

Functech

- $\mathcal{A}, \mathcal{A}'$ ab. cats, $F: \mathcal{A} \rightarrow \mathcal{A}'$ left exact
- $RF: D^+(\mathcal{A}) \rightarrow D^+(\mathcal{A}')$ "unique" (if exists)

- \mathcal{A} ab. cat, $C(\mathcal{A})$ ab. cat, $K(\mathcal{A}) \triangleleft \text{cat}$,
 $D(\mathcal{A}) = K(\mathcal{A})/W \triangleleft \text{cat}$

Def. Let $\mathcal{T}_1, \mathcal{T}_2 \triangleleft \text{cats}$. A functor $F: \mathcal{T}_1 \rightarrow \mathcal{T}_2$ is a functor of $\triangleleft \text{cats}$ if it is additive & commutes with $[1]$ & sends $\text{dist.} \triangleleft$ to $\text{dist.} \triangleleft$.

Lemma $F: \mathcal{A} \rightarrow \mathcal{A}'$ induces $C(F), K(F)$.

Pf. Define $F(A^i, d^i) = (F(A^i), F(d^i)), \dots$

- what do we want from RF ?
- if \mathcal{A} has enough injectives, let \mathcal{J} be subcat of \mathcal{A} of inj. objects
- $K^+(\mathcal{J}) \subseteq K^+(\mathcal{A})$ full subcat of cpx A^i s.t.
 $A^i \in \mathcal{J} \forall i, A^i = 0 \forall i < 0$.

Lemma $\frac{K^+(\mathcal{J})}{N \cap K^+(\mathcal{J})} \xrightarrow{\alpha} \frac{K^+(\mathcal{A})}{N \cap K^+(\mathcal{A})} = D^+(\mathcal{A})$ equiv.

- goal. Find univ. prop. for $RF: D^+(\mathcal{A}) \rightarrow D^+(\mathcal{A}')$

$$\begin{array}{ccccc} D^+(\mathcal{A}) & \xrightarrow{\alpha^{-1}} & \frac{K^+(\mathcal{J})}{N \cap K^+(\mathcal{J})} & \xrightarrow{\quad} & D^+(\mathcal{A}') \\ & \nearrow \text{not really} & \uparrow & \nearrow K^+(F) & \uparrow \\ & \text{unique?} & K^+(\mathcal{J}) & \xrightarrow{\quad} & K^+(\mathcal{A}') \\ & \text{it is only an} & & & \\ & \text{equiv. of cats} & & & \end{array}$$

Lemma $K^+(F)(N \cap K^+(\mathcal{J})) \subseteq \mathcal{N}_{K^+(A')}$

Recall Enough inj. $\Rightarrow \forall A^\bullet \in \text{ob } K^+(A)$
 \exists qis $A^\bullet \xrightarrow{\sim} I^\bullet$ with $I^\bullet \in K^+(\mathcal{J})$

$$K^+(F)(A) \rightarrow K^+F(I) \xrightarrow[\text{can.}]{\sim} RF(A^\bullet) \\ \parallel \swarrow \sim \\ RF(I) \xleftarrow{\sim} RF(\varphi)$$

Def. Let $F: A \rightarrow A'$ be left exact functor of ab cats.
 A right derived functor for F is
 a pair (T, φ) such that

- i) $T: D^+(A) \rightarrow D^+(A')$ is a
 functor of Δ cats
- ii) φ is a nat transformation

$$\varphi: Q' \circ K^+(F) \Rightarrow T \circ Q$$

$$\begin{array}{ccc} K^+(A) & \xrightarrow{K^+(F)} & K^+(A') \\ Q \downarrow & \searrow \varphi & \downarrow Q' \\ D^+(A) & \xrightarrow{T} & D^+(A') \end{array}$$

Such that for any $G: D^+(A) \rightarrow D^+(A')$ Δ cat functor,
 the natural map

$$\text{Hom}(T, G) \longrightarrow \text{Hom}(Q' \circ K^+(F), G \circ Q)$$

$$\gamma: T \Rightarrow G \mapsto \gamma \circ \varphi$$

is bijective.

