# Pixy formal semantics

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## 1 Term language

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\langle expr \rangle ::= \langle number \rangle
        \langle var \rangle
          \langle bool \rangle
          nil
          ? \langle expr \rangle
      if \langle expr \rangle then \langle expr \rangle else \langle expr \rangle
          \langle expr \rangle fby \langle expr \rangle
          \langle expr \rangle where \langle wheredecls \rangle
          next \langle expr \rangle
          \langle var \rangle ( \langle exprlist \rangle )
          \langle expr \rangle + \langle expr \rangle
           \langle expr \rangle - \langle expr \rangle
          \langle expr \rangle * \langle expr \rangle
          \langle expr \rangle / \langle expr \rangle
          ( \langle exprlist \rangle ( | \langle expr \rangle )? )
\langle exprlist \rangle ::= \langle expr \rangle, \langle exprlist \rangle \mid \langle expr \rangle
\langle varlist \rangle ::= \langle var \rangle , \langle varlist \rangle \mid \langle var \rangle
\langle wheredecl \rangle ::= \langle var \rangle = \langle expr \rangle
          \langle var \rangle ( \langle varlist \rangle ) = \langle expr \rangle
         (\langle varlist \rangle (|\langle var \rangle)?) = \langle expr \rangle
\langle wheredecls \rangle ::= \langle wheredecl \rangle; \langle wheredecls \rangle \mid \langle wheredecl \rangle
\langle bool \rangle ::= \mathtt{true} \mid \mathtt{false}
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### 2 Init rules

$$\frac{F;M \vdash E \overset{init}{\Rightarrow} S}{F;M \vdash ?E \overset{init}{\Rightarrow} S} [\text{Init} - \text{check}]$$

$$F;M \vdash L \overset{init}{\Rightarrow} S_{l}$$

$$F;M \vdash R \overset{init}{\Rightarrow} S_{r}$$

$$F;M \vdash R \overset{init}{\Rightarrow} S_{r}$$

$$F;M \vdash L \text{ fby } R \overset{init}{\Rightarrow} false, S_{l}, S_{r} [\text{Init} - \text{fby}]$$

$$F;M \vdash L \text{ finit} S_{c}$$

$$F;M \vdash T \overset{init}{\Rightarrow} S_{l}$$

$$F;M \vdash F \overset{init}{\Rightarrow} S_{l}$$

$$F;M \vdash F \overset{init}{\Rightarrow} S_{l}$$

$$F;M \vdash B \overset{init}{\Rightarrow} S_{r} [\text{Init} - \text{ite}]$$

$$\frac{F;M \vdash E \overset{init}{\Rightarrow} S}{F;M \vdash \text{next} E \overset{init}{\Rightarrow} false, S} [\text{Init} - \text{next}]$$

$$\frac{F;M;F_{f} \vdash E_{n} \overset{wheredecl}{\Rightarrow} F';M'}{F;M;F_{f} \vdash n = E;E_{n} \overset{wheredecl}{\Rightarrow} F';n \rightarrow (),M'} [\text{WhereInit} - \text{v} - \text{decl}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} M}{n,n_{s} \overset{wheredecltuple}{\Rightarrow} n \rightarrow (),M} [\text{WhereInit} - \text{tuple} - \text{decl} - \text{body}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} n \rightarrow ()}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{tail}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} M}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{empty}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} M}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{empty}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} M}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{empty}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} M_{s}}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{empty}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} M_{s}}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{empty}]$$

$$\frac{n_{s} \overset{wheredecltuple}{\Rightarrow} F';M'}{\Rightarrow} (\text{WhereInit} - \text{tuple} - \text{decl} - \text{empty}]$$

