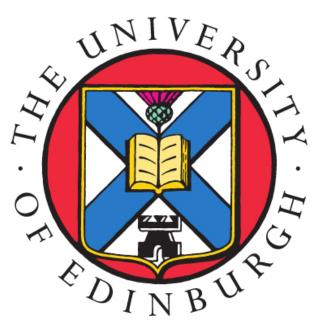


SEGA: Variance Reduction via Gradient Sketchin

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Problem setup

Optimization Problem:

$$\min_{x \in \mathbb{R}^n} F(x) = f(x) + R(x),$$

Oracle: sketched gradient for random matrix S

$$\zeta(S, x) \stackrel{\text{def}}{=} S^{\top} \nabla f(x) \in \mathbb{R}^b$$

Proximal operator of R is cheap.

R is not separable \rightarrow Coordinate descent fails.

How to do subspace descent with nonseparable prox function R?

Assumptions

M-smoothness:

 $f(x) \le f(y) + \langle \nabla f(y), x - y \rangle + \frac{1}{2}(x - y)^{\top} M(x - y)$ μ -strong convexity:

$$f(x) \ge f(y) + \langle \nabla f(y), x - y \rangle + \frac{\mu}{2} ||x - y||^2$$

→ natural for ERM with linear predictors

SEGA: Variance reduced coordinate descent

Learning gradient from sketches:

$$h^{k+1} = \arg\min_{h \in \mathbb{R}^n} ||h - h^k||^2$$

subject to $S_k^\top h = S_k^\top \nabla f(x^k)$.

Step in the unbiased direction

$$\left(\mathbf{E}\left[g^{k}\right] = \nabla f(x^{k})\right)$$
:

$$g^k = h^k + \theta_k Z_k(\nabla f(x^k) - h^k),$$

 $Z_k \stackrel{\mathrm{def}}{=} S_k \left(S_k^{ op} S_k \right)^{\dagger} S_k^{ op}.$

As $x_k \to x^*$, we have $g^k \to 0$, which is not the case for subspace descent.

Main Contributions

- SEGA with general analysis.
- Subspace SEGA.
- (Almost) recovered rates of CD.

Algorithm

Algorithm 1 SEGA: SkEtched GrAdient Method

- Initialize $x^0, h^0 \in \mathbb{R}^n$; $B \succ 0$; distribution \mathcal{D} ; stepsize $\alpha > 0$
- 2: for k = 0, 1, 2, ... do
- Sample $S_k \sim \mathcal{D}$
- $g^k = h^k + \theta_k B^{-1} Z_k (\nabla f(x^k) h^k)$
- $x^{k+1} = \operatorname{prox}_{\alpha R}(x^k \alpha g^k)$
- $h^{k+1} = h^k + B^{-1}Z_k(\nabla f(x^k) h^k)$
- 7: end for

Convergence result

Suppose that S is uniform distribution of standard coordinate basis vectors. Then with stepsize $\alpha = \Omega(\frac{1}{n\lambda_{\max}(M)})$ we have

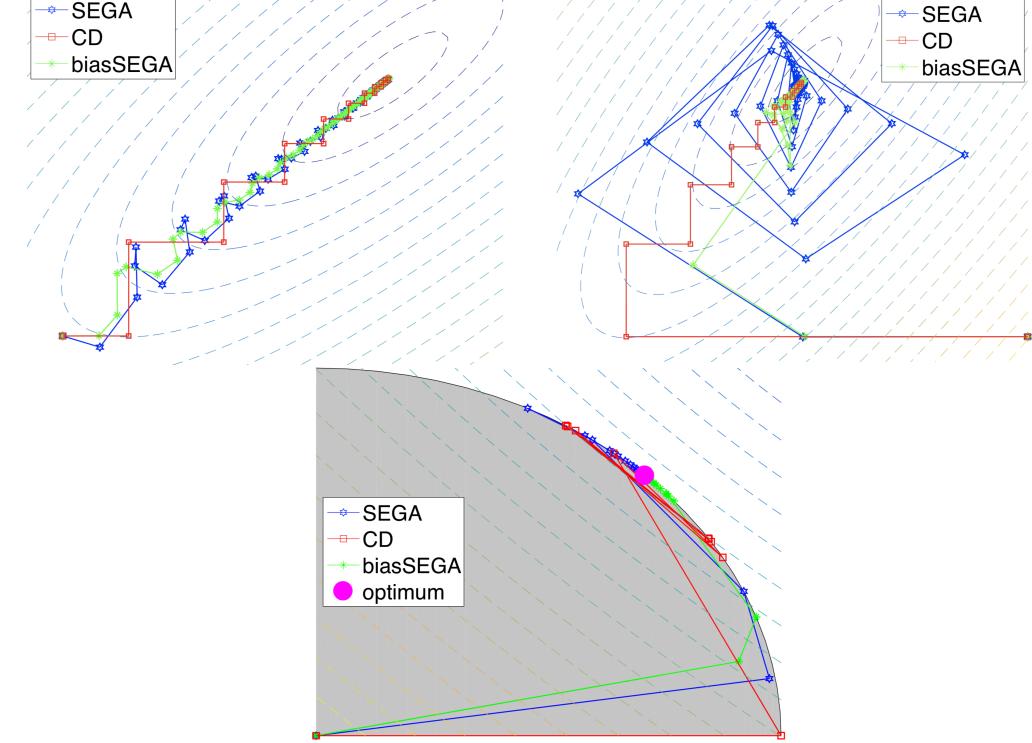
$$\mathbb{E}\Phi^{k+1} \le (1 - \alpha\mu)\Phi^k$$

