

t, r type variables

x, y, z, s variables

$\text{Typ}, \tau, \rho ::=$ type

- | t variable
- | $\tau_1 \rightarrow \tau_2$ function
- | $\forall (t.\tau)$ bind t in τ polymorphic
- | (τ) S

$\text{Exp}, e ::=$ expression

- | x variable
- | $\lambda(x : \tau)e$ bind x in e abstraction
- | $e_1(e_2)$ application
- | $\Lambda(t)e$ bind t in e type abstraction
- | $e[\tau]$ type application
- | (e) S

$\Delta ::=$ type formation hypothesis

- | **empty**
- | $\Delta, t \text{ type}$

$\Gamma ::=$ typing hypothesis

- | **empty**
- | $\Gamma, x : \tau$

$\Delta \vdash \tau \text{ type}$ type formation

$$\frac{t \in \Delta}{\Delta \vdash t \text{ type}} \quad \text{TYPE_VAR}$$

$$\frac{\Delta \vdash \tau_1 \text{ type} \quad \Delta \vdash \tau_2 \text{ type}}{\Delta \vdash \tau_1 \rightarrow \tau_2 \text{ type}} \quad \text{TYPE_ARR}$$

$$\frac{\Delta, t \text{ type} \vdash \tau \text{ type}}{\Delta \vdash \forall (t.\tau) \text{ type}} \quad \text{TYPE_ALL}$$

$\Delta \Gamma \vdash e : \tau$ typing

$$\frac{\Delta \vdash \tau \text{ type} \quad x : \tau \in \Gamma}{\Delta \Gamma \vdash x : \tau} \quad \text{EXP_VAR}$$

$$\frac{\Delta \vdash \tau_1 \text{ type} \quad \Delta \Gamma, x : \tau_1 \vdash e : \tau_2}{\Delta \Gamma \vdash \lambda(x : \tau_1)e : \tau_1 \rightarrow \tau_2} \quad \text{EXP_LAM}$$

$$\frac{\Delta \Gamma \vdash e_1 : \tau_2 \rightarrow \tau \quad \Delta \Gamma \vdash e_2 : \tau_2}{\Delta \Gamma \vdash e_1(e_2) : \tau} \quad \text{EXP_AP}$$

$$\frac{\Delta, t \text{ type} \Gamma \vdash e : \tau}{\Delta \Gamma \vdash \Lambda(t)e : \forall (t.\tau)} \quad \text{EXP_LAM}$$

$$\frac{\Delta \Gamma \vdash e : \forall (t.\tau') \quad \Delta \vdash \tau \text{ type}}{\Delta \Gamma \vdash e[\tau] : [\tau/t]\tau'} \quad \text{EXP_APP}$$

$e \text{ val}$

$\overline{\lambda(x : \tau)e \text{ val}}$ VAL_LAM

$\overline{\Lambda(t)e \text{ val}}$ VAL_LAM

$e_1 \mapsto e_2$

$\frac{e_2 \text{ val}}{\lambda(x : \tau_1)e(e_2) \mapsto [e_2/x]e}$ RED_LAM

$\frac{e_1 \mapsto e'_1}{e_1(e_2) \mapsto e'_1(e_2)}$ RED_AP1

$\frac{e_1 \text{ val} \quad e_2 \mapsto e'_2}{e_1(e_2) \mapsto e_1(e'_2)}$ RED_AP2

$\overline{\Lambda(t)e[\tau] \mapsto [\tau/t]e}$ RED_LAM

$\frac{e \mapsto e'}{e[\tau] \mapsto e'[\tau]}$ RED_APP

$\Delta \Gamma \vdash e_1 \equiv e_2 : \tau$

$\frac{\Delta \Gamma \vdash e : \tau}{\Delta \Gamma \vdash e \equiv e : \tau}$ EQ_REFL

$\frac{\Delta \Gamma \vdash e_1 \equiv e_2 : \tau}{\Delta \Gamma \vdash e_2 \equiv e_1 : \tau}$ EQ_COMM

$\frac{\Delta \Gamma \vdash e_1 \equiv e_2 : \tau \quad \Delta \Gamma \vdash e_2 \equiv e_3 : \tau}{\Delta \Gamma \vdash e_1 \equiv e_3 : \tau}$ EQ_TRANS

$\frac{\Delta \Gamma, x : \tau' \vdash e_1 \equiv e_2 : \tau}{\Delta \Gamma \vdash \lambda(x : \tau')e_1 \equiv \lambda(x : \tau')e_2 : \tau}$ EQ_LAM

$\frac{\Delta \Gamma \vdash e_1 \equiv e'_1 : \tau' \rightarrow \tau \quad \Delta \Gamma \vdash e_2 \equiv e'_2 : \tau'}{\Delta \Gamma \vdash e_1(e_2) \equiv e_1(e'_2) : \tau}$ EQ_AP0

$\frac{\Delta, t \text{ type } \Gamma \vdash e_1 \equiv e_2 : \tau}{\Delta \Gamma \vdash \Lambda(t)e_1 \equiv \Lambda(t)e_2 : \forall(t.\tau)}$ EQ_LAM

$\frac{\Delta \Gamma \vdash e_1 \equiv e_2 : \forall(t.\tau)}{\Delta \Gamma \vdash e_1[\tau'] \equiv e_2[\tau'] : [\tau'/t]\tau}$ EQ_APP0

$\frac{\Delta \Gamma, x : \tau_1 \vdash e_2 : \tau_2 \quad \Delta \Gamma \vdash e_1 : \tau_1}{\Delta \Gamma \vdash (\lambda(x : \tau_1)e_2)(e_1) \equiv [e_1/x]e_2 : \tau_2}$ EQ_AP

$\frac{\Delta, t \text{ type } \Gamma \vdash e : \tau \quad \Delta \vdash \rho \text{ type}}{\Delta \Gamma \vdash \Lambda(t)e[\rho] \equiv [\rho/t]e : \tau}$ EQ_APP

Definition rules: 24 good 0 bad

Definition rule clauses: 55 good 0 bad