



# CS 5541 – Computer Systems

"Based on lecture notes developed by Randal E. Bryant and David R. O'Hallaron in conjunction with their textbook "Computer Systems: A Programmer's Perspective"



# Module 2

## Machine Code

### Part 4 — Data

From: Computer Systems, Chapter 3

Instructor: James Yang  
<https://cs.wmich.edu/~zjiang>  
[zjiang.yang@wmich.edu](mailto:zjiang.yang@wmich.edu)

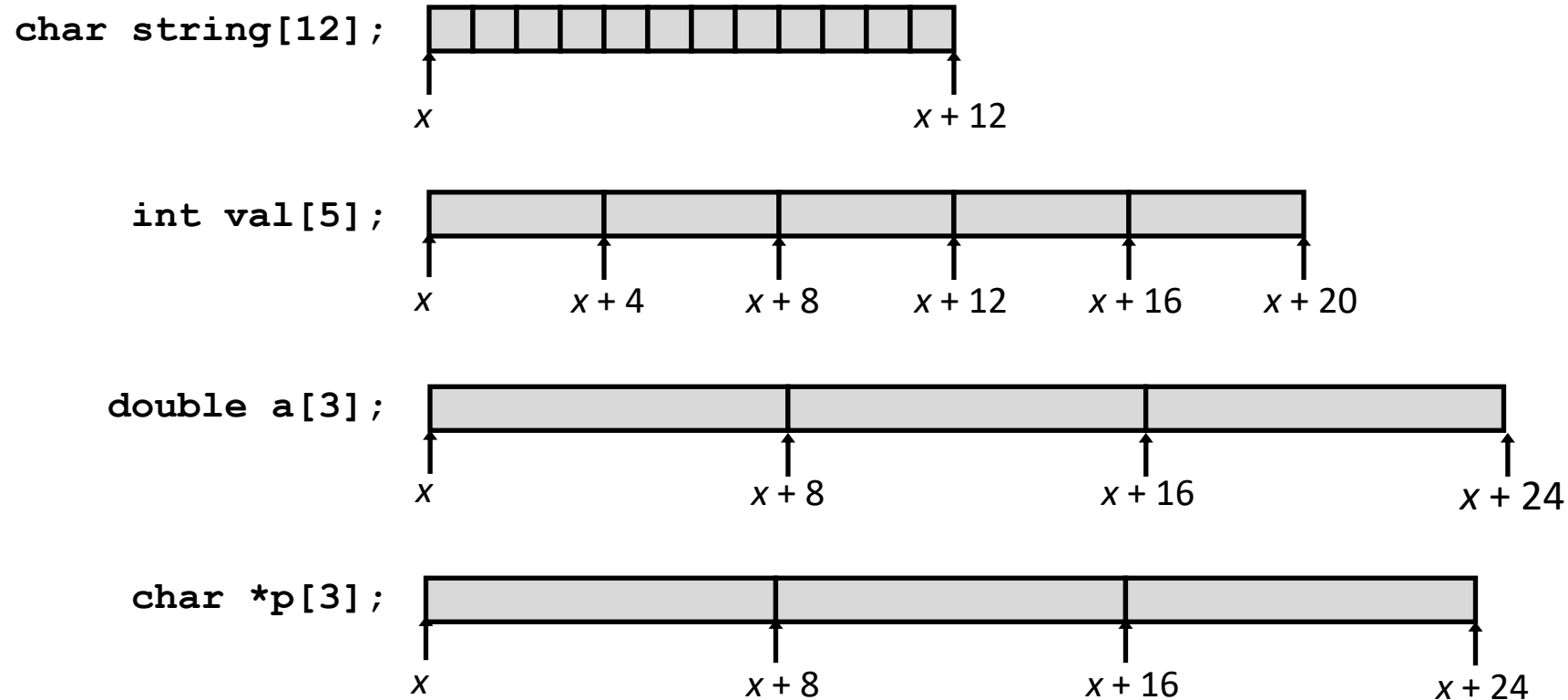


# Array Allocation

- Basic Principle

$T$   $A[L]$  ;

- Array of data type  $T$  and length  $L$
- Contiguously allocated region of  $L * \text{sizeof}(T)$  bytes in memory

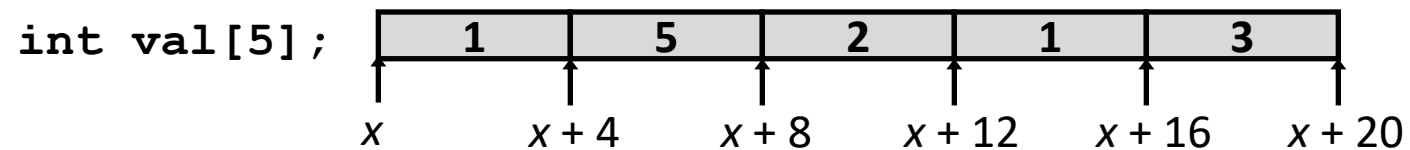


# Array Access

- Basic Principle

$T$  **A**[ $L$ ];

