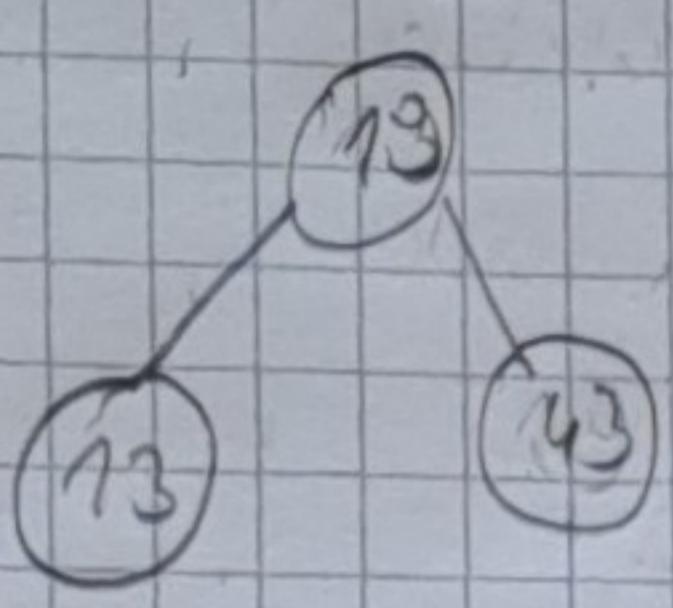
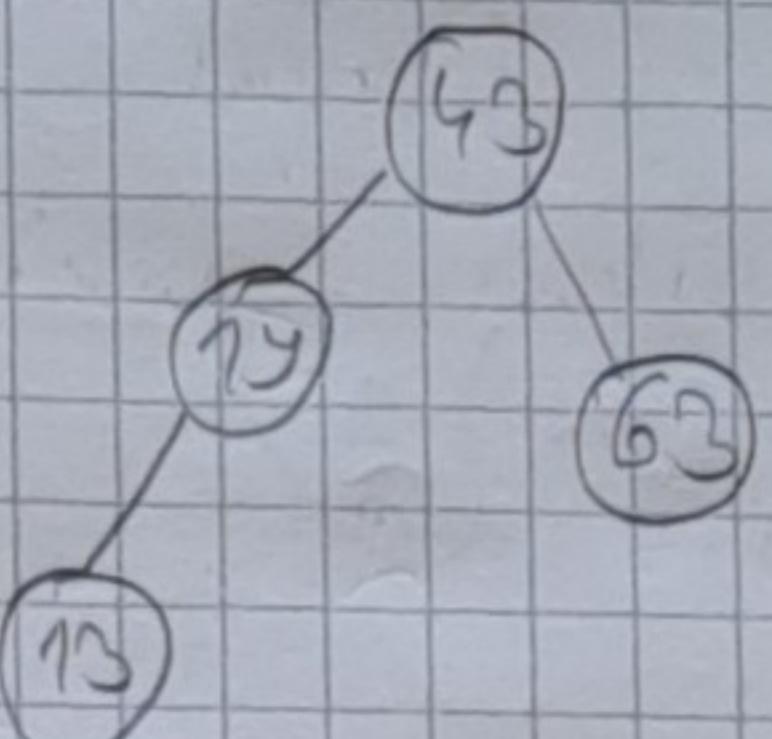


doosxán 3.

