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#### Basic

#### 1.1 template

```
| #pragma GCC optimize("Ofast, no-stack-
      protector, unroll-loops, fast-math, inline"
2 #define FOR(i, begin, end) for(int i = (
      begin), i## end = (end); i < i## end;
#define IFOR(i, begin, end) for(int i = (end
     ) - 1, i## begin = (begin); i >= i##
      _begin_; i--)
4 #define REP(i, n) FOR(i, 0, n)
5 #define IREP(i, n) IFOR(i, 0, n)
```

#### 1.2 vimre

```
| se nu ai hls et ru ic is sc cul
2 se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
3 syntax on
4 hi cursorline cterm=none ctermbg=89
5 set bg=dark
6 inoremap {<CR> {<CR>}<Esc>ko<tab>
```

#### **Data-Structure**

#### 2.1 wavelet-tree

```
i template < class T>
 struct wavelet tree {
   int n, log;
   vector<T> vals;
   vi sums;
   vector<ull> bits:
   void set_bit(int i, ull v) { bits[i >> 6]
        |= (v << (i & 63)); }
   int get_sum(int i) const { return sums[i
        >> 6] + __builtin_popcountll(bits[i >>
         6] & ((\overline{1}ULL << (\overline{i} \& 63)) - 1)); }
   wavelet tree(const vector<T>& v) : n(SZ(
        v)) {
     vals = sort_unique(_v);
     log = __lg(2 * vals.size() - 1);
     bits.resize((log * n + 64) >> 6, 0ULL);
     sums.resize(SZ(bits), 0);
     vi v(SZ(_v)), cnt(SZ(vals) + 1);
     REP(i, SZ(v)) {
       v[i] = lower_bound(ALL(vals), _v[i]) -
             vals.begin();
       cnt[v[i] + 1] += 1;
     partial sum(ALL(cnt) - 1, cnt.begin());
     REP(j, log) {
       for(int i : v) {
```

#### lazysegtree

```
template < class S,
         S (*e)(),
         S (*op)(S, S),
         class F.
         F (*id)(),
         S (*mapping)(F, S),
         F (*composition)(F, F)>
struct lazy_segtree {
 int n, size, log;
  vector<S> d; vector<F> lz;
  void update(int k) { d[k] = op(d[k << 1],
       d[k << 1 | 1]); }
  void all_apply(int k, F f) {
    d[k] = mapping(f, d[k]);
    if(k < size) lz[k] = composition(f, lz[k</pre>
         ]);
 void push(int k) {
    all_apply(k << 1, lz[k]);
    all_apply(k << 1 | 1, lz[k]);
   lz[k] = id();
```

int tmp =  $i \gg (\log - 1 - j);$ 

FOR(i, 1, SZ(sums)) sums[i] = sums[i -

cnt[pos]++;

1]);

j++) {

cnt\_a);

res <<= 1;

} else {

if(ab zeros > k) {

k -= ab\_zeros;

(log - j)]--;

T get kth(int a, int b, int k) {

if(j == log) return vals[res];

int cnt\_ia = get\_sum(n \* j + ia);

int cnt\_ib = get\_sum(n \* j + ib);

int ab\_zeros = (b - a) - (cnt\_b -

int cnt a = get sum(n \* j + a);

int cnt\_b = get\_sum(n \* j + b);

ib -= cnt\_ib - cnt\_ia;

a -= cnt a - cnt ia;

b -= cnt b - cnt ia;

res = (res << 1) | 1;

int pos = (tmp >> 1) << (log - j);</pre>

set bit(j \* n + cnt[pos], tmp & 1);

for(int i : v) cnt[(i >> (log - j)) << 24

1] + \_\_builtin\_popcountll(bits[i -

for(int j = 0, ia = 0, ib = n, res = 0;; 33

ia += (ib - ia) - (cnt\_ib - cnt\_ia);

a += (ib - a) - (cnt ib - cnt a);

b += (ib - b) - (cnt\_ib - cnt\_b);

28

38

39

52

53

```
lazy_segtree(int _n) : lazy_segtree(vector 81|
     <S>( n, e())) {}
lazy segtree(const vector<S>& v) : n(SZ(v) 82
  log = __lg(2 * n - 1), size = 1 << log;
                                                   template < class G> int max_right(int 1, G g
                                              84
  d.resize(size * 2, e());
  lz.resize(size, id());
  REP(i, n) d[size + i] = v[i];
  for(int i = size - 1; i; i--) update(i); 87
void set(int p, S x) {
  p += size;
  for(int i = log; i; --i) push(p >> i);
  d[p] = x:
  for(int i = 1; i <= log; ++i) update(p</pre>
       >> i);
                                              95
S get(int p) {
  p += size;
  for(int i = log; i; i--) push(p >> i);
  return d[p];
S prod(int 1, int r) {
                                              100
  if(1 == r) return e();
                                              101
  1 += size; r += size;
                                              102
  for(int i = log; i; i--) {
    if(((1 >> i) << i) != 1) push(1 >> i); 104
    if(((r >> i) << i) != r) push(r >> i);
  S sml = e(), smr = e();
                                              106
  while(1 < r) {</pre>
                                              107
    if(1 \& 1) sml = op(sml, d[1++]);
    if(r & 1) smr = op(d[--r], smr);
    1 >>= 1, r >>= 1;
                                              109
                                              110
  return op(sml, smr);
                                              111
                                              112
S all_prod() const { return d[1]; }
                                              113
void apply(int p, F f) {
                                              114
  p += size;
                                              115
  for(int i = log; i; i--) push(p >> i);
                                              116
  d[p] = mapping(f, d[p]);
                                              117
  for(int i = 1; i <= log; i++) update(p</pre>
       >> i):
void apply(int 1, int r, F f) {
                                              120
  if(1 == r) return;
                                              121
  1 += size; r += size;
                                              122
  for(int i = log; i; i--) {
    if(((1 >> i) << i) != 1) push(1 >> i); 124
    if(((r >> i) << i) != r) push((r - 1) | 125 | );
         >> i);
```

int 12 = 1, r2 = r;

1 >>= 1, r >>= 1;

if(1 & 1) all apply(1++, f);

if(r & 1) all\_apply(--r, f);

if(((1 >> i) << i) != 1) update(1 >> i

for(int i = 1; i <= log; i++) {</pre>

while(1 < r) {

1 = 12;

r = r2;

```
assert(0 <= 1 && 1 <= n && g(e()));
  if(1 == n) return n:
 1 += size;
  for(int i = log; i; i--) push(1 >> i);
 S sm = e();
    while(!(1 & 1)) 1 >>= 1;
    if(!g(op(sm, d[1]))) {
      while(1 < size) {</pre>
        push(1);
        1 <<= 1:
        if(g(op(sm, d[1]))) sm = op(sm, d[
             1++1);
      return 1 - size;
    sm = op(sm, d[1++]);
 } while((1 & -1) != 1);
 return n:
template < class G> int min_left(int r, G g)
  assert(0 <= r \& r <= n \& \& g(e()));
 if(r == 0) return 0;
  r += size;
  for(int i = log; i >= 1; i--) push((r -
      1) >> i);
 S sm = e();
  do {
    while(r > 1 && (r & 1)) r >>= 1;
    if(!g(op(d[r], sm))) {
      while(r < size) {</pre>
        push(r);
        r = r << 1 | 1;
        if(g(op(d[r], sm))) sm = op(d[r])
             --], sm);
      return r + 1 - size;
    sm = op(d[r], sm);
  } while((r & -r) != r);
  return 0:
```

if(((r >> i) << i) != r) update((r -</pre>

1) >> i);

#### 2.3 LiChao

```
1 struct LiChao { // min
   int n;
   vector<pll> seg;
   LiChao(int _n) : n(_n) {
     seg.assign(4 * n + 5, pll(0, INF));
   11 cal(pll line, ll x) { return line.F * x
         + line.S; }
```

```
void insert(int 1, int r, int id, pll line 26
  if(1 == r) {
    if(cal(line, 1) < cal(seg[id], 1)) seg 29</pre>
         [id] = line;
    return:
  int mid = (1 + r) / 2:
 if(line.F > seg[id].F) swap(line, seg[id
  if(cal(line, mid) <= cal(seg[id], mid))</pre>
    seg[id] = line;
    insert(1, mid, id * 2, seg[id]);
  else insert(mid + 1, r, id * 2 + 1, line
ll query(int 1, int r, int id, ll x) {
  if(x < 1 \mid | x > r) return INF;
 if(l == r) return cal(seg[id], x);
 int mid = (1 + r) / 2;
 11 \text{ val} = 0;
 if(x \le mid) val = query(1, mid, id * 2,
  else val = query(mid + 1, r, id * 2 + 1,
  return min(val, cal(seg[id], x));
```

```
1 struct DLX {
   int n, m, tot, ans;
   vi first, siz, L, R, U, D, col, row, stk;
   DLX(int n, int m) : n(n), m(m), tot(m)
     int sz = n * m;
     first = siz = L = R = U = D = col = row
          = stk = vi(sz);
     REP(i, m + 1) {
       L[i] = i - 1, R[i] = i + 1;
       U[i] = D[i] = i;
     L[0] = m, R[m] = 0;
   void insert(int r, int c) { // (r, c) is 1
     col[++tot] = c, row[tot] = r, ++siz[c];
     D[tot] = D[c], U[D[c]] = tot, U[tot] = c
          , D[c] = tot;
     if(!first[r]) first[r] = L[tot] = R[tot]
           = tot:
     else {
       L[R[tot] = R[first[r]]] = tot;
       R[L[tot] = first[r]] = tot;
   #define TRAV(i, X, j) for(i = X[j]; i != j
        ; i = X[i])
   void remove(int c) {
     int i, j;
```

2.4 DLX

```
2.5 sparse-table
1 template < class T, T (*op)(T, T)>
2 struct sparse_table {
                                                  32
    int n;
    vector<vector<T>> b;
    sparse_table(const vector<T>& a) : n(SZ(a) 34
      int \lg = \lg(n) + 1;
      b.resize(lg); b[0] = a;
                                                  36
      FOR(j, 1, lg) {
        b[j].resize(n - (1 << j) + 1);
        REP(i, n - (1 << j) + 1) b[j][i] = op( _{38}
             b[j - 1][i], b[j - 1][i + (1 << (j - 1)[i])
              - 1))]);
    T prod(int from, int to) {
      int \lg = \lg(to - from + 1);
      return op(b[lg][from], b[lg][to - (1 <<</pre>
           lg) + 1]);
17 };
```

L[R[c]] = L[c], R[L[c]] = R[c];

TRAV(i, D, c) TRAV(j, R, i) {

D[U[D[j]] = U[j]] = D[j];

TRAV(i, U, c) TRAV(j, L, i) {

if(!R[0]) return ans = dep, true;

TRAV(j, R, i) remove(col[j]);

if(dance(dep + 1)) return true;

TRAV(j, L, i) recover(col[j]);

TRAV(i, R, 0) if(siz[i] < siz[c]) c = i;

return vi(stk.begin() + 1, stk.begin() +

13

U[D[j]] = D[U[j]] = j;

L[R[c]] = R[L[c]] = c;

siz[col[j]]--;

void recover(int c) {

siz[col[j]]++;

bool dance(int dep) {

TRAV(i, D, c) {

int i, j, c = R[0];

stk[dep] = row[i];

if(!dance(1)) return {};

int i, j;

remove(c):

recover(c);

vi solve() {

return false:

ans);

2.6 static-range-lis

```
1 #define MEM(a, x, n) memset(a, x, sizeof(int 50
      ) * n)
2 using I = int*;
3 struct static_range_lis {
   int n, ps = 0;
   I invp, res monge, pool;
    vector<vector<pii>>> gry;
    static_range_lis(vi a) : n(SZ(a)), qry(n + 58
         1) {
      // a must be permutation of [0, n)
      pool = (I) malloc(sizeof(int) * n * 100) 60
      invp = A(n), res monge = A(n);
      REP(i, n) invp[a[i]] = i;
   inline I A(int x) { return pool + (ps += x
        ) - x; }
    void add query(int 1, int r) { gry[1].pb({ 64
        r, SZ(ans)}), ans.pb(r - 1); }
    void unit_monge_mult(I a, I b, I r, int n) 65
      if(n == 2){
        if(!a[0] && !b[0]) r[0] = 0, r[1] = 1; 68
        else r[0] = 1, r[1] = 0;
        return;
      if(n == 1) return r[0] = 0, void();
      int lps = ps, d = n / 2;
                                                 72
     I a1 = A(d), a2 = A(n - d), b1 = A(d),
          b2 = A(n - d);
      I mpa1 = A(d), mpa2 = A(n - d), mpb1 = A
          (d), mpb2 = A(n - d);
                                                 75
      int p[2] = {};
                                                 76
      REP(i, n) {
        if(a[i] < d) a1[p[0]] = a[i], mpa1[p</pre>
             [0]++] = i;
        else a2[p[1]] = a[i] - d, mpa2[p[1]++]
             = i;
      p[0] = p[1] = 0;
      REP(i, n) {
        if(b[i] < d) b1[p[0]] = b[i], mpb1[p</pre>
             [0]++] = i;
        else b2[p[1]] = b[i] - d, mpb2[p[1]++]
             = i;
      I c1 = A(d), c2 = A(n - d);
      unit monge mult(a1, b1, c1, d),
          unit_monge_mult(a2, b2, c2, n - d);
      I cpx = A(n), cpy = A(n), cqx = A(n),
          cqy = A(n);
      REP(i, d) cpx[mpa1[i]] = mpb1[c1[i]],
          cpy[mpa1[i]]=0;
      REP(i, n - d) cpx[mpa2[i]] = mpb2[c2[i]]
          ]], cpy[mpa2[i]]=1;
      REP(i, n) r[i] = cpx[i];
      REP(i, n) cqx[cpx[i]] = i, cqy[cpx[i]] =
            cpy[i];
      int hi = n, lo = n, his = 0, los = 0;
      REP(i, n)
        if(cqy[i] ^ (cqx[i] >= hi)) his--;
                                                 98
        while(hi > 0 && his < 0) {
                                                 99
         hi--;
                                                100
          if(cpy[hi] ^ (cpx[hi] > i)) his++;
```

```
while(lo > 0 && los <= 0) {
      if(cpy[lo] ^ (cpx[lo] >= i)) los++;
    if(los > 0 \&\& hi == lo) r[lo] = i;
    if(cqy[i] \land (cqx[i] >= lo)) los--;
  ps = 1ps;
void subunit_monge_mult(I a, I b, I c, int
  int lps = ps;
  I za = A(n), zb = A(n), res = A(n), vis
      = A(n), mpa = A(n), mpb = A(n), rb =
        A(n);
  MEM(vis, 0, n), MEM(mpa, -1, n), MEM(mpb
       , -1, n), MEM(rb, -1, n);
  int ca = n;
  IREP(i, n) if(a[i] != -1) vis[a[i]] = 1,
        za[--ca] = a[i], mpa[ca] = i;
  IREP(i, n) if(!vis[i]) za[--ca] = i;
  MEM(vis, -1, n);
  REP(i, n) if(b[i] != -1) vis[b[i]] = i;
  ca = 0:
  REP(i, n) if(vis[i] != -1) mpb[ca] = i,
       rb[vis[i]] = ca++;
  REP(i, n) if(rb[i] == -1) rb[i] = ca++;
  REP(i, n) zb[rb[i]] = i;
  unit_monge_mult(za, zb, res, n);
  MEM(c, -1, n);
  REP(i, n) if(mpa[i] != -1 && mpb[res[i]]
        != -1) c[mpa[i]] = mpb[res[i]];
  ps = lps;
void solve(I p, I ret, int n) {
  if(n == 1) return ret[0] = -1, void();
  int lps = ps, d = n / 2;
  I pl = A(d), pr = A(n - d);
  REP(i, d) pl[i] = p[i];
  REP(i, n - d) pr[i] = p[i + d];
  I vis = A(n); MEM(vis, -1, n);
  REP(i, d) vis[pl[i]] = i;
  I tl = A(d), tr = A(n - d), mpl = A(d),
       mpr = A(n - d);
  int ca = 0;
  REP(i, n) if(vis[i] != -1) mpl[ca] = i,
      tl[vis[i]] = ca++;
  ca = 0; MEM(vis, -1, n);
  REP(i, n - d) vis[pr[i]] = i;
  REP(i, n) if(vis[i] != -1) mpr[ca] = i,
       tr[vis[i]] = ca++;
  I vl = A(d), vr = A(n - d);
  solve(tl, vl, d), solve(tr, vr, n - d);
  I sl = A(n), sr = A(n);
  iota(sl, sl + n, 0); iota(sr, sr + n, 0)
  REP(i, d) sl[mpl[i]] = (vl[i] == -1 ? -1
        : mpl[vl[i]]);
  REP(i, n - d) sr[mpr[i]] = (vr[i] == -1)
      ? -1 : mpr[vr[i]]);
  subunit_monge_mult(sl, sr, ret, n);
  ps = lps;
vi solve() {
  solve(invp, res monge, n);
  vi fenw(n + 1);
```

attach(p, rgt, (rgt ? v->l : v->r));

if(!rt) attach(g, (g->r == p), v);

if(!p->is root()) push(g);

void all\_apply(Node\* v, F f) {

mapping(f, v->sum);

 $v \rightarrow lz = composition(f, v \rightarrow lz);$ 

if(!p->is\_root()) rotate((g->r == p)

== (p->r == v) ? p : v);

v->val = mapping(f, v->val), v->sum =

if(v->l != nullptr) all apply(v->l, v

if(v->r != nullptr) all apply(v->r, v

**if**(v->l != nullptr) v->l->rev ^= 1;

if(v->r != nullptr) v->r->rev ^= 1;

v->sum = reversal(v->sum);

 $v \rightarrow sum = op(v \rightarrow 1 \rightarrow sum, v \rightarrow sum)$ :