- Array of data type  $T$  and length  $L$
- Identifier **A** can be used as a pointer to array element 0: Type  $T^*$



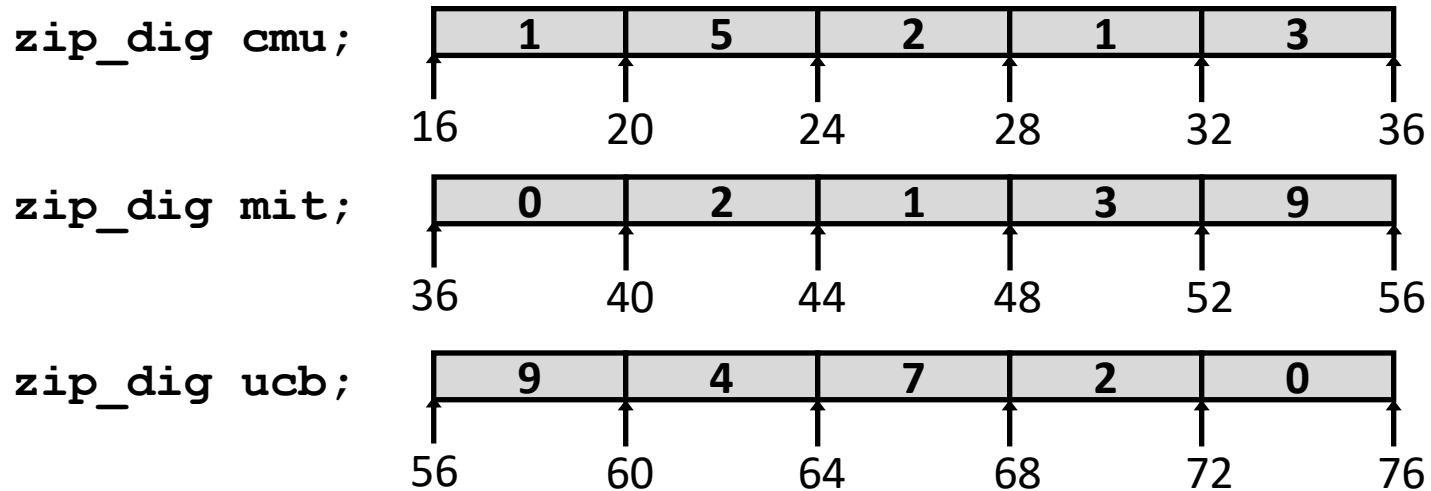
Reference	Type	Value
<code>val[4]</code>	<code>int</code>	3
<code>val</code>	<code>int *</code>	$x$
<code>val+1</code>	<code>int *</code>	$x+4$
<code>&amp;val[2]</code>	<code>int *</code>	$x+8$
<code>val[5]</code>	<code>int</code>	??
<code>*(val+1)</code>	<code>int</code>	5
<code>val + i</code>	<code>int *</code>	$x+4i$



# Array Example

```
#define ZLEN 5
typedef int zip_dig[ZLEN];

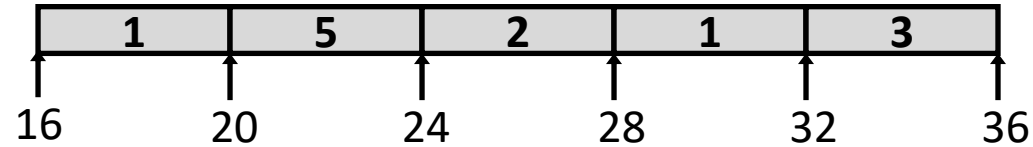
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration “`zip_dig cmu`” equivalent to “`int cmu[5]`”
- Example arrays were allocated in successive 20 byte blocks
  - Not guaranteed to happen in general

# Array Accessing Example

zip\_dig cmu;



```
int get_digit
(zip_dig z, int digit)
{
    return z[digit];
}
```

**IA32**

```
# %rdi = z
# %rsi = digit
movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register `%rdi` contains starting address of array
- Register `%rsi` contains array index
- Desired digit at `%rdi + 4*%rsi`
- Use memory reference `(%rdi,%rsi,4)`

# Array Loop Example

```
void zincr(zip_dig z) {  
    size_t i;  
    for (i = 0; i < ZLEN; i++)  
        z[i]++;  
}
```