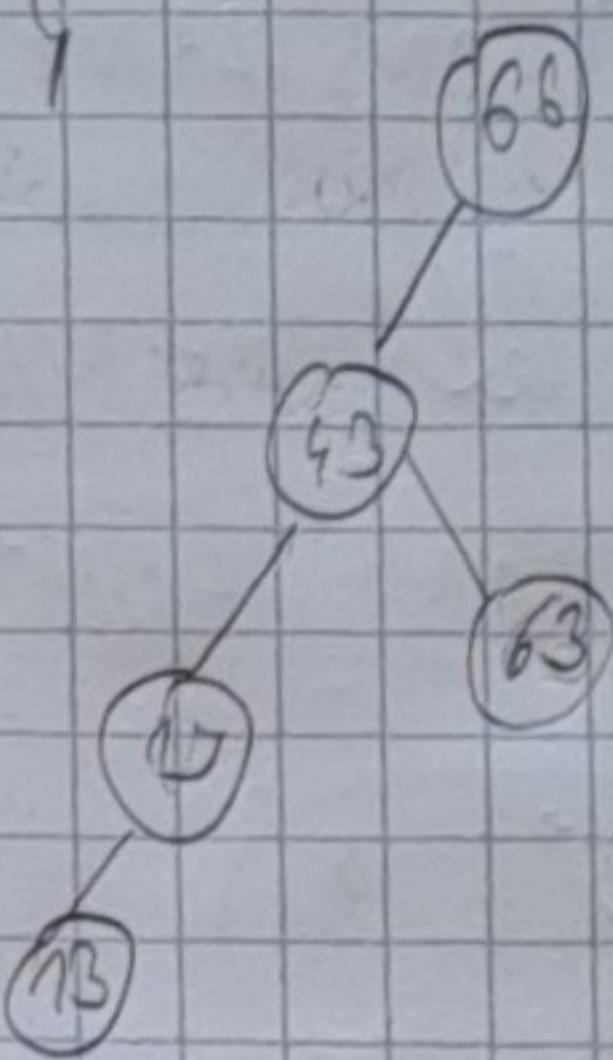
h = 2



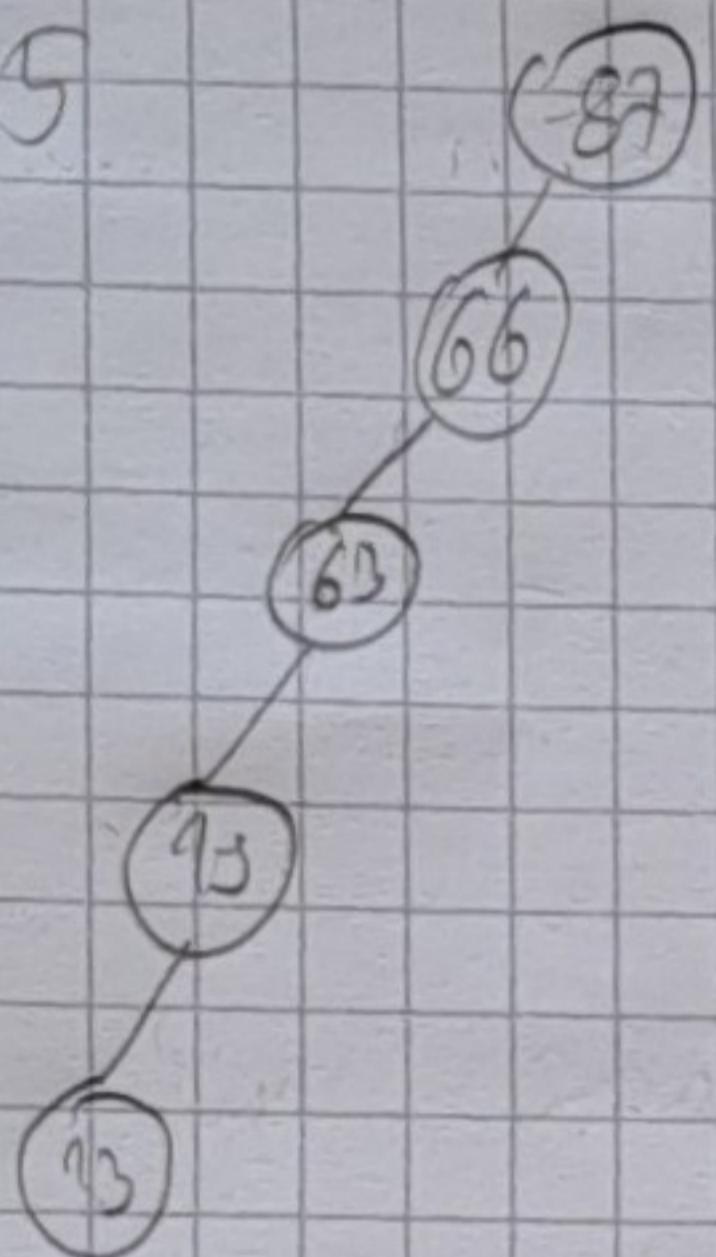
h = 3



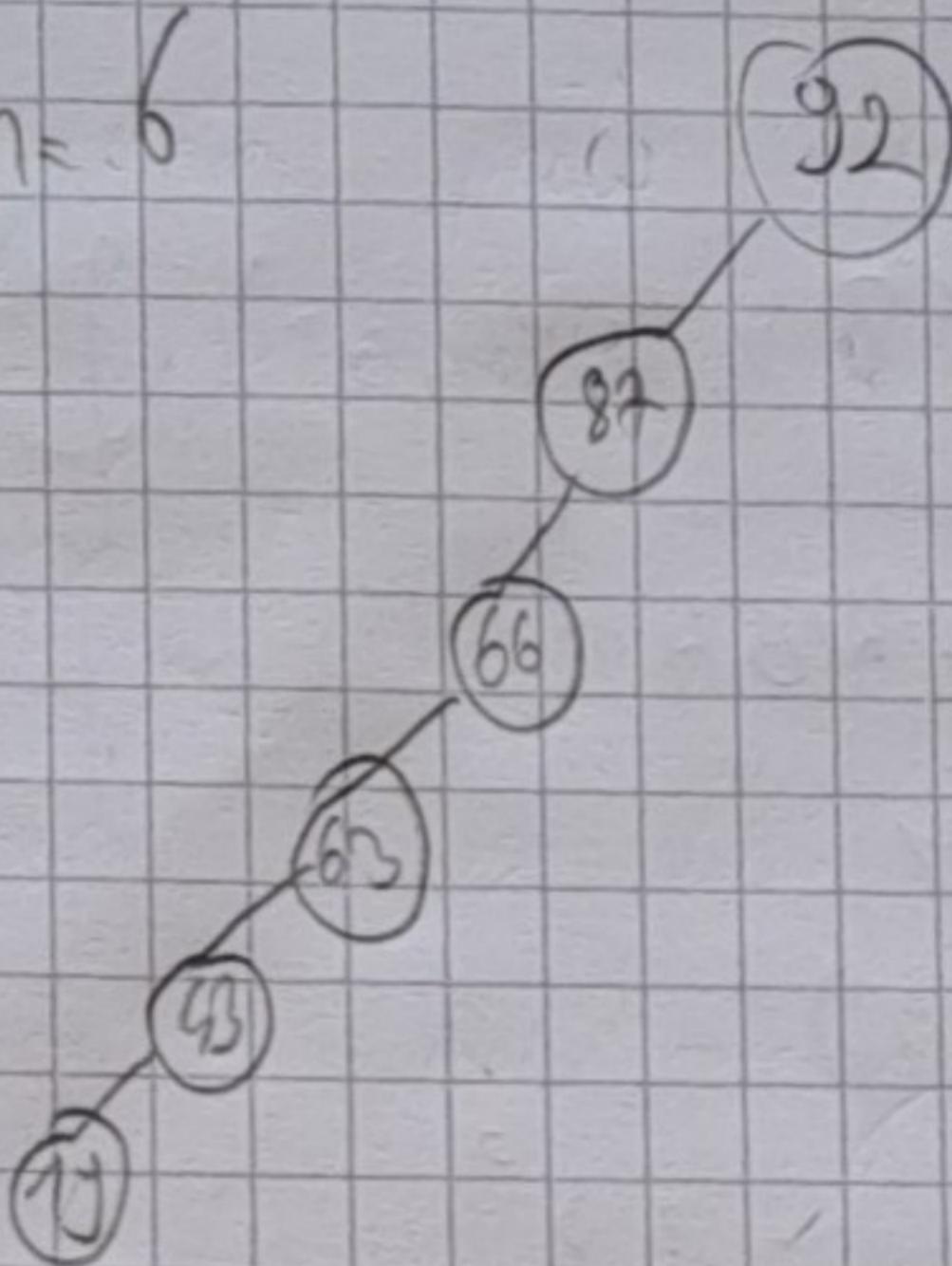
h = 4



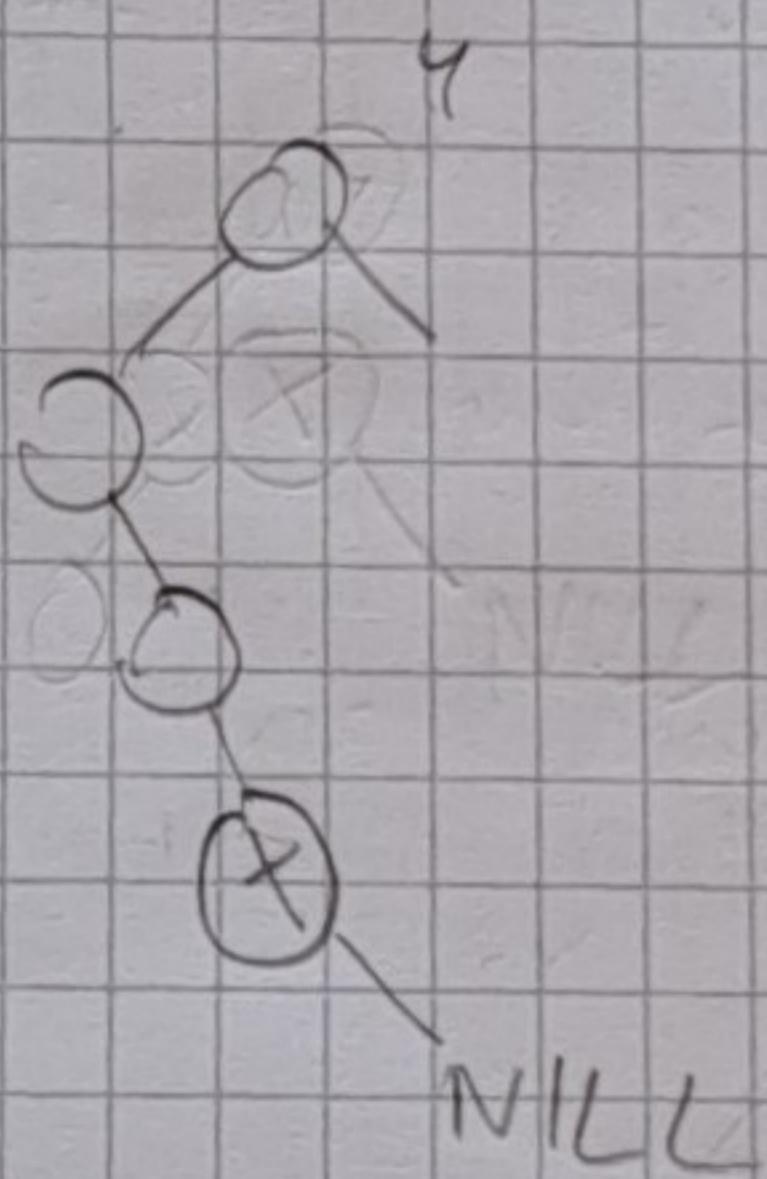
h = 5



h = 6



2.



Y mon 6it. predok so x. Aho. y nje predok so x  
pokazuj 2 kji me malji vlastu, oponovo mona vyzad  
x < 2 < y, a y nje predok so x.

Promostrov u Y. last more but predekt est X, ji obojek.  
biu, ono k i g. right biu predekt od X implicans  
X>Y. Pretpostrov no u myoj moynosti predekt od X  
tj. lyiu slyite ji tercetes predekt od X. Neho je  
2 tij moynosti predekt. Tvoje 2 more but u lyiu  
podstallu od Y, rts. mylians 2<Y, rts. ji  
kontroliyia 3 Pretpostrov oboji u Y objednani od X.

