Pixy formal semantics

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1 Term language

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\langle expr \rangle ::= \langle literal \rangle
   |\langle var \rangle|
         nil
         ? \langle expr \rangle
         if \langle expr \rangle then \langle expr \rangle else \langle expr \rangle
     \langle expr \rangle fby \langle expr \rangle
         \langle expr \rangle where \langle wheredecls \rangle
         next \langle expr \rangle
          \langle var \rangle ( \langle opargvals \rangle )
          \langle expr \rangle + \langle expr \rangle
          \langle expr \rangle - \langle expr \rangle
          \langle expr \rangle * \langle expr \rangle
          \langle expr \rangle / \langle expr \rangle
\langle opargvals \rangle ::= \langle expr \rangle, \langle opargvals \rangle \mid \langle expr \rangle
\langle opargslist \rangle ::= \langle var \rangle , \langle opargslist \rangle \mid \langle var \rangle
\langle wheredecl \rangle ::= \langle var \rangle = \langle expr \rangle
   |\langle var \rangle| ( \langle opargslist \rangle ) = \langle expr \rangle
\langle wheredecls \rangle ::= \langle wheredecl \rangle; \langle wheredecls \rangle \mid \langle wheredecl \rangle
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2 Init rules

$$\frac{F \vdash E \overset{init}{\Rightarrow} S}{F \vdash ?E \overset{init}{\Rightarrow} S} [\mathtt{Init} - \mathtt{check}]$$

$$\begin{split} F \vdash L & \stackrel{init}{\Rightarrow} S_l \\ F \vdash R & \stackrel{init}{\Rightarrow} S_r \\ \hline F \vdash L \text{ fby } R & \stackrel{init}{\Rightarrow} false, S_l, S_r \end{split} [\text{Init} - \text{fby}] \end{split}$$

$$\begin{split} F \vdash C &\overset{init}{\Rightarrow} S_c \\ F \vdash T &\overset{init}{\Rightarrow} S_t \\ F \vdash F &\overset{init}{\Rightarrow} S_f \\ \hline F \vdash &\text{if } C \text{ then } T \text{ else } F \overset{init}{\Rightarrow} S_c, S_t, S_f \end{split} [\text{Init} - \text{ite}]$$

$$\frac{F \vdash E \overset{init}{\Rightarrow} S}{F \vdash \text{next} \, E \overset{init}{\Rightarrow} false, S} [\texttt{Init} - \texttt{next}]$$

$$\begin{array}{c} F \vdash E \overset{init}{\Rightarrow} s \\ F \vdash E_n \overset{whereinit}{\Rightarrow} S; F' \\ \hline F \vdash n = E; E_n \overset{whereinit}{\Rightarrow} \langle nil, s \rangle \,, S; F' \end{array} [\text{WhereInit} - \text{v} - \text{decl}]$$

$$\frac{F \vdash E_n \overset{where init}{\Rightarrow} S; F'}{F \vdash f(A) = E; E_n \overset{where init}{\Rightarrow} S; f \to \langle A, E \rangle, F'} [\texttt{WhereInit} - \texttt{fn} - \texttt{decl}]$$

$$\frac{}{F \vdash^{whereinit} \emptyset; F} [\texttt{WhereInit} - \texttt{empty}]$$

$$\begin{split} F \vdash E & \stackrel{init}{\Rightarrow} S_e \\ F \vdash E_s & \stackrel{where init}{\Rightarrow} S \\ \hline F \vdash E & \text{where} \ E_s & \stackrel{init}{\Rightarrow} S_e, S \end{split} [\texttt{Init} - \texttt{where}] \end{split}$$

$$\begin{split} F \vdash A & \overset{init}{\Rightarrow} S \\ \frac{F \vdash A_n \overset{applyinit}{\Rightarrow} S_n}{F \vdash A, A_n \overset{applyinit}{\Rightarrow} S, S_n} [\texttt{ApplyInit} - \texttt{arg}] \end{split}$$

$$\frac{}{F \vdash^{applyinit} \underset{}{\Rightarrow} \emptyset} [\texttt{ApplyInit} - \texttt{empty}]$$

$$F(f) = \langle _, E \rangle$$
 $F \vdash E \stackrel{init}{\Rightarrow} S_e$
 $F \vdash A \stackrel{applyinit}{\Rightarrow} S$
 $F \vdash f(A) \stackrel{init}{\Rightarrow} S_e, S$
[Init - apply]
$$F \vdash f(A) \stackrel{init}{\Rightarrow} S_e, S$$

$$F \vdash f(A) \stackrel{init}{\Rightarrow} S_e, S$$
[Init - literal]
$$F \vdash L \stackrel{init}{\Rightarrow} S_l$$

$$F \vdash R \stackrel{init}{\Rightarrow} S_r$$

$$F \vdash L \stackrel{init}{\Rightarrow} S_r$$

$$F \vdash L \stackrel{init}{\Rightarrow} S_r$$
[Init - binop]

