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#### 1 Basic

#### 1.1 template

#### 1.2 vimrc

```
1 se nu ai hls et ru ic is sc cul
2 se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
3 syntax on
4 hi cursorline cterm=none ctermbg=89
5 set bg=dark
6 inoremap {<CR> {<CR>}<Esc>ko<tab>
```

#### 2 Data-Structure

#### 2.1 CDO

```
i void CDQ(int 1, int r) {
   if(1 + 1 == r) return;
   int mid = (1 + r) / 2;
   CDQ(1, mid), CDQ(mid, r);
   int i = 1:
   FOR(j, mid, r) {
     const Q& q = qry[j];
     while(i < mid && qry[i].x >= q.x) {
       if(qry[i].id == -1) fenw.add(qry[i].y,
             qry[i].w);
       i++:
     if(q.id >= 0) ans[q.id] += q.w * fenw.
          sum(q.y, sz - 1);
   FOR(p, 1, i) if (qry[p].id == -1) fenw.add(
        qry[p].y, -qry[p].w);
   inplace_merge(qry.begin() + 1, qry.begin()
         + mid, qry.begin() + r, [](const Q& a 12
        , const Q& b) {
     return a.x > b.x;
   });
```

#### 2.2 CHT

```
1 struct line t {
     mutable 1\overline{1} k, m, p;
    bool operator<(const line_t& o) const {</pre>
          return k < o.k; }</pre>
    bool operator<(ll x) const { return p < x;</pre>
   template < bool MAX >
  struct CHT : multiset<line t, less<>>> {
     const 11 INF = 1e18L;
     bool isect(iterator x, iterator y) {
       if(y == end()) return x->p = INF, 0;
       if(x->k == y->k) {
         x->p = (x->m > y->m ? INF : -INF);
         x \rightarrow p = floor div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y)
              ->k): // see Math
       return x - p >= y - p;
     void add line(ll k, ll m) {
      if(!MAX) k = -k, m = -m;
       auto z = insert(\{k, m, 0\}), y = z++, x =
       while(isect(y, z)) z = erase(z);
       if(x != begin() && isect(--x, y)) isect(
            x, y = erase(y));
       while((y = x) != begin() && (--x)->p >=
            y->p) isect(x, erase(y));
     11 get(11 x) {
       assert(!empty());
       auto 1 = *lower bound(x);
       return (1.k * x + 1.m) * (MAX ? +1 : -1)
30 };
```

#### 2.3 DLX

```
1 struct DLX {
   int n, m, tot, ans;
   vi first, siz, L, R, U, D, col, row, stk;
   DLX(int _n, int _m) : n(_n), m(_m), tot(_m
       ) {
     int sz = n * m:
      first = siz = L = R = U = D = col = row
          = stk = vi(sz);
     REP(i, m + 1) {
       L[i] = i - 1, R[i] = i + 1;
       U[i] = D[i] = i;
     L[0] = m, R[m] = 0;
   void insert(int r, int c) { // (r, c) is 1 10
     col[++tot] = c, row[tot] = r, ++siz[c];
     D[tot] = D[c], U[D[c]] = tot, U[tot] = c
          , D[c] = tot;
```

#### 2.4 lazvsegtree

if(!first[r]) first[r] = L[tot] = R[tot] 14|

L[R[tot] = R[first[r]]] = tot;

#define TRAV(i, X, j) for(i = X[j]; i != j

R[L[tot] = first[r]] = tot;

L[R[c]] = L[c], R[L[c]] = R[c];

TRAV(i, D, c) TRAV(j, R, i) {

D[U[D[j]] = U[j]] = D[j];

TRAV(i, U, c) TRAV(j, L, i) {

if(!R[0]) return ans = dep, true;

TRAV(j, R, i) remove(col[j]);

if(dance(dep + 1)) return true;

TRAV(j, L, i) recover(col[j]);

TRAV(i, R, 0) if(siz[i] < siz[c]) c = i;

return vi(stk.begin() + 1, stk.begin() +

U[D[j]] = D[U[j]] = j;

15

17

28

33

34

38

53

68

69

70

71

72

= tot:

; i = X[i])

void remove(int c) {

siz[col[j]]--;

void recover(int c) {

siz[col[j]]++;

bool dance(int dep) {

int i, j, c = R[0];

stk[dep] = row[i];

if(!dance(1)) return {};

L[R[c]] = R[L[c]] = c;

int i, j;

int i, j;

remove(c);

recover(c);

vi solve() {

return false;

ans);

TRAV(i, D, c) {

29

30

31

38

else {

```
if(k < size) lz[k] = composition(f, lz[k</pre>
       1);
void push(int k) {
  all apply(k << 1, lz[k]);
  all apply(k \ll 1 \mid 1, lz[k]);
  lz[k] = id();
lazy segtree(int n) : lazy segtree(vector
     <S>(_n, e())) {}
lazy segtree(const vector<S>& v) : n(SZ(v)
  log = __lg(2 * n - 1), size = 1 << log;
  d.resize(size * 2, e());
  lz.resize(size, id());
  REP(i, n) d[size + i] = v[i];
  for(int i = size - 1; i; i--) update(i);
void set(int p, S x) {
  p += size;
  for(int i = log; i; --i) push(p >> i);
  for(int i = 1; i <= log; ++i) update(p</pre>
       >> i):
S get(int p) {
  p += size;
  for(int i = log; i; i--) push(p >> i);
  return d[p];
S prod(int 1, int r) {
  if(1 == r) return e();
  1 += size; r += size;
  for(int i = log; i; i--) {
    if(((1 \Rightarrow i) \iff i) != 1) push(1 \Rightarrow i);
    if(((r >> i) << i) != r) push(r >> i);
  S sml = e(), smr = e();
  while(1 < r) {
    if(1 \& 1) sml = op(sml, d[1++]);
    if(r & 1) smr = op(d[--r], smr);
    1 >>= 1, r >>= 1;
  return op(sml, smr);
S all prod() const { return d[1]; }
void apply(int p, F f) {
  p += size;
  for(int i = log; i; i--) push(p >> i);
  d[p] = mapping(f, d[p]);
  for(int i = 1; i <= log; i++) update(p</pre>
       >> i);
void apply(int 1, int r, F f) {
  if(1 == r) return;
  1 += size; r += size;
  for(int i = log; i; i--) {
    if(((1 \Rightarrow i) \iff i) != 1) push(1 \Rightarrow i);
    if(((r \rightarrow i) << i) != r) push((r - 1)
         >> i):
    int 12 = 1, r2 = r;
    while(1 < r)  {
      if(1 & 1) all apply(1++, f);
```

if(r & 1) all\_apply(--r, f);

```
1 >>= 1, r >>= 1;
         1 = 12:
         r = r2;
       for(int i = 1; i <= log; i++) {</pre>
         if(((1 >> i) << i) != 1) update(1 >> i
         if(((r >> i) << i) != r) update((r -
             1) >> i);
82
     template < class G> int max_right(int 1, G g
       assert(0 <= 1 && 1 <= n && g(e()));
       if(1 == n) return n;
      1 += size:
       for(int i = log; i; i--) push(1 >> i);
      S sm = e();
       do {
         while(!(1 & 1)) 1 >>= 1;
         if(!g(op(sm, d[1]))) {
           while(1 < size) {</pre>
             push(1);
             1 <<= 1;
             if(g(op(sm, d[1]))) sm = op(sm, d[
                  1++]);
           return 1 - size;
         sm = op(sm, d[1++]);
       } while((1 & -1) != 1);
       return n:
     template < class G> int min left(int r, G g)
       assert(0 <= r && r <= n && g(e()));
      if(r == 0) return 0;
107
      r += size;
       for(int i = log; i >= 1; i--) push((r -
            1) \gg i);
       S sm = e();
110
       do {
         while(r > 1 \&\& (r \& 1)) r >>= 1;
         if(!g(op(d[r], sm))) {
           while(r < size) {</pre>
             push(r);
             r = r << 1 | 1;
             if(g(op(d[r], sm))) sm = op(d[r])
                  --1, sm);
           return r + 1 - size;
119
         sm = op(d[r], sm);
       } while((r & -r) != r);
122
       return 0;
125 };
  2.5 LCT
```

1 template < class S,</pre>

```
S (*e)(),
                                                         auto attach = [&](Node* p, bool side,
         S (*op)(S, S),
                                                              Node* c) {
         S (*reversal)(S),
                                                           (side ? p->r : p->1) = c:
         class F,
                                                           pull(p);
         F (*id)(),
                                                           if(c != nullptr) c->p = p;
         S (*mapping)(F, S),
         F (*composition)(F, F)>
                                                         Node *p = v - p, *g = p - p;
struct lazy lct {
                                                         bool rgt = (p->r == v);
 struct Node {
                                                         bool rt = p->is root();
   S \text{ val} = e(), \text{ sum} = e();
                                                         attach(p, rgt, (rgt ? v->l : v->r));
    F lz = id();
                                                         attach(v, !rgt, p);
    bool rev = false;
                                                         if(!rt) attach(g, (g->r == p), v);
    int sz = 1;
                                                         else v \rightarrow p = g;
                                                  60
    Node *1 = nullptr, *r = nullptr, *p =
                                                  61
         nullptr;
                                                       void splay(Node* v) {
    Node() {}
                                                  63
                                                         push(v);
    Node(const S& s) : val(s), sum(s) {}
                                                         while(!v->is_root()) {
    bool is_root() const { return p ==
                                                           auto p = v - p;
         nullptr || (p->1 != this && p->r !=
                                                           auto g = p->p;
         this); }
                                                           if(!p->is_root()) push(g);
                                                           push(p), push(v);
 int n;
                                                           if(!p->is_root()) rotate((g->r == p)
  vector<Node> a;
  lazy_lct() : n(0) {}
                                                           rotate(v);
  explicit lazy lct(int n) : lazy lct(
       vector<S>(_n, e())) {}
  explicit lazy lct(const vector<S>& v) : n(
                                                       void all apply(Node* v, F f) {
       SZ(v)) { REP(i, n) a.eb(v[i]); }
                                                        v->val = mapping(f, v->val), v->sum =
  Node* access(int u) {
                                                              mapping(f, v->sum);
    Node* v = &a[u];
                                                        v->lz = composition(f, v->lz);
    Node* last = nullptr;
    for(Node* p = v; p != nullptr; p = p -> p) 77
                                                       void push(Node* v) {
          splay(p), p->r = last, pull(last =
                                                         if(v->lz != id()) {
                                                           if(v->l != nullptr) all_apply(v->l, v
         p);
    splay(v);
                                                           if(v->r != nullptr) all_apply(v->r, v
    return last;
 void make root(int u) { access(u), a[u].
                                                           v \rightarrow lz = id();
       rev ^= 1, push(&a[u]); }
  void link(int u, int v) { make_root(v), a[
                                                         if(v->rev) {
                                                 83
                                                           swap(v->1, v->r);
       v].p = &a[u]; }
  void cut(int u) {
                                                           if(v->l != nullptr) v->l->rev ^= 1;
    access(u);
                                                           if(v->r != nullptr) v->r->rev ^= 1;
    if(a[u].1 != nullptr) a[u].1->p =
                                                           v->sum = reversal(v->sum);
         nullptr, a[u].1 = nullptr, pull(&a[u 88
                                                           v->rev = false;
 void cut(int u, int v) { make root(u), cut
                                                       void pull(Node* v) {
                                                        v \rightarrow sz = 1:
  bool is_connected(int u, int v) {
                                                         v \rightarrow sum = v \rightarrow val;
    if(u == v) return true;
                                                         if(v->1 != nullptr) {
    return access(u), access(v), a[u].p !=
                                                           push(v->1);
         nullptr;
                                                           v \rightarrow sum = op(v \rightarrow 1 \rightarrow sum, v \rightarrow sum);
                                                           v->sz += v->l->sz:
 int get lca(int u, int v) { return access(
       u), access(v) - &a[0]; }
                                                         if(v->r != nullptr) {
  void set(int u, const S& s) { access(u), a 100
                                                           push(v->r);
       [u].val = s, pull(&a[u]); }
                                                           v \rightarrow sum = op(v \rightarrow sum, v \rightarrow r \rightarrow sum);
 S get(int u) { return access(u), a[u].val; 102
                                                           v \rightarrow sz += v \rightarrow r \rightarrow sz;
                                                 103
  void apply(int u, int v, const F& f) {
       make root(u), access(v), all apply(&a[ 105 ] };
       v], f), push(&a[v]); }
 S prod(int u, int v) { return make_root(u)
       , access(v), a[v].sum; }
  void rotate(Node* v) {
```

#### 2.6 LiChao

== (p->r == v) ? p : v);

->1z);

```
1 struct LiChao { // min
    int n;
    vector<pll> seg;
    LiChao(int _n) : n(_n) {
      seg.assign(4 * n + 5, pll(0, INF));
    11 cal(pll line, ll x) { return line.F * x
          + line.S: }
    void insert(int 1, int r, int id, pll line
      if(1 == r) {
        if(cal(line, 1) < cal(seg[id], 1)) seg</pre>
             [id] = line:
13
      int mid = (1 + r) / 2;
      if(line.F > seg[id].F) swap(line, seg[id
      if(cal(line, mid) <= cal(seg[id], mid))</pre>
        seg[id] = line;
16
17
        insert(1, mid, id * 2, seg[id]);
18
19
       else insert(mid + 1, r, id * 2 + 1, line
           );
     11 query(int 1, int r, int id, ll x) {
      if (x < 1 \mid | x > r) return INF;
      if(1 == r) return cal(seg[id], x);
23
      int mid = (1 + r) / 2;
      11 \text{ val} = 0;
      if(x <= mid) val = query(1, mid, id * 2,</pre>
       else val = query(mid + 1, r, id * 2 + 1,
       return min(val, cal(seg[id], x));
28
29
30 };
```

#### 2.7 rect-add-rect-sum

```
i template < class Int. class T>
2 struct RectangleAddRectangleSum {
    struct AQ { Int xl, xr, yl, yr; T val; };
    struct SQ { Int xl, xr, yl, yr; };
    vector<AO> add gry;
    vector<SQ> sum qry;
    // A[x][y] += val for(x, y) in [xl, xr) *
    void add rectangle(Int xl, Int xr, Int yl,
          Int yr, T val) { add_qry.pb({xl, xr,
         yl, yr, val}); }
    // Get sum of A[x][y] for (x, y) in [xl, xr]
         ) * [yl, yr)
    void add_query(Int xl, Int xr, Int yl, Int
          yr) { sum_qry.pb({xl, xr, yl, yr}); }
    vector<T> solve() {
      vector<Int> ys;
12
      for(auto &a : add qry) ys.pb(a.yl), ys.
           pb(a.yr);
```

```
vs = sort unique(vs);
      const int Y = SZ(ys);
      vector<tuple<Int, int, int>> ops:
      REP(q, SZ(sum_qry)) {
        ops.eb(sum_qry[q].xl, 0, q);
        ops.eb(sum qry[q].xr, 1, q);
      REP(q, SZ(add_qry)) {
        ops.eb(add_qry[q].xl, 2, q);
        ops.eb(add_qry[q].xr, 3, q);
      sort(ALL(ops));
      fenwick<T> b00(Y), b01(Y), b10(Y), b11(Y
      vector<T> ret(SZ(sum_qry));
      for(auto o : ops) {
        int qtype = get<1>(o), q = get<2>(o);
        if(qtype >= 2) {
          const auto& query = add_qry[q];
          int i = lower_bound(ALL(ys), query.
              yl) - ys.begin();
          int j = lower_bound(ALL(ys), query.
              yr) - ys.begin();
          T x = get<0>(o);
          T yi = query.yl, yj = query.yr;
          if(qtype & 1) swap(i, j), swap(yi,
          b00.add(i, x * yi * query.val);
          b01.add(i, -x * query.val);
          b10.add(i, -yi * query.val);
          b11.add(i, query.val);
          b00.add(j, -x * yj * query.val);
          b01.add(j, x * query.val);
          b10.add(j, yj * query.val);
          b11.add(j, -query.val);
        } else {
          const auto& query = sum_qry[q];
          int i = lower_bound(ALL(ys), query.
               yl) - ys.begin();
          int j = lower_bound(ALL(ys), query.
              yr) - ys.begin();
          T x = get<0>(o);
          T yi = query.yl, yj = query.yr;
          if(qtype & 1) swap(i, j), swap(yi,
          ret[q] += b00.get(i - 1) + b01.get(i
                -1) * yi + b10.get(i - 1) * x 10
               + b11.get(i - 1) * x * yi;
          ret[q] = b00.get(j - 1) + b01.get(j
                -1) * yj + b10.get(j - 1) * x _{12}
              + b11.get(j - 1) * x * yj;
     return ret;
58 };
  2.8 rollback-dsu
1 struct RollbackDSU {
```

int n; vi sz, tag;

void init(int \_n) {

vector<tuple<int, int, int, int>> op;

## segtree-beats

n = n;

sz.assign(n, -1);

int leader(int x) {

while(sz[x] >= 0) x = sz[x];

x = leader(x), y = leader(y);

if(-sz[x] < -sz[y]) swap(x, y);

void add\_tag() { tag.pb(sz(op)); }

int z = tag.back(); tag.ppb();

int size(int x) { return -sz[leader(x);] }

auto [x, sx, y, sy] = op.back(); op.

