Pixy formal semantics

A. Finn Hackett, Reed Mullanix March 19, 2018

1 Term language

```
\langle expr \rangle ::= \langle number \rangle
           \langle var \rangle
           \langle bool \rangle
          nil
          ? \langle expr \rangle
          if \langle expr \rangle then \langle expr \rangle else \langle expr \rangle
           \langle expr \rangle fby \langle expr \rangle
           \langle expr \rangle where \langle wheredecls \rangle
          next \langle expr \rangle
           \langle var \rangle ( \langle exprlist \rangle )
           \langle expr \rangle + \langle expr \rangle
           \langle expr \rangle - \langle expr \rangle
           \langle expr \rangle * \langle expr \rangle
           \langle expr \rangle / \langle expr \rangle
           \langle expr \rangle \dots \langle expr \rangle
           ( \langle exprlist \rangle )
           len \langle expr \rangle
           [ \langle exprlist \rangle ] \langle expr \rangle
\langle exprlist \rangle ::= \langle expr \rangle , \langle exprlist \rangle \mid \langle expr \rangle
\langle \mathit{varlist} \rangle ::= \langle \mathit{var} \rangle , \langle \mathit{varlist} \rangle \mid \langle \mathit{var} \rangle
\langle wheredecl \rangle ::= \langle var \rangle = \langle expr \rangle
   |\langle var \rangle| ( \langle varlist \rangle ) = \langle expr \rangle
   | (\langle varlist \rangle) = \langle expr \rangle
\langle wheredecls \rangle ::= \langle wheredecl \rangle; \langle wheredecls \rangle \mid \langle wheredecl \rangle
\langle bool \rangle ::= true \mid false
```

2 Init rules

$$\frac{F;M\vdash E\overset{init}{\Rightarrow}S}{F;M\vdash R\overset{init}{\Rightarrow}S}[\text{Init-check}]$$

$$\frac{F;M\vdash L\overset{init}{\Rightarrow}S_l}{F;M\vdash R\overset{init}{\Rightarrow}S_r}$$

$$\frac{F;M\vdash L\overset{init}{\Rightarrow}S_r}{F;M\vdash L\text{ fby }R\overset{init}{\Rightarrow}false,S_l,S_r}[\text{Init-fby}]$$

$$\frac{F;M\vdash C\overset{init}{\Rightarrow}S_c}{F;M\vdash T\overset{init}{\Rightarrow}S_t}$$

$$\frac{F;M\vdash E\overset{init}{\Rightarrow}S_f}{F;M\vdash F\overset{init}{\Rightarrow}S_c,S_t,S_f}[\text{Init-ite}]$$

$$\frac{F;M\vdash E\overset{init}{\Rightarrow}S}{F;M\vdash E\overset{init}{\Rightarrow}S_c,S_t,S_f}[\text{Init-next}]$$

$$\frac{F;M\vdash E\overset{init}{\Rightarrow}S}{F;M;F_f;M_f\vdash E_n} \overset{wheredeel}{\Rightarrow}F';M'$$

$$\frac{F;M;F_f;M_f\vdash E_n\overset{wheredeel}{\Rightarrow}F';M'}{F;M;F_f;M_f\vdash E_n\overset{wheredeel}{\Rightarrow}F';M'}[\text{WhereInit-fn-decl}]$$

$$\frac{F;M;F_f;M_f\vdash E_n\overset{wheredeel}{\Rightarrow}F';M'}{F;M;F_f;M_f\vdash E_n\overset{wheredeel}{\Rightarrow}F';M'}[\text{WhereInit-empty-decl}]$$

$$\frac{F;M;F_f;M_f\vdash E_n\overset{wheredeel}{\Rightarrow}F;M}{F;M\vdash E_n\overset{wheredeel}{\Rightarrow}S}$$