for
$$\Phi^k = ||x^k - x^*||^2 + \sigma\alpha ||h^k - \nabla f(x^*)||^2$$
.

General convergence with arbitrary weighted norm and arbitrary sketching distribution is provided.

Algorithm behavior

Iterates evolution of SEGA, coordinate descent and biasSEGA (updates made using h^k instead of g^k).



Here R is the indicator function of the unit ball \Rightarrow CD does not converge!

Subspace SEGA

Suppose $f(x) = \phi(Ax)$.

Idea: Exploit that ∇f lies in known subspace. New update for h (and g):

$$h^{k+1} = \arg\min_{h \in \mathbb{R}^n} ||h - h^k||^2$$

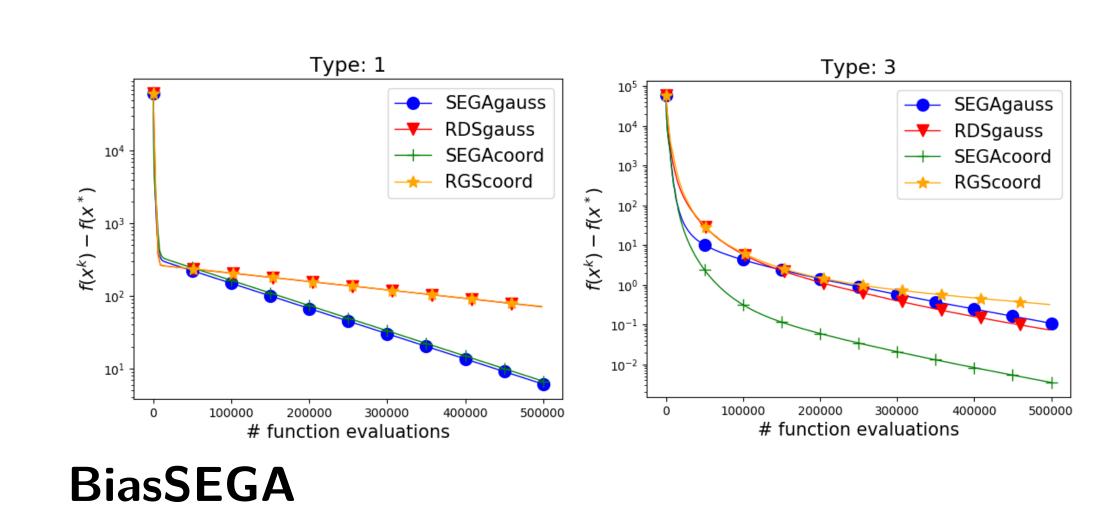
subject to $S_k^\top h = S_k^\top \nabla f(x^k)$
 $h \in \mathbf{Range}(A^\top)$

If S_k is sampled from columns of A^{\top} , we might achieve $\Omega(\frac{n}{d})$ speedup over the naive version of SEGA. $(A \in \mathbb{R}^{d \times n})$

