## Homework 4

# $6{\rm CCS3CFL}$ - Compilers & Formal Languages

## Finley Warman

## November 1, 2020

## Contents

Question 1	2
Question 2	2
Question 3	2
Question 4	3
Question 5 (Deleted)	3
Question 6	3
Question 7	3
Question 8	4
Question 9	4
Question 9	4

#### Question 1

Q: Give *all* the values and indicate which one is POSIX for how these expressions can recognise strings.

```
A: (ab + a) . (1 + b) matching ab.

• Sequ( Left(Sequ(Chr(a),Chr(b))), Left(Empty) ) (POSIX)

• Sequ( Right(Chr(a)), Right(Chr(b)) )

A: (aa + a)* matching aaa.

• Stars([ Left(Sequ(Chr(a), Chr(a))), Right(Chr(a)) ]) (POSIX)

• Stars([ Right(Chr(a)) ])

• Stars([ Left(Sequ(Chr(a), Chr(a))) ])
```

#### Question 2

Q: If a regular expression r does not contain any occurrence of  $\mathbf{0}$ , is it possible for L(r) to be empty?

Assuming that an expression can only be defined in terms of a valid alphabet (i.e. there is no 'unmatchable character'), and negation is not allowed, then L(r) cannot be empty. This is because excluding  ${\bf 0}$ , all expressions are defined recursively in terms of either  ${\bf 1}$  or the empty string, and if either of these is accepted then the language cannot be empty. If negation is allowed, then an example of an empty language would be  ${\bf r}$  .  ${\bf r}$  ., or the union of the language of  ${\bf r}$ , and its complement. (Which is always empty)

## Question 3

DIGIT

Q: Define tokens for a language with numbers, parenthesis, and operations. Can the given strings be lexed?

A: Token / Expression Defs:

= RANGE("0123456789")

```
START_DIGIT = RANGE("123456789")

NUMBER = DIGIT + (START_DIGIT . DIGIT*)

LPAREN = CHAR('(')
RPAREN = CHAR(')')

OPERATOR = CHAR('+') + CHAR('-') + CHAR('*')

LOWERCASE = RANGE("abcdefghijklmnopqrstuvwxyz")

ID = LOWERCASE . LOWERCASE*

LANG = ("num":NUMBER) + ("op":OPERATOR) + ("lp":LPAREN) + ("rp":RPAREN) + ("id":I
```

```
• (a+3)*b = YES, lp:(, id:a, op:+, num:3, rparen:), op:*, id:b
```

- )()++-33 = YES, rp:), lp:(, rp:), op:+, op:+, op:-, num:33
- (a/3)\*3 = NO, / is not an accepted token.

### Question 4

Q: Assuming r is nullable, show that 1+r+r.r == r.r holds.

```
1+r+r.r as a proper tree:
    = (((1+r)+r).r)
since r is nullable, (1+r) == r
    = (((r)+r).r) = ((r+r).r)
since (r+r) == r:
    = (r.r) = r.r
therefore
    1+r+r.r == r.r
```

### Question 5 (Deleted)

#### Question 6

```
A:

SEQ(
SEQ(CHAR('/'), CHAR('*')),
SEQ(
STAR(NOT(SEQ(CHAR('*'), CHAR('/')))),
SEQ(CHAR('*'), CHAR('/'))))
```

#### Question 7

Q: Simplify the given expression. Does simplification always preserve the meaning of a regular expression?

A:

```
Simplifying (0.(b.c))+((0.c)+1):
Using 0.r = 0:
    (0) + ((0)+1)
= 0 + (0 + 1)
Using 0+r = r:
    0 + (1)
= 0 + 1
= 1
```

The regex produced by simplification will be equivalent to its pre-simplified form, in that they will accept the same language.

However, the resulting expression may vary in *how* it matches a string, and thus (unless steps are taken to rectify this), the returned matching value may be different.

#### Question 8

Q: What is mkeps for the following expressions? A:

```
• (0.(b.c))+((0.c)+1) = Right(Right(Empty))

• (a+1).(1+1) = Sequ(Right(Empty), Left(Empty))

• a* = Stars(Nil)
```

### Question 9

Q: What is the purpose of the record regular expression in the Sulzmann & Lu Algorithm?

A: When tokenizing an expression (e.g. splitting into its component words), the record expression is used for classifying these tokens.

e.g. when lexing a block of code, we can produce a resulting expression of records which label each (notable) sub-expression with the token they matched.

### Question 9

Q: Define recursive functions atmostempty, somechars, infinitestrings. (Recalling nullable and zeroable).

```
somechars -
```

somechars(0): false
somechars(1): false
somechars(c): true

somechars(r1+r2): somechars(r1) || somechars(r2)
somechars(r1.r2): somechars(r1) || somechars(r2)

somechars(r\*): somechars(r)

#### infinitestrings -

infinitestrings(0): false
infinitestrings(1): false
infinitestrings(c): false

infinitestrings(r1+r2): infinitestrings(r1)  $\mid \mid$  infinitestrings(r2) infinitestrings(r1.r2): infinitestrings(r1)  $\mid \mid$  infinitestrings(r2)