$$\frac{F;M\vdash E_n\overset{init}{\Rightarrow}S}{F;M\vdash E_n\overset{wheredeel}{\Rightarrow}S}$$

$$\frac{F;M\vdash E_n\overset{wheredeel}{\Rightarrow}S}{F;M\vdash R_n \in E;E_n\overset{wheredeel}{\Rightarrow}S}}[\text{WhereInit-v-init}]$$

$$\frac{F;M\vdash E_n\overset{wheredeel}{\Rightarrow}S}{F;M\vdash R_n \in E;E_n\overset{wheredeel}{\Rightarrow}S}}[\text{WhereInit-v-init}]$$

$$\frac{}{F;M \vdash^{whereinit} \emptyset} [\texttt{WhereInit} - \texttt{empty} - \texttt{init}]$$

$$\begin{split} F; M; F'; M' \vdash E_s & \overset{wheredecl}{\Rightarrow} F''; M'' \\ F' &= F'' \\ M' &= M'' \\ F'; M' \vdash E & \overset{init}{\Rightarrow} S_e \\ \frac{F'; M' \vdash E_s & \overset{whereinit}{\Rightarrow} S}{F; M \vdash E \text{ where } E_s & \overset{init}{\Rightarrow} S_e, S} \end{split} [\texttt{Init-where}]$$

$$F; M \vdash E_a \overset{init}{\Rightarrow} S$$

$$\underbrace{A_n; F; M \vdash E_{a_n} \overset{applyinit}{\Rightarrow} S_n; M_i}_{A, A_n; F; M \vdash E_a, E_{a_n} \overset{applyinit}{\Rightarrow} S, S_n; A \rightarrow \langle E_a, M \rangle, M_i}_{\text{[ApplyInit - arg]}}$$

$$\frac{}{F;M \vdash \overset{applyinit}{\Rightarrow} \emptyset;\emptyset} [\texttt{ApplyInit} - \texttt{empty}]$$

$$F(f) = \langle A, E, F_i \rangle$$

$$A; F; M \vdash E_a \stackrel{applyinit}{\Rightarrow} S; M_i$$

$$\frac{F_i; M_i \vdash E \stackrel{init}{\Rightarrow} S_e}{F; M \vdash f(E_a) \stackrel{init}{\Rightarrow} S_e, S} [\texttt{Init} - \texttt{apply}]$$

$$\frac{E \in \{true, false\} \vee E \in \mathbb{R}}{F; M \vdash E \overset{init}{\Rightarrow} \emptyset} [\mathtt{Init-literal}]$$

$$\frac{M(E) = _}{F; M \vdash E \overset{init}{\Rightarrow} \emptyset} [\texttt{Init} - \texttt{id}]$$

$$\begin{split} F; M \vdash L \overset{init}{\Rightarrow} S_l \\ F; M \vdash R \overset{init}{\Rightarrow} S_r \\ F; M \vdash L &_ R \overset{init}{\Rightarrow} S_l, S_r \end{split} [\texttt{Init-binop}]$$