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bool merge(int x, int y) {

if(x == y) return false;

op.eb(x, sz[x], y, sz[y]);

sz[x] += sz[y]; sz[y] = x;

tag.clear():

return x;

return true;

void rollback() {

sz[x] = sx;

sz[y] = sy;

29

31 };

while(sz(op) > z) {

ppb();

```
1 struct segtree beats {
   static constexpr ll INF = numeric limits
        ll>::max() / 2.1;
   struct alignas(32) Node {
     11 \text{ sum} = 0, g1 = 0, 11 = 0;
     11 g2 = -INF, gc = 1, 12 = INF, 1c = 1,
          add = 0;
   11 n, log;
   vector<Node> v;
   segtree beats() {}
   segtree_beats(int _n) : segtree_beats(
        vector<ll>( n)) {}
   segtree beats(const vector<11>& vc) {
     n = 1, log = 0;
     while(n < SZ(vc)) n <<= 1, log++;
     v.resize(2 * n);
     REP(i, SZ(vc)) v[i + n].sum = v[i + n].
          g1 = v[i + n].11 = vc[i];
     for(ll i = n - 1; i; --i) update(i);
   void range_chmin(int 1, int r, 11 x) {
        inner_apply<1>(l, r, x); }
   void range_chmax(int 1, int r, 11 x) {
        inner_apply<2>(1, r, x); }
   void range_add(int 1, int r, 11 x) {
        inner_apply<3>(1, r, x); }
   void range_update(int 1, int r, 11 x) {
        inner_apply<4>(1, r, x); }
   11 range min(int 1, int r) { return
        inner_fold<1>(l, r); }
```

```
inner fold\langle 2 \rangle (1, r); \}
11 range_sum(int 1, int r) { return
                                              83
     inner_fold<3>(1, r); }
void update(int k) {
                                              84
  Node& p = v[k];
                                              85
  Node& 1 = v[k * 2];
  Node& r = v[k * 2 + 1];
  p.sum = 1.sum + r.sum;
  if(1.g1 == r.g1) {
    p.g1 = 1.g1;
    p.g2 = max(1.g2, r.g2);
    p.gc = 1.gc + r.gc;
  } else {
    bool f = 1.g1 > r.g1;
    p.g1 = f ? l.g1 : r.g1;
    p.gc = f ? 1.gc : r.gc;
    p.g2 = max(f ? r.g1 : l.g1, f ? l.g2 : 96
          r.g2);
  if(1.11 == r.11) {
    p.11 = 1.11;
    p.12 = min(1.12, r.12);
                                              101
    p.lc = 1.lc + r.lc:
                                              102
  } else {
                                              103
    bool f = 1.11 < r.11;
                                              104
    p.11 = f ? 1.11 : r.11;
    p.lc = f ? 1.lc : r.lc;
    p.12 = min(f ? r.11 : 1.11, f ? 1.12 : 106
          r.12);
                                              108
void push add(int k, ll x) {
  Node& p = v[k];
  p.sum += x << (log + builtin clz(k) -
                                              111
  p.g1 += x, p.11 += x;
                                              112
  if(p.g2 != -INF) p.g2 += x;
  if(p.12 != INF) p.12 += x;
                                              113
  p.add += x;
                                              115
void push_min(int k, ll x) {
                                             116
  Node& p = v[k];
                                             117
  p.sum += (x - p.g1) * p.gc;
                                             118
  if(p.11 == p.g1) p.11 = x;
                                             119
  if(p.12 == p.g1) p.12 = x;
                                             120
  p.g1 = x;
                                              121
void push max(int k, ll x) {
                                              122
  Node& p = v[k];
                                              123
  p.sum += (x - p.11) * p.1c;
                                              124
  if(p.g1 == p.11) p.g1 = x;
                                             125
  if(p.g2 == p.l1) p.g2 = x;
                                              126
  p.11 = x;
                                              127
                                              128
void push(int k) {
                                              129
  Node& p = v[k];
                                              130
  if(p.add != 0) {
                                             131
    push_add(k * 2, p.add);
                                             132
    push add(k * 2 + 1, p.add);
   p.add = 0;
                                              133
  if(p.g1 < v[k * 2].g1) push_min(k * 2, p 134
  if(p.11 > v[k * 2].11) push max(k * 2, p 136
```

11 range max(int 1, int r) { return

```
if(p.g1 < v[k * 2 + 1].g1) push_min(k *</pre>
       2 + 1, p.g1);
  if(p.11 > v[k * 2 + 1].11) push max(k *
       2 + 1, p.11);
void subtree chmin(int k, ll x) {
  if(v[k].g1 <= x) return;</pre>
  if(v[k].g2 < x) {
    push_min(k, x);
    return;
  push(k);
  subtree chmin(k * 2, x), subtree chmin(k
        * 2 + 1, x);
  update(k);
void subtree chmax(int k, ll x) {
  if(x <= v[k].l1) return;</pre>
  if(x < v[k].12) {
    push max(k, x);
    return;
  push(k);
  subtree_chmax(k * 2, x), subtree_chmax(k
        *^{2} + 1, x);
  update(k);
template<int cmd>
inline void _apply(int k, ll x) {
  if constexpr(cmd == 1) subtree chmin(k,
  if constexpr(cmd == 2) subtree_chmax(k,
  if constexpr(cmd == 3) push_add(k, x);
  if constexpr(cmd == 4) subtree chmin(k,
       x), subtree_chmax(k, x);
template<int cmd>
void inner_apply(int 1, int r, ll x) {
  if(1 == r) return;
  1 += n, r += n;
  for(int i = log; i >= 1; i--) {
    if(((1 >> i) << i) != 1) push(1 >> i);
    if(((r \rightarrow i) << i) != r) push((r - 1)
         >> i);
    int 12 = 1, r2 = r;
    while (1 < r) {
      if(1 & 1) apply<cmd>(1++, x);
      if(r \& 1) apply < cmd > (--r, x);
      1 >>= 1, r >>= 1;
    1 = 12, r = r2;
  for(int i = 1; i <= log; i++) {</pre>
    if(((1 >> i) << i) != 1) update(1 >> i
    if(((r >> i) << i) != r) update((r -
         1) >> i);
template<int cmd>
```

```
inline 11 e() {
       if constexpr(cmd == 1) return INF;
       if constexpr(cmd == 2) return -INF;
142
    template<int cmd>
     inline void op(ll& a, const Node& b) {
      if constexpr(cmd == 1) a = min(a, b.l1);
      if constexpr(cmd == 2) a = max(a, b.g1);
147
       if constexpr(cmd == 3) a += b.sum;
    template<int cmd>
    11 inner_fold(int 1, int r) {
      if(1 == r) return e<cmd>();
      1 += n, r += n;
       for(int i = log; i >= 1; i--) {
         if(((1 >> i) << i) != 1) push(1 >> i);
         if(((r >> i) << i) != r) push((r - 1)
              >> i):
156
157
       11 1x = e < cmd > (), rx = e < cmd > ();
       while (1 < r) {
         if(1 & 1) op<cmd>(lx, v[1++]);
         if(r & 1) op<cmd>(rx, v[--r]);
         1 >>= 1, r >>= 1;
161
162
       if constexpr(cmd == 1) lx = min(lx, rx);
      if constexpr(cmd == 2) lx = max(lx, rx);
      if constexpr(cmd == 3) lx += rx;
       return lx;
167
168 };
```

#### 2.10 segtree

```
1 template < class S, S (*e)(), S (*op)(S, S)>
2 struct segtree {
   int n, size, log;
   vector<S> st;
   void update(int v) { st[v] = op(st[v <<</pre>
        1], st[v << 1 | 1]); }
   segtree(int n) : segtree(vector<S>( n, e
        ())) {}
   segtree(const vector<S>& a): n(sz(a)) {
     log = lg(2 * n - 1), size = 1 << log;
     st.resize(size << 1, e());</pre>
     REP(i, n) st[size + i] = a[i];
     for(int i = size - 1; i; i--) update(i);
   void set(int p, S val) {
     st[p += size] = val;
     for(int i = 1; i <= log; ++i) update(p</pre>
          >> i);
   S get(int p) const {
     return st[p + size];
   S prod(int 1, int r) const {
     assert(0 <= 1 && 1 <= r && r <= n);
     S sml = e(), smr = e();
     1 += size, r += size;
     while(1 < r) {
       if(1 & 1) sml = op(sml, st[1++]);
```

```
1 >>= 1;
    r >>= 1:
  return op(sml, smr);
S all prod() const { return st[1]; }
template < class F> int max right(int 1, F f
  assert(0 <= 1 && 1 <= n && f(e()));
  if(1 = n) return n;
  1 += size;
  S sm = e();
    while(~1 & 1) 1 >>= 1;
    if(!f(op(sm, st[1]))) {
      while(1 < size) {</pre>
        1 <<= 1;
        if(f(op(sm, st[1]))) sm = op(sm,
             st[1++]);
      return 1 - size;
    sm = op(sm, st[1++]);
  } while((1 & -1) != 1);
  return n;
template < class F> int min left(int r, F f)
  assert(0 <= r \&\& r <= n \&\& f(e()));
  if(r == 0) return 0;
  r += size;
  S sm = e();
  do {
    while(r > 1 && (r & 1)) r >>= 1;
    if(!f(op(st[r], sm))) {
      while(r < size) {</pre>
        r = r << 1 | 1;
        if(f(op(st[r], sm))) sm = op(st[r
             --], sm);
      return r + 1 - size;
```

if(r & 1) smr = op(st[--r], smr);

#### 2.11 sparse-table

return 0:

sm = op(st[r], sm);

} while((r & -r) != r);

```
template < class T, T (*op)(T, T)>
struct sparse_table {
   int n;
   vector < vector < T >> b;
   sparse_table(const vector < T >& a) : n(SZ(a)
      ) {
   int lg = __lg(n) + 1;
   b.resize(lg); b[0] = a;
   FOR(j, 1, lg) {
      b[j].resize(n - (1 << j) + 1);
   }
}</pre>
```

```
REP(i, n - (1 << j) + 1) b[j][i] = op( 46|
                                                            while(1 < q1) DL(1++);</pre>
                                                            while(r > qr) DR(--r);
             b[j - 1][i], b[j - 1][i + (1 << (j 47)]
              - 1))]);
                                                            ans[id] = cnt:
11
12
                                                         return ans;
                                                   50
    T prod(int from, int to) {
13
                                                   51
      int \lg = \lg(to - from + 1);
                                                   52 };
      return op(b[lg][from], b[lg][to - (1 <<
           lg) + 1]);
16
17 };
```

2.12 static-range-inversion

#### 2.13 static-range-lis

```
1 struct static range inversion {
   int sz;
   vi a, L, R;
   vector<11> ans;
    static_range_inversion(vi _a) : a(_a) {
      _a = sort_unique(_a);
      REP(i, SZ(a)) a[i] = lower bound(ALL(a))
          , a[i]) - _a.begin();
      sz = SZ(_a);
    void add_query(int 1, int r) { L.push_back
         (1), R.push back(r); }
                                                 12
    vector<ll> solve() {
                                                 13
      const int q = SZ(L);
      const int B = max(1.0, SZ(a) / sqrt(q));
      vi ord(q);
      iota(ALL(ord), 0);
      sort(ALL(ord), [&](int i, int j) {
        if(L[i] / B == L[j] / B) {
          return L[i] / B & 1 ? R[i] > R[j] :
               R[i] < R[j];
        return L[i] < L[j];</pre>
      });
      ans.resize(q);
      fenwick<ll> fenw(sz + 1);
      11 cnt = 0;
      auto AL = [&](int i) {
        cnt += fenw.sum(0, a[i] - 1);
        fenw.add(a[i], +1);
      };
                                                 26
      auto AR = [&](int i) {
                                                 27
        cnt += fenw.sum(a[i] + 1, sz);
                                                 28
        fenw.add(a[i], +1);
                                                 29
      auto DL = [&](int i) {
        cnt -= fenw.sum(0, a[i] - 1);
                                                 30
        fenw.add(a[i], -1);
                                                 31
                                                 32
      auto DR = [&](int i) {
                                                 33
       cnt -= fenw.sum(a[i] + 1, sz);
        fenw.add(a[i], -1);
      int 1 = 0, r = 0;
      REP(i, q) {
        int id = ord[i], ql = L[id], qr = R[id 37
        while(l > ql) AL(--1);
        while(r < qr) AR(r++);</pre>
```

```
1 #define MEM(a, x, n) memset(a, x, sizeof(int
      ) * n)
using I = int*;
3 struct static range lis {
   int n, ps = 0;
   I invp, res_monge, pool;
   vector<vector<pii>>> qry;
   vi ans;
   static_range_lis(vi a) : n(SZ(a)), qry(n +
      // a must be permutation of [0, n)
     pool = (I) malloc(sizeof(int) * n * 100)
     invp = A(n), res_monge = A(n);
     REP(i, n) invp[a[i]] = i;
   inline I A(int x) { return pool + (ps += x
   void add_query(int 1, int r) { qry[1].pb({
        r, SZ(ans)}), ans.pb(r - 1); }
    void unit monge mult(I a, I b, I r, int n)
     if(n == 2){
       if(!a[0] && !b[0]) r[0] = 0, r[1] = 1;
       else r[0] = 1, r[1] = 0;
       return;
     if(n == 1) return r[0] = 0, void();
     int lps = ps, d = n / 2;
     I a1 = A(d), a2 = A(n - d), b1 = A(d),
          b2 = A(n - d);
     I mpa1 = A(d), mpa2 = A(n - d), mpb1 = A
          (d), mpb2 = A(n - d);
     int p[2] = {};
     REP(i, n) {
       if(a[i] < d) a1[p[0]] = a[i], mpa1[p]
            [0]++] = i;
       else a2[p[1]] = a[i] - d, mpa2[p[1]++]
     p[0] = p[1] = 0;
     REP(i, n) {
       if(b[i] < d) b1[p[0]] = b[i], mpb1[p]
            [0]++]=i;
       else b2[p[1]] = b[i] - d, mpb2[p[1]++]
             = i;
     I c1 = A(d), c2 = A(n - d);
     unit_monge_mult(a1, b1, c1, d),
          unit_monge_mult(a2, b2, c2, n - d);
     I cpx = A(n), cpy = A(n), cqx = A(n),
          cqy = A(n);
```

REP(i, d) cpx[mpa1[i]] = mpb1[c1[i]],

```
cpv[mpa1[i]]=0;
  REP(i, n - d) cpx[mpa2[i]] = mpb2[c2[i]]
      ]], cpy[mpa2[i]]=1;
  REP(i, n) r[i] = cpx[i];
  REP(i, n) cqx[cpx[i]] = i, cqy[cpx[i]] =
        cpv[i];
  int hi = n, lo = n, his = 0, los = 0;
  REP(i, n) {
    if(cqy[i] ^ (cqx[i] >= hi)) his--;
    while(hi > 0 && his < 0) {
      if(cpy[hi] ^ (cpx[hi] > i)) his++;
    while(lo > 0 && los <= 0) {</pre>
      lo--:
      if(cpy[lo] ^ (cpx[lo] >= i)) los++;
    if(los > 0 \&\& hi == lo) r[lo] = i;
    if(cqy[i] ^ (cqx[i] >= lo)) los--;
 ps = 1ps;
void subunit monge mult(I a, I b, I c, int 110
  int lps = ps;
 I za = A(n), zb = A(n), res = A(n), vis 113
      = A(n), mpa = A(n), mpb = A(n), rb = 114
  MEM(vis, 0, n), MEM(mpa, -1, n), MEM(mpb
       , -1, n), MEM(rb, -1, n);
  int ca = n;
 IREP(i, n) if(a[i] != -1) vis[a[i]] = 1,
       za[--ca] = a[i], mpa[ca] = i;
 IREP(i, n) if(!vis[i]) za[--ca] = i;
  MEM(vis, -1, n);
 REP(i, n) if(b[i] != -1) vis[b[i]] = i;
 ca = 0:
 REP(i, n) if(vis[i] != -1) mpb[ca] = i,
 rb[vis[i]] = ca++;
REP(i, n) if(rb[i] == -1) rb[i] = ca++;
  REP(i, n) zb[rb[i]] = i;
 unit_monge_mult(za, zb, res, n);
 MEM(c, -1, n):
 REP(i, n) if(mpa[i] != -1 && mpb[res[i]]
       != -1) c[mpa[i]] = mpb[res[i]];
void solve(I p, I ret, int n) {
 if(n == 1) return ret[0] = -1, void();
 int lps = ps, d = n / 2;
 I pl = A(d), pr = A(n - d);
 REP(i, d) pl[i] = p[i];
 REP(i, n - d) pr[i] = p[i + d];
 I vis = A(n); MEM(vis, -1, n);
  REP(i, d) vis[pl[i]] = i;
 I tl = A(d), tr = A(n - d), mpl = A(d),
       mpr = A(n - d);
  int ca = 0;
  REP(i, n) if(vis[i] != -1) mpl[ca] = i
      tl[vis[i]] = ca++;
  ca = 0; MEM(vis, -1, n);
  REP(i, n - d) vis[pr[i]] = i;
  REP(i, n) if(vis[i] != -1) mpr[ca] = i,
       tr[vis[i]] = ca++:
  I vl = A(d), vr = A(n - d);
```

```
solve(t1, v1, d), solve(tr, vr, n - d); 32
  I sl = A(n), sr = A(n);
  iota(sl, sl + n, 0); iota(sr, sr + n, 0)
  REP(i, d) sl[mpl[i]] = (vl[i] == -1 ? -1
       : mpl[vl[i]]);
  REP(i, n - d) sr[mpr[i]] = (vr[i] == -1)
       ? -1 : mpr[vr[i]]);
  subunit monge mult(sl, sr, ret, n);
  ps = lps;
vi solve() {
                                             42
  solve(invp, res_monge, n);
                                             43
  vi fenw(n + 1);
                                             44
  IREP(i, n) {
                                             45 int get_position(Node* v) { // 0-indexed
    if(res monge[i] != -1) {
      for(int p = res_monge[i] + 1; p <= n</pre>
           ; p += p & -p) fenw[p]++;
    for(auto& z : qry[i]){
      auto [id, c] = z;
      for(int p = id; p; p -= p & -p) ans[
                                            52
           cl -= fenw[p];
  free(pool);
  return ans:
```

