) \end{aligned}$$

7 9/26/2019

7.1 Spectral method for exact recovery of SSBM

Last time we showed regime for non-solvability of SSBM. Today we will see how a spectral method can be used to show solvability of exact recovery in SSBM.

Theorem 36

Exact recovery in $SSBM(n, 2, p = a\log n/n, q = b\log n/n)$ is efficiently solvable if $|\sqrt{a} - \sqrt{b}| > \sqrt{2}$ using a spectral method.

Algorithm:

- Form the modified adjacency matrix A' by adding self loops with probability p to the original adjacency matrix. Then $\mathbb{E}A' = n\frac{p+q}{2}\bar{\phi}_1\bar{\phi}_1^\top + n\frac{p-q}{2}\bar{\phi}_2\bar{\phi}_2^\top$ where

$$\bar{\phi}_1 = \frac{1}{\sqrt{n}}\mathbf{1} \qquad \bar{\phi}_2 = \frac{1}{\sqrt{n}}\begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \\ -1 \\ -1 \\ \vdots \\ -1 \end{bmatrix} \tag{191}$$

- Define $A = A' - n\frac{p+q}{2}\bar{\phi}_1\bar{\phi}_1^\top$
- Solve largest eigenvector problem: $A\phi = \lambda\phi$.
- Return labels $X_{spec}(i) = 1\mathbb{1}\{\phi(i) \geq 0\} + 2\mathbb{1}\{\phi(i) < 0\}$.

Define $\bar{\phi}$ and $\bar{\lambda}$ by

$$\mathbb{E}A = n\frac{p-q}{2}\bar{\phi}_2\bar{\phi}_2^\top := \bar{\lambda}\bar{\phi}\bar{\phi}^\top \tag{192}$$

Lemma 37

$\Pr[\|A - \bar{A}\|_2 \geq c_1\sqrt{\log n}] \leq c_2 n^{-3}$, *where* c_1 *and* c_2 *depend on* a *and* b.

Lemma 38 (*General version of above*)

Let A *be a symmetric zero-diagonal matrix with* $\{A_{ij} : i < j\}$ *independent,* $[0,1]$*-valued,* $\mathbb{E}A_{ij} \leq p$, $\frac{c_0 \log n}{n} \leq p \leq 1 - c_1$.

Then, for any $c > 0$, $\exists c' > 0$ *such that*

$$\Pr[\|A - \mathbb{E}A\|_2 \leq c'\sqrt{np}] \geq 1 - n^{-c} \tag{193}$$

Remark 39. The above result is different than what we have seen before. Davis-Kahan gives $\langle \phi, \bar{\phi} \rangle = 1 - o(1)$, Latala gives weaker bound beacuse of 4th moment requirement.

Instead, we will compare ϕ with $A\bar{\phi}/\bar{\lambda}$ instead of $\bar{\phi} = \bar{A}\bar{\phi}/\bar{\lambda}$.

Lemma 40

$\exists$ *constant* $C(a, b)$ *such that as* $n \to \infty$

$$\Pr\left[\min_{s\in\{\pm 1\}} \|s\phi - A\bar{\phi}/\bar{\lambda}\|_\infty \leq \frac{c}{\sqrt{n}\log\log n}\right] \geq 1 - \frac{c}{n^2} \tag{194}$$

Proof assuming lemma. Define events

$$\mathcal{E}_1 = \left\{\min_{i\in[1:n/2]} (A\bar{\phi}/\bar{\lambda})_i \geq \frac{2\varepsilon}{(a-b)\sqrt{n}}, \max_{i\in[n/2+1:n]} (A\bar{\phi}/\bar{\lambda})_i \leq \frac{-2\varepsilon}{(a-b)\sqrt{n}}\right\} \tag{195}$$

$$\mathcal{E}_2 = \left\{\min_{s\in\{\pm 1\}} \|s\phi - A\bar{\phi}/\bar{\lambda}\|_\infty \leq \frac{c}{\sqrt{n}\log\log n}\right\} \tag{196}$$

Claim: if $\Pr[\mathcal{E}_1 \cap \mathcal{E}_2] \to 1$, then problem solved.

$(A\bar{\phi}/\bar{\lambda})_i \sim B(n/2,p) - B(n/2,q)$ because $\bar{\phi}$ has its first $n/2$ entries $+1$ and second $n/2$ entries -1. Furthermore, since (see last time)

$$\Pr[B(n/2,p) - B(n/2,q) \geq O(1)] = n^{-\left(\frac{\sqrt{a}-\sqrt{b}}{\sqrt{2}}\right)^2 - o(1)} \tag{197}$$

Since we are in regime $\sqrt{a} - \sqrt{b} > \sqrt{2}$, by union bound

$$\Pr\left[\exists i : (A\bar{\phi}/\bar{\lambda})_i \leq \frac{2\varepsilon}{(a-b)\sqrt{n}}\right] \leq nn^{-1-\Omega(1)} = n^{-\Omega(1)} \tag{198}$$

