Assignment 4

Index Tuning – Selection

Database Tuning

New Group 8
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Notes

- Do not forget to run ANALYZE tablename after creating or changing a table.
- Use EXPLAIN ANALYZE for the query plans that you display in the report.

Experimental Setup

How do you send the queries to the database? How do you measure the execution time for a sequence of queries?

For our experiments we used the following hardware and software:

Component	Specs			
Processor Memory	i7-13700H 3.7-5.0 GHz 32 GiB			

Table 1: Hardware: Dell XPS 15 9530

Software	Version
OS	Ubuntu 22.04
Postgres	2.3.4
postgresql	42.7.3
MariaDB	10.6.16
mariadb-java-client	3.3.3
Java	18

Table 2: Software

Postgres was hosted on localhost, on which we also executed our experiments. The client was implemented in Java and gained access to our databases using the JDBC drivers listed above.

For the queries we used prepared statements, and the performance time was measured by determining the throughput of the queries in a 60-second time frame.

Clustering B⁺ Tree Index

Point Query Repeat the following query multiple times with different conditions for pubID.

```
SELECT * FROM Publ WHERE pubID = ?
```

Which conditions did you use?

For each query a random existing pubID was chosen.

Show the runtime results and compute the throughput.

Executed Queries: 2003130 Queries per Second: 33385.5 Query plan (for one of the queries):

Planning Time: 0.189 ms Execution Time: 0.067 ms

An index scan with the index $idx_clustering_pubid$ on the row pubID is performed, which make use of the B^+ tree index. Here the pubID is compared with a string.

Multipoint Query vs. Multipoint Query IN-Predicate – Low Selectivity Repeat the following query multiple times with different conditions for booktitle.

```
SELECT * FROM Publ WHERE booktitle = ?
SELECT * FROM Publ WHERE pubID IN (
    SELECT pubid FROM public.auth WHERE name IN (?, ?, ?))
```

Which conditions did you use?

To maintain low selectivity, we are using random non-empty book titles out of the dataset. Similarly, for the authors, we are choosing three random authors, but do not need to specify the "non-empty" condition, since an author is always given.

Show the runtime results and compute the throughput.

Multipoint Query

Executed Queries: 153010 Queries per Second: 2550.17

Multipoint Query IN-Predicate

Executed Queries: 83734 Queries per Second: 1395.57

Query plan (for one of the queries):

Multipoint Query

Similar to pubid, we again use a simple clustered index and compare it to a string, but now with the attribute booktitle.

Multipoint Query IN-Predicate

```
Nested Loop (cost=19.16..628.21 rows=72 width=112)
             (actual time=0.200..4.579 rows=294 loops=1)
  -> HashAggregate (cost=18.73..19.45 rows=72 width=23)
                     (actual time=0.171..0.229 rows=294 loops=1)
        Group Key: (auth.pubid)::text
        Batches: 1 Memory Usage: 77kB
        -> Index Scan using idx_clustering_name on auth
                              (cost=0.43..18.55 rows=72 width=23)
                              (actual time=0.041..0.114 rows=294 loops=1)"
                Index Cond: ((name)::text = ANY
                           ('{""William Kent"", ""Alfons Kemper"", ""Walid G. Aref""}'::text[]))"
  -> Index Scan using idx_clustering_pubid on publ
                           (cost=0.43..8.45 rows=1 width=112)
                           (actual time=0.014..0.014 rows=1 loops=294)
        Index Cond: ((pubid)::text = (auth.pubid)::text)
Planning Time: 0.824 ms
Execution Time: 4.625 ms
```

This query plan is significantly more complex because it involves a nested loop and a hash aggregate. The plan additionally uses a group key for the booktitle attribute. First, an index scan is performed on the name attribute of the auth table, comparing three different strings of type name, followed by a hash aggregate operation due to the use of the IN predicate. An index scan is then performed on the pubID attribute of the pubI table, comparing the group key of the inner selection. Finally, the results are linked using a nested loop.

Multipoint Query – High Selectivity Repeat the following query multiple times with different conditions for year.

```
SELECT * FROM Publ WHERE year = ?
```

Which conditions did you use?

Same as before, we are using random years from the dataset.

Show the runtime results and compute the throughput.

```
Executed Queries: 2422
Queries per Second: 40.37
```

Query plan (for one of the queries):

Planning Time: 0.054 ms Execution Time: 7.923 ms

Similar to the previous query plans, an index scan comparing a string is performed on the year attribute using the index idx_clustering_year.