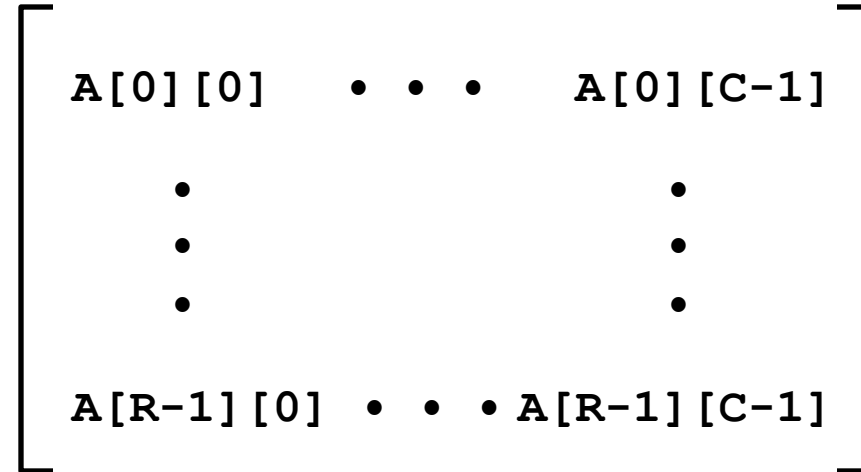
```
# %rdi = z  
movl    $0, %eax           # i = 0  
jmp      .L3               # goto middle  
.L4:                       # loop:  
    addl    $1, (%rdi,%rax,4) # z[i]++  
    addq    $1, %rax        # i++  
.L3:                       # middle  
    cmpq    $4, %rax        # i:4  
    jbe     .L4             # if <=, goto loop  
rep; ret
```

# Multidimensional (Nested) Arrays

- **Declaration**

`T A[R][C];`

- 2D array of data type  $T$
- $R$  rows,  $C$  columns
- Type  $T$  element requires  $K$  bytes