 $v \rightarrow sum = op(v \rightarrow sum, v \rightarrow r \rightarrow sum);$ 

attach(v, !rgt, p);

else  $v \rightarrow p = g$ ;

push(v);

void splay(Node\* v) {

auto p = v - p;

auto g = p->p;

rotate(v);

void push(Node\* v) {

**if**(v->1z != id()) {

->1z):

swap(v->1, v->r);

v->rev = false;

void pull(Node\* v) {

 $v \rightarrow sum = v \rightarrow val;$ 

push(v->1);

push(v->r);

if(v->l != nullptr) {

 $v \rightarrow sz += v \rightarrow 1 \rightarrow sz;$ 

if(v->r != nullptr) {

 $v \rightarrow sz += v \rightarrow r \rightarrow sz;$ 

 $v \rightarrow sz = 1;$ 

 $v \rightarrow lz = id();$ 

if(v->rev) {

while(!v->is root()) {

push(p), push(v);

```
IREP(i, n) {
                                                       if(res monge[i] != -1) {
                                                            , a[i]) - _a.begin();
           for(int p = res monge[i] + 1; p <= n</pre>
                                                       sz = SZ(a);
               ; p += p & -p) fenw[p]++;
                                                     void add_query(int 1, int r) { L.push_back 15
         for(auto& z : qry[i]){
                                                          (1), R.push back(r); }
           auto [id, c] = z;
                                                     vector<ll> solve() {
           for(int p = id; p; p -= p & -p) ans[
                                                       const int q = SZ(L);
                                                       const int B = max(1.0, SZ(a) / sqrt(q)); 18
               c] -= fenw[p];
110
                                                       vi ord(q);
111
                                                       iota(ALL(ord), 0);
112
       free(pool);
                                                       sort(ALL(ord), [&](int i, int j) {
                                                         if(L[i] / B == L[j] / B) {
      return ans;
113
                                                           return L[i] / B & 1 ? R[i] > R[j] :
114
115 };
                                                               R[i] < R[j];
                                                        return L[i] < L[j];</pre>
                                                       ans.resize(q);
         rollback-dsu
                                                       fenwick<ll> fenw(sz + 1);
                                                       11 cnt = 0;
                                                       auto AL = [&](int i) {
 1 struct RollbackDSU {
                                                         cnt += fenw.sum(0, a[i] - 1);
    int n; vi sz, tag;
                                                         fenw.add(a[i], +1);
    vector<tuple<int, int, int, int>> op;
    void init(int _n) {
                                                       auto AR = [&](int i) {
      n = n;
                                                         cnt += fenw.sum(a[i] + 1, sz);
      sz.assign(n, -1);
                                                         fenw.add(a[i], +1);
      tag.clear();
                                                       auto DL = [&](int i) {
     int leader(int x) {
                                                         cnt -= fenw.sum(0, a[i] - 1);
      while(sz[x] >= 0) x = sz[x];
                                                         fenw.add(a[i], -1);
      return x;
                                                       auto DR = [&](int i) {
     bool merge(int x, int y) {
                                                         cnt -= fenw.sum(a[i] + 1, sz);
      x = leader(x), y = leader(y);
                                                         fenw.add(a[i], -1);
      if(x == y) return false;
      if(-sz[x] < -sz[y]) swap(x, y);
                                                       int 1 = 0, r = 0;
      op.eb(x, sz[x], y, sz[y]);
                                                       REP(i, q) {
      sz[x] += sz[y]; sz[y] = x;
                                                         int id = ord[i], ql = L[id], qr = R[id
      return true:
```

#### 2.9 LCT

## 2.8 static-range-inversion

```
i struct static_range_inversion {
   int sz:
   vi a, L, R;
   vector<ll> ans:
   static range inversion(vi a) : a( a) {
     _a = sort_unique(_a);
```

int size(int x) { return -sz[leader(x);] }

auto [x, sx, y, sy] = op.back(); op.

void add\_tag() { tag.pb(sz(op)); }

int z = tag.back(); tag.ppb();

void rollback() {

while(sz(op) > z) {

ppb();

sz[x] = sx;

sz[y] = sy;

```
template < class S,
         S (*e)(),
         S (*op)(S, S),
         S (*reversal)(S),
         class F,
         F (*id)(),
         S (*mapping)(F, S),
         F (*composition)(F, F)>
struct lazy lct {
 struct Node {
```

while(1 > q1) AL(--1);

while(r < qr) AR(r++);</pre>

while(1 < q1) DL(1++);

while(r > qr) DR(--r);

ans[id] = cnt;

return ans;

```
S val = e(), sum = e();
  F lz = id();
  bool rev = false:
  int sz = 1;
                                             60
  Node *1 = nullptr, *r = nullptr, *p =
                                             61
       nullptr:
                                             62
  Node() {}
                                             63
  Node(const S& s) : val(s), sum(s) {}
  bool is root() const { return p ==
                                             65
       nullptr || (p->l != this && p->r !=
       this): }
};
                                             68
int n;
                                             69
vector<Node> a:
lazy_lct() : n(0) {}
explicit lazy lct(int n) : lazy lct(
                                             71
     vector<S>( n, e())) {}
                                             72
explicit lazy_lct(const vector<S>& v) : n(
     SZ(v)) { REP(i, n) a.eb(v[i]); }
Node* access(int u) {
  Node* v = &a[u];
  Node* last = nullptr:
  for(Node* p = v; p != nullptr; p = p->p)
        splay(p), p->r = last, pull(last =
  splay(v);
  return last:
void make_root(int u) { access(u), a[u].
     rev ^= 1, push(&a[u]); }
                                             82
void link(int u, int v) { make_root(v), a[
                                            83
     vl.p = &a[u]; }
void cut(int u) {
  access(u);
  if(a[u].1 != nullptr) a[u].1->p =
       nullptr, a[u].l = nullptr, pull(&a[u 88
void cut(int u, int v) { make_root(u), cut
bool is_connected(int u, int v) {
  if(u == v) return true;
  return access(u), access(v), a[u].p !=
       nullptr:
int get lca(int u, int v) { return access( 98
     u), access(v) - &a[0]; }
void set(int u, const S& s) { access(u), a
     [u].val = s, pull(&a[u]); }
                                            101
S get(int u) { return access(u), a[u].val; 102
void apply(int u, int v, const F& f) {
     make root(u), access(v), all apply(&a[ 105 ];
     v], f), push(&a[v]); }
S prod(int u, int v) { return make root(u)
     , access(v), a[v].sum; }
void rotate(Node* v) {
  auto attach = [&](Node* p, bool side,
       Node* c) {
    (side ? p->r : p->1) = c;
    if(c != nullptr) c->p = p;
```

Node \*p =  $v \rightarrow p$ , \*g =  $p \rightarrow p$ ;

bool rgt = (p->r == v);

bool rt = p->is root();

30

31

#### 2.10 segtree-beats

```
1 struct segtree_beats {
   static constexpr 11 INF = numeric limits<</pre>
         ll>::max() / 2.1;
    struct alignas(32) Node {
      11 \text{ sum} = 0, g1 = 0, 11 = 0;
      11 g2 = -INF, gc = 1, 12 = INF, 1c = 1,
           add = 0;
```

```
};
11 n, log;
vector<Node> v:
segtree_beats() {}
segtree_beats(int _n) : segtree_beats(
    vector<11>(_n)) {}
segtree beats(const vector<11>& vc) {
  n = 1, log = 0;
  while(n < SZ(vc)) n <<= 1, log++;</pre>
  v.resize(2 * n);
  REP(i, SZ(vc)) v[i + n].sum = v[i + n].
       g1 = v[i + n].l1 = vc[i];
  for(ll i = n - 1; i; --i) update(i);
void range_chmin(int 1, int r, 11 x) {
     inner apply<1>(1, r, x); }
void range_chmax(int 1, int r, 11 x) {
     inner_apply<2>(1, r, x); }
void range_add(int 1, int r, 11 x) {
     inner_apply<3>(1, r, x); }
void range_update(int 1, int r, 11 x) {
     inner_apply<4>(1, r, x); }
11 range_min(int 1, int r) { return
    inner_fold<1>(l, r); }
11 range max(int 1, int r) { return
    inner_fold<2>(1, r); }
11 range sum(int 1, int r) { return
     inner fold<3>(1, r);}
void update(int k) {
  Node& p = v[k];
  Node& 1 = v[k * 2];
  Node& r = v[k * 2 + 1];
  p.sum = 1.sum + r.sum;
  if(l.g1 == r.g1) {
    p.g1 = 1.g1;
    p.g2 = max(1.g2, r.g2);
    p.gc = 1.gc + r.gc;
  } else {
    bool f = 1.g1 > r.g1;
    p.g1 = f ? 1.g1 : r.g1;
    p.gc = f ? 1.gc : r.gc;
    p.g2 = max(f ? r.g1 : l.g1, f ? l.g2 :
          r.g2);
  if(1.11 == r.11) {
    p.11 = 1.11;
    p.12 = min(1.12, r.12);
    p.lc = 1.lc + r.lc;
  } else {
    bool f = 1.11 < r.11;</pre>
    p.11 = f ? 1.11 : r.11;
    p.lc = f ? 1.lc : r.lc;
    p.12 = min(f ? r.11 : 1.11, f ? 1.12 : 106
          r.12):
void push add(int k, ll x) {
  Node& p = v[k];
  p.sum += x << (log + __builtin_clz(k) -</pre>
       31):
  p.g1 += x, p.11 += x;
                                            112
  if(p.g2 != -INF) p.g2 += x;
  if(p.12 != INF) p.12 += x;
  p.add += x;
void push_min(int k, ll x) {
```

```
Node& p = v[k];
 p.sum += (x - p.g1) * p.gc;
                                             118
  if(p.l1 == p.g1) p.l1 = x;
                                             119
 if(p.12 == p.g1) p.12 = x;
 p.g1 = x:
                                             121
void push max(int k, ll x) {
                                             122
 Node& p = v[k]:
                                             123
 p.sum += (x - p.11) * p.1c;
                                             124
 if(p.g1 == p.11) p.g1 = x;
                                             125
 if(p.g2 == p.11) p.g2 = x;
 p.11 = x;
                                             127
                                             128
void push(int k) {
                                             129
  Node& p = v[k];
                                             130
  if(p.add != 0) {
                                             131
    push_add(k * 2, p.add);
                                             132
   push_add(k * 2 + 1, p.add);
   p.add = 0;
 if(p.g1 < v[k * 2].g1) push_min(k * 2, p 134
  if(p.11 > v[k * 2].11) push_max(k * 2, p 136)
  if(p.g1 < v[k * 2 + 1].g1) push min(k *
      2 + 1, p.g1);
  if(p.11 > v[k * 2 + 1].11) push max(k *
      2 + 1, p.11);
                                             141
                                             142
                                             143
void subtree_chmin(int k, ll x) {
                                             144
 if(v[k].g1 <= x) return;</pre>
  if(v[k].g2 < x) {
                                             146
   push_min(k, x);
                                             147
   return:
  push(k);
  subtree chmin(k * 2, x), subtree chmin(k 151
       * 2 + 1, x);
  update(k);
                                             153
                                             154
void subtree chmax(int k, ll x) {
 if(x <= v[k].l1) return;</pre>
  if(x < v[k].12) {
                                             157
    push max(k, x);
                                             158
   return;
  subtree_chmax(k * 2, x), subtree_chmax(k 162
       * 2 + 1, x);
  update(k);
template<int cmd>
inline void apply(int k, ll x) {
 if constexpr(cmd == 1) subtree chmin(k,
  if constexpr(cmd == 2) subtree chmax(k,
  if constexpr(cmd == 3) push add(k, x);
  if constexpr(cmd == 4) subtree chmin(k,
      x), subtree chmax(k, x);
template<int cmd>
void inner apply(int 1, int r, 11 x) {
 if(1 == r) return;
```

```
1 += n, r += n;
    for(int i = log; i >= 1; i--) {
     if(((1 >> i) << i) != 1) push(1 >> i):
     if(((r >> i) << i) != r) push((r - 1)
     int 12 = 1, r2 = r;
     while (1 < r) {
        if(1 & 1) _apply<cmd>(1++, x);
        if(r & 1) apply<cmd>(--r, x);
       1 >>= 1, r >>= 1;
     1 = 12, r = r2;
   for(int i = 1; i <= log; i++) {</pre>
     if(((1 >> i) << i) != 1) update(1 >> i
     if(((r >> i) << i) != r) update((r -
          1) >> i);
  template<int cmd>
 inline 11 e() {
   if constexpr(cmd == 1) return INF;
   if constexpr(cmd == 2) return -INF;
   return 0:
 template<int cmd>
 inline void op(11& a, const Node& b) {
   if constexpr(cmd == 1) a = min(a, b.l1);
   if constexpr(cmd == 2) a = max(a, b.g1);
   if constexpr(cmd == 3) a += b.sum;
 template<int cmd>
 11 inner fold(int 1, int r) {
   if(1 == r) return e<cmd>();
   1 += n, r += n;
   for(int i = log; i >= 1; i--) {
     if(((1 >> i) << i) != 1) push(1 >> i);
     if(((r >> i) << i) != r) push((r - 1)
          >> i):
   11 1x = e < cmd > (), rx = e < cmd > ();
   while (1 < r) {
     if(1 \& 1) op < cmd > (1x, v[1++]);
     if(r \& 1) op < cmd > (rx, v[--r]);
     1 >>= 1, r >>= 1:
   if constexpr(cmd == 1) lx = min(lx, rx);
   if constexpr(cmd == 2) lx = max(lx, rx);
   if constexpr(cmd == 3) lx += rx;
   return lx:
2.11 union-of-rectangles
```

12

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```
5 vector<int> vx, vy;
6 struct q { int piv, s, e, x; };
7 struct tree {
    vector<int> seg, tag;
    tree(int _n) : seg(_n * 16), tag(_n * 16)
    void add(int ql, int qr, int x, int v, int
          1, int r) {
      if(qr <= 1 || r <= q1) return;</pre>
      if(ql <= 1 && r <= qr) {
        tag[v] += x;
        if(tag[v] == 0) {
          if(1 != r) seg[v] = seg[2 * v] + seg
               [2 * v + 1];
           else seg[v] = 0;
        } else seg[v] = vx[r] - vx[1];
      } else {
        int mid = (1 + r) / 2;
        add(ql, qr, x, 2 * v, l, mid);
        add(q1, qr, x, 2 * v + 1, mid, r);
        if(tag[v] == 0 && 1 != r) seg[v] = seg
             [2 * v] + seg[2 * v + 1];
    int q() { return seg[1]; }
26 };
27 int main() {
    int n; cin >> n;
    vector\langle int \rangle x1(n), x2(n), y_(n), y2(n);
    for (int i = 0; i < n; i++) {</pre>
      cin \gg x1[i] \gg x2[i] \gg y_[i] \gg y2[i];
            // L R D U
      vx.pb(x1[i]), vx.pb(x2[i]);
      vy.pb(y_[i]), vy.pb(y2[i]);
    vx = sort unique(vx);
    vy = sort_unique(vy);
    vector<q> a(2 * n);
    REP(i, n) {
      x1[i] = lower_bound(ALL(vx), x1[i]) - vx
            .begin();
      x2[i] = lower_bound(ALL(vx), x2[i]) - vx
            .begin();
      y [i] = lower bound(ALL(vy), y [i]) - vy
      y2[i] = lower bound(ALL(vy), y2[i]) - vy
            .begin();
      a[2 * i] = {y_[i], x1[i], x2[i], +1};
      a[2 * i + 1] = \{y2[i], x1[i], x2[i],
    sort(ALL(a), [](q a, q b) { return a.piv <</pre>
          b.piv; });
    tree seg(n);
    11 \text{ ans} = 0;
    REP(i, 2 * n) {
      int j = i;
      while(j < 2 * n && a[i].piv == a[j].piv)</pre>
        seg.add(a[j].s, a[j].e, a[j].x, 1, 0,
             vx.size());
      if(a[i].piv + 1 != SZ(vy)) ans += 1LL *
            seg.q() * (vy[a[i].piv + 1] - vy[a[i
           ].piv]);
```

```
1 // 2
2 // 1 10 1 10
3 // 0 2 0 2
4 // ans = 84
```

i = j - 1;

cout << ans << "\n":

```
2.12 CHT
i struct line t {
   mutable 11 k, m, p;
   bool operator<(const line_t& o) const {</pre>
         return k < o.k; }</pre>
   bool operator<(ll x) const { return p < x;</pre>
 template < bool MAX >
  struct CHT : multiset<line_t, less<>>> {
   const ll INF = 1e18L;
   bool isect(iterator x, iterator y) {
     if(y == end()) return x->p = INF, 0;
     if(x->k == y->k) {
        x \rightarrow p = (x \rightarrow m \rightarrow y \rightarrow m ? INF : -INF);
     } else {
        x \rightarrow p = floor_div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y)
              ->k); // see Math
      return x->p >= y->p;
   void add line(ll k, ll m) {
     if(!MAX) k = -k, m = -m;
      auto z = insert(\{k, m, 0\}), y = z++, x =
      while(isect(y, z)) z = erase(z);
      if(x != begin() && isect(--x, y)) isect(
           x, y = erase(y));
      while((y = x) != begin() && (--x)->p >=
           y->p) isect(x, erase(y));
   11 get(11 x) {
      assert(!empty());
      auto 1 = *lower_bound(x);
     return (1.k * x + 1.m) * (MAX ? +1 : -1)
```

#### **if**(v->1) v->1->rev ^= 1; **if**(v->r) v->r->rev ^= 1; v->rev = false: Node\* merge(Node\* a, Node\* b) { **if**(!a | | !b) **return** (a ? a : b); push(a), push(b); **if**(a->pri > b->pri) { a->r = merge(a->r, b); pull(a); return a; } else { b->1 = merge(a, b->1); pull(b); return b; pair<Node\*, Node\*> split(Node\* v, int k) { if(!v) return {NULL, NULL}; push(v): $if(size(v->1) >= k) {$ auto p = split(v->1, k); if(p.first) p.first->p = NULL; $v \rightarrow 1 = p.second;$ pull(v); return {p.first, v}; } else { auto p = split(v->r, k - size(v->l) - 1)if(p.second) p.second->p = NULL; v->r = p.first; pull(v); return {v, p.second}; int get position(Node\* v) { // 0-indexed int k = (v->1 != NULL ? v->1->sz : 0);while(v->p != NULL) { **if**(v == v->p->r) { $if(v\rightarrow p\rightarrow 1 != NULL) k += v\rightarrow p\rightarrow 1\rightarrow sz;$ $v = v \rightarrow p;$ return k;

#### 2.14 VEB

#### **2.13** treap

```
struct Node {
    bool rev = false;
    int sz = 1, pri = rng();
    Node *1 = NULL, *r = NULL, *p = NULL;
    // TODO
}

void pull(Node*& v) {
    v->sz = 1 + size(v->l) + size(v->r);
    // TODO
}

void push(Node*& v) {
    if(v->rev) {
        swap(v->l, v->r);
}
```

```
1 template < int B, typename ENABLE = void>
    constexpr static int K = B / 2, R = (B +
         1) / 2, M = 1 << B, S = 1 << K, MASK = 65
          (1 << R) - 1;
    array<VEB<R>, S> child;
    VEB\langle K\rangle act = {};
    int mn = M, mx = -1;
    bool empty() { return mx < mn; }</pre>
    bool contains(int i) { return find next(i) 71
          == i; }
    int find_next(int i) { // >=
      if(i <= mn) return mn;</pre>
      if(i > mx) return M;
      int i = i \gg R, x = i \& MASK:
      int res = child[j].find next(x);
      if(res <= MASK) return (j << R) + res;</pre>
```

