

# Unit 1

## Proof

*Albert Sung*

# Tentative Teaching Plan for Part 2

- ❑ Instructor: Dr. Albert SUNG (replace Prof. Tommy Chow until further notice)
  - Office: G6354 (YEUNG)
  - Email: [albert.sung@cityu.edu.hk](mailto:albert.sung@cityu.edu.hk)
- ❑ One or two assignments
  - Due date for Assign 1: **Nov 2 (Tue) or later (TBC)**
- ❑ **Test: Nov 20** (Sat. of Week 12)
  - about one hour within the period **9:00 – 11:00 am**
  - Venues: **LT-5** and **LT-6**

# Outline of Unit 1

- 1.1 Why Proofs?
- 1.2 Direct Proofs
- 1.3 Indirect Proofs
- 1.4 Mathematical Induction

# Unit 1.1

## Why Proofs?

# What is a Proof?

- ❑ A proof is a **valid argument** that establishes the truth of a statement.
  - If the statement is about mathematical objects (integers, triangles, sets, etc.), then it is a mathematical proof.
- ❑ In mathematical proofs,
  - more than one rule of inference are often used in a step,
  - steps may be skipped, and
  - the rules of inference may not be explicitly stated.

# The Pigeonhole Principle

- ❑ Suppose that you have  $n$  pigeonholes.
- ❑ Suppose that you have  $m$  pigeons, where  $m > n$ .
- ❑ If you put the  $m$  pigeons into the  $n$  pigeonholes, some pigeonhole will have more than one pigeon in it.



- $n = 9$  pigeonholes
- $m = 10$  pigeons
- Some pigeonhole has more than one pigeon.

Is it true?

# Do We Need Proofs?

Mathematics consists in proving the most obvious thing in the least obvious way.



George Polya,  
a Hungarian  
mathematician

How about  
engineers?

**True love  
doesn't need  
proof.  
The eyes  
told what  
heart felt.**



Toba Beta,  
an Indonesian  
poet.

# Should EE Students Learn Proofs?

## ❑ My personal opinion:

Engineering students should learn to **discover**, **understand**, and **enjoy** proofs.



## ❑ Why?

- A way to convince oneself and others that a proposed engineering solution indeed works.
  - Network protocols, cryptographic protocols, database management, optimality of a (hardware/software) system, etc.
- A sign of understanding.
  - Problem solving relies on deep understanding of a problem.
- An intellectual challenge full of fun.
- An art for appreciation.



# Terminology

## ❑ **Definition**

- a precise description of a mathematical term (e.g., odd number).

## ❑ **Axiom**

- A statement assumed true without proof.
- Axioms form a basic building block from which all theorems are proved.

## ❑ **Theorem**

- a mathematical statement that is proved to be true using rigorous reasoning (i.e., rules of inference).

## ❑ **Lemma**

- a minor result whose purpose is to help in proving a theorem.
  - Very occasionally, some lemmas are very important on their own.

## ❑ **Corollary**

- a result whose (usually short) proof follows directly from a theorem.

# Forms of Theorems

- Many theorems assert that a property holds for all elements in a domain, such as the integers, the real numbers, the triangles, the sets.
  - The universal quantifier is, however, often omitted.

## □ Example:

“If  $x > y$ , where  $x$  and  $y$  are positive real numbers, then  $x^2 > y^2$ .”

can be written as the following universal statement:

“ $\forall x, y \in R_+$ , if  $x > y$ , then  $x^2 > y^2$ .”

# Unit 1.2

## Direct Proofs

# Direct Proofs

- ❑ A way of showing the truth of a statement by using established facts (e.g. definition, lemmas, theorems), rules of inference, and logical equivalences.
- ❑ Proving Existential Statements
  - Proof by **example**
- ❑ Proving Universal Statements
  - Proof by **exhaustion** (also called proof by cases)
  - Proof by **UG**

# Proving Existential Statements

- Consider an existential statement

$$\exists x \in D, Q(x).$$

Proof by example

Find an  $x$  in  $D$  that makes  $Q(x)$  true.

- Validity follows from Existential Generalization (EG).

# Disproving Universal Statements

- Consider a universal statement

$$\forall x \in D, Q(x).$$

- That it is false is equivalent to that its negation is true.

$$\exists x \in D, \sim Q(x).$$

Proof by counter-example

Find an  $x$  in  $D$  that makes  $Q(x)$  false.

