Handout Phase 2

Interpretation and Compilation

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Abstract Syntax

EE -> EE **;** EE | EE **:=** EE

| **num** | **id** | **bool** | **let** (**id** = EE)+ **in** EE **end**

| **new** EE | **<!>** EE

| **if** EE **then** EE **else** EE **end**

| **while** EE **do** EE **end**

| EE **binop** EE

| **unop** EE

Concrete Syntax

EM -> E(<;>EM)*	ASTSeq(E1,E2)
E -> EA(< == > EA)?	ASTEq(EA,EA)
EA -> T(<+>EA)*	ASTAdd(E1,E2)
T -> F ((<*>T)*	ASTMul(F,T)
| (<(>AL<)>)*	ASTApply(F,AL)
| <:=> E)	ASTAssign(F,E)

AL -> (EM(<,>EM)*)?

PL -> (id(<,>id)*)?

F -> **num** | **id** | **bool** | **let** (**id** = EM)+ **in** EM **end**

| <(> EM <)>	
| **new** F | <!> F	
| **if EM then EM else EM end**	ASTIf(EM,EM,EM)
| **while EM do EM end**	ASTWhile(EM,EM)

Basic operations

Arithmetic operations (on integer values)

E+E, E-E, E*E, E/E, -E

Relational operations

E==E, E>E, E<E, E<=E, E>=E

Logical operations (on boolean values)

E && E, E || E, ~E

IValues (schematic)

```
interface IValue { /* represents values */
void show();
}

// IValue eval(Environment env) { ... }

//Value constructors
VInt(n)
VBool(t)
VCell(value)
```

IValues (schematic)

```
class VInt implements IValue {
int v;
VInt(int v0) { v = v0; }
int getval() { return v;}
}
```

IValues (schematic)

```
class VCell implements IValue {
IValue v;
VCell(IValue v0) { v = v0; }
IValue get() { return v;}
void set(IValue v0) { v = v0;}
}
```

Interpreter with Dynamic Type Checking (idea)

```
class ASTAdd implements ASTNode {

IValue eval(Environment env) {
v1 = left.eval(env);
if (v1 instanceof  VInt) {
   v2 = right.eval(env)
   if (v2 instanceof VInt) {
        return new Vint((VInt)v1).getval()+((VInt)v2).getval())
}
throw TypeError("illegal arguments to + operator");
}
```

Examples

```
(new 3) := 6;;

let a = new 5 in a := !a + 1; !a end;;

let x = new 10
    s = new 0  in
while !x>0 do
    s := !s + !x ; x := !x – 1
end; !s
end;;
```