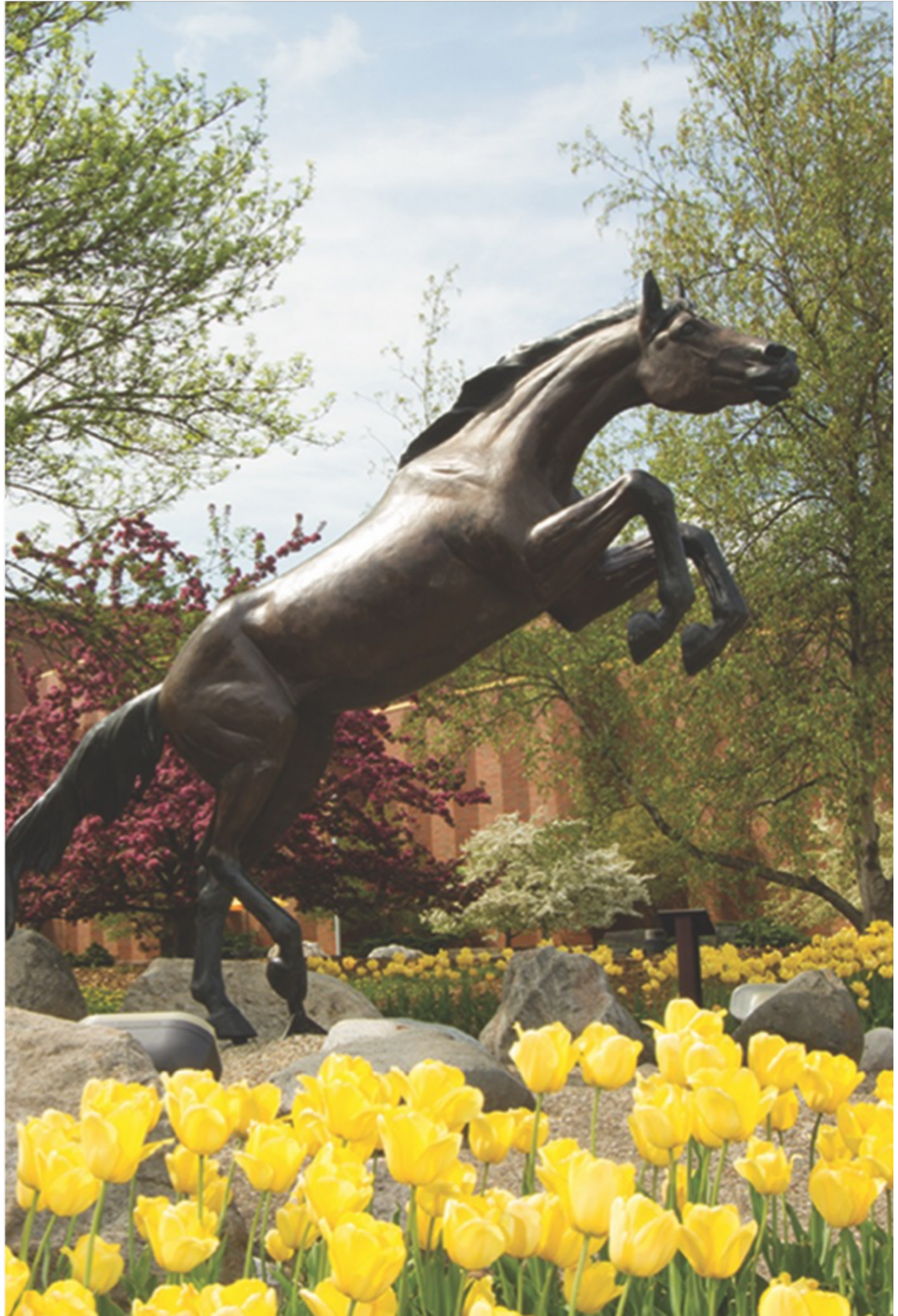
WESTERN MICHIGAN
UNIVERSITY

CS 5541 – Computer Systems

"Based on lecture notes developed by Randal E. Bryant and David R. O'Hallaron in conjunction with their textbook "Computer Systems: A Programmer's Perspective"



Module 4

Memory Allocation

Part 2 – Explicit and Segregated Free Lists

From: Computer Systems, Chapter 9, Section 9

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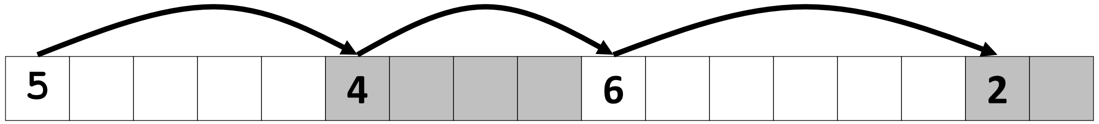
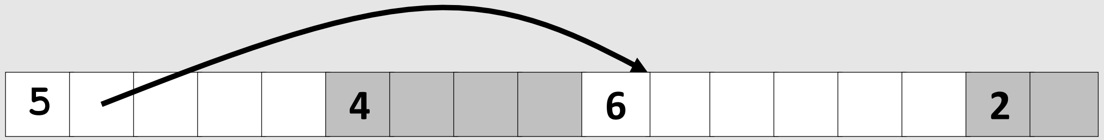
Keeping Track of Free Blocks