**2.14** treap

bool rev = false:

void pull(Node\*& v) {

void push(Node\*& v) {

swap(v->1, v->r):

v->rev = false:

push(a), push(b);

} else {

**if**(a->pri > b->pri) {

pull(a); return a;

pull(b); return b;

**if**(v->1) v->1->rev ^= 1;

if(v->r) v->r->rev ^= 1;

Node\* merge(Node\* a, Node\* b) {

a->r = merge(a->r, b);

b->1 = merge(a, b->1);

if(!v) return {NULL, NULL};

**if**(!a | | !b) **return** (a ? a : b);

go pair < Node\*, Node\* > split(Node\* v, int k) {

if(v->rev) {

int sz = 1, pri = rng();

Node \*1 = NULL, \*r = NULL, \*p = NULL;

 $v \rightarrow sz = 1 + size(v \rightarrow 1) + size(v \rightarrow r);$ 

1 struct Node {

// TODO

// TODO

17

#### 2.15 union-of-rectangles

 $if(size(v\rightarrow 1) \rightarrow = k) {$ 

 $v \rightarrow 1 = p.second;$ 

v->r = p.first;

while(v->p != NULL) {  $if(v == v \rightarrow p \rightarrow r)$ 

k++;

 $v = v \rightarrow p;$ 

return k:

} else {

auto  $p = split(v \rightarrow l, k)$ :

if(p.first) p.first->p = NULL;

pull(v); return {p.first, v};

if(p.second) p.second->p = NULL;

int k = (v->1 != NULL ? v->1->sz : 0);

 $if(v\rightarrow p\rightarrow 1 != NULL) k += v\rightarrow p\rightarrow 1\rightarrow sz;$ 

pull(v); return {v, p.second};

auto p = split(v->r, k - size(v->l) - 1)

```
2 // 1 10 1 10
3 // 0 2 0 2
 4 // ans = 84
5 vector<int> vx, vy;
6 struct q { int piv, s, e, x; };
  struct tree {
    vector<int> seg, tag;
    tree(int _n) : seg(_n * 16), tag(_n * 16)
    void add(int ql, int qr, int x, int v, int
          1, int r) {
      if(qr <= 1 || r <= q1) return;
      if(al <= 1 && r <= ar) {
        tag[v] += x;
        if(tag[v] == 0) {
           if(1 != r) seg[v] = seg[2 * v] + seg
                [2 * v + 1];
           else seg[v] = 0;
        } else seg[v] = vx[r] - vx[1];
      } else {
        int mid = (1 + r) / 2;
        add(q1, qr, x, 2 * v, 1, mid);
        add(ql, qr, x, 2 * v + 1, mid, r);
        if(tag[v] == 0 && 1 != r) seg[v] = seg
              [2 * v] + seg[2 * v + 1];
23
24
    int q() { return seg[1]; }
25
26 };
27 int main() {
   int n: cin >> n:
    vector\langle int \rangle x1(n), x2(n), y (n), y2(n);
    for (int i = 0; i < n; i++) {</pre>
```

```
cin >> x1[i] >> x2[i] >> y_[i] >> y2[i];
            // L R D U
       vx.pb(x1[i]), vx.pb(x2[i]);
33
      vy.pb(y_[i]), vy.pb(y2[i]);
34
    vx = sort unique(vx);
    vy = sort unique(vy);
    vector < a > a(2 * n):
    REP(i, n) {
      x1[i] = lower_bound(ALL(vx), x1[i]) - vx
            .begin():
      x2[i] = lower_bound(ALL(vx), x2[i]) - vx
            .begin();
      y_{[i]} = lower_bound(ALL(vy), y_{[i]}) - vy
            .begin();
      y2[i] = lower_bound(ALL(vy), y2[i]) - vy
      a[2 * i] = {y_[i], x1[i], x2[i], +1};
       a[2 * i + 1] = \{y2[i], x1[i], x2[i],
            -1};
45
    sort(ALL(a), [](q a, q b) { return a.piv <</pre>
          b.piv; });
    tree seg(n);
    11 \text{ ans} = 0:
    REP(i, 2 * n) {
      int j = i;
      while(j < 2 * n && a[i].piv == a[j].piv)</pre>
        seg.add(a[j].s, a[j].e, a[j].x, 1, 0,
             vx.size());
        j++;
53
54
      if(a[i].piv + 1 != SZ(vy)) ans += 1LL *
```

#### 2.16 **VEB**

57

58

].piv]);

cout << ans <<  $"\n"$ ;

i = j - 1;

```
i template<int B, typename ENABLE = void>
2 struct VEB {
    constexpr static int K = B / 2, R = (B +
         1) / 2, M = 1 << B, S = 1 << K, MASK =
          (1 << R) - 1:
    array<VEB<R>, S> child;
    VEB < K > act = {}:
    int mn = M, mx = -1:
    bool empty() { return mx < mn; }</pre>
    bool contains(int i) { return find next(i)
          == i; }
    int find_next(int i) { // >=
      if(i <= mn) return mn;</pre>
10
11
      if(i > mx) return M;
12
      int j = i \gg R, x = i \& MASK;
13
      int res = child[j].find next(x);
14
      if(res <= MASK) return (j << R) + res;</pre>
15
      j = act.find next(j + 1);
      return j >= S ? mx : (j << R) + child[j</pre>
           1.find next(0);
```

seg.q() \* (vy[a[i].piv + 1] - vy[a[i]

```
int find prev(int i) { // <=</pre>
   if(i >= mx) return mx;
   if(i < mn) return -1;</pre>
   int j = i \gg R, x = i \& MASK;
   int res = child[j].find prev(x);
   if(res >= 0) return (j << R) + res;
   j = act.find prev(j - 1);
   return j < 0 ? mn : (j << R) + child[j].</pre>
         find prev(MASK);
 void insert(int i) {
   if(i <= mn) {
      if(i == mn) return;
      swap(mn, i);
     if(i == M) mx = mn;
     if(i >= mx) return;
   } else if(i >= mx) {
     if(i == mx) return;
      swap(mx, i);
     if(i <= mn) return;</pre>
   int j = i >> R;
   if(child[j].empty()) act.insert(j);
   child[j].insert(i & MASK);
 void erase(int i) {
   if(i <= mn) {
      if(i < mn) return;</pre>
     i = mn = find next(mn + 1);
     if(i >= mx) {
       if(i > mx) mx = -1;
        return:
   } else if(i >= mx) {
     if(i > mx) return;
     i = mx = find prev(mx - 1);
     if(i <= mn) return;</pre>
   int j = i \gg R;
   child[j].erase(i & MASK);
   if(child[j].empty()) act.erase(j);
 void clear() {
   mn = M, mx = -1, act.clear();
   REP(i, S) child[i].clear();
template<int B>
struct VEB<B, enable if t<(B <= 6)>> {
 constexpr static int M = 1 << B;</pre>
 unsigned long long act = 0;
 bool empty() { return !act; }
 void clear() { act = 0; }
 bool contains(int i) { return find_next(i)
 void insert(int i) { act |= 1ULL << i; }</pre>
 void erase(int i) { act &= ~(1ULL << i); }</pre>
 int find next(int i) {
   ull tmp = act >> i;
   return (tmp ? i + builtin ctzll(tmp) :
         M);
 int find prev(int i) {
   ull tmp = act << (63 - i);
```

#### 2.17 wavelet-tree

struct wavelet\_tree {

vector<T> vals;

1 template < class T>

int n, log;

```
vi sums:
vector<ull> bits:
void set_bit(int i, ull v) { bits[i >> 6]
     |= (v << (i \& 63)); }
int get sum(int i) const { return sums[i
     >> 6] + __builtin_popcountll(bits[i >>
      6] \frac{1}{8} ((\frac{1}{1}ULL << (\frac{1}{1} & 63)) - 1)); }
wavelet_tree(const vector<T>& _v) : n(SZ(
  vals = sort unique( v);
  \log = -\lg(2 * vals.size() - 1);
  bits.resize((log * n + 64) >> 6, 0ULL);
  sums.resize(SZ(bits), 0);
  vi v(SZ(_v)), cnt(SZ(vals) + 1);
  REP(i, SZ(v)) {
    v[i] = lower_bound(ALL(vals), _v[i]) -
          vals.begin();
    cnt[v[i] + 1] += 1;
  partial sum(ALL(cnt) - 1, cnt.begin());
  REP(j, log) {
    for(int i : v) {
      int tmp = i \rightarrow (\log - 1 - j);
      int pos = (tmp >> 1) << (log - j);</pre>
      set bit(j * n + cnt[pos], tmp & 1);
      cnt[pos]++;
    for(int i : v) cnt[(i >> (log - j)) <<</pre>
          (log - j)]--;
  FOR(i, 1, SZ(sums)) sums[i] = sums[i -
       1] + __builtin_popcountll(bits[i -
       1]);
T get_kth(int a, int b, int k) {
  for(int j = 0, ia = 0, ib = n, res = 0;;
        j++) {
    if(j == log) return vals[res];
    int cnt ia = get sum(n * j + ia);
    int cnt_a = get_sum(n * j + a);
    int cnt b = get sum(n * j + b);
    int cnt_ib = get_sum(n * j + ib);
    int ab_zeros = (b - a) - (cnt_b -
         cnt a);
    if(ab_zeros > k) {
      res <<= 1;
      ib -= cnt ib - cnt ia:
      a -= cnt_a - cnt_ia;
      b -= cnt b - cnt ia:
      res = (res << 1) | 1;
```

```
54 };
```

int n, m; // 二分圖左右人數 (0 ~ n-1), (0

ia += (ib - ia) - (cnt ib - cnt ia); 42

a += (ib - a) - (cnt ib - cnt a):

b += (ib - b) - (cnt ib - cnt b);

### 3 Flow-Matching

53

#### 3.1 bipartite-matching

struct bipartite\_matching {

```
\sim m-1)
    vector<vi> g;
    vi lhs, rhs, dist; // i 與 Lhs[i] 配對 (
         Lhs[i] == -1 代表沒有配對)
    bipartite matching(int _n, int _m) : n(_n)
          m(_{m}), g(_{n}), lhs(_{n}, -1), rhs(_{m},
          -1), dist(_n) {}
    void add edge(int u, int v) { g[u].pb(v);
    void bfs() {
      queue<int> q;
      REP(i, n) {
    if(lhs[i] == -1) {
           q.push(i);
           dist[i] = 0;
        } else {
           dist[i] = -1;
      while(!q.empty()) {
        int u = q.front(); q.pop();
        for(auto v : g[u]) {
           if(rhs[v] != -1 && dist[rhs[v]] ==
             dist[rhs[v]] = dist[u] + 1;
             q.push(rhs[v]);
23
24
    bool dfs(int u) {
      for(auto v : g[u]) {
        if(rhs[v] == -1) {
29
           rhs[lhs[u] = v] = u;
           return true:
31
32
33
34
      for(auto v : g[u]) {
        if(dist[rhs[v]] == dist[u] + 1 && dfs(
              rhs[v])) {
           rhs[lhs[u] = v] = u;
           return true;
      return false;
```

#### 3.2 Dinic-LowerBound

int solve() {

```
1 template < class T>
2 struct DinicLowerBound {
    using Maxflow = Dinic<T>;
    int n;
    Maxflow d;
    vector<T> in;
    DinicLowerBound(int _n) : n(_n), d(_n + 2)
         , in( n) {}
    int add_edge(int from, int to, T low, T
         high) {
      assert(0 <= low && low <= high);
      in[from] -= low, in[to] += low;
      return d.add_edge(from, to, high - low);
12
    T flow(int s, int t) {
13
      T sum = 0;
      REP(i, n) {
        if(in[i] > 0) {
          d.add edge(n, i, in[i]);
          sum += in[i];
        if(in[i] < 0) d.add edge(i, n + 1, -in
             [i]);
22
      d.add edge(t, s, numeric limits<T>::max
      if(d.flow(n, n + 1) < sum) return -1;</pre>
      return d.flow(s, t);
25
```

#### 3.3 Dinic

```
int n;
vector<Edge> edges;
vector<vi>g:
vi cur, h; // h : level graph
Dinic(int _n) : n(_n), g(_n) {}
void add_edge(int u, int v, T c) {
 g[u].pb(SZ(edges));
  edges.eb(u, v, c);
 g[v].pb(SZ(edges));
  edges.eb(v, u, 0);
bool bfs(int s, int t) {
 h.assign(n, -1);
  queue<int> q;
 h[s] = 0;
  q.push(s);
  while(!q.empty()) {
    int u = q.front(); q.pop();
    for(int i : g[u]) {
      const auto& e = edges[i];
      int v = e.to;
      if(e.cap > 0 && h[v] == -1) {
        h[v] = h[u] + 1;
        if(v == t) return true;
        q.push(v);
  return false;
T dfs(int u, int t, T f) {
 if(u == t) return f;
 Tr = f:
  for(int& i = cur[u]; i < SZ(g[u]); ++i)</pre>
    int j = g[u][i];
    const auto& e = edges[j];
    int v = e.to;
    T c = e.cap;
    if(c > 0 && h[v] == h[u] + 1) {
      T = dfs(v, t, min(r, c));
      edges[j].cap -= a;
      edges[j ^ 1].cap += a;
      if((r -= a) == 0) return f;
 return f - r;
T flow(int s, int t, T f = INF) {
 T ans = 0;
  while(f > 0 && bfs(s, t)) {
    cur.assign(n, 0);
    T cur = dfs(s, t, f);
    ans += cur:
    f -= cur;
  return ans;
```

#### 3.4 Flow 建模

· Maximum/Minimum flow with lower bound / Circulation problem

- 1. Construct super source S and sink T.
- 2. For each edge (x, y, l, u), connect  $x \to y$  with capacity u-l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v) > 0, connect  $S \rightarrow v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
  - To maximize, connect  $t \rightarrow s$  with capacity  $\infty$  (skip this in circulation problem). and let f be the maximum flow from Sto T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
  - To minimize, let f be the maximum flow from S to T. Connect  $t \rightarrow s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, f' is the
- 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$ corresponds to the flow of edge e on the graph.
- · Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $u \in Y$  is chosen iff u is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1) if c > 0, otherwise con- $\operatorname{nect} y \to x \text{ with } (cost, cap) = (-c, 1)$
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v) > 0, connect 11  $S \to v$  with (cost, cap) = (0, d(v))
  - 5. For each vertex v with d(v) < 0, connect 13  $v \to T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the 15 flow C + K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're check- 19 ing answer T
  - 2. Construct a max flow model, let K be the sum 21
  - 3. Connect source  $s \to v, v \in G$  with capacity K 23
  - 4. For each edge (u, v, w) in G, connect  $u \rightarrow v$  24 and  $v \to u$  with capacity w
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with ca-26 pacity  $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$  27
  - 6. T is a valid answer if the maximum flow f < 28K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect 32  $u' \to v'$  with weight w(u, v).
  - 2. Connect  $v \rightarrow v'$  with weight  $2\mu(v)$ , where 34  $\mu(v)$  is the cost of the cheapest edge incident to 35

```
3. Find the minimum weight perfect matching on 37
```

- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise, create edge (v, t) with capacity  $-p_v$ . 41
  - 2. Create edge (u, v) with capacity w with w being  $_{42}$ the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit  $_{44}$ of a subset of projects.

• 0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y} 3.6 \quad \text{general-weighted-max-matching}$$

16

17

21

22

23

24

26

28

29

35

37

38

REP(it, 100) {

return ans;

shuffle(ALL(o), rng);

vis.assign(n, false);

+= dfs(i);

for(auto i : o) if(mate[i] == -1) ans

can be minimized by the mincut of the following graph:

- 1. Create edge (x, t) with capacity  $c_x$  and create edge (s, y) with capacity  $c_y$ .
- 2. Create edge (x, y) with capacity  $c_{xy}$ .
- 3. Create edge (x, y) and edge (x', y') with capacity  $c_{xyx'y'}$ .

