```
tmvar, x, y variables
 typ, \tau
                           ::=
                                                                                                                                   types
                                       \mathbf{nat}
                                                                                                                                         Natural numbers
                                                                                                                                         Function types
                                       \tau_1 \rightarrow \tau_2
                                                                                                                                   expressions
 exp, e
                                                                                                                                         Variables
                                       \boldsymbol{x}
                                                                                                                                         Zero
                                       \mathbf{z}
                                                                                                                                         Successor
                                      \mathbf{rec}\ e\{\mathbf{z} \to e_0; \mathbf{s}\ x \to e_1\} \quad \mathsf{bind}\ x \mathsf{ in }\ e_1
                                                                                                                                         Primitive recursion over nats
                                       \lambda(x:\tau)e
                                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;e
                                                                                                                                         Functions
                                                                                                                                         Application
 env, \Gamma
                                                                                                                                   typing environment
                                                                                                                                         empty
                            \Gamma, x : \tau
                                                                                                                                         cons
e val
                                                                                             \frac{}{\mathbf{z}\,\mathbf{val}} VAL_Z
                                                                                                              VAL_S
                                                                                             e val
                                                                                            \overline{\mathbf{s} \, e \, \mathbf{val}}
                                                                                                                  VAL_ABS
                                                                                   \overline{\lambda(x\!:\!\tau)e\,\mathbf{val}}
  \Gamma \vdash e : \tau
                                                                                 uniq \Gamma
                                                                                 \frac{x:\tau\ \mathbf{in}\ \Gamma}{\Gamma\vdash x:\tau} \quad \text{TYPING\_VAR}
                                                                                 \frac{\mathbf{uniq}\,\Gamma}{\Gamma\vdash\mathbf{z}:\mathbf{nat}}\quad {}^{\mathrm{TYPING\_Z}}
                                                                                  \Gamma \vdash e : \mathbf{nat}
                                                                                                                  TYPING_S
                                                                                \overline{\Gamma \vdash \mathbf{s} \, e : \mathbf{nat}}
                                                                  \Gamma \vdash e : \mathbf{nat}
                                                                  \Gamma \vdash e_0 : \tau
                                                                  \Gamma, x : \mathbf{nat} \vdash e_1 : \tau \to \tau
                                                                  e_1 val
                                                       \frac{e_1 \operatorname{val}}{\Gamma \vdash \operatorname{rec} e\{\mathbf{z} \to e_0; \mathbf{s} x \to e_1\} : \tau}
                                                                                                                                    {\tt TYPING\_REC}
                                                                 \frac{\Gamma, x : \tau_1 \vdash e : \tau_2}{\Gamma \vdash \lambda(x \colon \tau_1) e : \tau_1 \to \tau_2} \quad \text{TYPING\_ABS}
                                                                         \Gamma \vdash e_1 : \tau_1 \to \tau_2
                                                                       \frac{\Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2} \quad \text{TYPING\_APP}
  e \leadsto e'
                                                                                      \frac{e \leadsto e'}{\mathbf{s} \ e \leadsto \mathbf{s} \ e'} \quad \text{EVAL\_S}
```

$$\frac{e_1 \leadsto e_1'}{e_1 \ e_2 \leadsto e_1' \ e_2} \quad \text{EVAL\_APP\_LEFT}$$

$$\frac{e_1 \ \mathbf{val}}{e_1 \ e_2 \leadsto e_2'}$$

$$\frac{e_2 \leadsto e_2'}{e_1 \ e_2 \leadsto e_1 \ e_2'} \quad \text{EVAL\_APP\_RIGHT}$$

$$\frac{e_2 \ \mathbf{val}}{(\lambda(x \colon \tau) e_1) \ e_2 \leadsto e_1 \{ e_2/x \}} \quad \text{EVAL\_BETA}$$

$$\frac{e \leadsto e'}{\mathbf{rec} \ e\{ \mathbf{z} \to e_0; \ \mathbf{s} \ x \to e_1 \} \leadsto \mathbf{rec} \ e'\{ \mathbf{z} \to e_0; \ \mathbf{s} \ x \to e_1 \}} \quad \text{EVAL\_REC\_SCRUT}$$

$$\frac{(\mathbf{s} \ e) \ \mathbf{val}}{\mathbf{rec} \ (\mathbf{s} \ e) \{ \mathbf{z} \to e_0; \ \mathbf{s} \ x \to e_1 \} \leadsto e_0 } \quad \text{EVAL\_REC\_S}$$

Definition rules: 16 good 0 bad Definition rule clauses: 35 good 0 bad