```
void insert(int i) {
      if(i <= mn) {
        if(i == mn) return;
        swap(mn, i);
        if(i == M) mx = mn;
        if(i >= mx) return;
      } else if(i >= mx) {
        if(i == mx) return;
        swap(mx, i);
        if(i <= mn) return;</pre>
      int j = i \gg R;
      if(child[j].empty()) act.insert(j);
      child[j].insert(i & MASK);
    void erase(int i) {
      if(i <= mn) {
        if(i < mn) return;</pre>
        i = mn = find next(mn + 1);
        if(i >= mx) {
           if(i > mx) mx = -1;
           return;
      } else if(i >= mx) {
        if(i > mx) return;
        i = mx = find prev(mx - 1);
        if(i <= mn) return;</pre>
                                                   23
      int j = i >> R;
      child[j].erase(i & MASK);
                                                   25
      if(child[j].empty()) act.erase(j);
58
    void clear() {
      mn = M, mx = -1, act.clear();
      REP(i, S) child[i].clear();
  template<int B>
  struct VEB<B, enable if t<(B <= 6)>> {
                                                   33
    constexpr static int M = 1 << B;</pre>
    unsigned long long act = 0;
                                                   34
    bool empty() { return !act; }
                                                   35
    void clear() { act = 0; }
    bool contains(int i) { return find_next(i)
          == i: }
    void insert(int i) { act |= 1ULL << i; }</pre>
    void erase(int i) { act &= ~(1ULL << i); }</pre>
                                                  39
    int find next(int i) {
74
      ull tmp = act >> i;
75
      return (tmp ? i + builtin ctzll(tmp) :
```

j = act.find\_next(j + 1);

1.find next(0):

int find prev(int i) { // <=</pre>

j = act.find prev(j - 1);

find prev(MASK);

int  $j = i \gg R$ , x = i & MASK;

int res = child[j].find prev(x);

if(res >= 0) return (j << R) + res;

return j < 0 ? mn : (j << R) + child[j].</pre>

if(i >= mx) return mx;

if(i < mn) return -1;</pre>

return j >= S ? mx : (j << R) + child[j</pre>

#### 2.15 rect-add-rect-sum

```
i template < class Int, class T>
2 struct RectangleAddRectangleSum {
   struct AQ { Int xl, xr, yl, yr; T val; };
   struct SQ { Int xl, xr, yl, yr; };
   vector<AQ> add qry;
   vector<SQ> sum qry;
   // A[x][y] += val for(x, y) in [xl, xr) *
        [yl, yr)
   void add_rectangle(Int xl, Int xr, Int yl,
         Int yr, T val) { add_qry.pb({xl, xr,
        yl, yr, val}); }
   // Get sum of A[x][y] for (x, y) in [xl, xr]
        ) * [yl, yr)
   void add query(Int xl, Int xr, Int yl, Int
         yr) { sum_qry.pb({xl, xr, yl, yr}); }
   vector<T> solve() {
     vector<Int> ys;
     for(auto &a : add gry) ys.pb(a.yl), ys.
          pb(a.yr);
     ys = sort_unique(ys);
     const int Y = SZ(ys);
     vector<tuple<Int, int, int>> ops;
     REP(q, SZ(sum qry)) {
       ops.eb(sum qry[q].xl, 0, q);
       ops.eb(sum_qry[q].xr, 1, q);
     REP(q, SZ(add qry)) {
       ops.eb(add_qry[q].xl, 2, q);
       ops.eb(add_qry[q].xr, 3, q);
     sort(ALL(ops));
     fenwick\langle T \rangle b00(Y), b01(Y), b10(Y), b11(Y
      vector<T> ret(SZ(sum_qry));
     for(auto o : ops) {
       int qtype = get<1>(o), q = get<2>(o);
       if(qtype >= 2) {
          const auto& query = add_qry[q];
         int i = lower_bound(ALL(ys), query.
              yl) - ys.begin();
         int j = lower bound(ALL(ys), query.
              yr) - ys.begin();
         T x = get<0>(o);
         T yi = query.yl, yj = query.yr;
         if(qtype & 1) swap(i, j), swap(yi,
          b00.add(i, x * yi * query.val);
         b01.add(i, -x * query.val);
          b10.add(i, -yi * query.val);
         b11.add(i, query.val);
         b00.add(j, -x * yj * query.val);
         b01.add(j, x * query.val);
         b10.add(j, yj * query.val);
```

```
b11.add(j, -query.val);
 } else {
    const auto& query = sum qry[q];
    int i = lower_bound(ALL(ys), query.
        vl) - vs.begin();
    int j = lower bound(ALL(ys), query.
        yr) - ys.begin();
   T x = get<0>(o);
   T yi = query.yl, yj = query.yr;
   if(qtype & 1) swap(i, j), swap(yi,
    ret[q] += b00.get(i - 1) + b01.get(i
         -1) * yi + b10.get(i - 1) * x
        + b11.get(i - 1) * x * yi;
    ret[q] -= b00.get(j - 1) + b01.get(j
         -1) * yj + b10.get(j - 1) * x
        + b11.get(j - 1) * x * yj;
return ret;
```

#### 2.16 CDO

```
i void CDQ(int 1, int r) {
   if(1 + 1 == r) return;
   int mid = (1 + r) / 2;
   CDQ(1, mid), CDQ(mid, r);
   int i = 1;
   FOR(j, mid, r) {
     const Q& q = qry[j];
     while(i < mid && qry[i].x >= q.x) {
       if(qry[i].id == -1) fenw.add(qry[i].y,
             qry[i].w);
       i++;
     if(q.id >= 0) ans[q.id] += q.w * fenw.
          sum(q.y, sz - 1);
   FOR(p, 1, i) if(qry[p].id == -1) fenw.add(
        qry[p].y, -qry[p].w);
   inplace_merge(qry.begin() + 1, qry.begin()
         + mid, qry.begin() + r, [](const Q& a 51
        , const 0& b) {
     return a.x > b.x;
   });
```

#### 2.17 segtree

```
1 template < class S, S (*e)(), S (*op)(S, S)>
2 struct segtree {
   int n, size, log;
   vector<S> st;
   void update(int v) { st[v] = op(st[v <<</pre>
        1], st[v << 1 | 1]); }
   segtree(int _n) : segtree(vector<S>(_n, e
   segtree(const vector<S>& a): n(sz(a)) {
```

```
REP(i, n) st[size + i] = a[i];
  for(int i = size - 1; i; i--) update(i);
void set(int p, S val) {
  st[p += size] = val;
  for(int i = 1; i <= log; ++i) update(p</pre>
S get(int p) const {
  return st[p + size];
S prod(int 1, int r) const {
  assert(0 <= 1 && 1 <= r && r <= n);
  S sml = e(), smr = e();
  1 += size, r += size:
  while(1 < r)  {
    if(1 \& 1) sml = op(sml, st[1++]);
    if(r \& 1) smr = op(st[--r], smr);
   1 >>= 1;
    r >>= 1;
  return op(sml, smr);
S all_prod() const { return st[1]; }
template < class F> int max right(int 1, F f
    ) const {
  assert(0 <= 1 && 1 <= n && f(e()));
  if(1 == n) return n;
  1 += size;
  S sm = e();
    while(~1 & 1) 1 >>= 1;
    if(!f(op(sm, st[1]))) {
      while(1 < size) {</pre>
        1 <<= 1;
        if(f(op(sm, st[1]))) sm = op(sm,
             st[1++]);
      return 1 - size;
    sm = op(sm, st[1++]);
  } while((1 & -1) != 1);
  return n;
template < class F> int min left(int r, F f)
  assert(0 <= r && r <= n && f(e()));
                                              29
  if(r == 0) return 0;
                                              30
  r += size;
                                              31
  S sm = e();
                                              32
    while(r > 1 \&\& (r \& 1)) r >>= 1;
    if(!f(op(st[r], sm))) {
                                              35
      while(r < size) {</pre>
                                              36
        r = r << 1 | 1;
        if(f(op(st[r], sm))) sm = op(st[r
                                              38
             --], sm);
      return r + 1 - size;
                                              40
    sm = op(st[r], sm);
```

} while((r & -r) != r);

 $log = __lg(2 * n - 1), size = 1 << log;$ 

st.resize(size << 1, e());</pre>

# Flow-Matching

#### 3.1 KM

1 template < class T>

70 };

```
2 struct KM {
   static constexpr T INF = numeric limits<T</pre>
        >::max();
    int n, ql, qr;
    vector<vector<T>> w;
    vector<T> hl, hr, slk;
    vi fl, fr, pre, qu;
    vector<bool> v1, vr;
    KM(int n) : n(n), w(n, vector < T > (n, -INF))
         , hl(n), hr(n), slk(n), fl(n), fr(n),
        pre(n), qu(n), vl(n), vr(n) {}
   void add_edge(int u, int v, int x) { w[u][
         v] = x; } // 最小值要加負號
    bool check(int x) {
      v1[x] = 1;
      if(fl[x] != -1) return vr[qu[qr++] = fl[
      while (x != -1) swap (x, fr[fl[x] = pre[x
      return 0;
    void bfs(int s) {
     fill(ALL(slk), INF);
      fill(ALL(v1), 0), fill(ALL(vr), 0);
      ql = qr = 0, qu[qr++] = s, vr[s] = 1;
      while(true) {
       T d;
        while(al < ar) {</pre>
          for(int x = 0, y = qu[ql++]; x < n;
            if(!vl[x] \&\& slk[x] >= (d = hl[x]
                 + hr[y] - w[x][y])) {
              pre[x] = y;
              if(d) slk[x] = d;
              else if(!check(x)) return;
        REP(x, n) if(!vl[x] \&\& d > slk[x]) d =
              slk[x];
        REP(x, n) {
          if(vl[x]) hl[x] += d;
          else slk[x] -= d;
          if(vr[x]) hr[x] -= d;
        REP(x, n) if(!v1[x] \&\& !s1k[x] \&\& !
             check(x)) return;
   T solve() {
      fill(ALL(fl), -1);
      fill(ALL(fr), -1);
```

```
fill(ALL(hr), 0);
      REP(i, n) hl[i] = *max element(ALL(w[i]))
      REP(i, n) bfs(i);
      T ans = 0;
      REP(i, n) ans += w[i][fl[i]]; // i 跟 fl
           [i] 配對
      return ans;
51
```

#### 3.2 bipartite-matching

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```
struct bipartite_matching {
   int n, m; // 二分圖左右人數 (0 ~ n-1), (0
        \sim m-1)
   vector<vi> g;
   vi lhs, rhs, dist; // i 與 Lhs[i] 配對 (
        Lhs[i] == -1 代表沒有配對)
   bipartite matching(int n, int m) : n(n)
        , m(_m), g(_n), lhs(_n, -1), rhs(_m,
        -1), dist(_n) {}
   void add edge(int u, int v) { g[u].pb(v);
   void bfs() {
     queue<int> q;
     REP(i, n) {
       if(lhs[i] == -1) {
         q.push(i);
         dist[i] = 0;
       } else {
         dist[i] = -1;
     while(!q.empty()) {
       int u = q.front(); q.pop();
       for(auto v : g[u]) {
         if(rhs[v] != -1 && dist[rhs[v]] ==
           dist[rhs[v]] = dist[u] + 1;
           q.push(rhs[v]);
   bool dfs(int u) {
     for(auto v : g[u]) {
       if(rhs[v] == -1) {
         rhs[lhs[u] = v] = u;
         return true;
     for(auto v : g[u]) {
       if(dist[rhs[v]] == dist[u] + 1 && dfs(
            rhs[v])) {
         rhs[lhs[u] = v] = u;
         return true;
     return false;
   int solve() {
```

```
int ans = 0;
      while(true) {
        bfs():
       REP(i, n) if(lhs[i] == -1) aug += dfs( 16
       if(!aug) break;
       ans += aug:
      return ans;
53 };
  3.3 Dinic-LowerBound
```

```
1 template < class T>
 struct DinicLowerBound {
   using Maxflow = Dinic<T>;
   int n;
   Maxflow d:
   vector<T> in;
   DinicLowerBound(int _n) : n(_n), d(_n + 2)
   int add edge(int from, int to, T low, T
        high) {
     assert(0 <= low && low <= high);</pre>
     in[from] -= low, in[to] += low;
     return d.add edge(from, to, high - low);
   T flow(int s, int t) {
     T sum = 0;
     REP(i, n) {
       if(in[i] > 0) {
         d.add edge(n, i, in[i]);
         sum += in[i];
       if(in[i] < 0) d.add_edge(i, n + 1, -in</pre>
     d.add_edge(t, s, numeric_limits<T>::max
     if(d.flow(n, n + 1) < sum) return -1;</pre>
     return d.flow(s, t);
```

#### **3.4** MCMF

```
1 template < class S, class T>
2 class MCMF {
 public:
   struct Edge {
     int from, to;
     S cap;
     Edge(int u, int v, S x, T y) : from(u),
          to(v), cap(x), cost(y) {}
   const ll INF = 1E18L;
   int n;
```

```
vector<Edge> edges;
vector<vi> g;
vector<T> d:
vector<bool> ing;
vi pedge;
MCMF(int _n) : n(_n), g(_n), d(_n), inq(_n
     ), pedge(_n) {}
void add edge(int u, int v, S cap, T cost)
  g[u].pb(SZ(edges));
  edges.eb(u, v, cap, cost);
  g[v].pb(SZ(edges));
  edges.eb(v, u, 0, -cost);
bool spfa(int s, int t) {
  bool found = false;
  fill(ALL(d), INF);
  d[s] = 0;
  inq[s] = true;
  queue<int> q;
  q.push(s);
  while(!q.empty()) {
    int u = q.front(); q.pop();
    if(u == t) found = true;
    inq[u] = false;
    for(auto& id : g[u]) {
      const auto& e = edges[id];
      if(e.cap > 0 && d[u] + e.cost < d[e.</pre>
        d[e.to] = d[u] + e.cost;
        pedge[e.to] = id;
        if(!inq[e.to]) {
          q.push(e.to);
          inq[e.to] = true;
  return found;
pair<S, T> flow(int s, int t, S f = INF) {
 S cap = 0;
 T cost = 0;
  while(f > 0 && spfa(s, t)) {
   S \stackrel{\circ}{send} = f;
    int u = t;
    while(u != s) {
      const Edge& e = edges[pedge[u]];
      send = min(send, e.cap);
      u = e.from;
    u = t;
    while(u != s) {
     Edge& e = edges[pedge[u]];
      e.cap -= send;
      Edge& b = edges[pedge[u] ^ 1];
     b.cap += send;
      u = e.from;
    cap += send:
    f -= send;
    cost += send * d[t];
  return {cap, cost};
```

#### 3.5 minimum-general-weightedperfect-matching

// Minimum General Weiahted Matchina (

1 struct Graph {

33

```
Perfect Match) 0-base
    static const int MXN = 105;
    int n, edge[MXN][MXN];
    int match[MXN], dis[MXN], onstk[MXN];
    vector<int> stk;
    void init(int n) {
      for(int i=0; i<n; i++)</pre>
        for(int j=0; j<n; j++)</pre>
           edge[i][j] = 0;
    void add edge(int u, int v, int w) { edge[
         u][v] = edge[v][u] = w; }
    bool SPFA(int u){
      if(onstk[u]) return true;
      stk.push_back(u);
      onstk[u] = 1;
      for(int v=0; v<n; v++){</pre>
        if(u != v && match[u] != v && !onstk[v
           int m = match[v];
           if(dis[m] > dis[u] - edge[v][m] +
                edge[u][v]){
             dis[m] = dis[u] - edge[v][m] +
                  edge[u][v];
             onstk[v] = 1;
             stk.push back(v);
             if(SPFA(m)) return true;
             stk.pop back();
             onstk[v] = 0;
      onstk[u] = 0;
      stk.pop_back();
      return false;
34
    int solve() {
      for(int i = 0: i < n: i += 2) match[i] =</pre>
            i + 1, match[i+1] = i;
      while(true) {
        int found = 0;
        for(int i=0; i<n; i++) dis[i] = onstk[</pre>
             il = 0:
        for(int i=0; i<n; i++){</pre>
           stk.clear();
           if(!onstk[i] && SPFA(i)){
             found = 1;
             while(stk.size()>=2){
               int u = stk.back(); stk.pop_back
               int v = stk.back(); stk.pop back
               match[u] = v;
               match[v] = u;
        if(!found) break;
```

```
54
       int ans = 0;
       for(int i=0; i<n; i++) ans += edge[i][</pre>
            match[i]];
       return ans / 2;
57
58 } graph;
```

#### 3.6 Flow 建模

- · Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x, y, l, u), connect  $x \to y$  with
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \rightarrow v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t \rightarrow s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from Sto T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \rightarrow s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, f' is the
  - 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$ corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x, y) \in M$ ,  $x \to u$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1) if c > 0, otherwise con- $\operatorname{nect} y \to x \text{ with } (\cos t, cap) = (-c, 1)$
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v) > 0, connect  $S \rightarrow v$  with (cost, cap) = (0, d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C + K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer T
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v, v \in G$  with capacity K

- 4. For each edge (u, v, w) in G, connect  $u \to v$  20 and  $v \to u$  with capacity w
- 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$
- 6. T is a valid answer if the maximum flow f < 22K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u, v).
  - 2. Connect  $v \rightarrow v'$  with weight  $2\mu(v)$ , where 27  $\mu(v)$  is the cost of the cheapest edge incident to 28
  - 3. Find the minimum weight perfect matching on G'
- · Project selection problem
  - 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise, create edge (v, t) with capacity  $-p_v$ . 33
  - 2. Create edge (u, v) with capacity w with w being 34 the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit 36 of a subset of projects.
- 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + \mathcal{A}^{\mathcal{G}}_{41})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity  $c_x$  and create edge (s, y) with capacity  $c_y$ .
- 2. Create edge (x, y) with capacity  $c_{xy}$ .
- 3. Create edge (x, y) and edge (x', y') with capacity  $c_{xyx'y'}$ .