3. Y ji molety od X, rts. oboji oboji X musi but  
it lyiu it deino obite.

Abu ji X oboji slyite, ono k oboji oboji ji X  
myelling od Y, obojno u ja prethwint od

X (Y ji moynost modern X), rts. ji trudnyi 2.

No ji X lyiu slyite, ono k oboji oboji X

esthewint od Y, obojno Y ji esthewint od X (Y ji  
moynosti bgu modern X), rts. ji trudnyi 1.

Nic Nic  
X Y  
→ Pots se moli o lita obojno  
obu ji on moynosti pa vajok  
obu ji Y objednani od X.

4. *Nelw ji* X *dwor* na *wyjem* mu *Rosie*  
tree-successor 4. k-tu *wyedbeni* sol K. *Moldovian*  
ta dwu *dwo* ji *moyre* 2 h( tree *wuzdwo* oblic  
lyui , derwe *ritwne* myrie 20 *syalj* wru h+h)  
ta ta ji potdw O(2k+2h) < O(k+h)

Also ist das hier ein  
einfaches Beispiel für die  
Implementation eines  
Prototypen.

*Imomia* 2 magnusnot.

1. Bib twala kudu our in slow mo - mble hui  
mij Prestens reserv te, mij Guyé, oomme pnum  
tre - numin who a oldel man appedwil

2. Als worden we overleg stemmen met behulp van  
„achteren“ nu neigen modders i fryske re slak me oosteren  
de modders hja nu lje te.

Also down now in m<sup>th</sup> row, one in m-1<sup>th</sup> row  
now dry m-1 row tree - mulrow

$$\therefore O(m-1) = O(m)$$

6. Inset ; remove | sw pohor vt. put do  
āuso, two sw ; insert , remove se poci  
vitemi emporium. Przeciwny āuso ujednij grom  
obwolten, o to j' taj kuj w trój mowu āuz.