- **Method 1: *Implicit free list* using length—links all blocks**

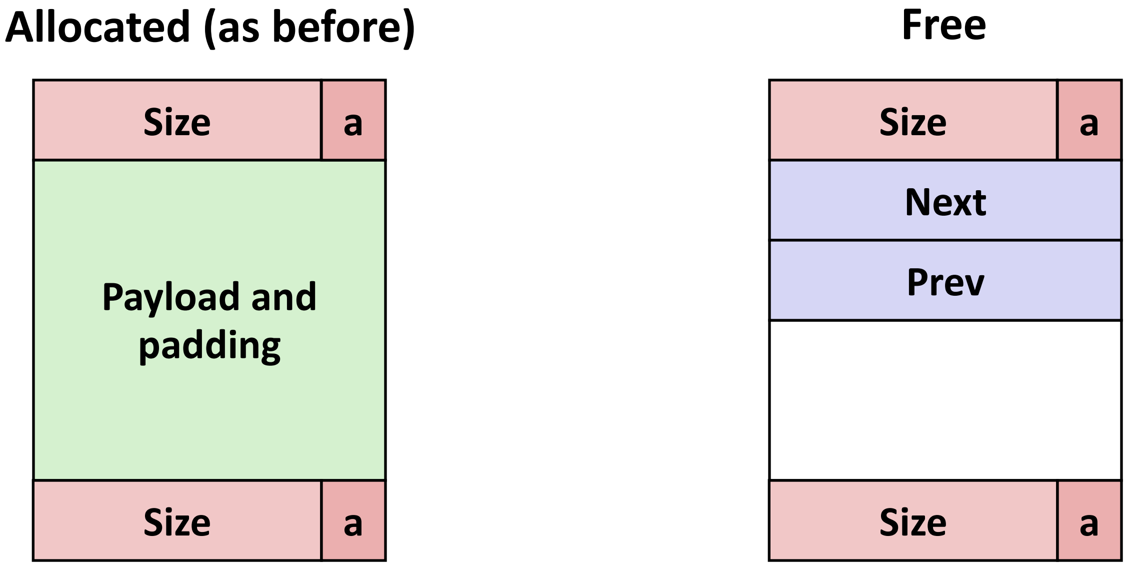


- **Method 2: *Explicit free list* among the free blocks using pointers**



- **Method 3: *Segregated free list***
  - Different free lists for different size classes

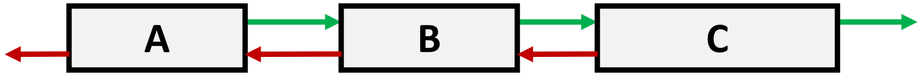
Explicit Free Lists



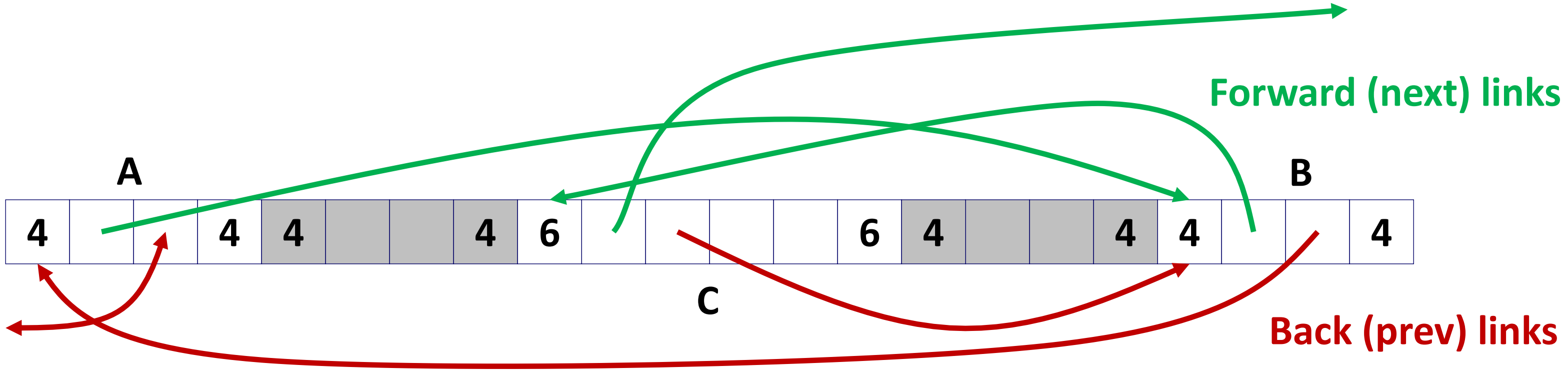
- **Maintain list(s) of *free* blocks, not *all* blocks**
  - The “next” free block could be anywhere
    - So we need to store forward/back pointers, not just sizes
  - Still need boundary tags for coalescing
  - Luckily we track only free blocks, so we can use payload area

