

PROF. DANILO CURVELO

### Quem sou eu?

### **Danilo Curvelo**

danilocurvelo@imd.ufrn.br

UFRN/IMD/CIVT/A216

Entusiasta da tecnologia Blockchain

### Carteiras:

**Bitcoin** bc1q8z8kkf24emy9577r7zmcyjc8t0ztk9xx49dsn5

**Ethereum** 0x6bE6ec844B5911476b03E5399147cA552702a471



# Quem são vocês?





PROF. **DANILO CURVELO** 

### Sobre o curso...

### **IMD0913**

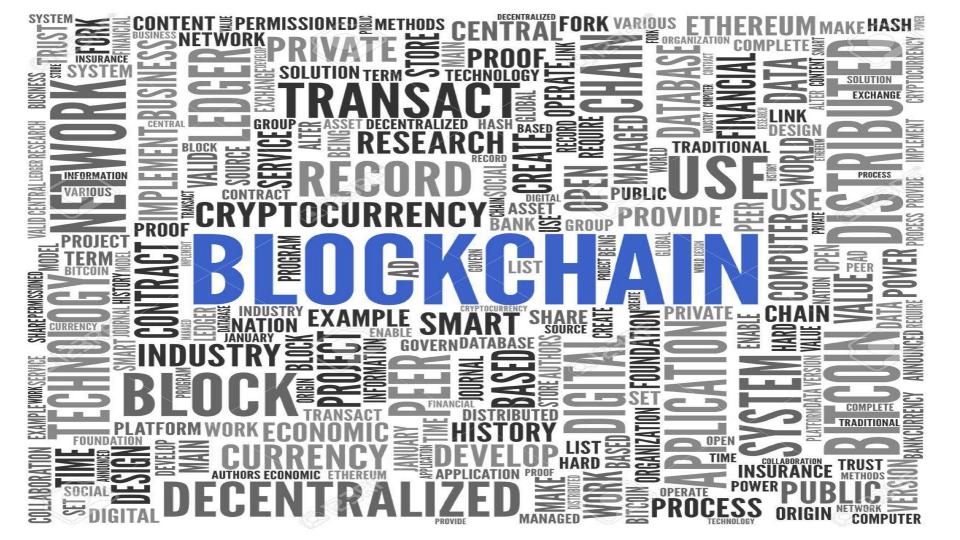
Blockchain e Aplicações Descentralizadas

60h - **24T56** 

CIVT/IMD A103

## **TECH TRENDS**

CompTIA	Gartner	Forbes
1. loT	1. Autonomous Things	1. Increased Automation
2. Al	2. Augmented Analytics	2. Blockchain
3. 5G	3. AI	3. Human/Al Collab
4. Serverless Computing	4. Digital Twins	4. loT
5. Blockchain	5. Edge Computing	5. VR/AR
6. Robotics	6. Immersive Technologies	6. Cybersecurity with ML/Al
7. Biometrics	7. Blockchain	7. Solutions to Tech Backslash
8. 3D Printing	8. Smart Spaces	8. Technology Convergence
9. VR/AR	9. Digital Ethics	
10. Drones	10. Quantum Computing	



# O que vamos aprender?

Tecnologia Blockchain

Fundamentos criptográficos

Consenso distribuído

**Smart Contracts** 

DApps

Estudos de caso:

Bitcoin

Ethereum



# O que não vamos aprender?

Economia;

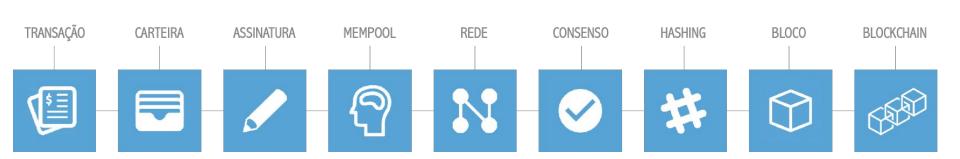
Investimentos;