#### general-matching

```
struct GeneralMaxMatch {
   int n;
   vector<pii> es:
   vi g, vis, mate; // i 與 mate[i] 配對 (
        mate[i] == -1 代表沒有匹配)
   GeneralMaxMatch(int n) : n(n), g(n, -1),
        mate(n, -1) {}
   bool dfs(int u) {
     if(vis[u]) return false;
     vis[u] = true;
     for(int ei = g[u]; ei != -1;) {
       auto [x, y] = es[ei]; ei = y;
       if(mate[x] == -1) {
          mate[mate[u] = x] = u;
          return true;
     for(int ei = g[u]; ei != -1;) {
       auto [x, y] = es[ei]; ei = y;
       int nu = mate[x];
       mate[mate[u] = x] = u;
       mate[nu] = -1;
       if(dfs(nu)) return true;
       mate[mate[nu] = x] = nu;
       mate[u] = -1:
     return false;
   void add_edge(int a, int b) {
     auto f = [&](int a, int b) {
       es.eb(b, g[a]);
       g[a] = SZ(es) - 1;
     f(a, b); f(b, a);
   int solve() {
     vi o(n); iota(ALL(o), 0);
     int ans = 0;
```

```
1 // 1-based QQ
 struct WeightGraph {
   static const int inf = INT MAX;
    static const int maxn = 514;
    struct edge {
     int u, v, w;
      edge() {}
      edge(int u, int v, int w): u(u), v(v), w
          (w) {}
    int n, n_x;
   edge g[maxn * 2][maxn * 2];
   int lab[maxn * 2];
   int match[maxn * 2], slack[maxn * 2], st[
        maxn * 2], pa[maxn * 2];
   int flo from[maxn * 2][maxn + 1], S[maxn *
         2], vis[maxn * 2];
   vector<int> flo[maxn * 2];
   queue<int> q;
   int e delta(const edge &e) { return lab[e.
        u] + lab[e.v] - g[e.u][e.v].w * 2; }
    void update slack(int u, int x) { if(!
        slack[x] \mid\mid e_delta(g[u][x]) < e_delta
         (g[slack[x]][x])) slack[x] = u; }
    void set slack(int x) {
      slack[x] = 0;
      REP(u, n) if(g[u + 1][x].w > 0 \&\& st[u +
           1] != x && S[st[u + 1]] == 0)
          update_slack(u + 1, x);
    void q push(int x) {
     if(x \le n) q.push(x);
      else REP(i, SZ(flo[x])) q push(flo[x][i
    void set st(int x, int b) {
      st[x] = b;
      if(x > n) REP(i, SZ(flo[x])) set st(flo[
          x][i], b);
    int get pr(int b, int xr) {
     int pr = find(ALL(flo[b]), xr) - flo[b].
          begin();
     if(pr % 2 == 1) {
        reverse(1 + ALL(flo[b]));
        return SZ(flo[b]) - pr;
     return pr;
   void set match(int u, int v) {
      match[u] = g[u][v].v;
```

```
if(u <= n) return;</pre>
      edge e = g[u][v];
      int xr = flo from[u][e.u], pr = get pr(u
      for(int i = 0; i < pr; ++i) set_match(</pre>
           flo[u][i], flo[u][i ^ 1]);
      set match(xr, v);
      rotate(flo[u].begin(), flo[u].begin() +
           pr, flo[u].end());
    void augment(int u, int v) {
      while(true) {
                                                 102
        int xnv = st[match[u]];
                                                 103
        set match(u, v);
        if(!xnv) return;
        set_match(xnv, st[pa[xnv]]);
        u = st[pa[xnv]], v = xnv;
    int get_lca(int u, int v) {
      static int t = 0;
      for(++t; u || v; swap(u, v)) {
        if(u == 0) continue;
        if(vis[u] == t) return u;
        vis[u] = t;
        if(u = st[match[u]]) u = st[pa[u]];
                                                 116
65
      return 0;
                                                 117
    void add blossom(int u, int lca, int v) {
      int b = n + 1;
      while(b <= n_x && st[b]) ++b;</pre>
      if(b > n x) n x++;
      lab[b] = S[b] = 0;
      match[b] = match[lca];
      flo[b].clear(); flo[b].pb(lca);
      for(int x = u, y; x != lca; x = st[pa[y
           ]]) flo[b].pb(x), flo[b].pb(y = st[ 125]
           match[x]]), q_push(y);
      reverse(1 + ALL(flo[b]));
      for(int x = v, y; x != lca; x = st[pa[y 128]]
           ]]) flo[b].pb(x), flo[b].pb(y = st[ 129]
           match[x]]), q_push(y);
      set st(b, b):
      REP(x, n_x) g[b][x + 1].w = g[x + 1][b]. 131
      REP(x, n) flo from[b][x + 1] = 0;
      REP(i, SZ(flo[b])) {
                                                 133
        int xs = flo[b][i];
        REP(x, n_x) if(g[b][x + 1].w == 0 | |
             e_delta(g[xs][x + 1]) < e_delta(g[136])
             b[x + 1]) g[b][x + 1] = g[xs][x
             + 1], g[x + 1][b] = g[x + 1][xs];
        REP(x, n) if(flo from[xs][x + 1])
             flo from [b][x + 1] = xs;
                                                 138
      set slack(b);
    void expand_blossom(int b) {
      REP(i, SZ(flo[b])) set st(flo[b][i], flo 141
      int xr = flo from[b][g[b][pa[b]].u], pr
                                                 143
           = get pr(b, xr);
                                                 144
      for(int i = 0; i < pr; i += 2) {</pre>
        int xs = flo[b][i], xns = flo[b][i +
```

```
pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set slack(xns);
                                             148
    q_push(xns);
                                             149
                                             150
  S[xr] = 1, pa[xr] = pa[b];
                                             151
  for(size t i = pr + 1; i < SZ(flo[b]);
                                             152
      ++i) {
    int xs = flo[b][i];
                                             153
   S[xs] = -1, set_slack(xs);
                                             154
                                             155
 st[b] = 0;
                                             156
bool on found edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
                                             157
  if(S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
                                             158
    int nu = st[match[v]];
                                             159
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
                                             160
 } else if(S[v] == 0) {
                                             161
    int lca = get_lca(u, v);
                                             162
    if(!lca) return augment(u,v), augment(
                                             163
         v,u), true;
    else add blossom(u, lca, v);
                                             165
                                             166
  return false;
                                             167
bool matching() {
 memset(S + 1, -1, sizeof(int) * n x);
  memset(slack + 1, 0, sizeof(int) * n_x); 170
  q = queue<int>();
  REP(x, n x) if(st[x + 1] == x + 1 & !
                                             172
      match[x + 1]) pa[x + 1] = 0, S[x +
                                             173
      1] = 0, q push(x + 1);
                                             174
  if(q.empty()) return false;
                                             175
  while(true) {
    while(!q.empty()) {
                                             176
      int u = q.front(); q.pop();
                                             177
      if(S[st[u]] == 1) continue;
                                             178
      for(int v = 1; v <= n; ++v)</pre>
                                             179
        if(g[u][v].w > 0 && st[u] != st[v
                                             180
                                             181
          if(e delta(g[u][v]) == 0) {
                                             182
            if(on_found_edge(g[u][v]))
                 return true;
                                             183
          } else update slack(u, st[v]);
                                             184
    int d = inf;
    for(int b = n + 1; b \le n x; ++b) if(
         st[b] == b && S[b] == 1) d = min(d)
         , lab[b] / 2);
    for(int x = 1; x <= n x; ++x) {
      if(st[x] == x && slack[x]) {
        if(S[x] == -1) d = min(d, e_delta(
             g[slack[x]][x]));
        else if(S[x] == 0) d = min(d,
             e_delta(g[slack[x]][x]) / 2);
    REP(u, n) {
     if(S[st[u + 1]] == 0) {
        if(lab[u + 1] <= d) return 0;</pre>
        lab[u + 1] -= d:
```

```
} else if(S[st[u + 1]] == 1) lab[u + 9|
                 1] += d;
         for(int b = n + 1; b <= n_x; ++b)
           if(st[b] == b) {
             if(S[st[b]] == 0) lab[b] += d * 2; 11
             else if(S[st[b]] == 1) lab[b] -= d
                                                  12
         q = queue<int>();
         for(int x = 1; x \leftarrow n_x; ++x)
           if(st[x] == x && slack[x] && st[
                slack[x]] != x && e_delta(g[
                slack[x]][x]) == 0)
             if(on_found_edge(g[slack[x]][x]))
                  return true;
         for(int b = n + 1; b <= n_x; ++b)</pre>
           if(st[b] == b && S[b] == 1 && lab[b]
                 == 0) expand blossom(b):
       return false:
     pair<ll, int> solve() {
       memset(match + 1, 0, sizeof(int) * n);
       n x = n:
       int n_matches = 0;
       11 \text{ tot weight = 0};
       for(int u = 0; u \leftarrow n; ++u) st[u] = u,
            flo[u].clear();
       int w max = 0;
       for(int u = 1; u <= n; ++u)</pre>
         for(int v = 1; v <= n; ++v) {</pre>
           flo from[u][v] = (u == v ? u : 0);
           w_max = max(w_max, g[u][v].w);
       for(int u = 1; u <= n; ++u) lab[u] =</pre>
       while(matching()) ++n_matches;
       for(int u = 1; u <= n; ++u)
         if(match[u] && match[u] < u)</pre>
           tot weight += g[u][match[u]].w;
       return make_pair(tot_weight, n_matches);
     void add edge(int u, int v, int w) { g[u][
          v].w = g[v][u].w = w; }
     void init(int _n) : n(_n) {
       REP(u, n) REP(v, n) g[u + 1][v + 1] =
            edge(u + 1, v + 1, 0);
186 };
   3.7 KM
```

```
1 template < class T>
2 struct KM {
   static constexpr T INF = numeric limits<T</pre>
         >::max();
    int n, ql, qr;
    vector<vector<T>> w;
    vector<T> hl. hr. slk:
    vi fl, fr, pre, qu;
    vector<bool> v1, vr;
```

185

```
KM(int n) : n(n), w(n, vector < T > (n, -INF))
     , hl(n), hr(n), slk(n), fl(n), fr(n),
     pre(n), qu(n), vl(n), vr(n) {}
void add_edge(int u, int v, int x) { w[u][
     v] = x; } // 最小值要加負號
bool check(int x) {
  vl[x] = 1;
  if(fl[x] != -1) return vr[qu[qr++] = fl[
  while(x != -1) swap(x, fr[fl[x] = pre[x
  return 0;
void bfs(int s) {
  fill(ALL(slk), INF);
  fill(ALL(v1), 0), fill(ALL(vr), 0);
  ql = qr = 0, qu[qr++] = s, vr[s] = 1;
  while(true) {
   T d;
    while(ql < qr) {</pre>
      for(int x = 0, y = qu[ql++]; x < n;
        if(!vl[x] \&\& slk[x] >= (d = hl[x]
             + hr[y] - w[x][y])) {
          pre[x] = y;
          if(d) slk[x] = d;
          else if(!check(x)) return;
    d = INF:
    REP(x, n) if(!vl[x] \&\& d > slk[x]) d =
          slk[x];
    REP(x, n) {
      if(vl[x]) hl[x] += d;
      else slk[x] -= d;
      if(vr[x]) hr[x] -= d;
    REP(x, n) if(!v1[x] && !s1k[x] && !
         check(x)) return;
T solve() {
  fill(ALL(fl), -1);
  fill(ALL(fr), -1);
  fill(ALL(hr), 0);
  REP(i, n) hl[i] = *max_element(ALL(w[i])
      );
  REP(i, n) bfs(i);
  T ans = 0;
  REP(i, n) ans += w[i][fl[i]]; // i 跟 fl
      [i] 配對
  return ans;
```

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#### 3.8 max-clique

```
1 template<int V>
2 struct max clique {
   using B = bitset<V>;
   int n = 0;
```

```
vector<B> g, buf;
    struct P {
      int idx, col, deg;
      P(int a, int b, int b) : idx(a), col(b),
            deg(c) {}
    max_clique(int _n) : n(_n), g(_n), buf(_n)
    void add edge(int a, int b) {
      assert(a != b);
      g[a][b] = g[b][a] = 1;
    vector<int> now, clique;
    void dfs(vector<P>& rem){
      if(SZ(clique) < SZ(now)) clique = now;</pre>
      sort(ALL(rem), [](P a, P b) { return a.
           deg > b.deg; });
      int max_c = 1;
      for(auto& p : rem){
        p.col = 0;
        while((g[p.idx] & buf[p.col]).any()) p
             .col++;
        max_c = max(max_c, p.idx + 1);
        buf[p.col][p.idx] = 1;
      REP(i, max_c) buf[i].reset();
      sort(ALL(rem), [&](P a, P b) { return a.
           col < b.col; });</pre>
      for(;SZ(rem); rem.pop_back()){
        auto& p = rem.back();
        if(SZ(now) + p.col + 1 <= SZ(clique))</pre>
             break:
        vector<P> nrem;
        B bs;
        for(auto& q : rem){
          if(g[p.idx][q.idx]){
            nrem.eb(q.idx, -1, 0);
            bs[q.idx] = 1;
        for(auto& q : nrem) q.deg = (bs & g[q.
             idx]).count();
        now.eb(p.idx);
        dfs(nrem):
        now.pop_back();
    vector<int> solve(){
      vector<P> remark;
      REP(i, n) remark.eb(i, -1, SZ(g[i]));
      dfs(remark);
      return clique;
51 };
  3.9 MCMF
i template < class S, class T>
2 class MCMF {
public:
    struct Edge {
      int from, to;
      S cap;
```

```
Edge(int u, int v, S x, T y) : from(u),
       to(v), cap(x), cost(y) {}
                                             71
                                             72
const 11 INF = 1E18L;
                                             73
int n:
                                             74 };
vector<Edge> edges;
vector<vi>g:
vector<T> d;
vector<bool> ing;
vi pedge;
MCMF(int _n) : n(_n), g(_n), d(_n), inq(_n
    ), pedge(_n) {}
void add edge(int u, int v, S cap, T cost)
  g[u].pb(SZ(edges));
  edges.eb(u, v, cap, cost);
  g[v].pb(SZ(edges));
  edges.eb(v, u, 0, -cost);
bool spfa(int s, int t) {
 bool found = false;
  fill(ALL(d), INF);
  d[s] = 0;
  inq[s] = true;
  queue<int> q;
  q.push(s);
  while(!q.empty()) {
    int u = q.front(); q.pop();
    if(u == t) found = true;
    inq[u] = false;
    for(auto& id : g[u]) {
      const auto& e = edges[id];
      if(e.cap > 0 && d[u] + e.cost < d[e.</pre>
           to]) {
        d[e.to] = d[u] + e.cost;
        pedge[e.to] = id;
       if(!inq[e.to]) {
          q.push(e.to);
          inq[e.to] = true;
                                             24
  return found;
pair<S, T> flow(int s, int t, S f = INF) {
 S cap = 0;
 T cost = 0:
                                             29
  while(f > 0 && spfa(s, t)) {
   S send = f;
    int u = t;
    while(u != s) {
                                             33
      const Edge& e = edges[pedge[u]];
                                             34
      send = min(send, e.cap);
     u = e.from;
    u = t;
    while(u != s) {
     Edge& e = edges[pedge[u]];
      e.cap -= send;
     Edge& b = edges[pedge[u] ^ 1];
                                             40
     b.cap += send:
     u = e.from;
    cap += send;
```

```
3.10 minimum-general-weighted-
          perfect-matching
1 struct Graph {
  // Minimum General Weighted Matching (
         Perfect Match) 0-base
    static const int MXN = 105;
    int n, edge[MXN][MXN];
    int match[MXN], dis[MXN], onstk[MXN];
    vector<int> stk;
    void init(int _n) {
     n = n;
      for(int i=0; i<n; i++)</pre>
        for(int j=0; j<n; j++)</pre>
          edge[i][j] = 0;
    void add_edge(int u, int v, int w) { edge[
        u][v] = edge[v][u] = w; }
    bool SPFA(int u){
     if(onstk[u]) return true;
      stk.push_back(u);
      onstk[u] = 1;
      for(int v=0; v<n; v++){</pre>
        if(u != v && match[u] != v && !onstk[v
          int m = match[v];
          if(dis[m] > dis[u] - edge[v][m] +
               edge[u][v]){
            dis[m] = dis[u] - edge[v][m] +
                 edge[u][v];
            onstk[v] = 1;
            stk.push_back(v);
            if(SPFA(m)) return true;
            stk.pop_back();
            onstk[v] = 0;
      onstk[u] = 0;
      stk.pop back();
      return false:
    int solve() {
      for(int i = 0; i < n; i += 2) match[i] =</pre>
            i + 1, match[i+1] = i;
      while(true) {
        int found = 0:
        for(int i=0; i<n; i++) dis[i] = onstk[</pre>
             i] = 0;
        for(int i=0; i<n; i++){</pre>
          stk.clear();
          if(!onstk[i] && SPFA(i)){
            found = 1;
            while(stk.size()>=2){
```

f -= send;

cost += send \* d[t];

return {cap, cost};

#### 4 Geometry

#### 4.1 closest-pair

```
| const 11 INF = 9e18L + 5:
  vector<P> a;
  sort(all(a), [](P a, P b) { return a.x < b.x
4 11 SQ(11 x) { return x * x; }
5 11 solve(int 1, int r) {
    if(1 + 1 == r) return INF;
    int m = (1 + r) / 2;
    11 \text{ midx} = a[m].x:
    11 d = min(solve(1, m), solve(m, r));
    inplace_merge(a.begin() + 1, a.begin() + m
         , a.begin() + r, [](P a, P b) {
      return a.y < b.y;</pre>
12
13
    vector<P> p;
    for(int i = 1; i < r; ++i) if(SQ(a[i].x -</pre>
         midx) < d) p.pb(a[i]);
    REP(i, sz(p)) {
      for(int j = i + 1; j < sz(p); ++j) {
          d = \min(d, SQ(p[i].x - p[j].x) + SQ(
               p[i].y - p[j].y));
        if(SQ(p[i].y - p[j].y) > d) break;
18
19
20
    return d; // 距離平方
21
```