Explicit Free Lists

- **Logically:**

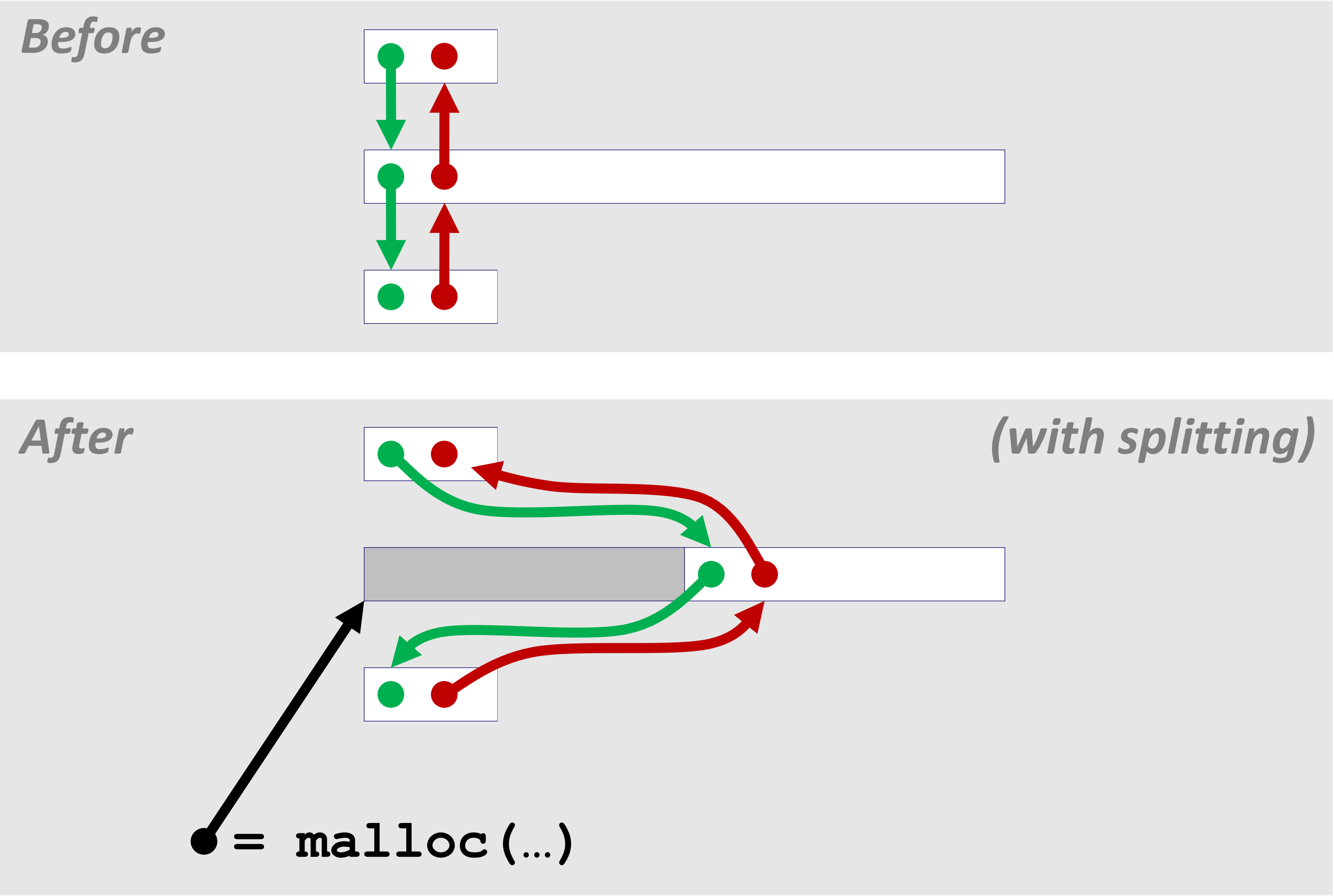


- **Physically: blocks can be in any order**



Allocating From Explicit Free Lists

conceptual graphic



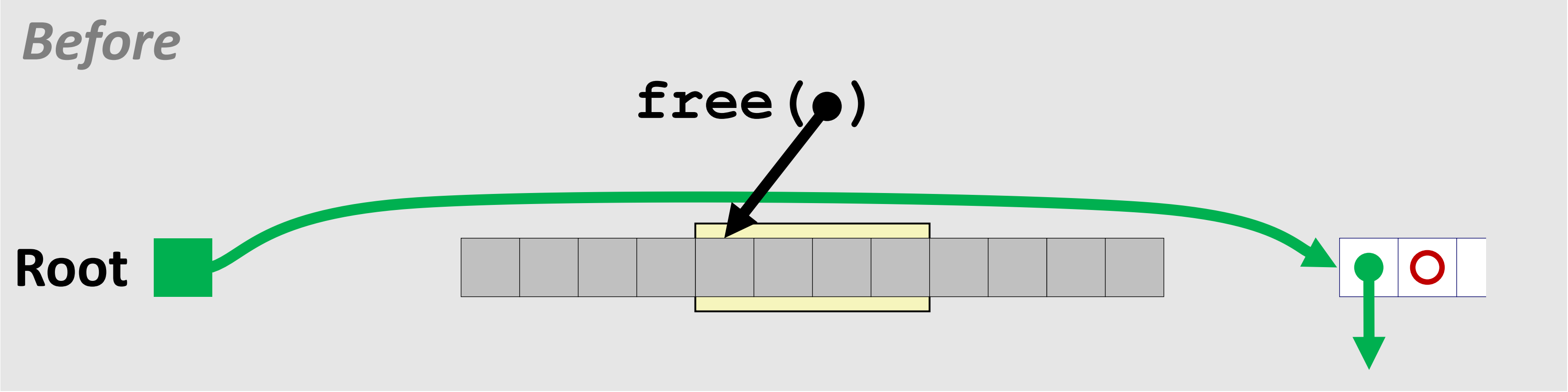
Freeing With Explicit Free Lists

- ***Insertion policy*: Where in the free list do you put a newly freed block?**
  - **LIFO (last-in-first-out) policy**
    - Insert freed block at the beginning of the free list