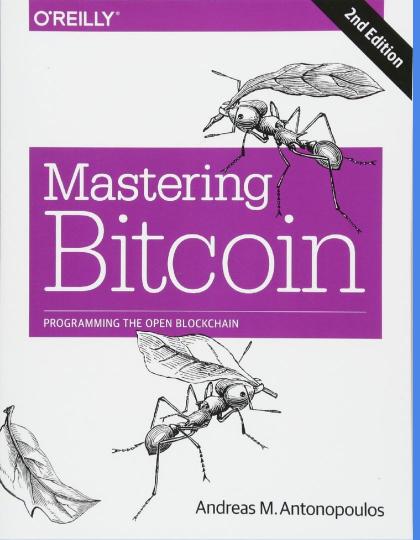
Mercado financeiro;

Como ficar rico com criptomoedas.



### ARQUITETURA DE UM **BLOCKCHAIN**

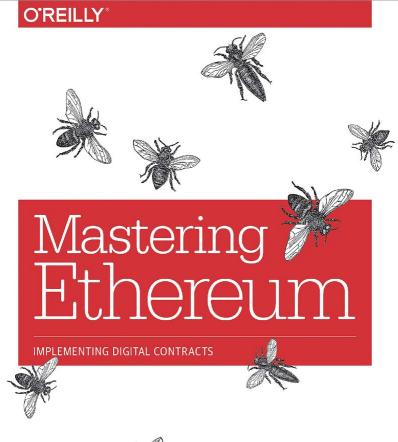




# MASTERING BITCOIN

# **Andreas Antonopoulos**

https://github.com/bitcoinbook/bitcoinbook

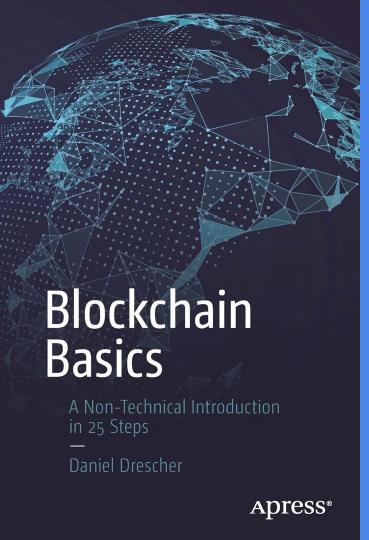




# MASTERING ETHEREUM

# Andreas Antonopoulos Gavin Wood

https://github.com/ethereumbook/ethereumbook



# BLOCKCHAIN BASICS

**Daniel Drescher** 

# A PEER-TO-PEER ELECTRONIC CASH SYSTEM

SATOSHI NAKAMOTO • OCTOBER 31, 2008

Abstract. A purely peer-to-peer version of electronic eash would allow online payments to be sent directly from one party to another without the main honefite are last if a trusted third harts and the solution but the main honefite are last if a trusted third harts against though a financial institution. Abstract. A purely peer-to-peer version of electronic cash would allow online payments to be sent directly from one party to another without directly from one party to another without party.

