# Module 2.5: Two Basic Proof Patterns MCIT Online - CIT592 - Professor Val Tannen

LECTURE NOTES



#### Proof of an "if-then" statement

Recall the statement

If m + n is even then m - n is even.

Logical structure:  $even(m+n) \Rightarrow even(m-n)$ .

What did we do?

We assumed the premise even(m + n).

Then,  $m+n=2\ell$ , for some integer  $\ell$  (by definition of "even").

Then,  $m = 2\ell - n$ .

Then,  $m - n = (2\ell - n) - n = 2\ell - 2n = 2(\ell - n)$ .

Then, we satisfied the definition of "even" (by taking  $k = \ell - n$ ).

We concluded even(m-n).

## The proof pattern for implication

You wish to prove  $P_1 \Rightarrow P_2$ 

#### Proof pattern.

assert the **premise**  $P_1$ 

(then derive/infer)

...logical/mathematical consequences ...

(until you can)

assert the **conclusion**  $P_2$ 

With all this you have proven  $P_1 \Rightarrow P_2$ .

## A proof with cases I

Recall the statement

If  $p = r \cdot s$  and p is prime, then one of r and s equals 1 and the other one equals p.

Logical structure:

$$(p = r \cdot s) \land prime(p) \Rightarrow (r = 1 \land s = p) \lor (s = 1 \land r = p).$$

What did we do to prove this one? To begin with, we have an implication.

We assumed the premise  $(p = r \cdot s) \land prime(p)$ 

Then,  $r \mid p$ 

Then, since p is prime, r = 1 or r = p.

Then, we proceeded in **two cases**.

## A proof with cases II

Because r = 1 or r = p we can continue in two cases.

In the first case we assume r = 1.

Therefore  $p = 1 \cdot s$ .

And thus s = p.

Hence,  $(r = 1 \land s = p) \lor (s = 1 \land r = p)$ .

In the second case we assume r = p.

Therefore  $p = p \cdot s$ .

And thus 1 = s.

Hence,  $(r = 1 \land s = p) \lor (s = 1 \land r = p)$ .

In both cases we have concluded  $(r = 1 \land s = p) \lor (s = 1 \land r = p)$ .

#### The by-cases proof pattern

Assuming  $P_1 \vee P_2$  you wish to prove  $P_3$ .

#### Proof pattern.

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assert P_1 \vee P_2
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**Case 1.** assert  $P_1$ .

... logical/mathematical consequences ...

assert  $P_3$ 

Case 2. assert  $P_2$ .

... logical/mathematical consequences ...

assert  $P_3$ 

Since in both cases we obtained  $P_3$ , we have proved it assuming  $P_1 \vee P_2$ .

## Some observations about by-cases

- 1. It generalizes easily to more than two cases. If we start from a disjunction of k statements, then we will have k cases.
- 2. The cases need not be *mutually exclusive*, as they were (almost) in our example:  $(r = 1) \lor (r = p)$ . We will give examples later in the course.
- 3. The disjunction that yields the cases need not appear as part of the assumptions in the original statement. In fact  $(r = 1) \lor (r = p)$  did not. You can see a more striking example in another segment in this module.