A similar argument handles the $i \in [n/2+1 : n]$ to conclude $\Pr[\mathcal{E}_1] \to 1$. The lemma handles $\mathcal{E}_2$. □

Proof of lemma. Choose ϕ such that $\phi^\top \bar{\phi} \geq 0$.

$$\|\phi - A\bar{\phi}/\bar{\lambda}\|_\infty \leq \|\phi - A\phi/\bar{\lambda}\|_\infty + \|A\phi/\bar{\lambda} - A\bar{\phi}/\bar{\lambda}\|_\infty \tag{199}$$

$$= \|\phi - \lambda/\bar{\lambda} \cdot \phi\|_\infty + \|\frac{A}{\bar{\lambda}}(\phi - \bar{\phi})\|_\infty \tag{200}$$

$$= \frac{|\lambda - \bar{\lambda}|}{\bar{\lambda}}\|\phi\|_\infty + \frac{1}{\bar{\lambda}}\|A(\phi - \bar{\phi})\|_\infty \tag{201}$$

Condition on event $\|A - \bar{A}\|_2 \lesssim \sqrt{\log n}$, by Davis-Kahan $|\lambda - \bar{\lambda}| \leq \|A - \mathbb{E}A\|_2 \lesssim \sqrt{\log n}$, and by definition $\bar{\lambda} \asymp \log n$, so the first term is bounded like $\frac{\|\phi\|_\infty}{\sqrt{\log n}}$.

The second term is more complicated. Define n auxiliary matrices (A delete row/col m)

$$(A_{ij}^{(m)}) = A_{ij}\delta_{i\neq m, j\neq m} \tag{202}$$

Let $\phi^{(m)}$ be the leading eigenvector of $A^{(m)}$ and note $(\phi^{(m)})^\top \bar{\phi} \geq 0$. We defined it like this so

$$(A(\phi - \bar{\phi}))_m = A_m(\phi - \bar{\phi}) = A_m(\phi - \phi^{(m)}) + A_m(\phi^{(m)} - \bar{\phi}) \tag{203}$$

where A_m is the mth row of A. Focusing on the first term for now:

$$|A_m(\phi - \phi^{(m)})| \leq \|A_m\|_2 \|\phi - \phi^{(m)}\|_2 \tag{204}$$

$$\leq \|A\|_{2\to\infty} \|\phi - \phi^{(m)}\|_2 \tag{205}$$

We're going to show the following:

$$\|A_m\|_2 \|\phi - \phi^{(m)}\|_2 \leq \sqrt{\log n}\|\phi\|_\infty \tag{206}$$

The intuition for this is that we want to first use Davis-Kahan for $\|\phi - \phi^{(m)}\|_2$,

$$\|A_m\|_2 = \|A - A^{(m)}\|_2 \leq \|A^{(m)} - A\|_F \leq \sqrt{2}\|A\|_{2\to\infty} =: \max_i \|A_i\|_2 \leq \|A\|_2 \tag{207}$$

$$\|A\|_{2\to\infty} \leq \|A - \bar{A}\|_{2\to\infty} + \|\bar{A}\|_{2\to\infty} \tag{208}$$

$$\lesssim \sqrt{\log n} + \frac{\log n}{\sqrt{n}} \lesssim \sqrt{\log n} \tag{209}$$

By Davis-Kahan

$$\min_{s\in\{\pm 1\}} \|s\phi - \phi^{(m)}\|_2 \lesssim \frac{\|A^{(m)} - A\|_2}{\bar{\lambda}} \lesssim \frac{1}{\sqrt{\log n}} \tag{210}$$

Here the maximum is attained at $s = 1$. To see this, recaall old davis-kahan to see

$$\min_s \|su - v\|_2 \lesssim \frac{\|A - B\|_2}{\max(\lambda_1(A) - \lambda_2(B), \lambda_1(B) - \lambda_2(A))} \tag{211}$$

$$\min_s \|su - v\|_2 \lesssim \frac{\|(A - B)u\|}{\text{max eigengap}} \tag{212}$$

Here is a new version of it we will need

$$\|\phi^{(m)} - \phi\|_2 \lesssim \frac{\|(A^{(m)} - A)\phi\|_2}{\bar{\lambda}} \tag{213}$$

$$\|(A^{(m)} - A)\phi\|_2 = \sqrt{\lambda^2|\phi_m|^2 + \sum_{i\neq m} A_{im}^2\phi_m^2} \leq |\phi_m|\sqrt{\lambda^2 + \|A\|_{2\to\infty}^2} \lesssim \bar{\lambda}|\phi_m| \tag{214}$$

$\|\phi^{(m)} - \phi\|_{\lesssim}|\phi_m| \lesssim \|\phi\|_\infty$.

□

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