ACM ICPC Team Reference -		2.8 sparse-table	4 4 4	5	Graph 5.1 2-SAT	9 9 9		6.18 Xor-Basis	13 13
NTHU	3	2.11 treap	5 5		5.3 HLD	9 10 10		6.21 數字	14
SplayTreap		3.1 bipartite-matching	5 5 6	6	6.1 Aliens	10	7	Misc 7.1 fast	14
Contents		3.5 general-matching	6 6 7		6.3 Chinese-Remainder	11 11 11		7.4 python	14 14 14
1 Basic 1.1 vimrc	2 2 2	3.8 max-clique	8 8		6.7 extgcd	11 11 12	8	String 8.1 AC	14 14 14
2.2 DLX	2 4 2 2 3 3 3 3	Geometry 4.1 closest-pair 4.2 convex-hull 4.3 point-in-convex-hull 4.4 point 4.5 polar-angle-sort	8 8 9 9 9		6.12 Linear-Sieve	12 12 12 12 12		8.3 LCP 8.4 manacher 8.5 rolling-hash 8.6 SAIS 8.7 SAM 8.8 smallest-rotation 8.9 Z	14 15 15 15 15

Basic

1.1 vimre

```
| se nu ai hls et ru ic is sc cul
se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
 svntax on
4 hi cursorline cterm=none ctermbg=89
set bg=dark
6 inoremap {<CR> {<CR>}<Esc>ko<tab>
```

Data-Structure

2.1 CHT

```
|| struct line t {
    mutable ll k, m, p;
    bool operator<(const line_t& o) const {</pre>
         return k < o.k; }</pre>
    bool operator<(11 x) const { return p < x;</pre>
 };
6 template < bool MAX>
  struct CHT : multiset<line t, less<>> {
    const ll INF = 1e18L:
    bool isect(iterator x, iterator y) {
      if(y == end()) return x \rightarrow p = INF, 0;
      if(x->k == y->k) {
        x->p = (x->m > y->m ? INF : -INF);
      } else {
         x \rightarrow p = floor div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y)
              ->k); // see Math
      return x->p >= y->p;
    void add_line(ll k, ll m) {
      if(!MAX) k = -k, m = -m;
      auto z = insert(\{k, m, 0\}), y = z++, x =
      while(isect(y, z)) z = erase(z);
      if(x != begin() && isect(--x, y)) isect(
           x, y = erase(y));
      while((y = x) != begin() && (--x)->p >=
           y->p) isect(x, erase(y));
    11 get(11 x) {
      assert(!empty());
      auto 1 = *lower_bound(x);
      return (1.k * x + 1.m) * (MAX ? +1 : -1)
30 };
```

```
| struct DLX {
    int n, m, tot, ans;
    vi first, siz, L, R, U, D, col, row, stk;
    DLX(int _n, int _m) : n(_n), m(_m), tot(_m
      int sz = n * m:
      first = siz = L = R = U = D = col = row
           = stk = vi(sz):
      REP(i, m + 1) {
        L[i] = i - 1, R[i] = i + 1;
        U[i] = D[i] = i;
      L[0] = m, R[m] = 0;
    void insert(int r, int c) {
      r++, c++;
      col[++tot] = c, row[tot] = r, ++siz[c];
      D[tot] = D[c], U[D[c]] = tot, U[tot] = c
           , D[c] = tot;
      if(!first[r]) first[r] = L[tot] = R[tot]
            = tot:
      else {
        L[R[tot] = R[first[r]]] = tot;
        R[L[tot] = first[r]] = tot;
    #define TRAV(i, X, j) for(i = X[j]; i != j 17
         ; i = X[i]
    void remove(int c) {
      int i, j;
      L[R[c]] = L[c], R[L[c]] = R[c];
      TRAV(i, D, c) TRAV(j, R, i) {
        D[U[D[j]] = U[j]] = D[j];
        siz[col[j]]--;
    void recover(int c) {
      int i, j;
      TRAV(i, U, c) TRAV(j, L, i) {
        U[D[j]] = D[U[j]] = j;
        siz[col[j]]++;
      L[R[c]] = R[L[c]] = c;
    bool dance(int dep) {
      if(!R[0]) return ans = dep, true;
      int i, j, c = R[0];
      TRAV(i, R, 0) if (siz[i] < siz[c]) c = i;
      remove(c);
      TRAV(i, D, c) {
        stk[dep] = row[i];
        TRAV(j, R, i) remove(col[j]);
        if(dance(dep + 1)) return true;
        TRAV(j, L, i) recover(col[j]);
      recover(c);
      return false;
    vi solve() {
      if(!dance(1)) return {};
      return vi(stk.begin() + 1, stk.begin() +
58 };
```

2.3 fast-set

39

40

```
1 // Can correctly work with numbers in range
       ΓΘ: MAXN1
2 // Supports all std::set operations in O(1)
      on random queries / dense arrays. O(
       log 64(N)) in worst case (sparce array). 51
3 // Count operation works in O(1) always.
4 template < uint MAXN>
5 class fast set {
  private:
   static const uint PREF = (MAXN <= 64 ? 0 :</pre>
                  MAXN <= 4096 ? 1 :
                  MAXN <= 262144 ? 1 + 64 :
                  MAXN <= 16777216 ? 1 + 64 +
                       4096 :
                  MAXN <= 1073741824 ? 1 + 64
                       + 4096 + 262144 : 227) +
    static constexpr ull lb(int x) {
      if(x == 64) return ULLONG_MAX;
      return (1ULL << x) - 1;
    static const uint SZ = PREF + (MAXN + 63)
         / 64 + 1;
    ull m[SZ] = \{0\};
    inline uint left(uint v) const { return (v
          - 62) * 64; }
    inline uint parent(uint v) const { return
         v / 64 + 62; }
    inline void setbit(uint v) { m[v >> 6] |=
         1ULL << (v & 63); }
   inline void resetbit(uint v) { m[v >> 6]
         &= ~(1ULL << (v & 63)); }
    inline uint getbit(uint v) const { return
         m[v >> 6] >> (v \& 63) \& 1; }
    inline ull childs_value(uint v) const {
         return m[left(v) >> 6]; }
    inline int left_go(uint x, const uint c)
         const {
      const ull rem = x \& 63;
      uint bt = PREF * 64 + x;
      ull num = m[bt >> 6] & lb(rem + c);
      if(num) return (x ^ rem) | lg(num);
      for(bt = parent(bt); bt > 62; bt =
           parent(bt)) {
        const ull rem = bt & 63:
        num = m[bt >> 6] & lb(rem);
        if(num) {
          bt = (bt ^ rem) | __lg(num);
          break;
      if(bt == 62) return -1;
      while(bt < PREF * 64) bt = left(bt) |</pre>
           __lg(m[bt - 62]);
      return bt - PREF * 64;
    inline int right_go(uint x, const uint c)
         const {
      const ull rem = x \& 63:
      uint bt = PREF * 64 + x;
      ull num = m[bt >> 6] \& \sim lb(rem + c);
      if(num) return (x ^ rem) |
           __builtin_ctzll(num);
```

```
for(bt = parent(bt); bt > 62; bt =
            parent(bt)) {
         const ull rem = bt & 63;
         num = m[bt >> 6] \& \sim lb(rem + 1);
           bt = (bt ^ rem) | builtin ctzll(
               num);
           break:
53
       if(bt == 62) return -1:
       while(bt < PREF * 64) bt = left(bt) |</pre>
            __builtin_ctzll(m[bt - 62]);
       return bt - PREF * 64;
59 public:
    fast_set() { assert(PREF != 228); setbit
     bool empty() const {return getbit(63);}
    void clear() { fill(m, m + SZ, 0); setbit
          (62); }
    bool count(uint x) const { return m[PREF +
          (x >> 6)] >> (x \& 63) \& 1; }
    void insert(uint x) { for(uint v = PREF *
          64 + x; !getbit(v); v = parent(v))
          setbit(v); }
     void erase(uint x) {
       if(!getbit(PREF * 64 + x)) return;
       resetbit(PREF * 64 + x);
67
       for(uint v = parent(PREF * 64 + x); v >
            62 && !childs value(v); v = parent(v
           )) resetbit(v);
    int find next(uint x) const { return
          right go(x, 0); } // >=
    int find_prev(uint x) const { return
          left_go(x, 1); } // <=
```

2.4 lazvsegtree

```
i template < class S.</pre>
            S (*e)(),
            S (*op)(S, S),
            class F,
            F (*id)(),
            S (*mapping)(F, S),
            F (*composition)(F, F)>
  struct lazy segtree {
    int n, size, log;
    vector<S> d; vector<F> lz;
    void update(int k) { d[k] = op(d[k << 1],</pre>
          d[k << 1 | 1]); }
    void all_apply(int k, F f) {
      d[k] = mapping(f, d[k]);
14
      if(k < size) lz[k] = composition(f, lz[k</pre>
           ]);
15
    void push(int k) {
16
      all apply(k << 1, lz[k]);
17
18
       all apply(k \ll 1 \mid 1, lz[k]);
19
      lz[k] = id();
```

```
for(int i = 1; i <= log; i++) {</pre>
lazy segtree(int n) : lazy segtree(vector
                                                      if(((1 >> i) << i) != 1) update(1 >> i
     <S>( n, e())) {}
lazy_segtree(const vector<S>& v) : n(sz(v) 83
                                                      if(((r >> i) << i) != r) update((r -</pre>
                                                           1) \gg i);
  log = lg(2 * n - 1), size = 1 << log;
  d.resize(size * 2, e());
  lz.resize(size, id());
                                                  template < class G > int max right(int 1, G g
  REP(i, n) d[size + i] = v[i];
                                                    assert(0 <= 1 && 1 <= n && g(e()));
  for(int i = size - 1; i; i--) update(i); 87
                                                    if(1 == n) return n;
void set(int p, S x) {
                                                    1 += size;
                                                                                               13
                                                    for(int i = log; i; i--) push(l >> i);
  p += size;
  for(int i = log; i; --i) push(p >> i);
                                                    S sm = e();
  for(int i = 1; i <= log; ++i) update(p</pre>
                                                      while(!(1 & 1)) 1 >>= 1;
                                                      if(!g(op(sm, d[1]))) {
       >> i):
                                                        while(1 < size) {</pre>
S get(int p) {
                                                          push(1):
                                                          1 <<= 1;
  p += size;
  for(int i = log; i; i--) push(p >> i);
                                                          if(g(op(sm, d[1]))) sm = op(sm, d[
                                                                                               21
  return d[p];
                                                               1++]);
                                                                                               23
S prod(int 1, int r) {
                                                        return 1 - size:
                                                                                               24
  if(1 == r) return e();
                                                                                               25
  1 += size; r += size;
                                                      sm = op(sm, d[1++]);
  for(int i = log; i; i--) {
                                                    } while((1 & -1) != 1);
    if(((1 >> i) << i) != 1) push(1 >> i); 104
                                                    return n:
    if(((r >> i) << i) != r) push(r >> i); 105
                                                  template < class G> int min_left(int r, G g)
  S sml = e(), smr = e();
                                                    assert(0 <= r && r <= n && g(e()));
  while(1 < r) {
    if(1 \& 1) sml = op(sml, d[1++]);
                                                    if(r == 0) return 0;
    if(r \& 1) smr = op(d[--r], smr);
                                                    r += size;
    1 >>= 1:
                                                    for(int i = log; i >= 1; i--) push((r -
    r >>= 1;
                                                         1) >> i);
                                                    S sm = e();
  return op(sml, smr);
                                                    do {
                                                                                               35
S all_prod() const { return d[1]; }
                                                      while(r > 1 && (r & 1)) r >>= 1;
void apply(int p, F f) {
                                                      if(!g(op(d[r], sm))) {
  p += size;
                                                        while(r < size) {</pre>
                                                                                                38
  for(int i = log; i; i--) push(p >> i);
                                                          push(r);
  d[p] = mapping(f, d[p]);
                                                          r = r << 1 | 1:
  for(int i = 1; i <= log; i++) update(p</pre>
                                                          if(g(op(d[r], sm))) sm = op(d[r])
       >> i);
                                                               --], sm);
void apply(int 1, int r, F f) {
                                             121
                                                        return r + 1 - size;
  if(l == r) return;
                                             122
  1 += size; r += size;
                                                      sm = op(d[r], sm);
  for(int i = log; i; i--) {
                                                    } while((r & -r) != r);
    if(((1 >> i) << i) != 1) push(1 >> i); 125
                                                    return 0;
    if(((r >> i) << i) != r) push((r - 1)
         >> i):
    int 12 = 1, r2 = r;
                                                2.5 rect-add-rect-sum
    while(1 < r)  {
      if(1 & 1) all_apply(1++, f);
      if(r & 1) all apply(--r, f);
      1 >>= 1;
                                              1 template < class Int, class T>
                                              2 struct RectangleAddRectangleSum {
      r >>= 1;
                                                  struct AQ { Int xl, xr, yl, yr; T val; };
    1 = 12;
                                                  struct SQ { Int xl, xr, yl, yr; };
    r = r2:
                                                  vector<AQ> add qry;
                                                  vector<SQ> sum_qry;
```

```
7 // A[x][y] += val for(x, y) in [xl, xr) *
        [vl. vr)
   void add rectangle(Int xl, Int xr, Int yl,
                                               58 };
         Int yr, T val) { add_qry.pb({xl, xr,
        yl, yr, val}); }
   // Get sum of A[x][y] for (x, y) in [xl, xr]
        ) * [vl. vr)
   void add query(Int xl, Int xr, Int yl, Int
         yr) { sum_qry.pb({xl, xr, yl, yr}); }
   vector<T> solve() {
     vector<Int> ys;
      for(auto &a : add_qry) ys.pb(a.yl), ys.
          pb(a.yr);
     ys = sort unique(ys);
     const int Y = SZ(ys);
     vector<tuple<Int, int, int>> ops;
     REP(q, SZ(sum qry)) {
       ops.eb(sum_qry[q].xl, 0, q);
       ops.eb(sum_qry[q].xr, 1, q);
                                                11
     REP(q, SZ(add_qry)) {
       ops.eb(add_qry[q].xl, 2, q);
       ops.eb(add_qry[q].xr, 3, q);
     sort(ALL(ops));
     fenwick<T> b00(Y), b01(Y), b10(Y), b11(Y
     vector<T> ret(SZ(sum_qry));
                                                19
     for(auto o : ops) {
       int qtype = get<1>(o), q = get<2>(o);
       if(qtype >= 2) {
          const auto& query = add_qry[q];
          int i = lower bound(ALL(ys), query.
              yl) - ys.begin();
          int j = lower bound(ALL(ys), query.
              yr) - ys.begin();
          T x = get<0>(o);
          T yi = query.yl, yj = query.yr;
                                                28
          if(qtype & 1) swap(i, j), swap(yi,
                                                29
          b00.add(i, x * yi * query.val);
          b01.add(i, -x * query.val);
          b10.add(i, -yi * query.val);
          b11.add(i, query.val);
          b00.add(j, -x * yj * query.val);
          b01.add(j, x * query.val);
          b10.add(j, yj * query.val);
          b11.add(j, -query.val);
       } else {
          const auto& query = sum_qry[q];
          int i = lower_bound(ALL(ys), query.
              yl) - ys.begin();
          int j = lower bound(ALL(ys), query.
              yr) - ys.begin();
          T x = get<0>(o);
          T yi = query.yl, yj = query.yr;
          if(qtype & 1) swap(i, j), swap(yi,
          ret[q] += b00.get(i - 1) + b01.get(i
               -1) * vi + b10.get(i - 1) * x
                                                -11
               + b11.get(i - 1) * x * yi;
          ret[q] = b00.get(j - 1) + b01.get(j
                                                13
               -1) * yj + b10.get(j - 1) * x
                                                14
              + b11.get(j - 1) * x * yj;
                                                15
```

```
2.6 rollback-dsu
```

