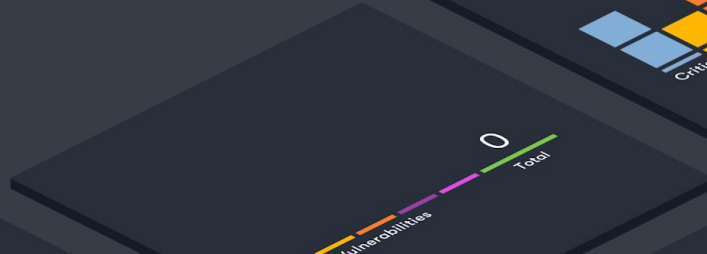


# Proof-of-Work Blockchains and Smart Contract Security

Martin Derka, Ph.D.  
Senior Research Engineer



# Talk structure

1. **Consensus and Proof-of-Work**
2. **Blockchain as a programmable platform**
3. **Possible security flaws**
4. **Research topics**

*Satoshi Nakamoto (2008)*

*Bitcoin: A Peer-to-Peer Electronic Cash System*

*Vitalik Buterin (2013)*

*Ethereum: The Ultimate Smart Contract and Decentralized Application Platform*

# Consensus and Proof-of-Work

# Consensus replacing banks

**Imagine that you have money at a bank...**



# Consensus replacing banks

**Imagine that you have money at a bank...**

You trust that the bank:

1. Handles the balances honestly
2. Does not prevent you from access
3. Does not make errors
4. Does not go out of business
5. Protects your personal data
6. Employs honest employees
7. ...



# Consensus replacing banks

**Is there another way of owning assets?**

**Consensus:**

General agreement. Everybody knows and agrees that you own assets.

# Developing a consensus

## Set up:

1. Students gather in a classroom.
2. Everyone maintains a personal copy of the classroom's ledger.
3. Proposed transactions are written on a board (in parallel).

**In order to achieve consensus**, we all need to execute the same transactions in our copies of the ledger.

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**In order to achieve consensus**, we all need to execute the same transactions in our copies of the ledger.

**Who chooses which transactions will be executed and in which order?**



# What can go wrong?

**Martin's original balance: \$1000**

## Transactions

\$700 from Martin to Alice  
\$700 from Martin to Bob

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3. Alice      +\$700
4. Reject Martin to Bob

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## Ledger 2

1. Martin    \$1000
2. Martin    -\$700
3. Bob        +\$700
4. Reject Martin to Alice
- ...

# Need to have a leader

**We still need a leader who will determine the truth.**  
(such as the order of transactions)

**One leader can be unfair.**  
(This is the bank)

**We want a different leader every single time.**

# Problems with marathons

**Choose leader by a marathon race.**



# Problems with marathons

**Choose leader by a marathon race.**



## **Problems:**

1. Authenticity of transactions
2. Privacy of transactions
3. Co-location of participants and transaction transmission
4. Learning the history of the ledger and trusting it
5. It will be very slow because organizing and running marathons takes time
6. What is the motivation to run marathons for others?

# Digital Proof-of-Work

**Marathon race** is analogous to what we call a **Proof-of-Work**:

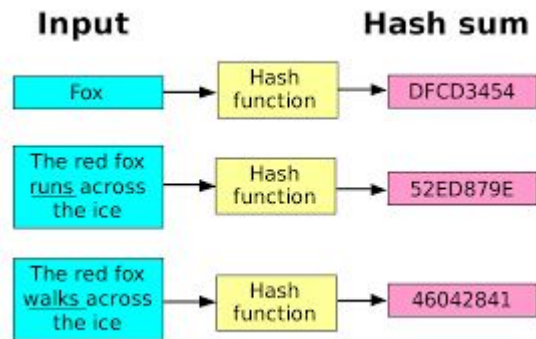
You do something that takes lots of effort. If you do it faster than anyone else, you become the leader.

In the digital world, the equivalent of **work** is **computing**.

**Proof-of-Work**: Spending computing power to solve some puzzle.

# Hashing

**Hash function** is a function that maps data of an arbitrary length to a short fixed-length representation.



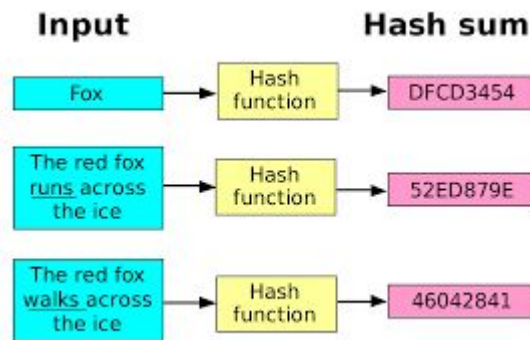


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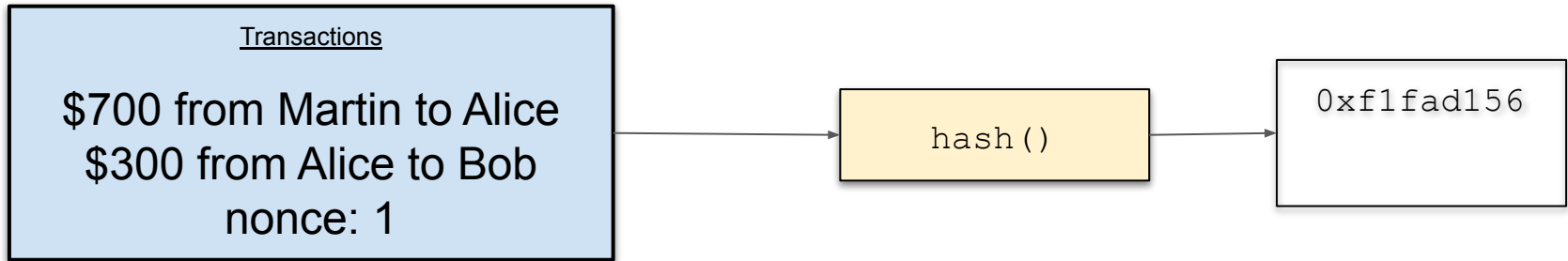
**Good hash functions are:**