$$F;M;F_f \vdash E_n \overset{where q deed}{\Rightarrow} F';M' \\ \hline F;M;F_f \vdash f(A) = E;E_n \overset{where q deed}{\Rightarrow} f \to \langle A,E,F_f \rangle, F';M' \\ \hline \\ \hline F;M;F_f;M_f \vdash \overset{where q deed}{\Rightarrow} F;M \\ \hline F;M \vdash E \overset{init}{\Rightarrow} s \\ F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash n = E;E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_n \overset{where q init}{\Rightarrow} S \\ \hline F;M \vdash E_s \overset{where q deed}{\Rightarrow} F'' \\ \hline F' = F'' \\ F' : M' \vdash E_s \overset{where q deed}{\Rightarrow} S_e \\ \hline F';M' \vdash E_s \overset{where q deed}{\Rightarrow} S \\ \hline F;M \vdash E \overset{init}{\Rightarrow} S_e \\ \hline F;M \vdash E \overset{init}{\Rightarrow} S_e, S \\ \hline A_n;F;M \vdash E_n \overset{apply init}{\Rightarrow} S_n;M_i \\ \hline A_n;F;M \vdash E_n \overset{apply init}{\Rightarrow} S_n;M_i \\ \hline A_n;F;M \vdash E_n \overset{apply init}{\Rightarrow} S_n;M_i \\ \hline \hline F;M \vdash \overset{apply init}{\Rightarrow} S_n;A \to \langle E_n,M \rangle,M_i \\ \hline \hline F;M \vdash \overset{apply init}{\Rightarrow} (\emptyset;\emptyset) \\ \hline \hline \left. \begin{array}{c} [ApplyInit-empty] \\ \hline \end{array} \right]$$

$$\begin{split} F(f) &= \langle A, E, F_i \rangle \\ A; F; M &\vdash E_a \overset{applyinit}{\Rightarrow} S; M_i \\ \frac{F_i; M_i \vdash E \overset{init}{\Rightarrow} S_e}{F; M \vdash f(E_a) \overset{init}{\Rightarrow} S_e, S} \text{[Init-apply]} \end{split}$$

$$\frac{E \in \{true, false\} \vee E \in \mathbb{R}}{F; M \vdash E \overset{init}{\Rightarrow} \emptyset} [\mathtt{Init-literal}]$$

$$rac{M(E) = \_}{F: M \vdash E \overset{init}{\Rightarrow} \emptyset} [\mathtt{Init} - \mathtt{id}]$$

$$\begin{split} F; M \vdash L &\overset{init}{\Rightarrow} S_l \\ F; M \vdash R &\overset{init}{\Rightarrow} S_r \\ \overline{F; M \vdash L \ \_ R \overset{init}{\Rightarrow} S_l, S_r} [\texttt{Init} - \texttt{binop}] \end{split}$$

$$\begin{split} F; M \vdash E &\overset{init}{\Rightarrow} S \\ \frac{F; M \vdash E_n \overset{inittuple}{\Rightarrow} S_n}{F; M \vdash E, E_n \overset{inittuple}{\Rightarrow} S, S_n} [\texttt{Tuple} - \texttt{state}] \end{split}$$

$$\frac{F; M \vdash E \overset{init}{\Rightarrow} S}{F; M \vdash \mid \ E \overset{inittuple}{\Rightarrow} S} [\texttt{Tuple} - \texttt{state} - \texttt{tail}]$$

$$\frac{}{F;M \vdash^{inittuple} \beta} [\texttt{Tuple} - \texttt{state} - \texttt{empty}]$$

$$\frac{F; M \vdash E_s \overset{inittuple}{\Rightarrow} S}{F; M \vdash (E_s) \overset{init}{\Rightarrow} S} [\texttt{Init} - \texttt{tuple}]$$

## 3 Data type sizing

#### 4 Evaluation rules

$$\begin{split} S; F; M \vdash E \Downarrow V; S' \\ \frac{V \neq nil}{S; F; M \vdash ?E \Downarrow true; S'} [\texttt{Eval} - \texttt{check} - \texttt{true}] \end{split}$$

$$\frac{S;F;M\vdash E\Downarrow nil;S'}{S;F;M\vdash E\Downarrow false;S'}[\texttt{Eval}-\texttt{check}-\texttt{false}]$$

$$S_c;F;M\vdash C\Downarrow true;S'_c$$

$$S_l;F;M\vdash T\Downarrow V;S'_l$$

$$S_l;F;M\vdash T\Downarrow nil;S'_l$$

$$S_c;S_l,S_l;F;M\vdash ifC\, then\, T\, else\, F\Downarrow V;S'_c,S'_l,S'_l\\ \hline S_l;F;M\vdash T\Downarrow nil;S'_l$$