(whiskering)



Cor A right der. fctor is unique up to a can. eq. if it exists.

- so if it exists, we call it the r.d.fctor R_F .

Thm Let $F: A \rightarrow A'$ left exact fctor, assume A has enough F -inj., let $\mathcal{I}_F \subseteq A$ be choice of full subcat of F -inj. Then the functor

$$\begin{array}{ccc} D^+(A) & \xleftarrow{\sim} & D^+(\mathcal{I}_F) \xrightarrow{\exists!} D^+(A') \\ & & \uparrow \quad \quad \uparrow \\ & & K^+(\mathcal{I}_F) \xrightarrow{K^+F} K^+(A') \end{array}$$

is a right der fctor.

Rmk Choosing different $\mathcal{I}_F' \neq \mathcal{I}_F$ gives you different functors.

Example X top.sp., \underline{Ab}_X sh. of ab. grp on X , $\Gamma: \underline{Ab}_X \rightarrow Ab$ left exact. $R\Gamma$ can be computed using injective Mod \mathbb{Z} . $\mathcal{I} \in \underline{Ab}_X$ is F -acyclic if $H^i(\mathcal{I}) = 0 \forall i$, $\mathcal{I}_\Gamma =$ all acyclics.

- we define left derived fctors the same way, but F right exact, $T: D^-(A) \rightarrow D^-(A')$, $\varphi: T \circ Q \Rightarrow Q' \circ K^+F$.

Remark Let $F: A \rightarrow A'$ exact fctor, then $K(F)$ sends qiso to qiso, induces $D(F)$ which restricts to RF on $D^+(A)$ and LF on $D^-(A)$, since A is made of F -inj / F -proj.

Remark Composition of left exact fctors is left exact.

Lemma Assume RF, RG, RH exist, where $H = G \circ F$. Then there is an induced nat transf $RH \Rightarrow RG \circ RF$

$$\begin{array}{ccccc} \text{Pf.} & K^+(A) & \xrightarrow{K^+F} & K^+(A') & \xrightarrow{K^+G} & K^+(A'') & K^+(A) \rightarrow K^+(A'') \\ & \downarrow & \swarrow \varphi & \downarrow & \swarrow \psi & \downarrow & \Rightarrow \downarrow \swarrow \varphi \circ \psi \downarrow \\ & D^+(A) & \xrightarrow{RF} & D^+(A') & \xrightarrow{RG} & D^+(A'') & D^+(A) \rightarrow D^+(A'') \end{array}$$

Prop Same assumptions, assume $\exists J \subseteq A$ full subset of F -inj and $J' \subseteq A'$ $\dashv \vdash$ G -inj so that

- 1) J, J' can be used to def RF, RG
- 2) $\forall I \in \text{ob } J, F(I) \in \text{ob } J'$.

Then RH exists $\wedge RH \Rightarrow RG \circ RF$ is equiv.

Cor Assume $F: A \rightarrow A'$ exact, $G: A' \rightarrow A''$ left exact, A, A' have enough inj.
Then $R(G \circ F) \Rightarrow RG \circ D^+(F)$ is equiv.

Application $X \xrightarrow{f} Y \xrightarrow{g} Z$ mor of
 sch, $h = g \circ f$, f affine, which gives
 $f_*: \mathcal{O}_{\text{coh}} X \rightarrow \mathcal{O}_{\text{coh}} Y$ exact.

$$\forall \mathcal{F} \in \mathcal{O}_{\text{coh}} X \subseteq C(\mathcal{O}_{\text{coh}} X) \rightarrow K^+(\mathcal{O}_{\text{coh}} X)$$

$$R h_* (\mathcal{F}) = R g_* (K^+(f_*) \mathcal{F}) = R g_* (f_* \mathcal{F})$$

$$\Rightarrow \forall i \quad R^i h_* \mathcal{F} \xrightarrow{\sim_{\text{can.}}} R^i g_* (f_* \mathcal{F})$$

$$H^i(X, \mathcal{F}) \xrightarrow{\sim_{\text{can.}}} H^i(Y, f_* \mathcal{F}).$$