	ACM ICPC		2.12 sparse-table	5 5	5	Graj	ph 2-SAT	11 11		6.18 Primes	
_	D 0		2.14 static-range-lis	6		5.2	BCC-tree			6.20 Xor-Basis	
Ί.	Team Reference	_	2.15 treap	6		5.3	centroid-tree			6.21 估計值	
_			2.16 union-of-rectangles	7		5.4	chromatic-number			6.22 定理	
	NTHU		2.17 wavelet-tree	7		5.5	HLD			6.23 數字	
	NIIIO					5.6	lowlink			6.24 歐幾里得類算法	
	I in 1-C4Tuese	3	· · · · · · · · · · · · · · · · ·	7		5.7	manhattan-mst			6.25 生成函數	
LinkCutTreap			3.1 bipartite-matching	7		5.8	SCC				
	•		3.2 Dinic-LowerBound	7		5.9	triangle-sum		7	Misc	17
			3.3 Dinic	8			_			7.1 fast	17
Contents			3.4 Flow 建模	8	6	Mat	h	13		7.2 for	17
						6.1	Aliens			7.3 next-combination	17
				0		6.2	Berlekamp-Massey			7.4 PBDS	17
1	Basic 1.1 vimrc	2	matching	8		6.3	Chinese-Remainder			7.5 python	18
-		2	and the second s	10		6.4	Combination			7.6 rng	18
				10		6.5	Determinant	14		7.7 rotate90	18
2	Data-Structure	2	3.10 minimum-general-weighted-	10		6.6	Discrete-Log	14		7.8 timer	18
	2.1 CDQ	2	perfect-matching	10		6.7	extgcd				
	2.2 CHT	2	periect matering	10		6.8	Floor-Sum			String	18
	2.3 DLX	2 4	Geometry	10		6.9	FWHT			8.1 AC	
	2.4 fast-set	2	4.1 closest-pair	10			Gauss-Jordan			8.2 KMP	
	2.5 lazysegtree	3	4.2 convex-hull				GCD-Convolution			8.3 LCP	
	2.6 LCT	3	4.3 half-plane	11			Int-Div			8.4 manacher	
	2.7 LiChao	4	4.4 min-enclosing-circle	11			Linear-Sieve	15		8.5 rolling-hash	
	2.8 rect-add-rect-sum	4	4.5 point-in-convex-hull	11			Miller-Rabin	15		8.6 SAIS	
	2.9 rollback-dsu	4	4.6 point	11			Min-of-Mod-of-Linear	15		8.7 SAM	
	2.10 segtree-beats	4	4.7 polar-angle-sort				Mod-Inv	15		8.8 smallest-rotation	
	2.11 segtree	5	4.8 定理	11		6.17	Pollard-Rho	16		8.9 Z	19

Basic

1.1 vimre

```
| se nu ai hls et ru ic is sc cul
2 se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
3 syntax on
4 hi cursorline cterm=none ctermbg=89
5 set bg=dark
6 inoremap {<CR> {<CR>}<Esc>ko<tab>
```

Data-Structure

CDO

```
void CDQ(int 1, int r) {
   if(1 + 1 == r) return;
   int mid = (1 + r) / 2;
   CDQ(1, mid), CDQ(mid, r);
   int i = 1;
   for(int j = mid; j < r; j++) {</pre>
     const Q& q = qry[j];
     while(i < mid && qry[i].x >= q.x) {
       if(qry[i].id == -1) fenw.add(qry[i].y,
             qry[i].w);
       i++:
     if(q.id >= 0) ans[q.id] += q.w * fenw.
          sum(q.y, sz - 1);
   for(int p = 1; p < i; p++) if(qry[p].id ==</pre>
         -1) fenw.add(qry[p].y, -qry[p].w);
   inplace_merge(qry.begin() + 1, qry.begin()
         + mid, qry.begin() + r, [](const Q& a
        , const Q& b) {
     return a.x > b.x;
   });
```

2.2 CHT

```
|| struct line t {
    mutable 11 k, m, p;
   bool operator<(const line_t& o) const {</pre>
         return k < o.k; }</pre>
   bool operator<(ll x) const { return p < x;</pre>
 };
6 template < bool MAX>
 struct CHT : multiset<line t, less<>>> {
   const ll INF = 1e18L:
   bool isect(iterator x, iterator y) {
      if(y == end()) return x->p = INF, 0;
      if(x->k == y->k) {
        x \rightarrow p = (x \rightarrow m \rightarrow y \rightarrow m ? INF : -INF);
```

```
x \rightarrow p = floor div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y) 34
              ->k): // see Math
       return x - p >= y - p;
     void add line(ll k, ll m) {
       if(!MAX) k = -k, m = -m;
       auto z = insert(\{k, m, 0\}), y = z++, x =
       while(isect(y, z)) z = erase(z);
       if(x != begin() && isect(--x, y)) isect(
            x, y = erase(y));
       while((y = x) != begin() && (--x)->p >=
            y->p) isect(x, erase(y));
    11 get(11 x) {
       assert(!empty());
       auto 1 = *lower bound(x);
       return (1.k * x + 1.m) * (MAX ? +1 : -1) 52
30 };
```

2.3 DLX

1 struct DLX {

21

29

```
int n, m, tot, ans;
vi first, siz, L, R, U, D, col, row, stk;
DLX(int _n, int _m) : n(_n), m(_m), tot(_m
    ) {
  int sz = n * m:
  first = siz = L = R = U = D = col = row
       = stk = vi(sz);
  REP(i, m + 1) {
    L[i] = i - 1, R[i] = i + 1;
    U[i] = D[i] = i;
  L[0] = m, R[m] = 0;
void insert(int r, int c) {
  r++, c++;
  col[++tot] = c, row[tot] = r, ++siz[c];
  D[tot] = D[c], U[D[c]] = tot, U[tot] = c
       . D[c] = tot:
  if(!first[r]) first[r] = L[tot] = R[tot] 11
        = tot;
    L[R[tot] = R[first[r]]] = tot;
    R[L[tot] = first[r]] = tot;
#define TRAV(i, X, j) for(i = X[j]; i != j
     ; i = X[i]
void remove(int c) {
  int i, j;
  L[R[c]] = L[c], R[L[c]] = R[c];
  TRAV(i, D, c) TRAV(j, R, i) {
    D[U[D[j]] = U[j]] = D[j];
    siz[col[j]]--;
void recover(int c) {
```

2.4 fast-set

int i, j;

remove(c);

recover(c);

vi solve() {

return false;

ans);

TRAV(i, D, c) {

38

57

58 };

TRAV(i, U, c) TRAV(j, L, i) {

if(!R[0]) return ans = dep, true;

TRAV(j, R, i) remove(col[j]);

if(dance(dep + 1)) return true;

return vi(stk.begin() + 1, stk.begin() +

TRAV(j, L, i) recover(col[j]);

TRAV(i, R, 0) if (siz[i] < siz[c]) c = i; 29

U[D[j]] = D[U[j]] = j;

siz[col[j]]++;

bool dance(int dep) {

int i, j, c = R[0];

stk[dep] = row[i];

if(!dance(1)) return {};

L[R[c]] = R[L[c]] = c;

```
1 // Can correctly work with numbers in range
       [0; MAXN]
2 // Supports all std::set operations in O(1)
      on random queries / dense arrays, O(
       log 64(N)) in worst case (sparce array). 51
3 // Count operation works in O(1) always.
                                                 52
4 template < uint MAXN>
                                                 53
5 class fast set {
6 private:
   static const uint PREF = (MAXN <= 64 ? 0 :</pre>
                  MAXN <= 4096 ? 1 :
                  MAXN <= 262144 ? 1 + 64 :
                                                 57
                  MAXN <= 16777216 ? 1 + 64 +
                  MAXN <= 1073741824 ? 1 + 64
                       + 4096 + 262144 : 227) +
    static constexpr ull lb(int x) {
     if(x == 64) return ULLONG MAX;
     return (1ULL << x) - 1;</pre>
   static const uint SZ = PREF + (MAXN + 63)
        / 64 + 1;
   ull m[SZ] = \{0\};
   inline uint left(uint v) const { return (v 65
          - 62) * 64; }
   inline uint parent(uint v) const { return
                                                67
        v / 64 + 62; }
   inline void setbit(uint v) { m[v >> 6] |=
        1ULL << (v & 63); }
   inline void resetbit(uint v) { m[v >> 6]
        &= ~(1ULL << (v & 63)); }
```

```
inline uint getbit(uint v) const { return
     m[v >> 6] >> (v & 63) & 1; }
inline ull childs value(uint v) const {
     return m[left(v) >> 6]; }
inline int left_go(uint x, const uint c)
     const {
  const ull rem = x \& 63;
  uint bt = PREF * 64 + x:
  ull num = m[bt >> 6] & lb(rem + c);
  if(num) return (x ^ rem) | __lg(num);
  for(bt = parent(bt); bt > 62; bt =
       parent(bt)) {
    const ull rem = bt & 63;
    num = m[bt >> 6] & lb(rem);
    if(num) {
      bt = (bt ^ rem) | __lg(num);
      break;
  if(bt == 62) return -1;
  while(bt < PREF * 64) bt = left(bt) |</pre>
       __lg(m[bt - 62]);
  return bt - PREF * 64;
inline int right go(uint x, const uint c)
     const {
  const ull rem = x \& 63;
  uint bt = PREF * 64 + x;
  ull num = m[bt >> 6] \& \sim lb(rem + c);
  if(num) return (x ^ rem) |
       __builtin_ctzll(num);
```

for(bt = parent(bt); bt > 62; bt =

bt = (bt ^ rem) | __builtin_ctzll(

while(bt < PREF * 64) bt = left(bt) |</pre>

fast set() { assert(PREF != 228); setbit

bool empty() const {return getbit(63);}

(x >> 6)] >> (x & 63) & 1; }

if(!getbit(PREF * 64 + x)) return;

void insert(uint x) { for(uint v = PREF *

64 + x; !getbit(v); v = parent(v))

for(uint v = parent(PREF * 64 + x); v >

62 && !childs value(v); v = parent(v

void clear() { fill(m, m + SZ, 0); setbit

bool count(uint x) const { return m[PREF +

builtin ctzll(m[bt - 62]);

parent(bt)) {

num);

if(bt == 62) **return** -1;

return bt - PREF * 64;

setbit(v); }

resetbit(PREF * 64 + x);

)) resetbit(v);

void erase(uint x) {

if(num) {

break;

const ull rem = bt & 63; num = $m[bt >> 6] \& \sim lb(rem + 1);$

31

```
int find next(uint x) const { return
         right go(x, 0); } // >=
    int find prev(uint x) const { return
         left_go(x, 1); } // <=
72 };
         lazysegtree
1 template < class S,</pre>
           S (*e)(),
            S (*op)(S, S),
            class F,
           F (*id)(),
           S (*mapping)(F, S),
           F (*composition)(F, F)>
  struct lazy_segtree {
    int n, size, log;
    vector<S> d; vector<F> lz;
    void update(int k) { d[k] = op(d[k << 1],</pre>
         d[k << 1 | 1]); }
    void all apply(int k, F f) {
      d[k] = mapping(f, d[k]);
      if(k < size) lz[k] = composition(f, lz[k</pre>
    void push(int k) {
      all_apply(k << 1, lz[k]);
      all apply(k \langle\langle 1 \mid 1, 1z[k]\rangle\rangle;
      lz[k] = id();
    lazy segtree(int n) : lazy segtree(vector
         <S>( n, e())) {}
    lazy segtree(const vector<S>& v) : n(sz(v) 83
      log = __lg(2 * n - 1), size = 1 << log;
      d.resize(size * 2, e());
      lz.resize(size, id());
      REP(i, n) d[size + i] = v[i];
```

for(int i = size - 1; i; i--) update(i);

for(int i = log; i; --i) push(p >> i);

for(int i = 1; i <= log; ++i) update(p</pre>

for(int i = log; i; i--) push(p >> i);

if(1 & 1) sml = op(sml, d[1++]);

if(r & 1) smr = op(d[--r], smr);

if(((1 >> i) << i) != 1) push(1 >> i); 104

if(((r >> i) << i) != r) push(r >> i); 105

void set(int p, S x) {

>> i);

p += size;

S get(int p) {

p += size;

return d[p];

S prod(int 1, int r) {

if(1 == r) return e();

1 += size; r += size;

S sml = e(), smr = e();

while(1 < r) {</pre>

for(int i = log; i; i--) {

```
1 >>= 1;
    r >>= 1;
                                             111
  return op(sml, smr);
                                              112
                                              113
S all prod() const { return d[1]; }
                                              114
void apply(int p, F f) {
                                              115
  p += size:
                                              116
  for(int i = log; i; i--) push(p >> i);
                                              117
  d[p] = mapping(f, d[p]);
                                              118
  for(int i = 1; i <= log; i++) update(p</pre>
                                              119
       >> i);
                                              120
void apply(int 1, int r, F f) {
                                              121
  if(1 == r) return:
                                              122
  1 += size; r += size;
                                              123
  for(int i = log; i; i--) {
                                              124
    if(((1 >> i) << i) != 1) push(1 >> i); 125
    if(((r \rightarrow i) << i) != r) push((r - 1)
                                             126
         >> i);
                                              127 };
    int 12 = 1, r2 = r;
    while(1 < r) {
      if(l & 1) all_apply(l++, f);
      if(r & 1) all_apply(--r, f);
      1 >>= 1:
      r >>= 1:
   1 = 12;
   r = r2;
  for(int i = 1; i <= log; i++) {
   if(((1 >> i) << i) != 1) update(1 >> i
    if(((r >> i) << i) != r) update((r -
         1) \gg i);
template < class G> int max_right(int 1, G g
  assert(0 <= 1 && 1 <= n && g(e()));
  if(1 == n) return n;
 1 += size:
  for(int i = log; i; i--) push(l >> i);
  S sm = e();
    while(!(1 & 1)) 1 >>= 1;
    if(!g(op(sm, d[1]))) {
      while(1 < size) {</pre>
        push(1);
        1 <<= 1:
        if(g(op(sm, d[1]))) sm = op(sm, d[
             1++]);
      return 1 - size;
    sm = op(sm, d[1++]);
                                              2.7
  } while((1 & -1) != 1);
  return n:
template < class G> int min left(int r, G g)
                                              30
  assert(0 <= r && r <= n && g(e()));
                                              31
  if(r == 0) return 0:
```

r += size;

```
for(int i = log; i >= 1; i--) push((r -
                                                 33
          1) >> i);
     S sm = e():
     do {
                                                 35
                                                  36
       while(r > 1 && (r & 1)) r >>= 1;
       if(!g(op(d[r], sm))) {
          while(r < size) {</pre>
                                                 37
            push(r);
                                                  38
            r = r << 1 | 1;
            if(g(op(d[r], sm))) sm = op(d[r])
                 --1, sm);
                                                 41
          return r + 1 - size;
                                                 42
       sm = op(d[r], sm);
     } while((r & -r) != r);
     return 0:
 2.6 LCT
1 template < class S.</pre>
          S (*e)(),
          S (*op)(S, S),
          S (*reversal)(S),
           class F,
          F (*id)(),
          S (*mapping)(F, S),
          F (*composition)(F, F)>
 struct lazy lct {
   struct Node {
     S \text{ val} = e(), \text{ sum} = e();
     F lz = id();
     bool rev = false;
     int sz = 1:
                                                 60
     Node *1 = nullptr, *r = nullptr, *p =
                                                 61
          nullptr;
                                                 62
     Node() {}
                                                 63
     Node(const S& s) : val(s), sum(s) {}
     bool is root() const { return p ==
          nullptr || (p->l != this && p->r !=
          this); }
   };
   int n;
   vector<Node> a;
   lazy lct() : n(0) {}
   explicit lazy_lct(int _n) : lazy_lct(
                                                 71
        vector<S>( n, e())) {}
                                                  72
   explicit lazy lct(const vector<S>& v) : n(
        SZ(v)) { REP(i, n) a.eb(v[i]); }
   Node* access(int u) {
     Node* v = &a[u];
                                                 75
     Node* last = nullptr;
     for(Node* p = v; p != nullptr; p = p->p)
           splay(p), p->r = last, pull(last =
                                                 78
          p);
     splay(v);
     return last;
   void make root(int u) { access(u), a[u].
                                                 81
```