3 Evaluation rules

$$\begin{split} S; F; M \vdash E \Downarrow V; S' \\ \frac{V \neq nil}{S; F; M \vdash ?E \Downarrow true; S'} [\texttt{Eval} - \texttt{check} - \texttt{true}] \end{split}$$

$$\frac{S; F; M \vdash E \Downarrow nil; S'}{S; F; M \vdash ?E \Downarrow false; S'} [\mathtt{Eval} - \mathtt{check} - \mathtt{false}]$$

$$S_{c}; F; M \vdash C \Downarrow true; S'_{c}$$

$$S_{t}; F; M \vdash T \Downarrow V; S'_{t}$$

$$S_{f}; F; M \vdash F \nsubseteq nil; S'_{f}$$

$$S_{c}, S_{t}, S_{f}; F; M \vdash \text{if } C \text{ then } T \text{ else } F \Downarrow V; S'_{c}, S'_{t}, S'_{f}$$
[Eval - ite - true]

$$S_{c}; F; M \vdash C \Downarrow false; S'_{c}$$

$$S_{t}; F; M \vdash T \Downarrow nil; S'_{t}$$

$$S_{f}; F; M \vdash F \Downarrow V; S'_{f}$$

$$S_{c}, S_{t}, S_{f}; F; M \vdash \text{if } C \text{ then } T \text{ else } F \Downarrow V; S'_{c}, S'_{t}, S'_{f} \text{ [Eval-ite-false]}$$

$$S_{l}; F; M \vdash E_{l} \Downarrow V; S'_{l}$$

$$V \neq nil$$

$$S_{r}; F; M \vdash E_{r} \nsubseteq nil; S'_{r}$$

$$false, S_{l}, S_{r}; F; M \vdash E_{l} \text{ fby } E_{r} \Downarrow V; true, S'_{l}, S'_{r} \text{ [Eval - fby - before]}$$

$$\frac{S_l; F; M \vdash E_l \Downarrow nil; S_l'}{S_r; F; M \vdash E_r \Downarrow nil; S_r'} \\ \frac{S_r; F; M \vdash E_r \Downarrow nil; S_r'}{false, S_l, S_r; F; M \vdash E_l \text{ fby } E_r \Downarrow nil; false, S_l', S_r'} [\texttt{Eval} - \texttt{fby} - \texttt{before} - \texttt{nil}]$$

$$\frac{S_l; F; M \vdash E_l \, \P \, nil; S_l'}{S_r; F; M \vdash E_r \, \Downarrow V; S_r'} \\ \frac{S_r; F; M \vdash E_r \, \Downarrow V; S_r'}{true, S_l, S_r; F; M \vdash E_l \, \text{fby} \, E_r \, \Downarrow V; true, S_l', S_r'} [\texttt{Eval} - \texttt{fby} - \texttt{after}]$$

$$\begin{split} &S_l; F; M \vdash E_l \, \P \, nil; S_l' \\ &S_r; F; M \vdash E_r \, \P \, nil; S_r' \\ &\overline{c, S_l, S_r; F; M \vdash E_l \, \text{fby} \, E_r \, \Downarrow \, nil; c, S_l', S_r'} [\texttt{Eval} - \texttt{fby} - \texttt{C}] \end{split}$$

$$\frac{S_n; F; M \vdash E_n \overset{names}{\Rightarrow} F'; M'}{\langle v, _ \rangle \,, S_n; F; M \vdash n = E; E_n \overset{names}{\Rightarrow} F'; n \to v, M'} [\texttt{WhereNames} - \texttt{v} - \texttt{decl}]$$

$$\frac{S; f \rightarrow \left\langle A, E \right\rangle, F; M \vdash E_v \overset{names}{\Rightarrow} F'; M'}{S; F; M \vdash f(A) = E; E_n \overset{names}{\Rightarrow} F'; M'} [\texttt{WhereNames} - \texttt{fn} - \texttt{decl}]$$

$$\frac{}{\emptyset;F;M\vdash \overset{names}{\Rightarrow}F';M'}[\texttt{WhereNames}-\texttt{empty}]$$

$$\frac{s; F; M \vdash E \Downarrow v; s'S_n; F; M \vdash E_n \overset{values}{\Rightarrow} S'_n}{\left<_, s\right>, S_n; F; M \vdash n = E; E_n \overset{values}{\Rightarrow} \left< v, s' \right>, S'_n} [\texttt{WhereVal} - \texttt{v} - \texttt{decl}]$$

$$\frac{S; F; M \vdash E_n \overset{volues}{>} S'}{S; F; M \vdash f(A) = E; E_n \overset{volues}{>} S'} [\texttt{WhereVal} - \texttt{fn} - \texttt{decl}]$$

$$\frac{0}{S; F; M \vdash f(A) = E; E_n \overset{volues}{>} S'} [\texttt{WhereVal} - \texttt{empty}]$$

$$\frac{S; F; M \vdash E_s \overset{names}{>} F_i; M_i}{S; F_i; M_i \vdash E \overset{volues}{>} S'} S_e; F_i; M_i \vdash E \overset{volues}{>} S' S_e; F_i; M_i \vdash E \overset{volues}{>} V; S'_e; S'} [\texttt{Eval} - \texttt{where}]$$