By through a financial institution. Digital signatures provide part of the solution, but the main benefits are lost if a trusted third party of the solution, but the main benefits are lost if a trusted third party of the solution to the double-chanding troublem without the main benefits are lost if a trusted third party of the solution to the double-chanding troublem without the payment of the solution to the double-chanding troublem without the payment of the solution to the double-chanding troublem without the party to another without the solution to the double-chanding troublem without the party to another without the solution to the double-chanding troublem to the solution to the double chanding troublem to the solution going through a financial institution. Digital signatures provide part of the solution, but the main benefits are lost y a trusted third party
is still required to prevent double-spending. We propose a solution to the double-spending problem using a peer-to-peer network handed without
is still required to prevent double-spending. We propose a solution to the double-spending forming a record that cannot be changed without
timestambs transactions by hashing them into an angoing chain of hash-based broad-af-reark is still required to prevent double-spending. We propose a solution to the double-spending problem using a peer-to-peer network. I ne network timestamps transactions by hashing them into an ongoing chain of hash-based proof-of-work, forming a record that cannot be changed without timestamps transactions by hashing them into an ongoing chain of hash-based proof-of-work, forming a record that it came from the largest timestamps transactions by hashing them into an ongoing chain of the commons of enemies introceed but transfer that it came from the largest proof-of-works are hard only corner as through the commons of enemies introceed but through that it came from the largest proof-of-works are hard only corner as through the commons of enemies introceed but through that it came from the largest proof-of-works are through the commons of enemies introceed but through the largest proof-of-works. timestamps transactions by hasting them into an ongoing chain of hast-based proof-of-work, forming a record inal cannot be changed without the largest redoing the proof-of-work. The longest chain not only serves as proof of the sequence of events witnessed, but proof that it came from the largest redoing the proof-of-work. The longest chain not only serves as proof of the sequence of events witnessed, but proof that it came from the are not conservation to attack the notwork. The value of controlled by nodes that are not conservation to attack the notwork. reasing the proof-of-work. The tongest chain not only serves as proof of the proof of CPU power. As long as a majority of CPU power is controlled by nodes that are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall are not cooperating to attack the network, they'll generate the lowest shall be not the lowest shall be not the lowest shall be not to be not the lowest shall be not the l pool of UPU power. As long as a majority of UPU power is controlled by nodes that are not cooperating to attack the network itself requires minimal structure. Messages are broadcast on a best effort basis, and nodes the longest chain and outpace attackers. The network itself requires minimal structure. the tongest chain and outpace anackers. The network usey requires minimal structure, Messages are oroaacast on a vest effort viasis, and nodes the tongest chain as proof of what happened while they were gone, can leave and rejoin the network at will, accepting the longest proof-of-work chain as proof of what happened while they were gone.

1. Introduction. Commerce on the Internet has come to rely almost exclusively on financial institutions are more absorbed to release a nearest object account. While the scarms records used asserted, for more absorbed institutions. 1. Introduction. Commerce on the Internet has come to rely almost extrained you financial sinstitutions are raised activated their parties to process electronic payments. While the system works well enough for most representational payments. While the system works well enough for most representations and suffers from the internet weeks work of the resurchanged Communication. serving as insured third parties to process electronic payments. While the system works well enough for most removed residence, field sudden from the indexent weaknesses of the trust based model. Completely universely and the system work works well as the system of the state of the system of the transections, it still utilier from the inherent weaknesses of the trust based model. Completely nonercoverable transactions are not really possible, since financial institutions cannot avoid