Non-Clustering B⁺ Tree Index

Note: Make sure the data is not physically ordered by the indexed attributes due to the clustering index that you created before.

Point Query Repeat the following query multiple times with different conditions for pubID.

```
SELECT * FROM Publ WHERE pubID = ?
```

Which conditions did you use?

For each query a random existing pubID was chosen.

Show the runtime results and compute the throughput.

Executed Queries: 2040997 Queries per Second: 34016.62 Query plan (for one of the queries):

Planning Time: 0.052 ms Execution Time: 0.034 ms

The query makes use of the B⁺ tree index on the pubID attribute, just as with the clustering index. The main difference is that in a clustered index the table is sorted by the index, which is not the case here.

Multipoint Query vs. Multipoint Query IN-Predicate – Low Selectivity Repeat the following query multiple times with different conditions for booktitle.

```
SELECT * FROM Publ WHERE booktitle = ?
SELECT * FROM Publ WHERE pubID IN (
    SELECT pubid FROM public.auth WHERE name IN (?, ?, ?))
```

Which conditions did you use?

To maintain low selectivity, we are using random non-empty book titles out of the dataset. Similarly, for the authors, we are choosing three random authors, but do not need to specify the "non-empty" condition, since an author is always given.

Show the runtime results and compute the throughput.

Multipoint Query

Executed Queries: 149266 Queries per Second: 2487.77

Multipoint Query IN-Predicate Executed Queries: 64079

Queries per Second: 1067.98

Query plan (for one of the queries):

Multipoint Query

Again an index scan is performed on the booktitle attribute of the publ table. As previously we use the index for the string comparison instead of the unsorted table.

Multipoint Query IN-Predicate

```
Nested Loop (cost=303.62..938.06 rows=75 width=112)
             (actual time=1.741..6.267 rows=294 loops=1)
  -> HashAggregate (cost=303.19..303.94 rows=75 width=23)
                     (actual time=1.702..1.752 rows=294 loops=1)
        Group Key: (auth.pubid)::text
        Batches: 1 Memory Usage: 77kB
        -> Bitmap Heap Scan on auth (cost=13.87..303.00 rows=75 width=23)
                                      (actual time=0.105..1.573 rows=294 loops=1)"
              Recheck Cond: ((name)::text = ANY
                          ('{""William Kent"",""Alfons Kemper"",""Walid G. Aref""}'::text[]))"
              Heap Blocks: exact=206
              -> Bitmap Index Scan on idx_clustering_name
                                        (cost=0.00..13.85 rows=75 width=0)
                                        (actual time=0.085..0.085 rows=294 loops=1)"
                    Index Cond: ((name)::text = ANY
                          ('{""William Kent"",""Alfons Kemper"",""Walid G. Aref""}'::text[]))"
  -> Index Scan using idx_clustering_pubid on publ
                                        (cost=0.43..8.45 rows=1 width=112)
                                        (actual time=0.015..0.015 rows=1 loops=294)
        Index Cond: ((pubid)::text = (auth.pubid)::text)
Planning Time: 0.677 ms
Execution Time: 6.316 ms
```

First, an index scan with the three chosen authors as in the condition is performed on the name index of the auth table. Then, similarly to the clustering index, the hash aggregate operation is performed, followed by an index scan on the pubID attribute of the publ table. Finally, the results are joined using a nested loop.

Multipoint Query – High Selectivity Repeat the following query multiple times with different conditions for year.

```
SELECT * FROM Publ WHERE year = ?
```

Which conditions did you use?

Same as before, we are using random years from the dataset.

Show the runtime results and compute the throughput.

First, a Bitmap Index Scan is performed on the year attribute of the publ table, followed by a Bitmap Heap Scan.

Non-Clustering Hash Index

Note: Make sure the data is not physically ordered by the indexed attributes due to the clustering index that you created before.

Point Query Repeat the following query multiple times with different conditions for pubID.

```
SELECT * FROM Publ WHERE pubID = ?
```

Which conditions did you use?

For each query a random existing pubID was chosen.

Show the runtime results and compute the throughput.

```
Executed Queries: 2618888
Queries per Second: 43648.13
```

Query plan (for one of the queries):

For this plan an index scan is again performed on the pubID attribute of the publ table. Similar to every other scan on pubID.

Multipoint Query vs. Multipoint Query IN-Predicate – Low Selectivity Repeat the following query multiple times with different conditions for booktitle.

```
SELECT * FROM Publ WHERE booktitle = ?

SELECT * FROM Publ WHERE pubID IN (
    SELECT pubid FROM public.auth WHERE name IN (?, ?, ?))
```

Which conditions did you use?

