Pixy formal semantics

A. Finn Hackett, Reed Mullanix March 19, 2018

1 Term language

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\langle expr \rangle ::= \langle literal \rangle
   |\langle var \rangle|
         nil
         ? \langle expr \rangle
         if \langle expr \rangle then \langle expr \rangle else \langle expr \rangle
     \langle expr \rangle fby \langle expr \rangle
         \langle expr \rangle where \langle wheredecls \rangle
         next \langle expr \rangle
          \langle var \rangle ( \langle opargvals \rangle )
          \langle expr \rangle + \langle expr \rangle
          \langle expr \rangle - \langle expr \rangle
          \langle expr \rangle * \langle expr \rangle
          \langle expr \rangle / \langle expr \rangle
\langle opargvals \rangle ::= \langle expr \rangle, \langle opargvals \rangle \mid \langle expr \rangle
\langle opargslist \rangle ::= \langle var \rangle , \langle opargslist \rangle \mid \langle var \rangle
\langle wheredecl \rangle ::= \langle var \rangle = \langle expr \rangle
   |\langle var \rangle| ( \langle opargslist \rangle ) = \langle expr \rangle
\langle wheredecls \rangle ::= \langle wheredecl \rangle; \langle wheredecls \rangle \mid \langle wheredecl \rangle
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2 Init rules

$$\frac{F \vdash E \overset{init}{\Rightarrow} S}{F \vdash ?E \overset{init}{\Rightarrow} S} [\mathtt{Init} - \mathtt{check}]$$

$$\begin{split} F \vdash L & \stackrel{init}{\Rightarrow} S_l \\ F \vdash R & \stackrel{init}{\Rightarrow} S_r \\ \hline F \vdash L \text{ fby } R & \stackrel{init}{\Rightarrow} false, S_l, S_r \end{split} [\text{Init} - \text{fby}] \end{split}$$

$$\begin{split} F \vdash C &\overset{init}{\Rightarrow} S_c \\ F \vdash T &\overset{init}{\Rightarrow} S_t \\ F \vdash F &\overset{init}{\Rightarrow} S_f \\ \hline F \vdash &\text{if } C \text{ then } T \text{ else } F \overset{init}{\Rightarrow} S_c, S_t, S_f \end{split} [\text{Init} - \text{ite}]$$

$$\frac{F \vdash E \overset{init}{\Rightarrow} S}{F \vdash \text{next} \, E \overset{init}{\Rightarrow} false, S} [\texttt{Init} - \texttt{next}]$$

$$\begin{array}{c} F \vdash E \overset{init}{\Rightarrow} s \\ F \vdash E_n \overset{whereinit}{\Rightarrow} S; F' \\ \hline F \vdash n = E; E_n \overset{whereinit}{\Rightarrow} \langle nil, s \rangle \,, S; F' \end{array} [\text{WhereInit} - \text{v} - \text{decl}]$$

$$\frac{F \vdash E_n \overset{where init}{\Rightarrow} S; F'}{F \vdash f(A) = E; E_n \overset{where init}{\Rightarrow} S; f \to \langle A, E \rangle, F'} [\texttt{WhereInit} - \texttt{fn} - \texttt{decl}]$$

$$\frac{}{F \vdash^{whereinit} \emptyset; F} [\texttt{WhereInit} - \texttt{empty}]$$

$$\begin{split} F \vdash E & \stackrel{init}{\Rightarrow} S_e \\ F \vdash E_s & \stackrel{where init}{\Rightarrow} S \\ \hline F \vdash E & \text{where} \ E_s & \stackrel{init}{\Rightarrow} S_e, S \end{split} [\texttt{Init} - \texttt{where}] \end{split}$$

$$\begin{split} F \vdash A & \overset{init}{\Rightarrow} S \\ \frac{F \vdash A_n \overset{applyinit}{\Rightarrow} S_n}{F \vdash A, A_n \overset{applyinit}{\Rightarrow} S, S_n} [\texttt{ApplyInit} - \texttt{arg}] \end{split}$$

$$\frac{}{F \vdash^{applyinit} \underset{}{\Rightarrow} \emptyset} [\texttt{ApplyInit} - \texttt{empty}]$$

$$F(f) = \langle _, E \rangle$$

$$F \vdash E \stackrel{init}{\Rightarrow} S_e$$

$$\frac{F \vdash A \stackrel{applyinit}{\Rightarrow} S}{F \vdash f(A) \stackrel{init}{\Rightarrow} S_e, S} [\texttt{Init-apply}]$$

$$\frac{F \vdash L \stackrel{init}{\Rightarrow} \emptyset}{F \vdash L \stackrel{init}{\Rightarrow} S_l} [\texttt{Init-literal}]$$

$$F \vdash L \stackrel{init}{\Rightarrow} S_l$$

$$\frac{F \vdash R \stackrel{init}{\Rightarrow} S_r}{F \vdash L \stackrel{init}{\Rightarrow} S_l, S_r} [\texttt{Init-binop}]$$