mediating disputes. The cost of mediation increases resourcions ones, financial institutions cannot avoid mediating disputes. The cost of the mediation increases a resourcion costs, financial mediations resourcion of the mediation of the m remarking are not really possible, since financial institutions cannot avoid mediating disputes. The cost of mediating increases transaction costs, financial minimum practical transaction size and exture off the minimum practical transaction size and exture off the model-like foreverse transactions are all them in a become of a dark to a contract contract to the law of a dark to a contract contract to the law of a dark to a contract contract to the law of a dark to a contract contract to the law of a dark to a contract contract to the law of a dark to a contract contract to the law of a dark to the contract contract to the law of a dark to the contract contract to the law of a dark to the contract to the law of a dark to the contract to the law of a dark to the contract to the law of a dark to the contract to the law of a dark to the contract to the law of a dark to the contract to the law of a dark mediates increases transaction costs, limiting the minimum practical transaction size and cutting off the properties of possibility for small causal transactions, and there is a broader cost in the loss of ability to make non-coversible properties. With the possibility of reversal, the need for trust spreads. Merchants the possibility of reversal, the need for trust spreads. Merchants the possibility of reversal, the need for trust spreads the properties of the properties of the possibility of the possibility of reversal trust the vector of the possibility of the possibility of reversal trust the vector of the possibility of the possibili posurers for non-reversible services. With the possibility of reversal, the need for rust spreads. Merchant must be vary of their customers, handling them for more information than they would otherwise need. A large state of the customers is a supplied to the customers and the supplied to the customers and the supplied to the customers and the supplied to the customers are supplied to the customers and the supplied to the customers are supplied to the customers and the supplied to the customers are supplied to the customers and the customers are supplied to the customer must be vary of their customers, hasding them for more information than they would otherwise need. A creating percentage of final si accepted as unavoidable. These costs and payment uncertainties can be avoided control to the cost of certain percentage of final is accepted as unavaidable. These costs and payment uncertainties can be avoided in person by using physical currency, but no mechanism exists to make payments over a communication of the proposal of the proposal property of the proposal property of the prop in person by using physical currency, but no mechanism exists to make payments over a communications channel submet a functed parts. ¶ What is needed is an electronic payment system based on cryptographic proof mesoned or over allegions one near actions exercise to constant allegions of some allegions one near actions exercise to constant allegions of the constant actions exercise to constant allegions of the constant actions exercise to constant allegions. channel without a trusted party. § What is needed is an electronic payment system based on cryptographic in interest of crust, allowing any two willing parties to transact directly with each other without the need for a trusted during the company of the company proof instead of trust, allowing any two stilling parties to transict directly with each other without the need for a trusted lined party. Transactions that are computationally impractical to receive would prove selfer for a trusted third party. Transaction that are computationally impractical to recene would protect selfers from from front fr from frank, and routine escrow mechanisms could easily be implemented to protect bayers. In this paper, we prepare a solution to the double-specifies problem using a person-poper distributed timestamp server to the proper of the solution propose a solution to the double-spending problem using a peer-to-peer distributed timestamp server to secure computational proof of the chromological order of transactions. The system is secure as long as secure under order order order order order (FSF) remove them these computations are consistent order ord generate comparational proof of the enumological order of infrastretions, the system is secure as using a conference of the enumological order of infrastretions, the system is secure as using a conference of the enumological order of infrastretions. I modes contentry control more GPU power than any cooperating group of attacker nodes, Z.