1. Irreversible
2. Unpredictable
3. Evenly distributed
4. Other colliding inputs cannot be guessed



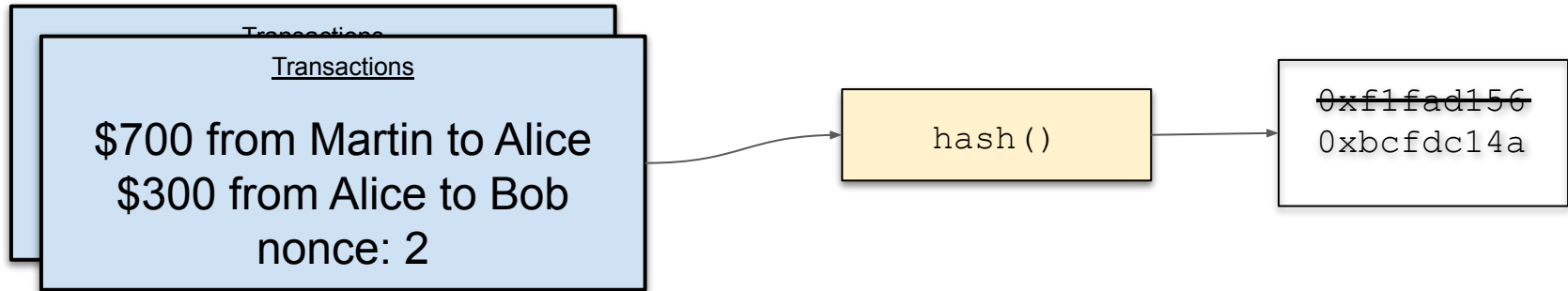
# Proof-of-Work puzzle

1. Take transactions that you want to append to the ledger.
2. Append a number (nonce) to them and hash the message
3. Find a nonce such that the produced hash starts with “00” (or with another restricted prefix)



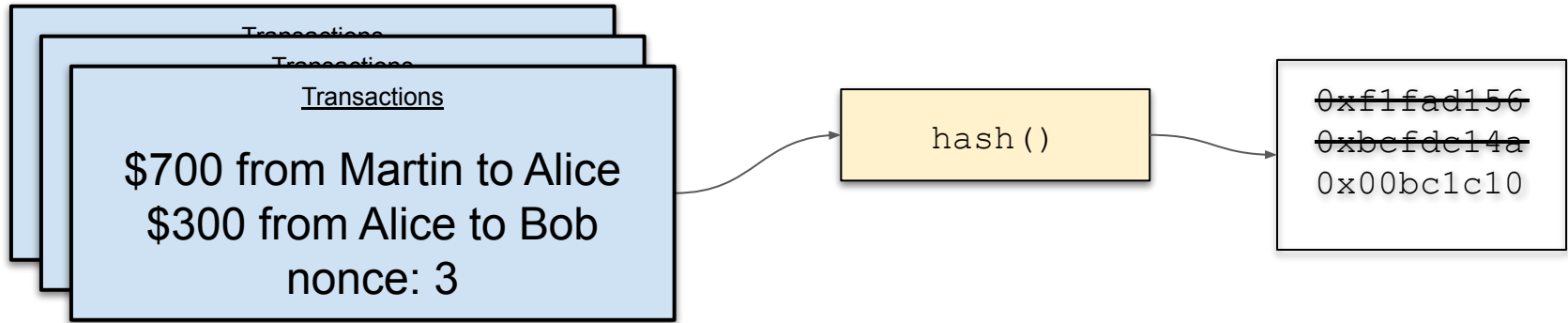
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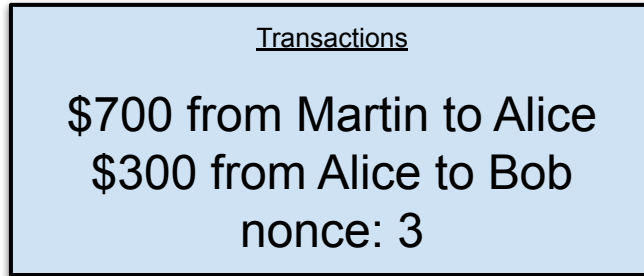
# Linking the blocks