    - ***Pro:*** simple and constant time
    - ***Con:*** studies suggest fragmentation is worse than address ordered
  - **Address-ordered policy**
    - Insert freed blocks so that free list blocks are always in address order:

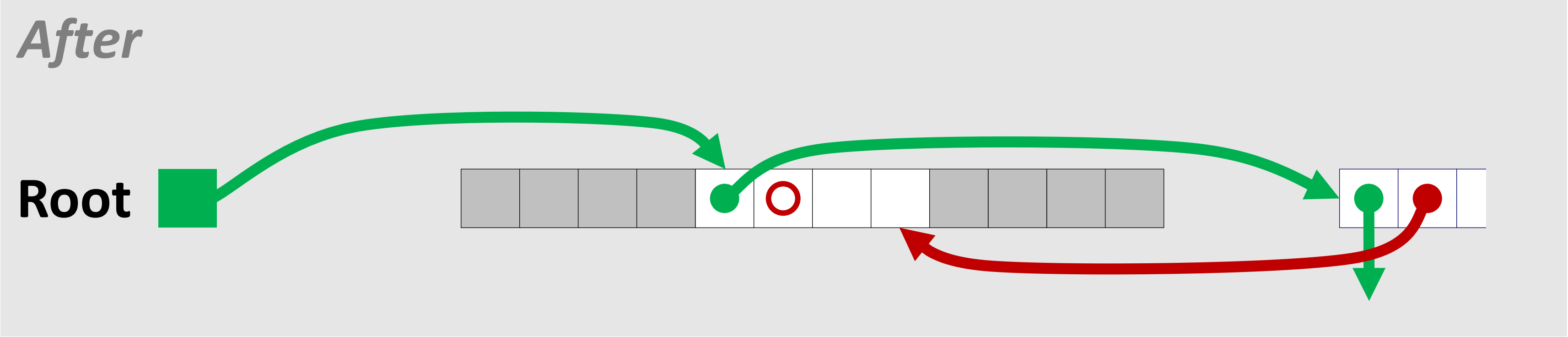
      addr(prev) < addr(curr) < addr(next)
    - ***Con:*** requires search
    - ***Pro:*** studies suggest fragmentation is lower than LIFO