return ret;

```
| struct RollbackDSU {
    int n; vi sz, tag;
    vector<tuple<int, int, int, int>> op;
    void init(int _n) {
      n = _n;
      sz.assign(n, -1);
      tag.clear();
    int leader(int x) {
      while(sz[x] >= 0) x = sz[x];
      return x:
12
    bool merge(int x, int y) {
      x = leader(x), y = leader(y);
      if(x == y) return false;
      if(-sz[x] < -sz[y]) swap(x, y);
      op.eb(x, sz[x], y, sz[y]);
      sz[x] += sz[y]; sz[y] = x;
      return true:
    int size(int x) { return -sz[leader(x);] }
    void add tag() { tag.pb(sz(op)); }
23
    void rollback() {
      int z = tag.back(); tag.ppb();
      while(sz(op) > z) {
        auto [x, sx, y, sy] = op.back(); op.
             ppb();
        sz[x] = sx:
        sz[y] = sy;
30
31 };
```

2.7 segtree

```
i template < class S, S (*e)(), S (*op)(S, S)>
2 struct segtree {
   int n, size, log;
   vector<S> st:
    void update(int v) { st[v] = op(st[v <<</pre>
         1], st[v << 1 | 1]); }
    segtree(int n) : segtree(vector<S>( n, e
         ())) {}
    segtree(const vector<S>& a): n(sz(a)) {
      log = __lg(2 * n - 1), size = 1 << log;
      st.resize(size << 1, e());
      REP(i, n) st[size + i] = a[i];
     for(int i = size - 1; i; i--) update(i);
   void set(int p, S val) {
      st[p += size] = val;
      for(int i = 1; i <= log; ++i) update(p</pre>
          >> i);
```

```
S get(int p) const {
  return st[p + size];
S prod(int 1, int r) const {
  assert(0 \le 1 \&\& 1 \le r \&\& r \le n);
  S sml = e(), smr = e();
  1 += size, r += size;
  while(1 < r) {
    if(1 \& 1) sml = op(sml, st[1++]);
    if(r & 1) smr = op(st[--r], smr);
    1 >>= 1:
    r >>= 1;
  return op(sml, smr);
S all prod() const { return st[1]; }
template < class F> int max right(int 1, F f
    ) const {
  assert(0 <= 1 && 1 <= n && f(e()));
  if(1 == n) return n;
  1 += size;
  S sm = e();
    while(~1 & 1) 1 >>= 1;
    if(!f(op(sm, st[1]))) {
      while(1 < size) {</pre>
        1 <<= 1;
        if(f(op(sm, st[1]))) sm = op(sm,
             st[l++]);
      return 1 - size;
    sm = op(sm, st[1++]);
  } while((1 & -1) != 1);
  return n:
template < class F> int min_left(int r, F f)
  assert(0 <= r \& r <= n \& f(e()));
  if(r == 0) return 0;
  r += size;
  S sm = e();
  do {
    while(r > 1 && (r & 1)) r >>= 1;
    if(!f(op(st[r], sm))) {
      while(r < size) {</pre>
        r = r << 1 | 1;
        if(f(op(st[r], sm))) sm = op(st[r])
             --], sm);
      return r + 1 - size;
    sm = op(st[r], sm);
  } while((r & -r) != r);
  return 0;
```

sparse-table

```
1 template < class T, T (*op)(T, T)>
2 struct sparse_table {
```

```
int n;
    vector<vector<T>> b;
    sparse table(const vector<T>& a) : n(sz(a) 41
      int lg = __lg(n) + 1;
      b.resize(lg); b[0] = a;
       for(int j = 1; j < lg; ++j) {</pre>
        b[j].resize(n - (1 << j) + 1);
        REP(i, n - (1 << j) + 1) b[j][i] = op( 46
             b[j - 1][i], b[j - 1][i + (1 << (j 47)]
              - 1))]);
11
12
    T prod(int from, int to) {
      int lg = __lg(to - from + 1);
      return op(b[lg][from], b[lg][to - (1 <<
           lg) + 1]);
17 };
```

static-range-inversion

```
i| struct static_range_inversion {
   int sz:
   vi a, L, R;
   vector<ll> ans;
   static_range_inversion(vi _a) : a(_a) {
      a = sort unique( a);
     REP(i, SZ(a)) a[i] = lower bound(ALL(a)
          , a[i]) - _a.begin();
     sz = SZ(a);
   void add_query(int 1, int r) { L.push_back
        (1), R.push back(r); }
   vector<ll> solve() {
     const int q = SZ(L);
     const int B = max(1.0, SZ(a) / sqrt(q));
     vi ord(q);
      iota(ALL(ord), 0);
      sort(ALL(ord), [&](int i, int j) {
       if(L[i] / B == L[j] / B) {
         return L[i] / B & 1 ? R[i] > R[j] :
              R[i] < R[j];
       return L[i] < L[j];</pre>
     });
      ans.resize(q);
      fenwick<ll> fenw(sz + 1);
     11 cnt = 0;
      auto AL = [&](int i) {
       cnt += fenw.sum(0, a[i] - 1);
       fenw.add(a[i], +1);
     };
      auto AR = [&](int i) {
       cnt += fenw.sum(a[i] + 1, sz);
       fenw.add(a[i], +1);
      auto DL = [&](int i) {
       cnt -= fenw.sum(0, a[i] - 1);
       fenw.add(a[i], -1);
     };
      auto DR = [&](int i) {
       cnt -= fenw.sum(a[i] + 1, sz);
```

2.10 static-range-lis

ans[id] = cnt;

return ans;

50

51

25

26

27

30

33

52 };

fenw.add(a[i], -1);

for(int i = 0; i < q; i++) {</pre>

while(1 > q1) AL(--1);

while(r < qr) AR(r++);

while(1 < q1) DL(1++);</pre>

while(r > qr) DR(--r);

int id = ord[i], ql = L[id], qr = R[id

52

int 1 = 0, r = 0:

```
1 struct static_range_lis {
   int n;
   vector<vector<pii>>> qry;
   vi invp, res_monge, ans;
   static range lis(vi a) : n(SZ(a)), qry(n +
         1), invp(n), res_monge(n) {
      // a must be permutation of [0, n)
     REP(i, n) invp[a[i]] = i;
   void add query(int 1, int r) { gry[1].pb({
        r, SZ(ans)}), ans.pb(r - 1); }
   void unit monge mult(vi& a, vi& b, vi& r)
      int n = SZ(a);
     if(n == 2){
       if(!a[0] && !b[0]) r[0] = 0, r[1] = 1;
       else r[0] = 1, r[1] = 0;
       return;
     if(n == 1) return r[0] = 0, void();
     int d = n / 2;
     vi a1(d), a2(n - d), b1(d), b2(n - d);
     vi mpa1(d), mpa2(n - d), mpb1(d), mpb2(n
           - d);
     int p[2] = {};
     REP(i, n) {
       if(a[i] < d) a1[p[0]] = a[i], mpa1[p</pre>
             [0]++] = i;
       else a2[p[1]] = a[i] - d, mpa2[p[1]++]
             = i;
     p[0] = p[1] = 0;
     REP(i, n) {
       if(b[i] < d) b1[p[0]] = b[i], mpb1[p]
                                                 82
             [0]++] = i;
       else b2[p[1]] = b[i] - d, mpb2[p[1]++]
             = i;
     vi c1(d), c2(n - d);
     unit_monge_mult(a1, b1, c1),
                                                 87
          unit_monge_mult(a2, b2, c2);
                                                88
     vi cpx(n), cpy(n), cqx(n), cqy(n);
     REP(i, d) cpx[mpa1[i]] = mpb1[c1[i]],
          cpy[mpa1[i]] = 0;
```

```
REP(i, n - d) cpx[mpa2[i]] = mpb2[c2[i]]
      ]], cpy[mpa2[i]] = 1;
  REP(i, n) r[i] = cpx[i];
  REP(i, n) cqx[cpx[i]] = i, cqy[cpx[i]] =
       cpy[i];
  int hi = n, lo = n, his = 0, los = 0;
  REP(i, n) {
   if(cqy[i] ^ (cqx[i] >= hi)) his--;
    while(hi > 0 && his < 0) {</pre>
     hi--:
      if(cpy[hi] ^ (cpx[hi] > i)) his++;
    while(lo > 0 && los <= 0) {
      if(cpy[lo] ^ (cpx[lo] >= i)) los++;
   if(los > 0 \&\& hi == lo) r[lo] = i;
   if(cqy[i] ^ (cqx[i] >= lo)) los--;
void subunit_monge_mult(vi& a, vi& b, vi&
    c) {
  int n = SZ(a);
  vi za(n), zb(n), res(n), vis(n), mpa(n,
       -1), mpb(n, -1), rb(n, -1);
  int ca = n;
  IREP(i, n) if(a[i] != -1) vis[a[i]] = 1,
       za[--ca] = a[i], mpa[ca] = i;
  IREP(i, n) if(!vis[i]) za[--ca] = i;
  fill(ALL(vis), -1);
  REP(i, n) if(b[i] != -1) vis[b[i]] = i;
  REP(i, n) if(vis[i] != -1) mpb[ca] = i,
       rb[vis[i]] = ca++;
  REP(i, n) if(rb[i] == -1) rb[i] = ca++;
  REP(i, n) zb[rb[i]] = i;
  unit_monge_mult(za, zb, res);
  fill(ALL(c), -1);
  REP(i, n) if(mpa[i] != -1 && mpb[res[i]]
       != -1) c[mpa[i]] = mpb[res[i]];
void solve(vi& p, vi& ret) {
 int n = SZ(p);
  if(n == 1) return ret[0] = -1, void();
 int d = n / 2;
  vi pl(d), pr(n - d);
  REP(i, d) pl[i] = p[i];
  REP(i, n - d) pr[i] = p[i + d];
  vi vis(n, -1);
  REP(i, d) vis[pl[i]] = i;
  vi tl(d), tr(n - d), mpl(d), mpr(n - d);
  int ca = 0;
  REP(i, n) if(vis[i] != -1) mpl[ca] = i,
      tl[vis[i]] = ca++;
  ca = 0;
  fill(ALL(vis), -1);
  REP(i, n - d) vis[pr[i]] = i;
  REP(i, n) if(vis[i] != -1) mpr[ca] = i,
      tr[vis[i]] = ca++;
  vi vl(d), vr(n - d);
  solve(t1, v1), solve(tr, vr);
  vi sl(n), sr(n);
  iota(ALL(sl), 0), iota(ALL(sr), 0);
  REP(i, d) sl[mpl[i]] = (vl[i] == -1 ? -1
       : mpl[vl[i]]);
```

```
REP(i, n - d) sr[mpr[i]] = (vr[i] == -1)
                                                           auto p = split(v->r, k - size(v->l) - 1) 38
            ? -1 : mpr[vr[i]]);
       subunit monge mult(sl, sr, ret);
                                                          if(p.second) p.second->p = NULL;
                                                          v->r = p.first;
     vi solve() {
                                                          pull(v); return {v, p.second};
       solve(invp, res monge);
       vi fenw(n + 1);
       IREP(i, n) {
                                                      int get position(Node* v) { // 0-indexed
         if(res monge[i] != -1) {
                                                        int k = (v->1 != NULL ? v->1->sz : 0);
           for(int p = res_monge[i] + 1; p <= n</pre>
                                                        while(v->p != NULL) {
                ; p += p & -p) fenw[p]++;
                                                          if(v == v \rightarrow p \rightarrow r) {
         for(auto& z : qry[i]){
                                                            if(v->p->l != NULL) k += v->p->l->sz;
           auto [id, c] = z;
           for(int p = id; p; p -= p & -p) ans[
                                                          v = v - p;
                c] -= fenw[p];
103
                                                        return k;
105
       return ans:
106
107 };
```