## Transactions

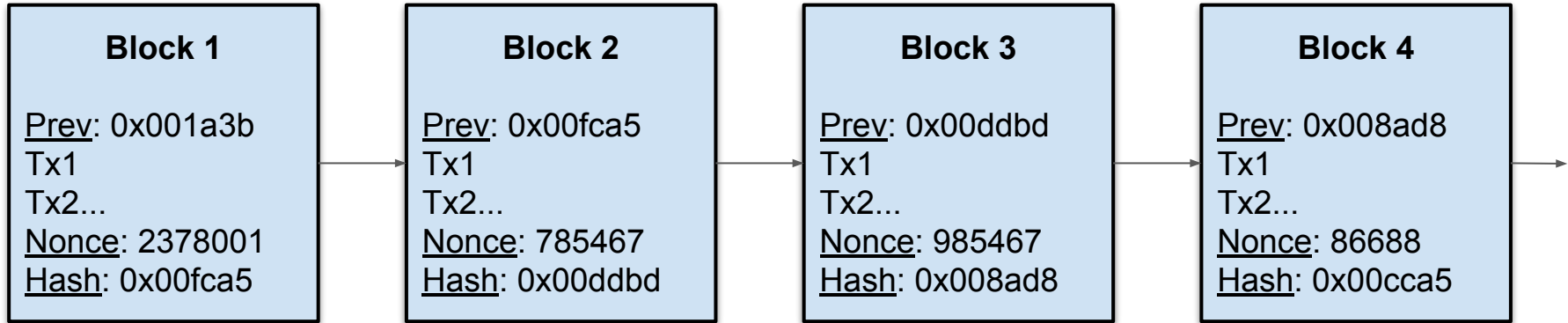
\$700 from Martin to Alice  
\$300 from Alice to Bob  
nonce: 3

= Block (almost)

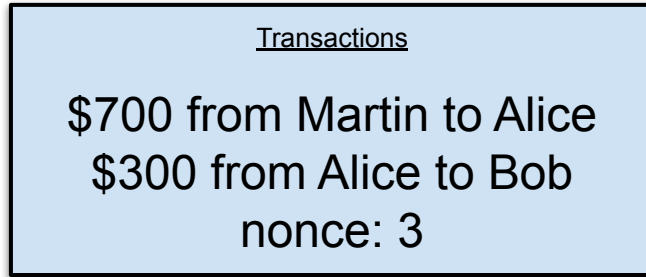
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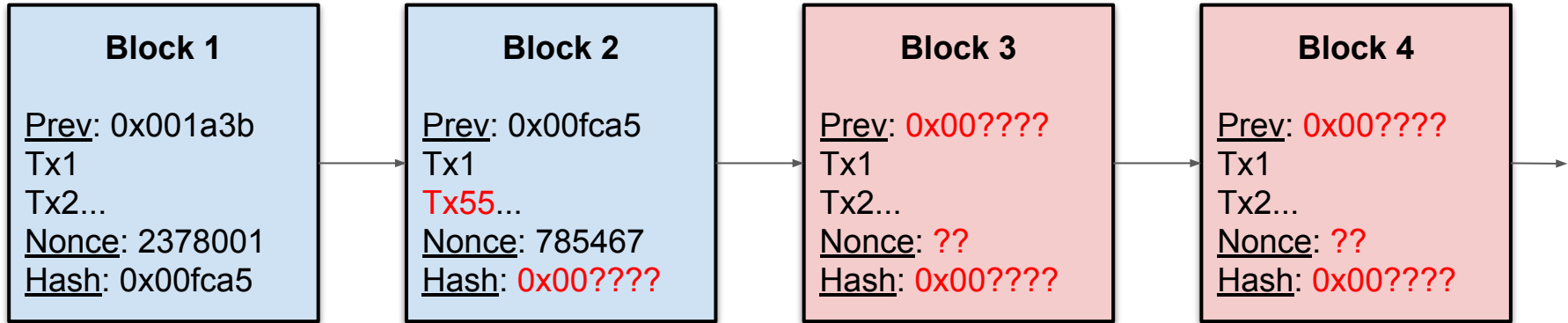
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# Proof-of-Work blockchain

1. Transactions and blocks are broadcasted over a P2P network.



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4. The first participant to succeed broadcasts the block to the network.
5. All participants prune the transactions and repeat.

# Proof-of-Work using hash functions

**Nonce** = An integer that complements the block for it to hash to a value smaller than some constant (given by the currently set **difficulty**).

**Proof-of-Work** = Guessing the right nonce.

**Proof-of-Work is hard to come up with, but quick to verify!**

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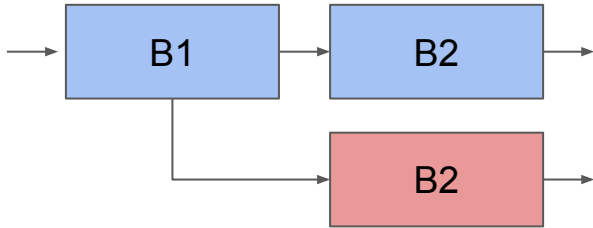
**Proof-of-Work** = Guessing the right nonce.

**Proof-of-Work is hard to come up with, but quick to verify!**

Providing the Proof-of-Work for transactions is also called **mining** (Bitcoin, Ethereum). The winner can append a special transaction that will award him with a **block reward** (“money” printed from the thin air).

# Forks in chains

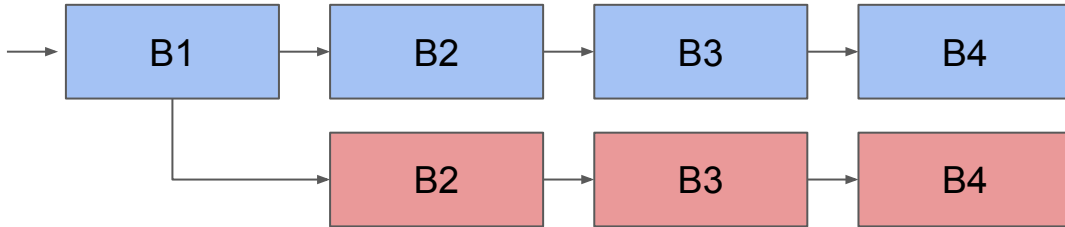
What happens if two miners arrive to a difference proof of work?