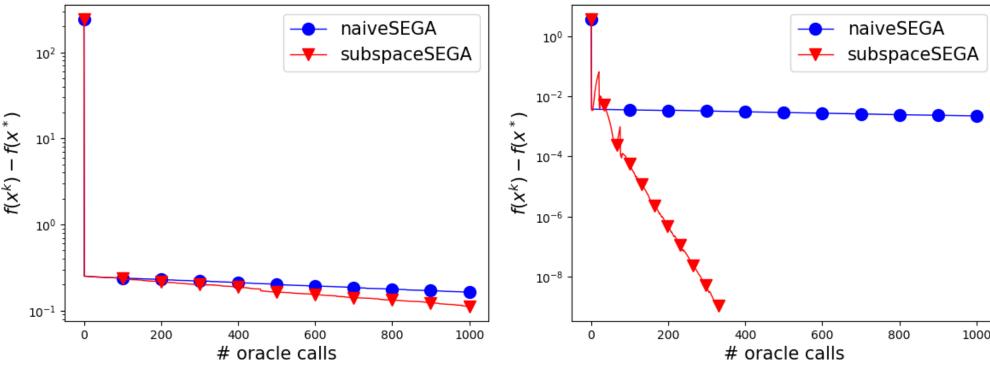
Experiments

Zeroth-order setting

- Directional gradient is $\Omega(n)$ times cheaper to full with forward diff
- Comparison with Random Direct Search [1]

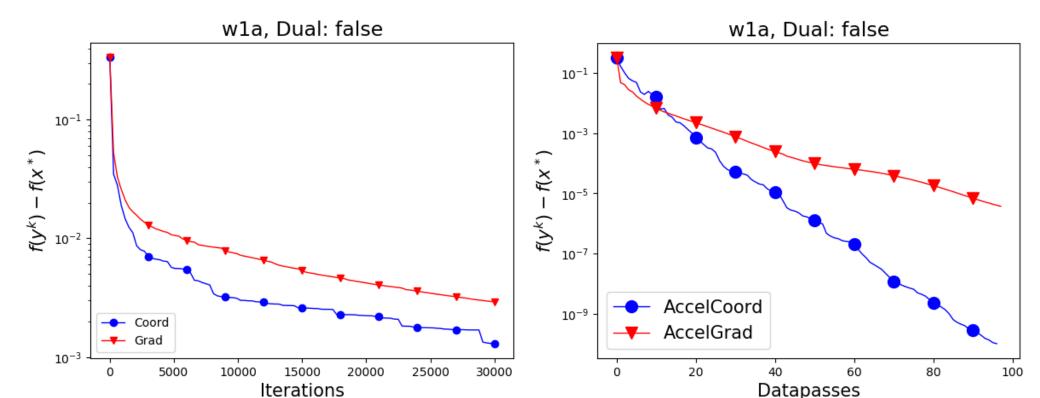


naiveSEGA



d/n = 0.01

Comparison to CD



SEGA at coordinate descent setup

- Sketches S are column submatrices of identity
- Probability vector $p: \mathbb{P}(e_i \in S) = p_i$
- Probability matrix $P: \mathbb{P}(e_i \in S, e_j \in S) = P_{i,j}$
- ESO vector v (for minibatching):

$$P \circ M \preceq \mathbf{Diag}(p \circ v)$$

Accelerated SEGA

Algorithm 2 ASEGA: Accelerated SEGA

1:
$$x^0=y^0=z^0\in\mathbb{R}^n$$
; $h^0\in\mathbb{R}^n$; S ; parameters $\alpha,\beta,\tau,\mu>0$

2: **for**
$$k = 1, 2, ...$$
 do

3:
$$x^k = (1 - \tau)y^{k-1} + \tau z^{k-1}$$

4: Sample
$$S_k$$
, and compute g^k, h^{k+1}

5:
$$y^k = x^k - \alpha p^{-1} \circ g^k$$

6:
$$z^k = \frac{1}{1+\beta\mu}(z^k + \beta\mu x^k - \beta g^k)$$

7: end for

Rates

Method	Complexity
Nonaccelerated,	$8.55 \cdot \frac{\mathbf{Tr}(M)}{\mu} \log \frac{1}{\epsilon}$
importance sampling,	$\frac{109}{\mu}$
Nonaccelerated,	$8.55 \cdot \left(\max_{i} \frac{v_i}{p_i \mu}\right) \log \frac{1}{\epsilon}$
arbitrary sampling	$\left(\frac{1110 \Lambda_l}{p_i \mu} \right) \frac{108}{\epsilon}$
Accelerated,	$9.8 \cdot \frac{\sum_{i} \sqrt{M_{ii}}}{\sqrt{\mu}} \log \frac{1}{\epsilon}$
importance sampling,	$\sqrt{\mu}$ $\sqrt{\mu}$ ϵ
Accelerated,	$9.8 \cdot \sqrt{\max_i \frac{v_i}{p_i^2 \mu} \log \frac{1}{\epsilon}}$
arbitrary sampling	$\int \int $

Up to constant factor same rates as CD [2,

References

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