2.12 union-of-rectangles

// 1 10 1 10

// 0 2 0 2

// ans = 84

2.11 treap

```
1 struct Node {
    bool rev = false;
    int sz = 1, pri = rng();
    Node *1 = NULL, *r = NULL, *p = NULL;
    // TODO
  void pull(Node*& v) {
    v \rightarrow sz = 1 + size(v \rightarrow l) + size(v \rightarrow r);
    // TODO
void push(Node*& v) {
    if(v->rev) {
       swap(v->1, v->r);
       if(v->1) v->1->rev ^= 1;
      if(v->r) v->r->rev ^= 1;
      v->rev = false;
  Node* merge(Node* a, Node* b) {
    if(!a | | !b) return (a ? a : b);
    push(a), push(b);
    if(a->pri > b->pri) {
       a \rightarrow r = merge(a \rightarrow r, b);
      pull(a); return a;
       b \rightarrow 1 = merge(a, b \rightarrow 1);
      pull(b); return b;
30 pair<Node*, Node*> split(Node* v, int k) {
    if(!v) return {NULL, NULL};
    push(v);
    if(size(v->1) >= k) {
       auto p = split(v \rightarrow l, k);
      if(p.first) p.first->p = NULL;
      v \rightarrow 1 = p.second:
       pull(v); return {p.first, v};
    } else {
```

```
vector<int> vx, vy;
struct q { int piv, s, e, x; };
struct tree {
  vector<int> seg, tag;
  tree(int _n) : seg(_n * 16), tag(_n * 16)
  void add(int ql, int qr, int x, int v, int
       1, int r) {
    if(qr <= 1 || r <= q1) return;
    if(ql <= 1 && r <= qr) {
      tag[v] += x;
      if(tag[v] == 0) {
        if(1 != r) seg[v] = seg[2 * v] + seg
             [2 * v + 1];
        else seg[v] = 0;
      } else seg[v] = vx[r] - vx[1];
    } else {
      int mid = (1 + r) / 2;
      add(ql, qr, x, 2 * v, 1, mid);
      add(ql, qr, x, 2 * v + 1, mid, r);
      if(tag[v] == 0 && 1 != r) seg[v] = seg
           [2 * v] + seg[2 * v + 1];
 int q() { return seg[1]; }
int main() {
  int n: cin >> n:
  vector<int> x1(n), x2(n), y_(n), y2(n);
  for (int i = 0; i < n; i++) {</pre>
    cin >> x1[i] >> x2[i] >> y_[i] >> y2[i];
         //LRDU
    vx.pb(x1[i]), vx.pb(x2[i]);
    vy.pb(y_[i]), vy.pb(y2[i]);
  vx = sort unique(vx);
  vy = sort_unique(vy);
  vector<q> a(2 * n);
```

```
REP(i, n) {
      x1[i] = lower_bound(ALL(vx), x1[i]) - vx 20
      x2[i] = lower_bound(ALL(vx), x2[i]) - vx 21
            .begin();
      y_{i} = lower_bound(ALL(vy), y_{i}) - vy
            .begin();
      y2[i] = lower_bound(ALL(vy), y2[i]) - vy
            .begin();
      a[2 * i] = {y_[i], x1[i], x2[i], +1};
      a[2 * i + 1] = \{y2[i], x1[i], x2[i],
                                                   28
           -1};
                                                   29
45
    sort(ALL(a), [](q a, q b) { return a.piv < 31</pre>
           b.piv; });
    tree seg(n);
                                                   33
    11 \text{ ans} = 0:
                                                   34
    REP(i, 2 * n) {
      int j = i;
      while(j < 2 * n && a[i].piv == a[j].piv)</pre>
        seg.add(a[j].s, a[j].e, a[j].x, 1, 0,
              vx.size());
        j++;
53
54
      if(a[i].piv + 1 != SZ(vy)) ans += 1LL *
           seg.q() * (vy[a[i].piv + 1] - vy[a[i 43
           ].piv]);
      i = j - 1;
57
    cout << ans << "\n";
```

3 Flow-Matching

3.1 bipartite-matching

```
struct bipartite matching {
   int n, m; // 二分圖左右人數 (0 ~ n-1), (0
        \sim m-1)
    vector<vi> g;
    vi lhs, rhs, dist; // i 與 Lhs[i] 配對 (
         Lhs[i] == -1 代表沒有配對)
    bipartite_matching(int _n, int _m) : n(_n)
         , m(_m), g(_n), lhs(_n, -1), rhs(_m,
         -1), dist(_n) {}
    void add_edge(int u, int v) { g[u].pb(v);
    void bfs() {
      queue<int> q;
      REP(i, n) {
       if(lhs[i] == -1) {
                                               11
          q.push(i);
                                               12
          dist[i] = 0;
                                               13
        } else {
          dist[i] = -1;
                                               15
16
      while(!q.empty()) {
        int u = q.front(); q.pop();
```

```
if(rhs[v] != -1 && dist[rhs[v]] ==
        dist[rhs[v]] = dist[u] + 1;
        q.push(rhs[v]);
bool dfs(int u) {
  for(auto v : g[u]) {
    if(rhs[v] == -1) {
      rhs[lhs[u] = v] = u;
      return true;
  for(auto v : g[u]) {
   if(dist[rhs[v]] == dist[u] + 1 && dfs(
        rhs[v])) {
      rhs[lhs[u] = v] = u;
      return true;
  return false;
int solve() {
 int ans = 0;
  while(true) {
   bfs();
   int aug = 0;
    REP(i, n) if(lhs[i] == -1) aug += dfs(
   if(!aug) break;
    ans += aug;
  return ans;
```

for(auto v : g[u]) {

3.2 Dinic-LowerBound

```
1 template < class T>
2 struct DinicLowerBound {
   using Maxflow = Dinic<T>:
   int n;
   Maxflow d;
   vector<T> in:
   DinicLowerBound(int _n): n(_n), d(_n + 2)
        , in(_n) {}
   int add edge(int from, int to, T low, T
        high) {
      assert(0 <= low && low <= high);
     in[from] -= low, in[to] += low;
     return d.add_edge(from, to, high - low);
   T flow(int s, int t) {
     T sum = 0;
     REP(i, n) {
       if(in[i] > 0) {
         d.add_edge(n, i, in[i]);
         sum += in[i];
```

```
if(in[i] < 0) d.add_edge(i, n + 1, -in 50</pre>
                                                      edges[i].cap -= a;
                                                      edges[j ^ 1].cap += a;
                                                      if((r -= a) == 0) return f:
d.add_edge(t, s, numeric_limits<T>::max
if(d.flow(n, n + 1) < sum) return -1;</pre>
                                                  return f - r;
return d.flow(s, t);
                                                T flow(int s, int t, T f = INF) {
                                                  T ans = 0;
                                                  while(f > 0 && bfs(s, t)) {
                                                    cur.assign(n, 0);
                                                    T cur = dfs(s, t, f);
  Dinic
                                                    ans += cur;
                                                    f -= cur;
```

```
i template < class T>
2 class Dinic {
 public:
   struct Edge {
     int from, to;
     Edge(int x, int y, T z) : from(x), to(y)
          , cap(z) {}
   constexpr T INF = 1e9;
   int n;
   vector<Edge> edges;
   vector<vi> g;
   vi cur, h; // h : Level graph
   Dinic(int _n) : n(_n), g(_n) {}
   void add edge(int u, int v, T c) {
     g[u].pb(sz(edges));
     edges.eb(u, v, c);
     g[v].pb(sz(edges));
     edges.eb(v, u, 0);
   bool bfs(int s, int t) {
     h.assign(n, -1);
     queue<int> q;
     h[s] = 0;
     q.push(s);
     while(!q.empty()) {
       int u = q.front(); q.pop();
       for(int i : g[u]) {
         const auto& e = edges[i];
         int v = e.to;
         if(e.cap > 0 && h[v] == -1) {
           h[v] = h[u] + 1;
           if(v == t) return true;
           q.push(v);
     return false;
   T dfs(int u, int t, T f) {
     if(u == t) return f;
     Tr = f;
     for(int& i = cur[u]; i < sz(g[u]); ++i)</pre>
       int j = g[u][i];
       const auto& e = edges[j];
       int v = e.to;
       T c = e.cap:
       if(c > 0 \&\& h[v] == h[u] + 1) {
```

T = dfs(v, t, min(r, c));

3.4 Flow 建模

return ans;

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1

```
4. For each vertex v with d(v)>0, connect 11 S\to v with (cost, cap)=(0,d(v)) 12
```

- 5. For each vertex v with d(v) < 0, connect 13 $v \rightarrow T$ with (cost, cap) = (0, -d(v)) 14
- 6. Flow from S to T, the answer is the cost of the 15 flow C+K
- · Maximum density induced subgraph
 - Binary search on answer, suppose we're checking answer T
 - 2. Construct a max flow model, let K be the sum 21 of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K_{23}
 - 4. For each edge (u, v, w) in G, connect $u \to v$ 24 and $v \to u$ with capacity w 25
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$ 27
 - 6. T is a valid answer if the maximum flow $f < \frac{28}{29}$
- · Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $\frac{32}{32}$ $u' \to v'$ with weight w(u, v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $_{34}$ $_{\mu}(v)$ is the cost of the cheapest edge incident to $_{35}$
 - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

```
\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})
3.6
```

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity c_x and create edge (s, y) with capacity c_y .
- 2. Create edge (x, y) with capacity c_{xy} .
- 3. Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

