

Session Types Course [Exercise Class 1]

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Notation: we use the symbol $\mathbf{0}$ to denote the process **inact** (*because I'm lazy*)

Definition A *prefix* is a process of one of the following forms

$$x[].P \quad x().P \quad x[y].P \quad x(y).P \quad x \triangleright \{\ell_i : P_i\} \quad x \triangleleft \ell_i : P_i$$

A process P is in **canonical form** if $P = (\nu x_1 y_1) \dots (\nu x_n y_n) (P_1 \mid \dots \mid P_m)$ with P_i (called **thread**) a prefix, and x_i or y_i occurring in a P_j for all i and some j in $\{1, \dots, m\}$.

Exercise 1. Prove that each process P is structurally equivalent to a process P' in canonical form.

Definition Let P and Q be processes. We say that P **reduces** to Q if there are P_1, \dots, P_n such that $P \equiv P_1 \rightarrow \dots \rightarrow P_n \equiv Q$ and that P is **reducible** if it reduces to a process Q .

Exercise 2. Check if the following processes reduce to $\mathbf{0}$.

1. $x[u].x[].\mathbf{0} \mid y(v).y[].\mathbf{0}$
2. $(\nu xy)(x[u].x[].\mathbf{0} \mid y(v).y[].\mathbf{0})$
3. $(\nu xy)(x[u].x(w).x[].\mathbf{0} \mid y(v).y[z].y[].\mathbf{0})$
4. $(\nu xy)(x[u].x(w).\mathbf{0} \mid y(v).y[z].\mathbf{0})$
5. $(\nu xy)(x[u].x[].\mathbf{0} \mid y[v].y[].\mathbf{0})$
6. $(\nu xy)(x[u].x(w).x[].\mathbf{0} \mid y(v).y(z).y[].\mathbf{0})$
7. $(\nu xy)(x(a).y[b].y().x[].\mathbf{0})$
8. $(\nu xy)(x(a).x().\mathbf{0})$
9. $(\nu xy)(x[u].x[].\mathbf{0} \mid y(v).y[].\mathbf{0} \mid a[b].a().\mathbf{0} \mid c(d).c[].)$
10. $(\nu x_1 y_1)(\nu x_2 y_2)(x_1[a].x_2(b).\mathbf{0} \mid y_2[c].y_1(d).\mathbf{0})$

Recall: a process is **linear** if each name occurs in at most a thread.

Exercise 3. Let P be a process. Prove that if P is linear, then if $P \rightarrow Q_1$ and $P \rightarrow Q_2$ with Q_1 and Q_2 irreducible, then $Q_1 \equiv Q_2$.

Definition A process P is **typable** if there is a typing derivation of a judgment of the form $\Gamma \vdash P$.

Exercise 4. Which of the processes in Exercise 2 are typable?

Processes			Structural Equivalence (Processes)
P, Q	$:=$	$\mathbf{0}$ inact	$P \mid \mathbf{0} \equiv P$
	$ $	$x[].P$ close	$P \mid Q \equiv Q \mid P$
	$ $	$x().P$ wait	$P \mid (Q \mid R) \equiv (P \mid Q) \mid R$
	$ $	$x[y].P$ send (y through x)	$(\nu x_1 x_2)(\nu y_1 y_2)P \equiv (\nu y_1 y_2)(\nu x_1 x_2)P$
	$ $	$x(y).P$ receive (y on x)	$(\nu xy)P_1 \mid P_2 \equiv (\nu xy)(P_1 \mid P_2)$
	$ $	$(\nu xy)P$ nu	<div style="border: 1px solid black; padding: 2px; display: inline-block;">with $x, y \in \text{fv}(P_2)$</div>
	$ $	$P \mid Q$ parallel	
Operational Semantics (Processes)			
Close:	$(\nu xy)(x[].P \mid y().Q) \rightarrow (\nu xy)(P \mid Q)$		
Com:	$(\nu xy)(a[x].P \mid b(y).Q) \rightarrow (\nu xy)(P \mid Q[a/b])$		
Par:	$P \mid Q \rightarrow P' \mid Q$	if $P \rightarrow P'$	
Res:	$(\nu xy)P \rightarrow (\nu xy)P'$	if $P \rightarrow P'$	
Struct:	$P \rightarrow Q$	if $P \equiv P' \rightarrow Q' \equiv Q$	

Figure 1: Syntax and semantics for processes

Types		Duality (for Types)
T, U	$:=$ close	
	$ $ wait	$\text{close} \perp \text{wait}$
	$ $ $(T) \blacktriangleleft U$	$(T) \blacktriangleleft U \perp [T] \blacktriangleleft V$ if $U \perp V$
	$ $ $[T] \blacktriangleleft U$	
Typing Rules		
$\text{T-Inact} \frac{}{\vdash \mathbf{0}}$	$\text{T-Par} \frac{\Gamma \vdash P \quad \Delta \vdash Q}{\Gamma, \Delta \vdash P \mid Q}$	$\text{T-Resr} \frac{\Gamma, x : T, y : T^\perp \vdash P \quad T \perp U}{\Gamma \vdash (\nu xy)P}$
	$\text{T-Close} \frac{\Gamma \vdash P}{\Gamma, x : \text{close} \vdash x[].P}$	$\text{T-Wait} \frac{\Gamma \vdash P}{\Gamma, x : \text{wait} \vdash x().P}$
	$\text{T-Send} \frac{\Gamma, x : U, y : T}{\Gamma, x : (T) \blacktriangleleft U \vdash x[y].P}$	$\text{T-Recv} \frac{\Gamma, x : U}{\Gamma, x : (T) \blacktriangleleft U \vdash x(y).P}$

Figure 2: Types