#### general-weighted-max-matching

```
1 // 1-based 00
2 struct WeightGraph {
   static const int inf = INT MAX;
   static const int maxn = 514;
   struct edge {
     int u, v, w;
     edge() {}
     edge(int u, int v, int w): u(u), v(v), w
   int n, n_x;
   edge g[maxn * 2][maxn * 2];
   int lab[maxn * 2];
   int match[maxn * 2], slack[maxn * 2], st[
        maxn * 2], pa[maxn * 2];
   int flo from[maxn * 2][maxn + 1], S[maxn *
         2], vis[maxn * 2];
   vector<int> flo[maxn * 2];
   queue<int> q;
   int e_delta(const edge &e) { return lab[e.
        u] + lab[e.v] - g[e.u][e.v].w * 2; }
   void update_slack(int u, int x) { if(!
        slack[x] || e_delta(g[u][x]) < e_delta</pre>
        (g[slack[x]][x])) slack[x] = u; }
   void set_slack(int x) {
```

```
slack[x] = 0;
  REP(u, n) if(g[u + 1][x].w > 0 \&\& st[u +
        1] != x && S[st[u + 1]] == 0)
       update_slack(u + 1, x);
void q_push(int x) {
  if(x <= n) q.push(x);
  else REP(i, SZ(flo[x])) q_push(flo[x][i
void set st(int x, int b) {
  st[x] = b;
  if(x > n) REP(i, SZ(flo[x])) set_st(flo[
       x][i], b);
int get pr(int b, int xr) {
  int pr = find(ALL(flo[b]), xr) - flo[b].
       begin();
  if(pr % 2 == 1) {
    reverse(1 + ALL(flo[b]));
    return SZ(flo[b]) - pr;
  return pr;
void set match(int u, int v) {
  match[u] = g[u][v].v;
  if(u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u
  for(int i = 0; i < pr; ++i) set_match(</pre>
       flo[u][i], flo[u][i ^ 1]);
  set match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() +
       pr, flo[u].end());
void augment(int u, int v) {
  while(true) {
                                             102
    int xnv = st[match[u]];
                                             103
    set_match(u, v);
    if(!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int get lca(int u, int v) {
  static int t = 0:
  for(++t; u || v; swap(u, v)) {
    if(u == 0) continue;
    if(vis[u] == t) return u;
    vis[u] = t;
    if(u = st[match[u]]) u = st[pa[u]];
                                             115
                                             116
  return 0;
                                             117
                                             118
void add blossom(int u, int lca, int v) {
  int b = n + 1;
  while(b <= n \times \&\& st[b]) ++b;
                                             121
  if(b > n_x) n_x++;
  lab[b] = S[b] = 0:
  match[b] = match[lca];
  flo[b].clear(); flo[b].pb(lca);
  for(int x = u, y; x != lca; x = st[pa[y
                                             124
       ]]) flo[b].pb(x), flo[b].pb(y = st[
                                            125
       match[x]]), q push(y);
                                             126
  reverse(1 + ALL(flo[b]));
```

127

```
for(int x = v, y; x != lca; x = st[pa[y 128]
       ]]) flo[b].pb(x), flo[b].pb(y = st[ 129]
       match[x]]), q_push(y);
  set_st(b, b);
  REP(x, n_x) g[b][x + 1].w = g[x + 1][b]. 131
       w = 0:
  REP(x, n) flo_from[b][x + 1] = 0;
  REP(i, SZ(flo[b])) {
                                             133
    int xs = flo[b][i];
                                             134
    REP(x, n_x) if(g[b][x + 1].w == 0 | |
                                             135
         e_delta(g[xs][x + 1]) < e_delta(g[ 136
         b][x + 1])) g[b][x + 1] = g[xs][x
         + 1], g[x + 1][b] = g[x + 1][xs];
    REP(x, n) if(flo_from[xs][x + 1])
                                             137
         flo_from[b][x + 1] = xs;
                                             138
                                             139
  set slack(b);
                                             140
void expand blossom(int b) {
  REP(i, SZ(flo[b])) set_st(flo[b][i], flo 141
  int xr = flo_from[b][g[b][pa[b]].u], pr
                                             143
       = get_pr(b, xr);
                                             144
  for(int i = 0; i < pr; i += 2) {</pre>
                                             145
    int xs = flo[b][i], xns = flo[b][i +
                                             146
                                             147
         1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
                                             148
    slack[xs] = 0, set_slack(xns);
                                             149
    q push(xns);
                                             150
                                             151
  S[xr] = 1, pa[xr] = pa[b];
                                             152
  for(size t i = pr + 1; i < SZ(flo[b]);</pre>
       ++i) {
    int xs = flo[b][i];
                                             154
    S[xs] = -1, set slack(xs);
                                             155
                                             156
  st[b] = 0;
bool on_found_edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
  if(S[v] == -1) {
                                             158
    pa[v] = e.u, S[v] = 1;
                                             159
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
                                             160
    S[nu] = 0, q push(nu);
                                             161
  } else if(S[v] == 0) {
                                             162
    int lca = get_lca(u, v);
    if(!lca) return augment(u,v), augment(
         v,u), true;
    else add blossom(u, lca, v);
  return false;
bool matching() {
                                             169
  memset(S + 1, -1, sizeof(int) * n_x);
                                             170
  memset(slack + 1, 0, sizeof(int) * n x);
                                             171
  q = queue<int>();
                                             172
  REP(x, n_x) if(st[x + 1] == x + 1 &&!
                                             173
       match[x + 1]) pa[x + 1] = 0, S[x +
                                             174
       1] = 0, q_{push}(x + 1);
                                             175
  if(q.empty()) return false;
  while(true) {
                                             176
    while(!q.empty()) {
                                             177
      int u = q.front(); q.pop();
                                             178
      if(S[st[u]] == 1) continue;
```

```
for(int v = 1; v \le n; ++v)
        if(g[u][v].w > 0 && st[u] != st[v
          if(e_delta(g[u][v]) == 0) {
            if(on_found_edge(g[u][v]))
                 return true;
          } else update_slack(u, st[v]);
    int d = inf;
    for(int b = n + 1; b \le n x; ++b) if(
        st[b] == b \&\& S[b] == 1) d = min(d)
         , lab[b] / 2);
    for(int x = 1; x <= n x; ++x) {
      if(st[x] == x && slack[x]) {
        if(S[x] == -1) d = min(d, e delta(
             g[slack[x]][x]));
        else if(S[x] == 0) d = min(d,
             e_delta(g[slack[x]][x]) / 2);
    REP(u, n) {
     if(S[st[u + 1]] == 0) {
        if(lab[u + 1] <= d) return 0;</pre>
        lab[u + 1] -= d;
      } else if(S[st[u + 1]] == 1) lab[u +
            1] += d;
    for(int b = n + 1; b <= n_x; ++b)
      if(st[b] == b) {
        if(S[st[b]] == 0) lab[b] += d * 2;
        else if(S[st[b]] == 1) lab[b] -= d
    q = queue<int>();
    for(int x = 1; x \leftarrow n_x; ++x)
     if(st[x] == x && slack[x] && st[
           slack[x]] != x && e_delta(g[
           slack[x]][x]) == 0
        if(on_found_edge(g[slack[x]][x]))
             return true;
    for(int b = n + 1; b <= n_x; ++b)
      if(st[b] == b && S[b] == 1 && lab[b]
            == 0) expand blossom(b);
  return false;
pair<ll, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n_x = n;
  int n matches = 0;
  11 \text{ tot weight = 0};
  for(int u = 0; u <= n; ++u) st[u] = u,
      flo[u].clear();
  int w max = 0;
  for(int u = 1; u \le n; ++u)
   for(int v = 1; v <= n; ++v) {
     flo from[u][v] = (u == v ? u : 0);
      w_max = max(w_max, g[u][v].w);
  for(int u = 1; u <= n; ++u) lab[u] =</pre>
  while(matching()) ++n matches;
  for(int u = 1; u \le n; ++u)
    if(match[u] && match[u] < u)</pre>
      tot_weight += g[u][match[u]].w;
```

```
return make pair(tot weight, n matches);
181
    void add edge(int u, int v, int w) { g[u][
         v].w = g[v][u].w = w; }
    void init(int _n) : n(_n) {
       REP(u, n) REP(v, n) g[u + 1][v + 1] =
            edge(u + 1, v + 1, 0);
186 };
```

#### general-matching

```
1 struct GeneralMaxMatch {
   int n:
   vector<pii> es;
   vi g, vis, mate; // i 與 mate[i] 配對 (
        mate[i] == -1 代表沒有匹配)
   GeneralMaxMatch(int n) : n(n), g(n, -1),
        mate(n, -1) {}
   bool dfs(int u) {
     if(vis[u]) return false;
     vis[u] = true;
     for(int ei = g[u]; ei != -1;) {
       auto [x, y] = es[ei]; ei = y;
       if(mate[x] == -1) {
         mate[mate[u] = x] = u;
         return true;
     for(int ei = g[u]; ei != -1;) {
       auto [x, y] = es[ei]; ei = y;
       int nu = mate[x]:
       mate[mate[u] = x] = u;
       mate[nu] = -1;
       if(dfs(nu)) return true;
       mate[mate[nu] = x] = nu;
       mate[u] = -1;
     return false:
   void add_edge(int a, int b) {
     auto f = [&](int a, int b) {
       es.eb(b, g[a]);
       g[a] = SZ(es) - 1;
     f(a, b); f(b, a);
   int solve() {
     vi o(n); iota(ALL(o), 0);
     int ans = 0;
     REP(it, 100) {
       shuffle(ALL(o), rng);
       vis.assign(n, false);
       for(auto i : o) if(mate[i] == -1) ans
            += dfs(i);
     return ans;
```

#### 3.9 Dinic

template < class T>

```
class Dinic {
public:
 struct Edge {
   int from, to;
   T cap;
    Edge(int x, int y, T z) : from(x), to(y)
        , cap(z) {}
  constexpr T INF = 1E9;
  int n;
  vector<Edge> edges;
 vector<vi> g;
  vi cur, h; // h : level graph
 Dinic(int _n) : n(_n), g(_n) {}
  void add edge(int u, int v, T c) {
   g[u].pb(SZ(edges));
   edges.eb(u, v, c);
    g[v].pb(SZ(edges));
    edges.eb(v, u, 0);
  bool bfs(int s, int t) {
   h.assign(n, -1);
   aueue<int> a:
   h[s] = 0;
    q.push(s);
    while(!q.empty()) {
     int u = q.front(); q.pop();
      for(int i : g[u]) {
       const auto& e = edges[i];
        int v = e.to;
       if(e.cap > 0 \&\& h[v] == -1) {
          h[v] = h[u] + 1;
         if(v == t) return true;
          q.push(v);
    return false;
  T dfs(int u, int t, T f) {
   if(u == t) return f;
   Tr = f:
    for(int& i = cur[u]; i < SZ(g[u]); ++i)</pre>
      int j = g[u][i];
      const auto& e = edges[j];
      int v = e.to:
      T c = e.cap;
      if(c > 0 \&\& h[v] == h[u] + 1) {
       T = dfs(v, t, min(r, c));
        edges[j].cap -= a;
        edges[j ^ 1].cap += a;
       if((r -= a) == 0) return f;
   return f - r;
 T flow(int s, int t, T f = INF) {
   T ans = 0;
    while(f > 0 && bfs(s, t)) {
      cur.assign(n, 0);
     T cur = dfs(s, t, f);
```

```
dfs(remark);
         ans += cur;
                                                          return clique;
         f -= cur;
                                                   50
                                                   51 };
       return ans;
67 };
```

#### 3.10 max-clique

63

66

42

```
1 template<int V>
2 struct max clique {
   using B = bitset<V>;
   int n = 0:
   vector<B> g, buf;
   struct P {
     int idx, col, deg;
     P(int a, int b, int b) : idx(a), col(b),
           deg(c) {}
   max_clique(int _n) : n(_n), g(_n), buf(_n)
   void add_edge(int a, int b) {
     assert(a != b);
     g[a][b] = g[b][a] = 1;
   vector<int> now, clique;
   void dfs(vector<P>& rem){
     if(SZ(clique) < SZ(now)) clique = now;</pre>
     sort(ALL(rem), [](P a, P b) { return a.
          deg > b.deg; });
     int max_c = 1;
     for(auto& p : rem){
       p.col = 0;
       while((g[p.idx] & buf[p.col]).any()) p
            .col++;
       max_c = max(max_c, p.idx + 1);
       buf[p.col][p.idx] = 1;
      REP(i, max_c) buf[i].reset();
     sort(ALL(rem), [&](P a, P b) { return a.
          col < b.col; });</pre>
     for(;SZ(rem); rem.pop back()){
       auto& p = rem.back();
       if(SZ(now) + p.col + 1 <= SZ(clique))</pre>
            break:
       vector<P> nrem;
       B bs;
       for(auto& q : rem){
          if(g[p.idx][q.idx]){
            nrem.eb(q.idx, -1, 0);
            bs[q.idx] = 1;
       for(auto\& q : nrem) q.deg = (bs \& g[q.
            idx]).count();
       now.eb(p.idx);
       dfs(nrem);
       now.pop_back();
   vector<int> solve(){
     vector<P> remark:
     REP(i, n) remark.eb(i, -1, SZ(g[i]));
```

# 4 Geometry

#### 4.1 point-in-convex-hull

```
i int point in convex hull(const vector<P>& a,
        P p) {
    // -1 ON, 0 OUT, +1 IN
    // 要先逆時針排序
    int n = SZ(a);
    if(btw(a[0], a[1], p) || btw(a[0], a[n -
         1], p)) return -1;
    int 1 = 0, r = n - 1;
    while(1 <= r) {</pre>
      int m = (1 + r) / 2;
      auto a1 = cross(a[m] - a[0], p - a[0]);
      auto a2 = cross(a[(m + 1) % n] - a[0], p
            - a[0]);
      if(a1 >= 0 && a2 <= 0) {
        auto res = cross(a[(m + 1) % n] - a[m
            1, p - a[m]);
        return res > 0 ? 1 : (res >= 0 ? -1 :
            0);
15
      if(a1 < 0) r = m - 1;
      else 1 = m + 1;
18
    return 0;
19 }
```

#### 4.2 half-plane

```
i typedef pair < double, double > pdd;
pdd operator-(pdd a,pdd b){return {a.F-b.F,a
       .S-b.S}:}
pdd operator+(pdd a,pdd b){return {a.F+b.F,a
       .S+b.S};}
4 pdd operator*(pdd a, double x){return {a.F*x,
       a.S*x};}
  double dot(pdd a,pdd b){return a.F*b.F+a.S*b
  double cross(pdd a,pdd b){return a.F*b.S-a.S
       *b.F:}
  struct bpmi{
    const double eps=1e-8;
    int n,m,id,1,r;
    pdd pt[55],q[1100];
    struct line{
      pdd x,y;
13
      double z;
      line(pdd _x,pdd _y):x(_x),y(_y){z=atan2(
           y.S,y.F);}
      line(){}
```

```
bool operator<(const line &a)const{</pre>
            return z<a.z;}</pre>
    }a[550],da[1005];
    pdd get_(line x,line y){
      pdd v=x.x-y.x;
      double d=cross(y.y,v)/cross(x.y,y.y);
      return x.x+x.y*d;
    void solve(){
      dq[l=r=1]=a[1];
      for(int i=2;i<=id;++i){</pre>
         while(l<r&&cross(a[i].y,q[r-1]-a[i].x)</pre>
              <=eps) --r;
         while(l<r&&cross(a[i].y,q[l]-a[i].x)<=</pre>
              eps) ++1;
        dq[++r]=a[i];
         if(fabs(cross(dq[r].y,dq[r-1].y))<=eps</pre>
             ){
           if(cross(dq[r].y,a[i].x-dq[r].x)>eps
                ) dq[r]=a[i];
        if(l<r) q[r-1]=get_(dq[r-1],dq[r]);</pre>
      while(l<r&&cross(dq[1].y,q[r-1]-dq[1].x)</pre>
            <=eps) --r;
      if(r-1<=1) return;</pre>
      q[r]=get_(dq[1],dq[r]);
    void cal(){
      double ans=0;
      q[r+1]=q[1];
      for(int i=1;i<=r;++i) ans+=cross(q[i],q[</pre>
      cout<<fixed<<setprecision(3)<<ans/2<<"\n</pre>
    void main_(){
      cin>>n;
      for(int x,y,i=0;i<n;++i){</pre>
         for(int i=0;i<m;++i) cin>>pt[i].F>>pt[
              i].S;
         pt[m]=pt[0];
         for(int i=0;i<m;++i) a[++id]=line(pt[i</pre>
              ],pt[i+1]-pt[i]);
      sort(a+1,a+1+id);
      solve();
      cal();
57 \valderyaya;
```

point

1 using P = pair<11, 11>;

a.Y + b.Y}; }

a.Y - b.Y}; }

.Y \* b}; }

```
5| P operator/(P a, 11 b) { return P{a.X / b, a 7|
       .Y / b}; }
  11 dot(P a, P b) { return a, X * b, X + a, Y *
       b.Y; }
  11 cross(P a, P b) { return a.X * b.Y - a.Y
       * b.X: }
  11 abs2(P a) { return dot(a, a); }
  double abs(P a) { return sqrt(abs2(a)); }
int sign(ll x) { return x < 0 ? -1 : (x == 0)
       ? 0 : 1); }
  int ori(P a, P b, P c) { return sign(cross(b 14 pdd min enclosing circle(vector<pdd> dots,
       - a, c - a)); }
12 bool collinear(P a, P b, P c) { return sign( 15
       cross(a - c, b - c)) == 0; }
13 bool btw(Pa, Pb, Pc) {
    if(!collinear(a, b, c)) return 0;
    return sign(dot(a - c, b - c)) <= 0;</pre>
17 bool seg intersect(P a, P b, P c, P d) {
    int a123 = ori(a, b, c);
    int a124 = ori(a, b, d);
    int a341 = ori(c, d, a);
    int a342 = ori(c, d, b);
    if(a123 == 0 && a124 == 0) {
      return btw(a, b, c) || btw(a, b, d) ||
           btw(c, d, a) || btw(c, d, b);
    return a123 * a124 <= 0 && a341 * a342 <=
         0;
  P intersect(P a, P b, P c, P d) {
    int a123 = cross(b - a, c - a);
    int a124 = cross(b - a, d - a);
    return (d * a123 - c * a124) / (a123 -
         a124):
33 struct line { P A, B; };
34 P vec(line L) { return L.B - L.A; }
35 P projection(P p, line L) { return L.A + vec
       (L) / abs(vec(L)) * dot(p - L.A, vec(L))
        / abs(vec(L)); }
        定理
```

- 皮克定理
  - 若一個多邊形的所有頂點都在整數點上,則該 多邊形的面積  $S = a + \frac{b}{3} - 1$  · 其中 a 為內部 格點數目, b 為邊上格點數目。

#### 4.5 min-enclosing-circle

```
i | pdd excenter(pdd x, pdd y, pdd z) {
                                                     #define f(x, y) (x*x+y*y)
2 P operator+(P a, P b) { return P{a.X + b.X,
                                                     auto [x1, y1] = x;
                                                    auto [x2, y2] = y;
3 P operator-(P a, P b) { return P{a.X - b.X,
                                                     auto [x3, y3] = z;
4 P operator*(P a, 11 b) { return P{a.X * b, a
                                                     double d1 = f(x2, y2) - f(x1, y1), d2 = f(
                                                         x3, y3) - f(x2, y2);
```

```
double fm = 2 * ((y3 - y2) * (x2 - x1) - (5) if(sign(cross(a, b)) == 0) return abs2(a)
         v2 - v1) * (x3 - x2));
    double ans_x = ((y3 - y2) * d1 - (y2 - y1)
         * d2) / fm;
    double ans y = ((x2 - x1) * d2 - (x3 - x2)
         * d1) / fm;
    return {ans_x, ans_y};
       double& r) {
    random_shuffle(ALL(dots));
    pdd C = dots[0];
    #define check(i, j) REP(i, j) if(abs(dots[
         i] - C) > r)
    check(i, SZ(dots)) {
      C = dots[i], r = 0;
      check(j, i) {
        C = (dots[i] + dots[j]) / 2.0;
        r = abs(dots[i] - C);
        check(k, j) {
          C = excenter(dots[i], dots[j], dots[ 12
          r = abs(dots[i] - C);
    #undef check
    return C;
32 }
```