3.5 general-matching

```
| struct GeneralMaxMatch {
| int n;
| vector<pii>| es;
| vi g, vis, mate; // i 與 mate[i] 配對 (
| mate[i] == -1 代表沒有匹配)
| GeneralMaxMatch(int n) : n(n), g(n, -1),
| mate(n, -1) {}
| bool dfs(int u) {
| if(vis[u]) return false;
| vis[u] = true;
| for(int ei = g[u]; ei != -1; ) {
| auto [x, y] = es[ei]; ei = y;
```

```
if(mate[x] == -1) {
     mate[mate[u] = x] = u;
      return true:
  for(int ei = g[u]; ei != -1; ) {
    auto [x, y] = es[ei]; ei = y;
   int nu = mate[x];
   mate[mate[u] = x] = u;
   mate[nu] = -1;
    if(dfs(nu)) return true;
   mate[mate[nu] = x] = nu;
   mate[u] = -1;
  return false:
void add_edge(int a, int b) {
  auto f = [&](int a, int b) {
   es.eb(b, g[a]);
   g[a] = sz(es) - 1;
  f(a, b); f(b, a);
int solve() {
 vi o(n);
  iota(all(o), 0);
 int ans = 0;
 REP(it, 100) {
   shuffle(all(o), rng);
   vis.assign(n, false);
   for(auto i : o) if(mate[i] == -1) ans
        += dfs(i);
  return ans;
```

```
1 // 1-based QQ
  struct WeightGraph {
    static const int inf = INT MAX;
    static const int maxn = 514:
    struct edge {
      int u, v, w;
      edge() {}
      edge(int u, int v, int w): u(u), v(v), w
           (w) {}
    int n, n x;
    edge g[maxn * 2][maxn * 2];
11
    int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[
12
         maxn * 2], pa[maxn * 2];
    int flo_from[maxn * 2][maxn + 1], S[maxn *
          2], vis[maxn * 2];
    vector<int> flo[maxn * 2];
15
    queue<int> q;
    int e delta(const edge &e) { return lab[e.
17
          u] + lab[e.v] - g[e.u][e.v].w * 2; }
    void update_slack(int u, int x) { if(!
          slack[x] || e delta(g[u][x]) < e delta</pre>
          (g[slack[x]][x])) slack[x] = u; }
```

general-weighted-max-matching

void set slack(int x) {

```
slack[x] = 0;
  REP(u, n) if(g[u + 1][x].w > 0 \&\& st[u +
       1] != x && S[st[u + 1]] == 0
       update slack(u + 1, x);
void q_push(int x) {
  if(x \le n) q.push(x);
  else REP(i, SZ(flo[x])) q_push(flo[x][i
void set_st(int x, int b) {
  if(x > n) REP(i, SZ(flo[x])) set st(flo[
       x][i], b);
int get_pr(int b, int xr) {
  int pr = find(ALL(flo[b]), xr) - flo[b].
  if(pr % 2 == 1) {
    reverse(1 + ALL(flo[b]));
    return SZ(flo[b]) - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if(u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u
  for(int i = 0; i < pr; ++i) set_match(</pre>
       flo[u][i], flo[u][i ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() +
       pr, flo[u].end());
void augment(int u, int v) {
  while(true) {
    int xnv = st[match[u]];
    set match(u, v);
    if(!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int get lca(int u, int v) {
  static int t = 0;
  for(++t; u || v; swap(u, v)) {
    if(u == 0) continue;
    if(vis[u] == t) return u;
    vis[u] = t;
    if(u = st[match[u]]) u = st[pa[u]];
                                            116
  return 0;
void add blossom(int u, int lca, int v) {
  int b = n + 1;
  while(b <= n_x && st[b]) ++b;</pre>
                                            121
  if(b > n x) n x++;
  lab[b] = S[b] = 0;
  match[b] = match[lca];
  flo[b].clear(); flo[b].pb(lca);
  for(int x = u, y; x != lca; x = st[pa[y
                                            124
       ]]) flo[b].pb(x), flo[b].pb(y = st[
                                            125
       match[x]]), q_push(y);
```

```
reverse(1 + ALL(flo[b]));
  for(int x = v, y; x != lca; x = st[pa[y 128]]
       ]]) flo[b].pb(x), flo[b].pb(y = st[ 129]
       match[x]]), q_push(y);
  set st(b, b);
  REP(x, n x) g[b][x + 1].w = g[x + 1][b]. 131
       w = 0;
  REP(x, n) flo from[b][x + 1] = 0;
  REP(i, SZ(flo[b])) {
                                              133
    int xs = flo[b][i];
                                              134
    REP(x, n_x) if(g[b][x + 1].w == 0 | |
                                             135
         e_delta(g[xs][x + 1]) < e_delta(g[ 136
         b][x + 1])) g[b][x + 1] = g[xs][x
         +1], g[x + 1][b] = g[x + 1][xs];
    REP(x, n) if(flo_from[xs][x + 1])
         flo from [b][x + 1] = xs;
                                              138
                                              139
  set_slack(b);
void expand_blossom(int b) {
  REP(i, SZ(flo[b])) set_st(flo[b][i], flo 141
       [b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr
                                             143
       = get_pr(b, xr);
                                              144
  for(int i = 0; i < pr; i += 2) {</pre>
                                              145
    int xs = flo[b][i], xns = flo[b][i +
                                             146
                                              147
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
                                              148
    slack[xs] = 0, set slack(xns);
                                              149
    q_push(xns);
                                              150
                                              151
  S[xr] = 1, pa[xr] = pa[b];
                                              152
  for(size_t i = pr + 1; i < SZ(flo[b]);</pre>
       ++i) {
    int xs = flo[b][i];
                                              154
    S[xs] = -1, set_slack(xs);
                                             155
                                             156
  st[b] = 0;
bool on found edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if(S[v] == -1) {
                                             158
    pa[v] = e.u, S[v] = 1;
                                              159
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
                                              160
    S[nu] = 0, q push(nu);
                                             161
  } else if(S[v] == 0) {
                                              162
    int lca = get lca(u, v);
    if(!lca) return augment(u,v), augment(
         v,u), true;
    else add blossom(u, lca, v);
  return false;
bool matching() {
                                              169
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x); 171
  q = queue<int>();
                                              172
  REP(x, n x) if(st[x + 1] == x + 1 &&!
                                             173
       match[x + 1]) pa[x + 1] = 0, S[x +
                                             174
       1] = 0, q push(x + 1);
  if(q.empty()) return false;
  while(true) {
                                             176
    while(!q.empty()) {
                                             177
      int u = q.front(); q.pop();
```

```
if(S[st[u]] == 1) continue;
      for(int v = 1; v \le n; ++v)
                                              180
        if(g[u][v].w > 0 && st[u] != st[v]
                                              181
             ]) {
                                              182
          if(e_delta(g[u][v]) == 0) {
            if(on found edge(g[u][v]))
                                              183
                 return true;
                                              184
          } else update slack(u, st[v]);
                                              185
    int d = inf:
    for(int b = n + 1; b <= n_x; ++b) if(</pre>
         st[b] == b \&\& S[b] == 1) d = min(d)
         , lab[b] / 2);
    for(int x = 1; x <= n_x; ++x) {
      if(st[x] == x && slack[x]) {
        if(S[x] == -1) d = min(d, e_delta(
             g[slack[x]][x]));
        else if(S[x] == 0) d = min(d,
             e_delta(g[slack[x]][x]) / 2);
    REP(u, n) {
      if(S[st[u + 1]] == 0) {
        if(lab[u + 1] <= d) return 0;</pre>
        lab[u + 1] -= d;
      } else if(S[st[u + 1]] == 1) lab[u +
            11 += d:
    for(int b = n + 1; b <= n x; ++b)
      if(st[b] == b) {
        if(S[st[b]] == 0) lab[b] += d * 2;
        else if(S[st[b]] == 1) lab[b] -= d
                                              13
    q = queue<int>();
    for(int x = 1; x \leftarrow n_x; ++x)
                                               15
      if(st[x] == x && slack[x] && st[
                                               16
           slack[x]] != x && e_delta(g[
                                               17
           slack[x]][x]) == 0
        if(on found edge(g[slack[x]][x]))
             return true;
                                               20
    for(int b = n + 1; b <= n_x; ++b)</pre>
                                               21
      if(st[b] == b && S[b] == 1 && lab[b]
            == 0) expand_blossom(b);
                                               23
                                               24
  return false;
                                               25
pair<11, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n x = n;
                                               27
  int n matches = 0;
                                               28
  11 tot weight = 0;
                                               29
  for(int u = 0; u <= n; ++u) st[u] = u,
       flo[u].clear();
                                               31
  int w max = 0;
                                               32
  for(int u = 1; u <= n; ++u)</pre>
                                               33
    for(int v = 1; v \le n; ++v) {
                                               34
      flo_from[u][v] = (u == v ? u : 0);
      w \max = \max(w \max, g[u][v].w);
                                               35
  for(int u = 1; u <= n; ++u) lab[u] =</pre>
                                               37
       w max:
                                               38
  while(matching()) ++n_matches;
                                               39
  for(int u = 1; u <= n; ++u)</pre>
    if(match[u] && match[u] < u)</pre>
```

3.7 KM

```
1 template < class T>
2 struct KM {
   static constexpr T INF = numeric limits<T</pre>
         >::max():
    int n, ql, qr;
    vector<vector<T>> w;
    vector<T> hl, hr, slk;
    vi fl, fr, pre, qu;
    vector<bool> v1, vr;
    KM(int n) : n(n), w(n, vector<T>(n, -INF))
         , hl(n), hr(n), slk(n), fl(n), fr(n),
         pre(n), qu(n), vl(n), vr(n) {}
    void add_edge(int u, int v, int x) { w[u][
         v] = x; } // 最小值要加負號
    bool check(int x) {
      vl[x] = 1:
      if(fl[x] != -1) return vr[qu[qr++] = fl[
           x]] = 1;
      while(x != -1) swap(x, fr[fl[x] = pre[x
      return 0;
    void bfs(int s) {
      fill(all(slk), INF);
      fill(all(v1), 0);
      fill(all(vr), 0);
      ql = qr = 0, qu[qr++] = s, vr[s] = 1;
      while(true) {
        while(ql < qr) {</pre>
          for(int x = 0, y = qu[ql++]; x < n;
            if(!vl[x] \&\& slk[x] >= (d = hl[x]
                 + hr[y] - w[x][y])) {
              pre[x] = y;
              if(d) slk[x] = d;
              else if(!check(x)) return;
        REP(x, n) if(!vl[x] \&\& d > slk[x]) d =
              slk[x];
        REP(x, n) {
          if(vl[x]) hl[x] += d;
          else slk[x] -= d;
          if(vr[x]) hr[x] -= d;
        REP(x, n) if(!v1[x] \&\& !s1k[x] \&\& !
             check(x)) return;
```

```
T solve() {
  fill(all(fl), -1);
  fill(all(fr), -1);
  fill(all(hr), 0);
  REP(i, n) hl[i] = *max_element(all(w[i]) 44
                                                 vector<int> solve(){
  REP(i, n) bfs(i);
                                                  vector<P> remark;
                                                   REP(i, n) remark.emplace_back(i, -1, SZ(
  T ans = 0;
  REP(i, n) ans += w[i][fl[i]]; // i 跟 fL
                                                   dfs(remark);
       [i] 配對
                                                   return clique;
  return ans;
                                            51 };
```

3.8 max-clique

MCMF

```
1 template<int V>
 struct max clique {
   using B = bitset<V>;
   int n = 0;
   vector<B> g, buf;
   struct P {
     int idx, col, deg;
     P(int a, int b, int b) : idx(a), col(b),
           deg(c) {}
   max_{clique}(int_n) : n(n), g(n), buf(n)
   void add edge(int a, int b) {
     assert(a != b);
     g[a][b] = g[b][a] = 1;
   vector<int> now, clique;
   void dfs(vector<P>& rem){
     if(SZ(clique) < SZ(now)) clique = now;</pre>
     sort(ALL(rem), [](P a, P b) { return a.
          deg > b.deg; });
     int max_c = 1;
     for(auto& p : rem){
       p.col = 0:
       while((g[p.idx] & buf[p.col]).any()) p
       max_c = max(max_c, p.idx + 1);
       buf[p.col][p.idx] = 1;
     REP(i, max c) buf[i].reset();
     sort(ALL(rem), [&](P a, P b) { return a.
          col < b.col; });</pre>
     for(; !rem.empty(); rem.pop_back()){
       auto& p = rem.back();
       if(now.size() + p.col + 1 <= clique.</pre>
            size()) break;
       vector<P> nrem;
       B bs;
       for(auto& q : rem){
         if(g[p.idx][q.idx]){
           nrem.emplace_back(q.idx, -1, 0);
           bs[q.idx] = 1;
```

```
class MCMF {
public:
 struct Edge {
   int from, to;
    S cap:
    T cost;
    Edge(int u, int v, S x, T y) : from(u),
         to(v), cap(x), cost(y) {}
  const ll INF = 1e18L;
  int n;
  vector<Edge> edges;
  vector<vi> g;
  vector<T> d;
  vector<bool> inq;
  MCMF(int _n) : n(_n), g(_n), d(_n), inq(_n
       ), pedge( n) {}
  void add_edge(int u, int v, S cap, T cost)
    g[u].pb(sz(edges));
    edges.eb(u, v, cap, cost);
    g[v].pb(sz(edges));
    edges.eb(v, u, 0, -cost);
  bool spfa(int s, int t) {
   bool found = false;
    fill(all(d), INF);
    d[s] = 0;
    inq[s] = true;
    queue<int> q;
    q.push(s);
    while(!q.empty()) {
      int u = q.front(); q.pop();
      if(u == t) found = true;
      inq[u] = false;
      for(auto& id : g[u]) {
        const auto& e = edges[id];
        if(e.cap > 0 && d[u] + e.cost < d[e.</pre>
             to]) {
          d[e.to] = d[u] + e.cost;
          pedge[e.to] = id;
          if(!inq[e.to]) {
            q.push(e.to);
```

for(auto& q : nrem) q.deg = (bs & g[q. 42]

idx]).count();

now.emplace back(p.idx);

dfs(nrem);

now.pop_back();

g[i]));

i template < class S, class T>

```
return found:
    pair<S, T> flow(int s, int t, S f = INF) {
      S cap = 0:
      T cost = 0;
                                                  30
      while(f > 0 && spfa(s, t)) {
                                                  31
        S send = f;
        int u = t;
                                                  33
        while(u != s) {
                                                  34
          const Edge& e = edges[pedge[u]];
                                                  35
          send = min(send, e.cap);
          u = e.from:
        u = t:
        while(u != s) {
          Edge& e = edges[pedge[u]];
          e.cap -= send;
          Edge& b = edges[pedge[u] ^ 1];
          b.cap += send:
          u = e.from;
        cap += send;
        f -= send:
        cost += send * d[t];
72
      return {cap, cost};
73
```

inq[e.to] = true;