# One Example is Enough

- It is easy to find an example to prove that

There exists positive integers  $a, b, c$  such that  
$$a^2 + b^2 = c^2.$$

- Euler's conjecture (1769):

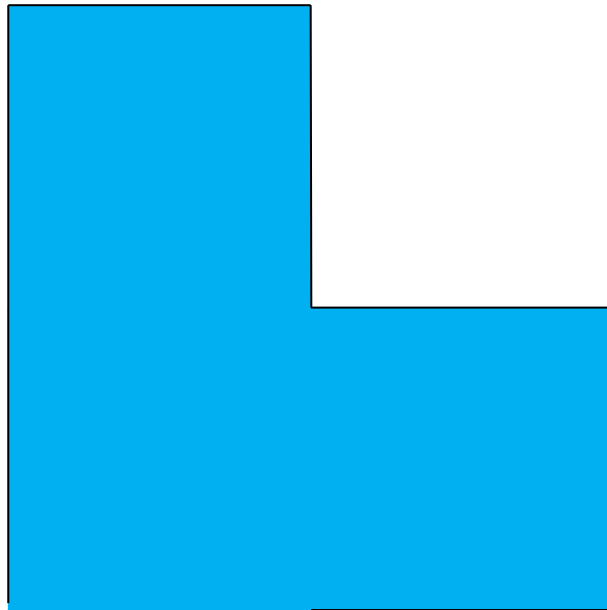
There does not exist positive integers  $a, b, c, d$  such that  
$$a^4 + b^4 + c^4 = d^4.$$

- It is disproved in 1986 by a counter-example:

$$2862440^4 + 15365639^4 + 18796760^4 = 20615673^4.$$

# Cutting Figures

- ❑ Congruent pieces: of the same shape and size, possibly rotated or flipped over.
- ❑ Prove that this figure can be cut into 2 congruent pieces.



*Too easy? How  
about cutting into  
4 congruent pieces?*



# Proving Universal Statements

- Consider a universal statement

$$\forall x \in D, Q(x).$$

- Proof by **Exhaustion** (also called Proof by **Cases**)

- 1) Split the domain  $D$  into a finite number of cases (i.e. subsets).
- 2) Check that the statement is true for each case (i.e.  $Q(x)$  for all  $x$  in each subset.)

- Proof by **Universal Generalization (UG)**

- 1) **Arbitrarily** pick an element  $x$  in  $D$ .
- 2) Show that  $x$  has the property  $Q$ .

## Two Examples (Proof by Exhaustion)

1. Prove that  $x^2 \leq 16$  for  $1 \leq x \leq 4$ ,  $x$  is an integer.

Solution:

○  $1^2 = 1 \leq 16, 2^2 = 4 \leq 16, 3^2 = 9 \leq 16, 4^2 = 16 \leq 16.$

*Q.E.D.*

2. Prove that  $\min(x, y) \leq \max(x, y)$ , where  $x, y \in R$ .

Solution:

○ Case 1:  $x \leq y$ . Then  $\min(x, y) = x \leq y = \max(x, y)$ .

○ Case 2:  $x > y$ . Then  $\min(x, y) = y \leq x = \max(x, y)$ .

*Q.E.D.*

# Even and Odd Integers

- ❑ Before we give an example to explain the proof method based on UG, we need the following:
- ❑ Definition
  - The integer  $n$  is **even** if there exists an integer  $k$  such that  $n = 2k$ , and
  - $n$  is **odd** if there exists an integer  $k$ , such that  $n = 2k + 1$ .
- Note that every integer is either even or odd and no integer is both even and odd.

# Example (Proof by UG)

**Theorem:** *The sum of any two even numbers is even.*

**Proof:** Suppose  $m$  and  $n$  are (*arbitrarily chosen*) even numbers. By the definition of even numbers,  $m = 2r$  and  $n = 2s$  for some integers  $r$  and  $s$ .

$$\begin{aligned} m + n &= 2r + 2s && \text{by substitution} \\ &= 2(r + s) && \text{by factoring out a 2.} \end{aligned}$$

Let  $t = r + s$ . Then

$$m + n = 2t \quad \text{where } t \text{ is an integer.}$$

Therefore,  $m + n$  is even.

Hence, the sum of **any** two even numbers is even.

*Q.E.D.*

# Unit 1.3

## Indirect Proofs

# Indirect Proofs (2 Major Types)

## Proof by Contradiction

- ❑ Also called **reductio ad absurdum**
  - (i.e., Reduction to the Absurd)
- ❑ Classic: Used in Socratic method (~400 BC)
  - By asking questions, Socrates revealed contradictions in other people's belief, showing that the belief is false.

## Proof by Contraposition

- ❑ Based on the logical equivalence between a conditional and its contrapositive.
  - See Unit 2.

$$p \rightarrow q \equiv \sim q \rightarrow \sim p$$

# Proof by Contradiction

To prove that  $p$  is true:

1. Assume that  $p$  is **false**.
2. With the above assumption, show that there is a **contradiction**.
3. Conclude that  $p$  is **true**.

**Contradiction rule:**

$$\sim p \rightarrow c$$

---

$$p$$

where  $c$  is a contradiction.

Why does it work?