To maintain low selectivity, we are using random non-empty book titles out of the dataset. Similarly, for the authors, we are choosing three random authors, but do not need to specify the "non-empty" condition, since an author is always given.

Show the runtime results and compute the throughput.

```
Multipoint Query
```

Executed Queries: 137115

Queries per Second: 2285.25

Multipoint Query IN-Predicate

Executed Queries: 114494 Queries per Second: 1908.23

Query plan (for one of the queries):

Multipoint Query

Unlike the B⁺ tree index, the hash index performs a Bitmap Index Scan on booktitle in combination with a Bitmap Heap Scan.

Multipoint Query IN-Predicate

```
Nested Loop (cost=298.12..892.90 rows=74 width=112)
             (actual time=1.236..4.378 rows=294 loops=1)
  -> HashAggregate (cost=298.12..298.86 rows=74 width=23)
                     (actual time=1.201..1.240 rows=294 loops=1)
        Group Key: (auth.pubid)::text
       Batches: 1 Memory Usage: 77kB
        -> Bitmap Heap Scan on auth (cost=12.58..297.94 rows=74 width=23)
                                      (actual time=0.038..1.110 rows=294 loops=1)"
              Recheck Cond: ((name)::text = ANY
                          ('{""William Kent"",""Alfons Kemper"",""Walid G. Aref""}'::text[]))"
              Heap Blocks: exact=206
              -> Bitmap Index Scan on idx_clustering_name
                            (cost=0.00..12.56 rows=74 width=0)
                            (actual time=0.023..0.023 rows=294 loops=1)"
                    Index Cond: ((name)::text = ANY
                             ('{""William Kent"", ""Alfons Kemper"", ""Walid G. Aref""}'::text[]))
  -> Index Scan using idx_clustering_pubid on publ
                             (cost=0.00..8.02 rows=1 width=112)
                             (actual time=0.010..0.010 rows=1 loops=294)
        Index Cond: ((pubid)::text = (auth.pubid)::text)
Planning Time: 0.410 ms
Execution Time: 4.417 ms
```

A Bitmap Heap Scan (including Bitmap Index Scan) is performed on the name attribute of the auth table, followed by hash aggregation. Then, an index scan is performed on pubID and finally the results are joined using a nested loop.

Multipoint Query – High Selectivity Repeat the following query multiple times with different conditions for year.

```
SELECT * FROM Publ WHERE year = ?
```

Which conditions did you use?

Same as before, we are using random years from the dataset.

Show the runtime results and compute the throughput.

Here another Bitmap Index Scan with a Bitmap Heap Scan is performed comparing the year.

Table Scan

Note: Make sure the data is not physically ordered by the indexed attributes due to the clustering index that you created before.

Point Query Repeat the following query multiple times with different conditions for pubID.

```
SELECT * FROM Publ WHERE pubID = ?
```

Which conditions did you use?

For each query a random existing pubID was chosen.

Show the runtime results and compute the throughput.

Executed Queries: 2152 Queries per Second: 35.87 Query plan (for one of the queries):

Here we are not using any index, because of this a parallel Seq Scan is performed on the publ table. Then Gather is used to collect the results from the two parallel workers, from the parallel Seq Scan. [1]

Multipoint Query vs. Multipoint Query IN-Predicate – Low Selectivity Repeat the following query multiple times with different conditions for booktitle.

```
SELECT * FROM Publ WHERE booktitle = ?
SELECT * FROM Publ WHERE pubID IN (
    SELECT pubid FROM public.auth WHERE name IN (?, ?, ?))
```

Which conditions did you use?

To maintain low selectivity, we are using random non-empty book titles out of the dataset. Similarly, for the authors, we are choosing three random authors, but do not need to specify the "non-empty" condition, since an author is always given.