#### 4.2 convex-hull

pdd operator-(pdd a,pdd b){return {a.F-b.F,a

#### 4.3 half-plane

i typedef pair < double, double > pdd;

```
.S-b.S};}
pdd operator+(pdd a,pdd b){return {a.F+b.F,a
       .S+b.S};}
4 pdd operator*(pdd a,double x){return {a.F*x,
      a.S*x};}
double dot(pdd a,pdd b){return a.F*b.F+a.S*b
 double cross(pdd a,pdd b){return a.F*b.S-a.S
      *b.F;}
 struct bpmj{
   const double eps=1e-8;
   int n,m,id,1,r;
   pdd pt[55],q[1100];
   struct line{
     pdd x,y;
     double z;
     line(pdd _x,pdd _y):x(_x),y(_y){z=atan2(
          y.S,y.F);}
     line(){}
     bool operator<(const line &a)const{</pre>
           return z<a.z;}</pre>
    }a[550],dq[1005];
   pdd get (line x,line y){
     pdd v=x.x-v.x;
     double d=cross(y.y,v)/cross(x.y,y.y);
     return x.x+x.y*d;
   void solve(){
     dq[l=r=1]=a[1];
     for(int i=2;i<=id;++i){</pre>
        while(l<r&&cross(a[i].y,q[r-1]-a[i].x)</pre>
             <=eps) --r;
        while(l<r&&cross(a[i].y,q[l]-a[i].x)<=</pre>
             eps) ++1;
        dq[++r]=a[i];
        if(fabs(cross(dq[r].y,dq[r-1].y))<=eps</pre>
            ){
          if(cross(dq[r].y,a[i].x-dq[r].x)>eps
              ) dq[r]=a[i];
        if(l<r) q[r-1]=get_(dq[r-1],dq[r]);</pre>
     while(1<r&&cross(dq[1].y,q[r-1]-dq[1].x)</pre>
           <=eps) --r;
     if(r-1<=1) return;</pre>
     q[r]=get_(dq[1],dq[r]);
   void cal(){
```

```
double ans=0;
    a[r+1]=a[1];
    for(int i=1;i<=r;++i) ans+=cross(q[i],q[</pre>
    cout<<fixed<<setprecision(3)<<ans/2<<"\n</pre>
  void main_(){
    cin>>n:
    for(int x,y,i=0;i<n;++i){</pre>
      cin>>m:
      for(int i=0;i<m;++i) cin>>pt[i].F>>pt[
      pt[m]=pt[0];
      for(int i=0;i<m;++i) a[++id]=line(pt[i</pre>
           ],pt[i+1]-pt[i]);
    sort(a+1,a+1+id);
    solve();
    cal();
}valderyaya;
```

#### 4.4 min-enclosing-circle

```
pdd excenter(pdd x, pdd y, pdd z) {
  #define f(x, y) (x*x+y*y)
  auto [x1, y1] = x;
  auto [x2, y2] = y;
  auto [x3, y3] = z;
  double d1 = f(x2, y2) - f(x1, y1), d2 = f(
      x3, y3) - f(x2, y2);
  double fm = 2 * ((y3 - y2) * (x2 - x1) - (
      y2 - y1) * (x3 - x2));
  double ans_x = ((y3 - y2) * d1 - (y2 - y1)
      * d2) / fm;
  double ans y = ((x2 - x1) * d2 - (x3 - x2)
      * d1) / fm;
  #undef f
  return {ans x, ans y};
pdd min_enclosing_circle(vector<pdd> dots,
    double& r) {
  random shuffle(ALL(dots));
  pdd C = dots[0];
  r = 0;
  #define check(i, j) REP(i, j) if(abs(dots[
      il - C) > r
  check(i, SZ(dots)) {
   C = dots[i], r = 0;
    check(j, i) {
     C = (dots[i] + dots[j]) / 2.0;
     r = abs(dots[i] - C);
      check(k, j) {
        C = excenter(dots[i], dots[j], dots[
            k]);
        r = abs(dots[i] - C);
  #undef check
  return C;
```

#### 4.5 point-in-convex-hull

32 }

```
int point in convex hull(const vector<P>& a,
       P p) {
   // -1 ON, 0 OUT, +1 IN
   // 要先逆時針排序
   int n = SZ(a);
   if(btw(a[0], a[1], p) || btw(a[0], a[n -
        1], p)) return -1;
   int 1 = 0, r = n - 1;
   while(1 <= r) {
     int m = (1 + r) / 2;
     auto a1 = cross(a[m] - a[0], p - a[0]);
     auto a2 = cross(a[(m + 1) % n] - a[0], p
           - a[0]);
     if(a1 >= 0 && a2 <= 0) {
       auto res = cross(a[(m + 1) % n] - a[m
            ], p - a[m]);
       return res > 0 ? 1 : (res >= 0 ? -1 :
            0);
     if(a1 < 0) r = m - 1;
     else l = m + 1;
   return 0;
```

#### 4.6 point

```
using P = pair<11, 11>;
2 P operator+(P a, P b) { return P{a.X + b.X,
       a.Y + b.Y; }
3 P operator-(P a, P b) { return P{a.X - b.X,
       a.Y - b.Y}; }
4 P operator*(P a, 11 b) { return P{a.X * b, a
       .Y * b; }
5 P operator/(P a, 11 b) { return P{a.X / b, a
       .Y / b}; }
6 11 dot(P a, P b) { return a.X * b.X + a.Y *
       b.Y; }
 7 | 11 cross(P a, P b) { return a.X * b.Y - a.Y
       * b.X; }
8 11 abs2(P a) { return dot(a, a); }
9 double abs(P a) { return sqrt(abs2(a)); }
| \text{int sign}(11 \ x) \ \{ \text{ return } x < 0 \ ? -1 : (x == 0) \} 
        ? 0 : 1); }
int ori(P a, P b, P c) { return sign(cross(b)
        - a, c - a)); }
12 bool collinear(P a, P b, P c) { return sign(
       cross(a - c, b - c)) == 0; }
13 bool btw(P a, P b, P c) {
   if(!collinear(a, b, c)) return 0;
   return sign(dot(a - c, b - c)) <= 0;</pre>
16 }
17 bool seg intersect(Pa, Pb, Pc, Pd) {
   int a123 = ori(a, b, c);
  int a124 = ori(a, b, d);
```

```
int a341 = ori(c, d, a);
    int a342 = ori(c, d, b);
    if(a123 == 0 && a124 == 0) {
      return btw(a, b, c) || btw(a, b, d) ||
           btw(c, d, a) || btw(c, d, b);
25
    return a123 * a124 <= 0 && a341 * a342 <=
         0:
28 P intersect(P a, P b, P c, P d) {
    int a123 = cross(b - a, c - a);
    int a124 = cross(b - a, d - a);
    return (d * a123 - c * a124) / (a123 -
         a124);
33 struct line { P A, B; };
34 P vec(line L) { return L.B - L.A; }
35 P projection(P p, line L) { return L.A + vec
       (L) / abs(vec(L)) * dot(p - L.A, vec(L))
        / abs(vec(L)); }
```

#### 4.7 polar-angle-sort

#### 4.8 定理

- 皮克定理
  - 若一個多邊形的所有頂點都在整數點上‧則該 多邊形的面積  $S=a+\frac{b}{2}-1$ ‧其中 a 為內部 格點數目‧b 為邊上格點數目。

#### 5 Graph

#### 5.1 2-SAT

```
struct two_sat {
   int n; SCC g;
   vector<bool> ans;
   two_sat(int _n) : n(_n), g(_n * 2) {}
   void add_or(int u, bool x, int v, bool y)
   {
      g.add_edge(2 * u + !x, 2 * v + y);
      g.add_edge(2 * v + !y, 2 * u + x);
   }
   bool solve() {
```

```
ans.resize(n);
      auto id = g.solve();
      REP(i, n) {
        if(id[2 * i] == id[2 * i + 1]) return
        ans[i] = (id[2 * i] < id[2 * i + 1]);
      return true:
17
18 };
```

#### 5.2 BCC-tree

```
i struct BlockCutTree {
   int n;
   vector<vi> g;
   vi dfn, low, stk;
   int cnt = 0, cur = 0;
   vector<pii> edges;
   BlockCutTree(int _n) : n(_n), g(_n), dfn(
        _n), low(_n) {}
   void ae(int u, int v) {
     g[u].pb(v);
     g[v].pb(u);
   void dfs(int x) {
     stk.pb(x);
     dfn[x] = low[x] = cur++;
     for(auto y : g[x]) {
       if(dfn[y] == -1) {
         dfs(y);
         low[x] = min(low[x], low[y]);
         if(low[y] == dfn[x]) {
           int v:
           do {
             v = stk.back(), stk.pop back();
             edges.eb(n + cnt, v);
           } while (v != y);
           edges.eb(x, n + cnt);
           cnt++;
       } else low[x] = min(low[x], dfn[y]);
   pair<int, vector<pii>>> work() {
     REP(i, n) {
       if(dfn[i] == -1) {
         stk.clear();
         dfs(i);
     return {cnt, edges};
```

#### 5.3 centroid-tree

```
1 pair<int, vector<vi>>> centroid tree(const
      vector<vi>& g) {
   int n = sz(g);
```

```
vi siz(n);
vector<bool> vis(n);
auto dfs sz = [&](auto f, int u, int p) ->
  siz[u] = 1;
  for(auto v : g[u]) {
    if(v == p || vis[v]) continue;
    f(f, v, u);
    siz[u] += siz[v];
auto find_cd = [&](auto f, int u, int p,
    int all) -> int {
  for(auto v : g[u]) {
    if(v == p || vis[v]) continue;
    if(siz[v] * 2 > all) return f(f, v, u,
          all):
  return u:
vector<vi> h(n);
auto build = [&](auto f, int u) -> int {
  dfs_sz(dfs_sz, u, -1);
  int cd = find cd(find cd, u, -1, siz[u])
  vis[cd] = true;
  for(auto v : g[cd]) {
    if(vis[v]) continue;
    int child = f(f, v);
   h[cd].pb(child);
  return cd;
int root = build(build, 0);
return {root, h};
```

#### 5.4 chromatic-number

```
1 // vi to(n);
2 // to[u] |= 1 << v;
3 // to[v] |= 1 << u;
 int chromatic number(vi g) {
    constexpr int MOD = 998244353;
    int n = SZ(g):
    vector < int > I(1 << n); I[0] = 1;
    FOR(s, 1, 1 \ll n) {
      int v = __builtin_ctz(s), t = s ^ (1 <<</pre>
           v);
      I[s] = (I[t] + I[t \& \sim g[v]]) \% MOD;
    auto f = I;
    FOR(k, 1, n + 1) {
      int sum = 0;
      REP(s, 1 << n) {
        if((__builtin_popcount(s) ^ n) & 1)
             sum -= f[s];
        else sum += f[s];
        sum = ((sum \% MOD) + MOD) \% MOD;
        f[s] = 1LL * f[s] * I[s] % MOD;
      if(sum != 0) return k;
```

#### 5.5 HLD

24 }

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40

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return 48763;

```
1 struct HLD {
                                                 57
   int n;
   vector<vi> g;
    vi siz, par, depth, top, tour, fi, id;
    sparse_table<pii, min> st;
    HLD(int _n) : n(_n), g(_n), siz(_n), par(
         _n), depth(_n), top(_n), fi(_n), id(_n 61
        ) {
      tour.reserve(n);
    void add_edge(int u, int v) {
     g[u].push back(v);
     g[v].push_back(u);
    void build(int root = 0) {
      par[root] = -1;
      top[root] = root;
      vector<pii> euler tour;
      euler_tour.reserve(2 * n - 1);
                                                 71
      dfs sz(root);
                                                 72
      dfs_link(euler_tour, root);
      st = sparse table<pii, min>(euler tour);
    int get_lca(int u, int v) {
     int L = fi[u], R = fi[v];
      if(L > R) swap(L, R);
      return st.prod(L, R).second;
    bool is_anc(int u, int v) {
      return id[u] <= id[v] && id[v] < id[u] +</pre>
                                                81
           siz[u];
    bool on_path(int a, int b, int x) {
      return (is ancestor(x, a) || is ancestor
           (x, b)) && is_ancestor(get_lca(a, b)
           , x);
```

return depth[u] + depth[v] - 2 \* depth[(

while(depth[top[u]] > d) u = par[top[u

int kth\_node\_on\_path(int a, int b, int k)

if  $(k < 0 \mid | k > fi + se)$  return -1;

if(k < fi) return kth anc(a, k);</pre>

return kth anc(b, fi + se - k);

return tour[id[u] + d - depth[u]];

int get\_dist(int u, int v) {

get\_lca(u, v))];

if(depth[u] < k) return -1;</pre>

int kth\_anc(int u, int k) {

int d = depth[u] - k;

int z = get\_lca(a, b);

int fi = depth[a] - depth[z];

int se = depth[b] - depth[z];

#### 5.6 lowlink

14

```
1 struct lowlink {
    int n, cnt = 0, tecc_cnt = 0, tvcc_cnt =
    vector<vector<pii>>> g;
    vector<pii> edges:
    vi roots, id, low, tecc id, tvcc id;
    vector<bool> is_bridge, is_cut,
         is tree edge;
    lowlink(int _n) : n(_n), g(_n), is_cut(_n,
          false), id(_n, -1), low(_n, -1) {}
    void add_edge(int u, int v) {
      g[u].eb(v, SZ(edges));
      g[v].eb(u, SZ(edges));
11
      edges.eb(u, v);
12
      is_bridge.pb(false);
      is tree edge.pb(false);
13
      tvcc id.pb(-1);
15
```

vector<pii> get path(int u, int v, bool

if(u == v && !include lca) return {};

if(depth[top[u]] > depth[top[v]]) swap

if(depth[u] > depth[v]) swap(u, v); // u

if(u != v || include\_lca) seg.eb(id[u] +

if(par[u] != -1) g[u].erase(find(ALL(g[u

if(siz[v] > siz[g[u][0]]) swap(v, g[u

void dfs link(vector<pii>& euler tour, int

top[v] = (v == g[u][0] ? top[u] : v);

include lca = true) {

while(top[u] != top[v]) {

seg.eb(id[top[v]], id[v]);

!include lca, id[v]);

vector<pii> seg;

v = par[top[v]];

is Lca

void dfs sz(int u) {

]), par[u]));

depth[v] = depth[u] + 1;

for(auto& v : g[u]) {

siz[u] += siz[v];

][0]);

fi[u] = SZ(euler\_tour);

euler\_tour.eb(depth[u], u);

dfs\_link(euler\_tour, v);

euler\_tour.eb(depth[u], u);

id[u] = SZ(tour);

for(auto v : g[u]) {

tour.pb(u);

return seg;

siz[u] = 1;

par[v] = u;

dfs sz(v);

```
void dfs(int u, int peid = -1) {
      static vi stk;
      static int rid:
      if(peid < 0) rid = cnt;</pre>
      if(peid == -1) roots.pb(u);
      id[u] = low[u] = cnt++;
      for(auto [v, eid] : g[u]) {
        if(eid == peid) continue;
        if(id[v] < id[u]) stk.pb(eid);</pre>
        if(id[v] >= 0) low[u] = min(low[u], id
             [v]);
        else {
          is_tree_edge[eid] = true;
          dfs(v, eid);
          low[u] = min(low[u], low[v]);
          if((id[u] == rid && id[v] != rid +
               1) || (id[u] != rid && low[v] >= 93 |;
                id[u])) {
            is cut[u] = true;
          if(low[v] >= id[u]) {
            while(true) {
              int e = stk.back();
              stk.pop_back();
              tvcc id[e] = tvcc cnt;
              if(e == eid) break;
            tvcc_cnt++;
    void build() {
      REP(i, n) if(id[i] < 0) dfs(i);</pre>
      REP(i, SZ(edges)) {
        auto [u, v] = edges[i];
        if(id[u] > id[v]) swap(u, v);
        is_bridge[i] = (id[u] < low[v]);</pre>
52
    vector<vi>two_ecc() { // 邊雙
      tecc cnt = 0;
      tecc id.assign(n, -1);
      vi stk;
      REP(i, n) {
        if(tecc_id[i] != -1) continue;
        tecc_id[i] = tecc_cnt;
        stk.pb(i);
        while(SZ(stk)) {
          int u = stk.back(); stk.pop_back();
          for(auto [v, eid] : g[u]) {
            if(tecc_id[v] >= 0 || is_bridge[
                 eid]) continue;
            tecc id[v] = tecc cnt;
            stk.pb(v);
        tecc_cnt++;
      vector<vi> comp(tecc_cnt);
      REP(i, n) comp[tecc id[i]].pb(i);
72
      return comp;
74
    vector<vi> bcc_vertices() { // 點雙
      vector<vi> comp(tvcc_cnt);
```

```
REP(i, SZ(edges)) {
    comp[tvcc_id[i]].pb(edges[i].first);
    comp[tvcc_id[i]].pb(edges[i].second);
  for(auto& v : comp) {
    sort(ALL(v));
   v.erase(unique(ALL(v)), v.end());
 REP(i, n) if(!SZ(g[i])) comp.pb({i});
  return comp;
vector<vi> bcc_edges() {
 vector<vi> ret(tvcc_cnt);
 REP(i, SZ(edges)) ret[tvcc_id[i]].pb(i);
  return ret;
```

#### manhattan-mst

template < class T> // [w, u, v]

```
vector<tuple<T, int, int>> manhattan mst(
    vector<T> xs, vector<T> ys) {
  const int n = SZ(xs);
  vi idx(n); iota(ALL(idx), 0);
  vector<tuple<T, int, int>> ret;
  REP(s, 2) {
   REP(t, 2) {
      auto cmp = [&](int i, int j) { return
           xs[i] + ys[i] < xs[j] + ys[j]; };
      sort(ALL(idx), cmp);
      map<T, int> sweep;
      for(int i : idx) {
        for(auto it = sweep.lower bound(-ys[
            i]); it != sweep.end(); it =
             sweep.erase(it)) {
          int j = it->second;
          if(xs[i] - xs[j] < ys[i] - ys[j])</pre>
          ret.eb(abs(xs[i] - xs[j]) + abs(ys
               [i] - ys[j]), i, j);
        sweep[-ys[i]] = i;
     swap(xs, ys);
    for(auto\& x : xs) x = -x;
  sort(ALL(ret));
 return ret:
```