$$S_l;F;M\vdash T\Downarrow nil;S'_l$$

$$S_l;F;M$$

$$\begin{split} &S_l; F; M \vdash E_l \, \P \, nil; S_l' \\ &\frac{S_r; F; M \vdash E_r \, \P \, nil; S_r'}{c, S_l, S_r; F; M \vdash E_l \, \text{fby} \, E_r \, \Downarrow \, nil; c, S_l', S_r'} [\text{Eval} - \text{fby} - \text{C}] \end{split}$$

$$\frac{S_n; F; M; F_f \vdash E_n \overset{names}{\Rightarrow} F'; M'}{\left< v, \_ \right>, S_n; F; M \vdash n = E; E_n \overset{names}{\Rightarrow} F'; n \rightarrow v, M'} [\texttt{WhereNames} - \texttt{v} - \texttt{decl}]$$

$$\frac{S_n; M \vdash n_n \overset{readtuple}{\Rightarrow} M'}{v, S_n; M \vdash n, n_n \overset{readtuple}{\Rightarrow} n \rightarrow v, M'} [\texttt{WhereNames} - \texttt{tuple} - \texttt{decl} - \texttt{read}]$$

$$\frac{}{v_s; M \vdash | \ n \overset{readtuple}{\Rightarrow} n \to (v_s), M} [\texttt{WhereNames} - \texttt{tuple} - \texttt{decl} - \texttt{read} - \texttt{tail}]$$

$$\frac{}{\emptyset ; M \vdash \overset{readtuple}{\Rightarrow} M} [\texttt{WhereNames} - \texttt{tuple} - \texttt{decl} - \texttt{read} - \texttt{empty}]$$

$$S; M \vdash n_s \overset{readtuple}{\Rightarrow} M_s \\ \frac{S_n; F; M; F_f \vdash E_n \overset{names}{\Rightarrow} F'; M'}{\left\langle S, \_ \right\rangle, S_n; F; M \vdash (n_s) = E; E_n \overset{names}{\Rightarrow} F'; M_s, M'} [\texttt{WhereNames} - \texttt{tuple} - \texttt{decl}]$$

$$\frac{S; F; M; F_f \vdash E_v \overset{names}{\Rightarrow} F'; M'}{S; F; M \vdash f(A) = E; E_n \overset{names}{\Rightarrow} f \to \left\langle A, E, F_f \right\rangle, F'; M'} [\texttt{WhereNames} - \texttt{fn} - \texttt{decl}]$$

$$\frac{1}{\emptyset; F; M; F_f \vdash \overset{names}{\Rightarrow} F'; M'} [\texttt{WhereNames} - \texttt{empty}]$$

$$\begin{array}{c} s; F; M \vdash E \Downarrow v; s' \\ S_n; F; M \vdash E_n \overset{values}{\Rightarrow} S_n' \\ \hline \langle \_, s \rangle, S_n; F; M \vdash n = E; E_n \overset{values}{\Rightarrow} \langle v, s' \rangle, S_n' \end{array} [ \texttt{WhereVal} - \texttt{v} - \texttt{decl} ] \end{array}$$

$$\frac{S; F; M \vdash E \Downarrow (V_s); S'}{S_n; F; M \vdash E_n \overset{values}{\Rightarrow} S'_n} [\texttt{WhereVal} - \texttt{tuple} - \texttt{decl}]$$

$$\frac{\langle \_, S \rangle \,, S_n; F; M \vdash (\_) = E; E_n \overset{values}{\Rightarrow} \langle V_s, S' \rangle \,, S'_n}$$

$$\frac{S; F; M \vdash E \Downarrow nil; S'}{\left\langle false, nil \right\rangle, S; F; M \vdash \text{next} E \Downarrow nil; \left\langle false, nil \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{before} - \texttt{nil}]$$

$$\frac{v \neq nil}{S; F; M \vdash E \Downarrow V; S'} \\ \frac{S; F; M \vdash E \Downarrow V; S'}{\left\langle true, v \right\rangle, S; F; M \vdash \text{next} \ E \Downarrow v; \left\langle true, V \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{after}]$$