3 Data type sizing

4 Evaluation rules

$$S_{l}; F; M \vdash E_{l} \ \, \forall v; S'_{r} \\ \hline true, S_{l}, S_{r}; F; M \vdash E_{r} \ \, \forall v; S'_{r} \\ \hline true, S_{l}, S_{r}; F; M \vdash E_{l} \ \, \text{fiby} \ \, E_{r} \ \, \forall v; true, S'_{l}, S'_{r} \\ \hline S_{l}; F; M \vdash E_{l} \ \, \text{fiby} \ \, E_{r} \ \, \forall nil; S'_{r} \\ \hline S_{r}; F; M \vdash E_{l} \ \, \text{fiby} \ \, E_{r} \ \, \forall nil; S'_{r} \\ \hline c, S_{l}, S_{r}; F; M \vdash E_{l} \ \, \text{fiby} \ \, E_{r} \ \, \forall nil; S'_{r} \\ \hline (v, _) \ \, , S_{n}; F; M \vdash E_{l} \ \, \text{fiby} \ \, E_{r} \ \, \forall nil; c, S'_{l}, S'_{r} \\ \hline (v, _) \ \, , S_{n}; F; M \vdash E_{n} \ \, \overset{names}{\Longrightarrow} F'; M' \\ \hline S_{l}; F; M \vdash H \ \, n = E; E_{n} \ \, \overset{names}{\Longrightarrow} F'; M' \\ \hline S_{l}; F; M \vdash H \ \, l = E; E_{n} \ \, \overset{names}{\Longrightarrow} F'; M' \\ \hline S_{l}; F; M \vdash H \ \, l = E; E_{n} \ \, \overset{names}{\Longrightarrow} F'; M' \\ \hline (where Names - empty) \\ \hline \frac{s; F; M \vdash E \ \, \forall v; s'S_{n}; F; M \vdash E_{n} \ \, \overset{names}{\Longrightarrow} S'_{n} \\ \hline (where Names - empty) \\ \hline \frac{s; F; M \vdash E \ \, \forall v; s'S_{n}; F; M \vdash E_{n} \ \, \overset{names}{\Longrightarrow} S'_{n} \\ \hline (where Val - v - decl) \\ \hline \frac{s; F; M \vdash E \ \, \forall v; s'S_{n}; F; M \vdash E_{n} \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash H \ \, \vdots \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \ \, \overset{values}{\Longrightarrow} S' \\ \hline S_{l}; F; M \vdash E \$$

$$\frac{S; F; M \vdash E \Downarrow V; S'}{V \neq nil} \frac{V \neq nil}{\left\langle false, nil \right\rangle, S; F; M \vdash \text{next} \, E \Downarrow nil; \left\langle true, V \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{before}]$$

$$\frac{S; F; M \vdash E \Downarrow nil; S'}{\left\langle false, nil \right\rangle, S; F; M \vdash \text{next} E \Downarrow nil; \left\langle false, nil \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{before} - \texttt{nil}]$$

$$\frac{v \neq nil}{S; F; M \vdash E \Downarrow V; S'} \\ \frac{S; F; M \vdash E \Downarrow V; S'}{\left\langle true, v \right\rangle, S; F; M \vdash \text{next} E \Downarrow v; \left\langle true, V \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{after}]$$

$$\frac{S; F; M \vdash E \Downarrow V; S'}{\left\langle true, nil \right\rangle, S; F; M \vdash \text{next} E \Downarrow V; \left\langle true, nil \right\rangle, S'} [\texttt{Eval} - \texttt{next} - \texttt{after} - \texttt{nil}]$$

$$\frac{S; F; M \vdash E \, \P \, nil; S'}{\langle c, v \rangle \, , S; F; M \vdash \text{next} \, E \, \P \, nil; \langle c, v \rangle \, , S'} [\texttt{Eval} - \texttt{next} - \texttt{C}]$$

$$\overline{\emptyset; F; M \vdash nil \Downarrow nil; \emptyset}$$
 [Eval - nil]

$$\overline{\emptyset; F; M \vdash nil \, \P \, nil; \emptyset} [\mathtt{Eval} - \mathtt{nil} - \mathtt{C}]$$

$$\frac{N \in \mathbb{R}}{\emptyset : F : M \vdash N \Downarrow N : \emptyset} [\mathtt{Eval} - \mathtt{num}]$$

$$\frac{N \in \mathbb{R}}{\emptyset; F; M \vdash N \, \P \, nil; \emptyset} [\mathtt{Eval} - \mathtt{num} - \mathtt{C}]$$

$$\frac{M(I) = V}{\emptyset; F; M \vdash I \Downarrow V; \emptyset} [\mathtt{Eval} - \mathtt{id}]$$

