Fantechi Gesterday -A ab.cat, ABB also AxB. - ABB means & Ce66A, A - Baf B S.L. Hon(ABBC)-> Hom (A,C) \* Hom(B,C) sending 2 -> (loi, l-j) byects - we want dual statement with AcfaxB40B -step1 A -> A @B -> coker(i)= :B1, - PICK (606 A) then B= Troj Hom (B', c) = { 2 & Hom (ABB, C) | 201=0} Ho~(B,C) -since 2 mm (0,20j) byeds, -op. bijects, so by co Joneda, equivalent functors => B byects rstep2 -fix coobt, want y -> (pop, you) byects. - explicit inverse o -for ( 12 A / ABB ) / = 2004 1 jo 24. -exercise: p-1= (4,24). Ruk. By induction A, D. - DAn - A, + -- x An. Bot untrue for infinite I s.l. tieI, A; 606t, even if DA; and TT to both exist. This { (ai) is T tier } The Ais { same, but } is the same of the s

- Rnoeth mays of subcut of Mode of fig. R-mod.
-> show arb D's or T's don't exist  $-A comm. sing, P = A [x_1 - x_n], B = P/I,$   $Q = P[y], B = Q/J, P \xrightarrow{\pi} B$ -claim 1 15 9156 J/J2-> SLQ/A 8 Q B -Step 1 - dfoeP s.L. fi(y)=1(fo), fo=0 (=> fi(y)=0 - what can be said of 3/32 wat I/I2 in that case? - SZP/A = DPdxi > SZQ/A = DQdxi DQdy Q=POPy DPy20.--J2=I2DIyDPy2D .... => J/J2 = I/I2 (D) By - Step Z -gen. cuse, let Q = P[z], L(i); y-fo, d: Q ->Q -then I/T2->SZPHEDB 3 quse J/J2->SIQHEDBB 3 (SO=> quso J/J2->SZQHEDBB 3 (SO=> quso

Recall X -> 7 & mor Sch/spect, A north, X, 9 f.y. /Spec 4 m> [I/I2->52P/IBPB] "unique up to giso" in C(hodg): ((QCohx) Def. Let & be cat, Q Shore We say that the localization of Ent Q 15 a functor & 326' s.t. 1) tyeQ, F(y) iso 11) + G: E' -> E" s.t. + qeQ, G(q) 15 150 1 3/ HI (1/21-> 6" s.t. G= HOF 111) F: 062 -> 0681 (8 byective - exercise: of localisation exists, it is unique up to canonical iso, Fle Fr Problems

41) does F: P -> E' exist? q11) if yes, how explicit is &'? ai) not in general due to size issues. an) usualy not at all.

The Let A abel cat. Then the localisation of K(A) at Q= {q-150m} exists, and we call it the desired cut. D(A) of A. Given A, Beeob ((t) = 06 K(t) = 06 D(t), any mor in DIA) can be written as

("500)" A morin or "sink" A B

("500)" A good C(A)

B

("A)

B

("A)

Rmk. If 7: A->t' exact, then "obvious"
D(7) really works

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-let A ab. cat s define full s obcat o ((A), k(A), D(A))

s,t. t a \in b, a, b \in \mathbb{Z}. 0 \notin s 0 \notin s
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-exercise: 1) I nut map A => TznA in ((A)

11) h'(d) (5 150 Hizn

11) if yiA -> B qiso, Tz(y) qiso

11) Tz induces functors on K(t), D(d)

-let C(u, o) (A) = full subcat of apt s.f A'=0

Vica, Then Tzn is left als of inclusion C(u, o)(A) = (A)

Fecall d, Babel cat, of enough injectives,

J: A > B left exact. Define Rifid > B.

Hizo s.t. HAEGER HOSTONI, >>.

my resu of A, RiF(A) = hi(F(Is > I, >-.))

-with assumptions ouchanged, we want?

IRT: D+(A) > D+(B), uniquely defermined

by demanding if (T\*eD+(A) is cpx) 1

(I" is inj the Z) 1 (Jus s.t. The there)

then RF(I\*) = (... > F(I") -> F(I"+1) -> ...)

Intermediate Step Let  $\mathcal{I}$  all injectures in  $\mathcal{I}$   $Ciy(A) \subseteq C^{\dagger}(A)$  full subcent -claims induces  $D^{\dagger}(A) \longrightarrow D^{\dagger}(A)$ . This is equiv. of cats.

-replace "enough injs" by "enough acyclics"

-note that injectives are acyclic for left ex. functors

Since to >> J' >> >> y" >> 0 injues in d,

HFiA > B lexact >> 0 >> t(5') >> F(5) >> F(5") -> 0

-look at Gelfand-Manin, or Chion derived cats from Kashinava-Shapira Sheaves on infels