Freeing With a LIFO Policy (Case 1)

conceptual graphic

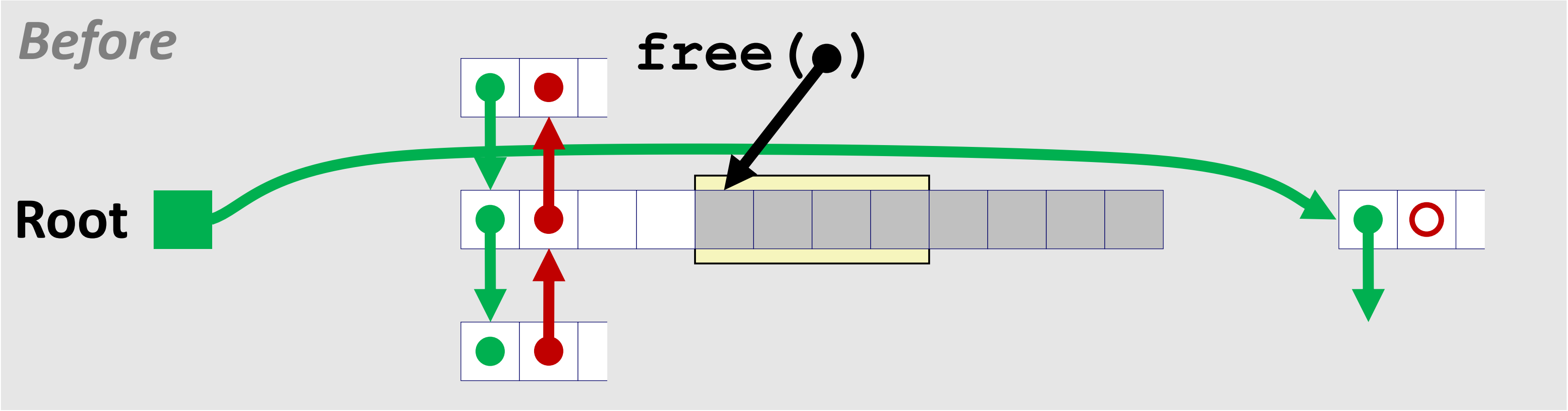


- Insert the freed block at the root of the list

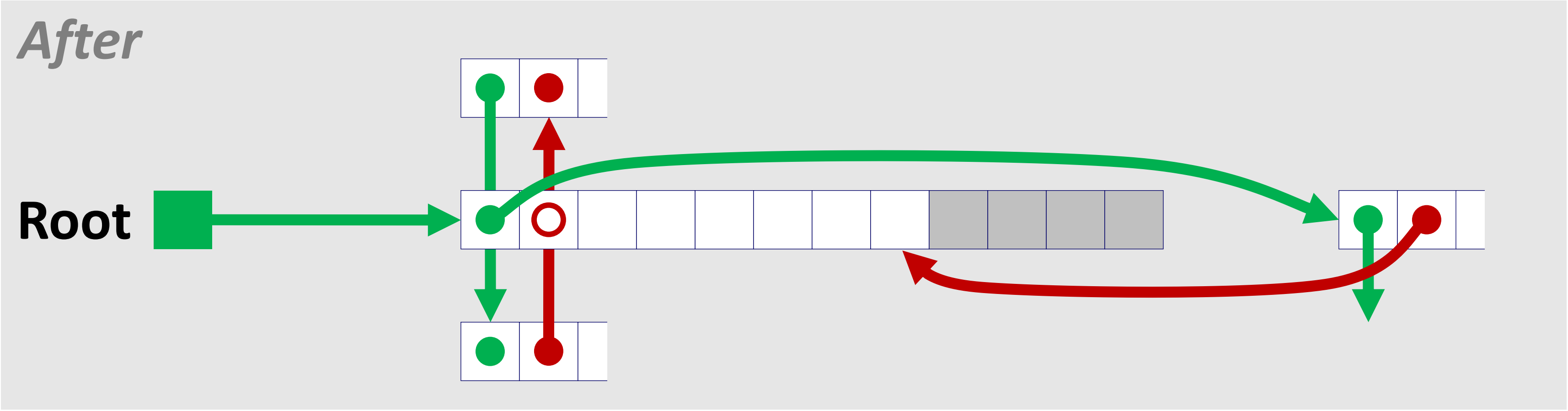


Freeing With a LIFO Policy (Case 2)