# Forks in chains

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**Basic assumption:** Proof of work is hard. Nobody can be the leader forever.

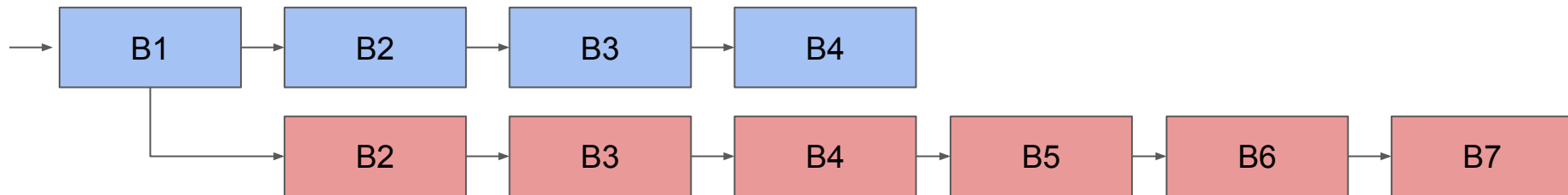




# Forks in chains

What happens if two miners arrive to a difference proof of work?

**Basic assumption:** Proof of work is hard. Nobody can be the leader forever.



**Solution:** Always trust the longest chain.

# Problems with our system

**Choose leader by Proof-of-Work in a P2P network.**

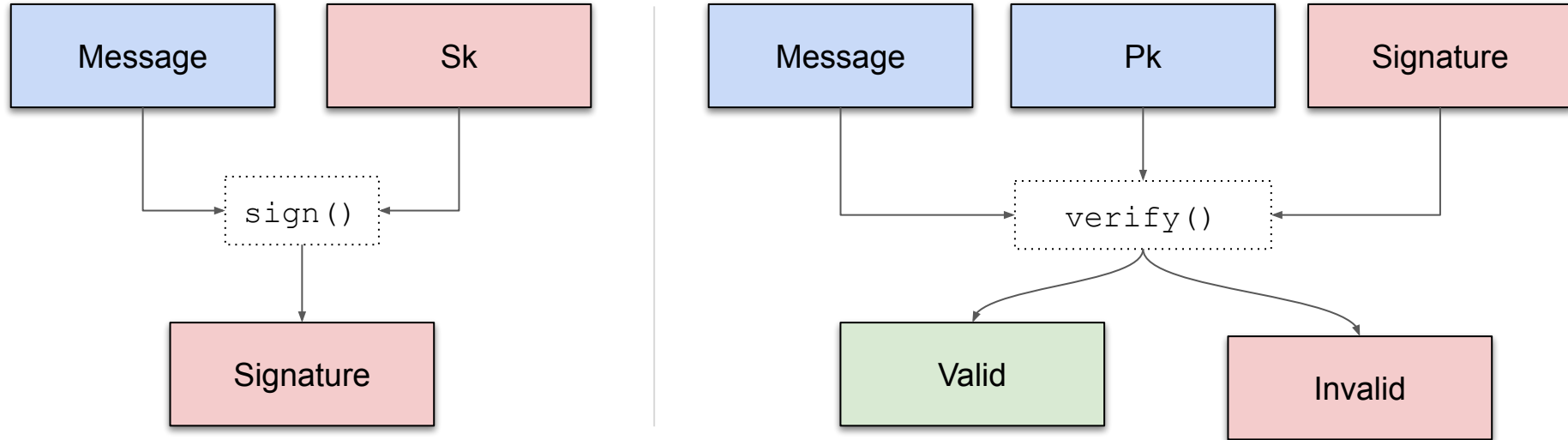
## **Problems:**

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# Cryptography in Blockchain

Public key = Account that identifies you (wallet address)

Private key = Cryptographic key that allows you to sign transactions



# Digital signature in blockchain

Public key = Account that identifies you (wallet address)  
Private key = Cryptographic key that allows you to sign transactions

## Block 1

Prev: 0x001a3b

Tx1

Tx2...

Nonce: 2378001

Hash: 0x00fca5

## Transaction:

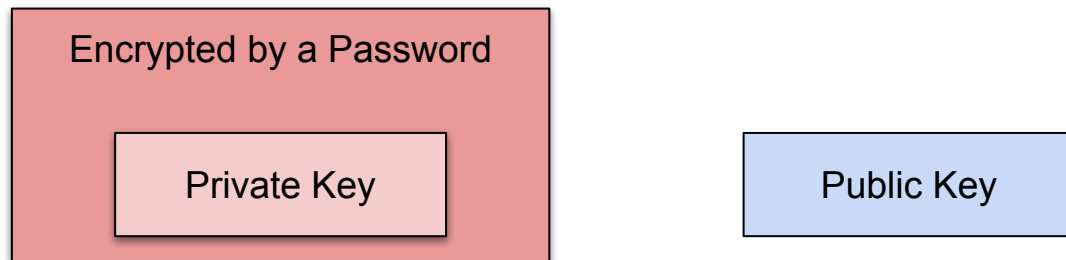
1. Source account (Pk)
2. Target account (Pk),
3. How much
4. Signature (using Sk of the source account)

# Digital signature in blockchain

Private-public key pairs can be easily generated locally (RSA, ECDSA).

