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Exam 1 CSCI 2600 Principles of Software October 6, 2015

- DO NOT OPEN THIS EXAM UNTIL TOLD TO DO SO!
- READ THROUGH THE ENTIRE EXAM BEFORE STARTING TO WORK.
- YOU ARE ALLOWED ONLY 2 "CHEAT" PAGES. NO OTHER MATERIAL IS ALLOWED.

This exam is worth 150 points.

Make sure you have 12 pages counting this one. There are 3 parts, each including multiple questions for a total of 13 questions. If you need more room for an answer than is provided, please use the back of the page and indicate that you have done so. If you re-do a question, please make clear what is your final answer.

Be clear and brief in your explanations—rambling and lengthy answers will be penalized. All questions have short answers.

The test is longer than those back tests, but you can do it! Relax.

The following is for the use of graders

2/40 3/55	
3/55	
TOTAL:/15	50

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Part I. Reasoning About Code

Note: All variables are integers.

Question 1. (5 pts) Order the conditions from strongest to weakest

- (a) y = x + 1 && x is even
- (b) y is odd
- (c) y is odd && x is even
- (d) y = 11 && x = 10

Answer: _ , _ , _ , _

- (a) result = 10
- (b) $5 \le \text{result} \le 10$
- (c) $0 \le \text{result} \le 11$
- (d) $7 \le \text{result} \le 10$

Answer: $_$, $_$, $_$, $_$

Question 2. (10 pts) Let $P \Rightarrow Q \Rightarrow R$, $S \Rightarrow T \Rightarrow U$ and $\{Q\} \text{ code } \{T\}$ be true. Circle the triples that are true.

- $a) \{ P \} code \{ S \}$
- $b) \{ P \} code \{ U \}$
- c) { R } code { S }
- $d) \left\{ \, R \, \right\} \, \textbf{code} \, \left\{ \, U \, \right\}$
- $e) \{Q\} code \{S\}$

Question 3. (5 pts) Compute the weakest precondition using backward reasoning. Fill in all intermediate conditions at the designated places.

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Question 4. (10 pts) Willy Wazoo thinks he can simplify the formula for computing the weakest precondition for if-then-else statements.

Willy claims wp("if (b) S1 else S2", Q) = wp(S1,Q) II wp(S2,Q) should work. Fill in the assignments to z with values of your choice and compute the precondition P according to Willy's formula. Use your choice to explain why Willy is wrong.

```
P: {
    if (x > 0) {
        z = _____
}
    else {
        z = _____
}
Q: { z > 0 }
```

Question 5. (13 pts) Prove that int gcd(int a, int b) computes the correct result if it terminates.

```
// requires: a > 0, b ≥ 0
// returns: the greatest common divisor of a and b.
int gcd(int a, int b) {
   int x = a;
   int y = b;
   int t;
   while (y != 0) {
        t = y;
        y = x%y; // x%y is the remainder in integer division
        x = t;
   }
   return x;
}
```

Question 6. (12 pts) Now prove that int gcd(int a, int b) terminates. Code is the same as in Question 5.

```
// requires: a > 0, b ≥ 0
// returns: the greatest common divisor of a and b.
int gcd(int a, int b) {
   int x = a;
   int y = b;
   int t;
   while (y != 0) {
        t = y;
        y = x%y; // x%y is the remainder in integer division
        x = t;
   }
   return x;
}
```

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Part II. Specifications

Question 7. (5 pts) TRUE/FALSE.

- a) (TRUE/FALSE) If specification A is stronger than specification B, then any implementation that satisfies B satisfies A as well.
- b) (TRUE/FALSE) If $P_A \Rightarrow P_B$ and $Q_B \Rightarrow Q_A$ then spec B is stronger than spec A. (P_A denotes the precondition of A and Q_A denotes the postcondition of A.)

Question 8. (15 pts) Below are 4 specifications of function double sqrt (double x).

Spec A: requires: $x \ge 0$

returns: y such that ly*y - xl < 0.001

Spec B: requires: $x \ge 0$

returns: y such that ly*y - xl < 0.0001

Spec C: returns: y such that $ly^*y - xl < 0.001$ if $x \ge 0$, and 0.0 if x < 0

Spec D: <u>returns</u>: y such that $|y^*y - x| < 0.001$ if $x \ge 0$ <u>throws</u>: IllegalArgumentException if x < 0

For each of the following pairs of specifications, circle the stronger specification, or circle "Neither" if the two specifications are incompatible or equivalent.

