

$CPA \iff CPA^D$

For traditional cryptosystems

# IND-CPA<sup>D</sup> (Li and Micciancio, 2020)

Why is that important?

- Approximate FHE schemes (like CKKS) leak the noise

$$\begin{cases} Enc_s(m) = (\vec{a}, \langle \vec{a}, \vec{s} \rangle + m + e) \\ Dec_s(\vec{a}, b) = b - \langle \vec{a}, \vec{s} \rangle \end{cases}$$

Toy approximate scheme

- Since we have  $Dec_s(Enc_s(m)) = m + e$ , then  $\mathcal{A}$  can
  1. Query  $\mathcal{O}$  for  $(\vec{a}_i, b_i)$  the encryptions of  $m_i$  and  $\tilde{m}_i$  the decryptions of  $(\vec{a}_i, b_i)$
  2. Solve  $As = b - \tilde{m}$  for  $s$  to recover the secret key