Show the runtime results and compute the throughput.

```
Multipoint Query
```

Executed Queries: 1808 Queries per Second: 30.13

Multipoint Query IN-Predicate

Executed Queries: 342 Queries per Second: 5.70

Query plan (for one of the queries):

Similar to the previous query, a parallel Seq Scan is performed on the publ table, thereafter the results are gathered from the parallel workers.

```
Gather (cost=45542.30..74465.15 rows=73 width=112)
        (actual time=155.433..236.956 rows=294 loops=1)
 Workers Planned: 2
 Workers Launched: 2
 -> Parallel Hash Semi Join (cost=44542.30..73457.85 rows=30 width=112)
                              (actual time=141.319..216.325 rows=98 loops=3)
       Hash Cond: ((publ.pubid)::text = (auth.pubid)::text)
        -> Parallel Seq Scan on publ
                (cost=0.00..27566.39 rows=513839 width=112)
                (actual time=0.383..40.797 rows=411071 loops=3)
        -> Parallel Hash
               (cost=44541.92..44541.92 rows=30 width=23)
                (actual time=139.036..139.036 rows=98 loops=3)
             Buckets: 1024 Batches: 1 Memory Usage: 72kB
             -> Parallel Seq Scan on auth
                    (cost=0.00..44541.92 rows=30 width=23)
                        (actual time=89.358..138.993 rows=98 loops=3)
                    Filter: ((name)::text = ANY
```

```
('{""William Kent"",""Alfons Kemper"",""Walid G. Aref""}'::text[]))"
Rows Removed by Filter: 1031636
```

Planning Time: 0.690 ms Execution Time: 236.986 ms

For the IN-Predicate query, a Seq Scan is performed on auth. A parallel Hash Semi Join joins the results from the parallel Hash, which we used the Seq Scan for, and a parallel Seq scan on publ. The results are then gathered from the parallel workers, as before.

Multipoint Query – High Selectivity Repeat the following query multiple times with different conditions for year.

```
SELECT * FROM Publ WHERE year = ?
```

Which conditions did you use?

Same as before, we are using random years from the dataset.

Show the runtime results and compute the throughput.

Executed Queries: 1276 Queries per Second: 21.27

Query plan (for one of the queries):

Filter: ((year)::text = '2000'::text)
Rows Removed by Filter: 394137

Planning Time: 0.117 ms Execution Time: 78.799 ms

A parallel Seq Scan is performed on the publ table and then the results from the two workers are gathered.

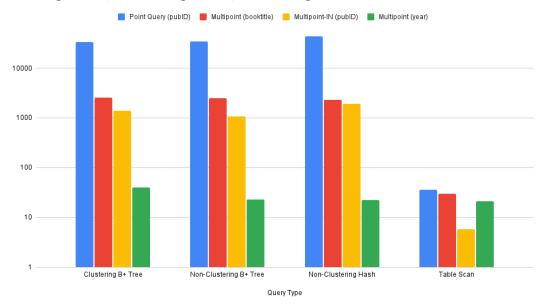
Discussion

Give the throughput of the query types and index types in queries/second.

	clustering	non-clust. B ⁺ tree	non-clust. hash	table scan
point (pubID)	33385.5	34016.62	43648.13	35.86
multipoint (booktitle)	2550.17	2487.77	2285.25	30.13
multipoint-IN (pubID)	1395.56	1067.98	1908.23	5.70
multipoint (year)	40.37	22.8	21.96	21.27

Discuss the runtime results for the different index types and the table scan. Are the results expected? Why (not)?

Clustering B+ Tree, Non-Clustering B+ Tree, Non-Clustering Hash und Table Scan



Firstly, we will look at the point query. The clustering and non-clustering B⁺ tree indexes have a similar throughput. The non-clustering hash index has the highest throughput, which is expected, as hash indexes are faster for point queries. As expected, the table scan has the lowest throughput, as there is no index that can be used.

Regarding the multipoint query with low selectivity, the B^+ indexes and the hash index performed somewhat similarly. The clustering B^+ index had the highest throughput, which makes sense since the clustering index sorts the data by the indexed attribute, which benefits ranged queries. Again, the table scan had the lowest throughput.

For the multipoint query with the IN-predicate, the hash index outperformed the B⁺ indexes. This outcome is quite interesting; from a naive perspective, one would expect the clustered B⁺ index to perform better than the hash index. However, since there are multiple exact matches involved in the IN-predicate, the hash index is faster. And just like with the previous queries, the table scan had the lowest throughput.

Lastly, for the multipoint query with high selectivity, the clustering B^+ index had the highest throughput, having almost double the throughput of all the other methods. This is expected, as the sorting of the data helps a lot with this type of query. The rest performed similarly, but the table scan still had the lowest throughput.

Time Spent on this Assignment

Time in hours per person:

• Florian Frauenschuh: 7

• Peter Lindner: **5.5**

• Alexander Weilert: 6.5

References

[1] May 2024. URL: https://www.postgresql.org/docs/current/how-parallel-query-works.html.