

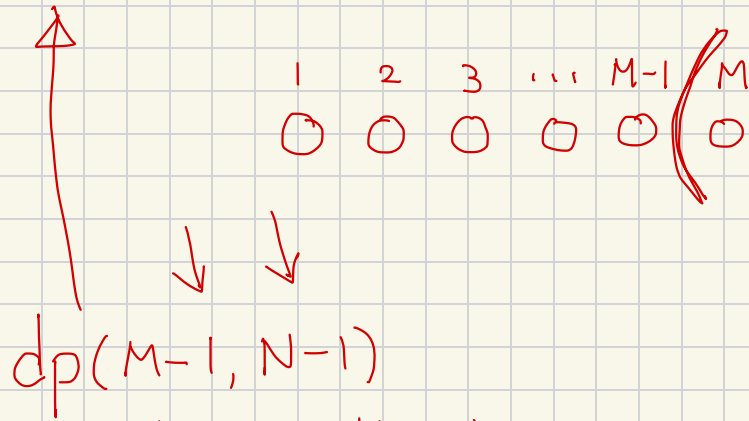
$$M \subset N$$

$M \subset N := M$ 개에서
 N 개 $\frac{N}{2}$ 를 넘는 가짓수.

$N \leq M$
 $\left[\begin{array}{c} \{G, B, R\} \\ \{R, G, B\} \end{array} \right]$
 $\rightarrow M P_N.$

\downarrow \downarrow \downarrow 순서없이
 $\underbrace{\begin{array}{ccc} G & B & R \\ R & G & B \end{array}}_M$ & 중복없이

$dp(M, N) := M \text{개 중 } N \text{개를 순서} \times \text{중복} \times \text{뽑는 가짓수}$



↳ M을 매칭에 포함

$dp(M-1, N)$

↳ M을 매칭에 미포함.

$dp(M, N)$ 을
 $dp(M-1, *)$ 으로
표현해보자!

$$\begin{array}{l}
 \checkmark \quad \boxed{
 \begin{array}{l}
 \cancel{M < N} \rightarrow 0. \\
 M = N \rightarrow 1. \\
 N = 0 \rightarrow 1. \\
 \cancel{N < 0} \rightarrow 0.
 \end{array}
 }
 \end{array}$$

$$M > N$$

$$\begin{array}{c}
 \downarrow \\
 \underline{M-1, N-1} \\
 \downarrow \\
 \underline{M-2, N-2} \\
 \vdots
 \end{array}$$

$$2_0 \rightarrow \textcircled{0}$$

$$\begin{array}{r}
 1 \\
 \hline
 0
 \end{array}
 \bigg/
 \begin{array}{r}
 2 (=M) \\
 \hline
 0
 \end{array}
 \quad N=1.$$

$${}_M C_N$$

$${}_2 C_1 = {}_1 C_0 + {}_1 C_1$$

$$(\boxed{0} \textcircled{0}) \quad (\textcircled{0} \boxed{0})$$

b task.

집한 어떤 두 집은 색이 같지 않도록 N 개의 집에 색칠하는 가짓수 (Easy)

집한 어떤 두 집은 색이 같지 ~~~~~ 가능한 모든 방법의 비용 총합. (Hard)

$$\begin{array}{l} \boxed{} \text{ cost}_1 \\ \boxed{} \text{ cost}_2 \\ \vdots \\ \text{cost}_K \end{array} \left. \vphantom{\begin{array}{l} \boxed{} \text{ cost}_1 \\ \boxed{} \text{ cost}_2 \\ \vdots \\ \text{cost}_K \end{array}} \right\} \begin{array}{l} \min_{i=1}^K \{ \text{cost}_i \} \\ \updownarrow \\ \sum_{i=1}^K \{ \text{cost}_i \} \end{array}$$

$$0 \leq c < 3.$$

$dp(i, c) := i$ 번째 집을 c 로 칠하는 방법의 수

($[1..i]$ 까지만 칠한 상태에서)

