

## Worksheet 26 - Group 1

### Worksheet Group 1 Members

Marc Clinedinst: clinedim@onid.oregonstate.edu  
Kelby Faessler: faesslek@onid.oregonstate.edu  
James Fitzwater: fitzwatj@onid.oregonstate.edu  
Tom Gariepy: gariepy@onid.oregonstate.edu  
Sean Reilly: reillys@onid.oregonstate.edu  
Joseph Struth: struthj@onid.oregonstate.edu

### Collaborators

Marc, Kelby, James, Tom, Sean, Joseph

### Worksheet 26: Ordered Bag Using a Sorted Array

In this assignment, we implement several different functions which control the behavior of an ordered bag using a sorted dynamic array. More specifically, we implement the `dynArrAddAt`, `orderedArrayContains`, and `orderedArrayRemove` functions. The `dynArrAddAt` function is by far the most complex of these functions. In any case, these functions are each defined below and are accompanied by comments where needed.

```
/*
```

```
    This function performs adds a value to an array at a particular index. If the value is
    inserted at the very beginning of the array or in the middle of the array, it
    adds the array at the indicated index and shifts the values to the right of that index
    one index to the right. This function begins by checking that the index is in the
    correct range--between 0 and the size of the array inclusive. The function then allocates
    a new chunk of memory for the current or resized capacity of the array. The function
    then copies the values before the given index into the new array, inserts the new value at
    the indicated index, and fills in the remaining values. The function frees the memory for
```

the existing data, and copies over the new data to the dynamic array. Finally, it updates the values for size and capacity.

\*/

```
void dynArrAddAt(struct DynArr *v, int index, TYPE newElement) {
    assert(index >= 0);
    assert(index <= sizeDynArr(v));

    TYPE *newData;
    int currentIndex,
        newCap;

    if (sizeDynArr(v) >= v->capacity) {
        newCap = v->capacity * 2;
    } else {
        newCap = v->capacity;
    }

    newData = malloc(sizeof(TYPE) * newCap);

    for (currentIndex = 0; currentIndex < index; currentIndex++) {
        newData[currentIndex] = v->data[currentIndex];
    }

    newData[index] = newElement;

    for (currentIndex; currentIndex < sizeDynArr(v); currentIndex++) {
        newData[currentIndex + 1] = v->data[currentIndex];
    }

    free(v->data);
```

```

    v->data = newData;
    v->capacity = newCap;
    v->size++;
}

```

```

/*

```

This function returns an integer value representing whether or not a particular target value appears in the dynamic array. It utilizes the binary search algorithm, which terminates in  $O(\log n)$  time; this is significantly faster than the linear search algorithm, which terminates in  $O(n)$  time. The binary search algorithm returns the index where the target value occurs or where the target value should be inserted to keep it sorted. The function compares the value at the index to the target value; if it's a match, the function returns 1 to indicate that the value was found. Otherwise, the value returns 0 to indicate that the target value was not found.

```

*/

```

```

int orderedArrayContains(struct DynArr *v, TYPE target) {
    int index = _binarySearch(v->data, sizeDynArr(v), target);
    return v->data[index] == target;
}

```

```

/*

```

This function removes an integer value from the dynamic array if it appears. The function utilizes the binary search algorithm to locate the index where the value might occur. If the target value appears at that index, then the value is removed and the function terminates. Otherwise, the dynamic array is left unchanged.

```

*/

```

```

void orderedArrayRemove(struct DynArr *v, TYPE target) {
    int index = _binarySearch(v->data, sizeDynArr(v), target);
    if (v->data[index] == target) {
        removeAtDynArr(v, index);
        return;
    }
}

```

```
}  
}
```

In addition to defining the functions above, we also had to provide the complexities for the various implementations of the dynamic array bag functions. A table listing the execution times is below.

|          | Dynamic Array Bag | Linked List Bag | Ordered Array Bag |
|----------|-------------------|-----------------|-------------------|
| Add      | $O(1)+$           | $O(1)$          | $O(n)$            |
| Contains | $O(n)$            | $O(n)$          | $O(\log n)$       |
| Remove   | $O(n)$            | $O(n)$          | $O(n)$            |

Finally, we had to respond to the following questions.

#### Short Answers

1. What is the algorithmic complexity of the binary search algorithm?

The algorithmic complexity of the binary search algorithm is  $O(\log n)$ .

2. Using your answer to the previous question, what is the algorithmic complexity of the method contains for an OrderedArrayBag?

Since the contains method uses the binary search function to locate the potential index for the target value, the contains function also executes in  $O(\log n)$  time.

3. What is the algorithmic complexity of the method addAt?

This function has to loop through every value in an array, so it has  $O(n)$  time.

4. Using your answer to the previous question, what is the algorithmic complexity of the method add for an OrderedArrayBag?

Although this function does use the binary search algorithm, which has  $O(\log n)$  complexity, it also uses the addAt function described above, which has  $O(n)$  complexity.  $O(n)$  dominates  $O(\log n)$ , so the function has  $O(n)$  complexity.

5. What is the complexity of the method removeAt? of the method remove?

The remove function uses the binary search algorithm to locate the index where a value should be removed; this has  $O(\log n)$  complexity. The remove function potentially creates a hole in the array, which must be filled—an  $O(n)$  operation. Since  $O(n)$  dominates  $O(\log n)$ , this function has  $O(n)$  complexity.

### **Piazza Discussion**

<https://piazza.com/class/ib2kus4hsie528?cid=173>