- a) A B Neither
- b) A C Neither
- c) A D Neither
- d) B C Neither
- e) B D Neither
- f) C D Neither

Question 9. (20 pts) Below is class ArrayList.	Name:the (simplified) Javadoc specification of the set method from
public E set(int	index, E element)
Replaces the element a	t the specified position in this list with the specified element.
	the element to replace ent to be stored at the specified position
Returns: The element prev	viously at the specified position
<pre>Throws:</pre>	andsException - if the index is out of range (index < 0
a) Convert the Javadoc	specification into a PoS specification:
requires:	
modifies:	
effects:	

b) Now, convert your PoS specification into a logical formula:

returns:

throws:

Name:	
-------	--

Part III. ADTs, Rep invariants, Abstraction functions, Representation exposure

Question 10. (10 pts) TRUE/FALSE

- a) (TRUE/FALSE) If every method in a Java class returns only immutable objects then the class is immutable.
 - b) (TRUE/FALSE) The representation invariant is part of the ADT specification.
 - c) (TRUE/FALSE) The abstraction function is part of the ADT specification.
- d) (TRUE/FALSE) The rep invariant must hold before and after every statement in every method.
 - e) (TRUE/FALSE) The abstraction function maps valid objects to abstract values.

Question 11. (10 pts) Recall the implementation of IntSet.contains we saw in class:

```
// requires: none
// modifies: none
// effects: none
// returns: true if x is in this IntSet, false otherwise
public boolean contains(int x) {
    int i = data.indexOf(x);
    if (i == -1)
        return false;
    // move-to-front optimization
    // speeds up repeated membership tests
    Integer y = data.elementAt(0);
    data.set(0,x);
    data.set(i,y);
}
```

The spec of ArrayList.contains is as follows:

```
// requires: none
// modifies: none
// effects: none
// returns: true if x is in this ArrayList, false otherwise
public boolean contains (Object x)
```

The representation of ArrayList is as follows:

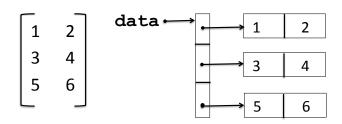
Is it possible to perform a move-to-front optimization in ArrayList.contains? Explain.

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Question 12. (15 pts) Willy Wazoo implemented a matrix with integer elements. (A matrix is a two-dimensional array arranged in rows and columns. An example and a graph elucidating Willy's choice of representation are given on the following page.)

```
public class Matrix {
    private List<List<Integer>> data; // the rep
    private int m; // number of rows
    private int n; // number of columns
    // Rep invariant:
    //returns: number of rows in this Matrix
    public int numRows() {
        return m;
    }
    //returns: number of columns in this Matrix
    public int numColums() {
        return n;
    //requires: 0 <= m,n and elts.size() = m*n and elts has no nulls
    public Matrix(List<Integer> elts, int m, int n) {
        data = new ArrayList<List<Integer>>();
        this.m = m;
        this.n = n;
         // Breaks the elts list into m n-element rows
        for (int i=0; i<m*n; i=i+n) {</pre>
              // Below, subList returns a view of the portion of elts from index i to i+n-1.
              // Modification to the object referred by elts is visible from row, and vice versa.
              List<Integer> row = elts.subList(i,i+n);
             data.add(row);
         }
    }
    //returns: j-th column in this Matrix
    //throws: IndexOutOfBoundsException if j < 0 || j >= numColumns()
    public List<Integer> getColumn(int j) {
        List<Integer> column = new ArrayList<Integer>();
         // Iterates over the m rows to retrieve the j-th element of each row
        for (int i=0; i<m; i++) {
             column.add(data.get(i).get(j));
         }
        return column;
    //returns: i-th row in this Matrix
    //throws: IndexOutOfBoundsException if i < 0 || i >= numRows()
    public List<Integer> getRow(int i) {
        return data.get(i);
    }
}
```

A 3-by-2 matrix and an illustration of Willy's representation of it.



- a) Write the rep invariant for Willy's implementation at the designated space on the previous page.
- b) There is a problem with Willy's implementation. State the problem and suggest a solution.

Name:

Question 13. (20 pts) Consider the IntTree and IntList interfaces specified below.

```
/** An IntTree is a complete binary tree with int nodes with the
 * following property: each node is larger than (or equal to) its
 * children.
 * NOTE: A complete binary tree is a binary tree such that every level
 * except possibly the last is completely filled and all nodes are as far
 * left as possible.
 */
interface IntTree {
  /** Adds val to this tree. Maintains the property. */
  bool insert(int val);
  /** Removes and returns the largest int in this tree.
   * Maintains the property.
   */
  int remove();
}
 * An IntList is a 1-based array list of ints.
 */
interface IntList {
  /** Adds val at the end of this list. */
  void add(int val);
  /** Replaces the element at specified index with val. */
  void set(int index, int val);
  /** Removes the element at specified index, moves any
   * subsequent elements to the left.
  void remove(int index);
  /** Returns the element at specified index */
  int get(int index);
  /** Returns the size of this list */
  int size();
}
Your task is to implement IntTree with an IntList. Below is a start:
     class MyIntTree implements IntTree {
          // Represents the tree
          private IntList rep;
```

Write a rep invariant, abstraction function and sketch the implementation.

- a) (5 pts) Rep invariant
- b) (5 pts) Abstraction function. (You may show your function pictorially.)

c) (10pts) In 1-2 sentences explain how would you implement insert and remove.