### TREE-INSERT(T,z)

1.  $x = T.\text{root}$
2.  $y = \text{NIL}$
3. while  $x \neq \text{NIL}$
4.  $y = x$
5. if  $z.\text{key} < x.\text{key}$
6.    $x = x.\text{left}$
7. else  $x = x.\text{right}$
8.  $z.\text{p} = y$
9. if  $y = \text{NIL}$
10.    $T.\text{root} = z$
11. else if  $z.\text{key} < y.\text{key}$
12.    $y.\text{left} = z$
13. else  $y.\text{right} = z$ .

### TREE-INSERT(T,z)

wintui o 6. Geolitter

### INORDER-WALK(T,x)

- i)  $x = \text{NIL}$   
return

INORDER-WALK(x.left+)

Print x,

INORDER-WALK(x.right)

## WORST CASE

Majority always be do good also insertion happens in  
 worst case possible. Tudo be worse because, whenever  
 we insert into now  $g_i$  in. Stabbing be tido power  
 has limited time. While  $p[i]$  be inserted in puts.  
 No doings, convex, const move. Best  $O(n^2)$   
 so when inserting now put in puts below it.  
 worst case. This resultant Algo will  $\Theta(n^2)$ .

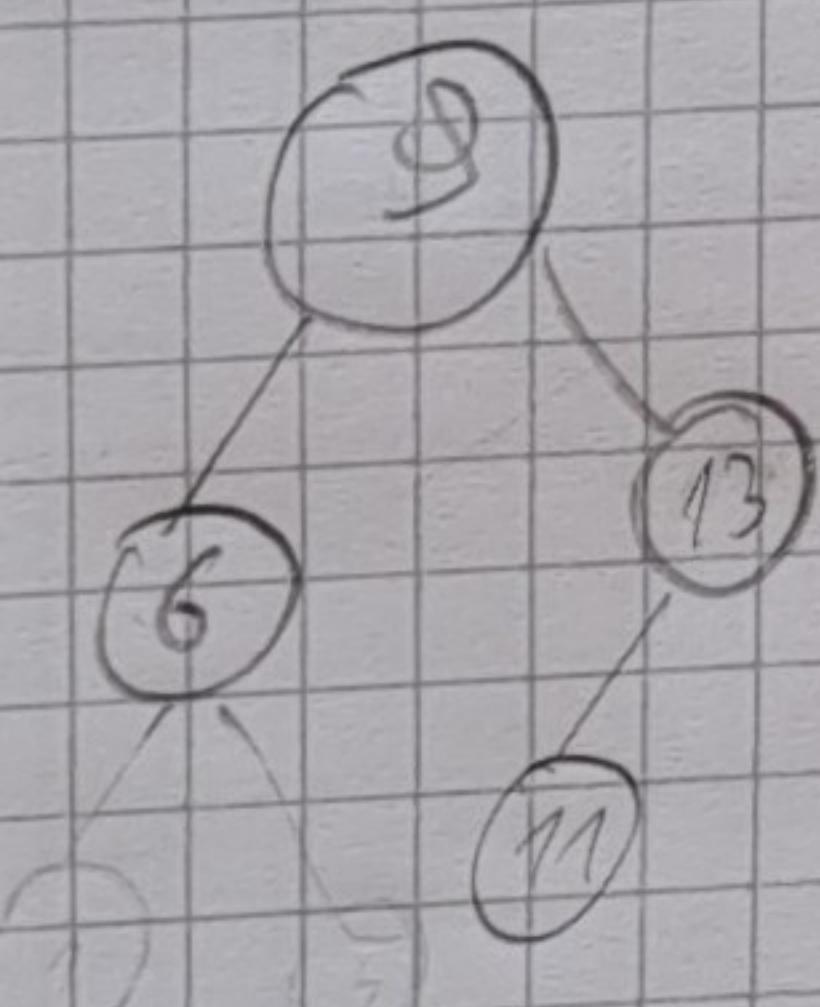
## BEST CASE

Majority always be just have amemo to determine which stabbing  
 adams, adams, visum stabbing  $g_i$  by. While  $p[i]$   
 works well so  $\Theta(n)$  it's resulting  $\Theta(n \log n)$ .

## 8. Organizing Elements BEST & WORST CASE

Npm.

BRISANDO 3 PA 6 MMAD



BRISANDO 6 PA 3 MMAD