conceptual graphic

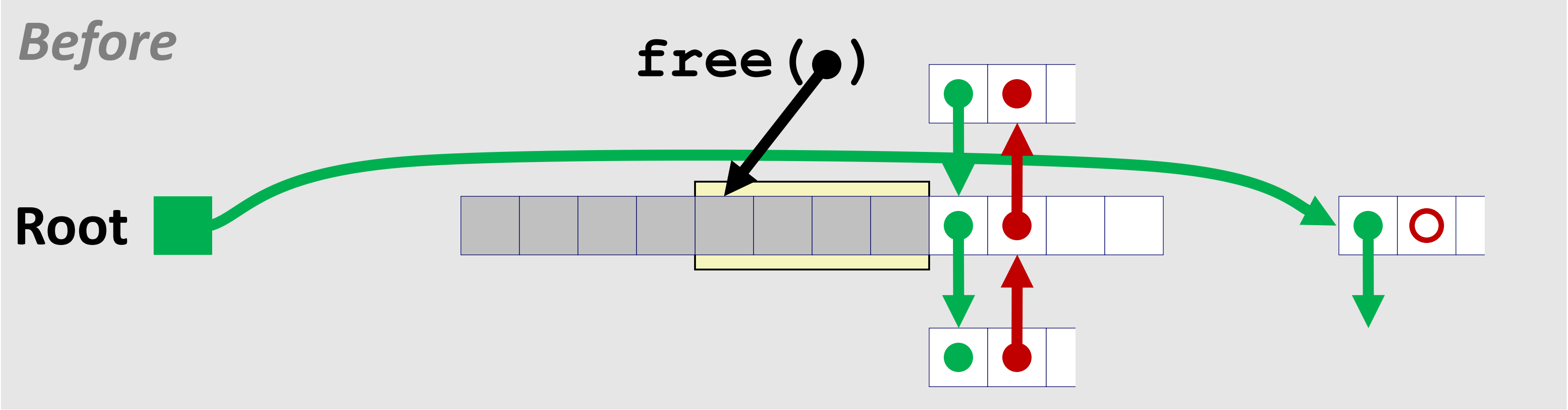


- **Splice out predecessor block, coalesce both memory blocks, and insert the new block at the root of the list**

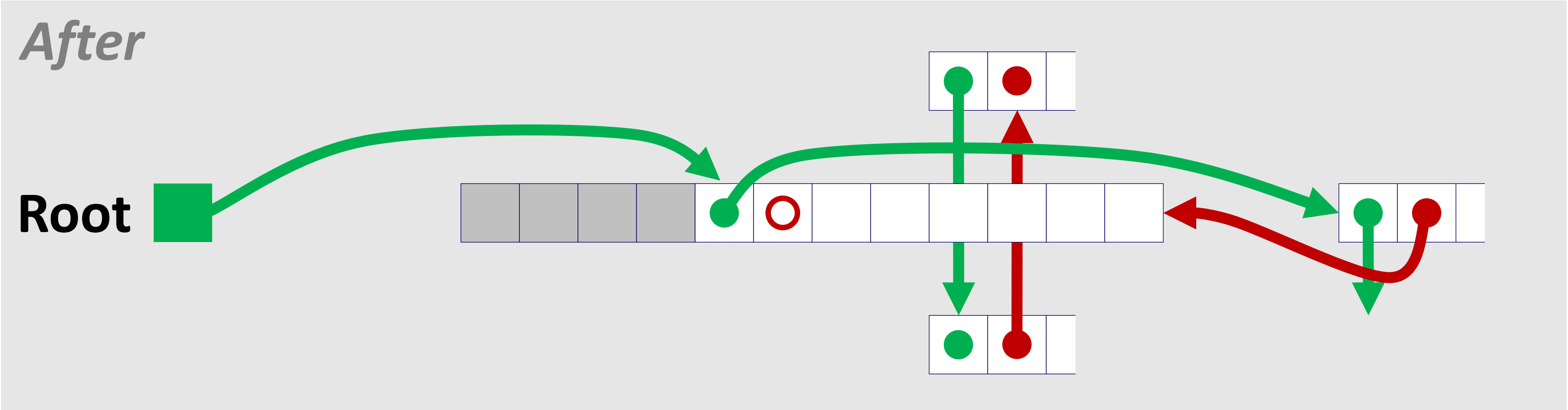


Freeing With a LIFO Policy (Case 3)