$$\frac{S; F; M \vdash E_s \overset{names}{=} F_i; M_i}{S; F_i; M_i \vdash E_s \overset{nales}{=} S'} S_e; F_i; M_i \vdash E_s \overset{volues}{=} S' S_e; F_i; M_i \vdash E \overset{volues}{=} S' S' S_e; F_i; M_i \vdash E \overset$$

$$\overline{\emptyset;F;M\vdash nil \, \P \, nil;\emptyset} [\mathtt{Eval} - \mathtt{nil} - \mathtt{C}]$$

$$\frac{N \in \mathbb{R}}{\emptyset; F; M \vdash N \Downarrow N; \emptyset} [\mathtt{Eval} - \mathtt{num}]$$

$$\frac{N \in \mathbb{R}}{\emptyset; F; M \vdash N \, \emptyset \, nil; \emptyset} [\mathtt{Eval} - \mathtt{num} - \mathtt{C}]$$

$$\frac{M(I) = V}{\emptyset \colon\! F \colon\! M \vdash I \Downarrow V \colon\! \emptyset} [\mathtt{Eval} - \mathtt{id}]$$

$$\frac{M(I) = V}{\emptyset \colon F \colon M \vdash I \, \emptyset \, nil \colon \emptyset} [\mathtt{Eval} - \mathtt{id} - \mathtt{C}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \Downarrow E; \emptyset} [\mathtt{Eval} - \mathtt{boolean}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \ \P \ nil; \emptyset} [\mathtt{Eval} - \mathtt{boolean} - \mathtt{C}]$$

$$S; F; M \vdash E \Downarrow V; S'$$

$$A_n; S_n; F; M \vdash E_n \stackrel{arg}{\Rightarrow} M_i; S'_n$$

$$A, A_n; S, S_n; F; M \vdash E, E_n \stackrel{arg}{\Rightarrow} A \rightarrow V, M_i; S', S'_n$$
[Apply - arg]

$$\frac{}{\emptyset; \emptyset; F; M \vdash \overset{arg}{\Rightarrow} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty}]$$

$$S; F; M \vdash E \ \ nil; S'$$

$$\frac{A_n; S_n; F; M \vdash E_n \overset{argC}{\Rightarrow} M_i; S'_n}{A, A_n; S, S_n; F; M \vdash E, E_n \overset{argC}{\Rightarrow} A \rightarrow nil, M_i; S', S'_n} [\texttt{Apply} - \texttt{arg} - \texttt{C}]$$

$$\frac{}{\emptyset; \emptyset; F; M \vdash \overset{argC}{\Rightarrow} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty} - \texttt{C}]$$

$$\begin{split} F(f) &= \langle A, E \rangle \\ A; S; F; M \vdash a \overset{arg}{\Rightarrow} M_i; S' \\ \frac{S_e; F; M_i \vdash E \Downarrow V; S'_e}{S_e, S; F; M \vdash f(a) \Downarrow V; S'_e, S'} [\texttt{Eval} - \texttt{apply}] \end{split}$$

$$\begin{split} F(f) &= \langle A, E \rangle \\ A; S; F; M \vdash a &\stackrel{argC}{\Rightarrow} M_i; S' \\ \frac{S_e; F; M_i \vdash E \ \mbox{\ofinite} \ nil; S'_e}{S_e, S; F; M \vdash f(a) \ \mbox{\ofinite} \ nil; S'_e, S'} [\mbox{Eval} - \mbox{apply} - \mbox{C}] \end{split}$$

$$\begin{split} &L,R \in \mathbb{R} \\ &\frac{V = L + R}{L + R} \underbrace{\text{Binop} - \text{plus}}_{} \end{split}$$

$$\frac{L,R \in \mathbb{R}}{V = L - R} \\ \frac{V = L - R}{L - R} \overset{binop}{\Rightarrow} V \\ [\texttt{Binop-minus}]$$

$$\frac{L,R \in \mathbb{R}}{V = L*R \atop L*R \overset{binop}{\Rightarrow} V}[\texttt{Binop-times}]$$

$$\frac{L,R \in \mathbb{R}}{V = L/R \atop L/R \overset{binop}{\Rightarrow} V} [\texttt{Binop-divide}]$$

$$S_{l}; F; M \vdash L_{e} \Downarrow L_{v}; S'_{l}$$

$$S_{r}; F; M \vdash R_{e} \Downarrow R_{v}; S'_{l}$$

$$\frac{L_{v} B R_{v} \overset{binop}{\Rightarrow} V}{S_{l}, S_{r}; F; M \vdash L_{e} B R_{e} \Downarrow V; S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{binop}]$$

$$\begin{split} &S_l; F; M \vdash L_e \ \$ \ nil; S_l' \\ &S_r; F; M \vdash R_e \ \$ \ nil; S_l' \\ &\overline{S_l, S_r; F; M \vdash L_e \ B \ R_e \ \$ \ nil; S_l', S_r'} [\texttt{Eval} - \texttt{binop} - \texttt{C}] \end{split}$$