3 Evaluation rules

$$S; F; M \vdash E \Downarrow V; S' \\ V \neq nil \\ S; F; M \vdash ?E \Downarrow true; S' [\texttt{Eval} - \texttt{check} - \texttt{true}]$$

$$\frac{S; F; M \vdash E \Downarrow nil; S'}{S; F; M \vdash ?E \Downarrow false; S'} [\texttt{Eval} - \texttt{check} - \texttt{false}]$$

$$S_c; F; M \vdash C \Downarrow true; S'_c \\ S_t; F; M \vdash T \Downarrow V; S'_t \\ S_f; F; M \vdash F \Downarrow nil; S'_f \\ \hline S_c, S_t, S_f; F; M \vdash \text{if } C \text{ then } T \text{ else } F \Downarrow V; S'_c, S'_t, S'_f [\texttt{Eval} - \texttt{ite} - \texttt{true}]$$

$$S_c; F; M \vdash C \Downarrow false; S'_c \\ S_t; F; M \vdash T \Downarrow nil; S'_t \\ S_f; F; M \vdash F \Downarrow V; S'_f \\ \hline S_c, S_t, S_f; F; M \vdash \text{if } C \text{ then } T \text{ else } F \Downarrow V; S'_c, S'_t, S'_f [\texttt{Eval} - \texttt{ite} - \texttt{false}]$$

$$S_c; F; M \vdash C \Downarrow nil; S'_c \\ S_t; F; M \vdash T \Downarrow nil; S'_t \\ S_f; F; M \vdash T \Downarrow nil; S'_f \\ \hline S_c, S_t, S_f; F; M \vdash \text{if } C \text{ then } T \text{ else } F \Downarrow nil; S'_c, S'_t, S'_f [\texttt{Eval} - \texttt{ite} - \texttt{nil}]$$

$$\begin{split} S_c; F; M \vdash C & \ \ \downarrow^C nil; S'_c \\ S_t; F; M \vdash T & \ \ \downarrow^C nil; S'_t \\ \frac{S_f; F; M \vdash F & \ \ \downarrow^C nil; S'_f }{S_c, S_t, S_f; F; M \vdash \text{if } C \text{ then } T \text{ else } F & \ \ \downarrow^C nil; S'_c, S'_t, S'_f \end{split} \text{[Eval-ite-C]}$$

$$S_{l}; F; M \vdash E_{l} \Downarrow V; S'_{l}$$

$$V \neq nil$$

$$S_{r}; F; M \vdash E_{r} \Downarrow nil; S'_{r}$$

$$false, S_{l}, S_{r}; F; M \vdash E_{l} \text{ fby } E_{r} \Downarrow V; true, S'_{l}, S'_{r} \text{[Eval - fby - before]}$$

$$S_{l}; F; M \vdash E_{l} \Downarrow nil; S'_{l}$$

$$S_{r}; F; M \vdash E_{r} \Downarrow nil; S'_{r}$$

$$\underbrace{S_{r}; F; M \vdash E_{l} \Downarrow nil; S'_{r}}_{false, S_{l}, S_{r}; F; M \vdash E_{l} \text{ fby } E_{r} \Downarrow nil; false, S'_{l}, S'_{r}}_{[Eval - fby - before - nil]}$$

$$\frac{S_{l}; F; M \vdash E_{l} \stackrel{C}{\Downarrow} nil; S'_{l}}{S_{r}; F; M \vdash E_{r} \Downarrow V; S'_{r}} [\texttt{Eval} - \texttt{fby} - \texttt{after}]$$

$$\frac{true, S_{l}, S_{r}; F; M \vdash E_{l} \, \texttt{fby} \, E_{r} \Downarrow V; true, S'_{l}, S'_{r}}{true, S_{l}, S_{r}; F; M \vdash E_{l} \, \texttt{fby} \, E_{r} \Downarrow V; true, S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{fby} - \texttt{after}]$$