RSA: `ssh-keygen`

ECDSA for Ethereum: `geth`, `MyEtherWallet`



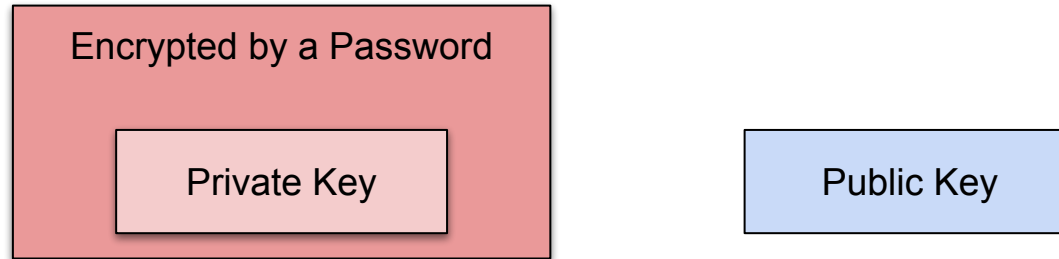
Everyone can obtain an account on blockchain!

# Digital signature in blockchain

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By replacing people's names with addresses, we introduced the concept of **pseudoanonymity**.

# Proof-of-Work Blockchain

**Choose leader by proof of work in a P2P network, use public keys as accounts, sign the transactions using private keys.**

## **Problems:**

- ✓ Authenticity of transactions
- ✓ Privacy of transactions
- ✓ Colocation of participants in a single room
- ✓ Learning the history of the ledger and trusting it
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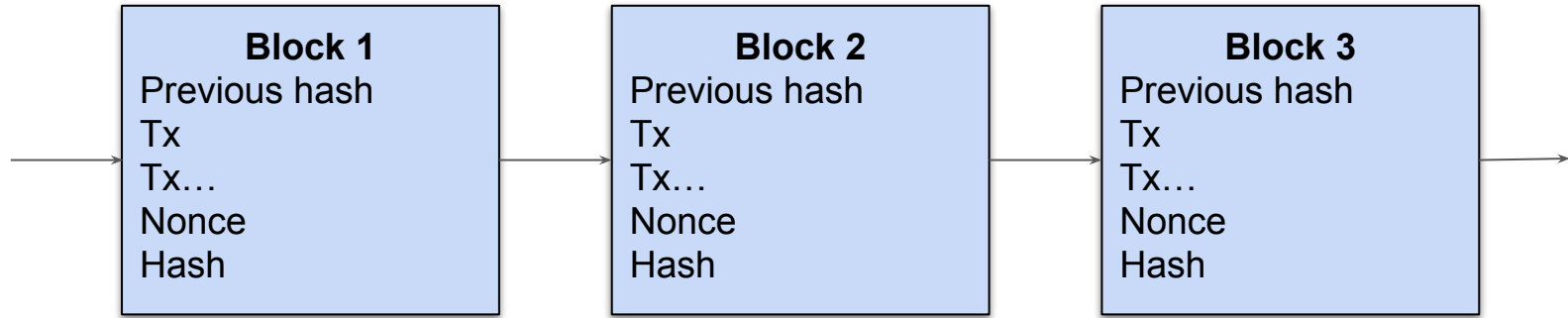
# **From Blockchain to Computational Platform**



# Computing on blockchain

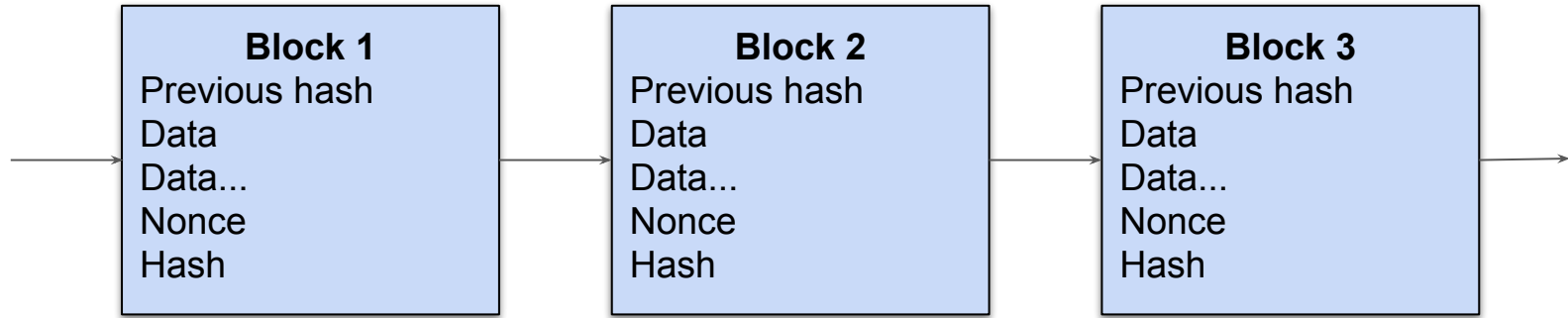
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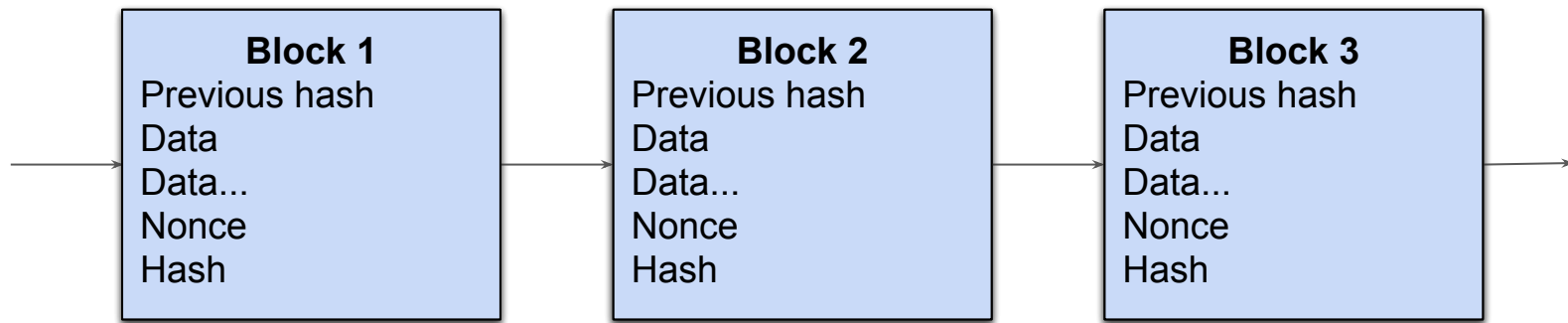
**Key idea:** Blockchain allows us to store data. This can be transactions or anything else!