$$\frac{S; F; M \vdash E \Downarrow V; S'}{\left\langle true, nil \right\rangle, S; F; M \vdash \text{next} \, E \Downarrow V; \left\langle true, nil \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{after} - \texttt{nil}]$$

$$\frac{S; F; M \vdash E \, \P \, nil; S'}{\left\langle c, v \right\rangle, S; F; M \vdash \text{next} \, E \, \P \, nil; \left\langle c, v \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{C}]$$

$$\overline{\emptyset; F; M \vdash nil \Downarrow nil; \emptyset} [\mathtt{Eval} - \mathtt{nil}]$$

$$\frac{}{\emptyset;F;M\vdash nil \, \P\, nil; \emptyset} [\mathtt{Eval} - \mathtt{nil} - \mathtt{C}]$$

$$\frac{N \in \mathbb{R}}{\emptyset; F; M \vdash N \Downarrow N; \emptyset} [\mathtt{Eval} - \mathtt{num}]$$

$$\frac{N \in \mathbb{R}}{\emptyset \colon F \colon M \vdash N \ \P \ nil \colon \emptyset} [\mathtt{Eval} - \mathtt{num} - \mathtt{C}]$$

$$\frac{M(I) = V}{\emptyset \colon\! F \colon\! M \vdash I \Downarrow V \colon\! \emptyset} [\mathtt{Eval} - \mathtt{id}]$$

$$\frac{M(I) = V}{\emptyset \colon F \colon M \vdash I \ \emptyset \ nil \colon \emptyset} [\mathtt{Eval} - \mathtt{id} - \mathtt{C}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \Downarrow E; \emptyset} [\mathtt{Eval} - \mathtt{boolean}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \ \P \ nil; \emptyset} [\mathtt{Eval} - \mathtt{boolean} - \mathtt{C}]$$

$$\frac{S; F; M \vdash E \Downarrow V; S'}{A_n; S_n; F; M \vdash E_n \overset{arg}{\Rightarrow} M_i; S'_n} [\texttt{Apply} - \texttt{arg}]$$

$$\frac{A_n; S, S_n; F; M \vdash E, E_n \overset{arg}{\Rightarrow} A \rightarrow V, M_i; S', S'_n}{A_n; S', S'_n} [\texttt{Apply} - \texttt{arg}]$$

$$\frac{}{\emptyset; \emptyset; F; M \vdash \overset{arg}{\Rightarrow} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty}]$$

$$S; F; M \vdash E \ \ nil; S' \\ \frac{A_n; S_n; F; M \vdash E_n \overset{argC}{\Rightarrow} M_i; S'_n}{A, A_n; S, S_n; F; M \vdash E, E_n \overset{argC}{\Rightarrow} A \rightarrow nil, M_i; S', S'_n} [\texttt{Apply} - \texttt{arg} - \texttt{C}]$$

$$\frac{}{\emptyset; \emptyset; F; M \vdash \overset{argC}{\Rightarrow} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty} - \texttt{C}]$$

$$\begin{split} F(f) &= \langle A, E, F_i \rangle \\ A; S; F; M \vdash a &\stackrel{arg}{\Rightarrow} M_i; S' \\ \frac{S_e; F_i; M_i \vdash E \Downarrow V; S'_e}{S_e, S; F; M \vdash f(a) \Downarrow V; S'_e, S'} [\texttt{Eval} - \texttt{apply}] \end{split}$$

$$\begin{split} F(f) &= \langle A, E, F_i \rangle \\ A; S; F; M \vdash a &\stackrel{argC}{\Rightarrow} M_i; S' \\ \frac{S_e; F_i; M_i \vdash E \ \ nil; S'_e}{S_e, S; F; M \vdash f(a) \ \ \ nil; S'_e, S'} [\texttt{Eval} - \texttt{apply} - \texttt{C}] \end{split}$$

$$\frac{L,R \in \mathbb{R}}{V = L + R \atop L + R \overset{binop}{\Rightarrow} V} [\texttt{Binop-plus}]$$

$$\frac{L,R \in \mathbb{R}}{V = L - R} \underbrace{\frac{V = L - R}{L - R}}_{\text{binop}} \text{[Binop-minus]}$$