$$S_{l}; F; M \vdash E_{l} \buildrel \bui$$

$$\frac{S_n; F; M \vdash E_n \overset{names}{\Rightarrow} F'; M'}{\langle v, _ \rangle, S_n; F; M \vdash n = E; E_n \overset{names}{\Rightarrow} F'; n \rightarrow v, M'} [\texttt{WhereNames} - \texttt{v} - \texttt{decl}]$$

$$\frac{S;f\rightarrow\left\langle A,E\right\rangle ,F;M\vdash E_{v}\overset{names}{\Rightarrow}F';M'}{S;F;M\vdash f(A)=E;E_{n}\overset{names}{\Rightarrow}F';M'}[\texttt{WhereNames}-\texttt{fn}-\texttt{decl}]$$

$$\frac{1}{\emptyset: F: M \vdash \overset{names}{\Rightarrow} F': M'} [\texttt{WhereNames} - \texttt{empty}]$$

$$\frac{s; F; M \vdash E \Downarrow v; s'S_n; F; M \vdash E_n \overset{values}{\Rightarrow} S'_n}{\langle _, s \rangle \,, S_n; F; M \vdash n = E; E_n \overset{values}{\Rightarrow} \langle v, s' \rangle \,, S'_n} [\texttt{WhereVal} - \texttt{v} - \texttt{decl}]$$

$$\frac{S; F; M \vdash E_n \overset{values}{\Rightarrow} S'}{S; F; M \vdash f(A) = E; E_n \overset{values}{\Rightarrow} S'} [\texttt{WhereVal} - \texttt{fn} - \texttt{decl}]$$

$$\frac{}{\emptyset;F;M\vdash \overset{values}{\Rightarrow}\emptyset}[\mathtt{WhereVal}-\mathtt{empty}]$$

$$\begin{split} S; F; M \vdash E_s \overset{names}{\Rightarrow} F_i; M_i \\ S; F_i; M_i \vdash E_s \overset{values}{\Rightarrow} S' \\ S_e; F_i; M_i \vdash E \Downarrow V; S'_e \\ \hline S_e, S; F; M \vdash E \text{ where } E_s \Downarrow V; S'_e, S' \end{split} [\texttt{Eval} - \texttt{where}]$$

$$\begin{split} S; F; M \vdash E_s \overset{names}{\Rightarrow} F_i; M_i \\ S; F_i; M_i \vdash E_s \overset{values}{\Rightarrow} S' \\ \frac{S_e; F_i; M_i \vdash E \ \psi \ nil; S'_e}{S_e, S; F; M \vdash E \ \text{where} \ E_s \overline{C} \psi nil; S'_e, S'} [\texttt{Eval} - \texttt{where} - \texttt{C}] \end{split}$$

$$S; F; M \vdash E \Downarrow V; S' \\ \frac{V \neq nil}{\left\langle false, nil \right\rangle, S; F; M \vdash \text{next}(E) \Downarrow nil; \left\langle true, V \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{before}]$$

$$\frac{S; F; M \vdash E \Downarrow nil; S'}{\left\langle false, nil \right\rangle, S; F; M \vdash \operatorname{next}(E) \Downarrow nil; \left\langle false, nil \right\rangle, S'} [\mathtt{Eval} - \mathtt{next} - \mathtt{before} - \mathtt{nil}]$$

$$\frac{v \neq nil}{S; F; M \vdash E \Downarrow V; S'} \\ \frac{S; F; M \vdash E \Downarrow V; S'}{\left\langle true, v \right\rangle, S; F; M \vdash \operatorname{next}(E) \Downarrow v; \left\langle true, V \right\rangle, S'} [\mathtt{Eval} - \mathtt{next} - \mathtt{after}]$$

$$\frac{S; F; M \vdash E \Downarrow V; S'}{\left\langle true, nil \right\rangle, S; F; M \vdash \operatorname{next}(E) \Downarrow V; \left\langle true, nil \right\rangle, S'} [\mathtt{Eval} - \mathtt{next} - \mathtt{after} - \mathtt{nil}]$$

$$\frac{S; F; M \vdash E \stackrel{C}{\Downarrow} nil; S'}{\left\langle c, v \right\rangle, S; F; M \vdash \operatorname{next}(E) \stackrel{C}{\Downarrow} nil; \left\langle c, v \right\rangle, S'} [\mathtt{Eval} - \mathtt{next} - \mathtt{C}]$$

$$\overline{\emptyset; F; M \vdash nil \Downarrow nil; \emptyset}$$
 [Eval - nil]

$$\frac{C}{\emptyset; F; M \vdash nil \stackrel{C}{\Downarrow} nil; \emptyset} [\mathtt{Eval} - \mathtt{nil} - \mathtt{C}]$$