#### 4.6 convex-hull

28

```
void convex hull(vector<P>& dots) {
   sort(ALL(dots));
   vector<P> ans(1, dots[0]);
   for(int it = 0; it < 2; it++, reverse(ALL(</pre>
        dots))) {
      for(int i = 1, t = SZ(ans); i < SZ(dots)</pre>
          ; ans.pb(dots[i++])) {
        while(SZ(ans) > t && ori(ans[SZ(ans) -
              2], ans.back(), dots[i]) < 0) {
          ans.ppb();
   ans.ppb();
   swap(ans, dots);
```

#### polar-angle-sort

```
1 bool cmp(P a, P b) {
  #define ng(k) (sign(k.Y) < 0 | (sign(k.Y)
         == 0 \&\& sign(k.X) < 0))
   int A = ng(a), B = ng(b);
  if(A != B) return A < B;</pre>
```

#### < abs2(b);</pre> return sign(cross(a, b)) > 0;

#### 4.8 closest-pair

```
| const | 11 INF = 9e18L + 5:
vector<P> a;
 11 SQ(11 x) { return x * x; }
 ll solve(int l, int r) {
   if(1 + 1 == r) return INF;
    int m = (1 + r) / 2;
    11 \text{ midx} = a[m].x;
    11 d = min(solve(1, m), solve(m, r));
    inplace_merge(a.begin() + 1, a.begin() + m
        , a.begin() + r, [](P a, P b) {
      return a.y < b.y;</pre>
    });
    vector<P> p;
    for(int i = 1; i < r; ++i) if(SQ(a[i].x -</pre>
        midx) < d) p.pb(a[i]);
    REP(i, sz(p)) {
15
      for(int j = i + 1; j < sz(p); ++j) {
          d = min(d, SQ(p[i].x - p[j].x) + SQ(
              p[i].y - p[j].y));
        if(SQ(p[i].y - p[j].y) > d) break;
19
20
   return d; // 距離平方
```

## Graph

#### 5.1 centroid-tree

```
| pair<int, vector<vi>>> centroid_tree(const
       vector<vi>& g) {
    int n = sz(g);
    vi siz(n);
    vector<bool> vis(n);
    auto dfs_sz = [&](auto f, int u, int p) ->
          void {
      siz[u] = 1;
      for(auto v : g[u]) {
        if(v == p || vis[v]) continue;
        f(f, v, u);
        siz[u] += siz[v];
11
12
    auto find_cd = [&](auto f, int u, int p,
         int all) -> int {
      for(auto v : g[u]) {
        if(v == p || vis[v]) continue;
        if(siz[v] * 2 > all) return f(f, v, u,
              all);
```

#### 5.2 chromatic-number

```
1 // vi to(n);
2 // to[u] |= 1 << v;
3 // to[v] |= 1 << u;
4 int chromatic_number(vi g) {
   constexpr int MOD = 998244353;
   int n = SZ(g);
   vector<int> I(1 << n); I[0] = 1;</pre>
   FOR(s, 1, 1 << n) {
     int v = __builtin_ctz(s), t = s ^ (1 <<</pre>
     I[s] = (I[t] + I[t \& \sim g[v]]) \% MOD;
    auto f = I;
   FOR(k, 1, n + 1) {
     int sum = 0;
     REP(s, 1 << n) {
        if(( builtin popcount(s) ^ n) & 1)
             sum -= f[s];
        else sum += f[s];
        sum = ((sum % MOD) + MOD) % MOD;
        f[s] = 1LL * f[s] * I[s] % MOD;
      if(sum != 0) return k;
   return 48763;
```

#### 5.3 count-bridge-online

```
last visit.assign(n, 0);
      iota(ALL(dsu cc), 0);
      dsu 2ecc = dsu cc:
      bridges = 0;
int find 2ecc(int v) {
      if(v == -1) return -1;
      return dsu 2ecc[v] == v ? v : dsu 2ecc[v 78
           ] = find 2ecc(dsu 2ecc[v]);
  int find cc(int v) {
      v = find 2ecc(v);
      return dsu_cc[v] == v ? v : dsu_cc[v] =
           find cc(dsu cc[v]);
  void make root(int v) {
      v = find 2ecc(v):
      int root = v, child = -1;
      while(v != -1) {
          int p = find_2ecc(par[v]);
          par[v] = child;
          dsu_cc[v] = root;
          child = v;
          v = p;
      dsu_cc_size[root] = dsu_cc_size[child];
  void merge_path(int a, int b) {
      ++lca_iteration;
      vector<int> path a, path b;
      int lca = -1;
      while(lca == -1) {
          if(a != -1) {
              a = find_2ecc(a);
              path a.push back(a);
              if(last visit[a] ==
                   lca iteration){
                  lca = a:
                  break;
              last visit[a] = lca iteration;
              a = par[a];
          if(b != -1) {
              b = find_2ecc(b);
              path b.push back(b);
              if(last visit[b] ==
                   lca iteration){
                  lca = b:
                  break;
              last visit[b] = lca iteration;
              b = par[b];
      for(int v : path a) {
          dsu 2ecc[v] = lca;
          if(v == lca) break;
          --bridges:
      for(int v : path b) {
          dsu 2ecc[v] = 1ca:
          if(v == lca) break;
          --bridges:
```

# 5.4 2-SAT

72 void add edge(int a, int b) {

if(a == b) return;

++bridges;

make root(a);

} else merge\_path(a, b);

if(ca != cb)

a = find 2ecc(a), b = find 2ecc(b);

par[a] = dsu cc[a] = b;

int ca = find\_cc(a), cb = find\_cc(b);

if(dsu cc size[ca] > dsu cc size[cb

dsu\_cc\_size[cb] += dsu\_cc\_size[a];

1) swap(a, b), swap(ca, cb);

21

22

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#### 5.5 lowlink

```
1 struct lowlink {
   int n, cnt = 0, tecc cnt = 0, tvcc cnt =
    vector<vector<pii>>> g;
    vector<pii> edges:
    vi roots, id, low, tecc_id, tvcc_id;
    vector<bool> is bridge, is cut,
         is tree edge;
    lowlink(int _n) : n(_n), g(_n), is_cut(_n,
          false), id(_n, -1), low(_n, -1) {}
    void add_edge(int u, int v) {
      g[u].eb(v, SZ(edges));
      g[v].eb(u, SZ(edges));
      edges.eb(u, v);
      is_bridge.pb(false);
13
      is tree edge.pb(false);
      tvcc_id.pb(-1);
14
15
    void dfs(int u, int peid = -1) {
      static vi stk;
```

```
static int rid;
  if(peid < 0) rid = cnt;</pre>
  if(peid == -1) roots.pb(u);
  id[u] = low[u] = cnt++;
  for(auto [v, eid] : g[u]) {
   if(eid == peid) continue;
   if(id[v] < id[u]) stk.pb(eid);</pre>
    if(id[v] >= 0) low[u] = min(low[u], id
      is tree edge[eid] = true;
      dfs(v, eid);
      low[u] = min(low[u], low[v]);
      if((id[u] == rid && id[v] != rid +
           1) || (id[u] != rid && low[v] >=
            id[u])) {
        is cut[u] = true;
      if(low[v] >= id[u]) {
        while(true) {
          int e = stk.back();
          stk.pop_back();
          tvcc_id[e] = tvcc_cnt;
          if(e == eid) break;
        tvcc_cnt++;
void build() {
  REP(i, n) if(id[i] < 0) dfs(i);</pre>
  REP(i, SZ(edges)) {
   auto [u, v] = edges[i];
if(id[u] > id[v]) swap(u, v);
    is bridge[i] = (id[u] < low[v]);</pre>
vector<vi>two_ecc() { // 邊雙
  tecc cnt = 0:
  tecc_id.assign(n, -1);
  vi stk;
  REP(i, n) {
   if(tecc_id[i] != -1) continue;
    tecc_id[i] = tecc_cnt;
    stk.pb(i);
    while(SZ(stk)) {
      int u = stk.back(); stk.pop back();
      for(auto [v, eid] : g[u]) {
        if(tecc_id[v] >= 0 || is_bridge[
             eid]) continue;
        tecc_id[v] = tecc_cnt;
        stk.pb(v);
    tecc cnt++;
  vector<vi> comp(tecc_cnt);
  REP(i, n) comp[tecc_id[i]].pb(i);
  return comp;
vector<vi> bcc_vertices() { // 點雙
  vector<vi> comp(tvcc_cnt);
  REP(i, SZ(edges)) {
    comp[tvcc_id[i]].pb(edges[i].first);
```

```
comp[tvcc_id[i]].pb(edges[i].second);
  for(auto& v : comp) {
    sort(ALL(v));
    v.erase(unique(ALL(v)), v.end());
 REP(i, n) if(!SZ(g[i])) comp.pb({i});
 return comp:
vector<vi> bcc_edges() {
 vector<vi> ret(tvcc cnt);
 REP(i, SZ(edges)) ret[tvcc_id[i]].pb(i);
 return ret;
```

#### manhattan-mst

```
i template < class T> // [w, u, v]
vector<tuple<T, int, int>> manhattan mst(
      vector<T> xs, vector<T> ys) {
   const int n = SZ(xs);
   vi idx(n); iota(ALL(idx), 0);
   vector<tuple<T, int, int>> ret;
   REP(s, 2) {
     REP(t, 2) {
       auto cmp = [&](int i, int j) { return
            xs[i] + ys[i] < xs[j] + ys[j]; };
       sort(ALL(idx), cmp);
       map<T, int> sweep;
       for(int i : idx) {
         for(auto it = sweep.lower_bound(-ys[
              i]); it != sweep.end(); it =
              sweep.erase(it)) {
           int j = it->second;
           if(xs[i] - xs[j] < ys[i] - ys[j])
                break;
           ret.eb(abs(xs[i] - xs[j]) + abs(ys
                [i] - ys[j]), i, j);
         sweep[-ys[i]] = i;
       swap(xs, ys);
     for(auto\& x : xs) x = -x;
   sort(ALL(ret));
   return ret;
```

#### 5.7 SCC

```
1 struct SCC {
   int n;
   vector<vi> g, h;
   SCC(int _n) : n(_n), g(_n), h(_n) {}
   void add_edge(int u, int v) {
     g[u].pb(v);
     h[v].pb(u);
```

```
vi solve() { // 回傳縮點的編號
  vi id(n), top;
  top.reserve(n):
  function<void(int)> dfs1 = [&](int u) {
   for(auto v : g[u]) if(id[v] == 0) dfs1 35
   top.pb(u);
  REP(v, n) if(id[v] == 0) dfs1(v);
  fill(ALL(id), -1);
  function<void(int, int)> dfs2 = [&](int
      u, int x) {
   id[u] = x;
   for(auto v : h[u]) if(id[v] == -1)
        dfs2(v, x);
 for(int i = n - 1, cnt = 0; i >= 0; --i)
   int u = top[i];
   if(id[u] == -1) dfs2(u, cnt++);
                                            48
  return id;
```

, x);

int get dist(int u, int v) {

int kth anc(int u, int k) {

int d = depth[u] - k;

int z = get lca(a, b);

int fi = depth[a] - depth[z];

int se = depth[b] - depth[z];

get lca(u, v))];

if(depth[u] < k) return -1;</pre>

return depth[u] + depth[v] - 2 \* depth[(

while(depth[top[u]] > d) u = par[top[u

int kth node on path(int a, int b, int k)

 $if(k < 0 \mid | k > fi + se)$  return -1;

vector<pii> get\_path(int u, int v, bool

if(u == v && !include lca) return {};

if(depth[top[u]] > depth[top[v]]) swap

if(depth[u] > depth[v]) swap(u, v); // u

if(u != v || include lca) seg.eb(id[u] +

if(par[u] != -1) g[u].erase(find(ALL(g[u

if(siz[v] > siz[g[u][0]]) swap(v, g[u

void dfs\_link(vector<pii>& euler\_tour, int

top[v] = (v == g[u][0] ? top[u] : v);

if(k < fi) return kth\_anc(a, k);</pre>

return kth anc(b, fi + se - k);

include lca = true) {

while(top[u] != top[v]) {

seg.eb(id[top[v]], id[v]);

!include lca, id[v]);

(u, v);

v = par[top[v]];

vector<pii> seg;

return seg;

siz[u] = 1;

void dfs\_sz(int u) {

par[v] = u;

dfs sz(v);

]), par[u]));

depth[v] = depth[u] + 1;

for(auto& v : g[u]) {

siz[u] += siz[v];

1[0]);

fi[u] = SZ(euler tour);

euler tour.eb(depth[u], u);

dfs link(euler tour, v); euler\_tour.eb(depth[u], u);

id[u] = SZ(tour);

for(auto v : g[u]) {

tour.pb(u);

57

71

81

return tour[id[u] + d - depth[u]];

#### 5.8 HLD

```
1 struct HLD {
   int n;
   vector<vi> g;
   vi siz, par, depth, top, tour, fi, id;
   sparse table<pii, min> st;
   HLD(int _n) : n(_n), g(_n), siz(_n), par(
         _n), depth(_n), top(_n), fi(_n), id(_n
                                                61
      tour.reserve(n);
   void add_edge(int u, int v) {
     g[u].push back(v);
     g[v].push_back(u);
   void build(int root = 0) {
     par[root] = -1;
     top[root] = root:
     vector<pii> euler_tour;
      euler tour.reserve(2 * n - 1);
     dfs sz(root):
     dfs link(euler_tour, root);
     st = sparse table<pii, min>(euler tour);
   int get_lca(int u, int v) {
     int L = fi[u], R = fi[v];
     if(L > R) swap(L, R);
     return st.prod(L, R).second;
   bool is_anc(int u, int v) {
     return id[u] <= id[v] && id[v] < id[u] +</pre>
           siz[u];
   bool on path(int a, int b, int x) {
     return (is ancestor(x, a) || is ancestor
          (x, b)) && is_ancestor(get_lca(a, b)
```

#### planar

```
1 struct FringeOpposedSubset {
    deque<int> left, right;
    FringeOpposedSubset() = default;
    FringeOpposedSubset(int h) : left{h},
         right() {}
 6 template<typename T>
  void extend(T& a, T& b, bool rev = false) {
    rev ? a.insert(a.begin(), b.rbegin(), b.
         : a.insert(a.end(), b.begin(), b.end()
10 }
struct Fringe {
    deque<FringeOpposedSubset> FOPs;
    Fringe(int h) : FOPs{{h}} {}
    bool operator<(const Fringe& o) const {</pre>
      return std::tie(FOPs.back().left.back(),
15
            FOPs.front().left.front()) <</pre>
           std::tie(o.FOPs.back().left.back(),
               o.FOPs.front().left.front());
    void merge(Fringe& o) {
18
      o.merge t alike edges();
      merge_t_opposite_edges_into(o);
      if (FOPs.front().right.empty())
        o.align_duplicates(FOPs.back().left.
             front());
      else
        make_onion_structure(o);
       if (o.FOPs.front().left.size()) FOPs.
            push front(o.FOPs.front());
26
    void merge t alike edges() {
      FringeOpposedSubset ans;
      for (auto& FOP: FOPs) {
        if (!FOP.right.empty()) throw
             runtime error("Exception");
         extend(ans.left, FOP.left);
31
32
33
      FOPs = {ans};
34
    void merge t opposite edges into(Fringe& o
36
       while (FOPs.front().right.emptv() &&
              FOPs.front().left.front() > o.
                   FOPs.front().left.back()) {
         extend(o.FOPs.front().right, FOPs.
              front().left);
         FOPs.pop front();
39
40
41
    void align duplicates(int dfs h) {
      if (FOPs.front().left.back() == dfs h) {
         FOPs.front().left.pop_back();
         swap side();
    void swap side() {
      if (FOPs.front().left.empty() ||
           (!FOPs.front().right.empty() &&
            FOPs.front().left.back() > FOPs.
                front().right.back())) {
```

```
swap(FOPs.front().left, FOPs.front().
             right);
53
    void make onion structure(Fringe& o) {
      auto low = &FOPs.front().left, high = &
           FOPs.front().right;
      if (FOPs.front().left.front() >= FOPs.
           front().right.front())
        swap(low, high);
      if (o.FOPs.front().left.back() < low->
           front())
                                                107
        throw runtime_error("Exception");
      if (o.FOPs.front().left.back() < high->
           front()) {
        extend(*low, o.FOPs.front().left, true
        extend(*high, o.FOPs.front().right,
             true);
                                                114
        o.FOPs.front().left.clear();
                                                115
        o.FOPs.front().right.clear();
                                                116
66
    auto lr condition(int deep) const {
      bool L = !FOPs.front().left.empty() &&
           FOPs.front().left.front() >= deep;
      bool R = !FOPs.front().right.empty() &&
           FOPs.front().right.front() >= deep; 121
      return make_pair(L, R);
                                                122
    void prune(int deep) {
      auto [left, right] = lr_condition(deep);
      while (!FOPs.empty() && (left || right)) 124
        if (left) FOPs.front().left.pop front
        if (right) FOPs.front().right.
                                                127
             pop front();
        if (FOPs.front().left.empty() && FOPs. 129
             front().right.empty())
          FOPs.pop front();
        else
          swap_side();
        if (!FOPs.empty()) tie(left, right) =
             lr_condition(deep);
  unique ptr<Fringe> get merged fringe(deque<
       unique_ptr<Fringe>>& upper) {
    if (upper.empty()) return nullptr;
    sort(upper.begin(), upper.end(), [](auto& 141
         a, auto& b) { return *a < *b; });
    for (auto it = next(upper.begin()); it !=
         upper.end(); ++it)
                                                143
      upper.front()->merge(**it);
                                                144
    return move(upper.front());
  void merge_fringes(vector<deque<unique_ptr</pre>
       Fringe>>>& fringes, int deep) {
    auto mf = get_merged_fringe(fringes.back()
         );
    fringes.pop_back();
                                                150
    if (mf) {
                                                151
      mf->prune(deep);
```

```
if (mf->FOPs.size()) fringes.back().
        push back(move(mf));
struct Edge {
 int from, to:
  Edge(int from, int to) : from(from), to(to 160 int main() {
  bool operator==(const Edge& o) const {
    return from == o.from && to == o.to;
};
struct Graph {
 int n = 0:
  vector<vector<int>> neighbor;
  vector<Edge> edges;
  void add edge(int from, int to) {
   if (from == to) return;
    edges.emplace back(from, to);
    edges.emplace_back(to, from);
 void build() {
    sort(edges.begin(), edges.end(), [](
         const auto& a, const auto& b) {
      return a.from < b.from || (a.from == b</pre>
           .from && a.to < b.to);
    edges.erase(unique(edges.begin(), edges.
        end()), edges.end());
    n = 0;
    for (auto& e : edges) n = max(n, max(e.
         from, e.to) + 1);
    neighbor.resize(n);
    for (auto& e : edges) neighbor[e.from].
        push back(e.to);
vector<int> Deeps;
vector<deque<unique ptr<Fringe>>> fringes;
bool dfs(int x, int parent = -1) {
  for (int y : g.neighbor[x]) {
    if (y == parent) continue;
    if (Deeps[y] < 0) { // tree edge</pre>
      fringes.push back({});
      Deeps[v] = Deeps[x] + 1;
      if (!dfs(y, x)) return false;
    } else if (Deeps[x] > Deeps[y]) { //
      fringes.back().push back(make unique<
          Fringe>(Deeps[v]));
 try {
    if (fringes.size() > 1) merge fringes(
         fringes, Deeps[parent]);
  } catch (const exception& e) {
    return false:
 return true;
bool is planar() {
 Deeps.assign(g.n, -1);
 for (int i = 0; i < g.n; ++i) {
```

```
fringes.clear();
155
       Deeps[i] = 0;
156
       if (!dfs(i)) return false;
157
    return true;
158
159 }
     int n, m, u, v;
     cin >> n >> m;
     for (int i = 0; i < m; ++i) {</pre>
164
       cin >> u >> v;
165
       g.add_edge(u, v);
166
     g.build();
     cout << (is_planar() ? "YES" : "NO") <<</pre>
          endl:
     return 0;
170
```