}

3.10 minimum-general-weightedperfect-matching

```
1 struct Graph {
    // Minimum General Weighted Matching (
         Perfect Match) 0-base
    static const int MXN = 105;
    int n, edge[MXN][MXN];
    int match[MXN], dis[MXN], onstk[MXN];
    vector<int> stk:
    void init(int _n) {
      for(int i=0; i<n; i++)</pre>
        for(int j=0; j<n; j++)</pre>
           edge[i][j] = 0;
11
12
    void add_edge(int u, int v, int w) { edge[
13
         u|[v] = edge[v][u] = w;
    bool SPFA(int u){
      if(onstk[u]) return true;
      stk.push back(u);
      onstk[u] = 1;
      for(int v=0; v<n; v++){</pre>
        if(u != v && match[u] != v && !onstk[v
             ]){
20
           int m = match[v];
           if(dis[m] > dis[u] - edge[v][m] +
                edge[u][v]){
```

```
dis[m] = dis[u] - edge[v][m] +
                  edge[u][v];
            onstk[v] = 1;
            stk.push_back(v);
            if(SPFA(m)) return true;
            stk.pop back();
            onstk[v] = 0;
      onstk[u] = 0;
      stk.pop_back();
      return false:
    int solve() {
      for(int i = 0; i < n; i += 2) match[i] =</pre>
            i + 1, match[i+1] = i;
      while(true) {
        int found = 0:
        for(int i=0; i<n; i++) dis[i] = onstk[</pre>
        for(int i=0; i<n; i++){</pre>
          stk.clear();
          if(!onstk[i] && SPFA(i)){
            found = 1;
            while(stk.size()>=2){
               int u = stk.back(); stk.pop back
               int v = stk.back(); stk.pop_back
               match[u] = v;
               match[v] = u;
        if(!found) break;
      int ans = 0;
      for(int i=0; i<n; i++) ans += edge[i][</pre>
           match[i]];
      return ans / 2;
58 } graph;
```

Geometry

4.1 closest-pair

```
1 const 11 INF = 9e18L + 5:
vector<P> a;
 sort(all(a), [](P a, P b) { return a.x < b.x
       ; });
4 11 SQ(11 x) { return x * x; }
5 11 solve(int 1, int r) {
   if(1 + 1 == r) return INF;
    int m = (1 + r) / 2;
    11 \text{ midx} = a[m].x;
    11 d = min(solve(1, m), solve(m, r));
    inplace_merge(a.begin() + 1, a.begin() + m
         , a.begin() + r, [](Pa, Pb) {
      return a.y < b.y;</pre>
```

4.2 convex-hull

4.3 point-in-convex-hull

```
i int point in convex hull(const vector<P>& a,
       P p) {
   // -1 ON, 0 OUT, +1 IN
   // 要先逆時針排序
   int n = sz(a);
   if(btw(a[0], a[1], p) || btw(a[0], a[n -
       1], p)) return -1;
   int 1 = 0, r = n - 1;
   while(1 <= r) {</pre>
     int m = (1 + r) / 2;
     auto a1 = cross(a[m] - a[0], p - a[0]);
     auto a2 = cross(a[(m + 1) \% n] - a[0], p
           - a[0]);
     if(a1 >= 0 && a2 <= 0) {
       auto res = cross(a[(m + 1) % n] - a[m
            ], p - a[m]);
       return res > 0 ? 1 : (res >= 0 ? -1 :
     if(a1 < 0) r = m - 1;
     else l = m + 1;
   return 0;
```

4.4 point

```
1 using P = pair<11, 11>;
 2 P operator+(P a, P b) { return P{a.X + b.X,
       a.Y + b.Y; }
 3 P operator-(P a, P b) { return P{a.X - b.X,
        a.Y - b.Y}; }
  P operator*(P a, 11 b) { return P{a.X * b, a
        .Y * b}; }
   P operator/(P a, 11 b) { return P{a.X / b, a
        .Y / b}; }
  11 dot(P a, P b) { return a.X * b.X + a.Y *
   11 cross(P a, P b) { return a.X * b.Y - a.Y
        * b.X: }
   11 abs2(P a) { return dot(a, a); }
   double abs(P a) { return sqrt(abs2(a)); }
   int sign(ll x) { return x < 0 ? -1 : (x == 0
        ? 0 : 1); }
   int ori(P a, P b, P c) { return sign(cross(b
        - a, c - a)); }
12 bool collinear(P a, P b, P c) { return sign(
        cross(a - c, b - c)) == 0; }
13 bool btw(Pa, Pb, Pc) {
    if(!collinear(a, b, c)) return 0;
     return sign(dot(a - c, b - c)) <= 0:
bool seg_intersect(P a, P b, P c, P d) {
    int a123 = ori(a, b, c);
     int a124 = ori(a, b, d);
     int a341 = ori(c, d, a);
     int a342 = ori(c, d, b);
     if(a123 == 0 && a124 == 0) {
       return btw(a, b, c) || btw(a, b, d) ||
            btw(c, d, a) || btw(c, d, b);
    return a123 * a124 <= 0 && a341 * a342 <=
          0;
   P intersect(P a, P b, P c, P d) {
    int a123 = cross(b - a, c - a);
    int a124 = cross(b - a, d - a);
     return (d * a123 - c * a124) / (a123 -
          a124):
33 struct line { P A, B; };
34 P vec(line L) { return L.B - L.A; }
35 P projection(P p, line L) { return L.A + vec
        (L) / abs(vec(L)) * dot(p - L.A, vec(L)) <sub>12</sub>
        / abs(vec(L)); }
```

4.5 polar-angle-sort

5 Graph

7 }

5.1 **2-SAT**

```
1 struct two sat {
    int n; SCC g;
    vector<bool> ans;
    two_sat(int _n) : n(_n), g(_n * 2) {}
    void add or(int u, bool x, int v, bool y)
      g.add_edge(2 * u + !x, 2 * v + y);
      g.add edge(2 * v + !y, 2 * u + x);
    bool solve() {
      ans.resize(n);
      auto id = g.solve();
      REP(i, n) {
        if(id[2 * i] == id[2 * i + 1]) return
        ans[i] = (id[2 * i] < id[2 * i + 1]);
      return true:
17
18 };
```

5.2 centroid-tree

```
19
| pair<int, vector<vi>>> centroid_tree(const
      vector<vi>& g) {
                                                 21
   int n = sz(g);
                                                 22
   vi siz(n);
   vector<bool> vis(n);
   auto dfs sz = [&](auto f, int u, int p) ->
                                                 25
          void {
                                                 26
     siz[u] = 1;
     for(auto v : g[u]) {
       if(v == p || vis[v]) continue;
       f(f, v, u);
                                                 29
       siz[u] += siz[v];
                                                 30
   auto find_cd = [&](auto f, int u, int p,
        int all) -> int {
     for(auto v : g[u]) {
                                                 33
       if(v == p || vis[v]) continue;
                                                 34
       if(siz[v] * 2 > all) return f(f, v, u,
             all);
     return u;
                                                 37
   vector<vi> h(n);
   auto build = [&](auto f, int u) -> int {
     dfs_sz(dfs_sz, u, -1);
     int cd = find cd(find cd, u, -1, siz[u])
     vis[cd] = true;
```

5.3 HLD

1 struct HLD {

12

```
int n;
vector<vi> g;
vi siz, par, depth, top, tour, fi, id;
sparse_table<pii, min> st;
HLD(int _n) : n(_n), g(_n), siz(_n), par(
     _n), depth(_n), top(_n), fi(_n), id(_n
  tour.reserve(n);
void add edge(int u, int v) {
  g[u].push_back(v);
  g[v].push back(u);
void build(int root = 0) {
  par[root] = -1;
  top[root] = root;
  vector<pii> euler tour;
  euler tour.reserve(2 * n - 1);
  dfs sz(root);
  dfs link(euler tour, root);
  st = sparse_table<pii, min>(euler_tour);
int get lca(int u, int v) {
  int L = fi[u], R = fi[v];
  if(L > R) swap(L, R);
  return st.prod(L, R).second;
bool is anc(int u, int v) {
  return id[u] <= id[v] && id[v] < id[u] +</pre>
        siz[u];
bool on_path(int a, int b, int x) {
  return (is ancestor(x, a) || is ancestor
       (x, b)) && is_ancestor(get_lca(a, b)
       , x);
int get dist(int u, int v) {
  return depth[u] + depth[v] - 2 * depth[(
       get lca(u, v))];
int kth_anc(int u, int k) {
  if(depth[u] < k) return -1;</pre>
  int d = depth[u] - k;
  while(depth[top[u]] > d) u = par[top[u
      11:
  return tour[id[u] + d - depth[u]];
int kth node on path(int a, int b, int k)
```

```
int z = get_lca(a, b);
  int fi = depth[a] - depth[z];
  int se = depth[b] - depth[z];
  if(k < 0 \mid | k > fi + se) return -1;
  if(k < fi) return kth anc(a, k);</pre>
  return kth anc(b, fi + se - k);
vector<pii> get path(int u, int v, bool
     include lca = true) {
  if(u == v && !include_lca) return {};
  vector<pii> seg;
  while(top[u] != top[v]) {
    if(depth[top[u]] > depth[top[v]]) swap
    seg.eb(id[top[v]], id[v]);
    v = par[top[v]];
  if(depth[u] > depth[v]) swap(u, v); // u
  if(u != v || include_lca) seg.eb(id[u] +
        !include_lca, id[v]);
  return seg;
void dfs sz(int u) {
 if(par[u] != -1) g[u].erase(find(all(g[u
       ]), par[u]));
  siz[u] = 1;
  for(auto& v : g[u]) {
    par[v] = u;
    depth[v] = depth[u] + 1;
    dfs_sz(v);
    siz[u] += siz[v];
    if(siz[v] > siz[g[u][0]]) swap(v, g[u
        ][0]);
void dfs_link(vector<pii>& euler_tour, int
  fi[u] = sz(euler_tour);
  id[u] = sz(tour);
  euler tour.eb(depth[u], u);
  tour.pb(u);
  for(auto v : g[u]) {
    top[v] = (v == g[u][0] ? top[u] : v);
    dfs_link(euler_tour, v);
    euler tour.eb(depth[u], u);
```