$$\frac{N \in \mathbb{R}}{\emptyset; F; M \vdash N \Downarrow N; \emptyset} [\mathtt{Eval} - \mathtt{num}]$$

$$\frac{N \in \mathbb{R}}{\emptyset; F; M \vdash N \Downarrow nil; \emptyset} [\mathtt{Eval} - \mathtt{num} - \mathtt{C}]$$

$$\frac{M(I) = V}{\emptyset; F; M \vdash I \Downarrow V; \emptyset} [\mathtt{Eval} - \mathtt{id}]$$

$$\frac{M(I) = V}{\emptyset; F; M \vdash I \Downarrow nil; \emptyset} [\mathtt{Eval} - \mathtt{id} - \mathtt{C}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \Downarrow E; \emptyset} [\mathtt{Eval} - \mathtt{boolean}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \Downarrow nil; \emptyset} [\mathtt{Eval} - \mathtt{boolean} - \mathtt{C}]$$

$$S; F; M \vdash E \Downarrow V; S'$$

$$A_n; S_n; F; M \vdash E_n \overset{arg}{\Rightarrow} M_i; S'_n$$

$$A, A_n; S, S_n; F; M \vdash E, E_n \overset{arg}{\Rightarrow} A \rightarrow V, M_i; S', S'_n$$
[Apply - arg]

$$\frac{1}{\emptyset : \emptyset : F : M \vdash^{arg}_{\Rightarrow} \emptyset : \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty}]$$

$$S; F; M \vdash E \overset{C}{\Downarrow} nil; S' \\ \frac{A_n; S_n; F; M \vdash E_n \overset{argC}{\Rightarrow} M_i; S'_n}{A, A_n; S, S_n; F; M \vdash E, E_n \overset{argC}{\Rightarrow} A \rightarrow nil, M_i; S', S'_n} [\texttt{Apply} - \texttt{arg} - \texttt{C}]$$

$$\frac{}{\emptyset; \emptyset; F; M \vdash \overset{argC}{\Rightarrow} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty} - \texttt{C}]$$

$$\begin{split} F(f) &= \langle A, E \rangle \\ A; S; F; M \vdash a \overset{arg}{\Rightarrow} M_i; S' \\ \frac{S_e; F; M_i \vdash E \Downarrow V; S'_e}{S_e, S; F; M \vdash f(a) \Downarrow V; S'_e, S'} [\texttt{Eval} - \texttt{apply}] \end{split}$$

$$\begin{split} F(f) &= \langle A, E \rangle \\ A; S; F; M \vdash a &\overset{argC}{\Rightarrow} M_i; S' \\ \frac{S_e; F; M_i \vdash E \ \Downarrow nil; S'_e}{S_e, S; F; M \vdash f(a) \ \Downarrow nil; S'_e, S'} [\texttt{Eval} - \texttt{apply} - \texttt{C}] \end{split}$$

$$\begin{split} &L,R \in \mathbb{R} \\ &\frac{V = L + R}{L + R} \underbrace{[\texttt{Binop} - \texttt{plus}]} \\ &L + R \overset{binop}{\Rightarrow} V \end{split}$$

$$\begin{split} &L,R \in \mathbb{R} \\ &\frac{V = L - R}{L - R} \underbrace{\begin{bmatrix} \text{Binop} - \text{minus} \end{bmatrix}} \\ &\stackrel{binop}{\Rightarrow} V \end{split}$$

$$\frac{L,R \in \mathbb{R}}{V = L/R \atop L/R \overset{binop}{\Rightarrow} V} [\texttt{Binop-divide}]$$

$$\begin{split} S_{l}; F; M \vdash L_{e} \Downarrow L_{v}; S'_{l} \\ S_{r}; F; M \vdash R_{e} \Downarrow R_{v}; S'_{l} \\ \frac{L_{v} \ B \ R_{v} \overset{binop}{\Rightarrow} V}{S_{l}, S_{r}; F; M \vdash L_{e} \ B \ R_{e} \Downarrow V; S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{binop}] \end{split}$$

$$\begin{split} S_{l}; F; M \vdash L_{e} & \underset{C}{\Downarrow} nil; S'_{l} \\ \frac{S_{r}; F; M \vdash R_{e} & \underset{C}{\Downarrow} nil; S'_{l}}{S_{l}, S_{r}; F; M \vdash L_{e} & B & R_{e} & \underset{C}{\Downarrow} nil; S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{binop} - \texttt{C}] \end{split}$$