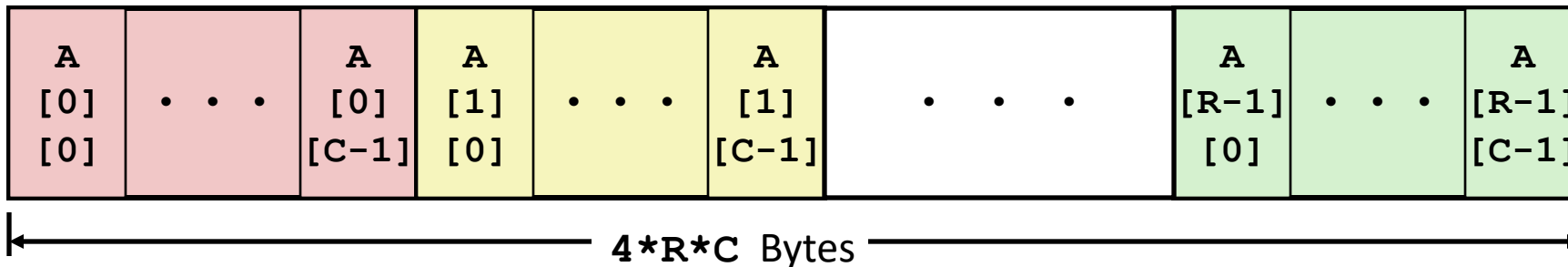
- **Array Size**

- $R * C * K$  bytes

- **Arrangement**

- Row-Major Ordering

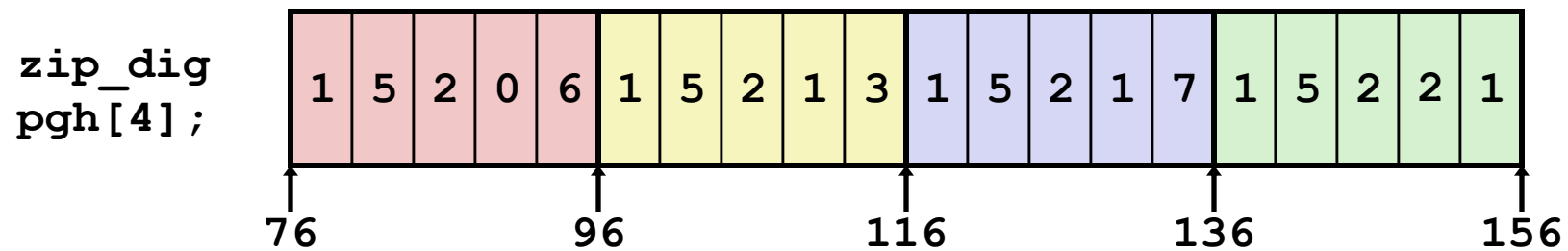
`int A[R][C];`





# Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```



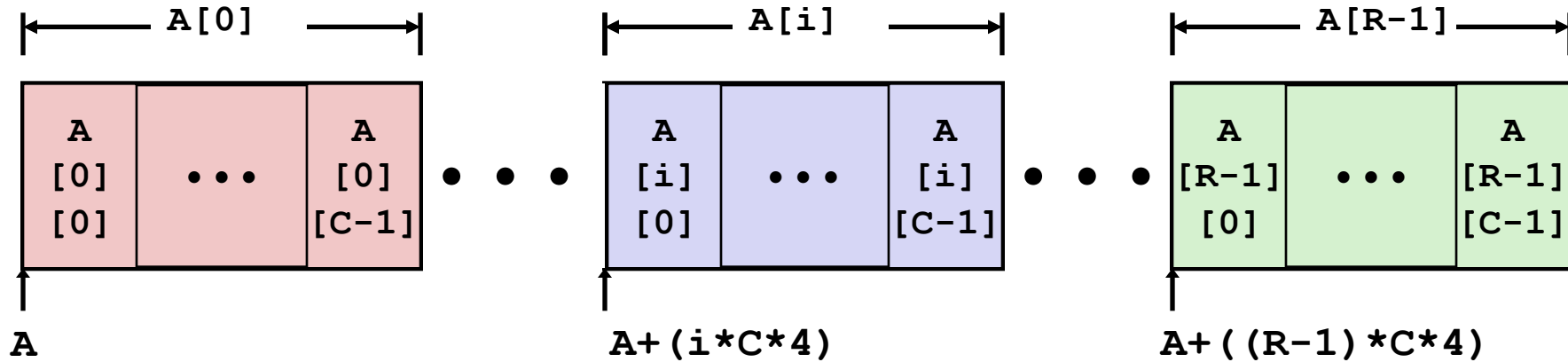
- “zip\_dig pgh[4]” equivalent to “int pgh[4][5]”
  - Variable **pgh**: array of 4 elements, allocated contiguously
  - Each element is an array of 5 **int**’s, allocated contiguously
- “Row-Major” ordering of all elements in memory

# Nested Array Row Access

- **Row Vectors**

- $\mathbf{A}[i]$  is array of  $C$  elements
- Each element of type  $T$  requires  $K$  bytes
- Starting address  $\mathbf{A} + i * (C * K)$

`int A[R][C];`

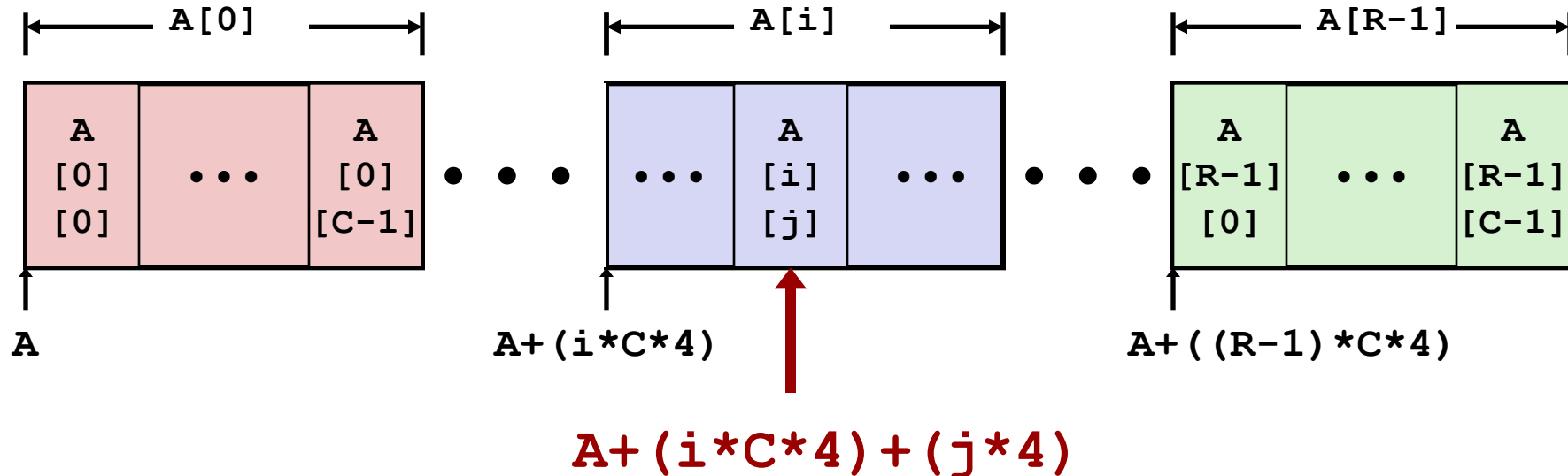


# Nested Array Element Access

- **Array Elements**

- $A[i][j]$  is element of type  $T$ , which requires  $K$  bytes
- Address  $A + i * (C * K) + j * K = A + (i * C + j) * K$