rev ^= 1, push(&a[u]); }

```
void link(int u, int v) { make root(v), a[
     vl.p = &a[u]; }
void cut(int u) {
  access(u);
  if(a[u].1 != nullptr) a[u].1->p =
       nullptr, a[u].1 = nullptr, pull(&a[u
void cut(int u, int v) { make root(u), cut
bool is connected(int u, int v) {
  if(u == v) return true;
  return access(u), access(v), a[u].p !=
       nullptr:
int get lca(int u, int v) { return access(
     u), access(v) - &a[0]; }
void set(int u, const S& s) { access(u), a
     [u].val = s, pull(&a[u]); }
S get(int u) { return access(u), a[u].val;
void apply(int u, int v, const F& f) {
     make_root(u), access(v), all_apply(&a[
     v], f), push(&a[v]); }
S prod(int u, int v) { return make root(u)
     , access(v), a[v].sum; }
void rotate(Node* v) {
  auto attach = [&](Node* p, bool side,
       Node* c) {
    (side ? p - > r : p - > 1) = c;
    pull(p);
    if(c != nullptr) c->p = p;
  Node *p = v->p, *g = p->p;
  bool rgt = (p->r == v);
  bool rt = p->is root();
  attach(p, rgt, (rgt ? v->l : v->r));
  attach(v, !rgt, p);
  if(!rt) attach(g, (g->r == p), v);
  else v \rightarrow p = g;
void splay(Node* v) {
  push(v);
  while(!v->is_root()) {
    auto p = v->p;
    auto g = p \rightarrow p;
    if(!p->is root()) push(g);
    push(p), push(v);
    if(!p->is root()) rotate((g->r == p)
         == (p->r == v) ? p : v);
    rotate(v);
void all apply(Node* v, F f) {
  v->val = mapping(f, v->val), v->sum =
       mapping(f, v->sum);
  v \rightarrow lz = composition(f, v \rightarrow lz);
void push(Node* v) {
  if(v->lz != id()) {
    if(v->l != nullptr) all_apply(v->l, v
    if(v->r != nullptr) all apply(v->r, v
         ->1z);
    v \rightarrow lz = id():
```

```
if(v->rev) {
                                                                         while(1 < n) 1 <<= 1, r <<= 1;
           swap(v->1, v->r);
                                                                         while(1 < r) {
           if(v \rightarrow l != nullptr) v \rightarrow l \rightarrow rev ^= 1:
                                                                           int c = (1 + r) / 2:
           if(v->r != nullptr) v->r->rev ^= 1;
                                                                           T xl = xs[1 - n], xr = xs[r - 1 - n],
           v->sum = reversal(v->sum);
                                                                                 xc = xs[c - n];
           v->rev = false:
                                                                           Line& f = fs[i]:
                                                                           if(f(x1) \leftarrow g(x1) \&\& f(xr) \leftarrow g(xr))
      void pull(Node* v) {
                                                                           if(f(xl) >= g(xl) \&\& f(xr) >= g(xr)) {
        v \rightarrow sz = 1;
                                                                                   f = g; return; }
        v \rightarrow sum = v \rightarrow val:
                                                                           if(f(xc) > g(xc)) swap(f, g);
        if(v->l != nullptr) {
                                                                           if(f(x1) > g(x1)) i = 2 * i, r = c;
           push(v->1);
                                                                           else i = 2 * i + 1, l = c;
           v \rightarrow sum = op(v \rightarrow 1 \rightarrow sum, v \rightarrow sum);
           v \rightarrow sz += v \rightarrow 1 \rightarrow sz;
                                                                     T get(T x) {
        if(v->r != nullptr) {
                                                                        int i = index(x);
           push(v->r);
                                                                        T res = INF;
                                                                        for(i += n; i; i >>= 1) res = min(res,
           v \rightarrow sum = op(v \rightarrow sum, v \rightarrow r \rightarrow sum);
           v \rightarrow sz += v \rightarrow r \rightarrow sz;
                                                                               fs[i](x));
103
                                                                         return res;
104
105 };
                                                                51 };
```

2.7 LiChao

```
i template < class T>
  struct LiChao {
    static constexpr T INF = numeric limits<T</pre>
         >::max();
    struct Line {
      T a, b:
      Line(T a, T b) : a(a), b(b) {}
      T operator()(T x) const { return a * x +
    int n;
    vector<Line> fs;
    vector<T> xs;
    LiChao(const vector<T>& xs ) {
      xs = sort_unique(xs_);
      n = SZ(xs);
      fs.assign(2 * n, Line(T(0), INF));
    int index(T x) const { return lower bound(
         ALL(xs), x) - xs.begin(); }
    void add_line(T a, T b) { update(a, b, 0,
         n); }
    // [xl, xr) ax+b
    void add segment(T xl, T xr, T a, T b) {
      int l = index(xl), r = index(xr);
      update(a, b, l, r);
23
    void update(T a, T b, int 1, int r) {
      Line g(a, b);
      for(1 += n, r += n; 1 < r; 1 >>= 1, r
           >>= 1){
        if(1 & 1) descend(g, 1++);
        if(r & 1) descend(g, --r);
29
    void descend(Line g, int i) {
      int l = i, r = i + 1;
```

2.8 rect-add-rect-sum

i template < class Int, class T>

```
struct RectangleAddRectangleSum {
  struct AQ { Int xl, xr, yl, yr; T val; };
  struct SQ { Int xl, xr, yl, yr; };
  vector<AQ> add qry;
  vector<SQ> sum qry;
 // A[x][y] += val for(x, y) in [xl, xr) *
  void add_rectangle(Int xl, Int xr, Int yl,
       Int yr, T val) { add_qry.pb({xl, xr,
      yl, yr, val}); }
  // Get sum of A[x][y] for (x, y) in [xl, xr]
      ) * [yl, yr)
  void add query(Int xl, Int xr, Int yl, Int
       yr) { sum_qry.pb({xl, xr, yl, yr}); }
  vector<T> solve() {
    vector<Int> vs:
    for(auto &a : add_qry) ys.pb(a.yl), ys.
        pb(a.yr);
    ys = sort_unique(ys);
    const int Y = SZ(ys);
    vector<tuple<Int, int, int>> ops;
    REP(q, SZ(sum_qry)) {
     ops.eb(sum_qry[q].xl, 0, q);
     ops.eb(sum_qry[q].xr, 1, q);
    REP(q, SZ(add qry)) {
     ops.eb(add_qry[q].xl, 2, q);
     ops.eb(add_qry[q].xr, 3, q);
    sort(ALL(ops));
    fenwick<T> b00(Y), b01(Y), b10(Y), b11(Y
    vector<T> ret(SZ(sum_qry));
    for(auto o : ops) {
      int qtype = get<1>(o), q = get<2>(o);
      if(qtype >= 2) {
```

```
b10.add(i, -yi * query.val);
  b11.add(i, query.val);
  b00.add(j, -x * yj * query.val);
  b01.add(j, x * query.val);
  b10.add(j, yj * query.val);
  b11.add(j, -query.val);
                                          1 struct segtree_beats {
} else {
                                             static constexpr 11 INF = numeric limits<</pre>
  const auto& query = sum qry[q];
  int i = lower_bound(ALL(ys), query.
       yl) - ys.begin();
  int j = lower_bound(ALL(ys), query.
```

yr) - ys.begin(); T x = get<0>(o);T yi = query.yl, yj = query.yr; if(qtype & 1) swap(i, j), swap(yi,

ret[q] += b00.get(i - 1) + b01.get(i 10

const auto& query = add gry[q];

T yi = query.yl, yj = query.yr;

b00.add(i, x * yi * guery.val);

b01.add(i, -x * query.val);

if(qtype & 1) swap(i, j), swap(yi,

vl) - vs.begin():

yr) - ys.begin();

T x = get<0>(o);

int i = lower_bound(ALL(ys), query.

int j = lower_bound(ALL(ys), query.

-1) * yi + b10.get(i - 1) * x+ b11.get(i - 1) * x * yi; ret[q] -= b00.get(j - 1) + b01.get(j 12 $-1) * yj + b10.get(j - 1) * x _{13}$ + b11.get(j - 1) * x * yj;

return ret;

54

2.9 rollback-dsu

```
1 struct RollbackDSU {
    int n: vi sz. tag:
    vector<tuple<int, int, int, int>> op;
    void init(int n) {
      n = _n;
      sz.assign(n, -1);
      tag.clear();
    int leader(int x) {
      while(sz[x] >= 0) x = sz[x];
                                                 27
      return x;
11
12
    bool merge(int x, int y) {
      x = leader(x), y = leader(y);
      if(x == y) return false;
      if(-sz[x] < -sz[y]) swap(x, y);
      op.eb(x, sz[x], y, sz[y]);
      sz[x] += sz[y]; sz[y] = x;
      return true;
19
    int size(int x) { return -sz[leader(x);] } 38
    void add_tag() { tag.pb(sz(op)); }
```

```
void rollback() {
      int z = tag.back(); tag.ppb();
      while(sz(op) > z) {
        auto [x, sx, y, sy] = op.back(); op.
             ppb();
        sz[x] = sx;
        sz[y] = sy;
31 };
```

2.10 segtree-beats

24

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32

33

```
ll>::max() / 2.1;
struct alignas(32) Node {
  11 \text{ sum} = 0, g1 = 0, 11 = 0;
  11 g2 = -INF, gc = 1, 12 = INF, 1c = 1,
       add = 0:
11 n, log;
vector<Node> v;
segtree beats() {}
segtree_beats(int _n) : segtree_beats(
     vector<11>( n)) {}
segtree_beats(const vector<11>& vc) {
  n = 1, log = 0;
  while(n < SZ(vc)) n <<= 1, log++;
  v.resize(2 * n);
  REP(i, SZ(vc)) v[i + n].sum = v[i + n].
       g1 = v[i + n].l1 = vc[i];
  for(ll i = n - 1; i; --i) update(i);
void range_chmin(int l, int r, ll x) {
     inner_apply<1>(l, r, x); }
void range chmax(int 1, int r, 11 x) {
     inner_apply<2>(1, r, x); }
void range_add(int 1, int r, 11 x) {
     inner apply \langle 3 \rangle (1, r, x); \}
void range_update(int 1, int r, 11 x) {
     inner apply<4>(1, r, x);}
11 range min(int 1, int r) { return
     inner_fold<1>(l, r); }
11 range_max(int 1, int r) { return
     inner_fold<2>(1, r); }
11 range_sum(int 1, int r) { return
     inner fold<3>(1, r); }
void update(int k) {
  Node& p = v[k]:
  Node& 1 = v[k * 2];
  Node& r = v[k * 2 + 1];
  p.sum = 1.sum + r.sum;
  if(1.g1 == r.g1) {
   p.g1 = 1.g1;
    p.g2 = max(1.g2, r.g2);
    p.gc = 1.gc + r.gc;
  } else {
    bool \dot{f} = 1.g1 > r.g1;
    p.g1 = f ? l.g1 : r.g1;
    p.gc = f ? 1.gc : r.gc;
    p.g2 = max(f ? r.g1 : 1.g1, f ? 1.g2 :
```

```
if(1.11 == r.11) {
    p.11 = 1.11:
    p.12 = min(1.12, r.12);
    p.lc = 1.lc + r.lc;
  } else {
    bool f = 1.11 < r.11;</pre>
    p.l1 = f ? l.l1 : r.l1:
    p.lc = f ? 1.lc : r.lc;
    p.12 = min(f ? r.11 : 1.11, f ? 1.12 : 106
          r.12);
void push add(int k, ll x) {
  Node& p = v[k];
 p.sum += x << (log + builtin clz(k) -
                                             111
 p.g1 += x, p.l1 += x;
  if(p.g2 != -INF) p.g2 += x;
 if(p.12 != INF) p.12 += x;
                                             113
 p.add += x;
void push_min(int k, ll x) {
  Node& p = v[k]:
 p.sum += (x - p.g1) * p.gc;
 if(p.11 == p.g1) p.11 = x;
 if(p.12 == p.g1) p.12 = x;
 p.g1 = x;
                                             121
void push max(int k, ll x) {
                                             122
  Node& p = v[k];
                                             123
 p.sum += (x - p.11) * p.1c;
                                             124
  if(p.g1 == p.11) p.g1 = x;
                                             125
 if(p.g2 == p.l1) p.g2 = x;
 p.11 = x;
                                             127
                                             128
void push(int k) {
                                             129
  Node& p = v[k];
                                             130
 if(p.add != 0) {
                                             131
    push_add(k * 2, p.add);
    push add(k * 2 + 1, p.add);
    p.add = 0;
  if(p.g1 < v[k * 2].g1) push min(k * 2, p 134
  if(p.11 > v[k * 2].11) push max(k * 2, p 136
       .11);
  if(p.g1 < v[k * 2 + 1].g1) push_min(k *</pre>
       2 + 1, p.g1);
  if(p.11 > v[k * 2 + 1].11) push_max(k *
       2 + 1, p.11);
void subtree chmin(int k, ll x) {
  if(v[k].g1 <= x) return;</pre>
 if(v[k].g2 < x) {
    push min(k, x);
    return;
  push(k):
  subtree_chmin(k * 2, x), subtree_chmin(k 151
        * 2 + 1, x);
  update(k);
                                             154
void subtree_chmax(int k, ll x) {
```

```
if(x <= v[k].l1) return;</pre>
  if(x < v[k].12) {
    push max(k, x);
    return;
                                            159
                                             160
  push(k);
  subtree chmax(k * 2, x), subtree chmax(k 162
        * 2 + 1, x);
  update(k);
                                             164
                                             165
template<int cmd>
                                             166
inline void apply(int k, ll x) {
                                             167
  if constexpr(cmd == 1) subtree_chmin(k,
  if constexpr(cmd == 2) subtree_chmax(k,
       x);
  if constexpr(cmd == 3) push add(k, x);
  if constexpr(cmd == 4) subtree_chmin(k,
       x), subtree chmax(k, x);
template<int cmd>
                                              i template < class S, S (*e)(), S (*op)(S, S)>
void inner_apply(int l, int r, ll x) {
  if(1 == r) return;
 1 += n, r += n;
  for(int i = log; i >= 1; i--) {
    if(((1 >> i) << i) != 1) push(1 >> i);
    if(((r >> i) << i) != r) push((r - 1)
         >> i):
    int 12 = 1, r2 = r;
    while (1 < r) {
      if(1 & 1) _apply<cmd>(1++, x);
      if(r & 1) _apply<cmd>(--r, x);
                                             12
     1 >>= 1, r >>= 1;
   1 = 12, r = r2;
  for(int i = 1; i <= log; i++) {</pre>
   if(((1 >> i) << i) != 1) update(1 >> i
    if(((r >> i) << i) != r) update((r -</pre>
        1) \gg i);
template<int cmd>
inline 11 e() {
 if constexpr(cmd == 1) return INF;
  if constexpr(cmd == 2) return -INF;
  return 0;
                                             28
template<int cmd>
inline void op(11& a, const Node& b) {
  if constexpr(cmd == 1) a = min(a, b.l1);
                                             32
  if constexpr(cmd == 2) a = max(a, b.g1); 33
  if constexpr(cmd == 3) a += b.sum;
template<int cmd>
11 inner fold(int 1, int r) {
  if(1 == r) return e<cmd>();
 1 += n, r += n;
  for(int i = log; i >= 1; i--) {
    if(((1 >> i) << i) != 1) push(1 >> i); 40
    if(((r >> i) << i) != r) push((r - 1)
```

```
11 1x = e < cmd > (), rx = e < cmd > ();
       while (1 < r) {
         if(1 & 1) op<cmd>(lx, v[1++]);
         if(r & 1) op<cmd>(rx, v[--r]);
         1 >>= 1, r >>= 1:
       if constexpr(cmd == 1) lx = min(lx, rx);
       if constexpr(cmd == 2) lx = max(lx, rx);
       if constexpr(cmd == 3) lx += rx;
       return lx:
168 };
```