$$\frac{L,R \in \mathbb{R}}{V = L*R \atop L*R \overset{binop}{\Rightarrow} V}[\texttt{Binop-times}]$$

$$\frac{L,R \in \mathbb{R}}{V = L/R \atop L/R \overset{binop}{\Rightarrow} V} [\texttt{Binop-divide}]$$

$$\begin{split} S_{l}; F; M \vdash L_{e} \Downarrow L_{v}; S'_{l} \\ S_{r}; F; M \vdash R_{e} \Downarrow R_{v}; S'_{l} \\ \frac{L_{v} \ B \ R_{v} \overset{binop}{\Rightarrow} V}{S_{l}, S_{r}; F; M \vdash L_{e} \ B \ R_{e} \Downarrow V; S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{binop}] \end{split}$$

$$\begin{split} &S_{l}; F; M \vdash L_{e} \ \P \ nil; S'_{l} \\ &\frac{S_{r}; F; M \vdash R_{e} \ \P \ nil; S'_{l}}{S_{l}, S_{r}; F; M \vdash L_{e} \ B \ R_{e} \ \P \ nil; S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{binop} - \texttt{C}] \end{split}$$

$$\begin{aligned} &E_n \overset{len}{\Rightarrow} N \\ &\frac{C = N + 1}{-, E_n \overset{len}{\Rightarrow} C} [\texttt{Len} - \texttt{count}] \end{aligned}$$

$$\mathop{\underset{\Rightarrow}{len}} 0 [\mathtt{Len} - \mathtt{count} - \mathtt{0}]$$

$$\begin{split} S; F; M \vdash E \Downarrow (V_s); S' \\ \frac{V_s \overset{len}{\Rightarrow} N}{S; F; M \vdash \text{len} \, E \Downarrow N; S'} [\texttt{Eval} - \texttt{len}] \end{split}$$

$$\frac{S; F; M \vdash E \, \P \, nil; S'}{S; F; M \vdash \operatorname{len} E \, \P \, nil; S'} [\mathtt{Eval} - \mathtt{len} - \mathtt{C}]$$

$$S; F; M \vdash E \Downarrow V; S'$$

$$S_n; F; M \vdash E_n \overset{tuple cons}{\Rightarrow} V_n; S'_n$$

$$S, S_n; F; M \vdash E, E_n \overset{tuple cons}{\Rightarrow} V, V_n; S', S'_n$$
[Tuple - construct]

$$\frac{}{\emptyset; F; M \vdash \overset{tuplecons}{\Rightarrow} \emptyset; \emptyset} [\texttt{Tuple} - \texttt{construct} - \texttt{empty}]$$

$$\frac{S; F; M \vdash E \overset{tuplecons}{\Rightarrow} (V_s); S'}{S; F; M \vdash \mid E \overset{tuplecons}{\Rightarrow} V_s; S'} [\texttt{Tuple} - \texttt{construct} - \texttt{tail}]$$

$$\frac{S; F; M \vdash E_s \overset{tuplecons}{\Rightarrow} V_s; S'}{S; F; M \vdash (E_s) \Downarrow (V_s); S'} [\texttt{Eval} - \texttt{tuple}]$$

$$S; F; M \vdash E \, \$ \, nil; S' \\ \frac{S_n; F; M \vdash E_n \overset{tupleconsC}{\Longrightarrow} S'_n}{S, S_n; F; M \vdash E, E_n \overset{tupleconsC}{\Longrightarrow} S', S'_n} [\texttt{Tuple} - \texttt{construct} - \texttt{C}]$$

$$\frac{}{\emptyset;F;M \vdash^{tupleconsC} \emptyset} [\texttt{Tuple} - \texttt{construct} - \texttt{empty} - \texttt{C}]$$

$$\frac{S; F; M \vdash E \ \P \ nil; S'}{S; F; M \vdash |E| \overset{tupleconsC}{\Rightarrow} S'} [\texttt{Tuple} - \texttt{construct} - \texttt{tail} - \texttt{C}]$$

$$\frac{S; F; M \vdash E_s \stackrel{tuple consC}{\Rightarrow} S'}{S; F; M \vdash (E_s) \ \$ \ nil; S'} [\texttt{Eval} - \texttt{tuple} - \texttt{C}]$$