# Computing on blockchain

**Key idea:** Blockchain allows us to store data. This can be transactions or anything else!

**Use the “Data” to store instructions of the program!**



# Computing on blockchain

**Ethereum Virtual Machine (EVM):** Stack-based virtual machine described by V. Buterin in 2013 and G. Wood in 2014 - 2019.

**Smart contract:** Program for the EVM.

1. Smart contracts expose API that can be called within transactions by Ethereum accounts

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**Smart contract:** Program for the EVM.

1. Smart contracts expose API that can be called within transactions by Ethereum accounts
2. Smart contracts define what EVM opcodes should be executed by miners in order to update their chain state

Executing each opcode costs **gas** which is used to reward the miners for using their computational power for the PoW and for maintaining the network state

# Example of a smart contract

```
1.  contract BasicToken {
2.
3.      mapping(address => uint256) balances;
4.
5.      function transfer(address _to, uint256 _value)
6.      public returns (bool) {
7.          balance.balances[msg.sender] -= _value;
8.          balances[_to] += _value;
9.          return true;
10.     }
11.
12.     ...
```

# Examples of Vulnerabilities



# Live Demo

**Re-entrancy:** A method is called **re-entrant** if it can be interrupted in the middle of its execution and then safely called again ("re-entered") before its previous invocations complete their executions.

[Live demo in Remix](#)

# Vulnerability Categories in Ethereum

Solidity	EVM	Blockchain
<b>Reentrancy</b>	Immutable bugs	Unpredictable state
Call to unknown	Ether lost in transfer	Generating randomness
Gasless send	Stack size limit	Time constraints
Exception disorders	Integer overflows	<b>Transaction ordering</b>
Type casts		Network throughput
		<b>Keeping secrets</b>

Equitable dividend-sharing game with  
payment in ethernic currency

Provable fairness The dividend-sharing game is simple Open  
Source Contracts Provide Support

All Wagers

 703610ETH

58914bets

1. A ponzi scheme deployed on July 27, 2019.
2. On September 26, it held ~49,518 ETH.
3. Vulnerabilities were found and reported in September 2019.
4. The contract is currently empty.