5.4 lowlink

```
| struct lowlink {
   int n, cnt = 0, tecc cnt = 0, tvcc cnt =
   vector<vector<pii>>> g;
   vector<pii> edges;
   vi roots, id, low, tecc_id, tvcc_id;
   vector<bool> is_bridge, is_cut,
        is_tree_edge;
   lowlink(int _n) : n(_n), g(_n), is_cut(_n,
         false), id(_n, -1), low(_n, -1) {}
   void add_edge(int u, int v) {
```

```
g[u].eb(v, sz(edges));
  g[v].eb(u, sz(edges));
  edges.eb(u, v):
  is_bridge.pb(false);
  is_tree_edge.pb(false);
  tvcc id.pb(-1);
void dfs(int u, int peid = -1) {
  static vi stk:
  static int rid;
  if(peid < 0) rid = cnt;</pre>
  if(peid == -1) roots.pb(u);
  id[u] = low[u] = cnt++;
  for(auto [v, eid] : g[u]) {
   if(eid == peid) continue;
   if(id[v] < id[u]) stk.pb(eid);</pre>
    if(id[v] >= 0) {
     low[u] = min(low[u], id[v]);
    } else {
      is_tree_edge[eid] = true;
      dfs(v, eid);
      low[u] = min(low[u], low[v]);
      if((id[u] == rid && id[v] != rid +
          1) || (id[u] != rid && low[v] >= 94
           id[u])) {
        is_cut[u] = true;
     if(low[v] >= id[u]) {
        while(true) {
         int e = stk.back();
          stk.pop_back();
          tvcc_id[e] = tvcc_cnt;
          if(e == eid) break;
        tvcc cnt++;
void build() {
 REP(i, n) if(id[i] < 0) dfs(i);
  REP(i, sz(edges)) {
    auto [u, v] = edges[i];
   if(id[u] > id[v]) swap(u, v);
   is_bridge[i] = (id[u] < low[v]);</pre>
vector<vi> two ecc() { // 邊雙
  tecc cnt = 0;
  tecc_id.assign(n, -1);
  vi stk;
  REP(i, n) {
   if(tecc id[i] != -1) continue;
    tecc id[i] = tecc cnt;
    stk.pb(i);
    while(sz(stk)) {
     int u = stk.back(); stk.pop_back();
      for(auto [v, eid] : g[u]) {
        if(tecc id[v] >= 0 || is bridge[
             eid]) {
          continue;
```

tecc_id[v] = tecc_cnt;

stk.pb(v);

```
tecc cnt++;
      vector<vi> comp(tecc_cnt);
      REP(i, n) comp[tecc_id[i]].pb(i);
    vector<vi> bcc vertices() { // 點雙
      vector<vi> comp(tvcc_cnt);
      REP(i, sz(edges)) {
        comp[tvcc_id[i]].pb(edges[i].first);
        comp[tvcc id[i]].pb(edges[i].second);
      for(auto& v : comp) {
        sort(all(v)):
        v.erase(unique(all(v)), v.end());
      REP(i, n) if(g[i].empty()) comp.pb({i});
      return comp;
    vector<vi> bcc edges() {
      vector<vi> ret(tvcc cnt);
      REP(i, sz(edges)) ret[tvcc id[i]].pb(i);
      return ret:
96 };
```

6 Math

33

36 37 };

cnt += 1;

return id;

6.1 Aliens

```
1 template < class Func, bool MAX>
  ll Aliens(ll l, ll r, int k, Func f) {
    while(l < r) {</pre>
      11 \text{ m} = 1 + (r - 1) / 2;
      auto [score, op] = f(m);
      if(op == k) return score + m * k * (MAX
           ? +1 : -1);
      if(op < k) r = m;
      else 1 = m + 1;
    return f(1).first + 1 * k * (MAX ? +1 :
          -1);
11 }
```

5.5 SCC

26

27

```
1 struct SCC {
    int n:
    vector<vi> g, h;
    SCC(int _n) : n(_n), g(_n), h(_n) {}
    void add edge(int u, int v) {
      g[u].pb(v);
      h[v].pb(u);
    vi solve() { // 回傳縮點的編號
      vi id(n), top;
      top.reserve(n);
      #define GO if(id[v] == 0) dfs1(v);
      function<void(int)> dfs1 = [&](int u) {
        id[u] = 1;
        for(auto v : g[u]) GO;
        top.pb(u);
      REP(v, n) GO:
      fill(all(id), -1);
      function<void(int, int)> dfs2 = [&](int
           u, int x) {
        id[u] = x;
        for(auto v : h[u]) {
          if(id[v] == -1) {
            dfs2(v, x);
25
        }
      for(int i = n - 1, cnt = 0; i >= 0; --i) 23
        int u = top[i];
        if(id[u] == -1) {
```

dfs2(u, cnt);

6.2 Berlekamp-Massey

```
1 // - [1, 2, 4, 8, 16] \rightarrow (1, [1, -2])
2 // - [1, 1, 2, 3, 5, 8] -> (2, [1, -1, -1])
 998244352]) (mod 998244353)
 4 // - [] -> (0, [1])
 5 // - [0, 0, 0] -> (0, [1])
 6 // - [-2] -> (1, [1, 2])
 7 template < class T>
  pair<int, vector<T>> BM(const vector<T>& S)
    using poly = vector<T>;
    int N = SZ(S);
    poly C rev{1}, B{1};
    int L = 0, m = 1;
13
    T b = 1;
    auto adjust = [](poly C, const poly &B, T
         d, T b, int m) -> poly {
      C.resize(max(SZ(C), SZ(B) + m));
      Ta = d / b:
17
      REP(i, SZ(B)) C[i + m] -= a * B[i];
18
      return C;
19
    REP(n, N) {
20
      T d = S[n];
      REP(i, L) d += C_{rev}[i + 1] * S[n - 1 -
           i];
      if(d == 0) m++;
      else if (2 * L <= n) {
25
        poly Q = C rev;
        C rev = adjust(C rev, B, d, b, m);
        L = n + 1 - L, B = Q, b = d, m = 1;
```

```
} else C_rev = adjust(C_rev, B, d, b, m
           ++);
   return {L, C_rev};
33 // Calculate x^N \b f(x)
34 // Complexity: $0(K^2 \Log N)$ ($K$: dea. of
| (4, [1, -1, -1]) | \rightarrow [2, 3] 
2)
37 template < class T>
38 vector<T> monomial mod polynomial(long long
       N, const vector<T> &f_rev) {
    assert(!f rev.empty() && f rev[0] == 1);
    int K = SZ(f_rev) - 1;
    if(!K) return {};
    int D = 64 - builtin clzll(N);
    vector<T> ret(K, 0);
    ret[0] = 1;
    auto self conv = [](vector<T> x) -> vector
         <T> {
      int d = SZ(x);
      vector<T> ret(d * 2 - 1);
      REP(i, d) {
        ret[i * 2] += x[i] * x[i];
        REP(j, i) ret[i + j] += x[i] * x[j] *
      return ret;
    for(int d = D; d--;) {
      ret = self_conv(ret);
      for(int i = 2 * K - 2; i >= K; i--) {
        REP(j, k) ret[i - j - 1] -= ret[i] *
            f rev[i + 1];
      ret.resize(K);
      if (N >> d & 1) {
        vector<T> c(K);
        c[0] = -ret[K - 1] * f_rev[K];
        for(int i = 1; i < K; i++) c[i] = ret[</pre>
            i - 1] - ret[K - 1] * f rev[K - i
            ];
        ret = c;
    return ret;
  // Guess k-th element of the sequence,
       assuming linear recurrence
71 template < class T>
72 T guess kth term(const vector<T>& a, long
      long k) {
    assert(k >= 0);
    if(k < 1LL * SZ(a)) return a[k];</pre>
    auto f = BM<T>(a).second;
    auto g = monomial mod polynomial<T>(k, f);
    T ret = 0;
    REP(i, SZ(g)) ret += g[i] * a[i];
    return ret:
```

6.3 Chinese-Remainder

```
1 // (rem, mod) {0, 0} for no solution
2 pair<11, 11> crt(11 r0, 11 m0, 11 r1, 11 m1)
   r0 = (r0 \% m0 + m0) \% m0;
    r1 = (r1 \% m1 + m1) \% m1;
    if(m0 < m1) swap(r0, r1), swap(m0, m1);</pre>
   if(m0 % m1 == 0) {
     if(r0 % m1 != r1) return {0, 0};
   11 g, im, qq;
   g = ext_gcd(m0, m1, im, qq);
   11 u1 = (m1 / g);
   if((r1 - r0) % g) return {0, 0};
   11 x = (r1 - r0) / g % u1 * im % u1;
   r0 += x * m0;
   m0 *= u1:
   if(r0 < 0) r0 += m0;
    return {r0, m0};
```

6.4 Combination

6.5 Determinant

```
1 T det(vector<vector<T>> a) {
   int n = SZ(a);
   T ret = 1;
   REP(i, n) {
     int idx = -1;
     for(int j = i; j < n; j++) {</pre>
       if(a[j][i] != 0) {
         idx = j;
         break;
     if(idx == -1) return 0;
     if(i != idx) {
       ret *= T(-1);
       swap(a[i], a[idx]);
     ret *= a[i][i];
     T inv = T(1) / a[i][i];
     REP(j, n) a[i][j] *= inv;
```

6.6 Discrete-Log

6.7 extgcd

6.8 Floor-Sum

6.9 FWHT

```
1 #define ppc __builtin_popcount
2 template < class T, class F>
3 void fwht(vector<T>& a, F f) {
    int n = SZ(a);
    assert(ppc(n) == 1);
    for(int i = 1; i < n; i <<= 1) {</pre>
      for(int j = 0; j < n; j += i << 1) {</pre>
        REP(k, i) f(a[j + k], a[i + j + k]);
10
11 }
12 template < class T>
void or_transform(vector<T>& a, bool inv) {
       fwht(a, [\&](T\& x, T\& y) \{ y += x * (inv) \}
       ? -1 : +1); }) }
14 template < class T>
15 void and transform(vector<T>& a, bool inv) {
        fwht(a, [\&](T\& x, T\& y) \{ x += y * (inv) \}
        ? -1 : +1); }); }
16 template < class T>
void xor transform(vector<T>& a, bool inv) {
    fwht(a, [](T& x, T& y) {
      T z = x + y;
      y = x - y;
      x = z;
22
    if(inv) {
      T z = T(1) / T(SZ(a));
      for(auto& x : a) x *= z;
25
26
27 }
28 template < class T>
  vector<T> convolution(vector<T> a, vector<T>
    assert(SZ(a) == SZ(b));
    transform(a, false), transform(b, false);
    REP(i, SZ(a)) a[i] *= b[i];
33
    transform(a, true);
34
    return a;
35 }
36 template < class T>
  vector<T> subset convolution(const vector<T</pre>
       >& f, const vector<T>& g) {
    assert(SZ(f) == SZ(g));
    int n = SZ(f);
    assert(ppc(n) == 1);
    const int lg = __lg(n);
    vector<vector<T>> fhat(lg + 1, vector<T>(n
         )), ghat(fhat);
    REP(i, n) fhat[ppc(i)][i] = f[i], ghat[ppc
         (i)][i] = g[i];
    REP(i, lg + 1) or_transform(fhat[i], false
         ), or_transform(ghat[i], false);
    vector<vector<T>> h(lg + 1, vector<T>(n));
    REP(m, n) REP(i, lg + 1) REP(j, i + 1) h[i]
         ][m] += fhat[j][m] * ghat[i - j][m];
```

```
REP(i, lg + 1) or_transform(h[i], true);
vector<T> res(n);
REP(i, n) res[i] = h[ppc(i)][i];
return res;
```

i int GaussJordan(vector<vector<ld>>& a) {

Gauss-Jordan

```
// -1 no sol, 0 inf sol
int n = SZ(a);
REP(i, n) assert(SZ(a[i]) == n + 1);
REP(i, n) {
  int p = i;
  REP(j, n) {
    if(j < i && abs(a[j][j]) > EPS)
    if(abs(a[j][i]) > abs(a[p][i])) p = j;
  REP(j, n + 1) swap(a[i][j], a[p][j]);
  if(abs(a[i][i]) <= EPS) continue;</pre>
  REP(j, n) {
    if(i == j) continue;
    ld delta = a[j][i] / a[i][i];
    FOR(k, i, n + 1) a[j][k] -= delta * a[
         i][k];
bool ok = true;
REP(i, n) {
  if(abs(a[i][i]) <= EPS) {</pre>
    if(abs(a[i][n]) > EPS) return -1;
    ok = false;
}
return ok;
```

6.11 Int-Div

```
return a/b - ((a^b) < 0 && a%b != 0);
4 ll ceil div(ll a, ll b) {
  return a/b + ((a^b) > 0 && a^b != 0);
```

6.12 Linear-Sieve

```
| vi primes, least = {0, 1}, phi, mobius;
void LinearSieve(int n) {
   least = phi = mobius = vi(n + 1);
   for(int i = 2; i <= n; i++) {
     if(!least[i]) {
       least[i] = i;
       primes.pb(i);
```

```
phi[i] = i - 1;
        mobius[i] = -1;
      for(auto j : primes) {
        if(i * j > n) break;
        least[i * j] = j;
        if(i % j == 0) {
          mobius[i * j] = 0;
          phi[i * j] = phi[i] * j;
          break:
        } else {
          mobius[i * j] = -mobius[i];
          phi[i * j] = phi[i] * phi[j];
23
```