#### 5.10 BCC-tree

1 struct BlockCutTree {

int n;

```
vector<vi> g;
    vi dfn, low, stk;
    int cnt = 0, cur = 0;
    vector<pii> edges;
    BlockCutTree(int _n) : n(_n), g(_n), dfn(
         _n), low(_n) {}
    void ae(int u, int v) {
      g[u].pb(v);
      g[v].pb(u);
    void dfs(int x) {
      stk.pb(x);
      dfn[x] = low[x] = cur++;
      for(auto y : g[x]) {
        if(dfn[y] == -1) {
           dfs(y);
           low[x] = min(low[x], low[y]);
           if(low[y] == dfn[x]) {
             int v;
             do {
              v = stk.back(), stk.pop_back();
22
23
               edges.eb(n + cnt, v);
             } while (v != y);
24
25
             edges.eb(x, n + cnt);
             cnt++:
        } else low[x] = min(low[x], dfn[y]);
28
29
30
    pair<int, vector<pii>> work() {
31
32
      REP(i, n) {
        if(dfn[i] == -1) {
33
34
           stk.clear();
35
           dfs(i);
37
      return {cnt, edges};
38
39
40 };
```

#### 5.11 triangle-sum

```
1|// Three vertices a < b < cconnected by</pre>
        three edges \{a, b\}, \{a, c\}, \{b, c\}. Find
         xa * xb * xc over all trianales.
 1 int triangle_sum(vector<array<int, 2>> edges
        , vi x) {
    int n = SZ(x);
    vi deg(n);
    vector<vi> g(n);
     for(auto& [u, v] : edges) {
      if(u > v) swap(u, v);
       deg[u]++, deg[v]++;
     REP(i, n) g[i].reserve(deg[i]);
     for(auto [u, v] : edges) {
      if(deg[u] > deg[v]) swap(u, v);
      g[u].pb(v);
13
14
15
    vi val(n);
      int128 ans = 0;
     REP(a, n) {
       for(auto b : g[a]) val[b] = x[b];
       for(auto b : g[a]) {
        11 \text{ tmp} = 0;
         for(auto c : g[b]) tmp += val[c];
         ans += _{int128(tmp)} * x[a] * x[b];
      for(auto b : g[a]) val[b] = 0;
24
25
    return ans % mod:
27 }
```

#### Math

#### 6.1 Min-of-Mod-of-Linear

```
1 \mid // \mid min\{Ax + B \pmod{M} \mid 0 \le x \le N\}
2 int min_of_mod_of_linear(int n, int m, int a
       , int b) {
    11 v = floor sum(n, m, a, b);
    int l = -1, r = m - 1;
     while (r - 1 > 1) {
      int k = (1 + r) / 2;
      if(floor_sum(n, m, a, b + (m - 1 - k)) <
            v + n) r = k;
       else 1 = k;
10
    return r;
11 }
```

#### 6.2 Gauss-Jordan

```
i int GaussJordan(vector<vector<ld>>& a) {
   // -1 no sol, 0 inf sol
   int n = SZ(a);
   REP(i, n) assert(SZ(a[i]) == n + 1);
```

```
REP(i, n) {
  int p = i;
  REP(j, n) {
    if(j < i && abs(a[j][j]) > EPS)
         continue:
    if(abs(a[j][i]) > abs(a[p][i])) p = j;
  REP(j, n + 1) swap(a[i][j], a[p][j]);
  if(abs(a[i][i]) <= EPS) continue;</pre>
  REP(j, n) {
    if(i == j) continue;
    ld delta = a[j][i] / a[i][i];
    FOR(k, i, n + 1) a[j][k] -= delta * a[
         i][k];
bool ok = true:
REP(i, n) {
  if(abs(a[i][i]) <= EPS) {</pre>
    if(abs(a[i][n]) > EPS) return -1;
    ok = false;
return ok:
```

#### 6.3 Miller-Rabin

```
| bool is prime(ll n, vector<ll> x) {
   11 d = n - 1;
   d >>= builtin ctzll(d);
   for(auto a : x) {
     if(n <= a) break;</pre>
     11 t = d, y = 1, b = t;
     while(b) {
       if(b \& 1) y = i128(y) * a % n;
       a = i128(a) * a % n;
       b >>= 1;
     while(t != n - 1 && y != 1 && y != n -
       y = i128(y) * y % n;
       t <<= 1;
     if(y != n - 1 && t % 2 == 0) return 0;
   return 1;
 bool is_prime(ll n) {
   if(n <= 1) return 0;
   if(n % 2 == 0) return n == 2;
   if(n < (1LL << 30)) return is_prime(n, {2,</pre>
   return is_prime(n, {2, 325, 9375, 28178,
        450775, 9780504, 1795265022});
```

#### 6.4 Floor-Sum

```
| | // sum_{i} = 0 ^{n - 1} floor((ai + b) / c)
      in O(a + b + c + n)
 11 floor sum(ll n, ll a, ll b, ll c) {
   assert(0 <= n && n < (1LL << 32));
   assert(1 <= c && c < (1LL << 32));
   ull ans = 0:
   if(a < 0) {
     ull a2 = (a \% c + c) \% c;
     ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a)
         ) / c);
   if(b < 0) {
     ull b2 = (b \% c + c) \% c;
     ans -= 1ULL * n * ((b2 - b) / c);
   ull N = n, C = c, A = a, B = b;
   while(true) {
     if(A >= C) {
       ans += N * (N - 1) / 2 * (A / C);
     if(B >= C) {
       ans += N * (B / C);
       B %= C;
     ull y max = A * N + B;
     if(y_max < C) break;</pre>
     N = y \max / C, B = y \max % C;
     swap(C, A);
   return ans;
 6.5 Discrete-Log
```

#### 6.6 Xor-Basis

```
template<int B>
struct xor_basis {
    using T = long long;
    bool zero = false, change = false;
```

```
int cnt = 0;
array < T, B > p = {};
vector<T> d:
void insert(T x) {
  IREP(i, B) {
    if(x >> i \& 1) {
      if(!p[i]) {
        p[i] = x, cnt++;
        change = true;
        return;
      } else x ^= p[i];
  if(!zero) zero = change = true;
T get min() {
  if(zero) return 0;
  REP(i, B) if(p[i]) return p[i];
T get_max() {
  T ans = 0;
  IREP(i, B) ans = max(ans, ans ^ p[i]);
  return ans;
T get kth(long long k) {
  if(k == 1 && zero) return 0;
  k -= zero;
  if(k >= (1LL << cnt)) return -1;
  update();
  T ans = 0;
  REP(i, SZ(d)) if(k \gg i \& 1) ans ^= d[i
  return ans;
bool contains(T x) {
  if(x == 0) return zero;
  IREP(i, B) if(x \gg i \& 1) \times ^= p[i];
  return x == 0;
void merge(const xor basis& other) { REP(i
     , B) if(other.p[i]) insert(other.p[i])
void update() {
  if(!change) return;
  change = false;
  d.clear();
  REP(j, B) IREP(i, j) if(p[j] \gg i \& 1) p
       [j] ^= p[i];
  REP(i, B) if(p[i]) d.pb(p[i]);
```

#### 6.7 數字

· Bernoulli numbers

 $B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$   $\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^{x}-1} = \sum_{j=0}^{\infty} B_n \frac{x^n}{n!}.$ 

$$S_m(n) = \sum_{k=1}^n k^m$$

$$\frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of *n* distinct elements into exactly *k* groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^{n}$$

$$x^{n} = \sum_{i=0}^{n} S(n,i)(x)_{i}$$

· Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

· Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ . E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)

$$E(n,0) = E(n, n-1) = 1$$
  

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

少方和

$$\begin{split} &\sum_{k=1}^n k^3 = (\frac{n(n+1)}{2})^2 \\ &\sum_{k=1}^n k^4 = \frac{1}{30}(6n^5 + 15n^4 + 10n^3 - n) \\ &\sum_{k=1}^n k^5 = \frac{1}{12}(2n^6 + 6n^5 + 5n^4 - n^2) \\ &\sum_{k=1}^n k^6 = \frac{1}{42}(6n^7 + 21n^6 + 21n^5 - 7n^3 + n) \\ &\text{General form:} \end{split}$$

$$\sum_{k=1}^{n} k^{p} = \frac{1}{p+1} (n \sum_{i=1}^{p} (n+1)^{i} - \sum_{i=2}^{p} {i \choose p+1} \sum_{k=1}^{n} k^{p+1-i})$$

#### 6.8 Primes

1 /\* 12721 13331 14341 75577 123457 222557 556679 999983 1097774749 1076767633 100102021 999997771 1001010013 1000512343 987654361 999991231 999888733 98789101 987777733 999991921 1010101333 1010102101 100000000039 100000000000037 2305843009213693951 4611686018427387847 9223372036854775783 18446744073709551557 \*/

#### 6.9 **Determinant**

```
1 T det(vector<vector<T>> a) {
   int n = SZ(a);
   T ret = 1;
   REP(i, n) {
     int idx = -1;
     FOR(j, i, n) if(a[j][i] != 0) {
       idx = j;
       break:
     if(idx == -1) return 0;
     if(i != idx) {
       ret *= T(-1);
       swap(a[i], a[idx]);
     ret *= a[i][i];
     T inv = T(1) / a[i][i];
     REP(j, n) a[i][j] *= inv;
     FOR(j, i + 1, n) {
       T \times = a[j][i];
       if(x == 0) continue;
       FOR(k, i, n) {
         a[j][k] -= a[i][k] * x;
     }
   return ret;
```

#### 6.10 extgcd

```
1 // ax + by = qcd(a, b)
2 11 ext gcd(11 a, 11 b, 11& x, 11& y) {
   if(b == 0) {
     x = 1, y = 0;
     return a;
   ll x1, y1;
   11 g = ext_gcd(b, a % b, x1, y1);
   x = y1, y = x1 - (a / b) * y1;
   return g;
```

#### 6.11 NTT

```
1 \mid const \ 11 \ mod = (119 << 23) + 1, \ root = 62;
      // = 998244353
2 // For p < 2^30 there is also e.g. 5 << 25,
      7 << 26, 479 << 21
3 // and 483 << 21 (same root). The last two
      are > 10^9.
4 typedef vector<11> v1;
 void ntt(vl &a) {
   int n = SZ(a), L = 31 - builtin clz(n);
   static vl rt(2, 1);
   for(static int k = 2, s = 2; k < n; k *=
        2, s++) {
     rt.resize(n);
```

```
ll z[] = \{1, mod_pow(root, mod >> s, mod_23]
        )};
    FOR(i, k, 2 * k) rt[i] = rt[i / 2] * z[i]
         & 1] % mod;
 vi rev(n):
 REP(i, n) rev[i] = (rev[i / 2] | (i \& 1)
       << L) / 2:
 REP(i, n) if (i < rev[i]) swap(a[i], a[rev 29</pre>
       [i]]);
 for(int k = 1; k < n; k *= 2)
    for(int i = 0; i < n; i += 2 * k) REP(j,</pre>
      11 z = rt[j + k] * a[i + j + k] % mod,
            &ai = a[i + j];
      a[i + j + k] = ai - z + (z > ai ? mod
          : 0):
      ai += (ai + z >= mod ? z - mod : z);
vl conv(const vl &a, const vl &b) {
 if(a.empty() || b.empty()) return {};
 int s = SZ(a) + SZ(b) - 1, B = 32 -
       builtin clz(s), n = 1 \ll B;
 11 \text{ inv} = \text{mod pow(n, mod - 2, mod)};
  vl L(a), R(b), out(n);
 L.resize(n), R.resize(n);
  ntt(L), ntt(R);
 REP(i, n) out[-i & (n - 1)] = inv * L[i] %
       mod * R[i] % mod;
  return {out.begin(), out.begin() + s};
```

25

#### **6.12** Poly

21

```
1 template<int mod>
 struct Poly {
                                                55
   vector<ll> a;
   Poly() {}
   Poly(int n) : a(n) {}
   Poly(const vector<ll>& a) : a(a) {}
   Poly(const initializer list<ll>& a) : a(
   int size() const { return SZ(a); }
                                                61
   void resize(int n) { a.resize(n); }
   void shrink() {
     while(size() && a.back() == 0) a.ppb();
   11 at(int idx) const {
     return idx >= 0 && idx < size() ? a[idx]</pre>
           : 0;
   11& operator[](int idx) { return a[idx]; }
   friend Poly operator+(const Poly& a, const
         Poly& b) {
      Poly c(max(SZ(a), SZ(b)));
     REP(i, SZ(c)) c[i] = (a.at(i) + b.at(i))
           % mod:
     return c;
   friend Poly operator-(const Poly& a, const 76
         Poly& b) {
```

```
Poly c(max(SZ(a), SZ(b)));
   REP(i, SZ(c)) c[i] = (a.at(i) - b.at(i)
       + mod) % mod;
                                               78
   return c;
friend Poly operator*(Poly a, Poly b) {
   return Poly(conv(a.a, b.a)); // see NTT.
                                               83
                                               84
friend Poly operator*(ll a, Poly b) {
   REP(i, SZ(b)) (b[i] *= a) %= mod;
                                               85
   return b;
                                               86
                                               87
friend Poly operator*(Poly a, ll b) {
                                               88
   REP(i, SZ(a)) (a[i] *= b) %= mod;
                                               89
   return a:
Poly& operator+=(Poly b) { return (*this)
     = (*this) + b; }
Poly& operator -= (Poly b) { return (*this)
     = (*this) - b; }
Poly& operator*=(Poly b) { return (*this)
     = (*this) * b; }
Poly& operator*=(ll b) { return (*this) =
      (*this) * b; }
#define MSZ if(m == -1) m = size();
Poly mulxk(int k) const {
   auto b = a:
  b.insert(b.begin(), k, 0);
   return Poly(b);
                                              101
Poly modxk(int k) const {
                                              102
  k = min(k, size());
                                              103
   return Poly(vector<ll>(a.begin(), a.
       begin() + k));
                                              105
                                              106
Poly divxk(int k) const {
   if(size() <= k) return Poly();</pre>
   return Poly(vector<ll>(a.begin() + k, a.
                                              108
       end()));
                                              110
Poly deriv() const {
                                              111
  if(!SZ(a)) return Poly();
                                              112
   Poly c(size() - 1);
   REP(i, size() - 1) c[i] = (i + 1LL) * a[114]
       i + 1 % mod;
                                              115
   return c;
                                              116
                                              117
Poly integr() const {
                                              118
   Poly c(size() + 1);
   REP(i, size()) c[i + 1] = a[i] * mod pow 119
        (i+1, mod-2, mod) % mod;
                                              120
   return c;
                                              121
                                              122
Poly inv(int m = -1) const { MSZ;
                                              123
  Poly x{mod_pow(a[0], mod-2, mod)};
                                              124
   int k = 1;
                                              125
   while(k < m) {</pre>
                                              126
    k *= 2;
                                              127
     x = (x * (Polv{2} - modxk(k) * x)).
                                              128
          modxk(k);
                                              129
                                              130
   return x.modxk(m);
                                              131
                                              132
Polv log(int m = -1) const { MSZ;
```

```
return (deriv() * inv(m)).integr().modxk
       (m);
Poly exp(int m = -1) const { MSZ;
  Poly x{1};
  int k = 1;
  while(k < m) {</pre>
   k *= 2;
    x = (x * (Poly{1} - x.log(k) + modxk(k))
        ))).modxk(k);
  return x.modxk(m);
Poly pow(ll k, int m = -1) const { MSZ;
 if(k == 0) {
    Poly b(m); b[0] = 1;
    return b:
 int s = 0, sz = size();
  while(s < sz && a[s] == 0) s++;</pre>
  if(s == sz) return *this;
 if(m > 0 \&\& s >= (sz + k - 1) / k)
       return Poly(m);
  if(s * k >= m) return Poly(m);
  return (((divxk(s) * mod pow(a[s], mod
       -2, mod)).log(m) * (k % mod)).exp(m)
        * mod pow(a[s], k, mod)).mulxk(s *
       k).modxk(m);
bool has_sqrt() const {
 if(size() == 0) return true;
 int x = 0;
  while(x < size() && a[x] == 0) x++;</pre>
 if(x == size()) return true;
 if(x % 2 == 1) return false:
 11 y = a[x];
  return (y == 0 \mid \mid mod_pow(y, (mod-1)/2,
       mod) == 1):
Poly sqrt(int m = -1) const { MSZ;
 if(size() == 0) return Poly();
 int x = 0;
  while (x < size() \&\& a[x] == 0) x++;
  if(x == size()) return Poly(size());
  Poly f = divxk(x);
  Poly g({mod sqrt(f[0], mod)});
  11 \text{ inv2} = \text{mod pow}(2, \text{mod-2}, \text{mod});
  for(int i = 1; i < m; i *= 2) {</pre>
    g = (g + f.modxk(i * 2) * g.inv(i * 2)
         ) * inv2;
  return g.modxk(m).mulxk(x / 2);
Poly shift(ll c) const {
 int n = size();
  Poly b(*this);
 11 f = 1;
  REP(i, n) {
    (b[i] *= f) %= mod;
    (f *= i + 1) \% = mod:
  reverse(ALL(b.a));
  Polv exp cx(vector<ll>(n, 1));
  FOR(i, 1, n) \exp cx[i] = \exp cx[i - 1] *
        c % mod * mod pow(i, mod-2, mod) %
```

```
b = (b * exp cx).modxk(n);
       reverse(ALL(b.a));
       (f *= mod pow(n, mod-2, mod)) %= mod;
       11 z = mod_pow(f, mod-2, mod);
       IREP(i, n) {
         (b[i] *= z) \% = mod;
         (z *= i) \% = mod;
141
       return b;
142
     Poly mulT(Poly b) const {
       int n = SZ(b);
       if(!n) return Poly();
       reverse(ALL(b.a));
       return ((*this) * b).divxk(n - 1);
     vector<ll> eval(vector<ll> x) const {
       if(size() == 0) return vector<ll>(SZ(x),
       const int n = max(SZ(x), size());
       vector<Poly> q(4 * n);
       vector<ll> ans(SZ(x));
       x.resize(n);
       function < void(int, int, int) > build =
            [&](int p, int 1, int r) {
         if(r - 1 == 1) q[p] = Poly{1, mod - x[}
              1]};
         else {
           int m = (1 + r) / 2;
           build(2 * p, 1, m), build(2 * p + 1,
           q[p] = q[2 * p] * q[2 * p + 1];
162
       build(1, 0, n);
       function<void(int, int, int, const Poly</pre>
            &)> work = [&](int p, int l, int r,
            const Poly& num) {
         if(r - l == 1) {
           if(1 < SZ(ans)) ans[1] = num.at(0);
         } else {
           int m = (1 + r) / 2;
           work(2 * p, 1, m, num.mulT(q[2 * p +
                 1]).modxk(m - 1));
           work(2 * p + 1, m, r, num.mulT(q[2 *
                 p]).modxk(r - m));
172
       };
       work(1, 0, n, mulT(q[1].inv(n)));
       return ans;
176 };
```