`int A[R][C];`

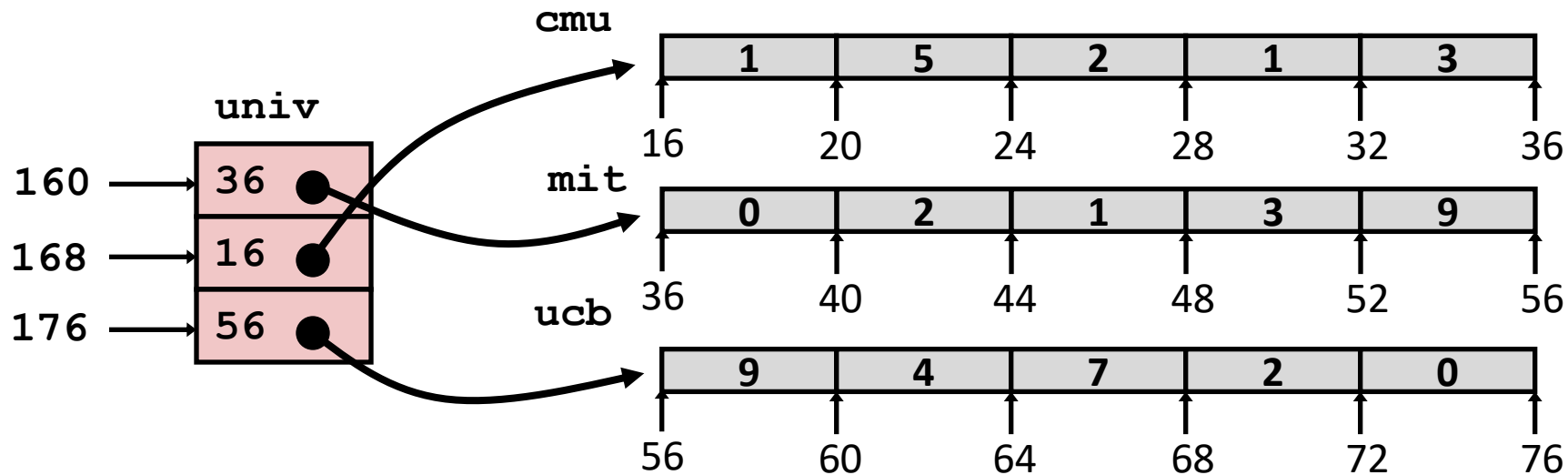


# Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig mit = { 0, 2, 1, 3, 9 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3  
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
  - 8 bytes
- Each pointer points to array of `int`'s



# Element Access in Multi-Level Array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```

```
salq    $2, %rsi          # 4*digit
addq    univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl    (%rsi), %eax       # return *p
ret
```

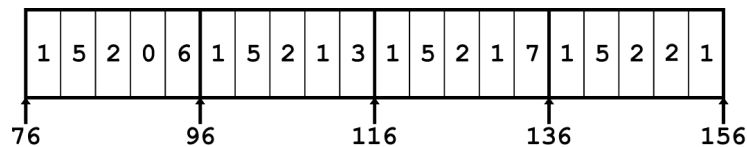
- **Computation**

- Element access **Mem[Mem[univ+8\*index]+4\*digit]**
- Must do two memory reads
  - First get pointer to row array
  - Then access element within array

# Array Element Accesses

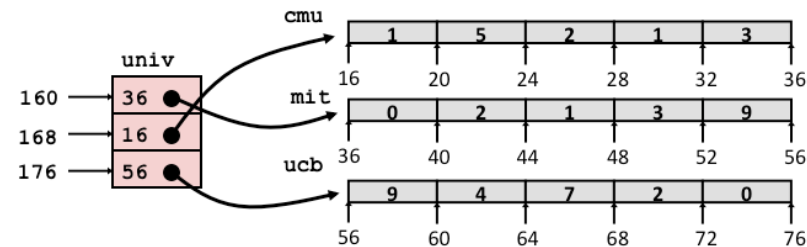
## Nested array

```
int get_pgh_digit
(size_t index, size_t digit)
{
    return pgh[index][digit];
}
```



## Multi-level array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



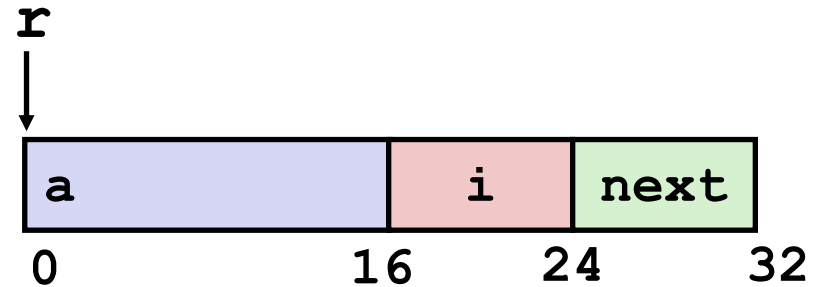
Accesses looks similar in C, but address computations very different:

$\text{Mem}[\text{pgh} + 20 * \text{index} + 4 * \text{digit}]$      $\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{digit}]$



# Structure Representation

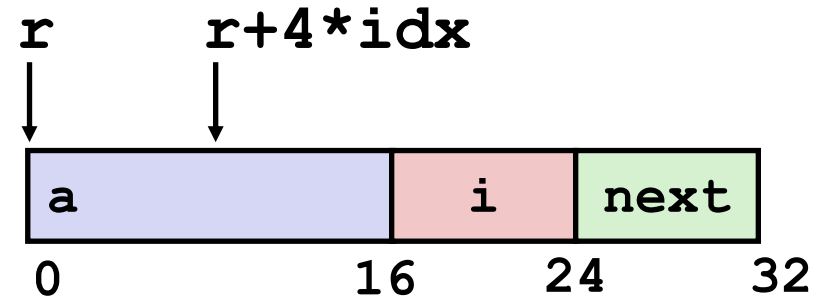
```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



- **Structure represented as block of memory**
  - Big enough to hold all of the fields
- **Fields ordered according to declaration**
  - Even if another ordering could yield a more compact representation
- **Compiler determines overall size + positions of fields**
  - Machine-level program has no understanding of the structures in the source code

# Generating Pointer to Structure Member

```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



## • Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as  $r + 4 * idx$

```
int *get_ap  
(struct rec *r, size_t idx)  
{  
    return &r->a[idx];  
}
```

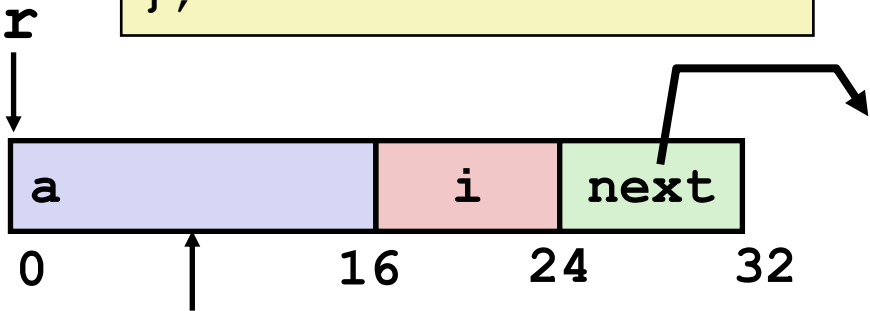
```
# r in %rdi, idx in %rsi  
leaq (%rdi,%rsi,4), %rax  
ret
```

# Following Linked List

- C Code

```
void set_val
(struct rec *r, int val)
{
    while (r) {
        int i = r->i;
        r->a[i] = val;
        r = r->next;
    }
}
```

```
struct rec {
    int a[3];
    int i;
    struct rec *next;
};
```



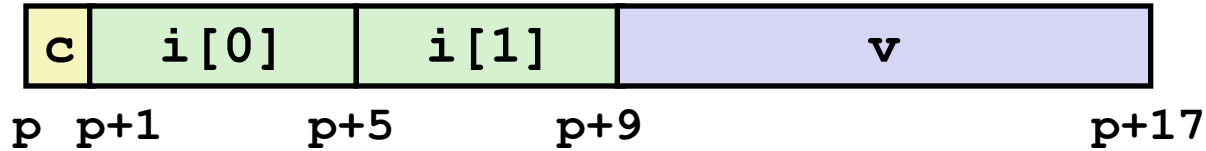
Element i

Register	Value
%rdi	r
%rsi	val

```
.L11:                                # loop:
    movslq    16(%rdi), %rax          # i = M[r+16]
    movl      %esi, (%rdi,%rax,4)    # M[r+4*i] = val
    movq      24(%rdi), %rdi         # r = M[r+24]
    testq     %rdi, %rdi             # Test r
    jne       .L11                  # if !=0 goto loop
```

