

Improving Run Length Encoding through preprocessing

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Introduction

Basics

Design

Analysis

Implementation

Evaluation and Discussion

Introduction - A Bit of History

- ▶ rise of multimedia
- ▶ rise of the World Wide Web
- ▶ ever increasing data transfer
- ▶ compress to save storage space & to handle new types and volumes of data

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- ▶ massive and rapid increasing data transfer
- ▶ compress to lower transmission cost / time
- ▶ compress to handle increasing resolution, fidelity, dynamic range
- ▶ compression for cold archiving

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Huffman Encoding

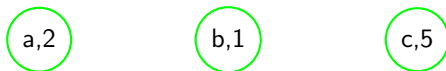


Figure: Example Huffman tree with 3 leaf nodes.

Huffman Encoding

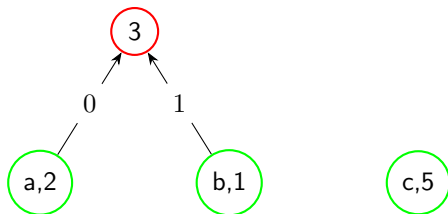


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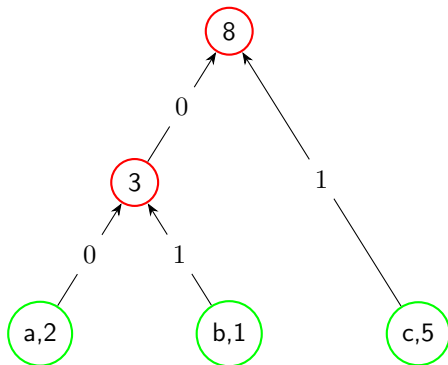


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