#### 6.13 Simplex

```
* Description: Solves a general linear
    maximization problem: maximize $c^T x$
    subject to $Ax \le b$, $x \qe 0$.
* Returns -inf if there is no solution, inf
     if there are arbitrarily good
    solutions, or the maximum value of $c^T
     x$ otherwise.
```

```
4 * The input vector is set to an optimal $x$ 52
        (or in the unbounded case, an
       arbitrary solution fulfilling the
       constraints).
  * Numerical stability is not guaranteed.
       For better performance, define
       variables such that x = 0 is viable.
                                                57
  * vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
  * vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
  * T val = LPSolver(A, b, c).solve(x);
  * Time: O(NM * \#pivots), where a pivot may
        be e.g. an edge relaxation. O(2^n) in
       the general case.
  * 將最小化改成最大化 -> 去除等式 -> 去除大
       於等於 -> 去除自由參數·將 x1 用 x1-x3
 typedef double T; // long double, Rational,
      double + mod<P>...
 typedef vector<T> vd:
 typedef vector<vd> vvd;
 struct LP {
   const T eps = 1e-8, inf = 1/.0;
                                                73
   #define MP make pair
   #define ltj(X) if(s == -1 \mid | MP(X[j], N[j])
         < MP(X[s],N[s])) s=j
   int m, n;
   vi N, B;
   vvd D:
   LP(const vvd& A, const vd& b, const vd& c)
         : m(SZ(b)), n(SZ(c)), N(n+1), B(m), D
        (m+2, vd(n+2)) {
     REP(i, m) REP(j, n) D[i][j] = A[i][j];
     REP(i, m) { B[i] = n+i; D[i][n] = -1; D[i][n] = -1
          i|[n+1] = b[i];
     REP(j, n) \{ N[j] = j; D[m][j] = -c[j]; \}
     N[n] = -1; D[m+1][n] = 1;
   void pivot(int r, int s) {
     T *a = D[r].data(), inv = 1 / a[s];
     REP(i, m + 2) if(i != r \&\& abs(D[i][s])
          > eps) {
       T *b = D[i].data(), inv2 = b[s] * inv;
       REP(j, n + 2) b[j] -= a[j] * inv2;
       b[s] = a[s] * inv2;
     REP(j, n + 2) if(j != s) D[r][j] *= inv;
     REP(i, m + 2) if(i != r) D[i][s] *= -inv
     D[r][s] = inv;
     swap(B[r], N[s]);
   bool simplex(int phase) {
     int x = m + phase - 1;
     while(true) {
       int s = -1;
       REP(j, n + 1) if(N[j] != -phase) ltj(D
            [x]);
       if(D[x][s] >= -eps) return true;
       int r = -1;
       REP(i, m) {
```

if(D[i][s] <= eps) continue;</pre>

```
], B[r])) r = i;
    if(r == -1) return false;
    pivot(r, s);
T solve(vd &x) {
  int r = 0;
  FOR(i, 1, m) if(D[i][n+1] < D[r][n+1]) r
  if(D[r][n+1] < -eps) {
    pivot(r, n);
    if(!simplex(2) || D[m+1][n+1] < -eps)</pre>
         return -inf:
    REP(i, m) if(B[i] == -1) {
      int s = 0;
      FOR(j, 1, n + 1) ltj(D[i]);
      pivot(i, s);
  bool ok = simplex(1); x = vd(n);
  REP(i, m) if(B[i] < n) \times [B[i]] = D[i][n]
  return ok ? D[m][n+1] : inf;
```

#### 6.14 Triangle

```
1 // Counts x. v >= 0 such that Ax + Bv <= C.
       Requires A, B > 0. Runs in Log time.
2 // Also representable as sum {0 <= x <= C /
       A) floor((C - Ax) / B + 1).
3 ll count triangle(ll A, ll B, ll C) {
      if(C < 0) return 0:
      if(A < B) swap(A, B);
      11 m = C / A, k = A / B;
      11 h = (C - m * A) / B + 1;
      return m * (m + 1) / 2 * k + (m + 1) * h
            + count triangle(B, A - k * B, C -
           B * (k * m + h));
11 // Counts 0 \le x \le RA, 0 \le y \le RB such that
         Ax + Bv \leftarrow C. Requires A. B > 0.
12 ll count triangle rectangle intersection(ll
       A, 11 B, 11 C, 11 RA, 11 RB) {
      if(C < 0 || RA <= 0 || RB <= 0) return
      if(C >= A * (RA - 1) + B * (RB - 1))
           return RA * RB:
      return count_triangle(A, B, C) -
           count_triangle(A, B, C - A * RA) -
           count triangle(A, B, C - B * RB);
```

#### 6.15 Chinese-Remainder

```
if(r == -1 \mid \mid MP(D[i][n+1] \mid D[i][s \mid \mid \mid \mid /\mid (rem, mod) \{0, 0\} for no solution
     ], B[i]) < MP(D[r][n+1] / D[r][s 2|pair<ll, ll> crt(ll r0, ll m0, ll r1, ll m1)
                                              r0 = (r0 \% m0 + m0) \% m0;
                                              r1 = (r1 \% m1 + m1) \% m1;
                                              if(m0 < m1) swap(r0, r1), swap(m0, m1);</pre>
                                              if(m0 \% m1 == 0) {
                                                if(r0 % m1 != r1) return {0, 0};
                                              11 g, im, qq;
                                              g = ext gcd(m0, m1, im, qq);
                                              ll u1 = (m1 / g);
                                              if((r1 - r0) % g) return {0, 0};
                                             11 x = (r1 - r0) / g % u1 * im % u1;
                                              r0 += x * m0:
                                              m0 *= u1;
                                             if(r0 < 0) r0 += m0;
                                         16
                                              return {r0, m0};
```

#### 6.16 Pollard-Rho

```
1 | void PollardRho(map<11, int>& mp, 11 n) {
    if(n == 1) return;
    if(is_prime(n)) return mp[n]++, void();
    if(n % 2 == 0) {
       mp[2] += 1;
      PollardRho(mp, n / 2);
       return:
    11 \times 2, y = 2, d = 1, p = 1;
    #define f(x, n, p) ((i128(x) * x % n + p)
         % n)
     while(1) {
      if(d != 1 && d != n) {
12
        PollardRho(mp, d);
13
        PollardRho(mp, n / d);
15
         return;
      p += (d == n);
      x = f(x, n, p), y = f(f(y, n, p), n, p);
18
19
      d = gcd(abs(x - y), n);
20
    #undef f
21
22 }
24 vector<ll> get divisors(ll n) {
    if(n == 0) return {};
    map<ll, int> mp;
    PollardRho(mp, n);
27
    vector<pair<ll, int>> v(ALL(mp));
    vector<ll> res;
29
     auto f = [&](auto f, int i, ll x) -> void
30
      if(i == SZ(v)) {
31
32
         res.pb(x);
33
         return;
34
       for(int j = v[i].second; ; j--) {
        f(f, i + 1, x);
        if(i == 0) break:
        x *= v[i].first;
```

```
0 };
1 f(f, 0, 1);
2 sort(ALL(res));
3 return res;
4 }
```

#### 6.17 Mod-Sqrt

```
1 // return -1 if sqrt DNE
2 | 11 mod_sqrt(11 a, 11 mod) {
   a %= mod;
   if(mod == 2 || a < 2) return a;
   if(mod_pow(a, (mod-1)/2, mod) != 1) return
   11 b = 1;
   while (mod_pow(b, (mod-1)/2, mod) == 1) b
   int m = mod-1, e = __builtin_ctz(m);
   m >>= e;
   11 x = mod_pow(a, (m-1)/2, mod);
   11 y = a * x % mod * x % mod;
   x = x * a % mod;
   11 z = mod_pow(b, m, mod);
   while(y != 1) {
     int j = 0;
     11 t = y;
     while(t != 1) t = t * t % mod, j++;
     z = mod pow(z, 1LL << (e - j - 1), mod);
     x = x*z\%mod, z = z*z\%mod, y = y*z\%mod;
   return min(x, mod-x); // neg is $mod-x$
```

#### 6.18 Combination

#### 6.19 Mod-Inv

```
int inv(int a) {
   if(a < N) return inv[a];
   if(a == 1) 1;
   return (MOD - 1LL * (MOD / a) * inv(MOD %
        a) % MOD) % MOD;</pre>
```

#### **6.20 FWHT**

```
| #define ppc builtin popcount
 template < class T, class F>
  void fwht(vector<T>& a, F f) {
   int n = SZ(a);
   assert(ppc(n) == 1);
   for(int i = 1; i < n; i <<= 1) {</pre>
     for(int j = 0; j < n; j += i << 1) {
       REP(k, i) f(a[j + k], a[i + j + k]);
 template < class T>
 void or_transform(vector<T>& a, bool inv) {
      fwht(a, [\&](T\& x, T\& y) { y += x * (inv)}
      ? -1 : +1); }) }
 template < class T>
 void and transform(vector<T>& a, bool inv) {
       fwht(a, [\&](T\& x, T\& y) { x += y * (inv)}
       ? -1 : +1); }); }
 template < class T>
  void xor_transform(vector<T>& a, bool inv) {
   fwht(a, [](T& x, T& y) {
```

```
Tz = x + y;
      y = x - y;
      x = z;
    });
    if(inv) {
      Tz = T(1) / T(SZ(a));
      for(auto& x : a) x *= z;
27 }
  template < class T>
  vector<T> convolution(vector<T> a, vector<T>
    assert(SZ(a) == SZ(b));
    transform(a, false), transform(b, false);
    REP(i, SZ(a)) a[i] *= b[i];
    transform(a, true);
    return a;
  template < class T>
  vector<T> subset convolution(const vector<T</pre>
       >& f, const vector<T>& g) {
    assert(SZ(f) == SZ(g));
    int n = SZ(f);
    assert(ppc(n) == 1);
    const int lg = __lg(n);
    vector<vector<T>> fhat(lg + 1, vector<T>(n 24
```

)), ghat(fhat);

#### 6.21 Aliens

#### 6.22 Berlekamp-Massey

```
1 // - [1, 2, 4, 8, 16] -> (1, [1, -2])
2 // - [1, 1, 2, 3, 5, 8] -> (2, [1, -1, -1])
3 // - [0, 0, 0, 0, 1] -> (5, [1, 0, 0, 0, 0,
       998244352]) (mod 998244353)
  // - [] -> (0, [1])
5 // - [0, 0, 0] -> (0, [1])
6 // - [-2] -> (1, [1, 2])
7 template < class T>
  pair<int, vector<T>> BM(const vector<T>& S)
    using poly = vector<T>;
    int N = SZ(S);
    poly C_rev{1}, B{1};
    int L = 0, m = 1;
    T b = 1;
    auto adjust = [](poly C, const poly &B, T
         d, T b, int m) -> poly {
      C.resize(max(SZ(C), SZ(B) + m));
      Ta = d / b:
17
      REP(i, SZ(B)) C[i + m] -= a * B[i];
      return C;
19
    REP(n, N) {
20
      T d = S[n];
      REP(i, L) d += C_{rev}[i + 1] * S[n - 1 -
           i];
      if(d == 0) m++;
      else if (2 * L <= n) {</pre>
        poly Q = C_rev;
```

```
++);
    return {L, C rev};
31 }
33 // Calculate x^N \b f(x)
34 // Complexity: \$0(K^2 \setminus \log N)\$ (\$K\$: deg. of
35 // (4, [1, -1, -1]) -> [2, 3]
2)
37 template < class T>
38 vector<T> monomial mod polynomial(long long
       N, const vector<T> &f rev) {
     assert(!f_rev.empty() && f_rev[0] == 1);
    int K = SZ(f rev) - 1;
    if(!K) return {};
    int D = 64 - __builtin_clzll(N);
42
    vector<T> ret(K, 0);
43
     ret[0] = 1;
     auto self conv = [](vector<T> x) -> vector
          <T> {
       int d = SZ(x);
       vector<T> ret(d * 2 - 1);
       REP(i, d) {
        ret[i * 2] += x[i] * x[i];
         REP(j, i) ret[i + j] += x[i] * x[j] *
51
52
       return ret;
53
     for(int d = D; d--;) {
       ret = self conv(ret);
       for(int i = 2 * K - 2; i >= K; i--) {
        REP(j, k) ret[i - j - 1] -= ret[i] *
             f_{rev}[j + 1];
58
59
       ret.resize(K);
       if (N >> d & 1) {
61
        vector<T> c(K);
        c[0] = -ret[K - 1] * f rev[K];
         for(int i = 1; i < K; i++) c[i] = ret[</pre>
             i - 1] - ret[K - 1] * f rev[K - i
64
         ret = c;
65
66
    return ret;
70 // Guess k-th element of the sequence,
       assuming linear recurrence
71 template < class T>
72 T guess kth term(const vector<T>& a, long
       long k) {
     assert(k >= 0);
    if(k < 1LL * SZ(a)) return a[k];</pre>
     auto f = BM<T>(a).second;
     auto g = monomial mod polynomial<T>(k, f);
    T ret = 0:
    REP(i, SZ(g)) ret += g[i] * a[i];
79
    return ret:
```

C rev = adjust(C\_rev, B, d, b, m);

L = n + 1 - L, B = 0, b = d, m = 1;

} else C rev = adjust(C rev, B, d, b, m

#### 定理 6.23

· Cramer's rule

$$ax + by = e$$

$$cx + dy = f$$

$$x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Burnside's Lemma

Let us calculate the number of necklaces of n pearls. where each pearl has m possible colors. Two necklaces are symmetric if they are similar after rotating them. There are n ways to change the position of a necklace, because we can rotate it  $0, 1, \ldots, n_1$  steps clockwise. If the number of steps is 0, all  $m^n$  necklaces remain the same, and if the number of steps is 1, only the m necklaces where each pearl has the same color remain the same. More generally, when the number of steps is k, a total of  $m^{\gcd(k,n)}$  necklaces remain the same. The reason for this is that blocks of pearls of size gcd(k, n)will replace each other. Thus, according to Burnside's lemma, the number of necklaces is  $\sum_{i=0}^{n-1} \frac{m^{\gcd(i,n)}}{n}$ . For example, the number of necklaces of length 4 with 3 colors is  $\frac{3^4+3+3^2+3}{4}=24$ 

Lindstr□m-Gessel-Viennot Lemma

定義

 $\omega(P)$  表示 P 這條路徑上所有邊的邊權之積。( 路徑 計數時,可以將邊權都設為1)(事實上,邊權可以為 生成函數) e(u,v) 表示 u 到 v 的 \*\* 每一條 \*\* 路徑 P 的  $\omega(P)$  之和 · 即  $e(u,v) = \sum \omega(P)$  。 起點

集合 A,是有向無環圖點集的一個子集,大小為 n。 終點集合 B, 也是有向無環圖點集的一個子集, 大小 也為  $n \cdot -$ 組  $A \rightarrow B$  的不相交路徑  $S : S_i$  是一條從  $A_i$  到  $B_{\sigma(S)_i}$  的路徑 ( $\sigma(S)$  是一個排列)·對於任 何  $i \neq j \cdot S_i$  和  $S_j$  沒有公共頂點  $\circ t(\sigma)$  表示排列  $\sigma$ 

$$M = \begin{bmatrix} e(A_1, B_1) & e(A_1, B_2) & \cdots & e(A_1, B_n) \\ e(A_2, B_1) & e(A_2, B_2) & \cdots & e(A_2, B_n) \\ \vdots & \vdots & \ddots & \vdots \\ e(A_n, B_1) & e(A_n, B_2) & \cdots & e(A_n, B_n) \end{bmatrix}$$

$$\det(M) = \sum_{S:A \to B} (-1)^{t(\sigma(S))} \prod_{i=1}^n \omega(S_i)$$

其中  $\sum_{S:A\to B}$  表示滿足上文要求的  $A\to B$  的每一組 不相交路徑 S。

· Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is  $|\det(L_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- · Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i, j) \in E$ , otherwise  $d_{ij}=-d_{ji}.$   $\frac{rank(D)}{2}$  is the maximum matching on G.

- · Cayley's Formula
  - Given a degree sequence  $d_1, d_2, \ldots, d_n$ for each labeled vertices, there are
  - $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees. Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1, 2, \ldots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$
- Erd□s-Gallai theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + \cdots + d_n$  is even

and 
$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$
 holds for every  $1 < k < n$ .