#### 5.8 SCC

22

```
1 struct SCC {
   int n;
   vector<vi> g, h;
   SCC(int _n) : n(_n), g(_n), h(_n) {}
   void add edge(int u, int v) {
     g[u].pb(v);
```

```
h[v].pb(u);
8
    vi solve() { // 回傳縮點的編號
      vi id(n), top;
      top.reserve(n);
      function<void(int)> dfs1 = [&](int u) {
        id[u] = 1;
        for(auto v : g[u]) if(id[v] == 0) dfs1
        top.pb(u);
      REP(v, n) if(id[v] == 0) dfs1(v);
      fill(ALL(id), -1);
      function<void(int, int)> dfs2 = [&](int
           u, int x) {
        id[u] = x;
        for(auto v : h[u]) if(id[v] == -1)
             dfs2(v, x);
      for(int i = n - 1, cnt = 0; i >= 0; --i)
        int u = top[i];
        if(id[u] == -1) dfs2(u, cnt++);
      return id;
29 };
```

#### triangle-sum

17

20

21

22

23

24

26

27

28

```
1 // Three vertices a < b < cconnected by
       three edges \{a, b\}, \{a, c\}, \{b, c\}. Find
        xa * xb * xc over all triangles.
int triangle sum(vector<array<int, 2>> edges
       , vi x) {
    int n = SZ(x);
    vi deg(n);
    vector<vi> g(n);
    for(auto& [u, v] : edges) {
      if(u > v) swap(u, v);
      deg[u]++, deg[v]++;
    REP(i, n) g[i].reserve(deg[i]);
    for(auto [u, v] : edges) {
      if(deg[u] > deg[v]) swap(u, v);
      g[u].pb(v);
14
    vi val(n);
     int128 ans = 0;
    REP(a, n) {
      for(auto b : g[a]) val[b] = x[b];
      for(auto b : g[a]) {
        11 \text{ tmp} = 0;
        for(auto c : g[b]) tmp += val[c];
        ans += int128(tmp) * x[a] * x[b];
24
      for(auto b : g[a]) val[b] = 0;
25
    return ans % mod;
27 }
```

#### Math

#### 6.1 Aliens

```
1 template < class Func, bool MAX>
2 11 Aliens(11 1, 11 r, int k, Func f) {
   while(1 < r) {</pre>
     11 m = 1 + (r - 1) / 2;
      auto [score, op] = f(m);
     if(op == k) return score + m * k * (MAX
          ? +1 : -1);
     if(op < k) r = m;
     else l = m + 1;
   return f(1).first + 1 * k * (MAX ? +1 :
         -1);
```

#### 6.2 Berlekamp-Massey

```
1 / / - [1, 2, 4, 8, 16] \rightarrow (1, [1, -2])
2 // - [1, 1, 2, 3, 5, 8] -> (2, [1, -1, -1])
3 // - [0, 0, 0, 0, 1] -> (5, [1, 0, 0, 0, 0,
       998244352]) (mod 998244353)
 4 // - [] -> (0, [1])
 5 // - [0, 0, 0] -> (0, [1])
 6 // - [-2] -> (1, [1, 2])
7 template < class T>
  pair<int, vector<T>> BM(const vector<T>& S)
    using poly = vector<T>;
    int N = SZ(S);
    poly C_rev{1}, B{1};
12
    int L = 0, m = 1;
13
    T b = 1;
     auto adjust = [](poly C, const poly &B, T
          d, T b, int m) -> poly {
      C.resize(max(SZ(C), SZ(B) + m));
16
      Ta = d / b;
17
      REP(i, SZ(B)) C[i + m] -= a * B[i];
      return C;
18
19
20
     REP(n, N) {
21
      T d = S[n];
      REP(i, L) d += C rev[i + 1] * S[n - 1 -
           i];
      if(d == 0) m++;
       else if (2 * L <= n) {
25
        poly Q = C_rev;
26
         C_rev = adjust(C_rev, B, d, b, m);
        L = n + 1 - L, B = Q, b = d, m = 1;
27
      } else C_rev = adjust(C_rev, B, d, b, m
28
           ++);
29
    return {L, C_rev};
31 }
33 // Calculate x^N \b f(x)
34 // Complexity: \$0(K^2 \setminus \log N)\$ (\$K\$: deg. of
```

 $|35|//(4, [1, -1, -1]) \rightarrow [2, 3]$ 

```
2)
37 template < class T>
38 vector<T> monomial_mod_polynomial(long long
       N, const vector(T> &f rev) {
    assert(!f_rev.empty() && f_rev[0] == 1);
    int K = SZ(f rev) - 1;
    if(!K) return {};
    int D = 64 - __builtin_clzll(N);
    vector<T> ret(K, 0);
    ret[0] = 1;
    auto self_conv = [](vector<T> x) -> vector
         <T> {
      int d = SZ(x);
      vector<T> ret(d * 2 - 1);
      REP(i, d) {
        ret[i * 2] += x[i] * x[i];
        REP(j, i) ret[i + j] += x[i] * x[j] *
      return ret;
    for(int d = D: d--:) {
      ret = self conv(ret);
      for(int i = 2 * K - 2; i >= K; i--) {
        REP(j, k) ret[i - j - 1] -= ret[i] *
             f rev[j + 1];
      ret.resize(K);
      if (N >> d & 1) {
        vector<T> c(K);
        c[0] = -ret[K - 1] * f rev[K];
        for(int i = 1; i < K; i++) c[i] = ret[</pre>
             i - 1] - ret[K - 1] * f rev[K - i
        ret = c;
    return ret;
  // Guess k-th element of the sequence,
       assumina linear recurrence
71 template < class T>
72 T guess kth term(const vector<T>& a, long
       long k) {
    assert(k >= 0);
    if(k < 1LL * SZ(a)) return a[k];</pre>
    auto f = BM<T>(a).second;
    auto g = monomial mod polynomial<T>(k, f);
    T ret = 0:
    REP(i, SZ(g)) ret += g[i] * a[i];
    return ret;
```

#### 6.3 Chinese-Remainder

```
if(m0 < m1) swap(r0, r1), swap(m0, m1);
if(m0 % m1 == 0) {
    if(r0 % m1 != r1) return {0, 0};
}
ll g, im, qq;
g = ext_gcd(m0, m1, im, qq);
ll u1 = (m1 / g);
if((r1 - r0) % g) return {0, 0};
ll x = (r1 - r0) / g % u1 * im % u1;
r0 += x * m0;
m0 *= u1;
if(r0 < 0) r0 += m0;
return {r0, m0};
}</pre>
```

#### **6.4** Combination

#### 6.5 Determinant

```
1 T det(vector<vector<T>> a) {
   int n = SZ(a);
   T ret = 1;
   REP(i, n) {
     int idx = -1;
     FOR(j, i, n) if(a[j][i] != 0) {
       idx = j;
       break;
     if(idx == -1) return 0:
     if(i != idx) {
       ret *= T(-1);
       swap(a[i], a[idx]);
     ret *= a[i][i];
     T inv = T(1) / a[i][i];
     REP(j, n) a[i][j] *= inv;
     FOR(j, i + 1, n) {
       T x = a[j][i];
       if(x == 0) continue;
       FOR(k, i, n) {
         a[j][k] -= a[i][k] * x;
   return ret:
```

#### 6.6 Discrete-Log

#### 6.7 extgcd

```
1  // ax + by = gcd(a, b)
2  ll ext_gcd(ll a, ll b, ll& x, ll& y) {
3    if(b == 0) {
4        x = 1, y = 0;
5        return a;
6    }
7   ll x1, y1;
8   ll g = ext_gcd(b, a % b, x1, y1);
9   x = y1, y = x1 - (a / b) * y1;
10   return g;
11 }
```

#### 6.8 Floor-Sum

```
1 // sum \{i = 0\}^{n - 1} floor((ai + b) / c)
       in O(a + b + c + n)
2 11 floor sum(11 n, 11 a, 11 b, 11 c) {
    assert(0 <= n && n < (1LL << 32));
    assert(1 <= c && c < (1LL << 32));
    ull ans = 0:
    if(a < 0) {
      ull a2 = (a \% c + c) \% c;
      ans -= 1ULL * n * (n - 1) / 2 * ((a2 - a
           ) / c);
      a = a2:
    if(b < 0) {
      ull b2 = (b \% c + c) \% c;
      ans -= 1ULL * n * ((b2 - b) / c);
      b = b2;
    ull N = n, C = c, A = a, B = b;
    while(true) {
      if(A >= C) {
        ans += N * (N - 1) / 2 * (A / C);
        A %= C:
21
      if(B >= C) {
```

1 #define ppc builtin popcount

template < class T, class F>
void fwht(vector<T>& a, F f) {

#### **6.9 FWHT**

int n = SZ(a);

```
assert(ppc(n) == 1);
    for(int i = 1; i < n; i <<= 1) {</pre>
      for(int j = 0; j < n; j += i << 1) {
        REP(k, i) f(a[j + k], a[i + j + k]);
11
12 template < class T>
void or_transform(vector<T>& a, bool inv) {
       fwht(a, [\&](T\& x, T\& y) \{ y += x * (inv) \}
       ? -1 : +1); }) }
14 template < class T>
void and transform(vector<T>& a, bool inv) {
        fwht(a, [\&](T\& x, T\& y) \{ x += y * (inv) \}
        ? -1 : +1); }); }
16 template < class T>
17 void xor_transform(vector<T>& a, bool inv) {
    fwht(a, [](T& x, T& y) {
      Tz = x + y;
      y = x - y;
21
      x = z;
22
    if(inv) {
23
      Tz = T(1) / T(SZ(a));
24
25
      for(auto& x : a) x *= z;
26
27 }
28 template < class T>
29 vector<T> convolution(vector<T> a, vector<T>
    assert(SZ(a) == SZ(b));
    transform(a, false), transform(b, false);
    REP(i, SZ(a)) a[i] *= b[i];
33
    transform(a, true);
34
    return a:
35 }
36 template < class T>
  vector<T> subset convolution(const vector<T</pre>
       >& f, const vector<T>& g) {
    assert(SZ(f) == SZ(g));
    int n = SZ(f):
39
    assert(ppc(n) == 1);
40
41
    const int lg = __lg(n);
    vector<vector<T>> fhat(lg + 1, vector<T>(n
         )), ghat(fhat);
```

#### 6.10 Gauss-Jordan

```
i int GaussJordan(vector<vector<ld>>& a) {
   // -1 no sol, 0 inf sol
   int n = SZ(a);
   REP(i, n) assert(SZ(a[i]) == n + 1);
   REP(i, n) {
     int p = i;
     REP(j, n) {
       if(j < i && abs(a[j][j]) > EPS)
            continue;
       if(abs(a[j][i]) > abs(a[p][i])) p = j;
     REP(j, n + 1) swap(a[i][j], a[p][j]);
     if(abs(a[i][i]) <= EPS) continue;</pre>
     REP(j, n) {
       if(i == j) continue;
       ld delta = a[j][i] / a[i][i];
       FOR(k, i, n + 1) a[j][k] -= delta * a[
   bool ok = true;
   REP(i, n) {
     if(abs(a[i][i]) <= EPS) {</pre>
       if(abs(a[i][n]) > EPS) return -1;
       ok = false;
   return ok;
```

#### 6.11 GCD-Convolution

```
6.12 Int-Div
```

for(int p = 5, d = 4, i = 1; p <= N; p +=

static void zeta transform(vector<T>& a) {

for(auto p : prime enumerate(n)) {

for(int i = 1; i \* p <= n; i++) {

static void mobius transform(vector<T>& a)

for(auto p : prime enumerate(n)) {

for(int i = n / p; i > 0; i--) {

static void zeta\_transform(vector<T>& a) {

for(auto p : prime\_enumerate(n)) {

for(int i = n / p; i > 0; i--) {

static void mobius transform(vector<T>& a)

for(auto p : prime enumerate(n)) {

for(int i = 1; i \* p <= n; i++) {</pre>

vector<T> gcd convolution(const vector<T>& a

multiple\_transform::zeta\_transform(f);

multiple transform::zeta transform(g):

multiple transform::mobius transform(f);

while(SZ(ret) && ret.back() > N) ret.

d = 6 - d, i++) {

if(sieve[i]) {

ret.pb(p);

pop back();

struct divisor\_transform {
 template < class T>

template < class T>

int n = a.size() - 1;

int n = a.size() - 1;

struct multiple transform {

int n = a.size() - 1;

int n = a.size() - 1;

// lcm: multiple -> divisor

template < class T>

return f:

auto f = a, g = b;

a[i] -= a[i \* p];

, const vector<T>& b) {

assert(a.size() == b.size());

REP(i, SZ(f)) f[i] \*= g[i];

a[i] += a[i \* p];

template < class T>

template < class T>

a[i \* p] -= a[i];

a[i \* p] += a[i];

return ret;

#### 6.13 Linear-Sieve

```
vi primes, least = {0, 1}, phi, mobius;
void LinearSieve(int n) {
    least = phi = mobius = vi(n + 1):
    mobius[1] = 1:
    for(int i = 2; i <= n; i++) {</pre>
      if(!least[i]) {
        least[i] = i;
        primes.pb(i);
        phi[i] = i - 1;
        mobius[i] = -1;
      for(auto j : primes) {
        if(i * j > n) break;
        least[i * j] = j;
        if(i % j == 0) {
          mobius[i * j] = 0;
          phi[i * j] = phi[i] * j;
          break;
        } else {
          mobius[i * j] = -mobius[i];
          phi[i * j] = phi[i] * phi[j];
23
24
25 }
```

#### 6.14 Miller-Rabin

```
| bool is_prime(ll n, vector<ll> x) {
    11 d = n - 1:
    d >>= __builtin_ctzll(d);
    for(auto a : x) {
      if(n <= a) break;</pre>
      11 t = d, v = 1, b = t;
      while(b) {
        if(b & 1) y = i128(y) * a % n;
        a = i128(a) * a % n;
        b >>= 1:
      while(t != n - 1 && y != 1 && y != n -
        y = i128(y) * y % n;
       t <<= 1;
      if(y != n - 1 && t % 2 == 0) return 0;
17
   return 1;
```

```
20 bool is_prime(11 n) {
21    if(n <= 1) return 0;
22    if(n % 2 == 0) return n == 2;
23    if(n < (1LL << 30)) return is_prime(n, {2, 7, 61});
24    return is_prime(n, {2, 325, 9375, 28178, 450775, 9780504, 1795265022});
25 }</pre>
```

#### 6.15 Min-of-Mod-of-Linear

#### **6.16** Mod-Inv

```
int inv(int a) {
   if(a < N) return inv[a];
   if(a == 1) 1;
   return (MOD - 1LL * (MOD / a) * inv(MOD %
        a) % MOD) % MOD;
}

vi mod_inverse(int m, int n = -1) {
   assert(n < m);
   if(n == -1) n = m - 1;
   vi inv(n + 1);
   inv[0] = inv[1] = 1;
   for(int i = 2; i <= n; i++) inv[i] = m - 1
   LL * (m / i) * inv[m % i] % m;
   return inv;
}</pre>
```

#### 6.17 Mod-Sqrt

```
1  // return -1 if sqrt DNE
2  ll mod_sqrt(ll a, ll mod) {
3    a %= mod;
4    if(mod == 2 || a < 2) return a;
5    if(mod_pow(a, (mod-1)/2, mod) != 1) return
-1;
6    ll b = 1;
7    while(mod_pow(b, (mod-1)/2, mod) == 1) b
++;
8    int m = mod-1, e = __builtin_ctz(m);
9    m >>= e;
```

```
11 \times = \text{mod pow}(a, (m-1)/2, \text{mod});
11 \ y = a * x \% \ mod * x \% \ mod;
x = x * a % mod:
11 z = mod pow(b, m, mod);
while(y != 1) {
  int i = 0:
  11 t = v;
  while(t != 1) t = t * t % mod, i++:
  z = mod pow(z, 1LL << (e - j - 1), mod);
  x = x*z\%mod, z = z*z\%mod, y = y*z\%mod;
  e = j;
return min(x, mod-x); // neg is $mod-x$
```

#### 6.18 NTT

```
// = 998244353
_{2} // For p < 2^30 there is also e.g. 5 << 25,
      7 << 26, 479 << 21
3 // and 483 << 21 (same root). The last two
      are > 10^9.
  typedef vector<ll> v1:
  void ntt(vl &a) {
   int n = SZ(a), L = 31 - builtin clz(n);
    static vl rt(2, 1);
    for(static int k = 2, s = 2; k < n; k *=
        2, s++) {
      rt.resize(n);
     |1| z[] = \{1, mod pow(root, mod >> s, mod\}
      FOR(i, k, 2 * k) rt[i] = rt[i / 2] * z[i 3]
           & 1] % mod;
    vi rev(n);
    REP(i, n) rev[i] = (rev[i / 2] | (i \& 1)
         << L) / 2;
    REP(i, n) if (i < rev[i]) swap(a[i], a[rev</pre>
         [i]]);
    for(int k = 1; k < n; k *= 2)
     for(int i = 0; i < n; i += 2 * k) REP(j,
        11 z = rt[j + k] * a[i + j + k] % mod,
             &ai = a[i + j];
        a[i + j + k] = ai - z + (z > ai ? mod
            : 0);
        ai += (ai + z >= mod ? z - mod : z):
23 vl conv(const vl &a, const vl &b) {
    if(a.empty() || b.empty()) return {};
    int s = SZ(a) + SZ(b) - 1, B = 32 -
         __builtin_clz(s), n = 1 << B;
    11 inv = mod_pow(n, mod - 2, mod);
    vl L(a), R(b), out(n);
    L.resize(n), R.resize(n);
    ntt(L), ntt(R);
    REP(i, n) out[-i & (n - 1)] = inv * L[i] %
          mod * R[i] % mod;
    ntt(out):
    return {out.begin(), out.begin() + s};
```