visions. We define an electronic coin ava chain of digital signatures. Each owner transfers the coin to

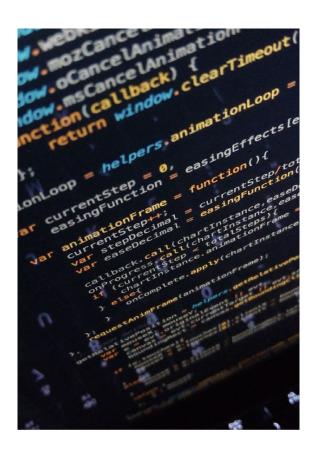
Morkle Treessure, with only the root included in the block's hash. Old blocks can then be compacted by another of the tree. The interior hashes do not need to be stored, § A block header with another of the tree. The interior hashes do not need to be stored, § A block header with a subject to the property of the store of the tree of the tree of the store subling off branches of the tree. The interior hashes do not need to be stored,  $\frac{\pi}{2}$  A block header with no manuscions would be about 80 bytes. If we suppose blocks are generated every 10 minutes, 80 bytes, 6g = 24 + 25 = 4 × ML near or with commuter sustained evaluation with 97 H or 8 AM are of state used Maconical Part of the Section 10 of th prinsections would be about 80 bytes. If we suppose blocks are generated every 10 minutes, 30 bytes © 6°24 at 23 to 24.23 Bpcf years With computer systems typically ediling with 2GB or RAM as of 2008, and Movement and the system of 1 bright new years of control to the a consideration of the block household and proposed of the block household and the system of the system o \*355 = 4.2MB per year With computer systems spically selling with 2GB of RAM as of 2008, and Moorks are predicting current growth of 1.2GB per year, storage should not be a problem even if the block backbeckness, and the same to be two in overones. Simulified Pavenant Verification. It is usuable to verify navarents without property to overone a simulation of the property of th Law predicting current growth of 1.2GH per year, storage should not be a problem even if the black headers and the begin in memory 8. Simplified Payment Verification. It is possible to verify payments which produces the black headers of the must be kept in memory. 8. Simplified Payment Verification. It is possible to verify payment without mining afful network node. A near only needs to keep a copy of the block leaders of the longest proof-fickness which the near only needs to keep a copy of the block leaders of the longest proof-fickness which which the near or the consequence nearests revolves until hex reconstructed by her the burness closer, and absoint running a full network node. A user only needs to keep a copy of the block headers of the longest proof-of-work chain, which he can get by querying network nodes until he's convinced he has the longest chain, and other the Verlack household follow the removement to the klock it's intercomposed in Placestry, shock the streams can for the Verlack household follow the removement to the klock it's intercomposed in Placestry, shock the streams can find the verlack of the placestry of the placestr chain, which he can get by querying setwork nodes until he's convinced he has the longest chain, and obtain the Merkle branch fashing the transaction to the block it's timestamped in. He can't check the transaction to the block it's timestamped in. He can't check the transaction to the block it's timestamped in. He can't check the transaction to the block it's timestamped in. He can't check the transaction to the block it's timestamped in. 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He can't check the transaction for himself, but by finking it to a place in the chain, he can see that a network node has accepted it, and block adoled after a further resuffers the measured but a proposed by a Ca such the swiffersion is optished as followed. himself, but by linking it to a place in the claim, he can see that a network node has accepted it, and blocks and the claim of the control o added after it further confirm the network has accepted it. ¶ As such, the verification is reliable as long as horses under control the network, but is more vulnerable if the network is overpowered by an attacker. While the control produce of the recognition for the models of the control of honest modes control the network, but is more vulnerable if the network is overpowered by an attacker. While relative framework modes can verify transactions for themselves, the simplified method can be fooled by an attacker. While the substitute of the substitute network nodes can verify transactions for themselves, the simplified method can be fooled by an attacker's fabricated transactions for as long as the attacker can continue to overpower the network. One strategy in the fabrication of the strategy of the s tabricated transactions for as long as the attacker can continue to overprower the network. One strategy to protect against this would be to accept alerts from network nodes when they detect an invalid block, prompting the now's ordered protection or double of the first black and absorbed transactions to confirm the invasculations. protect against this would be to accept alerts from network nodes when they detect an invalid block, prompting the user's software to download the full block and alerted transactions to confirm the inconsistency. Business the user's software to download the full block and alerted transactions to confirm the inconsistency. Businesses that receive frequent payments will probably still want to run their own nodes for more independent security and ordere confirmation. O Combinions and Sections Value Abbassach is consist to consider a least section. that receive frequent payments will probably still want to run their own nodes for more interpretent security and quirker verification. 9, Combining and Splitting Value. Additionship is would be possible to handle delicitation to worship to make a consequent proportion of the provided to contract the contract of the proportion of the provided to contract the contract of the proportion of the provided to contract the provided to contract the proportion of the provided to the provided to contract the proportion of the provided to the prov and quicker verification. 9, Combining and Splitting Value. Although it would be possible to landice coins of individually, it would be unwickly to make a separate transaction for every cent in a transfer. To allow value of the control and contro individually, it would be unweekly to make a separate transaction for every cent in a transfer. It is addownable for a separate framework of the control of be split and combined, transactions contain multiple inputs and outputs, Normally there will no entere a single input from a larger previous transaction or multiple inputs combining smaller amounts, and at most and the next of the combined of the container and one sentential the closure if some back in the smaller & it closured. input from a larger prevaust transaction or multiple inputs combining smaller amounts, and at most two objects one for the payment, and our returning the change, if any, back to the sender, § I should be noted to force where a remembring the change, if any, back to the sender, § I should be noted to force where a remembring the change, if any, back to the sender, § I should be noted to the force of the sender, § I should be noted to the force of the sender, § I should be noted to the force of the sender, § I should be noted to the force of the sender, § I should be noted to the force of the sender, § I should be noted to the sender of the sender outpute one for the payment, and one returning the entange, it and, tooks to the sential. If a success or onest that fair-out, where a transaction depends on several transactions, and those transactions depend on many onere a transaction depends on several transactions, and under transactions (septents us many) not banking model achieves a level of privacy by limiting access to