Gale–Ryser theorem

A pair of sequences of nonnegative integers  $a_1 >$  $\cdots \geq a_n$  and  $b_1, \ldots, b_n$  is bigraphic if and only if

Fulkerson–Chen–Anstee theorem

A sequence  $(a_1, b_1), \ldots, (a_n, b_n)$  of nonnegative integer pairs with  $a_1 \ge \cdots \ge a_n$  is digraphic if and only

$$\begin{split} & \text{if } \sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) - \\ & \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n. \end{split}$$

M□bius inversion formula

· Spherical cap

```
h^2)/6 = \pi r^3 (2 + \cos \theta) (1 - \cos \theta)^2/3.
- Area = 2\pi rh = \pi (a^2 + h^2) = 2\pi r^2 (1
```

- A portion of a sphere cut off by a plane.

h: height of the cap,  $\theta$ :  $\arcsin(a/r)$ . - Volume =  $\pi h^2(3r - h)/3 = \pi h(3a^2 + h)$ 

#### **6.24** Int-Div

```
1 | 11 floor_div(11 a, 11 b) {
   return a/b - ((a^b) < 0 && a%b != 0);
4 ll ceil_div(ll a, ll b) {
   return a/b + ((a^b) > 0 && a%b != 0);
```

#### 6.25 生成函數

• Ordinary Generating Function  $A(x) = \sum_{i \geq 0} a_i x^i$ 

• Exponential Generating Function A(x) $\sum_{i>0} \frac{a_i}{i!} x_i$ 

```
 -A(x) + B(x) \Rightarrow a_n + b_n 
-A^{(k)}(x) \Rightarrow a_{n+k_n} 
-A(x)B(x) \Rightarrow \sum_{i=0}^{n} \binom{n}{i} a_i b_{n-i} 
-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n}^{n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1}^{47} a_{i_2}^{48} 
-xA(x) \Rightarrow na_n
```

· Special Generating Function

```
 - (1+x)^n = \sum_{i\geq 0} {n \choose i} x^i 
 - \frac{1}{(1-x)^n} = \sum_{i\geq 0} {n \choose i-1} x^i
```

#### 6.26 GCD-Convolution

```
1 // 2, 3, 5, 7, ...
vector<int> prime enumerate(int N) {
   vector<bool> sieve(N / 3 + 1, 1);
   for(int p = 5, d = 4, i = 1, sqn = sqrt(N)
                                               61
       ; p \leftarrow sqn; p += d = 6 - d, i++) {
      if(!sieve[i]) continue;
      for(int q = p * p / 3, r = d * p / 3 + (
          d * p % 3 == 2), s = 2 * p; q < SZ(
           sieve); q += r = s - r) sieve[q] =
    for(int p = 5, d = 4, i = 1; p <= N; p +=
         d = 6 - d, i++) {
      if(sieve[i]) {
```

```
ret.pb(p);
- r: sphere radius, a: radius of the base of the cap, |a|
                                     13
                                          }
                                           while(SZ(ret) && ret.back() > N) ret.
                                                pop_back();
                                          return ret;
                                      16 }
                                      17 struct divisor transform {
                                      18
                                          template < class T>
                                          static void zeta_transform(vector<T>& a) {
                                            int n = a.size() - 1;
                                             for(auto p : prime_enumerate(n)) {
                                              for(int i = 1; i * p <= n; i++) {
                                                a[i * p] += a[i];
                                      25
                                           template < class T>
                                           static void mobius transform(vector<T>& a)
                                             int n = a.size() - 1;
                                             for(auto p : prime_enumerate(n)) {
                                              for(int i = n / p; i > 0; i--) {
                                                a[i * p] -= a[i];
                                           static void zeta_transform(vector<T>& a) {
                                            int n = a.size() - 1;
                                             for(auto p : prime_enumerate(n)) {
                                              for(int i = n / p; i > 0; i--) {
                                                 a[i] += a[i * p];
                                           template < class T>
                                           static void mobius_transform(vector<T>& a)
                                             int n = a.size() - 1;
                                             for(auto p : prime_enumerate(n)) {
                                             for(int i = 1; i * p <= n; i++) {</pre>
                                                a[i] -= a[i * p];
                                      57 // lcm: multiple -> divisor
                                      58 template < class T>
                                        vector<T> gcd convolution(const vector<T>& a
                                             , const vector<T>& b) {
                                           assert(a.size() == b.size());
                                           auto f = a, g = b;
                                           multiple_transform::zeta_transform(f);
                                      62
                                          multiple transform::zeta transform(g);
                                     64
                                          REP(i, SZ(f)) f[i] *= g[i];
                                          multiple_transform::mobius_transform(f);
                                     65
                                          return f:
```

#### 歐幾里得類算法

•  $m = \lfloor \frac{an+b}{c} \rfloor$ • Time complexity:  $O(\log n)$ 

phi[i \* j] = phi[i] \* phi[j]; 24 25 }

# $f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), \\ 0, \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ - h(c, c-b-1, a, m-1)), \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), \\ 0, & \\ nm(m+1) - 2g(c, c-b-1, a, m-1) - \frac{3}{2} \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), \end{cases} \end{split}$$

#### 6.28 Linear-Sieve

```
vi primes, least = {0, 1}, phi, mobius;
 void LinearSieve(int n) {
   least = phi = mobius = vi(n + 1);
   mobius[1] = 1;
   for(int i = 2; i <= n; i++) {
     if(!least[i]) {
       least[i] = i;
       primes.pb(i);
       phi[i] = i - 1;
       mobius[i] = -1;
     for(auto j : primes) {
       if(i * j > n) break;
       least[i * j] = j;
       if(i % j == 0) {
         mobius[i * j] = 0;
         phi[i * j] = phi[i] * j;
         break:
         mobius[i * j] = -mobius[i];
```

- Estimation
  - The number of divisors of n is at most around 100 for n < 5e4,500 for n < 1e7,2000 for | |inline ull rng() { n < 1e10, 200000 for n < 1e19.
- The number of ways of writing n as a sum of  $\frac{1}{3}$ positive integers, disregarding the order of the  $\frac{1}{2}$   $\tilde{Q} \sim \tilde{Q} >> 9$ ; summands. [1,1,2,3,5,7,11,15,22,30 for 5] return Q & 0xFFFFFFFULL;  $n = 0 \sim 9,627 \text{ for } n = 20, \sim 2e5 \text{ for } _{6}$  $a \geq c \lor b \geq c \begin{cases} n = 50, \sim 2e8 \text{ for } n = 100. \\ \text{Total number of partitions of } n < 0 \lor a = 0 \end{cases}$  distinct elements: B(n) = 1001, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, **97.85**70, **421**3597,  $27644437, 190899322, \ldots$

## Misc

#### **7.1 PBDS**

```
a > c \lor b > c
| #include <ext/pb ds/assoc container.hpp>
2 using namespace __gnu_pbds;
 tree<ll, null_type, less<ll>, rb_tree_tag,
      tree order statistics node update> st;
 // find_by_order order_of_key
5 // float128 t
 for(int i = bs._Find_first(); i < bs.size();</pre>
       i = bs. Find next(i));
```

#### 7.2 python

```
I | from decimal import Decimal, getcontext
getcontext().prec = 1000000000
a = pow(Decimal(2), 82589933) - 1
```

#### timer

```
i clock_t T1 = clock();
2 double getCurrentTime() { return (double) (
      clock() - T1) / CLOCKS_PER_SEC; }
```

#### 7.4 next-combination

```
1 // Example: 1 -> 2, 4 -> 8, 12(1100) ->
      17(10001)
2 11 next_combination(11 comb) {
\frac{1}{3} 11 x = comb & -comb, y = comb + x;
   return ((comb & ~y) / x >> 1) | y;
```

#### 7.5 rng

```
static ull 0 = 48763;
Q ^= Q << 7;
```

```
i inline char gc() {
  static const size_t sz = 65536;
   static char buf[sz];
   static char *p = buf, *end = buf;
   if(p == end) end = buf + fread(buf, 1, sz,
         stdin), p = buf;
   return *p++;
```

#### **7.7** rotate90

```
| vector<vector<T>> rotate90(const vector
      vector<T>>& a) {
     int n = sz(a), m = sz(a[0]);
     vector<vector<T>> b(m, vector<T>(n));
     REP(i, n) REP(j, m) b[j][i] = a[i][m - 1]
          - j];
     return b;
```

# String

#### smallest-rotation

```
string small_rot(string s) {
   int n = SZ(s), i = 0, j = 1;
   s += s;
   while(i < n && j < n) {</pre>
     int k = 0:
      while (k < n \&\& s[i + k] == s[j + k]) k
```

#### $if(s[i + k] \le s[j + k]) j += k + 1;$ **else** i += k + 1; j += (i == j); int ans = i < n ? i : j;</pre> return s.substr(ans, n);

#### 8.2 AC

```
i template<int ALPHABET = 26, char MIN_CHAR =</pre>
 2 struct ac_automaton {
    struct Node {
       int fail = 0, cnt = 0;
       array<int, ALPHABET> go{};
    vector<Node> node;
    vi que;
    int new_node() { return node.eb(), SZ(node
         ) - 1; }
    Node& operator[](int i) { return node[i];
     ac_automaton() { new_node(); } // reserve
     int insert(const string& s) {
       int p = 0;
       for(char c : s) {
        int v = c - MIN_CHAR;
        if(node[p].go[v] == 0) node[p].go[v] =
               new_node();
        p = node[p].go[v];
       node[p].cnt++;
      return p;
21
     void build() {
       que.reserve(SZ(node)); que.pb(0);
24
       REP(i, SZ(que)) {
        int u = que[i];
26
         REP(j, ALPHABET) {
           if(node[u].go[j] == 0) node[u].go[j]
                 = node[node[u].fail].go[j];
           else {
            int v = node[u].go[j];
             node[v].fail = (u == 0 ? u : node[
                  node[u].fail].go[j]);
             que.pb(v);
35
36 };
```

#### 8.3 Z

```
1 // abacbaba -> [0, 0, 1, 0, 0, 3, 0, 1]
vi z_algorithm(const vi& a) {
  int n = SZ(a);
   vi z(n); int j = 0;
   FOR(i, 1, n) {
```

#### 8.4 rolling-hash

```
const ll M = 911382323, mod = 972663749;
ll Get(vector<11>& h, int l, int r) {
    if(!!) return h[r]; // p[i] = M^i % mod
    ll ans = (h[r] - h[l - 1] * p[r - 1 + 1])
        % mod;
    return (ans + mod) % mod;
}

vector<ll> Hash(string s) {
    vector<ll> ans(SZ(s));
    ans[0] = s[0];
    for(int i = 1; i < SZ(s); i++) ans[i] = (
        ans[i - 1] * M + s[i]) % mod;
    return ans;
}</pre>
```

#### 8.5 hash61

```
| const | 11 M30 = (1LL << 30) - 1;
  const 11 M31 = (1LL << 31) - 1;</pre>
  const ll M61 = (1LL << 61) - 1;</pre>
  ull modulo(ull x){
   ull xu = x \gg 61;
    ull xd = x \& M61;
    ull res = xu + xd;
    if(res >= M61) res -= M61;
    return res;
ull mul(ull a, ull b){
   ull au = a >> 31, ad = a & M31;
    ull bu = b \gg 31, bd = b \& M31;
    ull mid = au * bd + ad * bu;
    ull midu = mid >> 30;
    ull midd = mid & M30;
    return modulo(au * bu * 2 + midu + (midd
         << 31) + ad * bd);
```

#### 8.6 LCP

```
vi lcp(const vi& s, const vi& sa) {
  int n = SZ(s), h = 0;
  vi rnk(n), lcp(n - 1);
  REP(i, n) rnk[sa[i]] = i;
  REP(i, n) {
  h -= (h > 0);
}
```

lcp[rnk[i] - 1] = h;

return lcp;

**8.7 SAIS** 

1 // mississippi

if(rnk[i] == 0) continue;

for(; i + h < n && i + h < n; h++) if(s[</pre>

j + h] != s[i + h]) break;

int j = sa[rnk[i] - 1];

```
// 10 7 4 1 0 9 8 6 3 5 2
vi SAIS(string a) {
  int n = SZ(a), m = *max_element(ALL(a)) +
  vi pos(m + 1), x(m), sa(n), val(n), lms;
  for(auto c : a) pos[c + 1]++;
  REP(i, m) pos[i + 1] += pos[i];
  vector<bool> s(n);
  IREP(i, n - 1) s[i] = a[i] != a[i + 1] ? a
       [i] < a[i + 1] : s[i + 1];
  auto ind = [&](const vi& ls){
    fill(ALL(sa), -1);
    auto L = [&](int i) { if(i >= 0 && !s[i
         ]) sa[x[a[i]]++] = i; };
    auto S = [\&](int i) \{ if(i >= 0 \&\& s[i]) \}
          sa[--x[a[i]]] = i; };
    REP(i, m) x[i] = pos[i + 1];
    IREP(i, SZ(ls)) S(ls[i]);
    REP(i, m) x[i] = pos[i];
    L(n - 1);
    REP(i, n) L(sa[i] - 1);
    REP(i, m) x[i] = pos[i + 1];
    IREP(i, n) S(sa[i] - 1);
  auto ok = [&](int i) { return i == n || (!)
       s[i - 1] && s[i]); };
  auto same = [&](int i,int j) {
      if(a[i++] != a[j++]) return false;
    } while(!ok(i) && !ok(j));
    return ok(i) && ok(j);
  FOR(i, 1, n) if(ok(i)) lms.pb(i);
  ind(lms);
  if(SZ(lms)) {
    int p = -1, w = 0;
    for(auto v : sa) if(v && ok(v)) {
      if(p != -1 && same(p, v)) w--;
      val[p = v] = w++;
    auto b = lms;
    for(auto& v : b) v = val[v];
    b = SAIS(b);
    for(auto& v : b) v = lms[v];
    ind(b);
  return sa;
```

#### 8.8 KMP

```
1  // abacbaba -> [0, 0, 1, 0, 0, 1, 2, 3]
2  vi  KMP(const vi& a) {
3    int n = SZ(a);
4   vi  k(n);
5   FOR(i, 1, n) {
6    int j = k[i - 1];
7   while(j > 0 && a[i] != a[j]) j = k[j - 1];
8   j += (a[i] == a[j]);
8   k[i] = j;
9  }
10  }
11  return k;
}
```

#### 8.9 wildcard-pattern-matching

```
1 // 0 <= i <= n - m に対し、s[i, i + m) == t
      かどうか
2 // abc*b*a***a
3 // *b*a
4 // 10111011
5 template < class T, class U = modint998244353 > 29
 vector<bool> wildcard_matching(const vector< 30</pre>
      T> &s, const vector<T> &t, T wildcard) { 31
     const int n = s.size(), m = t.size();
     vector<U> s1(n), s2(n), s3(n), t1(m), t2
          (m), t3(m);
     REP(i, n) {
          s1[i] = s[i] == wildcard ? 0 : s[i]
              == 0 ? wildcard : s[i];
          s2[i] = s1[i] * s1[i], s3[i] = s2[i]
                * s1[i];
     REP(j, m) {
          t1[j] = t[m - 1 - j] == wildcard ? 0
               : t[m - 1 - j] == 0 ? wildcard
               : t[m - 1 - j];
          t2[j] = t1[j] * t1[j], t3[j] = t2[j]
                * t1[j];
     vector<U> u13 = convolution(s1, t3);
     vector<U> u22 = convolution(s2, t2);
     vector<U> u31 = convolution(s3, t1);
     vector<bool> res(n - m + 1);
     REP(i, n - m + 1) res[i] = u13[i + m -
          1] - 2 * u22[i + m - 1] + u31[i + m
          - 1] == 0;
     return res;
```

```
int u = sz++;
    SA[u].first_pos = SA[u].len = SA[p].len +
    SA[u].cnt = 1;
    while(p != -1 && SA[p].go[c] == 0) {
      SA[p].go[c] = u;
      p = SA[p].link;
12
    if(p == -1) {
      SA[u].link = 0;
15
      return u;
    int q = SA[p].go[c];
    if(SA[p].len + 1 == SA[q].len) {
      SA[u].link = q;
      return u:
    int x = sz++;
    SA[x] = SA[q];
    SA[x].cnt = 0;
    SA[x].len = SA[p].len + 1;
```

SA[q].link = SA[u].link = x;

while(p != -1 && SA[p].go[c] == q) {

4 void sa\_init() { SA[0].link = -1, SA[0].len

= 0, sz = 1;

int sa extend(int p, int c) {

#### 8.11 manacher

return u;

SA[p].go[c] = x;

p = SA[p].link;

#### 8.10 SAM

# ACM ICPC Judge Test NTHU Cocacolastic

#### C++ Resource Test

```
#include <bits/stdc++.h>
using namespace std;

namespace system_test {

const size_t KB = 1024;
const size_t MB = KB * 1024;
const size_t GB = MB * 1024;
```

```
chrono::duration<double> diff = end -
10 size t block size, bound;
                                                          begin;
  void stack size dfs(size t depth = 1) {
                                                     return diff.count():
   if (depth >= bound)
                                                   void runtime_error_1() {
    int8_t ptr[block_size]; // 若無法編譯將
                                                    // Segmentation fault
         block size 改成常數
                                                    int *ptr = nullptr;
    memset(ptr, 'a', block_size);
                                                     *(ptr + 7122) = 7122;
    cout << depth << endl;</pre>
                                                 42 }
   stack_size_dfs(depth + 1);
                                                   void runtime_error_2() {
                                                    // Segmentation fault
  void stack_size_and_runtime_error(size_t
                                                    int *ptr = (int *)memset;
       block size, size t bound = 1024) {
                                                     *ptr = 7122;
    system test::block size = block size;
                                                 48
    system_test::bound = bound;
    stack size dfs();
                                                   void runtime_error_3() {
                                                    // munmap_chunk(): invalid pointer
                                                    int *ptr = (int *)memset;
  double speed(int iter num) {
                                                     delete ptr;
    const int block_size = 1024;
                                                 54
    volatile int A[block_size];
    auto begin = chrono::high resolution clock
                                                   void runtime_error_4() {
         ::now();
                                                    // free(): invalid pointer
    while (iter num--)
                                                    int *ptr = new int[7122];
      for (int j = 0; j < block_size; ++j)
                                                     ptr += 1;
                                                     delete[] ptr;
    auto end = chrono::high resolution clock::
         now();
```

```
63 void runtime_error_5() {
    // maybe illegal instruction
    int a = 7122, b = 0;
    cout << (a / b) << endl;</pre>
  void runtime error 6() {
    // floating point exception
    volatile int a = 7122, b = 0;
    cout << (a / b) << endl;</pre>
73 }
  void runtime error 7() {
    // call to abort.
    assert(false);
78 }
80 } // namespace system test
82 #include <sys/resource.h>
void print_stack_limit() { // only work in
       Linux
    struct rlimit 1;
    getrlimit(RLIMIT STACK, &1);
    cout << "stack_size = " << l.rlim_cur << "</pre>
          byte" << endl;</pre>
87 }
```