# Why is Contradiction Rule Valid?

## □ By truth table

premises			conclusion
$p$	$\sim p$	$c$	$\sim p \rightarrow c$
T	F	F	T
F	T	F	F

There is only one critical row in which the premise is true, and in this row the conclusion is also true. Hence this form of argument is valid.

## □ By showing that it is a tautology

$$(\sim p \rightarrow c) \rightarrow p \equiv (p \vee c) \rightarrow p$$

$$\equiv p \rightarrow p$$

$$\equiv \sim p \vee p$$

$$\equiv \mathbf{t}$$



## Example (Proof by Contradiction)

**Theorem:** *There is no greatest integer.*

**Proof:** We prove by contradiction. Suppose there is a greatest integer  $N$ . Then  $N \geq k$  for all integer  $k$ .

Let  $M = N + 1$ . Now  $M$  is an integer and  $M > N$ .

Therefore,  $N$  is not a greatest integer.

We have reached a contradiction.

Hence, the statement is true.

*Q.E.D.*

# Proof by Contraposition

□ This method is based on

$$p \rightarrow q \equiv \sim q \rightarrow \sim p$$

To prove that  $p \rightarrow q$  is true:

1. Assume  $\sim q$  is true.
  2. Show that  $\sim p$  is true.
  3. Conclude that  $p \rightarrow q$ .
- } This shows that  $\sim q \rightarrow \sim p$  is true.

## Example (Proof by Contraposition)

**Theorem:** *For all integer  $n$ , if  $n^2$  is even, then  $n$  is even.*

**Proof:** Suppose  $n$  is not even. Then  $n = 2k + 1$  for some integer  $k$ .

$$n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1$$

Let  $t = 2k^2 + 2k$ , which is an integer.

Then  $n^2 = 2t + 1$ .

Therefore,  $n^2$  is odd (i.e., not even).

Hence, the statement is proved.

*Q.E.D.*

# If-and-Only-If Proof

- “ $P$  if and only if  $Q$ ” (or simply  $P$  iff  $Q$ ) can be split up into the two parts:
  - 1) The “only if” part:  $P \rightarrow Q$
  - 2) The “if” part:  $Q \rightarrow P$
- Each part is usually proved separately.

□ Let  $E$  denote the equation  $x^2 + px + q = 0$ . Prove that  $E$  has two distinct real roots iff  $p^2 - 4q > 0$ .

***Solution:***

1) (if part) If  $p^2 - 4q > 0$ , by the quadratic formula, there are two distinct roots:  $\frac{-p + \sqrt{p^2 - 4q}}{2}$  and  $\frac{-p - \sqrt{p^2 - 4q}}{2}$ .

2) (only if part) Suppose  $E$  has two distinct real roots.

Denote them by  $\alpha$  and  $\beta$ , where  $\alpha \neq \beta$ . Then,

$$x^2 + px + q = (x - \alpha)(x - \beta) = x^2 - (\alpha + \beta)x + \alpha\beta.$$

Comparing coefficients,  $p = -(\alpha + \beta)$  and  $q = \alpha\beta$ .

$$\text{Thus, } p^2 - 4q = (\alpha + \beta)^2 - 4\alpha\beta = (\alpha - \beta)^2 > 0.$$

*Q.E.D.*

# The Pigeonhole Principle (revisited)

- There are  $m$  pigeons and  $n$  pigeonholes, where  $m > n$ .
- Some pigeonhole will have more than one pigeon.



**Theorem:** Let  $m$  objects be distributed into  $n$  bins. If  $m > n$ , then some bin contains more than one object.

**Theorem:** Let  $m$  objects be distributed into  $n$  bins. If  $m > n$ , then some bin contains more than one object.

**Proof:**

Assume that every bin contains no more than one object. We want to prove  $m \leq n$ . (proof by contraposition)

Let  $x_i$  be the number of objects in bin  $i$ .

By assumption,  $x_i \leq 1$ .

Since  $m$  is the number of objects, we have

$$m = \sum_{i=1}^n x_i \leq \sum_{i=1}^n 1 = n.$$

Hence,  $m \leq n$ , as required.

*Q.E.D.*

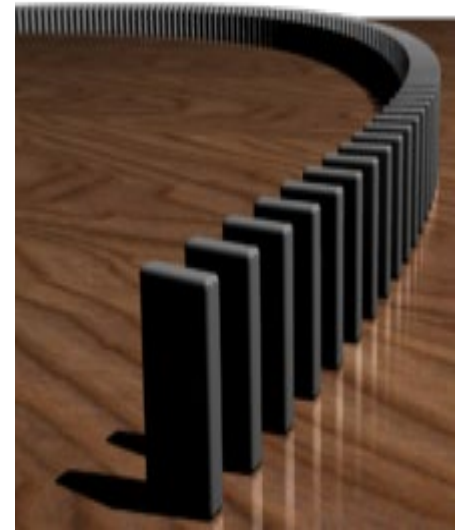
## Unit 1.4

### Mathematical Induction



# Mathematical Induction

- ❑ Mathematical induction can be used to prove statements that assert that  $P(n)$  is true for all positive integers  $n$ , where  $P(n)$  is a propositional function.
- ❑ A proof by induction contains two parts:
  - i. **Base case:** Show that  $P(1)$  is true.
  - ii. **Induction step:** Show that for all positive integers  $k$ , if  $P(k)$  is true, then  $P(k + 1)$  is also true.