# Simplified Code

```
contract FairWin {

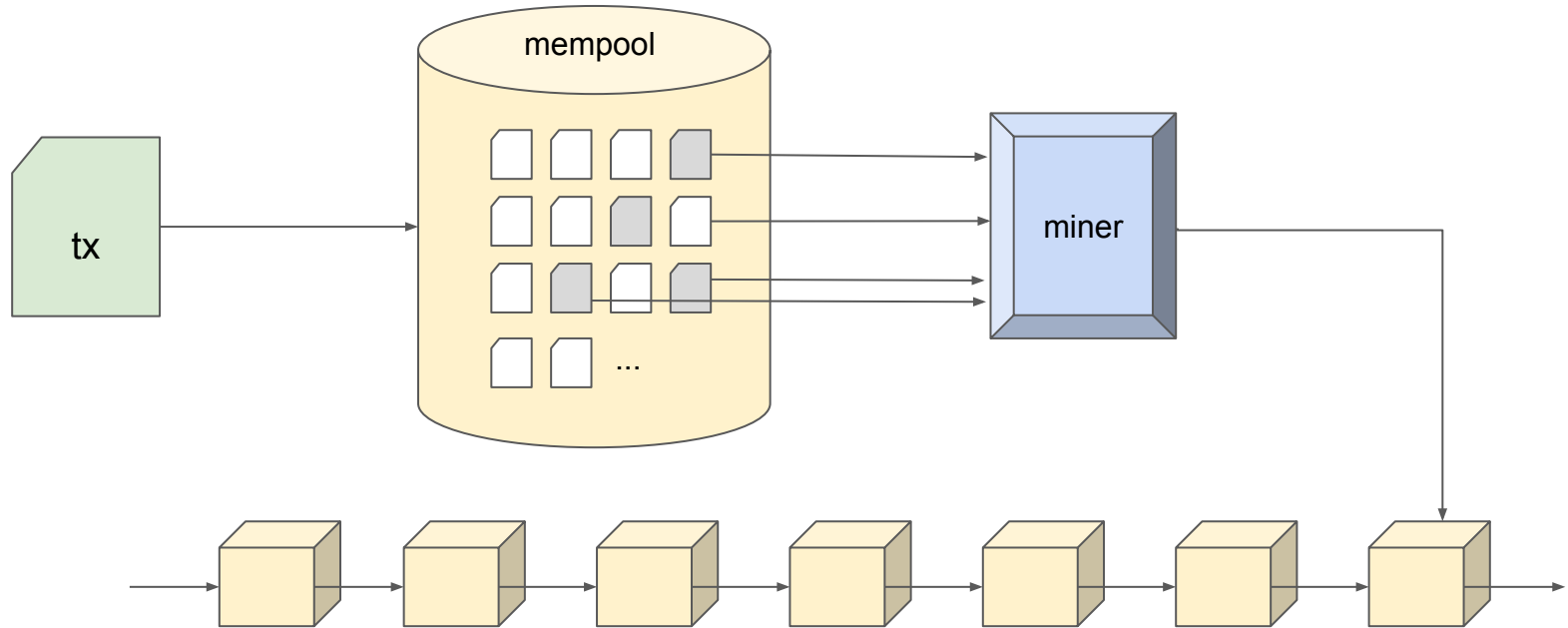
    mapping (string => address) investors;
    mapping (string => uint256) investments;

    function invest(string inviteCode, ...) payable {
        if (investors[inviteCode] == 0x0) investors[inviteCode] = address;
        investments[inviteCode] = investments[inviteCode].add(msg.value);

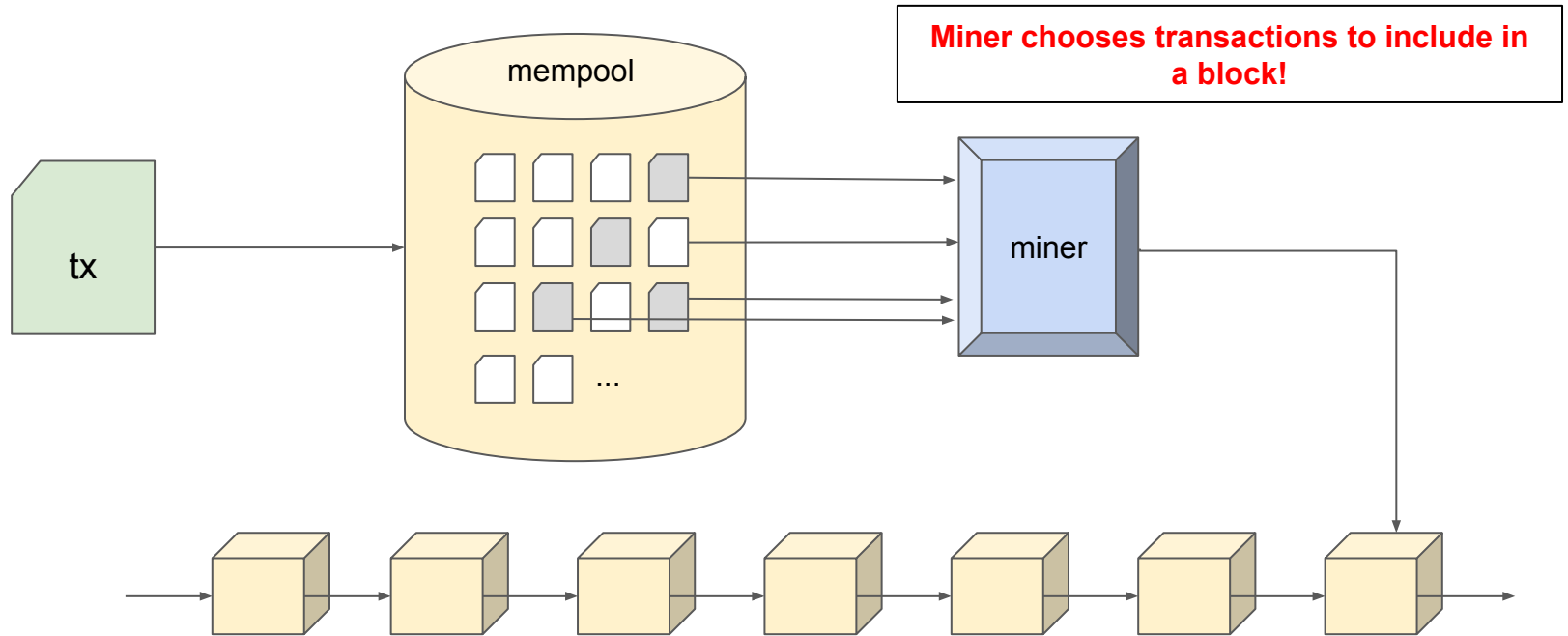
        //... other logic
    }

    function exit(string inviteCode) {
        require(investors[inviteCode] == msg.sender);
        investors[inviteCode].send(investments[inviteCode]);
    }
}
```