2.11 segtree

```
2 struct segtree {
   int n, size, log;
   vector<S> st;
   void update(int v) { st[v] = op(st[v <</pre>
        1], st[v << 1 | 1]); }
   segtree(int _n) : segtree(vector<S>(_n, e
        ())) {}
   segtree(const vector<S>& a): n(sz(a)) {
     log = __lg(2 * n - 1), size = 1 << log;
     st.resize(size << 1, e());
     REP(i, n) st[size + i] = a[i];
     for(int i = size - 1; i; i--) update(i);
   void set(int p, S val) {
     st[p += size] = val;
     for(int i = 1; i <= log; ++i) update(p</pre>
          >> i);
   S get(int p) const {
     return st[p + size];
   S prod(int 1, int r) const {
     assert(0 <= 1 && 1 <= r && r <= n);
     S sml = e(), smr = e();
     1 += size, r += size;
     while(1 < r) {</pre>
       if(1 \& 1) sml = op(sml, st[1++]);
       if(r \& 1) smr = op(st[--r], smr);
       1 >>= 1;
       r >>= 1:
     return op(sml, smr);
   S all_prod() const { return st[1]; }
    template < class F> int max right(int 1, F f
        ) const {
     assert(0 <= 1 && 1 <= n && f(e()));
     if(1 == n) return n;
     1 += size;
     S sm = e();
     do {
       while(~1 & 1) 1 >>= 1;
       if(!f(op(sm, st[1]))) {
          while(1 < size) {</pre>
           1 <<= 1;
```

```
st[1++]);
           return 1 - size;
         sm = op(sm, st[1++]);
       } while((1 & -1) != 1);
       return n:
     template < class F > int min_left(int r, F f)
       assert(0 <= r && r <= n && f(e()));
      if(r == 0) return 0;
      r += size;
      S sm = e();
       do {
         while(r > 1 && (r & 1)) r >>= 1;
         if(!f(op(st[r], sm))) {
           while(r < size) {</pre>
             r = r \leftrightarrow 1 \mid 1;
             if(f(op(st[r], sm))) sm = op(st[r
                  --], sm);
64
           return r + 1 - size;
         sm = op(st[r], sm);
      } while((r & -r) != r);
       return 0;
68
69
70 };
```

if(f(op(sm, st[1]))) sm = op(sm,

2.12 sparse-table

```
i template < class T, T (*op)(T, T)>
2 struct sparse table {
    int n;
    vector<vector<T>> b;
     sparse table(const vector<T>& a) : n(sz(a)
       int lg = __lg(n) + 1;
       b.resize(lg); b[0] = a;
       for(int j = 1; j < lg; ++j) {</pre>
         b[j].resize(n - (1 << j) + 1);
         REP(i, n - (1 << j) + 1) b[j][i] = op(
              b[j - 1][i], b[j - 1][i + (1 << (j - 1)[i])
                - 1))]);
11
12
13
    T prod(int from, int to) {
      int lg = __lg(to - from + 1);
return op(b[lg][from], b[lg][to - (1 <</pre>
14
            lg) + 1]);
16
17 };
```

2.13 static-range-inversion

```
1 struct static range inversion {
2 int sz;
```

```
vi a, L, R;
vector<ll> ans;
static range inversion(vi a) : a( a) {
  _a = sort_unique(_a);
  REP(i, SZ(a)) a[i] = lower_bound(ALL(_a)
       , a[i]) - a.begin();
  sz = SZ(_a);
void add query(int 1, int r) { L.push back
     (1), R.push_back(r); }
vector<ll> solve() {
  const int q = SZ(L);
  const int B = max(1.0, SZ(a) / sqrt(q)); 15
  vi ord(q);
  iota(ALL(ord), 0);
  sort(ALL(ord), [&](int i, int j) {
    if(L[i] / B == L[j] / B) {
      return L[i] / B & 1 ? R[i] > R[j] :
           R[i] < R[j];
    return L[i] < L[j];</pre>
  ans.resize(q);
  fenwick<ll> fenw(sz + 1);
  11 \text{ cnt} = 0:
  auto AL = [&](int i) {
    cnt += fenw.sum(0, a[i] - 1);
    fenw.add(a[i], +1);
  auto AR = [&](int i) {
    cnt += fenw.sum(a[i] + 1, sz);
    fenw.add(a[i], +1);
  auto DL = [&](int i) {
    cnt -= fenw.sum(0, a[i] - 1);
    fenw.add(a[i], -1);
  auto DR = [&](int i) {
    cnt -= fenw.sum(a[i] + 1, sz);
    fenw.add(a[i], -1);
  int 1 = 0, r = 0;
  for(int i = 0; i < q; i++) {</pre>
    int id = ord[i], ql = L[id], qr = R[id
    while(1 > q1) AL(--1);
    while(r < qr) AR(r++);</pre>
    while(1 < q1) DL(1++);
    while(r > qr) DR(--r);
    ans[id] = cnt;
  return ans;
```

2.14 static-range-lis

```
#define MEM(a, x, n) memset(a, x, sizeof(int
        ) * n)
using I = int*;
struct static_range_lis {
   int n, ps = 0;
   I invp, res_monge, pool;
```

```
vector<vector<pii>>> gry;
static range lis(vi a) : n(SZ(a)), qry(n +
     1) {
  // a must be permutation of [0, n)
  pool = (I) malloc(sizeof(int) * n * 100) 60
  invp = A(n), res monge = A(n);
  REP(i, n) invp[a[i]] = i;
inline I A(int x) { return pool + (ps += x
    ) - x; }
void add_query(int 1, int r) { qry[1].pb({ 64
    r, SZ(ans)}), ans.pb(r - 1); }
void unit_monge_mult(I a, I b, I r, int n) 65
  if(n == 2){
    if(!a[0] && !b[0]) r[0] = 0, r[1] = 1; 68
    else r[0] = 1, r[1] = 0;
    return;
  if(n == 1) return r[0] = 0, void();
  int lps = ps, d = n / 2;
 I a1 = A(d), a2 = A(n - d), b1 = A(d),
      b2 = A(n - d);
  I mpa1 = A(d), mpa2 = A(n - d), mpb1 = A
       (d), mpb2 = A(n - d);
  int p[2] = {};
                                             76
  REP(i, n) {
   if(a[i] < d) a1[p[0]] = a[i], mpa1[p
         [0]++] = i;
    else a2[p[1]] = a[i] - d, mpa2[p[1]++]
         = i:
  p[0] = p[1] = 0;
  REP(i, n) {
    if(b[i] < d) b1[p[0]] = b[i], mpb1[p]
         [0]++] = i;
    else b2[p[1]] = b[i] - d, mpb2[p[1]++]
         = i;
  I c1 = A(d), c2 = A(n - d);
  unit_monge_mult(a1, b1, c1, d),
      unit monge mult(a2, b2, c2, n - d);
  I cpx = A(n), cpy = A(n), cqx = A(n),
      cqy = A(n);
  REP(i, d) cpx[mpa1[i]] = mpb1[c1[i]],
       cpv[mpa1[i]]=0;
  REP(i, n - d) cpx[mpa2[i]] = mpb2[c2[i]]
       ]], cpy[mpa2[i]]=1;
  REP(i, n) r[i] = cpx[i];
  REP(i, n) cqx[cpx[i]] = i, cqy[cpx[i]] =
  int hi = n, lo = n, his = 0, los = 0;
  REP(i, n)
    if(cqy[i] ^ (cqx[i] >= hi)) his--;
                                             98
    while(hi > 0 && his < 0) {</pre>
                                             99
                                             100
      if(cpy[hi] ^ (cpx[hi] > i)) his++;
    while(lo > 0 && los <= 0) {
                                            103
     lo--;
      if(cpy[lo] ^ (cpx[lo] >= i)) los++;
                                            105
    if(los > 0 && hi == lo) r[lo] = i:
                                            106
    if(cqy[i] ^ (cqx[i] >= lo)) los--;
                                            107
```

```
ps = lps;
void subunit_monge_mult(I a, I b, I c, int 110
  int lps = ps:
                                            112
  I za = A(n), zb = A(n), res = A(n), vis 113
      = A(n), mpa = A(n), mpb = A(n), rb = 114
  MEM(vis, 0, n), MEM(mpa, -1, n), MEM(mpb
       , -1, n), MEM(rb, -1, n);
  int ca = n;
  IREP(i, n) if(a[i] != -1) vis[a[i]] = 1,
        za[--ca] = a[i], mpa[ca] = i;
  IREP(i, n) if(!vis[i]) za[--ca] = i;
  MEM(vis, -1, n);
  REP(i, n) if(b[i] != -1) vis[b[i]] = i;
  REP(i, n) if(vis[i] != -1) mpb[ca] = i,
      rb[vis[i]] = ca++;
  REP(i, n) if(rb[i] == -1) rb[i] = ca++;
  REP(i, n) zb[rb[i]] = i;
  unit_monge_mult(za, zb, res, n);
  MEM(c, -1, n);
  REP(i, n) if(mpa[i] != -1 && mpb[res[i]]
       != -1) c[mpa[i]] = mpb[res[i]];
  ps = lps;
                                             12
                                             13
void solve(I p, I ret, int n) {
  if(n == 1) return ret[0] = -1, void();
  int lps = ps, d = n / 2;
                                             16
  I pl = A(d), pr = A(n - d);
                                             17
  REP(i, d) pl[i] = p[i];
                                             18
  REP(i, n - d) pr[i] = p[i + d];
  I vis = A(n); MEM(vis, -1, n);
  REP(i, d) vis[pl[i]] = i;
  I tl = A(d), tr = A(n - d), mpl = A(d),
       mpr = A(n - d);
  int ca = 0;
  REP(i, n) if(vis[i] != -1) mpl[ca] = i,
       tl[vis[i]] = ca++;
  ca = 0; MEM(vis, -1, n);
  REP(i, n - d) vis[pr[i]] = i;
  REP(i, n) if(vis[i] != -1) mpr[ca] = i,
       tr[vis[i]] = ca++;
  I vl = A(d), vr = A(n - d);
  solve(tl, vl, d), solve(tr, vr, n - d);
  I sl = A(n), sr = A(n);
  iota(sl, sl + n, 0); iota(sr, sr + n, 0)
  REP(i, d) sl[mpl[i]] = (vl[i] == -1 ? -1
        : mpl[vl[i]]);
  REP(i, n - d) sr[mpr[i]] = (vr[i] == -1)
      ? -1 : mpr[vr[i]]);
  subunit monge mult(sl, sr, ret, n);
  ps = lps;
vi solve() {
  solve(invp, res_monge, n);
  vi fenw(n + 1):
  IREP(i, n) {
    if(res monge[i] != -1) {
      for(int p = res monge[i] + 1; p <= n</pre>
                                            47
           ; p += p & -p) fenw[p]++;
    for(auto& z : qry[i]){
```

```
auto [id, c] = z;
           for(int p = id; p; p -= p & -p) ans[
                 cl -= fenw[pl;
       free(pool);
       return ans;
  2.15 treap
 1 struct Node {
    bool rev = false;
    int sz = 1, pri = rng();
    Node *1 = NULL, *r = NULL, *p = NULL;
    // TODO
  void pull(Node*& v) {
    v\rightarrow sz = 1 + size(v\rightarrow l) + size(v\rightarrow r);
    // TODO
void push(Node*& v) {
    if(v->rev) {
       swap(v->1, v->r);
       if(v\rightarrow 1) v\rightarrow 1\rightarrow rev ^= 1;
       if(v->r) v->r->rev ^= 1;
       v->rev = false;
Node* merge(Node* a, Node* b) {
    if(!a | | !b) return (a ? a : b);
    push(a), push(b);
    if(a->pri > b->pri) {
       a->r = merge(a->r, b);
       pull(a); return a;
     } else {
       b\rightarrow 1 = merge(a, b\rightarrow 1);
       pull(b); return b;
30 pair<Node*, Node*> split(Node* v, int k) {
    if(!v) return {NULL, NULL};
     push(v);
     if(size(v->1) >= k) {
       auto p = split(v->1, k);
       if(p.first) p.first->p = NULL;
       v \rightarrow 1 = p.second:
       pull(v); return {p.first, v};
     } else {
       auto p = split(v \rightarrow r, k - size(v \rightarrow l) - 1)
       if(p.second) p.second->p = NULL;
       v \rightarrow r = p.first;
       pull(v); return {v, p.second};
45 int get_position(Node* v) { // 0-indexed
    int k = (v->1 != NULL ? v->1->sz : 0);
```

while(v->p != NULL) {

 $if(v == v \rightarrow p \rightarrow r)$ {

if($v \rightarrow p \rightarrow l != NULL$) k += $v \rightarrow p \rightarrow l \rightarrow sz$;

```
v = v - p;
return k;
```

union-of-rectangles

```
2 // 1 10 1 10
3 // 0 2 0 2
4 // ans = 84
  vector<int> vx, vy;
    struct q { int piv, s, e, x; };
   struct tree {
        vector<int> seg, tag;
        tree(int _n) : seg(_n * 16), tag(_n * 16)
         void add(int ql, int qr, int x, int v, int
                         1, int r) {
              if(qr <= 1 | | r <= q1) return;
              if(ql <= 1 && r <= qr) {
                    tag[v] += x;
                    if(tag[v] == 0) {
                          if(1 != r) seg[v] = seg[2 * v] + seg
                                       [2 * v + 1];
                           else seg[v] = 0;
                    } else seg[v] = vx[r] - vx[1];
              } else {
                    int mid = (1 + r) / 2;
                    add(q1, qr, x, 2 * v, 1, mid);
                    add(ql, qr, x, 2 * v + 1, mid, r);
                    if(tag[v] == 0 \&\& 1 != r) seg[v] = seg
                                  [2 * v] + seg[2 * v + 1];
        int q() { return seg[1]; }
    };
   int main() {
        int n; cin >> n;
         vector<int> x1(n), x2(n), y_(n), y2(n);
         for (int i = 0; i < n; i++) {
              cin >> x1[i] >> x2[i] >> y_[i] >> y2[i];
                               // L R D U
              vx.pb(x1[i]), vx.pb(x2[i]);
              vy.pb(y_[i]), vy.pb(y2[i]);
        vx = sort unique(vx);
         vy = sort unique(vy);
         vector\langle q \rangle a(2 * n);
         REP(i, n) {
              x1[i] = lower_bound(ALL(vx), x1[i]) - vx
                             .begin();
              x2[i] = lower_bound(ALL(vx), x2[i]) - vx 26
                             .begin();
              y_{i} = lower_{i} = lower_{i
                             .begin();
              y2[i] = lower_bound(ALL(vy), y2[i]) - vy 29
                             .begin();
              a[2 * i] = {y_[i], x1[i], x2[i], +1};
              a[2 * i + 1] = \{y2[i], x1[i], x2[i],
                             -1};
```

```
sort(ALL(a), [](q a, q b) { return a.piv < 33
      b.piv; });
tree seg(n):
11 \text{ ans} = 0;
REP(i, 2 * n)  {
  int i = i:
  while(j < 2 * n && a[i].piv == a[j].piv)</pre>
    seg.add(a[j].s, a[j].e, a[j].x, 1, 0,
        vx.size());
    j++;
  if(a[i].piv + 1 != SZ(vy)) ans += 1LL *
       seg.q() * (vy[a[i].piv + 1] - vy[a[i 44
       ].piv]);
  i = j - 1;
cout << ans << "\n";
```

wavelet-tree

struct wavelet_tree {

template < class T>

int n, log;

```
vector<T> vals;
vi sums;
vector<ull> bits;
inline void set bit(int i, ull v) { bits[i
      >> 6] |= (v << (i & 63)); }
inline int get_sum(int i) const { return
     sums[i >> 6] + __builtin_popcountll(
bits[i >> 6] & ((1ULL << (i & 63)) -</pre>
wavelet_tree(const vector<T>& _v) : n(SZ(
      v)) {
  vals = sort unique( v);
  log = lg(2 * vals.size() - 1);
  bits.resize((log * n + 64) >> 6, 0ULL);
  sums.resize(SZ(bits), 0);
  vi v(SZ(_v)), cnt(SZ(vals) + 1);
  REP(i, SZ(v)) {
    v[i] = lower bound(ALL(vals), _v[i]) -
           vals.begin();
    cnt[v[i] + 1] += 1;
  partial sum(ALL(cnt) - 1, cnt.begin());
  REP(j, log) {
    for(int i : v) {
      int tmp = i >> (log - 1 - j);
      int pos = (tmp >> 1) << (log - j);</pre>
      set bit(j * n + cnt[pos], tmp & 1);
      cnt[pos]++;
    for(int i : v) cnt[(i >> (log - j)) <<</pre>
           (log - j)]--;
  for(int i = 1; i < (int) sums.size(); i</pre>
       ++) sums[i] = sums[i - 1] +
       __builtin_popcountll(bits[i - 1]);
T get_kth(int a, int b, int k) {
```

```
for(int j = 0, ia = 0, ib = n, res = 0;; 26
            j++) {
        if(j == log) return vals[res];
                                                 28
        int cnt_ia = get_sum(n * j + ia);
                                                 29
        int cnt a = get sum(n * j + a);
                                                 30
        int cnt b = get sum(n * j + b);
                                                 31
        int cnt ib = get_sum(n * j + ib);
                                                 32
        int ab zeros = (b - a) - (cnt b -
                                                 33
             cnt a);
                                                 34
        if(ab zeros > k) {
          res <<= 1:
          ib -= cnt ib - cnt ia;
          a -= cnt_a - cnt_ia;
          b -= cnt b - cnt ia;
        } else {
          res = (res << 1) | 1;
          k -= ab zeros:
          ia += (ib - ia) - (cnt_ib - cnt_ia);
          a += (ib - a) - (cnt ib - cnt a);
          b += (ib - b) - (cnt_ib - cnt_b);
54 };
```