# Structures & Alignment

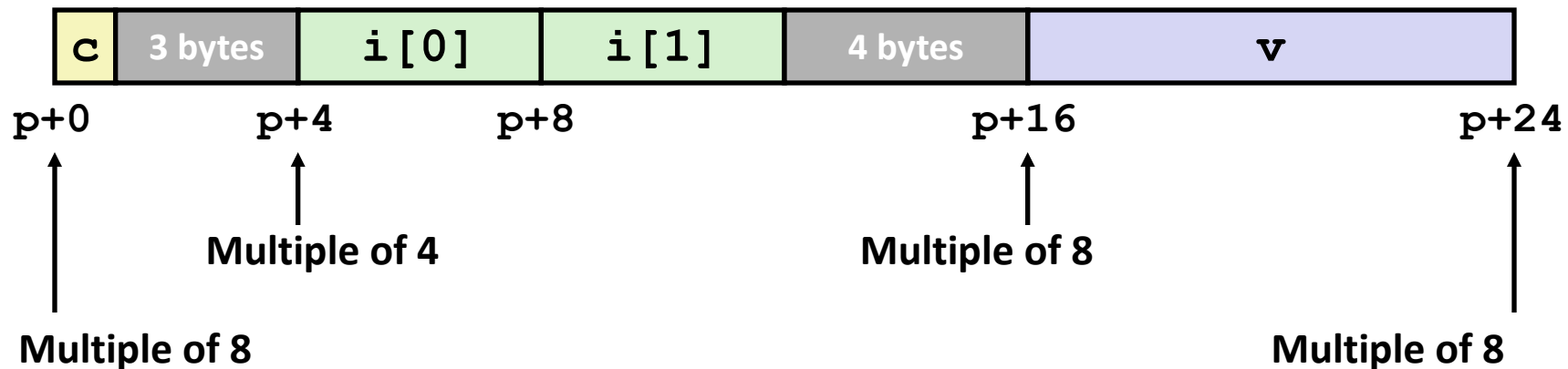
- Unaligned Data



```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```

- Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



# Alignment Principles

- **Aligned Data**
  - Primitive data type requires K bytes
  - Address must be multiple of K
  - Required on some machines; advised on x86-64
- **Motivation for Aligning Data**
  - Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
    - Inefficient to load or store datum that spans quad word boundaries
    - Virtual memory trickier when datum spans 2 pages
- **Compiler**
  - Inserts gaps in structure to ensure correct alignment of fields

# Specific Cases of Alignment (x86-64)

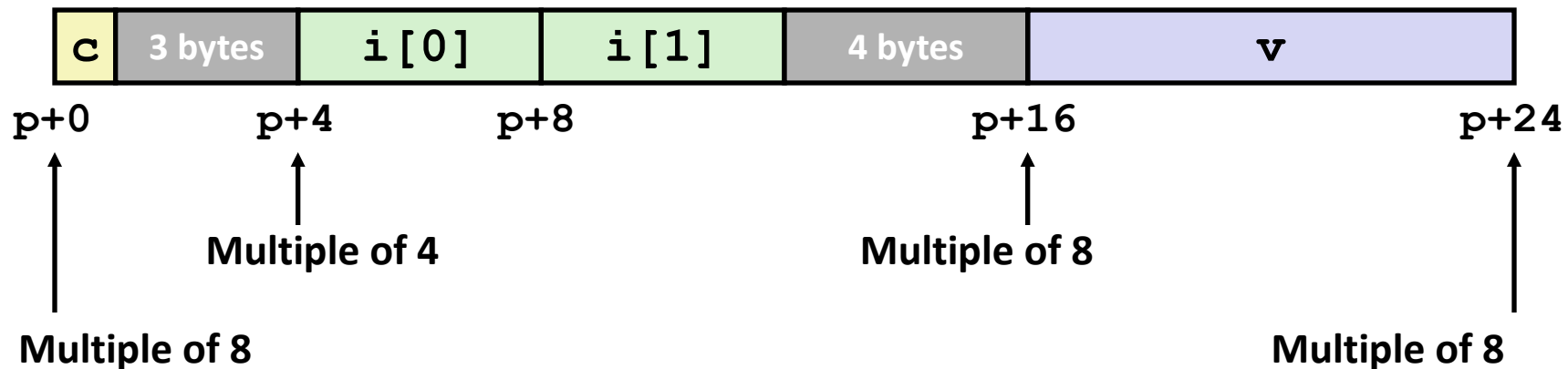
- **1 byte: `char`, ...**
  - no restrictions on address
- **2 bytes: `short`, ...**
  - lowest 1 bit of address must be 0<sub>2</sub>
- **4 bytes: `int`, `float`, ...**
  - lowest 2 bits of address must be 00<sub>2</sub>
- **8 bytes: `double`, `long`, `char *`, ...**
  - lowest 3 bits of address must be 000<sub>2</sub>
- **16 bytes: `long double` (GCC on Linux)**
  - lowest 4 bits of address must be 0000<sub>2</sub>



# Satisfying Alignment with Structures

- **Within structure:**
  - Must satisfy each element's alignment requirement
- **Overall structure placement**
  - Each structure has alignment requirement K
    - K = Largest alignment of any element
  - Initial address & structure length must be multiples of K
- **Example:**
  - K = 8, due to **double** element

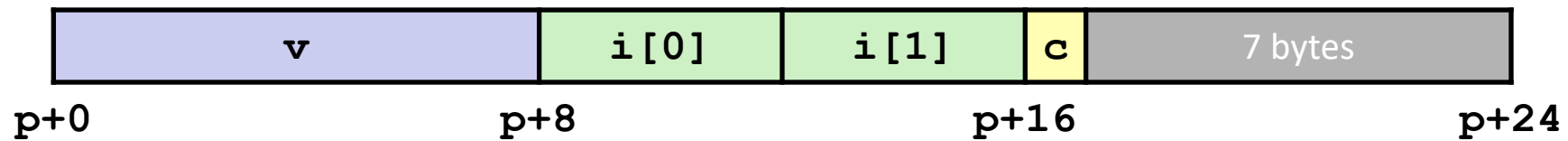
```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```