#### 6.19 Pollard-Rho

```
void PollardRho(map<11, int>& mp, 11 n) {
   if(n == 1) return;
   if(is_prime(n)) return mp[n]++, void();
   if(n \% 2 == 0) {
     mp[2] += 1;
     PollardRho(mp, n / 2);
     return;
   11 x = 2, y = 2, d = 1, p = 1;
   #define f(x, n, p) ((i128(x) * x % n + p)
        % n)
   while(1) {
     if(d != 1 && d != n) {
       PollardRho(mp, d);
       PollardRho(mp, n / d);
       return:
     p += (d == n);
     x = f(x, n, p), y = f(f(y, n, p), n, p);
     d = \_gcd(abs(x - y), n);
   #undef f
 vector<ll> get divisors(ll n) {
   if(n == 0) return {};
   map<11. int> mp:
   PollardRho(mp, n);
   vector<pair<ll, int>> v(ALL(mp));
   vector<ll> res:
   auto f = [&](auto f, int i, ll x) -> void
     if(i == SZ(v)) {
       res.pb(x);
       return:
     for(int j = v[i].second; ; j--) {
       f(f, i + 1, x);
       if(j == 0) break;
       x *= v[i].first;
   f(f, 0, 1);
   sort(ALL(res));
   return res;
```

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#### **6.20** Poly

```
template<int mod>
struct Poly {
  vector<ll> a;
  Poly() {}
  Poly(int n) : a(n) {}
  Poly(const vector<ll>& _a) : a(_a) {}
  Poly(const initializer list<ll>& a) : a(
  int size() const { return SZ(a): }
  void resize(int n) { a.resize(n); }
  void shrink() {
```

```
while(size() && a.back() == 0) a.ppb(); 64
12
   11 at(int idx) const {
      return idx >= 0 \&\& idx < size() ? a[idx] 66
    11& operator[](int idx) { return a[idx]; }
    friend Poly operator+(const Poly& a. const
          Polv& b) {
      Poly c(max(SZ(a), SZ(b)));
      REP(i, SZ(c)) c[i] = (a.at(i) + b.at(i))
            % mod;
      return c;
                                                  74
    friend Poly operator-(const Poly& a, const 76
          Polv& b) {
      Poly c(max(SZ(a), SZ(b)));
      REP(i, SZ(c)) c[i] = (a.at(i) - b.at(i)
           + mod) % mod;
      return c;
                                                  80
    friend Poly operator*(Poly a, Poly b) {
      return Poly(conv(a.a, b.a)); // see NTT. 83
    friend Poly operator*(ll a, Poly b) {
                                                  85
      REP(i, SZ(b)) (b[i] *= a) %= mod;
                                                  86
      return b:
                                                  87
                                                  88
    friend Poly operator*(Poly a, 11 b) {
                                                  89
      REP(i, SZ(a)) (a[i] *= b) %= mod;
      return a;
    Poly& operator+=(Poly b) { return (*this)
                                                  93
         = (*this) + b; }
    Poly& operator -= (Poly b) { return (*this)
                                                  95
         = (*this) - b; }
    Poly& operator*=(Poly b) { return (*this)
         = (*this) * b; }
    Poly& operator*=(ll b) { return (*this) =
         (*this) * b; }
    #define MSZ if(m == -1) m = size();
    Poly mulxk(int k) const {
      auto b = a:
      b.insert(b.begin(), k, 0);
                                                 100
      return Poly(b);
                                                 101
                                                 102
    Poly modxk(int k) const {
                                                 103
      k = min(k, size());
                                                 104
      return Poly(vector<ll>(a.begin(), a.
                                                 105
           begin() + k);
                                                 106
                                                 107
    Poly divxk(int k) const {
      if(size() <= k) return Poly();</pre>
      return Poly(vector<ll>(a.begin() + k, a.
                                                 109
           end()));
                                                 110
                                                 111
    Poly deriv() const {
                                                 112
      if(!SZ(a)) return Poly();
                                                 113
      Polv c(size() - 1):
                                                 114
      REP(i, size() - 1) c[i] = (i + 1LL) * a[115]
           i + 1 % mod;
      return c;
                                                 117
                                                 118
    Polv integr() const {
      Poly c(size() + 1);
```

```
REP(i, size()) c[i + 1] = a[i] * mod pow
       (i+1, mod-2, mod) % mod;
  return c:
Poly inv(int m = -1) const { MSZ;
  Poly x{mod pow(a[0], mod-2, mod)};
  int k = 1;
  while(k < m) {</pre>
   k *= 2:
   x = (x * (Poly{2} - modxk(k) * x)).
         modxk(k);
 return x.modxk(m);
Poly log(int m = -1) const { MSZ;
  return (deriv() * inv(m)).integr().modxk
Poly exp(int m = -1) const { MSZ;
 Poly x{1};
 int k = 1;
  while(k < m) {</pre>
   k *= 2;
   x = (x * (Poly{1} - x.log(k) + modxk(k))
        ))).modxk(k);
 return x.modxk(m);
Poly pow(ll k, int m = -1) const { MSZ;
 if(k == 0) {
   Poly b(m); b[0] = 1;
    return b;
 int s = 0, sz = size();
  while(s < sz && a[s] == 0) s++;</pre>
  if(s == sz) return *this;
 if(m > 0 \&\& s > = (sz + k - 1) / k)
       return Poly(m);
  if(s * k >= m) return Poly(m);
  return (((divxk(s) * mod_pow(a[s], mod
       -2, mod).log(m) * (k % <math>mod).exp(m)
        * mod_pow(a[s], k, mod)).mulxk(s *
       k).modxk(m);
bool has_sqrt() const {
 if(size() == 0) return true;
 int x = 0;
  while(x < size() && a[x] == 0) x++;
 if(x == size()) return true;
 if(x % 2 == 1) return false;
  11 v = a[x];
  return (y == 0 \mid \mid mod pow(y, (mod-1)/2,
       mod) == 1);
Poly sqrt(int m = -1) const { MSZ;
 if(size() == 0) return Poly();
  int x = 0:
  while(x < size() && a[x] == 0) x++;
  if(x == size()) return Poly(size());
  Polv f = divxk(x):
  Poly g({mod sqrt(f[0], mod)});
  11 \text{ inv2} = \text{mod pow}(2, \text{mod-2}, \text{mod});
  for(int i = 1: i < m: i *= 2) {</pre>
   g = (g + f.modxk(i * 2) * g.inv(i * 2)
         ) * inv2:
```

```
return g.modxk(m).mulxk(x / 2);
                                                  176 };
121
     Polv shift(ll c) const {
       int n = size();
       Poly b(*this);
124
       11 f = 1:
       REP(i, n) {
         (b[i] *= f) \%= mod;
         (f *= i + 1) \% = mod;
129
130
       reverse(ALL(b.a));
       Poly exp cx(vector<ll>(n, 1));
       FOR(i, 1, n) exp_cx[i] = exp_cx[i - 1] *
            c % mod * mod_pow(i, mod-2, mod) %
       b = (b * exp_cx).modxk(n);
       reverse(ALL(b.a));
       (f *= mod_pow(n, mod-2, mod)) %= mod;
       11 z = mod pow(f, mod-2, mod);
137
       IREP(i, n) {
138
         (b[i] *= z) %= mod;
         (z *= i) %= mod;
139
140
141
       return b;
142
     Poly mulT(Poly b) const {
       int n = SZ(b);
       if(!n) return Poly();
       reverse(ALL(b.a));
       return ((*this) * b).divxk(n - 1);
     vector<ll> eval(vector<ll> x) const {
       if(size() == 0) return vector<ll>(SZ(x),
       const int n = max(SZ(x), size());
       vector<Poly> q(4 * n);
       vector<ll> ans(SZ(x));
154
       x.resize(n);
       function < void(int, int, int) > build =
            [&](int p, int l, int r) {
         if(r - 1 == 1) q[p] = Poly{1, mod - x[}
              1]};
         else {
           int m = (1 + r) / 2;
           build(2 * p, 1, m), build(2 * p + 1,
                 m, r);
           q[p] = q[2 * p] * q[2 * p + 1];
       };
       build(1, 0, n);
       function<void(int, int, int, const Poly</pre>
            &)> work = [&](int p, int 1, int r,
            const Poly& num) {
         if(r - 1 == 1) {
           if(1 < SZ(ans)) ans[1] = num.at(0);
         } else {
           int m = (1 + r) / 2;
           work(2 * p, 1, m, num.mulT(q[2 * p +
                 1]).modxk(m - 1));
           work(2 * p + 1, m, r, num.mulT(q[2 *
                 p]).modxk(r - m));
171
172
       };
173
       work(1, 0, n, mulT(q[1].inv(n)));
       return ans:
```

#### 6.21 Primes

```
1 /* 12721 13331 14341 75577 123457 222557
      556679 999983 1097774749 1076767633
      100102021 999997771 1001010013
      1000512343 987654361 999991231 999888733
       98789101 987777733 999991921 1010101333
       1010102101 1000000000039
      1000000000000037 2305843009213693951
      4611686018427387847 9223372036854775783
      18446744073709551557 */
```

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#### 6.22 Simplex

```
* Description: Solves a general linear
     maximization problem: maximize $c^T x$
     subject to $Ax \le b$, $x \qe 0$.
* Returns -inf if there is no solution, inf
      if there are arbitrarily good
     solutions, or the maximum value of $c^T
      x$ otherwise.
* The input vector is set to an optimal $x$
      (or in the unbounded case, an
     arbitrary solution fulfilling the
     constraints).
 * Numerical stability is not quaranteed.
     For better performance, define
     variables such that $x = 0$ is viable.
 * vvd^{T}A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
 * vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
 * T val = LPSolver(A, b, c).solve(x);
 * Time: O(NM * \#pivots), where a pivot may
      be e.g. an edge relaxation. O(2^n) in
     the general case.
* 將最小化改成最大化 -> 去除等式 -> 去除大
     於等於 -> 去除自由變數,將 x1 用 x1-x3
typedef double T: // Long double, Rational,
    double + mod<P>...
typedef vector<T> vd:
typedef vector<vd> vvd;
struct LP {
 const T eps = 1e-8, inf = 1/.0;
  #define MP make pair
 #define ltj(X) if(s == -1 || MP(X[j],N[j])
       < MP(X[s],N[s])) s=j
 int m, n;
 vi N, B;
 vvd D;
 LP(const vvd& A, const vd& b, const vd& c)
      (m+2, vd(n+2)) {
```

```
REP(i, m) REP(j, n) D[i][j] = A[i][j];
      i|[n+1] = b[i];
      REP(j, n) \{ N[j] = j; D[m][j] = -c[j]; \}
29
      N[n] = -1; D[m+1][n] = 1;
    void pivot(int r, int s) {
      T *a = D[r].data(), inv = 1 / a[s];
32
      REP(i, m + 2) if(i != r \&\& abs(D[i][s])
           > ens) {
        T *b = D[i].data(), inv2 = b[s] * inv;
        REP(j, n + 2) b[j] -= a[j] * inv2;
        b[s] = a[s] * inv2;
      REP(j, n + 2) if(j != s) D[r][j] *= inv;
                                               12 ll count triangle rectangle intersection(ll
      REP(i, m + 2) if(i != r) D[i][s] *= -inv
      D[r][s] = inv;
      swap(B[r], N[s]);
    bool simplex(int phase) {
      int x = m + phase - 1;
      while(true) {
        int s = -1:
        REP(j, n + 1) if(N[j] != -phase) ltj(D
        if(D[x][s] >= -eps) return true;
        int r = -1:
        REP(i, m) {
          if(D[i][s] <= eps) continue;</pre>
          if(r == -1 || MP(D[i][n+1] / D[i][s
               ], B[i]) < MP(D[r][n+1] / D[r][s | template < int B>
               ], B[r]) r = i;
        if(r == -1) return false;
        pivot(r, s);
57
    T solve(vd &x) {
      int r = 0;
      FOR(i, 1, m) if(D[i][n+1] < D[r][n+1]) r
      if(D[r][n+1] < -eps) {
        pivot(r, n);
        if(!simplex(2) || D[m+1][n+1] < -eps)</pre>
             return -inf;
        REP(i, m) if(B[i] == -1) {
          int s = 0;
          FOR(j, 1, n + 1) ltj(D[i]);
          pivot(i, s);
      bool ok = simplex(1); x = vd(n);
      REP(i, m) if(B[i] < n) \times [B[i]] = D[i][n]
      return ok ? D[m][n+1] : inf;
73
74 };
  6.23 Triangle
```

#### 6.24 Xor-Basis

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```
2 struct xor basis {
    using T = long long;
    bool zero = false, change = false;
    int cnt = 0;
    array < T, B > p = {};
    vector<T> d;
    void insert(T x) {
      IREP(i, B) {
        if(x >> i & 1) {
           if(!p[i]) {
            p[i] = x, cnt++;
            change = true;
            return;
          } else x ^= p[i];
      if(!zero) zero = change = true;
      get min() {
      if(zero) return 0;
22
      REP(i, B) if(p[i]) return p[i];
23
      get max() {
24
      T ans = 0:
      IREP(i, B) ans = max(ans, ans ^ p[i]);
      return ans;
28
      get_kth(long long k) {
      k++;
31
      if(k == 1 && zero) return 0;
      k -= zero;
      if(k >= (1LL << cnt)) return -1;
      update();
      T ans = 0;
```

2 // Also representable as sum {0 <= x <= C /

return m \* (m + 1) / 2 \* k + (m + 1) \* h

+ count\_triangle(B, A - k \* B, C -

A) floor((C - Ax) / B + 1).

3 11 count triangle(11 A, 11 B, 11 C) {

11 h = (C - m \* A) / B + 1;

B \* (k \* m + h));

return RA \* RB:

11 // Counts  $\theta \leftarrow x \leftarrow RA$ ,  $\theta \leftarrow y \leftarrow RB$  such that

A, 11 B, 11 C, 11 RA, 11 RB) {

return count\_triangle(A, B, C) -

 $Ax + By \leftarrow C$ . Requires A, B > 0.

if(C < 0 || RA <= 0 || RB <= 0) return

count\_triangle(A, B, C - A \* RA) -

count\_triangle(A, B, C - B \* RB);

if(C >= A \* (RA - 1) + B \* (RB - 1))

**if**(C < 0) **return** 0;

if(A < B) swap(A, B);

11 m = C / A, k = A / B;

: m(SZ(b)), n(SZ(c)), N(n+1), B(m),  $D \mid // Counts x$ ,  $y \ge 0$  such that Ax + By < C. Requires A, B > 0. Runs in log time.