Flow-Matching

53

24

3.1 bipartite-matching

```
1 struct bipartite_matching {
   int n, m; // 二分圖左右人數 (0 ~ n-1), (0
        \sim m-1)
    vector<vi> g;
    vi lhs, rhs, dist; // i 與 Lhs[i] 配對 (
         Lhs[i] == -1 代表沒有配對)
    bipartite_matching(int _n, int _m) : n(_n)
         , m(_m), g(_n), lhs(_n, -1), rhs(_m,
         -1), dist(_n) {}
    void add edge(int u, int v) { g[u].pb(v);
    void bfs() {
      aueue<int> a:
      REP(i, n) {
        if(lhs[i] == -1) {
          q.push(i);
          dist[i] = 0;
        } else {
          dist[i] = -1;
      while(!q.empty()) {
        int u = q.front(); q.pop();
        for(auto v : g[u]) {
          if(rhs[v] != -1 && dist[rhs[v]] ==
               -1) {
            dist[rhs[v]] = dist[u] + 1;
            q.push(rhs[v]);
22
23
```

```
rhs[lhs[u] = v] = u;
      return true:
  for(auto v : g[u]) {
   if(dist[rhs[v]] == dist[u] + 1 && dfs(
        rhs[v])) {
      rhs[lhs[u] = v] = u;
      return true;
 return false;
int solve() {
 int ans = 0:
  while(true) {
   bfs();
   int aug = 0;
    REP(i, n) if(lhs[i] == -1) aug += dfs(
   if(!aug) break;
    ans += aug;
  return ans;
```

bool dfs(int u) {

for(auto v : g[u]) {

if(rhs[v] == -1) {

3.2 Dinic-LowerBound

```
i template < class T>
2 struct DinicLowerBound {
    using Maxflow = Dinic<T>;
    int n;
    Maxflow d;
    vector<T> in;
    DinicLowerBound(int _n) : n(_n), d(_n + 2)
          , in(_n) {}
    int add edge(int from, int to, T low, T
          high) {
       assert(0 <= low && low <= high);
       in[from] -= low, in[to] += low;
11
       return d.add_edge(from, to, high - low);
12
    T flow(int s, int t) {
13
14
      T sum = 0;
       REP(i, n) {
        if(in[i] > 0) {
           d.add_edge(n, i, in[i]);
18
           sum += in[i];
19
20
         if(in[i] < 0) d.add_edge(i, n + 1, -in</pre>
              [i]);
2.1
       d.add_edge(t, s, numeric_limits<T>::max
22
       if(d.flow(n, n + 1) < sum) return -1;</pre>
23
24
       return d.flow(s, t);
25
26 };
```

3.3 Dinic

```
1 template < class T>
2 class Dinic {
 public:
   struct Edge {
     int from, to;
     T cap;
     Edge(int x, int y, T z) : from(x), to(y)
          , cap(z) {}
   constexpr T INF = 1e9;
   int n;
   vector<Edge> edges;
   vector<vi> g;
   vi cur, h; // h : level graph
   Dinic(int _n) : n(_n), g(_n) {}
   void add edge(int u, int v, T c) {
     g[u].pb(sz(edges));
     edges.eb(u, v, c);
     g[v].pb(sz(edges));
     edges.eb(v, u, 0);
   bool bfs(int s, int t) {
     h.assign(n, -1);
     queue<int> q;
     h[s] = 0;
     q.push(s);
     while(!q.empty()) {
       int u = q.front(); q.pop();
       for(int i : g[u]) {
         const auto& e = edges[i];
         int v = e.to;
         if(e.cap > 0 \&\& h[v] == -1) {
           h[v] = h[u] + 1;
           if(v == t) return true;
           q.push(v);
     return false;
   T dfs(int u, int t, T f) {
     if(u == t) return f;
     Tr = f;
     for(int& i = cur[u]; i < sz(g[u]); ++i)</pre>
       int j = g[u][i];
        const auto& e = edges[j];
       int v = e.to:
       T c = e.cap;
       if(c > 0 \&\& h[v] == h[u] + 1) {
         T = dfs(v, t, min(r, c));
         edges[j].cap -= a;
         edges[j ^ 1].cap += a;
         if((r -= a) == 0) return f;
     return f - r;
   T flow(int s, int t, T f = INF) {
     T ans = 0;
     while(f > 0 && bfs(s, t)) {
       cur.assign(n, 0);
       T cur = dfs(s, t, f);
```

```
ans += cur;
  f -= cur;
return ans;
```

3.4 Flow 建模

- · Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \rightarrow v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \rightarrow s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from Sto T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \rightarrow s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise con- $\operatorname{nect} y \to x \text{ with } (\cos t, \operatorname{cap}) = (-c, 1)$
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v) > 0, connect 11 $S \to v$ with (cost, cap) = (0, d(v))
 - 5. For each vertex v with d(v) < 0, connect 13 $v \to T$ with (cost, cap) = (0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the 15 flow C + K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're check- 19 ing answer T
 - 2. Construct a max flow model, let K be the sum 21of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K 23

```
4. For each edge (u, v, w) in G, connect u \to v 24
   and v \to u with capacity w
5. For v \in G, connect it with sink v \to t with ca-26
   pacity K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v) 27
6. T is a valid answer if the maximum flow f < 28
```

· Minimum weight edge cover

K|V|

- 1. For each $v \in V$ create a copy v', and connect 32 $u' \to v'$ with weight w(u, v).
- 2. Connect $v \rightarrow v'$ with weight $2\mu(v)$, where 34 $\mu(v)$ is the cost of the cheapest edge incident to 35
- 3. Find the minimum weight perfect matching on 37
- · Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - 2. Create edge (u, v) with capacity w with w being $_{42}$ the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit $\frac{1}{44}$ of a subset of projects.
- 0/1 quadratic programming

```
\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'}) \\ \mathbf{3.6}
```

general-weighted-max-matching

for(auto i : o) if(mate[i] == -1) ans

return false:

void add_edge(int a, int b) {

shuffle(all(o), rng);

vis.assign(n, false);

+= dfs(i):

es.eb(b, g[a]);

f(a, b); f(b, a);

iota(all(o), 0);

int solve() {

int ans = 0;

return ans;

1 // 1-based 00

16

REP(it, 100) {

vi o(n);

g[a] = sz(es) - 1;

auto f = [&](int a, int b) {

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity c_x and create edge (s, y) with capacity c_y .
- 2. Create edge (x, y) with capacity c_{xy} .
- 3. Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

general-matching

```
1 struct GeneralMaxMatch {
                                                12
   int n;
   vector<pii> es;
   vi g, vis, mate; // i 與 mate[i] 配對 (
        mate[i] == -1 代表沒有匹配)
   GeneralMaxMatch(int n) : n(n), g(n, -1),
                                                15
        mate(n, -1) {}
   bool dfs(int u) {
                                                17
     if(vis[u]) return false;
     vis[u] = true;
     for(int ei = g[u]; ei != -1; ) {
       auto [x, y] = es[ei]; ei = y;
       if(mate[x] == -1) {
         mate[mate[u] = x] = u;
                                                20
          return true;
                                                21
     for(int ei = g[u]; ei != -1; ) {
                                                22
       auto [x, y] = es[ei]; ei = y;
                                                23
       int nu = mate[x];
                                                24
       mate[mate[u] = x] = u;
                                                25
       mate[nu] = -1;
       if(dfs(nu)) return true;
                                                26
       mate[mate[nu] = x] = nu;
       mate[u] = -1;
```

```
struct WeightGraph {
  static const int inf = INT MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge() {}
    edge(int u, int v, int w): u(u), v(v), w
  int n, n x;
  edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[
       maxn * 2], pa[maxn * 2];
  int flo from[maxn * 2][maxn + 1], S[maxn *
        2], vis[maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) { return lab[e.
       u] + lab[e.v] - g[e.u][e.v].w * 2; }
  void update_slack(int u, int x) { if(!
       slack[x] || e_delta(g[u][x]) < e_delta</pre>
       (g[slack[x]][x])) slack[x] = u; }
  void set slack(int x) {
    slack[x] = 0;
    REP(u, n) if(g[u + 1][x].w > 0 \&\& st[u +
          1] != x && S[st[u + 1]] == 0)
         update_slack(u + 1, x);
  void q_push(int x) {
    if(x <= n) q.push(x);
    else REP(i, SZ(flo[x])) q_push(flo[x][i
  void set st(int x, int b) {
    st[x] = b;
```

```
if(x > n) REP(i, SZ(flo[x])) set_st(flo[
       x|[i], b);
int get_pr(int b, int xr) {
  int pr = find(ALL(flo[b]), xr) - flo[b].
       begin();
  if(pr % 2 == 1) {
    reverse(1 + ALL(flo[b]));
    return SZ(flo[b]) - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if(u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u
  for(int i = 0; i < pr; ++i) set_match(</pre>
       flo[u][i], flo[u][i ^ 1]);
  set match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() +
       pr, flo[u].end());
void augment(int u, int v) {
  while(true) {
                                             102
    int xnv = st[match[u]];
    set match(u, v);
    if(!xnv) return;
    set match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int get_lca(int u, int v) {
  static int t = 0;
  for(++t; u || v; swap(u, v)) {
                                             112
    if(u == 0) continue;
    if(vis[u] == t) return u;
    vis[u] = t;
    if(u = st[match[u]]) u = st[pa[u]];
                                             116
  return 0;
                                             117
void add blossom(int u, int lca, int v) {
  int b = n + 1;
  while(b <= n_x && st[b]) ++b;</pre>
  if(b > n x) n x++;
  lab[b] = S[b] = 0;
  match[b] = match[lca];
  flo[b].clear(); flo[b].pb(lca);
  for(int x = u, y; x != lca; x = st[pa[y
       ]]) flo[b].pb(x), flo[b].pb(y = st[ 125]
       match[x]]), q_push(y);
  reverse(1 + ALL(flo[b]));
  for(int x = v, y; x != lca; x = st[pa[y
       ]]) flo[b].pb(x), flo[b].pb(y = st[ 129]
       match[x]]), q push(y);
  set st(b, b);
  REP(x, n_x) g[b][x + 1].w = g[x + 1][b]. 131
  REP(x, n) flo_from[b][x + 1] = 0;
                                             132
  REP(i, SZ(flo[b])) {
    int xs = flo[b][i]:
    REP(x, n_x) if(g[b][x + 1].w == 0 | |
         e delta(g[xs][x + 1]) < e delta(g[ 136
         b][x + 1])) g[b][x + 1] = g[xs][x
```

```
+ 1], g[x + 1][b] = g[x + 1][xs];
    REP(x, n) if(flo from[xs][x + 1])
                                              137
         flo from [b][x + 1] = xs;
                                              138
                                              139
  set slack(b);
void expand blossom(int b) {
  REP(i, SZ(flo[b])) set st(flo[b][i], flo 141
                                              142
  int xr = flo_from[b][g[b][pa[b]].u], pr
                                              143
       = get_pr(b, xr);
                                              144
  for(int i = 0; i < pr; i += 2) {</pre>
                                              145
    int xs = flo[b][i], xns = flo[b][i +
                                              146
                                              147
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
                                              148
    slack[xs] = 0, set_slack(xns);
                                              149
    q_push(xns);
                                              150
                                              151
  S[xr] = 1, pa[xr] = pa[b];
                                              152
  for(size_t i = pr + 1; i < SZ(flo[b]);</pre>
       ++i) {
    int xs = flo[b][i];
                                              154
    S[xs] = -1, set slack(xs);
                                              155
                                              156
  st[b] = 0;
bool on found edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if(S[v] == -1) {
                                              158
    pa[v] = e.u, S[v] = 1;
                                              159
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
                                              160
    S[nu] = 0, q_push(nu);
                                              161
  } else if(S[v] == 0) {
                                              162
    int lca = get lca(u, v);
                                              163
    if(!lca) return augment(u,v), augment(
                                              164
         v,u), true;
    else add_blossom(u, lca, v);
                                              166
                                              167
  return false;
                                              168
bool matching() {
  memset(S + 1, -1, sizeof(int) * n x);
  memset(slack + 1, 0, sizeof(int) * n_x);
                                              171
  q = queue<int>();
                                              172
  REP(x, n x) if(st[x + 1] == x + 1 \&\& !
                                              173
       match[x + 1]) pa[x + 1] = 0, S[x +
                                              174
       1] = 0, q push(x + 1);
                                              175
  if(q.empty()) return false;
  while(true) {
                                              176
    while(!q.empty()) {
                                              177
      int u = q.front(); q.pop();
                                              178
      if(S[st[u]] == 1) continue;
                                              179
      for(int v = 1; v \leftarrow n; ++v)
                                              180
        if(g[u][v].w > 0 && st[u] != st[v
                                              181
              1) {
          if(e_delta(g[u][v]) == 0) {
             if(on found_edge(g[u][v]))
                                              183
                  return true:
                                              184
          } else update_slack(u, st[v]);
                                              185
                                              186 };
    int d = inf;
    for(int b = n + 1; b \le n x; ++b) if(
```

st[b] == b && S[b] == 1) d = min(d)

```
, lab[b] / 2);
    for(int x = 1; x <= n x; ++x) {
      if(st[x] == x && slack[x]) {
        if(S[x] == -1) d = min(d, e_delta(
             g[slack[x]][x]));
        else if(S[x] == 0) d = min(d,
             e_delta(g[slack[x]][x]) / 2);
    REP(u, n) {
      if(S[st[u + 1]] == 0) {
        if(lab[u + 1] <= d) return 0;</pre>
        lab[u + 1] -= d;
      } else if(S[st[u + 1]] == 1) lab[u +
            11 += d:
    for(int b = n + 1; b \le n \times +b)
      if(st[b] == b) {
        if(S[st[b]] == 0) lab[b] += d * 2;
        else if(S[st[b]] == 1) lab[b] -= d
    q = queue<int>();
    for(int x = 1; x \leftarrow n x; ++x)
                                              15
      if(st[x] == x && slack[x] && st[
                                              16
           slack[x]] != x && e_delta(g[
                                              17
           slack[x]][x]) == 0)
                                              18
        if(on_found_edge(g[slack[x]][x]))
                                              19
             return true;
    for(int b = n + 1; b \le n \times +b)
                                              21
      if(st[b] == b && S[b] == 1 && lab[b]
                                              22
                                              23
            == 0) expand blossom(b);
                                              24
                                              25
  return false;
pair<11, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
                                              27
  n x = n;
  int n_matches = 0;
                                              28
  11 \text{ tot_weight} = 0;
                                              29
  for(int u = 0; u <= n; ++u) st[u] = u,
                                              31
       flo[u].clear();
  int w max = 0;
                                              32
  for(int u = 1; u <= n; ++u)</pre>
                                              33
                                              34
    for(int v = 1; v <= n; ++v) {</pre>
      flo from[u][v] = (u == v ? u : 0);
      w \max = \max(w \max, g[u][v].w);
  for(int u = 1; u <= n; ++u) lab[u] =</pre>
                                               38
  while(matching()) ++n_matches;
  for(int u = 1; u <= n; ++u)
    if(match[u] && match[u] < u)</pre>
      tot weight += g[u][match[u]].w;
  return make pair(tot_weight, n_matches);
void add edge(int u, int v, int w) { g[u][
    v].w = g[v][u].w = w; }
void init(int _n) : n(_n) {
  REP(u, n) REP(v, n) g[u + 1][v + 1] =
       edge(u + 1, v + 1, 0);
                                               48
                                              49
```