Mathematical induction can be informally illustrated by reference to the sequential effect of falling dominoes (from Wikipedia)

# Examples

□ Prove that for all positive integers  $n$ ,

$$1 + 2 + \cdots + n = \frac{n(n+1)}{2}.$$

***Solution:***

- 1) (Base case) Since  $1 = \frac{1(1+1)}{2}$ , the statement is true for  $n = 1$ .
- 2) (Induction step) Assume the statement is true for  $n = k$  (where  $k$  is an arbitrary value), i.e.,

$$1 + 2 + \cdots + k = \frac{k(k+1)}{2}.$$

Consider the case where  $n = k + 1$ .

$$1 + 2 + \cdots + k + (k + 1) = \frac{k(k+1)}{2} + (k + 1) = \frac{(k+1)(k+2)}{2}.$$

Therefore, the statement is true for  $n = k + 1$ .

Hence, by induction, it is true for all positive integers.

□ Prove that for all positive integers  $n$ ,

$$1^2 + 2^2 + \cdots + n^2 = \frac{n(n+1)(2n+1)}{6}.$$

**Solution:**

- 1) (Base case) Since  $1^2 = \frac{1(1+1)(2+1)}{6}$ , the statement is true for  $n = 1$ .  
2) (Induction step) Assume the statement is true for  $n = k$ , i.e.,

$$1^2 + 2^2 + \cdots + k^2 = \frac{k(k+1)(2k+1)}{6}.$$

Consider the case where  $n = k + 1$ .

$$\begin{aligned} 1^2 + 2^2 + \cdots + k^2 + (k+1)^2 &= \frac{k(k+1)(2k+1)}{6} + (k+1)^2 \\ &= \frac{k(k+1)(2k+1) + 6(k+1)^2}{6} = \frac{(k+1)[k(2k+1) + 6(k+1)]}{6} \\ &= \frac{(k+1)[2k^2 + 7k + 6]}{6} = \frac{(k+1)(k+2)(2k+3)}{6} \end{aligned}$$

Therefore, the statement is true for  $n = k + 1$ .

Hence, by induction, it is true for all positive integers.

# Example: Summing a Geometric Series

□ Let  $r$  be a fixed real number. Prove that for all integers  $n$ ,

$$1 + r + r^2 + \cdots + r^n = \frac{1 - r^{n+1}}{1 - r}.$$

Solution:

- 1) (Base case) Since  $1 + r = \frac{1-r^2}{1-r}$ , the statement is true for  $n = 1$ .
- 2) (Induction step) Assume the statement is true for  $n = k$ , i.e.,

$$1 + r + r^2 + \cdots + r^k = \frac{1 - r^{k+1}}{1 - r}.$$

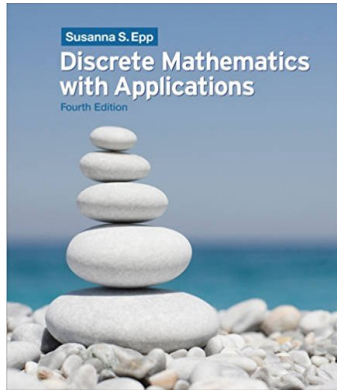
Consider the case where  $n = k + 1$ .

$$\begin{aligned} 1 + r + r^2 + \cdots + r^k + r^{k+1} &= \frac{1 - r^{k+1}}{1 - r} + r^{k+1} \\ &= \frac{1 - r^{k+1} + (r^{k+1} - r^{k+2})}{1 - r} = \frac{1 - r^{k+2}}{1 - r}. \end{aligned}$$

Therefore, the statement is true for  $n = k + 1$ .

Hence, by induction, it is true for all positive integers.

# Recommended Reading



- Sections 4.1-4.7, 5.2, Susanna S. Epp, *Discrete Mathematics with Applications*, 4<sup>th</sup> ed., Brooks Cole, 2010.

# Appendix (optional)

## Pythagoras Theorem

# An Art for Appreciation

## ❑ An interesting demo:

- <https://www.youtube.com/watch?v=CAkMUdeB06o>  
(<1 min.)

## ❑ How to prove it?

- [https://www.youtube.com/watch?v=BNCj-K2hd\\_k](https://www.youtube.com/watch?v=BNCj-K2hd_k)  
(~4 min.)