$[i+1] \sim [n]$ 까지는 아직 고려 안 함

$dp(i, 0)$

$[1 \sim i]$ 색칠할 때

$\boxed{i} = i$ 번째 집을 빨간색으로 칠한 케이스들에서 최소 비용

$dp(i, 1) = 1 \quad \boxed{i} := i$ 번째 집 파랑

$dp(i, 2) = 2 \quad \boxed{i} := i$ 번째 집 초록

$c_{i,0} = c_i := i$ 번째 집을 칠하는 비용

$c_{i,1} = c_i$

$c_{i,2} = c_i$

$$\boxed{0} = 0.$$

$$\boxed{0} = 0.$$

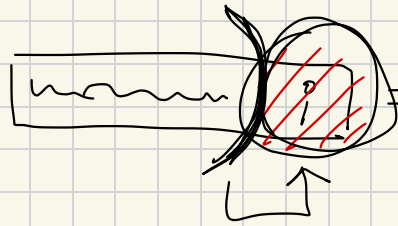
$$\boxed{0} = 0.$$

$$dp(i, 0) = \min \{ dp(i-1, 2), dp(i-1, 1) \} + c_i$$

$$\boxed{i} = \min \{ \boxed{i-1}, \boxed{i-1} \} + c_i$$

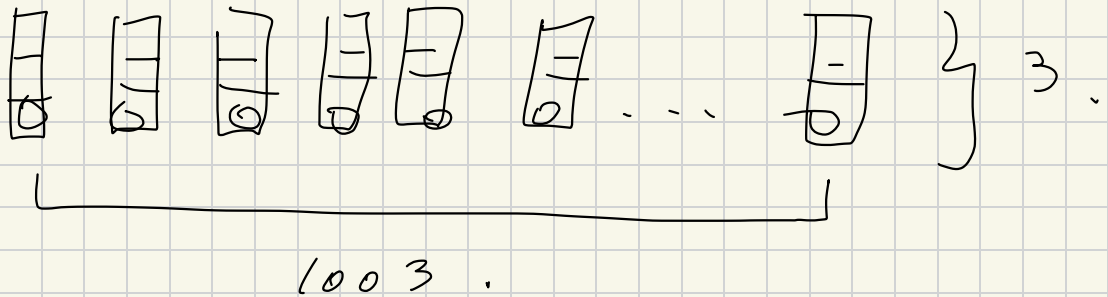
$$\boxed{i} = \min \{ \boxed{i-1}, \boxed{i-1} \} + c_i$$

$dp(i, 0) :=$



\Rightarrow 현재까지 가능한
비용 최소화.

$$\underline{dp(i, 0) = \min \{ dp(i-1, 1), dp(i-1, 2) \} + C_{i,0} .}$$



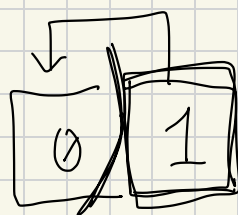
$$\underline{dp(i,1)} := \text{[Diagram: A horizontal rectangle with a wavy line inside, crossed out with blue diagonal lines and a blue 'X' mark.]}$$

$$dp(i,1) = \min \{ dp(i-1,0) + C_{i,1}, dp(i-1,2) + C_{i,1} \}$$

↓

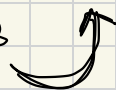
$$ans = \min \{ \underline{dp(n,0)}, \underline{dp(n,1)}, \underline{dp(n,2)} \}$$

solve solve solve,



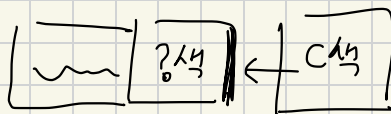
RGB

0



$a-1$

$a.$




```
int best = 1e9;
```

```
for (int j=0; j<3; j++) {
```

```
    if (j!=c) {  $\leftarrow$  dp[i-1][j]
```

```
        best = min(best, solve(i-1, j));
```

```
    }
```

```
}
```

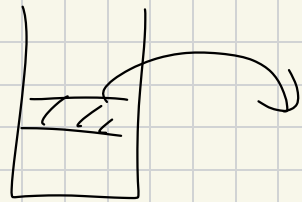
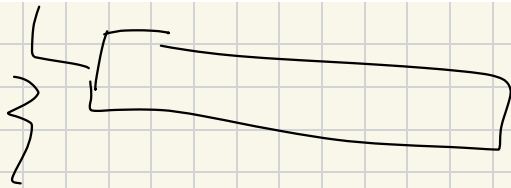
```
return dp[a][c] = best + cost[a][c];
```

~~best = 1e9~~

```
int solve(3, 1), }
```

t=1e9; \rightarrow int best = 1e9; ✓





$\min(\{ \dots \}) \dots \}$

DP [LIS. ✓
- LCS. ✓
- knapsack. ✓]