# Back to Proof-of-Work

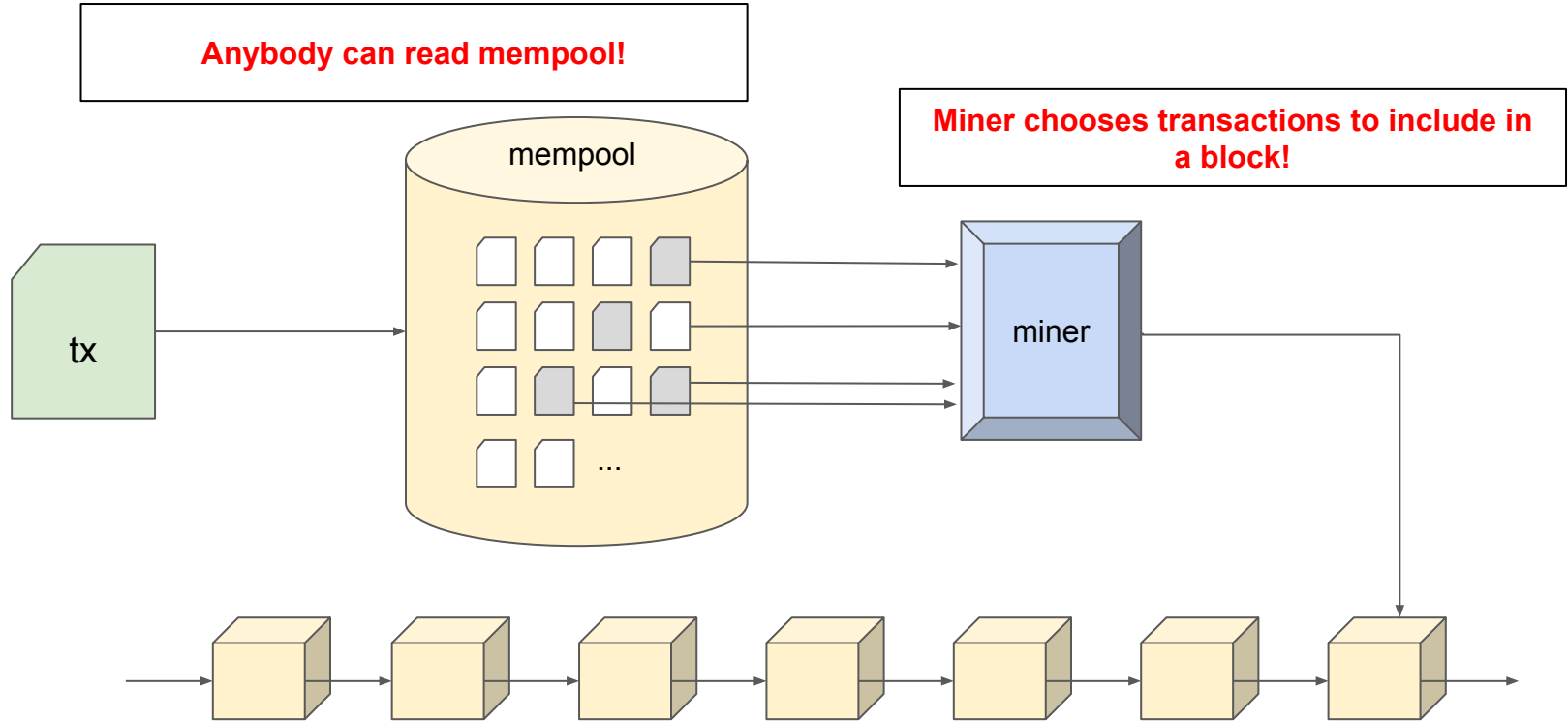


# Back to Proof-of-Work





# Back to Proof-of-Work



# Research Topics

# Research Areas

## Scaling blockchain

- Proof-of-Work takes time to produce (it cannot be too easy!)
- The number of transactions per minute can be low
- Proof-of-Work is far from being green

## Possible solutions

- Different type of consensus (Proof-of-Stake, Proof-of-Authority, Algorand)
- Side channels (plasma, chain sharding)
- Combination of on-chain and off-chain approaches

# Research Areas

## Smart contract and platform implementation

- Every opcode execution costs gas
  - *Perez, Livshits (2019): Broken Metre: Attacking Resource Metering in EVM*
- Traditional algorithms are often insufficient
  - Long iterations are costly
  - Updates to storage is expensive
  - Array lists are not usable
  - There is no primitive for generating (pseudo)random numbers
- Blockchain is not a database
  - Unpredictable availability, unpredictable rollbacks

# Research Areas

## **Tools for smart contract analysis**

- Static analysis and symbolic execution (Mythril, Oyente, Slither, Securify)
- Formal verification tools (K-Framework)
- Compiler verification tools
- EVM verification tools
- New programming languages

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## **Verification of off-chain computations**

- SNARKs
- STARKs
- BulletProofs

# Conclusions

- **Blockchain is NOT a new technology**
  - Orchestration of known techniques in a new way
  - Using it is prone to making mistake
- **Social revolution**
  - Redistribution of wealth
  - Means of spreading access to services
  - New field for tech business
  - DeFi (decentralized finance)



# Conclusions

- **Popularization of computer science**
  - LaTeX is cool
  - Turing completeness is cool
  - Randomness, probability and cryptography is cool
  - Algebra and game theory is cool



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**Thank you!**