3.7 KM

```
i template < class T>
2 struct KM {
   static constexpr T INF = numeric limits<T</pre>
        >::max():
   int n, ql, qr;
   vector<vector<T>> w;
   vector<T> hl, hr, slk;
   vi fl, fr, pre, qu;
   vector<bool> v1, vr;
   KM(int n) : n(n), w(n, vector<T>(n, -INF))
         , hl(n), hr(n), slk(n), fl(n), fr(n),
         pre(n), qu(n), vl(n), vr(n) {}
   void add_edge(int u, int v, int x) { w[u][
         v] = x; } // 最小值要加負號
    bool check(int x) {
     vl[x] = 1;
     if(fl[x] != -1) return vr[qu[qr++] = fl[
          x]] = 1;
      while(x != -1) swap(x, fr[fl[x] = pre[x
     return 0;
    void bfs(int s) {
     fill(all(slk), INF);
     fill(all(vl), 0);
     fill(all(vr), 0);
     ql = qr = 0, qu[qr++] = s, vr[s] = 1;
      while(true) {
       T d;
       while(ql < qr) {</pre>
         for(int x = 0, y = qu[ql++]; x < n;
           if(!vl[x] \&\& slk[x] >= (d = hl[x]
                 + hr[y] - w[x][y])) {
              pre[x] = y;
              if(d) slk[x] = d;
              else if(!check(x)) return;
       REP(x, n) if(!v1[x] \&\& d > s1k[x]) d =
             slk[x];
        REP(x, n) {
         if(vl[x]) hl[x] += d;
         else slk[x] -= d;
         if(vr[x]) hr[x] -= d;
       REP(x, n) if(!v1[x] && !s1k[x] && !
            check(x)) return;
   T solve() {
     fill(all(fl), -1);
     fill(all(fr), -1);
     fill(all(hr), 0);
     REP(i, n) hl[i] = *max element(all(w[i])
     REP(i, n) bfs(i);
     T ans = 0;
     REP(i, n) ans += w[i][fl[i]]; // i 跟 fl
          [i] 配對
     return ans;
```

3.8 max-clique

52 } 53 };

```
1 template<int V>
 struct max clique {
   using B = bitset<V>;
   int n = 0;
   vector<B> g, buf;
   struct P {
     int idx, col, deg;
     P(int a, int b, int b) : idx(a), col(b),
           deg(c) {}
   max clique(int n) : n( n), g( n), buf( n)
   void add edge(int a, int b) {
     assert(a != b);
     g[a][b] = g[b][a] = 1;
   vector<int> now, clique;
   void dfs(vector<P>& rem){
     if(SZ(clique) < SZ(now)) clique = now;</pre>
     sort(ALL(rem), [](P a, P b) { return a.
          deg > b.deg; });
     int max_c = 1;
     for(auto& p : rem){
       p.col = 0:
       while((g[p.idx] & buf[p.col]).any()) p
        \max c = \max(\max c, p.idx + 1);
       buf[p.col][p.idx] = 1;
     REP(i, max_c) buf[i].reset();
     sort(ALL(rem), [&](P a, P b) { return a.
          col < b.col; });</pre>
     for(; !rem.empty(); rem.pop_back()){
        auto& p = rem.back();
       if(now.size() + p.col + 1 <= clique.</pre>
             size()) break;
       vector<P> nrem;
       for(auto& q : rem){
         if(g[p.idx][q.idx]){
           nrem.emplace_back(q.idx, -1, 0);
           bs[q.idx] = \overline{1};
       for(auto& q : nrem) q.deg = (bs & g[q.
            idx]).count();
       now.emplace back(p.idx);
       dfs(nrem);
       now.pop_back();
   vector<int> solve(){
     vector<P> remark;
     REP(i, n) remark.emplace back(i, -1, SZ(
          g[i]));
     dfs(remark);
     return clique;
```

3.9 MCMF

class MCMF {

public:

i template < class S, class T>

51 };

```
struct Edge {
  int from, to;
  S cap;
  T cost;
  Edge(int u, int v, S x, T y) : from(u),
       to(v), cap(x), cost(y) {}
                                             72
const ll INF = 1e18L;
                                             73
int n;
vector<Edge> edges;
vector<vi>g;
vector<T> d;
vector<bool> ing;
vi pedge;
MCMF(int _n) : n(_n), g(_n), d(_n), inq(_n
     ), pedge(_n) {}
void add edge(int u, int v, S cap, T cost)
  g[u].pb(sz(edges));
  edges.eb(u, v, cap, cost);
  g[v].pb(sz(edges));
  edges.eb(v, u, 0, -cost);
bool spfa(int s, int t) {
  bool found = false;
  fill(all(d), INF);
  d[s] = 0;
  inq[s] = true;
  queue<int> q;
  q.push(s);
  while(!q.empty()) {
    int u = q.front(); q.pop();
    if(u == t) found = true;
    inq[u] = false;
    for(auto& id : g[u]) {
      const auto& e = edges[id];
      if(e.cap > 0 && d[u] + e.cost < d[e.</pre>
           to]) {
        d[e.to] = d[u] + e.cost;
        pedge[e.to] = id;
        if(!inq[e.to]) {
          q.push(e.to);
          inq[e.to] = true;
                                             23
                                             24
  return found;
pair<S, T> flow(int s, int t, S f = INF) {
 S cap = 0;
  T cost = 0:
                                             30
  while(f > 0 && spfa(s, t)) {
                                             31
    S send = f:
    int u = t;
                                             33
    while(u != s) {
```

```
const Edge& e = edges[pedge[u]];
send = min(send, e.cap);
u = e.from;
}
u = t;
while(u != s) {
Edge& e = edges[pedge[u]];
e.cap -= send;
Edge& b = edges[pedge[u] ^ 1];
b.cap += send;
u = e.from;
}
cap += send;
f -= send;
cost += send * d[t];
}
return {cap, cost};
}
```

3.10 minimum-general-weightedperfect-matching

// Minimum General Weighted Matching (

Perfect Match) 0-base

1 struct Graph {

```
static const int MXN = 105;
int n, edge[MXN][MXN];
int match[MXN], dis[MXN], onstk[MXN];
vector<int> stk;
void init(int n) {
  for(int i=0; i<n; i++)</pre>
    for(int j=0; j<n; j++)</pre>
      edge[i][j] = 0;
void add edge(int u, int v, int w) { edge[
u][v] = edge[v][u] = w; }
bool SPFA(int u){
  if(onstk[u]) return true;
  stk.push back(u);
  onstk[u] = 1:
  for(int v=0; v<n; v++){</pre>
    if(u != v && match[u] != v && !onstk[v
      int m = match[v];
      if(dis[m] > dis[u] - edge[v][m] +
           edge[u][v]){
        dis[m] = dis[u] - edge[v][m] +
             edge[u][v];
        onstk[v] = 1;
        stk.push back(v);
        if(SPFA(m)) return true;
        stk.pop_back();
        onstk[v] = 0;
  onstk[u] = 0;
  stk.pop_back();
  return false;
```

```
int solve() {
      for(int i = 0; i < n; i += 2) match[i] =</pre>
             i + 1, match[i+1] = i;
       while(true) {
38
         int found = 0;
         for(int i=0; i<n; i++) dis[i] = onstk[</pre>
         for(int i=0: i<n: i++){</pre>
           stk.clear();
           if(!onstk[i] && SPFA(i)){
             found = 1:
             while(stk.size()>=2){
               int u = stk.back(); stk.pop_back
               int v = stk.back(); stk.pop_back
               match[u] = v;
               match[v] = u;
         if(!found) break;
      int ans = 0:
      for(int i=0; i<n; i++) ans += edge[i][</pre>
            match[i]];
      return ans / 2;
57
58 } graph;
```

4 Geometry

4.1 closest-pair

```
| const | 11 | INF = 9e18L + 5;
vector<P> a:
  sort(all(a), [](P a, P b) { return a.x < b.x
4 11 SQ(11 x) { return x * x; }
5 11 solve(int 1, int r) {
    if(1 + 1 == r) return INF;
    int m = (1 + r) / 2;
    11 \text{ midx} = a[m].x;
    11 d = min(solve(1, m), solve(m, r));
    inplace_merge(a.begin() + 1, a.begin() + m
         , a.begin() + r, [](P a, P b) {
      return a.y < b.y;</pre>
11
12
13
    vector<P> p;
    for(int i = 1; i < r; ++i) if(SQ(a[i].x -</pre>
         midx) < d) p.pb(a[i]);
    REP(i, sz(p)) {
15
      for(int j = i + 1; j < sz(p); ++j) {
          d = min(d, SQ(p[i].x - p[j].x) + SQ(
               p[i].y - p[j].y));
         if(SQ(p[i].y - p[j].y) > d) break;
18
19
20
    return d; // 距離平方
```

4.2 convex-hull

```
void convex hull(vector<P>& dots) {
   sort(all(dots));
   vector<P> ans(1, dots[0]);
   for(int it = 0; it < 2; it++, reverse(all(</pre>
        dots))) {
     for(int i = 1, t = sz(ans); i < sz(dots)</pre>
          ; ans.pb(dots[i++])) {
        while(sz(ans) > t && ori(ans[sz(ans) -
             2], ans.back(), dots[i]) < 0) {
          ans.ppb();
   ans.ppb();
   swap(ans, dots);
```

half-plane

```
i typedef pair < double, double > pdd;
pdd operator-(pdd a,pdd b){return {a.F-b.F,a
      .S-b.S};}
pdd operator+(pdd a,pdd b){return {a.F+b.F,a
       .S+b.S};}
 pdd operator*(pdd a,double x){return {a.F*x,
5 double dot(pdd a,pdd b){return a.F*b.F+a.S*b
6 double cross(pdd a,pdd b){return a.F*b.S-a.S
      *b.F;}
 struct bpmj{
   const double eps=1e-8;
   int n,m,id,1,r;
   pdd pt[55],q[1100];
   struct line{
     pdd x,y;
     double z;
     line(pdd _x,pdd _y):x(_x),y(_y){z=atan2(
          y.S,y.F);}
     line(){}
     bool operator<(const line &a)const{</pre>
           return z<a.z;}</pre>
   }a[550],dq[1005];
   pdd get_(line x,line y){
     pdd v=x.x-v.x;
     double d=cross(y.y,v)/cross(x.y,y.y);
     return x.x+x.y*d;
   void solve(){
     dq[l=r=1]=a[1];
     for(int i=2;i<=id;++i){</pre>
        while(l<r&&cross(a[i].y,q[r-1]-a[i].x)</pre>
             <=eps) --r;
       while(l<r&&cross(a[i].y,q[1]-a[i].x)<=</pre>
             eps) ++1;
       dq[++r]=a[i];
        if(fabs(cross(dq[r].y,dq[r-1].y))<=eps 19</pre>
            ) {
          --r;
```

```
if(cross(dq[r].y,a[i].x-dq[r].x)>eps 22
             ) dq[r]=a[i];
      if(l<r) q[r-1]=get_(dq[r-1],dq[r]);</pre>
    while(l<r&&cross(dq[1].y,q[r-1]-dq[1].x)</pre>
         <=eps) --r;
    if(r-1<=1) return:</pre>
    q[r]=get_(dq[1],dq[r]);
  void cal(){
    double ans=0;
    q[r+1]=q[1];
    for(int i=1;i<=r;++i) ans+=cross(q[i],q[</pre>
    cout<<fixed<<setprecision(3)<<ans/2<<"\n</pre>
 void main_(){
    cin>>n:
    for(int x,y,i=0;i<n;++i){</pre>
      for(int i=0;i<m;++i) cin>>pt[i].F>>pt[
           i].S;
      pt[m]=pt[0];
      for(int i=0;i<m;++i) a[++id]=line(pt[i</pre>
           ],pt[i+1]-pt[i]);
    sort(a+1,a+1+id);
    solve();
    cal();
}valderyaya;
```

4.4 min-enclosing-circle

```
pdd excenter(pdd x, pdd y, pdd z) {
  #define f(x, y) (x*x+y*y)
  auto [x1, y1] = x;
  auto [x2, y2] = y;
  auto [x3, y3] = z;
  double d1 = f(x2, y2) - f(x1, y1), d2 = f(
       x3, y3) - f(x2, y2);
  double fm = 2 * ((y3 - y2) * (x2 - x1) - (
       y2 - y1) * (x3 - x2));
  double ans_x = ((y3 - y2) * d1 - (y2 - y1)
       * d2) / fm;
  double ans_y = ((x2 - x1) * d2 - (x3 - x2)
       * d1) / fm;
  #undef f
  return {ans x, ans y};
pdd min_enclosing_circle(vector<pdd> dots,
    double& r) {
  random shuffle(ALL(dots));
  pdd C = dots[0];
  r = 0;
  #define check(i, j) REP(i, j) if(abs(dots[
       i] - C) > r
  check(i, SZ(dots)) {
    C = dots[i], r = 0;
    check(j, i) {
```

```
C = (dots[i] + dots[j]) / 2.0;
        r = abs(dots[i] - C);
        check(k, j) {
          C = excenter(dots[i], dots[j], dots[
          r = abs(dots[i] - C);
      }
    #undef check
    return C;
32 }
```

point-in-convex-hull

23

24

28

29

```
int point in convex hull(const vector<P>& a,
        P p) {
    // -1 ON, 0 OUT, +1 IN
    // 要先逆時針排序
    int n = sz(a);
    if(btw(a[0], a[1], p) || btw(a[0], a[n -
        1], p)) return -1;
    int 1 = 0, r = n - 1;
    while(1 <= r) {</pre>
      int m = (1 + r) / 2;
      auto a1 = cross(a[m] - a[0], p - a[0]);
      auto a2 = cross(a[(m + 1) % n] - a[0], p
           - a[0]);
      if(a1 >= 0 && a2 <= 0) {
        auto res = cross(a[(m + 1) % n] - a[m
            ], p - a[m]);
        return res > 0 ? 1 : (res >= 0 ? -1 :
             0);
      if(a1 < 0) r = m - 1;
      else 1 = m + 1:
17
    return 0;
```

4.6 point

```
using P = pair<11, 11>;
2 P operator+(P a, P b) { return P{a.X + b.X,
      a.Y + b.Y; }
3 P operator-(P a, P b) { return P{a.X - b.X,
      a.Y - b.Y}; }
4 P operator*(P a, 11 b) { return P{a.X * b, a
      .Y * b}; }
5 P operator/(P a, 11 b) { return P{a.X / b, a
      .Y / b}; }
6 11 dot(P a, P b) { return a.X * b.X + a.Y *
      b.Y; }
7 | 11 cross(P a, P b) { return a.X * b.Y - a.Y
      * b.X; }
8 11 abs2(P a) { return dot(a, a); }
9 double abs(P a) { return sqrt(abs2(a)); }
int sign(ll x) { return x < 0 ? -1 : (x == 0
       ? 0 : 1); }
```

```
int ori(P a, P b, P c) { return sign(cross(b)
        - a, c - a)); }
12 bool collinear(P a, P b, P c) { return sign(
       cross(a - c, b - c)) == 0; }
13 bool btw(Pa, Pb, Pc) {
    if(!collinear(a, b, c)) return 0;
    return sign(dot(a - c, b - c)) <= 0;</pre>
16 }
17 bool seg intersect(Pa, Pb, Pc, Pd) {
    int a123 = ori(a, b, c);
    int a124 = ori(a, b, d);
    int a341 = ori(c, d, a);
    int a342 = ori(c, d, b);
    if(a123 == 0 && a124 == 0) {
      return btw(a, b, c) || btw(a, b, d) ||
           btw(c, d, a) || btw(c, d, b);
24
    return a123 * a124 <= 0 && a341 * a342 <=
28 P intersect(P a, P b, P c, P d) {
    int a123 = cross(b - a, c - a);
    int a124 = cross(b - a, d - a);
    return (d * a123 - c * a124) / (a123 -
          a124);
33 struct line { P A, B; };
34 P vec(line L) { return L.B - L.A; }
35 P projection(P p, line L) { return L.A + vec
       (L) / abs(vec(L)) * dot(p - L.A, vec(L))
        / abs(vec(L)); }
```

4.7 polar-angle-sort

```
1 bool cmp(P a, P b) {
   #define ng(k) (sign(k.Y) < 0 | (sign(k.Y)
         == 0 \&\& sign(k.X) < 0))
   int A = ng(a), B = ng(b);
   if(A != B) return A < B;</pre>
   if(sign(cross(a, b)) == 0) return abs2(a)
        < abs2(b):
   return sign(cross(a, b)) > 0;
```