```
REP(i, SZ(d)) if(k \gg i \& 1) ans ^= d[i
           ];
      return ans:
    bool contains(T x) {
      if(x == 0) return zero;
      IREP(i, B) if(x \gg i \& 1) \times ^= p[i];
      return x == 0:
    void merge(const xor_basis& other) { REP(i
         , B) if(other.p[i]) insert(other.p[i])
    void update() {
      if(!change) return;
      change = false;
      d.clear();
      REP(j, B) IREP(i, j) if(p[j] \gg i \& 1) p
           [j] ^= p[i];
      REP(i, B) if(p[i]) d.pb(p[i]);
52 };
```

#### 估計值

- Estimation
  - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
  - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1, 1, 2, 3, 5, 7, 11, 15, 22, 30 for  $n = 0 \sim 9,627 \text{ for } n = 20, \sim 2e5 \text{ for }$
  - $\begin{array}{ll} n & \text{distinct} & \text{elements:} & B(n) & = & \det(M) = \sum_{S:A \rightarrow B} (-1)^{t(\sigma(S))} \prod_{i=1}^n \omega(S_i) \\ 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975, 678570, 4213597, \\ 27644437, 190899322, \ldots \end{array}$ - Total number of partitions

#### 定理 6.26

· Cramer's rule

$$ax + by = e$$

$$cx + dy = f$$

$$x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Burnside's Lemma

Let us calculate the number of necklaces of n pearls. where each pearl has m possible colors. Two necklaces are symmetric if they are similar after rotating them. There are n ways to change the position of a necklace, because we can rotate it  $0, 1, \ldots, n_1$  steps clockwise.

If the number of steps is 0, all  $m^n$  necklaces remain the same, and if the number of steps is 1, only the m necklaces where each pearl has the same color remain the same. More generally, when the number of steps is k, a total of  $m^{\gcd(k,n)}$  necklaces remain the same. The reason for this is that blocks of pearls of size gcd(k, n)will replace each other. Thus, according to Burnside's lemma, the number of necklaces is  $\sum_{i=0}^{n-1} \frac{m^{\gcd(i,n)}}{n}$ For example, the number of necklaces of length  $\stackrel{n}{4}$  with 3 colors is  $\frac{3^4+3+3^2+3}{4}=24$ 

· Lindstr□m-Gessel-Viennot Lemma

定義

 $\omega(P)$  表示 P 這條路徑上所有邊的邊權之積。( 路徑 計數時,可以將邊權都設為1)(事實上,邊權可以為 生成函數) e(u,v) 表示 u 到 v 的 \*\* 每一條 \*\* 路徑 P的  $\omega(P)$  之和 '即  $e(u,v) = \sum\limits_{P:u \to v} \omega(P)$  。 起點 集合 A · 是有向無環圖點集的一個子集 · 大小為 n 。 終點集合 B, 也是有向無環圖點集的一個子集, 大小 也為  $n \cdot -$ 組  $A \rightarrow B$  的不相交路徑  $S : S_i$  是一條從  $A_i$  到  $B_{\sigma(S)_i}$  的路徑 ( $\sigma(S)$  是一個排列),對於任 何  $i \neq i \cdot S_i$  和  $S_i$  沒有公共頂點。 $t(\sigma)$  表示排列  $\sigma$ 

$$M = \begin{bmatrix} e(A_1, B_1) & e(A_1, B_2) & \cdots & e(A_1, B_n) \\ e(A_2, B_1) & e(A_2, B_2) & \cdots & e(A_2, B_n) \end{bmatrix}$$

$$\vdots & \vdots & \ddots & \vdots \\ e(A_n, B_1) & e(A_n, B_2) & \cdots & e(A_n, B_n) \end{bmatrix}$$

$$\det(M) = \sum_{S:A\to B} (-1)^{t(\sigma(S))} \prod_{i=1}^{n} \omega(S_i)$$

其中  $\sum_{S:A\to B}$  表示滿足上文要求的  $A\to B$  的每一組 不相交路徑 S。

· Kirchhoff's Theorem

的逆序對個數。

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is
- $|\det(\tilde{L}_{11})|$ . The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i, j) \in E$ , otherwise  $d_{ij}=-d_{ji}.\;rac{rank(D)}{2}$  is the maximum match-

- · Cayley's Formula
  - Given a degree sequence  $d_1, d_2, \ldots, d_n$ for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.

- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1, 2, \ldots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .
- Erd□s–Gallai theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + \cdots + d_n$  is even

and 
$$\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$
 holds for every  $1 \le k \le n$ .

· Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq$  $\cdots \geq a_n$  and  $b_1, \ldots, b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$ 

· Fulkerson-Chen-Anstee theorem

A sequence  $(a_1, b_1), \ldots, (a_n, b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only  $\text{if } \sum_{i=1}^{n} a_i = \sum_{i=1}^{n} b_i \text{ and } \sum_{i=1}^{k} a_i \leq \sum_{i=1}^{k} \min(b_i, k-1) +$  $\sum_{i=1}^{n} \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$ 

M□bius inversion formula

$$\begin{array}{lll} -f(n) &=& \sum_{d\mid n} g(d) & \Leftrightarrow & g(n) &=& \\ & \sum_{d\mid n} \mu(d) f(\frac{n}{d}) \\ -f(n) &=& \sum_{n\mid d} g(d) & \Leftrightarrow & g(n) &=& \\ & \sum_{n\mid d} \mu(\frac{d}{n}) f(d) & & & \end{array}$$

- Spherical cap
  - A portion of a sphere cut off by a plane. - r: sphere radius, a: radius of the base of the cap,
  - h: height of the cap,  $\theta$ : arcsin(a/r).
  - Volume =  $\pi h^2 (3r h)/3 = \pi h (3a^2 + h)$  $h^2)/6 = \pi r^3 (2 + \cos \theta)(1 - \cos \theta)^2/3.$ - Area =  $2\pi rh = \pi (a^2 + h^2) = 2\pi r^2(1 - \cos \theta)^2/3$

#### 6.27 數字

Bernoulli numbers

$$\begin{split} B_0 - 1, B_1^{\pm} &= \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ \sum_{j=0}^m \binom{m+1}{j} B_j &= 0, \text{ EGF is } B(x) = \frac{x}{\epsilon^x - 1} = \\ \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}. \end{split}$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

 Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} &S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = \\ &S(n,n) = 1 \\ &S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n \end{split}$$

· Pentagonal number theorem

 $x^n = \sum_{i=0}^n S(n,i)(x)_i$ 

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

· Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

· Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1), k+1$  j:s s.t.  $\pi(j) \ge j, k$  j:s s.t.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{i=0}^{k} (-1)^{i} {n+1 \choose i} (k+1-j)^{n}$$

次方和

$$\sum_{k=1}^{n} k^3 = \left(\frac{n(n+1)}{2}\right)^2$$
$$\sum_{k=1}^{n} k^4 = \frac{1}{30} (6n^5 + 15n^4 + 10n^3 - n)$$

$$\sum_{k=1}^{n} k^5 = \frac{1}{12} (2n^6 + 6n^5 + 5n^4 - n^2)$$

$$\sum_{k=1}^{n} k^{6} = \frac{1}{42} (6n^{7} + 21n^{6} + 21n^{5} - 7n^{3} + n)$$

$$\sum_{k=1}^{n} k^{p} = \frac{1}{p+1} (n \sum_{i=1}^{p} (n+1)^{i} - \sum_{i=2}^{p} {i \choose p+1} \sum_{k=1}^{n} k^{p+1-i})$$

#### 歐幾里得類算法 6.28

#### • $m = \lfloor \frac{an+b}{c} \rfloor$

#### • Time complexity: $O(\log n)$

# $f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor$

```
g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                                      = \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), \\ 0, \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)), \\ -h(c, c-b-1, a, m-1)), \end{cases} 
h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^{2}
                                                     \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \end{cases}
```

```
+2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
+2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
 nm(m+1) - 2g(c, c-b-1, a, m-1)
\left| -2f(c,c-b-1,a,m-1) - f(a,b,c,n) \right| // find_by_order order_of_key
```

 $+h(a \bmod c, b \bmod c, c, n)$ 

#### 6.29 生成函數

```
• Ordinary Generating Function A(x) = \sum_{i>0} a_i x^i
```

```
 \begin{array}{lll} -A(rx)\Rightarrow r^na_n & & \\ -A(x)+B(x)\Rightarrow a_n+b_n & & \\ -A(x)B(x)\Rightarrow \sum_{i=0}^n a_ib_{n-i} & & \\ -A(x)^k\Rightarrow \sum_{i=1}^n a_ib_{n-i} & & \\ -A(x)^k\Rightarrow \sum_{i=1}^n a_ib_{n-i} & & \\ \end{array} 
- xA(x)' \Rightarrow na_n
- \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i
```

• Exponential Generating Function A(x) $\sum_{i\geq 0} \frac{a_i}{i!} x_i$ 

```
\begin{array}{lll} -A(x)+B(x)\Rightarrow a_{n}+b_{n} & \text{ in }\\ -A^{(k)}(x)\Rightarrow a_{n}+k_{n}\\ -A(x)B(x)\Rightarrow \sum_{i=0}^{n}\binom{n}{i}a_{i}b_{n-i}\\ -A(x)^{k}\Rightarrow \sum_{i_{1}+i_{2}+\cdots+i_{k}=n}\binom{n}{i_{1},i_{2},\ldots,i_{k}}a_{i_{1}} \stackrel{k}{a} & \text{ inline ull rng() }\\ -xA(x)\Rightarrow na_{n} & \text{ static ull } \mathbb{Q}=48\\ & \text{ of } n=0 <<7; \end{array}
```

· Special Generating Function

```
 - (1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i 
 - \frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{n}{n-1} x^i
```

#### Misc

#### 7.1 gc

```
i inline char gc() {
                                                                                    static const size_t sz = 65536;
                                                                                    static char buf[sz];
                                                                                    static char *p = buf, *end = buf;
= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) & \text{if (p = end) end = buf + fread(buf, 1, sz,} \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \\ 0, & n < 0 \vee \omega \\ 0, & \text{return *p++}; \end{cases}
```

```
7.2 next-combination
```

```
return ((comb & ~y) / x >> 1) | y;
```

#### **7.3 PBDS**

```
i #include <ext/pb ds/assoc container.hpp>
2 using namespace gnu pbds;
3 tree<11, null type, less<11>, rb tree tag,
      tree order statistics node update> st;
 // __float128_t
 for(int i = bs. Find first(); i < bs.size();</pre>
       i = bs. Find next(i));
```

#### 7.4 python

```
| from decimal import Decimal, getcontext
a = pow(Decimal(2), 82589933) - 1
```

#### 7.5 rng

```
static ull Q = 48763;
Q ^= Q << 7;
Q ^= Q >> 9;
return Q & 0xFFFFFFFFULL;
```

```
7.6 rotate90
```

```
i vector<vector<T>> rotate90(const vector
      vector<T>>& a) {
     int n = sz(a), m = sz(a[0]);
     vector<vector<T>> b(m, vector<T>(n));
     REP(i, n) REP(j, m) b[j][i] = a[i][m - 1]
         - j];
     return b;
```

#### timer

```
1 clock t T1 = clock();
2 double getCurrentTime() { return (double) (
      clock() - T1) / CLOCKS PER SEC; }
```

| template < int ALPHABET = 26, char MIN CHAR =

#### String

#### 8.1 AC

13

```
'a'>
2 struct ac_automaton {
   struct Node {
     int fail = 0, cnt = 0;
     array<int, ALPHABET> go{};
   vector<Node> node;
   vi que;
   int new node() { return node.eb(), SZ(node
        ) - 1; }
   Node& operator[](int i) { return node[i];
   ac_automaton() { new_node(); } // reserve
   int insert(const string& s) {
     int p = 0:
     for(char c : s) {
       int v = c - MIN CHAR;
       if(node[p].go[v] == 0) node[p].go[v] =
            new_node();
       p = node[p].go[v];
     node[p].cnt++;
     return p;
   void build() {
     que.reserve(SZ(node)); que.pb(0);
     REP(i, SZ(que)) {
       int u = que[i];
       REP(j, ALPHABET) {
         if(node[u].go[j] == 0) node[u].go[j]
               = node[node[u].fail].go[j];
         else {
           int v = node[u].go[j];
           node[v].fail = (u == 0 ? u : node[
                node[u].fail].go[j]);
```

```
que.pb(v);
35
```

#### 8.2 hash61

```
| const 11 M30 = (1LL << 30) - 1;
2 const 11 M31 = (1LL << 31) - 1;
3 const 11 M61 = (1LL << 61) - 1;
4 ull modulo(ull x){
    ull xu = x \gg 61;
    ull xd = x \& M61;
    ull res = xu + xd;
    if(res >= M61) res -= M61;
    return res;
ull mul(ull a, ull b){
    ull au = a >> 31, ad = a & M31;
    ull bu = b \gg 31, bd = b \& M31;
    ull mid = au * bd + ad * bu;
    ull midu = mid >> 30;
    ull midd = mid & M30;
16
    return modulo(au * bu * 2 + midu + (midd
         << 31) + ad * bd);
```

#### 8.3 KMP

```
1 // abacbaba -> [0, 0, 1, 0, 0, 1, 2, 3]
2 vi KMP(const vi& a) {
   int n = SZ(a);
   vi k(n);
   FOR(i, 1, n) {
     int j = k[i - 1];
     while(j > 0 && a[i] != a[j]) j = k[j -
     j += (a[i] == a[j]);
     k[i] = j;
  return k;
```

#### **8.4** LCP

```
i vi lcp(const vi& s, const vi& sa) {
   int n = SZ(s), h = 0;
   vi rnk(n), lcp(n - 1);
   REP(i, n) rnk[sa[i]] = i;
   REP(i, n) {
     h -= (h > 0);
     if(rnk[i] == 0) continue;
     int j = sa[rnk[i] - 1];
     for(; j + h < n && i + h < n; h++) if(s[</pre>
          j + h] != s[i + h]) break;
```

```
8.5 manacher

| // Length: (z[i] - (i & 1)) / 2 * 2 + (i & 1) / 2 * 2 + (i & 1) / 2 * 2 + (i & 2) / 2
```

lcp[rnk[i] - 1] = h;

return lcp;

12

13 }

#### 8.6 rolling-hash

```
const ll M = 911382323, mod = 972663749;
ll Get(vector<11>& h, int l, int r) {
    if(!1) return h[r]; // p[i] = M^i % mod
    ll ans = (h[r] - h[1 - 1] * p[r - 1 + 1])
        % mod;
    return (ans + mod) % mod;
}

vector<11> Hash(string s) {
    vector<11> ans(SZ(s));
    ans[0] = s[0];
    for(int i = 1; i < SZ(s); i++) ans[i] = (
        ans[i - 1] * M + s[i]) % mod;
    return ans;
}</pre>
```

#### **8.7 SAIS**

```
1  // mississippi
2  // 10 7 4 1 0 9 8 6 3 5 2
3 vi SAIS(string a) {
4    int n = SZ(a), m = *max_element(ALL(a)) +
        1;
5    vi pos(m + 1), x(m), sa(n), val(n), lms;
6    for(auto c : a) pos[c + 1]++;
7    REP(i, m) pos[i + 1] += pos[i];
8    vector<bool> s(n);
9    IREP(i, n - 1) s[i] = a[i] != a[i + 1] ? a
        [i] < a[i + 1] : s[i + 1];
10    auto ind = [&](const vi& ls){
11    fill(ALL(sa), -1);</pre>
```

```
auto L = [&](int i) { if(i >= 0 && !s[i
           ]) sa[x[a[i]]++] = i; };
       auto S = [\&](int i) \{ if(i >= 0 \&\& s[i]) \}
            sa[--x[a[i]]] = i; };
      REP(i, m) x[i] = pos[i + 1];
      IREP(i, SZ(ls)) S(ls[i]);
      REP(i, m) x[i] = pos[i];
      L(n - 1);
      REP(i, n) L(sa[i] - 1);
      REP(i, m) x[i] = pos[i + 1];
      IREP(i, n) S(sa[i] - 1);
    auto ok = [&](int i) { return i == n || (!)
         s[i - 1] && s[i]); };
    auto same = [&](int i,int j) {
        if(a[i++] != a[j++]) return false;
      } while(!ok(i) && !ok(j));
      return ok(i) && ok(j);
    FOR(i, 1, n) if(ok(i)) lms.pb(i);
    ind(lms);
    if(SZ(lms)) {
      int p = -1, w = 0;
      for(auto v : sa) if(v && ok(v)) {
        if(p != -1 && same(p, v)) w--;
        val[p = v] = w++;
      auto b = lms;
      for(auto& v : b) v = val[v];
      b = SAIS(b);
      for(auto& v : b) v = lms[v];
      ind(b);
42
    return sa;
```

#### 8.8 SAM

```
il// cnt 要先用 bfs 往回推, 第一次出現的位置是
       state.first_pos - |S| + 1
 struct Node { int go[26] = {}, len, link,
      cnt, first_pos; };
 Node SA[N]; int sz;
 void sa_init() { SA[0].link = -1, SA[0].len
      = 0, sz = 1; 
 int sa_extend(int p, int c) {
   int u = sz++;
   SA[u].first_pos = SA[u].len = SA[p].len +
       1:
   SA[u].cnt = 1;
   while(p != -1 && SA[p].go[c] == 0) {
     SA[p].go[c] = u;
     p = SA[p].link;
   if(p == -1) {
     SA[u].link = 0;
     return u;
   int q = SA[p].go[c];
   if(SA[p].len + 1 == SA[q].len) {
     SA[u].link = q;
     return u;
```

#### 8.9 smallest-rotation

```
string small_rot(string s) {
   int n = SZ(s), i = 0, j = 1;
   s += s;
   while(i < n && j < n) {
      int k = 0;
      while(k < n && s[i + k] == s[j + k]) k
      ++;
   if(s[i + k] <= s[j + k]) j += k + 1;
   else i += k + 1;
   j += (i == j);
   }
   int ans = i < n ? i : j;
   return s.substr(ans, n);
}</pre>
```

#### 8.10 Z

# ACM ICPC Judge Test NTHU LinkCutTreap

#### C++ Resource Test

```
#include <bits/stdc++.h>
using namespace std;

namespace system_test {
   const size_t KB = 1024;
   const size_t MB = KB * 1024;
   const size_t GB = MB * 1024;
}
```

```
chrono::duration<double> diff = end -
10 size t block size, bound;
                                                          begin;
  void stack size dfs(size t depth = 1) {
                                                     return diff.count():
   if (depth >= bound)
                                                   void runtime_error_1() {
    int8_t ptr[block_size]; // 若無法編譯將
                                                     // Segmentation fault
         block size 改成常數
                                                     int *ptr = nullptr;
    memset(ptr, 'a', block_size);
                                                     *(ptr + 7122) = 7122;
    cout << depth << endl;</pre>
                                                 42 }
    stack_size_dfs(depth + 1);
                                                   void runtime_error_2() {
                                                     // Segmentation fault
  void stack_size_and_runtime_error(size_t
                                                     int *ptr = (int *)memset;
       block size, size t bound = 1024) {
                                                     *ptr = 7122;
    system test::block size = block size;
                                                 48 }
    system_test::bound = bound;
    stack size dfs();
                                                   void runtime_error_3() {
                                                     // munmap_chunk(): invalid pointer
                                                     int *ptr = (int *)memset;
  double speed(int iter num) {
                                                     delete ptr;
    const int block_size = 1024;
                                                 54
    volatile int A[block_size];
    auto begin = chrono::high resolution clock
                                                   void runtime_error_4() {
         ::now();
                                                     // free(): invalid pointer
    while (iter num--)
                                                     int *ptr = new int[7122];
      for (int j = 0; j < block_size; ++j)</pre>
                                                     ptr += 1;
                                                     delete[] ptr;
    auto end = chrono::high resolution clock::
```

```
63 void runtime error 5() {
    // maybe illegal instruction
    int a = 7122, b = 0;
    cout << (a / b) << endl;</pre>
  void runtime error 6() {
    // floating point exception
    volatile int a = 7122, b = 0;
    cout << (a / b) << endl;
73 }
  void runtime error 7() {
    // call to abort.
    assert(false);
78 }
80 } // namespace system test
82 #include <sys/resource.h>
void print_stack_limit() { // only work in
       Linux
    struct rlimit 1;
    getrlimit(RLIMIT STACK, &1);
    cout << "stack_size = " << l.rlim_cur << "</pre>
          byte" << endl;</pre>
87 }
```