conceptual graphic

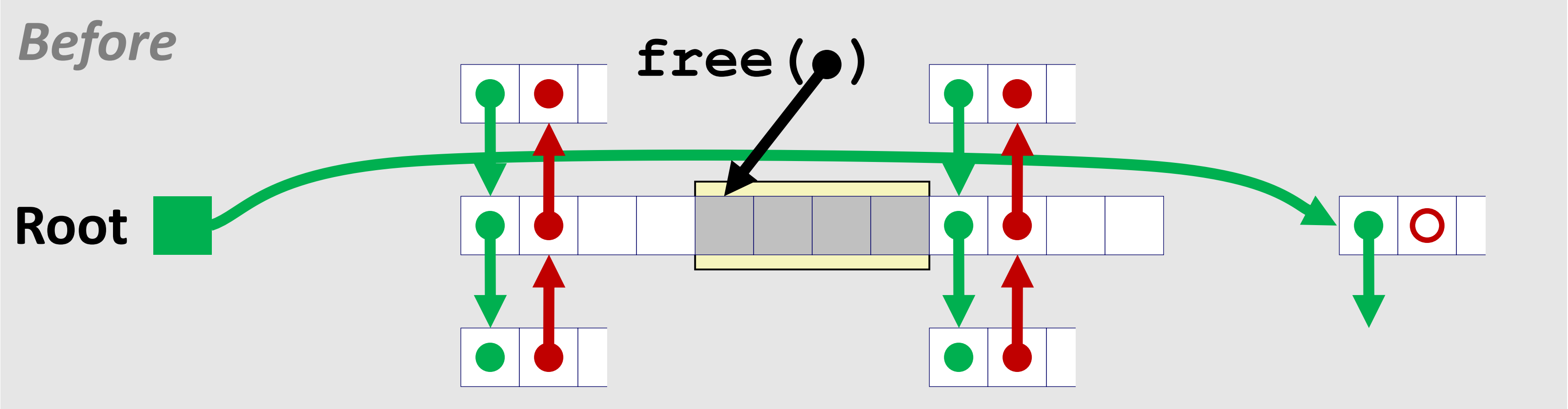


- **Splice out successor block, coalesce both memory blocks and insert the new block at the root of the list**

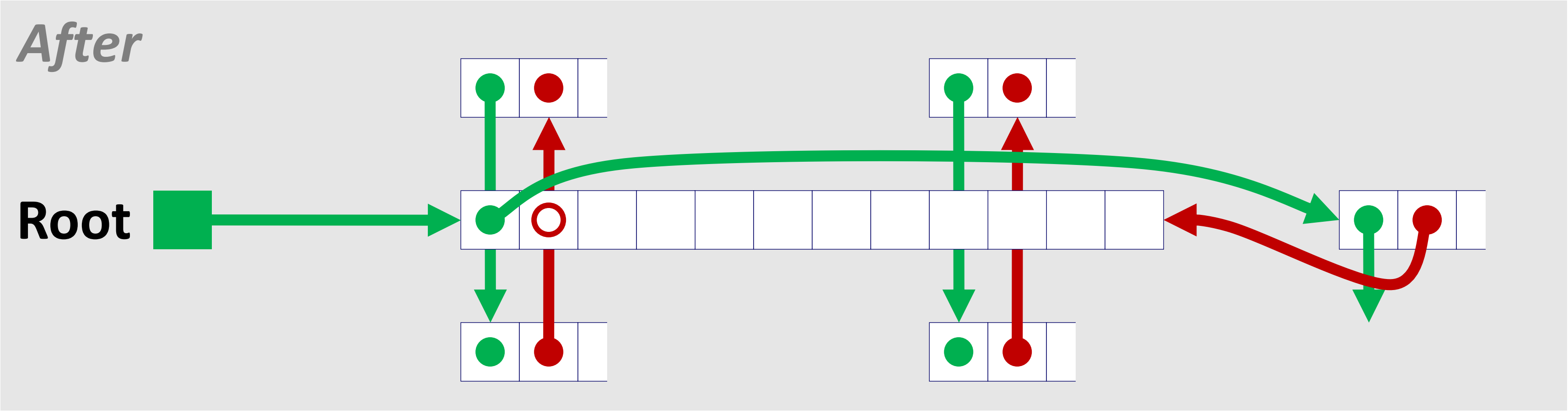


Freeing With a LIFO Policy (Case 4)

conceptual graphic



- **Splice out predecessor and successor blocks, coalesce all 3 memory blocks and insert the new block at the root of the list**



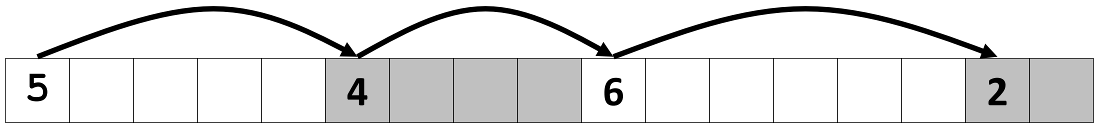
Explicit List Summary

- **Comparison to implicit list:**
  - Allocate is linear time in number of ***free*** blocks instead of ***all*** blocks
    - ***Much faster*** when most of the memory is full
  - Slightly more complicated allocate and free since needs to splice blocks in and out of the list
  - Some extra space for the links (2 extra words needed for each block)
    - Does this increase internal fragmentation?