$$\frac{M(I) = V}{\emptyset; F; M \vdash I \, \emptyset \, nil; \emptyset} [\mathtt{Eval} - \mathtt{id} - \mathtt{C}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \Downarrow E; \emptyset} [\mathtt{Eval} - \mathtt{boolean}]$$

$$\frac{E \in \{true, false\}}{\emptyset; F; M \vdash E \ \P \ nil; \emptyset} [\mathtt{Eval} - \mathtt{boolean} - \mathtt{C}]$$

$$S; F; M \vdash E \Downarrow V; S'$$

$$A_n; S_n; F; M \vdash E_n \overset{arg}{\Rightarrow} M_i; S'_n$$

$$A, A_n; S, S_n; F; M \vdash E, E_n \overset{arg}{\Rightarrow} A \rightarrow V, M_i; S', S'_n$$
[Apply - arg]

$$\frac{}{\emptyset; \emptyset; F; M \overset{arg}{\mapsto} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty}]$$

$$S; F; M \vdash E \ \ nil; S'$$

$$\frac{A_n; S_n; F; M \vdash E_n \overset{argC}{\Rightarrow} M_i; S'_n}{A, A_n; S, S_n; F; M \vdash E, E_n \overset{argC}{\Rightarrow} A \rightarrow nil, M_i; S', S'_n} [\texttt{Apply} - \texttt{arg} - \texttt{C}]$$

$$\frac{}{\emptyset; \emptyset; F; M \vdash \overset{argC}{\Rightarrow} \emptyset; \emptyset} [\texttt{Apply} - \texttt{arg} - \texttt{empty} - \texttt{C}]$$

$$F(f) = \langle A, E, F_i
angle \ A; S; F; M \vdash a \overset{arg}{\Rightarrow} M_i; S' \ \frac{S_e; F_i; M_i \vdash E \Downarrow V; S'_e}{S_e, S; F; M \vdash f(a) \Downarrow V; S'_o, S'} [ext{Eval} - ext{apply}]$$

$$\begin{split} F(f) &= \langle A, E, F_i \rangle \\ A; S; F; M \vdash a &\stackrel{argC}{\Rightarrow} M_i; S' \\ \frac{S_e; F_i; M_i \vdash E \ \P \ nil; S'_e}{S_e, S; F; M \vdash f(a) \ \P \ nil; S'_e, S'} [\texttt{Eval} - \texttt{apply} - \texttt{C}] \end{split}$$

$$L,R \in \mathbb{R} \ rac{V = L + R}{L + R \overset{binop}{\Rightarrow} V} [\mathtt{Binop-plus}]$$

$$\frac{L,R \in \mathbb{R}}{\frac{V = L - R}{L - R} \overset{binop}{\Rightarrow} V}[\mathtt{Binop-minus}]$$

$$\frac{L,R \in \mathbb{R}}{V = L/R \atop L/R \overset{binop}{\Rightarrow} V} [\texttt{Binop-divide}]$$

$$\begin{split} S_{l}; F; M \vdash L_{e} \Downarrow L_{v}; S'_{l} \\ S_{r}; F; M \vdash R_{e} \Downarrow R_{v}; S'_{l} \\ \frac{L_{v} \ B \ R_{v} \overset{binop}{\Rightarrow} V}{S_{l}, S_{r}; F; M \vdash L_{e} \ B \ R_{e} \Downarrow V; S'_{l}, S'_{r}} [\texttt{Eval} - \texttt{binop}] \end{split}$$

$$\begin{split} &S_l; F; M \vdash L_e \ \$ \ nil; S_l' \\ &\frac{S_r; F; M \vdash R_e \ \$ \ nil; S_l'}{S_l, S_r; F; M \vdash L_e \ B \ R_e \ \$ \ nil; S_l', S_r'} [\texttt{Eval} - \texttt{binop} - \texttt{C}] \end{split}$$

$$\frac{a}{\emptyset; F; M \vdash \text{len } E}$$