# Meeting Overall Alignment Requirement

- For largest alignment requirement K
- Overall structure must be multiple of K

```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} *p;
```

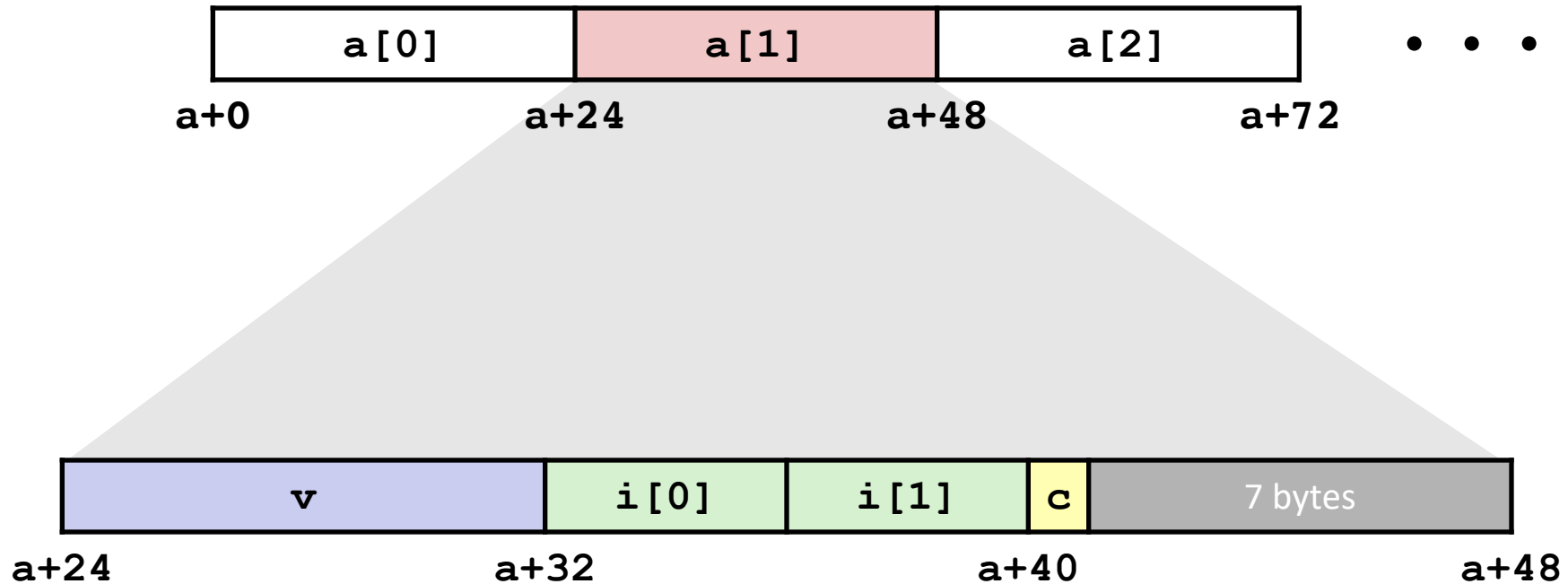


Multiple of K=8

# Arrays of Structures

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} a[10];
```



# Saving Space

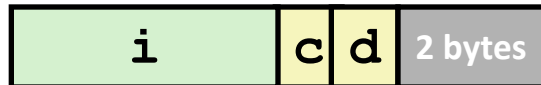
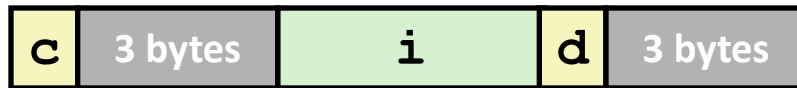
- Put large data types first

```
struct S4 {  
    char c;  
    int i;  
    char d;  
} *p;
```



```
struct S5 {  
    int i;  
    char c;  
    char d;  
} *p;
```

- Effect (K=4)



# Module 2 (Part 4)

## Summary

- **Arrays**
  - Elements packed into contiguous region of memory
  - Use index arithmetic to locate individual elements
- **Structures**
  - Elements packed into single region of memory
  - Access using offsets determined by compiler
  - Possible require internal and external padding to ensure alignment
- **Combinations**
  - Can nest structure and array code arbitrarily