# Requisitos

Lógica de Programação (Python 3+)

HTML+CSS+JS

**REST-APIs** 



2022.2 Prof. **Danilo Curvelo** 

# Tecnologias

















2022.2 Prof. **Danilo Curvelo** 

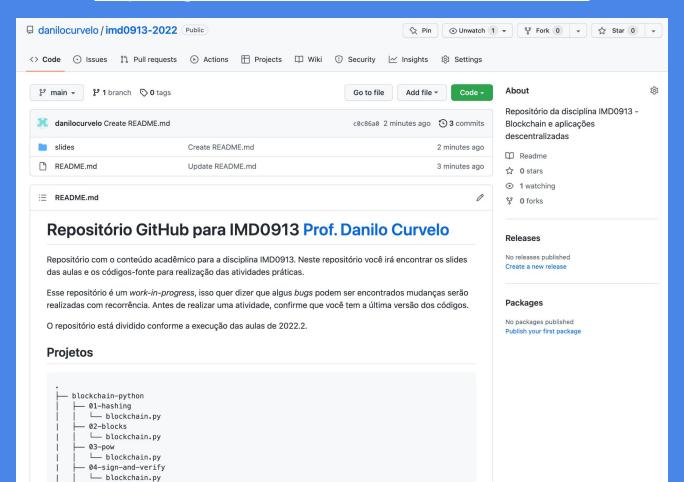
# Avaliação

Atividades práticas

Prova teórica \*

Projeto final

### https://github.com/danilocurvelo/imd0913-2022







Classrooms / imd0913-2022.2

### imd0913-2022.2

imd0913

Assignments 0

As Students 0

(\*) TAs and Admins 1

Settings

#### **Assignments**



#### Create an assignment to get started.

Create an individual assignment to generate an assignment repository for each student to work from. Or, create a group assignment and have students work collaboratively in groups from team repositories.

Create an assignment

Learn more about individual and group assignments.

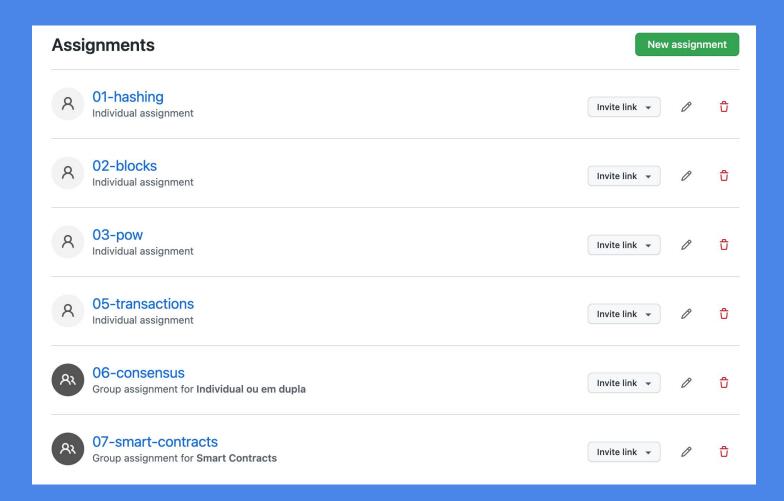


#### Need to teach Git & GitHub fundamentals?

The Classroom team has created an assignment for you to use to teach your students the fundamentals of Git & GitHub.