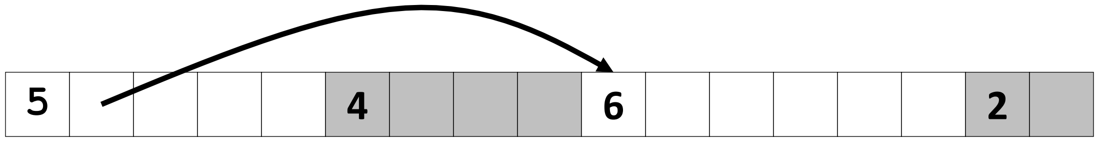
- **Most common use of linked lists is in conjunction with segregated free lists**
  - Keep multiple linked lists of different size classes, or possibly for different types of objects

Keeping Track of Free Blocks

- **Method 1: *Implicit list* using length—links all blocks**



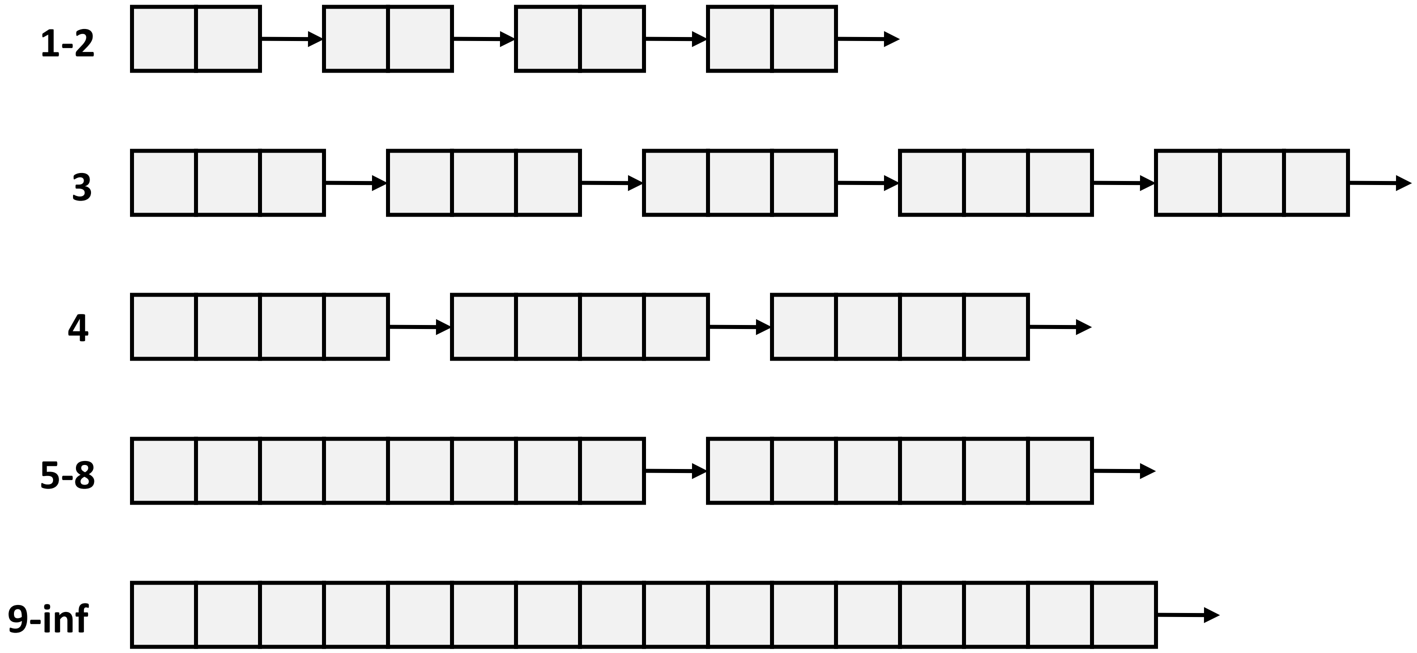
- **Method 2: *Explicit list* among the free blocks using pointers**



- **Method 3: *Segregated free list***
  - Different free lists for different size classes

Segregated List (Seglist) Allocators

- Each ***size class*** of blocks has its own free list



- Often have separate classes for each small size
- For larger sizes: One class for each two-power size

Seglist Allocator

- **Given an array of free lists, each one for some size class**

- **To allocate a block of size *n*:**
  - Search appropriate free list for block of size *m* > *n*
  - If an appropriate block is found:
    - Split block and place fragment on appropriate list (optional)
  - If no block is found, try next larger class
  - Repeat until block is found

- **If no block is found:**
  - Request additional heap memory from OS (using **`sbrk()`**)
  - Allocate block of *n* bytes from this new memory
  - Place remainder as a single free block in largest size class.

Seglist Allocator (cont.)

- **To free a block:**
  - Coalesce and place on appropriate list (optional)

- **Advantages of seglist allocators**
  - Higher throughput
    - log time for power-of-two size classes
  - Better memory utilization
    - First-fit search of segregated free list approximates a best-fit search of entire heap.
    - Extreme case: Giving each block its own size class is equivalent to best-fit.

Module 4 (Part 2)

Summary

- Describe Explicit Free Lists
- Describe boundary tags
- Describe the process of freeing memory with Explicit Free Lists
- Describe Segregated Free Lists