6.13 Miller-Rabin

1 bool is prime(ll n, vector<ll> x) {

```
ll d = n - 1;
 d >>= builtin ctzll(d);
 for(auto a : x) {
   if(n <= a) break;</pre>
   11 t = d, y = 1, b = t;
   while(b) {
     if(b \& 1) y = i128(y) * a % n;
     a = i128(a) * a % n;
     b >>= 1;
   while(t != n - 1 && v != 1 && v != n -
     y = i128(y) * y % n;
     t <<= 1;
   if(y != n - 1 && t % 2 == 0) return
        false;
 return true;
bool is prime(ll n) {
 if(n <= 1) return false;</pre>
 if(n % 2 == 0) return n == 2;
 if(n < (1LL << 30)) return is prime(n, {2,
       7, 61});
 return is_prime(n, {2, 325, 9375, 28178,
      450775, 9780504, 1795265022});
```

6.14 Min-of-Mod-of-Linear

```
1 // \min\{Ax + B \pmod{M} \mid \emptyset \le x \le N\}
int min_of_mod_of_linear(int n, int m, int a
       , int b) {
   11 v = floor_sum(n, m, a, b);
   int l = -1, r = m - 1;
    while (r - 1 > 1) {
     int k = (1 + r) / 2;
```

```
if(floor sum(n, m, a, b + (m - 1 - k)) < 32
                                                         res.pb(x);
           v + n) r = k;
                                                33
                                                         return;
     else 1 = k:
                                                       for(int j = v[i].second; ; j--) {
   return r;
                                                        f(f, i + 1, x);
                                                        if(j == 0) break;
                                                        x *= v[i].first;
 6.15 Mod-Inv
                                                    f(f, 0, 1);
                                                    sort(all(res));
                                                    return res;
i int inv(int a) {
   if(a < N) return inv[a];</pre>
   if(a == 1) 1;
   return (MOD - 1LL * (MOD / a) * inv(MOD %
                                                  6.17 Primes
        a) % MOD) % MOD;
 vi mod_inverse(int m, int n = -1) {
                                                 1 /* 12721 13331 14341 75577 123457 222557
   assert(n < m);</pre>
   if(n == -1) n = m - 1;
                                                       556679 999983 1097774749 1076767633
   vi inv(n + 1);
                                                       100102021 999997771 1001010013
```

16

19

20

23

24

6.16 Pollard-Rho

if(i == sz(v)) {

return inv;

inv[0] = inv[1] = 1;

```
1 | void PollardRho(map<11, int>& mp, 11 n) {
   if(n == 1) return;
    if(is prime(n)) return mp[n]++, void();
    if(n % 2 == 0) {
      mp[2] += 1;
      PollardRho(mp, n / 2);
    11 x = 2, y = 2, d = 1, p = 1;
    #define f(x, n, p) ((i128(x) * x % n + p)
    while(true) {
      if(d!= 1 && d!= n) {
        PollardRho(mp, d);
        PollardRho(mp, n / d);
        return;
      p += (d == n);
      x = f(x, n, p), y = f(f(y, n, p), n, p);
                                                 18
      d = \_gcd(abs(x - y), n);
20
21
    #undef f
  vector<ll> get_divisors(ll n) {
    if(n == 0) return {};
    map<ll, int> mp;
    PollardRho(mp, n);
    vector<pair<ll, int>> v(all(mp));
    vector<11> res;
    auto f = [\&](auto f, int i, ll x) \rightarrow void
```

for(int i = 2; i <= n; i++) inv[i] = m - 1</pre>

LL * (m / i) * inv[m % i] % m;

6.18 Xor-Basis

```
1 template<int B>
2 struct xor basis {
    using T = long long;
    bool zero = false, change = false;
    int cnt = 0;
    array < T, B > p = {};
    vector<T> d;
    void insert(T x) {
      IREP(i, B) {
        if(x >> i & 1) {
          if(!p[i]) {
            p[i] = x, cnt++;
            change = true;
            return:
          } else x ^= p[i];
      if(!zero) zero = change = true;
    T get_min() {
      if(zero) return 0;
      REP(i, B) if(p[i]) return p[i];
    T get_max() {
      T ans = 0:
      IREP(i, B) ans = max(ans, ans ^ p[i]);
      return ans;
28
29
      get_kth(long long k) {
      k++;
      if(k == 1 && zero) return 0;
      k -= zero;
```

1000512343 987654361 999991231 999888733

1000000000000037 2305843009213693951

4611686018427387847 9223372036854775783

1010102101 1000000000039

18446744073709551557 */

98789101 987777733 999991921 1010101333

```
if(k >= (1LL << cnt)) return -1;
      update();
      T ans = 0:
      REP(i, SZ(d)) if(k \gg i \& 1) ans ^= d[i
      return ans;
    bool contains(T x) {
      if(x == 0) return zero;
      IREP(i, B) if(x >> i \& 1) x ^= p[i];
      return x == 0;
    void merge(const xor_basis& other) { REP(i
         , B) if(other.p[i]) insert(other.p[i])
    void update() {
      if(!change) return;
      change = false;
      d.clear();
      REP(j, B) IREP(i, j) if(p[j] >> i & 1) p
           [j] ^= p[i];
      REP(i, B) if(p[i]) d.pb(p[i]);
52 };
```

估計值

- Estimation
 - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for
 - n < 1e10, 200000 for n < 1e19.

 The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1, 1, 2, 3, 5, 7, 11, 15, 22, 30 for $n = 0 \sim 9,627 \text{ for } n = 20, \sim 2e5 \text{ for }$ $n = 50, \sim 2e8 \text{ for } n = 100.$
 - Total number of partitions of $\begin{array}{lll} n & \text{distinct} & \text{elements:} & B(n) & = & \det(M) = \sum_{S:A \rightarrow B} (-1)^{t(\sigma(S))} \prod_{i=1}^n \omega(S_i) \\ 1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597, & \\ 27644437,190890322 & & \\ \end{array}$ 27644437, 190899322,

定理 6.20

· Cramer's rule

$$ax + by = e \Rightarrow x = \frac{ed - bf}{ad - bc}$$
$$cx + dy = f \Rightarrow y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Burnside's Lemma

Let us calculate the number of necklaces of n pearls, where each pearl has m possible colors. Two necklaces

are symmetric if they are similar after rotating them. There are n ways to change the position of a necklace, because we can rotate it $0, 1, \ldots, n_1$ steps clockwise. If the number of steps is 0, all m^n necklaces remain the same, and if the number of steps is 1, only the m necklaces where each pearl has the same color remain the same. More generally, when the number of steps is k, a total of $m^{\gcd(k,n)}$ necklaces remain the same. The reason for this is that blocks of pearls of size gcd(k, n)will replace each other. Thus, according to Burnside's lemma, the number of necklaces is $\sum_{i=0}^{n-1} \frac{m^{\gcd(i,n)}}{n}$. For example, the number of necklaces of length 4 with 3 colors is $\frac{3^4+3+3^2+3}{4}=24$

Lindstr□m–Gessel–Viennot Lemma

定義

 $\omega(P)$ 表示 P 這條路徑上所有邊的邊權之積。(路徑 計數時,可以將邊權都設為1)(事實上,邊權可以為 生成函數) e(u,v) 表示 u 到 v 的 ** 每一條 ** 路徑 生成函數) e(u,v) 农小 u ±3 v +3 v +3 v +4 v +2 v +4 v

集合 A,是有向無環圖點集的一個子集,大小為 n。 終點集合 B, 也是有向無環圖點集的一個子集, 大小 也為 $n \circ -$ 組 $A \to B$ 的不相交路徑 $S : S_i$ 是一條從 A_i 到 $B_{\sigma(S)_i}$ 的路徑 ($\sigma(S)$ 是一個排列),對於任 何 $i \neq j$ S_i 和 S_i 沒有公共頂點 $t(\sigma)$ 表示排列 σ 的逆序對個數。

$$M = \begin{bmatrix} e(A_1, B_1) & e(A_1, B_2) & \cdots & e(A_1, B_n) \\ e(A_2, B_1) & e(A_2, B_2) & \cdots & e(A_2, B_n) \\ \vdots & \vdots & \ddots & \vdots \\ e(A_n, B_1) & e(A_n, B_2) & \cdots & e(A_n, B_n) \end{bmatrix}$$

$$\det(M) = \sum_{S: A \to B} (-1)^{t(\sigma(S))} \prod_{i=1}^{n} \omega(S_i)$$

其中 $\sum_{S:A \to B}$ 表示滿足上文要求的 $A \to B$ 的每一組 不相交路徑 S。

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- · Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i, j) \in E$, otherwise $d_{ij}=-d_{ji}.$ $\frac{rank(D)}{2}$ is the maximum matching on G.

· Cayley's Formula

- Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there are

 $\begin{array}{c} (n-2)! \\ (d_1-1)!(d_2-1)!\cdots(d_n-1)! \end{array} \text{ spanning trees.} \\ - \text{ Let } T_{n,k} \text{ be the number of labeled forests on} \\ \end{array}$ n vertices with k components, such that vertex $1, 2, \ldots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

· Erd□s-Gallai theorem

A sequence of nonnegative integers $d_1 > \cdots > d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + \cdots + d_n$ is even

and
$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$
 holds for every $1 \le k \le n$.

Gale–Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq$ $\cdots \geq a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$ holds

· Fulkerson-Chen-Anstee theorem

A sequence $(a_1, b_1), \ldots, (a_n, b_n)$ of nonnegative integer pairs with $a_1 \ge \cdots \ge a_n$ is digraphic if and only if $\sum_{i=1}^{n} a_i = \sum_{i=1}^{n} b_i$ and $\sum_{i=1}^{k} a_i \le \sum_{i=1}^{k} \min(b_i, k-1) + \sum_{i=1}^{k} a_i \le \sum_{i=1}^{k} \min(b_i, k-1)$ $\sum_{i=1}^{n} \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$

• M bius inversion formula

- Spherical cap
 - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap,
 - h: height of the cap, θ : $\arcsin(a/r)$. Volume = $\pi h^2(3r h)/3 = \pi h(3a^2 + h^2)$
 - $h^2)/6 = \pi r^3 (2 + \cos \theta) (1 \cos \theta)^2 / 3.$ Area = $2\pi rh = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos \theta)^2 / 3.$

6.21 數字

· Bernoulli numbers

$$\begin{split} B_0 - 1, B_1^{\pm} &= \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ \sum_{j=0}^m \binom{m+1}{j} B_j &= 0, \text{ EGF is } B(x) = \frac{x}{e^x - 1} = \\ \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}. \end{split}$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

 Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} &S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = \\ &S(n,n) = 1 \\ &S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n \end{split}$$

Pentagonal number theorem

 $x^n = \sum_{i=0}^n S(n,i)(x)_i$

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

· Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1), k+1$ j:s s.t. $\pi(j) \ge j, k$ j:s s.t.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{i=0}^{k} (-1)^{i} {n+1 \choose i} (k+1-j)^{n}$$

次方和

$$\sum_{k=1}^{n} k^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$\sum_{k=1}^{n} k^{4} = \frac{1}{30} (6n^{5} + 15n^{4} + 10n^{3} - n)$$

$$\sum_{k=1}^{n} k^5 = \frac{1}{12} (2n^6 + 6n^5 + 5n^4 - n^2)$$

$$\sum_{k=1}^{n} k^{6} = \frac{1}{42} (6n^{7} + 21n^{6} + 21n^{5} - 7n^{3} + n)$$

$$\sum_{k=1}^{n} k^{p} = \frac{1}{p+1} (n \sum_{i=1}^{p} (n+1)^{i} - \sum_{i=2}^{p} {i \choose p+1} \sum_{k=1}^{n} k^{p+1-i})$$

歐幾里得類算法

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity: $O(\log n)$

```
f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor
 g(a, b, c, n) = \sum_{i=0}^{n} i \lfloor \frac{ai + b}{c} \rfloor
                                        = \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} & \textbf{7.2 next-combination} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & 2 \end{cases} / \begin{cases} \textbf{Example: 1 -> 2, 4 -> 8, 12(1100) -> 17(10001)} \\ 11 \text{ next\_combination(11 comb) } \end{cases}
h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^{2}
                                                      \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \end{cases} 
                                                       +h(a \bmod c, b \bmod c, c, n)
                                                       +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                                                       +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
```

6.23 牛成函數

• Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$

```
\begin{array}{l}
-A(rx) \Rightarrow r^n a_n \\
-A(x) + B(x) \Rightarrow a_n + b_n \\
-A(x) B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\
-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}
\end{array}

- xA(x)' \Rightarrow na_n
- \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i
```

