

MUTUAL EXCLUSION WITHIN RING TOPOLOGY (1)

Simple method:

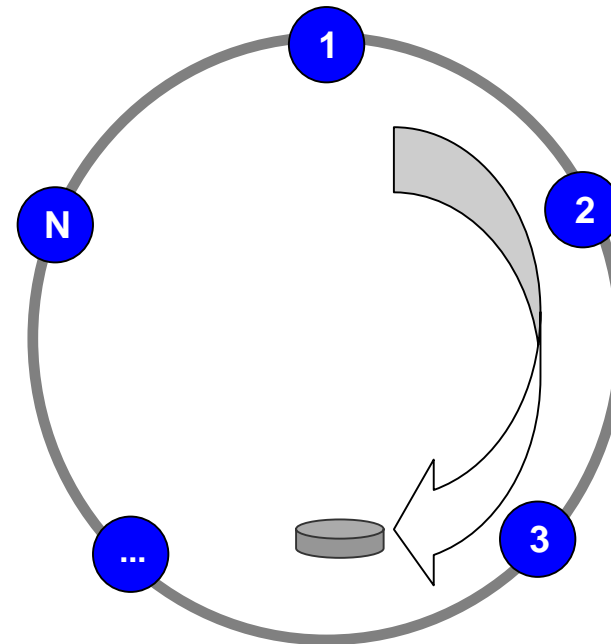
- token passing

Failure model:

- transient link failures

Possible defects:

- token loss



Ping-Pong algorithm (Misra 1983)

System model:

- $P_0 \dots P_{N-1} \rightarrow N$ unknown
- knowledge of P_i 's neighbor identities not required
- uniform algorithm (no distinguished process)
- unreliable FIFO channels

Two tokens:

- *ping* – circulating **privilege**, granting entry to the critical section
- *pong* – circulating after ping
- aim of each is to detect a loss of the other

Assumption:

- only one token gets lost in one round

MUTUAL EXCLUSION WITHIN RING TOPOLOGY (3)

Initialization:

- $\#ping := +1$
- $\#pong := -1$
- $m_i := 0$

m_i is used by P_i to remember the last token passed through

Invariant:

- $\#ping + \#pong = 0$

MUTUAL EXCLUSION WITHIN RING TOPOLOGY (4)

P_i :

```
when received PING(#ping) do
    if  $m_i == \#ping$  then    // pong is lost, regenerate it
        regenerate()

when received PONG(#pong) do
    . . .    // similar to above

when meeting PING and PONG do    // new incarnation of tokens
    incarnate()

when sending PING(#ping) forward do
     $m_i := \#ping$ 

when sending PONG(#pong) do
    . . .    // similar to above
```

MUTUAL EXCLUSION WITHIN RING TOPOLOGY (5)

P_i :

```
when received PING(#ping) do
    if  $m_i == \#ping$  then      // pong is lost, regenerate it
        regenerate(#ping)

when received PONG(#pong) do
    . . .                    // similar to above

when meeting PING and PONG do // new incarnation of tokens
    incarnate(#ping)

when sending PING(#ping) forward do
     $m_i := \#ping$ 

when sending PONG(#pong) do
    . . .                    // similar to above
```

MUTUAL EXCLUSION WITHIN RING TOPOLOGY (6)

P_i :

. . .

regenerate(x) // restore lost PING or (and) PONG

```
{  
    #ping := |x|  
    #pong := -#ping  
}
```

incarnate(x) // increment absolute value of PING and PONG

```
{  
    #ping := |x| + 1  
    #pong := -#ping  
}
```

Values of #ping, #pong

- values can be bounded (e.g. for comparisons)
- their absolute values may increase at each P_i
- it is sufficient to increment them modulo $N+1$

Channel requirements:

- we need FIFO channels
- what would happen if tokens could overpass each other in channel ...?
- ... it will be considered as a loss of the other token – regeneration
- effect : 2 PINGs (or PONGs)
- can we detect such inappropriate regeneration and delete unnecessary token ?
- which one ?

MUTUAL EXCLUSION WITHIN RING TOPOLOGY (8)

Old token deletion:



```
when received PING(#ping) do
    if |#ping| < |mi| then
        // some old PING has arrived - delete it
    else if mi == #ping then
        {
            . . .
```

(be aware that it can happen that somewhere in the ring the old ping comes before the new ping)

Problems:

- What about the mutual exclusion in this case?
- What about bounding modulo $N+1$?
- How can we protect the system from loosing 2 tokens?

BIBLIOGRAPHY

- [1]  Jayadev Misra, "Detecting Termination of Distributed Computations Using Markers"
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Montreal, Canada, August 17-19, 1983.
- [2]  Michel Raynal, "Distributed Algorithms and Protocols", John Wiley & Sons, 1988.