4.8 定理

- 皮克定理
 - 若一個多邊形的所有頂點都在整數點上,則該 多邊形的面積 $S = a + \frac{b}{3} - 1$ · 其中 a 為內部 格點數目, b 為邊上格點數目。

Graph

5.1 2-SAT

```
1 struct two sat {
    int n; SCC g;
    vector<bool> ans:
    two_sat(int _n) : n(_n), g(_n * 2) {}
    void add_or(int u, bool x, int v, bool y)
      g.add_edge(2 * u + !x, 2 * v + y);
      g.add edge(2 * v + !y, 2 * u + x);
    bool solve() {
      ans.resize(n);
      auto id = g.solve();
      REP(i, n) {
        if(id[2 * i] == id[2 * i + 1]) return
        ans[i] = (id[2 * i] < id[2 * i + 1]);
      return true;
18 };
```

5.2 BCC-tree

```
i struct BlockCutTree {
   int n;
   vector<vi> g;
   vi dfn, low, stk;
   int cnt = 0, cur = 0;
   vector<pii> edges;
   BlockCutTree(int _n) : n(_n), g(_n), dfn(
        n), low( n) {}
   void ae(int u, int v) {
     g[u].pb(v);
     g[v].pb(u);
   void dfs(int x) {
     stk.pb(x);
     dfn[x] = low[x] = cur++;
     for(auto y : g[x]) {
       if(dfn[y] == -1) {
         dfs(y);
         low[x] = min(low[x], low[y]);
         if(low[y] == dfn[x]) {
           int v;
           do {
             v = stk.back(), stk.pop_back();
             edges.eb(n + cnt, v);
           } while (v != y);
           edges.eb(x, n + cnt);
           cnt++:
       } else low[x] = min(low[x], dfn[y]);
   pair<int, vector<pii>> work() {
     REP(i, n) {
       if(dfn[i] == -1) {
         stk.clear();
         dfs(i);
     return {cnt, edges};
```

centroid-tree

int n = sz(g);

vector<vi>& g) {

40 };

```
vi siz(n);
vector<bool> vis(n);
auto dfs_sz = [&](auto f, int u, int p) ->
      void {
  siz[u] = 1;
  for(auto v : g[u]) {
  if(v == p || vis[v]) continue;
    f(f, v, u);
    siz[u] += siz[v];
auto find_cd = [&](auto f, int u, int p,
     int all) -> int {
  for(auto v : g[u]) {
  if(v == p || vis[v]) continue;
    if(siz[v] * 2 > all) return f(f, v, u,
          all);
  return u;
vector<vi> h(n);
auto build = [&](auto f, int u) -> int {
  dfs_sz(dfs_sz, u, -1);
  int cd = find cd(find cd, u, -1, siz[u])
  vis[cd] = true;
  for(auto v : g[cd]) {
    if(vis[v]) continue;
    int child = f(f, v);
    h[cd].pb(child);
  return cd;
int root = build(build, 0);
return {root, h};
```

| pair<int, vector<vi>>> centroid_tree(const

5.4 chromatic-number

```
1 // vi to(n);
2 // to[u] |= 1 << v;
3 // to[v] |= 1 << u;
 int chromatic number(vi g) {
    constexpr int MOD = 998244353;
    int n = SZ(g);
    vector<int> I(1 << n); I[0] = 1;
    for(int s = 1; s < (1 << n); s++) {</pre>
      int v = __builtin_ctz(s), t = s ^ (1 <<</pre>
           v);
     I[s] = (I[t] + I[t \& \sim g[v]]) \% MOD;
    auto f = I;
   for(int k = 1; k <= n; k++) {</pre>
```

```
int sum = 0;
      REP(s, 1 \ll n) {
        if(( builtin popcount(s) ^ n) & 1)
             sum -= f[s];
        else sum += f[s];
        sum = ((sum % MOD) + MOD) % MOD;
        f[s] = 1LL * f[s] * I[s] % MOD;
      if(sum != 0) return k;
                                                 49
22
    return 48763:
24 }
```

5.5 HLD

26

29

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39

int d = depth[u] - k;

while(depth[top[u]] > d) u = par[top[u

return tour[id[u] + d - depth[u]];

```
1 struct HLD {
   int n;
   vector<vi> g;
   vi siz, par, depth, top, tour, fi, id;
   sparse_table<pii, min> st;
   HLD(int _n) : n(_n), g(_n), siz(_n), par(
        _n), depth(_n), top(_n), fi(_n), id(_n 61
     tour.reserve(n);
   void add edge(int u, int v) {
     g[u].push back(v);
     g[v].push_back(u);
   void build(int root = 0) {
     par[root] = -1;
     top[root] = root;
     vector<pii> euler tour;
     euler tour.reserve(2 * n - 1);
     dfs_sz(root);
                                                 72
     dfs link(euler tour, root);
     st = sparse table<pii, min>(euler tour);
    int get_lca(int u, int v) {
     int L = fi[u], R = fi[v];
     if(L > R) swap(L, R);
     return st.prod(L, R).second;
   bool is_anc(int u, int v) {
     return id[u] <= id[v] && id[v] < id[u] +</pre>
           siz[u];
                                                 83
   bool on path(int a, int b, int x) {
     return (is ancestor(x, a) || is ancestor
          (x, b)) && is ancestor(get lca(a, b)
          , x);
   int get dist(int u, int v) {
     return depth[u] + depth[v] - 2 * depth[(
          get_lca(u, v))];
   int kth_anc(int u, int k) {
     if(depth[u] < k) return -1;</pre>
```

```
vector<pii> get_path(int u, int v, bool
     include lca = true) {
  if(u == v && !include lca) return {};
  vector<pii> seg;
  while(top[u] != top[v]) {
   if(depth[top[u]] > depth[top[v]]) swap
    seg.eb(id[top[v]], id[v]);
   v = par[top[v]];
  if(depth[u] > depth[v]) swap(u, v); // u
  if(u != v || include_lca) seg.eb(id[u] +
       !include_lca, id[v]);
  return seg:
void dfs_sz(int u) {
 if(par[u] != -1) g[u].erase(find(all(g[u
      ]), par[u]));
  siz[u] = 1;
  for(auto& v : g[u]) {
   par[v] = u;
    depth[v] = depth[u] + 1;
    dfs sz(v);
    siz[u] += siz[v];
   if(siz[v] > siz[g[u][0]]) swap(v, g[u
        ][0]);
void dfs_link(vector<pii>& euler_tour, int
```

fi[u] = sz(euler tour);

euler_tour.eb(depth[u], u);

dfs link(euler tour, v);

euler tour.eb(depth[u], u);

top[v] = (v == g[u][0] ? top[u] : v);

id[u] = sz(tour);

for(auto v : g[u]) {

tour.pb(u);

int kth node on path(int a, int b, int k)

 $if(k < 0 \mid | k > fi + se)$ return -1;

if(k < fi) return kth anc(a, k);</pre>

return kth anc(b, fi + se - k);

int z = get lca(a, b);

int fi = depth[a] - depth[z];

int se = depth[b] - depth[z];

5.6 lowlink

```
1 | struct lowlink {
   int n, cnt = 0, tecc_cnt = 0, tvcc_cnt =
   vector<vector<pii>>> g;
   vector<pii> edges;
   vi roots, id, low, tecc_id, tvcc_id;
   vector<bool> is bridge, is cut,
        is tree edge;
```

```
lowlink(int _n) : n(_n), g(_n), is_cut(_n,
      false), id(_n, -1), low(_n, -1) {}
void add edge(int u, int v) {
  g[u].eb(v, sz(edges));
 g[v].eb(u, sz(edges));
  edges.eb(u, v);
  is bridge.pb(false);
  is tree edge.pb(false);
  tvcc_id.pb(-1);
void dfs(int u, int peid = -1) {
  static vi stk;
  static int rid;
  if(peid < 0) rid = cnt;</pre>
  if(peid == -1) roots.pb(u);
  id[u] = low[u] = cnt++;
  for(auto [v, eid] : g[u]) {
    if(eid == peid) continue;
    if(id[v] < id[u]) stk.pb(eid);</pre>
    if(id[v] >= 0) {
      low[u] = min(low[u], id[v]);
    } else {
      is_tree_edge[eid] = true;
      dfs(v, eid);
      low[u] = min(low[u], low[v]);
      if((id[u] == rid && id[v] != rid +
           1) || (id[u] != rid && low[v] >=
            id[u])) {
        is_cut[u] = true;
      if(low[v] >= id[u]) {
        while(true) {
          int e = stk.back();
          stk.pop_back();
          tvcc id[e] = tvcc cnt;
          if(e == eid) break;
        tvcc_cnt++;
void build() {
  REP(i, n) if(id[i] < 0) dfs(i);</pre>
 REP(i, sz(edges)) {
    auto [u, v] = edges[i];
    if(id[u] > id[v]) swap(u, v);
    is_bridge[i] = (id[u] < low[v]);</pre>
vector<vi> two ecc() { // 邊雙
  tecc_cnt = 0;
  tecc id.assign(n, -1);
  vi stk;
  REP(i, n) {
    if(tecc id[i] != -1) continue;
    tecc_id[i] = tecc_cnt;
    stk.pb(i);
    while(sz(stk)) {
      int u = stk.back(); stk.pop_back();
      for(auto [v, eid] : g[u]) {
        if(tecc id[v] >= 0 || is bridge[
             eid]) {
                                             19
          continue;
                                             20
```

```
tecc_id[v] = tecc_cnt;
       stk.pb(v);
   tecc_cnt++;
 vector<vi> comp(tecc cnt);
 REP(i, n) comp[tecc id[i]].pb(i);
 return comp;
vector<vi> bcc_vertices() { // 點雙
 vector<vi> comp(tvcc cnt);
 REP(i, sz(edges)) {
   comp[tvcc_id[i]].pb(edges[i].first);
   comp[tvcc id[i]].pb(edges[i].second);
  for(auto& v : comp) {
   sort(all(v));
   v.erase(unique(all(v)), v.end());
 REP(i, n) if(g[i].empty()) comp.pb({i});
 return comp;
vector<vi> bcc edges() {
 vector<vi> ret(tvcc cnt);
 REP(i, sz(edges)) ret[tvcc id[i]].pb(i);
 return ret;
```

manhattan-mst

```
template < class T> // [w, u, v]
vector<tuple<T, int, int>> manhattan mst(
     vector<T> xs, vector<T> ys) {
    const int n = SZ(xs);
    vi idx(n);
    iota(ALL(idx), 0);
    vector<tuple<T, int, int>> ret;
    REP(s, 2) {
        REP(t, 2) {
            auto cmp = [&](int i, int j) {
                 return xs[i] + ys[i] < xs[j]</pre>
                  + ys[j]; };
            sort(ALL(idx), cmp);
            map<T, int> sweep;
            for(int i : idx) {
                for(auto it = sweep.
                     lower_bound(-ys[i]); it
                     != sweep.end(): it =
                     sweep.erase(it)) {
                    int j = it->second;
                    if(xs[i] - xs[j] < ys[i]
                           - ys[j]) break;
                    ret.eb(abs(xs[i] - xs[j
                         ]) + abs(ys[i] - ys[
                         j]), i, j);
                sweep[-ys[i]] = i;
            swap(xs, ys);
        for(auto &x : xs) x = -x;
```

```
5.8 SCC
```

sort(ALL(ret));

return ret:

1 struct SCC {

int n;

24

25

26

27

33

35

37 };

26 }

```
vector<vi> g, h;
SCC(int _n) : n(_n), g(_n), h(_n) {}
void add_edge(int u, int v) {
 g[u].pb(v);
 h[v].pb(u);
vi solve() { // 回傳縮點的編號
 vi id(n), top;
 top.reserve(n);
 #define GO if(id[v] == 0) dfs1(v);
 function<void(int)> dfs1 = [&](int u) {
   id[u] = 1;
   for(auto v : g[u]) GO;
   top.pb(u);
 REP(v, n) GO;
 fill(all(id), -1);
 function<void(int, int)> dfs2 = [&](int
      u, int x) {
   id[u] = x:
   for(auto v : h[u]) {
     if(id[v] == -1) {
       dfs2(v, x);
   }
 for(int i = n - 1, cnt = 0; i >= 0; --i)
   int u = top[i];
   if(id[u] == -1) {
      dfs2(u, cnt);
      cnt += 1;
 return id:
```

triangle-sum

```
1 // Three vertices a < b < cconnected by</pre>
       three edges {a, b}, {a, c}, {b, c}. Find 8
        xa * xb * xc over all triangles.
2 int triangle_sum(vector<array<int, 2>> edges
      , vi x) {
   int n = SZ(x);
   vi deg(n);
    vector<vector<int>> g(n);
    for(auto& [u, v] : edges) {
     if(u > v) swap(u, v);
```

```
deg[u]++, deg[v]++;
     REP(i, n) g[i].reserve(deg[i]);
     for(auto [u, v] : edges) {
12
       if(deg[u] > deg[v]) swap(u, v);
13
       g[u].pb(v);
14
15
    vi val(n);
      int128 ans = 0;
16
17
     REP(a, n) {
       for(auto b : g[a]) val[b] = x[b];
       for(auto b : g[a]) {
        11 \text{ tmp} = 0;
21
         for(auto c : g[b]) tmp += val[c];
22
         ans += _{int128(tmp)} * x[a] * x[b];
23
24
       for(auto b : g[a]) val[b] = 0;
25
    return ans % mod:
27
```

Math

6.1 Aliens

11 |

12

13

```
1 template < class Func, bool MAX>
2 11 Aliens(11 1, 11 r, int k, Func f) {
    while(1 < r) {</pre>
      11 m = 1 + (r - 1) / 2;
      auto [score, op] = f(m);
      if(op == k) return score + m * k * (MAX
           ? +1 : -1);
      if(op < k) r = m;
      else 1 = m + 1;
    return f(1).first + 1 * k * (MAX ? +1 :
         -1);
11 }
```

6.2 Berlekamp-Massey

```
1 // - [1, 2, 4, 8, 16] -> (1, [1, -2])
2 // - [1, 1, 2, 3, 5, 8] -> (2, [1, -1, -1])
3 // - [0, 0, 0, 0, 1] -> (5, [1, 0, 0, 0, 0,
       998244352]) (mod 998244353)
4 // - [] -> (0, [1])
5 // - [0, 0, 0] -> (0, [1])
6 // - [-2] -> (1, [1, 2])
 template<class T>
 pair<int, vector<T>> BM(const vector<T>& S)
   using poly = vector<T>;
   int N = SZ(S);
   poly C_rev{1}, B{1};
   int L = 0, m = 1;
   T b = 1:
   auto adjust = [](poly C, const poly &B, T
        d, T b, int m) -> poly {
```

```
C.resize(max(SZ(C), SZ(B) + m));
      Ta = d / b;
      REP(i, SZ(B)) C[i + m] -= a * B[i];
      return C;
    REP(n, N) {
      T d = S[n];
      REP(i, L) d += C rev[i + 1] * S[n - 1 -
      if(d == 0) m++;
      else if (2 * L <= n) {</pre>
        poly Q = C_rev;
        C_rev = adjust(C_rev, B, d, b, m);
        L = n + 1 - L, B = Q, b = d, m = 1;
      } else C_rev = adjust(C_rev, B, d, b, m
           ++);
    return {L, C_rev};
33 // Calculate x^N \not f(x)
34 // Complexity: $0(K^2 \log N)$ ($K$: deg. of
35 // (4, [1, -1, -1]) -> [2, 3]
37 template < class T>
38 vector<T> monomial mod polynomial(long long
       N, const vector<T> &f_rev) {
    assert(!f rev.empty() && f rev[0] == 1);
    int K = SZ(f_rev) - 1;
    if(!K) return {};
    int D = 64 - builtin clzll(N);
    vector<T> ret\overline{(K, 0)};
    ret[0] = 1;
    auto self_conv = [](vector<T> x) -> vector
         <T> {
      int d = SZ(x);
      vector<T> ret(d * 2 - 1);
      REP(i, d) {
        ret[i * 2] += x[i] * x[i];
        REP(j, i) ret[i + j] += x[i] * x[j] *
      return ret;
    for(int d = D; d--;) {
      ret = self conv(ret);
      for(int i = 2 * K - 2; i >= K; i--) {
        REP(j, k) ret[i - j - 1] -= ret[i] *
            f rev[i + 1];
      ret.resize(K);
      if (N >> d & 1) {
        vector<T> c(K);
        c[0] = -ret[K - 1] * f_rev[K];
        for(int i = 1; i < K; i++) c[i] = ret[</pre>
            i - 1] - ret[K - 1] * f_rev[K - i
        ret = c:
    return ret:
```