• Exponential Generating Function A(x) = $\sum_{i\geq 0} \frac{a_i}{i!} x_i$

```
\begin{array}{lll} -A(x)+B(x)\Rightarrow a_{n}+b_{n} & \textbf{7.5 rng} \\ -A^{(k)}(x)\Rightarrow a_{n+k_{n}} & \\ -A(x)B(x)\Rightarrow \sum_{i=0}^{n}\binom{n}{i}a_{i}b_{n-i} & \\ -A(x)^{k}\Rightarrow \sum_{i_{1}+i_{2}+\cdots+i_{k}=n}\binom{n}{i_{1},i_{2},\ldots,i_{k}}a_{i_{1}}a_{i} & \text{inline ull rng()} \\ -xA(x)\Rightarrow na_{n} & \text{2} & \text{static ull Q = 48} \end{array}
```

· Special Generating Function

```
 - (1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i 
 - \frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{n}{n-1} x^i
```

Misc

7.1 fast

```
#pragma GCC optimize("Ofast, no-stack-
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        protector, unroll-loops, fast-math, inline
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                              inline char gc() {
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                            static const size t sz = 65536;
= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)}{2} + \left\lfloor \frac{b}{c} \right\rfloor \cdot (n+1) & \text{3} \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \\ 0, & n < 0 \lor 6 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \\ \text{static char buf[sz];} \\
```

```
2 11 next_combination(11 comb) {
x = 11 \times x = comb \times -comb, y = comb + x;
   return ((comb & ~y) / x >> 1) | y;
```

7.3 PBDS

```
i #include <ext/pb_ds/assoc_container.hpp>
                                 2 using namespace gnu pbds;
nm(m+1)-2g(c,c-b-1,a,m-1) 3 tree-tl, null_type, less<11>, rb_tree_tag,
-2f(c,c-b-1,a,m-1)-f(a,b,c,n),
                                      tree_order_statistics_node_update> st;
                                   // find_by_order order_of_key
                                  // __float128_t
                                   for(int i = bs._Find_first(); i < bs.size();</pre>
                                        i = bs._Find_next(i));
```

7.4 python

```
1 from decimal import Decimal, getcontext
getcontext().prec = 1000000000
a = pow(Decimal(2), 82589933) - 1
```

7.5 rng

```
static ull Q = 48763;
0 ^= 0 << 7;
Q ^= Q >> 9;
return Q & 0xFFFFFFFFULL;
```

7.6 rotate90

```
i vector<vector<T>> rotate90(const vector
      vector<T>>& a) {
     int n = sz(a), m = sz(a[0]);
     vector<vector<T>> b(m, vector<T>(n));
     REP(i, n) REP(j, m) b[j][i] = a[i][m - 1]
     return b;
```

timer

```
1 clock t T1 = clock();
2 double getCurrentTime() { return (double) (
      clock() - T1) / CLOCKS PER SEC; }
```

String

8.1 AC

```
1 template < int ALPHABET = 26, char MIN CHAR =</pre>
      'a'>
 struct ac_automaton {
   struct Node {
     int fail = 0, cnt = 0;
     array<int, ALPHABET> go{};
   vector<Node> node;
   vi que;
   int new node() { return node.eb(), SZ(node
       ) - 1; }
   Node& operator[](int i) { return node[i];
   ac_automaton() { new_node(); }
   int insert(const string& s) {
     int p = 0:
     for(char c : s) {
       int v = c - MIN_CHAR;
       if(node[p].go[v] == 0) node[p].go[v] =
            new node();
       p = node[p].go[v];
     node[p].cnt++;
     return p;
   void build() {
     que.reserve(SZ(node)); que.pb(0);
     REP(i, SZ(que)) {
       int u = que[i];
       REP(j, ALPHABET) {
         if(node[u].go[j] == 0) node[u].go[j]
               = node[node[u].fail].go[j];
           int v = node[u].go[j];
           node[v].fail = (u == 0 ? u : node[
                node[u].fail].go[j]);
```

```
que.pb(v);
35
```

8.2 KMP

```
ı|// abacbaba -> [0, 0, 1, 0, 0, 1, 2, 3]
vi KMP(const vi& a) {
   int n = SZ(a);
   vi k(n):
   for(int i = 1; i < n; ++i) {</pre>
     int j = k[i - 1];
     while(j > 0 && a[i] != a[j]) j = k[j -
     j += (a[i] == a[j]);
     k[i] = j;
   return k;
```

8.3 LCP

```
i vi lcp(const vi& s, const vi& sa) {
    int n = SZ(s), h = 0;
    vi rnk(n), lcp(n - 1);
    REP(i, n) rnk[sa[i]] = i;
    REP(i, n) {
      h -= (h > 0);
      if(rnk[i] == 0) continue;
      int j = sa[rnk[i] - 1];
      for(; j + h < n && i + h < n; h++) if(s[</pre>
         j + h] != s[i + h]) break;
      lcp[rnk[i] - 1] = h;
11
   }
12
   return lcp;
```

8.4 manacher

```
ı|// length: (z[i] - (i & 1)) / 2 * 2 + (i &
      1)
vi manacher(string t) {
    string s = "\&";
    for(char c : t) s.pb(c), s.pb('%');
    int 1 = 0, r = 0;
    vi z(sz(s));
    REP(i, sz(s)) {
      z[i] = r > i ? min(z[2 * 1 - i], r - i)
      while(s[i + z[i]] == s[i - z[i]]) z[i
      if(z[i] + i > r) r = z[i] + 1, l = i;
11
   return z;
```

8.5 rolling-hash

```
const ll M = 911382323, mod = 972663749;
ll Get(vector<ll>% h, int l, int r) {
    if(!l) return h[r]; // p[i] = M^i % mod
    ll ans = (h[r] - h[l - 1] * p[r - 1 + 1])
        % mod;
    return (ans + mod) % mod;
}

vector<ll> Hash(string s) {
    vector<ll> Hash(string s) {
    vector<ll> ans[0] = s[0];
    for(int i = 1; i < sz(s); i++) ans[i] = (
        ans[i - 1] * M + s[i]) % mod;
    return ans;
}</pre>
```

8.6 SAIS

```
1 // mississippi
2 // 10 7 4 1 0 9 8 6 3 5 2
3 vi SAIS(string a) {
   #define QQ(i, n) for(int i = (n); i >= 0;
   int n = sz(a), m = *max element(all(a)) +
   vi pos(m + 1), x(m), sa(n), val(n), lms;
   for(auto c : a) pos[c + 1]++;
   REP(i, m) pos[i + 1] += pos[i];
   vector<bool> s(n);
   QQ(i, n - 2) s[i] = a[i] != a[i + 1] ? a[i]
        ] < a[i + 1] : s[i + 1];
   auto ind = [&](const vi& ls){
     fill(all(sa), -1);
     auto L = [\&](int i) \{ if(i >= 0 \&\& !s[i]) \}
          ]) sa[x[a[i]]++] = i; };
     auto S = [&](int i) { if(i >= 0 && s[i])
           sa[--x[a[i]]] = i; };
     REP(i, m) x[i] = pos[i + 1];
     QQ(i, sz(ls) - 1) S(ls[i]);
     REP(i, m) x[i] = pos[i];
     L(n - 1);
     REP(i, n) L(sa[i] - 1);
     REP(i, m) x[i] = pos[i + 1];
     QQ(i, n - 1) S(sa[i] - 1);
   auto ok = [&](int i) { return i == n || (! 32|)
        s[i - 1] && s[i]); };
   auto same = [&](int i,int j) {
     do {
       if(a[i++] != a[j++]) return false;
     } while(!ok(i) && !ok(j));
     return ok(i) && ok(j);
   for (int i = 1; i < n; i++) if (ok(i)) lms.
        pb(i);
   ind(lms);
   if(sz(lms)) {
     int p = -1, w = 0;
     for(auto v : sa) if(v && ok(v)) {
       if(p != -1 && same(p, v)) w--;
       val[p = v] = w++;
```

```
}
auto b = lms;
for(auto& v : b) v = val[v];
b = SAIS(b);
for(auto& v : b) v = lms[v];
ind(b);
}
return sa;
}
```

8.7 SAM

```
ı|// cnt 要先用 bfs 往回推, 第一次出現的位置是
       state.first_pos - |S| + 1
 struct Node { int go[26], len, link, cnt,
      first_pos; };
 Node SA[N]; int sz;
 void sa_init() { SA[0].link = -1, SA[0].len
      = 0, sz = 1; 
 int sa_extend(int p, int c) {
   int u = sz++;
   SA[u].first_pos = SA[u].len = SA[p].len +
   SA[u].cnt = 1;
   while(p != -1 && SA[p].go[c] == 0) {
     SA[p].go[c] = u;
     p = SA[p].link;
   if(p == -1) {
     SA[u].link = 0;
     return u;
   int q = SA[p].go[c];
   if(SA[p].len + 1 == SA[q].len) {
     SA[u].link = q;
     return u;
   int x = sz++;
   SA[x] = SA[q];
   SA[x].cnt = 0;
   SA[x].len = SA[p].len + 1;
   SA[q].link = SA[u].link = x;
   while(p != -1 && SA[p].go[c] == q) {
     SA[p].go[c] = x;
     p = SA[p].link;
   return u;
```

8.8 smallest-rotation

```
string small_rot(string s) {
   int n = sz(s), i = 0, j = 1;
   s += s;

while(i < n && j < n) {
   int k = 0;
   while(k < n && s[i + k] == s[j + k]) k
   ++;

if(s[i + k] <= s[j + k]) j += k + 1;
   else i += k + 1;</pre>
```

```
9     if(i == j) j++;
10     }
11     int ans = i < n ? i : j;
12     return s.substr(ans, n);
13 }</pre>
```

8.9 Z

```
1  // abacbaba -> [0, 0, 1, 0, 0, 3, 0, 1]
2  vi z_algorithm(const vi& a) {
3   int n = sz(a);
4  vi z(n);
5  for(int i = 1, j = 0; i < n; ++i) {
6   if(i <= j + z[j]) z[i] = min(z[i - j], j + z[j] - i);
7  while(i + z[i] < n && a[i + z[i]] == a[z [i]]) z[i]++;
8  if(i + z[i] > j + z[j]) j = i;
9  }
10  return z;
11 }
```

ACM ICPC Judge Test NTHU SplayTreap

C++ Resource Test

```
#include <bits/stdc++.h>
using namespace std;

namespace system_test {

const size_t KB = 1024;
const size_t MB = KB * 1024;
const size_t GB = MB * 1024;
```

```
chrono::duration<double> diff = end -
10 size t block size, bound;
                                                          begin;
  void stack size dfs(size t depth = 1) {
                                                     return diff.count():
   if (depth >= bound)
                                                   void runtime_error_1() {
    int8_t ptr[block_size]; // 若無法編譯將
                                                     // Segmentation fault
         block size 改成常數
                                                     int *ptr = nullptr;
    memset(ptr, 'a', block_size);
                                                     *(ptr + 7122) = 7122;
    cout << depth << endl;</pre>
                                                 42 }
    stack_size_dfs(depth + 1);
                                                   void runtime_error_2() {
                                                     // Segmentation fault
  void stack_size_and_runtime_error(size_t
                                                     int *ptr = (int *)memset;
       block size, size t bound = 1024) {
                                                     *ptr = 7122;
    system test::block size = block size;
                                                 48
    system_test::bound = bound;
    stack size dfs();
                                                   void runtime_error_3() {
                                                     // munmap_chunk(): invalid pointer
                                                     int *ptr = (int *)memset;
  double speed(int iter num) {
                                                     delete ptr;
    const int block_size = 1024;
                                                 54
    volatile int A[block_size];
    auto begin = chrono::high resolution clock
                                                   void runtime_error_4() {
         ::now();
                                                     // free(): invalid pointer
    while (iter_num--)
                                                     int *ptr = new int[7122];
      for (int j = 0; j < block_size; ++j)</pre>
                                                     ptr += 1;
                                                     delete[] ptr;
    auto end = chrono::high resolution clock::
         now();
```

```
63 void runtime error 5() {
    // maybe illegal instruction
    int a = 7122, b = 0;
    cout << (a / b) << endl;</pre>
  void runtime error 6() {
    // floating point exception
    volatile int a = 7122, b = 0;
    cout << (a / b) << endl;
73 }
  void runtime error 7() {
    // call to abort.
    assert(false);
78 }
  } // namespace system test
82 #include <sys/resource.h>
void print_stack_limit() { // only work in
       Linux
    struct rlimit 1;
    getrlimit(RLIMIT STACK, &1);
    cout << "stack_size = " << l.rlim_cur << "</pre>
          byte" << endl;</pre>
87 }
```