Use starter assignment

Learn more



### 01-hashing

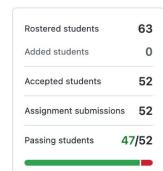
8 Individual assignment - Deadline Passed • Active

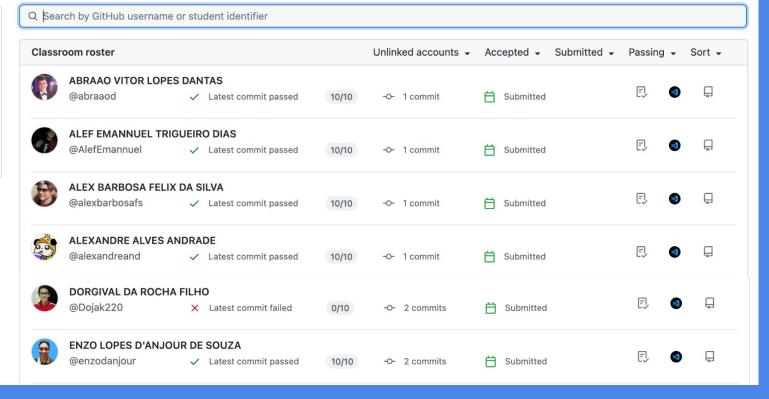


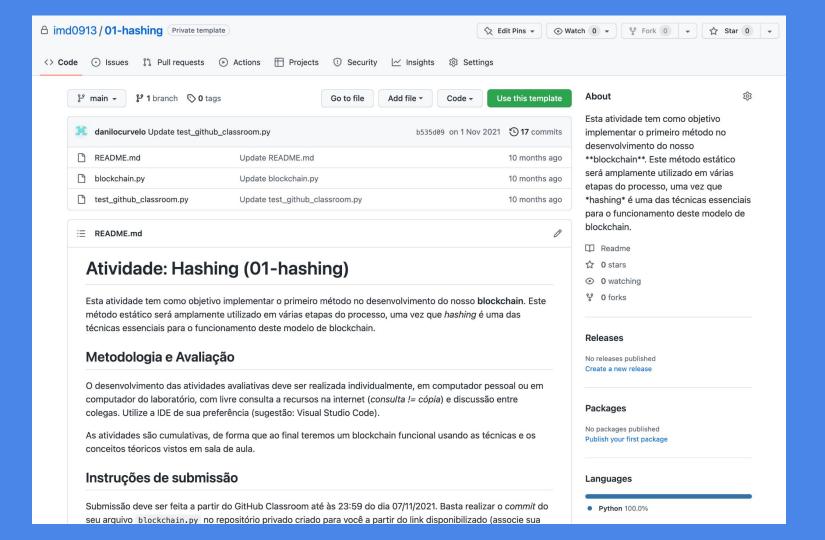
https://classroom.git [











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blockchain.py 01-hashing X blockchain.py 02-blocks
                                                                       blockchain.solution.py
                                                                                                   blockchain.py (
             01-hashing ▶ 🕏 blockchain.py ▶ ...
                    class Blockchain(object):
cks
                        @staticmethod
y 02-bl... 1
                        def generateHash(data):
nplete
                            # Implemente agui seu método para retornar a string referente ao hash SHA256 do argume
                            # Confira a documentação do hashlib: https://docs.python.org/3/library/hashlib.html
                            # Note que o argumento passado pode ser um objeto, portanto serialize o argumento ante
                            # Dica: Use o json.dumps() do módulo json.
                    # Testando sua implementação: espera-se um retorno True.
                    var1 = {
                                'nome': "Jon Snow",
                                'idade': 18,
                    expected_hash1 = "4145c81419ee987c94f741936c3277e9b281e2ffc9faa3edb5693128e1ee65c1"
                    var1 hash = Blockchain.generateHash(var1)
                    print(f'Dados: {var1}')
                    print(f'Hash gerado: {var1_hash}')
                    print(f'Hash esperado: {expected_hash1}')
                    print(f'Iguais? {expected_hash1==var1_hash}\n')
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```

#### [danilo@imd ~]\$ for months in 8 9 10 11 12; do cal \$months 2022; done

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Legenda:

■ Aula

Feriado

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September 2022

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October 2022

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