6.3 Chinese-Remainder

```
1 // (rem, mod) {0, 0} for no solution
 pair<ll, ll> crt(ll r0, ll m0, ll r1, ll m1)
   r0 = (r0 \% m0 + m0) \% m0;
   r1 = (r1 \% m1 + m1) \% m1;
   if(m0 < m1) swap(r0, r1), swap(m0, m1);</pre>
   if(m0 \% m1 == 0) {
     if(r0 % m1 != r1) return {0, 0};
   11 g, im, qq;
   g = ext_gcd(m0, m1, im, qq);
   11 u1 = (m1 / g);
   if((r1 - r0) % g) return {0, 0};
   11 x = (r1 - r0) / g % u1 * im % u1;
   r0 += x * m0;
   m0 *= u1:
   if(r0 < 0) r0 += m0;
   return {r0, m0};
```

6.4 Combination

6.5 Determinant

```
1 T det(vector<vector<T>> a) {
2 int n = SZ(a);
```

```
T ret = 1;
     REP(i, n) {
       int idx = -1:
       for(int j = i; j < n; j++) {</pre>
         if(a[j][i] != 0) {
           idx = j;
           break;
       if(idx == -1) return 0;
       if(i != idx) {
         ret *= T(-1);
         swap(a[i], a[idx]);
       ret *= a[i][i];
       T \text{ inv} = T(1) / a[i][i];
       REP(j, n) a[i][j] *= inv;
       for(int j = i + 1; j < n; j++) {</pre>
         T \times = a[j][i];
         if(x == 0) continue;
         for(int k = i; k < n; k++) {</pre>
23
           a[j][k] -= a[i][k] * x;
25
       }
26
27
28
    return ret;
```

6.6 Discrete-Log

6.7 extgcd

```
1  // ax + by = gcd(a, b)
2  ll ext_gcd(ll a, ll b, ll& x, ll& y) {
3   if(b == 0) {
4      x = 1, y = 0;
5      return a;
6   }
7   ll x1, y1;
8   ll g = ext_gcd(b, a % b, x1, y1);
9   x = y1, y = x1 - (a / b) * y1;
10   return g;
11 }
```

6.8 Floor-Sum

```
1 / / sum_{i} = 0 ^{n - 1} floor((ai + b) / c)
       in O(a + b + c + n)
2 11 floor sum(11 n, 11 a, 11 b, 11 c) {
    assert(0 <= n && n < (1LL << 32));
    assert(1 <= c && c < (1LL << 32));
    ull ans = 0:
    if(a < 0) {
      ull a2 = (a \% c + c) \% c;
       ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a
           ) / c);
     if(b < 0) {
      ull b2 = (b \% c + c) \% c;
       ans -= 1ULL * n * ((b2 - b) / c);
    ull N = n, C = c, A = a, B = b;
     while(true) {
         ans += N * (N - 1) / 2 * (A / C);
        A %= C:
      if(B >= C) {
         ans += N * (B / C);
        B %= C;
      ull y max = A * N + B;
      if(y_max < C) break;</pre>
      N = y \max / C, B = y \max % C;
      swap(C, A);
31
    return ans;
32 }
```

6.9 FWHT

```
| #define ppc __builtin_popcount
2 template < class T, class F>
3 void fwht(vector<T>& a, F f) {
    int n = SZ(a);
    assert(ppc(n) == 1);
    for(int i = 1; i < n; i <<= 1) {</pre>
      for(int j = 0; j < n; j += i << 1) {</pre>
        REP(k, i) f(a[j + k], a[i + j + k]);
11 }
12 template < class T>
void or transform(vector<T>& a, bool inv) {
       fwht(a, [\&](T\& x, T\& y) \{ y += x * (inv) \}
       ? -1 : +1); }) }
14 template < class T>
void and_transform(vector<T>& a, bool inv) {
        fwht(a, [\&](T\& x, T\& y) \{ x += y * (inv) \}
        ? -1 : +1); }); }
16 template < class T>
void xor transform(vector<T>& a, bool inv) {
    fwht(a, [](T& x, T& y) {
     Tz = x + y;
```

```
y = x - y;
     x = z;
    });
   if(inv) {
     Tz = T(1) / T(SZ(a));
      for(auto& x : a) x *= z;
28 template < class T>
29 vector<T> convolution(vector<T> a, vector<T>
    assert(SZ(a) == SZ(b));
    transform(a, false), transform(b, false);
    REP(i, SZ(a)) a[i] *= b[i];
    transform(a, true);
    return a:
36 template < class T>
  vector<T> subset convolution(const vector<T</pre>
       >& f, const vector<T>& g) {
    assert(SZ(f) == SZ(g));
    int n = SZ(f);
    assert(ppc(n) == 1);
    const int lg = __lg(n);
    vector<vector<T>> fhat(lg + 1, vector<T>(n
         )), ghat(fhat);
    REP(i, n) fhat[ppc(i)][i] = f[i], ghat[ppc
         (i)|[i] = g[i];
    REP(i, lg + 1) or_transform(fhat[i], false 11
        ), or transform(ghat[i], false);
    vector<vector<T>> h(lg + 1, vector<T>(n)); 13
    REP(m, n) REP(i, lg + 1) REP(j, i + 1) h[i]
         ][m] += fhat[j][m] * ghat[i - j][m];
    REP(i, lg + 1) or_transform(h[i], true);
    vector<T> res(n);
    REP(i, n) res[i] = h[ppc(i)][i];
    return res;
```

6.10 Gauss-Jordan

```
i int GaussJordan(vector<vector<ld>>& a) {
   // -1 no sol. 0 inf sol
   int n = SZ(a);
   REP(i, n) assert(SZ(a[i]) == n + 1);
   REP(i, n) {
     int p = i;
     REP(j, n) {
       if(j < i && abs(a[j][j]) > EPS)
       if(abs(a[j][i]) > abs(a[p][i])) p = j;
     REP(j, n + 1) swap(a[i][j], a[p][j]);
     if(abs(a[i][i]) <= EPS) continue;</pre>
     REP(j, n) {
       if(i == j) continue;
       ld delta = a[j][i] / a[i][i];
       FOR(k, i, n + 1) a[j][k] -= delta * a[
            i][k];
   bool ok = true;
   REP(i, n) {
```

```
if(abs(a[i][i]) <= EPS) {
    if(abs(a[i][n]) > EPS) return -1;
    ok = false;
}

return ok;
}
```

6.11 GCD-Convolution

if(!sieve[i]) continue;

vector<int> ret{2, 3};

vector<int> prime enumerate(int N) {

vector<bool> sieve(N / 3 + 1, 1);

for(int p = 5, d = 4, i = 1, sqn = sqrt(N)

; $p \le sqn$; p += d = 6 - d, i++) {

for(int q = p * p / 3, r = d * p / 3 + (

d * p % 3 == 2), s = 2 * p; q < SZ(

sieve); q += r = s - r) sieve[q] =

1 // 2, 3, 5, 7, ...

```
for(int p = 5, d = 4, i = 1; p <= N; p +=
         d = 6 - d, i++) {
      if(sieve[i]) {
        ret.pb(p);
    while(SZ(ret) && ret.back() > N) ret.
         pop_back();
    return ret;
  struct divisor transform {
    template < class T>
    static void zeta_transform(vector<T>& a) {
      int n = a.size() - 1;
      for(auto p : prime enumerate(n)) {
        for(int i = 1; i * p <= n; i++) {</pre>
          a[i * p] += a[i];
    template < class T>
    static void mobius_transform(vector<T>& a)
      int n = a.size() - 1;
      for(auto p : prime enumerate(n)) {
        for(int i = n / p; i > 0; i--) {
          a[i * p] -= a[i];
36 };
  struct multiple transform {
    template < class T>
    static void zeta transform(vector<T>& a) {
      int n = a.size() - 1;
      for(auto p : prime_enumerate(n)) {
        for(int i = n / p; i > 0; i--) {
          a[i] += a[i * p];
```

```
static void mobius transform(vector<T>& a)
      int n = a.size() - 1;
      for(auto p : prime enumerate(n)) {
        for(int i = 1; i * p <= n; i++) {</pre>
          a[i] -= a[i * p];
55
57 // lcm: multiple -> divisor
58 template < class T>
59 vector<T> gcd convolution(const vector<T>& a
       , const vector<T>& b) {
    assert(a.size() == b.size());
    auto f = a, g = b;
    multiple transform::zeta_transform(f);
    multiple transform::zeta transform(g);
    REP(i, SZ(f)) f[i] *= g[i];
    multiple transform::mobius_transform(f);
    return f;
```

6.12 Int-Div

template < class T>

6.13 Linear-Sieve

```
vi primes, least = {0, 1}, phi, mobius;
void LinearSieve(int n) {
   least = phi = mobius = vi(n + 1);
    mobius[1] = 1;
    for(int i = 2; i <= n; i++) {</pre>
      if(!least[i]) {
        least[i] = i;
        primes.pb(i);
        phi[i] = i - 1;
        mobius[i] = -1;
      for(auto j : primes) {
        if(i * j > n) break;
        least[i * j] = j;
        if(i % j == 0) {
          mobius[i * j] = 0;
          phi[i * j] = phi[i] * j;
          break:
        } else {
          mobius[i * j] = -mobius[i];
          phi[i * j] = phi[i] * phi[j];
23
24
25 }
```

6.14 Miller-Rabin

```
1 bool is prime(ll n, vector<ll> x) {
   11 d = n - 1;
    d >>= builtin ctzll(d);
    for(auto a : x) {
      if(n <= a) break;</pre>
      11 t = d, y = 1, b = t;
      while(b) {
        if(b & 1) y = i128(y) * a % n;
        a = i128(a) * a % n;
        b >>= 1;
      while(t != n - 1 && y != 1 && y != n -
        y = i128(y) * y % n;
14
        t <<= 1;
      if(y != n - 1 && t % 2 == 0) return
           false:
    return true;
20 bool is prime(ll n) {
   if(n <= 1) return false;</pre>
    if(n \% 2 == 0) return n == 2;
    if(n < (1LL << 30)) return is_prime(n, {2,</pre>
    return is_prime(n, {2, 325, 9375, 28178,
         450775, 9780504, 1795265022});
25 }
```

6.15 Min-of-Mod-of-Linear

6.16 Mod-Inv

```
int inv(int a) {
   if(a < N) return inv[a];
   if(a == 1) 1;
   return (MOD - 1LL * (MOD / a) * inv(MOD %
        a) % MOD) % MOD;
}

vi mod_inverse(int m, int n = -1) {
   assert(n < m);
   if(n == -1) n = m - 1;</pre>
```

```
vi inv(n + 1);
inv[0] = inv[1] = 1;
for(int i = 2; i <= n; i++) inv[i] = m - 1
     LL * (m / i) * inv[m % i] % m;
return inv;
```

6.17 Pollard-Rho

```
1 | void PollardRho(map<11, int>& mp, 11 n) {
   if(n == 1) return;
   if(is prime(n)) return mp[n]++, void();
   if(n % 2 == 0) {
     mp[2] += 1;
     PollardRho(mp, n / 2);
     return;
   11 x = 2, y = 2, d = 1, p = 1;
   #define f(x, n, p) ((i128(x) * x % n + p)
        % n)
   while(true) {
     if(d!= 1 && d!= n) {
       PollardRho(mp, d);
       PollardRho(mp, n / d);
       return;
     p += (d == n);
     x = f(x, n, p), y = f(f(y, n, p), n, p);
     d = \underline{gcd(abs(x - y), n)};
   #undef f
 vector<ll> get divisors(ll n) {
   if(n == 0) return {};
   map<ll, int> mp;
   PollardRho(mp, n);
   vector<pair<ll, int>> v(all(mp));
   vector<ll> res:
   auto f = [&](auto f, int i, ll x) -> void
     if(i == sz(v)) {
       res.pb(x);
       return;
     for(int j = v[i].second; ; j--) {
       f(f, i + 1, x);
       if(i == 0) break;
       x *= v[i].first;
   f(f, 0, 1);
   sort(all(res));
   return res;
```

6.18 Primes

```
1 /* 12721 13331 14341 75577 123457 222557
      556679 999983 1097774749 1076767633
      100102021 999997771 1001010013
      1000512343 987654361 999991231 999888733 23
       98789101 987777733 999991921 1010101333 24
       1010102101 10000000000039
      1000000000000037 2305843009213693951
      4611686018427387847 9223372036854775783
      18446744073709551557 */
```

6.19 Triangle

```
1 // Counts x, y >= 0 such that <math>Ax + By <= C.
       Requires A, B > 0. Runs in log time.
2 // Also representable as sum_{0 <= x <= C /
       A) floor((C - Ax) / B + 1).
3 11 count triangle(11 A, 11 B, 11 C) {
      if(C < 0) return 0;
      if(A < B) swap(A, B);
      11 m = C / A, k = A / B;
      11 h = (C - m * A) / B + 1;
      return m * (m + 1) / 2 * k + (m + 1) * h 43
            + count_triangle(B, A - k * B, C - 44
           B * (k * m + h));
  // Counts 0 \leftarrow x \leftarrow RA, 0 \leftarrow y \leftarrow RB such that 46
        Ax + By \leftarrow C. Requires A, B > 0.
12 11 count_triangle_rectangle_intersection(11
       A, 11 B, 11 C, 11 RA, 11 RB) {
      if(C < 0 || RA <= 0 || RB <= 0) return
      if(C >= A * (RA - 1) + B * (RB - 1))
           return RA * RB:
      return count_triangle(A, B, C) -
           count triangle(A, B, C - A * RA) -
            count_triangle(A, B, C - B * RB);
```

6.20 Xor-Basis

```
1 template<int B>
 struct xor_basis {
   using T = long long;
   bool zero = false, change = false;
   int cnt = 0;
   array<T, B > p = \{\};
   vector<T> d;
   void insert(T x) {
     IREP(i, B) {
       if(x >> i & 1) {
         if(!p[i]) {
           p[i] = x, cnt++;
           change = true;
           return;
         } else x ^= p[i];
     if(!zero) zero = change = true;
```

```
T get min() {
  if(zero) return 0;
  REP(i, B) if(p[i]) return p[i];
T get_max() {
  T ans = 0:
  IREP(i, B) ans = max(ans, ans ^ p[i]);
  return ans:
T get_kth(long long k) {
  k++;
  if(k == 1 && zero) return 0;
  k -= zero:
  if(k >= (1LL << cnt)) return -1;
  update();
  T ans = 0:
  REP(i, SZ(d)) if(k \gg i \& 1) ans ^= d[i
  return ans:
bool contains(T x) {
  if(x == 0) return zero;
  IREP(i, B) if(x \gg i \& 1) \times ^= p[i];
  return x == 0:
void merge(const xor_basis& other) { REP(i
    , B) if(other.p[i]) insert(other.p[i])
    ; }
void update() {
  if(!change) return;
  change = false;
  d.clear();
  REP(j, B) IREP(i, j) if(p[j] \gg i \& 1) p
       [j] ^= p[i];
  REP(i, B) if(p[i]) d.pb(p[i]);
```

6.21 估計值

28

37

38

48

51

- Estimation
 - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for
 - $\begin{array}{l} n < 1e10,\,200000\, {\rm for}\, n < 1e19. \\ \ \, {\rm The}\, \, {\rm number}\, \, {\rm of}\, \, {\rm ways}\, \, {\rm of}\, \, {\rm writing}\, \, n \, \, {\rm as}\, \, {\rm a}\, \, {\rm sum}\, \, {\rm of} \end{array}$ positive integers, disregarding the order of the summands. 1, 1, 2, 3, 5, 7, 11, 15, 22, 30 for $n = 0 \sim 9,627 \text{ for } n = 20, \sim 2e5 \text{ for }$ $\underline{n} = 50, \sim 2e8 \text{ for } n = 100.$
 - Total number of partitions B(n)n distinct elements: n distinct elements: $B(n) = \underbrace{\sharp \varphi}_{1,1,2,5,15,52,203,877,4140,21147,115975,678570,42135977,B}$ 27644437, 190899322,

6.22 定理

· Cramer's rule

$$ax + by = e \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$cx + dy = f \Rightarrow y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

Burnside's Lemma

Let us calculate the number of necklaces of n pearls, where each pearl has m possible colors. Two necklaces are symmetric if they are similar after rotating them. There are n ways to change the position of a necklace, because we can rotate it $0, 1, \ldots, n_1$ steps clockwise. If the number of steps is 0, all m^n necklaces remain the same, and if the number of steps is 1, only the m necklaces where each pearl has the same color remain the same. More generally, when the number of steps is k, a total of $m^{\gcd(k,n)}$ necklaces remain the same. The reason for this is that blocks of pearls of size gcd(k, n)will replace each other. Thus, according to Burnside's lemma, the number of necklaces is $\sum_{i=0}^{n-1} \frac{m^{\gcd(i,n)}}{n}$. For example, the number of necklaces of length 4 with 3 colors is $\frac{3^4+3+3^2+3}{4} = 24$

Lindstr□m–Gessel–Viennot Lemma

 $\omega(P)$ 表示 P 這條路徑上所有邊的邊權之積。(路徑 計數時,可以將邊權都設為1)(事實上,邊權可以為 生成函數) e(u,v) 表示 u 到 v 的 ** 每一條 ** 路徑 P 的 $\omega(P)$ 之和 · 即 $e(u,v) = \sum \omega(P)$ 。 起點 集合 A,是有向無環圖點集的一個子集,大小為 n。 終點集合 B, 也是有向無環圖點集的一個子集, 大小

也為 $n \circ -$ 組 $A \to B$ 的不相交路徑 $S : S_i$ 是一條從 A_i 到 $B_{\sigma(S)_i}$ 的路徑 $(\sigma(S)$ 是一個排列) · 對於任 何 $i \neq j \cdot S_i$ 和 S_i 沒有公共頂點。 $t(\sigma)$ 表示排列 σ 的逆序對個數。

$$M = \begin{bmatrix} e(A_1, B_1) & e(A_1, B_2) & \cdots & e(A_1, B_n) \\ e(A_2, B_1) & e(A_2, B_2) & \cdots & e(A_2, B_n) \\ \vdots & \vdots & \ddots & \vdots \\ e(A_n, B_1) & e(A_n, B_2) & \cdots & e(A_n, B_n) \end{bmatrix}$$

$$\det(M) = \sum_{S:A \to B} (-1)^{t(\sigma(S))} \prod_{i=1}^n \omega(S_i)$$

表示滿足上文要求的 $A \rightarrow B$ 的每一組 不相交路徑 S。

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is
- $|\det(\tilde{L}_{11})|$. The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

· Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum match-

· Cayley's Formula

- Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there are
- $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees. Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \ldots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

• Erd□s-Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + \cdots + d_n$ is even

and
$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$$
 holds for every $1 \le k \le n$.

• Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \ge \cdots \ge a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^n a_i \leq \sum_{i=1}^n \min(b_i,k) \text{ holds}$ for every $1 \le k \le n$.

· Fulkerson-Chen-Anstee theorem

A sequence $(a_1, b_1), \ldots, (a_n, b_n)$ of nonnegative integer pairs with $a_1 \ge \cdots \ge a_n$ is digraphic if and only

$$\begin{aligned} & \text{if } \sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1) + \\ & \sum_{i=k+1}^n \min(b_i, k) \text{ holds for every } 1 \leq k \leq n. \end{aligned}$$

M□bius inversion formula

$$\begin{array}{llll} -f(n) &=& \sum_{d\mid n} g(d) & \Leftrightarrow & g(n) &=& \\ & \sum_{d\mid n} \mu(d) f(\frac{n}{d}) & \\ -f(n) &=& \sum_{n\mid d} g(d) & \Leftrightarrow & g(n) &=& \\ & \sum_{n\mid u} \frac{d}{d} f(d) & & & \end{array}$$

Spherical cap

- A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : arcsin(a/r).
- Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta) (1 \cos \theta)^2/3$. Area = $2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos \theta)^2/3$.

6.23 數字

· Bernoulli numbers

$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = \\ S(n,n) &= 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

• Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

· Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

· Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1), k+1$ j:s s.t. $\pi(j) \ge j, k$ j:s s.t.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

$$\sum_{k=1}^{n} k^3 = \left(\frac{n(n+1)}{2}\right)^2$$

$$\sum_{k=1}^{n} k^4 = \frac{1}{30} (6n^5 + 15n^4 + 10n^3 - n)$$

$$\sum_{k=1}^{n} k^5 = \frac{1}{12} (2n^6 + 6n^5 + 5n^4 - n^2)$$

$$\sum_{k=1}^{n} k^6 = \frac{1}{42} (6n^7 + 21n^6 + 21n^5 - 7n^3 + n)$$
General form:

$$\begin{array}{lll} \sum_{k=1}^{n} k^{p} & = & \frac{1}{p+1} (n \sum_{i=1}^{p} (n + 1)^{i} & - \\ \sum_{i=2}^{p} \binom{i}{p+1} \sum_{k=1}^{n} k^{p+1-i}) \end{array}$$

歐幾里得類算法

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity: $O(\log n)$

Misc

7.1 **fast**

```
f(a, b, c, n) = \sum_{i=1}^{n} \lfloor \frac{ai + b}{c} \rfloor
                                                                                                                      | #pragma GCC optimize("Ofast, no-stack-
                                                                                                                                     protector, unroll-loops, fast-math, inline
                                                                                                                      2 inline char gc() {
                        = \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) & \text{3} \\ +f(a \bmod c, b \bmod c, c, n), & a \ge c \lor b \\ 0, & n < 0 \lor 5 \end{cases}
= \begin{cases} \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) & \text{3} \\ +f(a \bmod c, b \bmod c, c, n), & a \ge c \lor b \\ n < 0 \lor 5 \end{cases}
= \begin{cases} \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) & \text{3} \\ +f(a \bmod c, b \bmod c, c, n), & a \ge c \lor b \\ -f(a \bmod c, b \bmod c, c, n), & a \ge c \lor b \end{cases}
                                                                                                                            static const size_t sz = 65536;
                                                                                                                              static char *p = buf, *end = buf;
                                                                                                                              if(p == end) end = buf + fread(buf, 1, sz,
                                                                                                                                          stdin), p = buf;
                                                                                                                              return *p++;
  g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                        = \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & \textbf{7.2} \ ^{0} \textbf{Yor} \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases}
                                                                                                                      | #define FOR(i, begin, end) for(int i = (
h(a, b, c, n) = \sum_{i=1}^{n} \lfloor \frac{ai + b}{c} \rfloor^{2}
                                                                                                                                      begin), i##_end_ = (end); i < i##_end_;
                                 \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \end{cases}
                                                                                                                      #define IFOR(i, begin, end) for(int i = (end
                                                                                                                                     ) - 1, i##_begin_ = (begin); i >= i##
                                                                                                                                      begin ; i--)
                                  +h(a \bmod c, b \bmod c, c, n)
                                                                                                                     3 #define REP(i, n) FOR(i, 0, n)
                         = \begin{cases} +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), \end{cases}
                                                                                                                     4 #define IREP(i, n) IFOR(i, 0, n)
                                                                                                                          a \ge c \lor b \ge c
                                                                                                                           n < 0 \lor a = 0
                                    nm(m+1) - 2g(c, c-b-1, a, m-1)
                                   nm(m+1) - 2g(c, c - c - 1, a, m-1) - f(a, b, c, n), 7therwise next-combination
```

6.25 生成函數

- Ordinary Generating Function $A(x) = \sum_{i \geq 0} a_i x^i$
- Exponential Generating Function A(x) $\sum_{i>0} \frac{a_i}{i!} x_i$

$$-A(x) + B(x) \Rightarrow a_n + b_n$$

$$-A^{(k)}(x) \Rightarrow a_{n \neq k_n}$$

$$-A(x)B(x) \Rightarrow \sum_{i=0}^{k_n} {n \choose i} a_i b_{n-i}$$

$$-A(x)^k \Rightarrow \sum_{i+1+i_2+\dots+i_k=n} {n \choose i_1,i_2,\dots,i_k} a_{i_1}$$

$$-xA(x) \Rightarrow na_n$$

· Special Generating Function

```
 - (1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i 
 - \frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{n}{i-1} x^i
```

```
1 // Example: 1 -> 2, 4 -> 8, 12(1100) ->
```

7.4 PBDS

```
i #include <ext/pb ds/assoc container.hpp>
using namespace __gnu_pbds;
tree<11, null_type, less<11>, rb_tree_tag,
   tree_order_statistics_node_update> st;
4 // find by order order of key
5 // __float128_t
6 for(int i = bs._Find_first(); i < bs.size();</pre>
       i = bs. Find next(i));
```

7.5 python

```
1 | from decimal import Decimal, getcontext
2 | getcontext().prec = 1000000000
3 | getcontext().Emax = 999999999
4 | a = pow(Decimal(2), 82589933) - 1
```

7.6 rng

```
inline ull rng() {
    static ull Q = 48763;
    Q ^= Q << 7;
    Q ^= Q >> 9;
    return Q & 0xfffffffffflll;
}
```

7.7 rotate90

7.8 timer

8 String

8.1 AC

```
int insert(const string& s) {
      int p = 0;
      for(char c : s) {
        int v = c - MIN_CHAR;
        if(node[p].go[v] == 0) node[p].go[v] =
              new node();
        p = node[p].go[v];
      node[p].cnt++;
      return p;
    void build() {
      que.reserve(SZ(node)); que.pb(0);
      REP(i, SZ(que)) {
        int u = que[i];
        REP(j, ALPHABET) {
          if(node[u].go[j] == 0) node[u].go[j]
                = node[node[u].fail].go[j];
            int v = node[u].go[j];
            node[v].fail = (u == 0 ? u : node[
                 node[u].fail].go[j]);
             que.pb(v);
33
```

8.2 KMP

```
1  // abacbaba -> [0, 0, 1, 0, 0, 1, 2, 3]
2  vi KMP(const vi& a) {
3    int n = SZ(a);
4    vi k(n);
5    for(int i = 1; i < n; ++i) {
6        int j = k[i - 1];
7        while(j > 0 && a[i] != a[j]) j = k[j - 1];
8        j += (a[i] == a[j]);
8        k[i] = j;
9        k[i] = j;
10    }
11    return k;
12 }
```

8.3 LCP

8.4 manacher

13 }

8.5 rolling-hash

8.6 SAIS

```
1 // mississippi
2 // 10 7 4 1 0 9 8 6 3 5 2
3 vi SAIS(string a) {
   #define QQ(i, n) for(int i = (n); i >= 0;
        i--)
   int n = sz(a), m = *max_element(all(a)) +
        1:
    vi pos(m + 1), x(m), sa(n), val(n), lms;
    for(auto c : a) pos[c + 1]++;
    REP(i, m) pos[i + 1] += pos[i];
    vector<bool> s(n);
    QQ(i, n - 2) s[i] = a[i] != a[i + 1] ? a[i]
        ] < a[i + 1] : s[i + 1];
    auto ind = [&](const vi& ls){
     fill(all(sa), -1);
      auto L = [\&](int i) \{ if(i >= 0 \&\& !s[i] \}
          ]) sa[x[a[i]]++] = i; };
```

```
auto S = [\&](int i) \{ if(i >= 0 \&\& s[i]) \}
            sa[--x[a[i]]] = i; };
      REP(i, m) x[i] = pos[i + 1];
      QQ(i, sz(ls) - 1) S(ls[i]);
      REP(i, m) x[i] = pos[i];
      L(n - 1):
19
      REP(i, n) L(sa[i] - 1);
      REP(i, m) x[i] = pos[i + 1];
      QQ(i, n - 1) S(sa[i] - 1);
22
23
    auto ok = [&](int i) { return i == n || (!)
         s[i - 1] && s[i]); };
    auto same = [&](int i,int j) {
25
        if(a[i++] != a[j++]) return false;
26
      } while(!ok(i) && !ok(j));
      return ok(i) && ok(j);
28
29
    for(int i = 1; i < n; i++) if(ok(i)) lms.
         pb(i);
    ind(lms);
    if(sz(lms)) {
      int p = -1, w = 0;
      for(auto v : sa) if(v && ok(v)) {
        if(p != -1 \&\& same(p, v)) w--;
        val[p = v] = w++;
      auto b = lms;
      for(auto& v : b) v = val[v];
      b = SAIS(b);
      for(auto& v : b) v = lms[v];
41
42
      ind(b);
43
    return sa;
```

8.7 SAM

```
ı | // cnt 要先用 bfs 往回推, 第一次出現的位置是
        state.first pos - |S| + 1
2 struct Node { int go[26], len, link, cnt,
       first pos; };
 3 Node SA[N]; int sz;
  void sa_init() { SA[0].link = -1, SA[0].len
       = 0, sz = 1; 
  int sa_extend(int p, int c) {
    int u = sz++;
    SA[u].first_pos = SA[u].len = SA[p].len +
    SA[u].cnt = 1;
    while(p != -1 && SA[p].go[c] == 0) {
      SA[p].go[c] = u;
10
11
      p = SA[p].link;
12
13
    if(p == -1) {
14
      SA[u].link = 0;
15
      return u;
16
    int q = SA[p].go[c];
18
    if(SA[p].len + 1 == SA[q].len) {
      SA[u].link = q;
20
      return u;
```

```
22    int x = sz++;
23    SA[x] = SA[q];
24    SA[x].cnt = 0;
25    SA[x].len = SA[p].len + 1;
26    SA[q].link = SA[u].link = x;
27    while(p != -1 && SA[p].go[c] == q) {
28     SA[p].go[c] = x;
29    p = SA[p].link;
30    }
31    return u;
32 }
```

8.8 smallest-rotation

8.9 Z

```
1  // abacbaba -> [0, 0, 1, 0, 0, 3, 0, 1]
2  vi z_algorithm(const vi& a) {
3    int n = sz(a);
4    vi z(n);
5    for(int i = 1, j = 0; i < n; ++i) {
6        if(i <= j + z[j]) z[i] = min(z[i - j], j + z[j] - i);
7    while(i + z[i] < n && a[i + z[i]] == a[z [i]]) z[i]++;
8    if(i + z[i] > j + z[j]) j = i;
9   }
10   return z;
11 }
```

ACM ICPC Judge Test NTHU LinkCutTreap

C++ Resource Test

```
#include <bits/stdc++.h>
using namespace std;

namespace system_test {
   const size_t KB = 1024;
   const size_t MB = KB * 1024;
   const size_t GB = MB * 1024;
}
```

```
chrono::duration<double> diff = end -
10 size t block size, bound;
                                                          begin;
  void stack size dfs(size t depth = 1) {
                                                     return diff.count():
   if (depth >= bound)
                                                   void runtime_error_1() {
    int8_t ptr[block_size]; // 若無法編譯將
                                                     // Segmentation fault
         block size 改成常數
                                                     int *ptr = nullptr;
    memset(ptr, 'a', block_size);
                                                     *(ptr + 7122) = 7122;
    cout << depth << endl;</pre>
                                                 42 }
    stack_size_dfs(depth + 1);
                                                   void runtime_error_2() {
                                                     // Segmentation fault
  void stack_size_and_runtime_error(size_t
                                                     int *ptr = (int *)memset;
       block size, size t bound = 1024) {
                                                     *ptr = 7122;
    system test::block size = block size;
                                                 48 }
    system_test::bound = bound;
    stack size dfs();
                                                   void runtime_error_3() {
                                                     // munmap_chunk(): invalid pointer
                                                     int *ptr = (int *)memset;
  double speed(int iter num) {
                                                     delete ptr;
    const int block_size = 1024;
                                                 54
    volatile int A[block_size];
    auto begin = chrono::high resolution clock
                                                   void runtime_error_4() {
         ::now();
                                                     // free(): invalid pointer
    while (iter num--)
                                                     int *ptr = new int[7122];
      for (int j = 0; j < block_size; ++j)</pre>
                                                     ptr += 1;
                                                     delete[] ptr;
    auto end = chrono::high resolution clock::
```

```
63 void runtime error 5() {
    // maybe illegal instruction
    int a = 7122, b = 0;
    cout << (a / b) << endl;</pre>
  void runtime error 6() {
    // floating point exception
    volatile int a = 7122, b = 0;
    cout << (a / b) << endl;
73 }
  void runtime error 7() {
    // call to abort.
    assert(false);
78 }
80 } // namespace system test
82 #include <sys/resource.h>
void print_stack_limit() { // only work in
       Linux
    struct rlimit 1;
    getrlimit(RLIMIT STACK, &1);
    cout << "stack_size = " << l.rlim_cur << "</pre>
          byte" << endl;</pre>
87 }
```