

## 一、pareto前沿

目标之间存在冲突

MOEAs多目标演化算法

- 目的是为了找到Pareto front
- 冲突举例: XM-1

$$\begin{aligned} \text{Minimize } f_1(\mathbf{x}) &= \frac{1}{2}x_1x_2\cdots x_{M-1}(1+g(\mathbf{x}_M)), \\ \text{Minimize } f_2(\mathbf{x}) &= \frac{1}{2}x_1x_2\cdots (1-x_{M-1})(1+g(\mathbf{x}_M)), \end{aligned}$$

x比y好: 那就是在每个目标上都更好 (或至少一样weakly dominates)

- $x$  weakly dominates  $y$ , denoted as  $x \succsim y$ , if
$$\forall i \in \{1, 2, \dots, m\}: f_i(x) \geq f_i(y)$$
- $x$  dominates  $y$ , denoted as  $x \succ y$ , if
$$\forall i \in \{1, 2, \dots, m\}: f_i(x) \geq f_i(y) \text{ and } \exists i \in \{1, 2, \dots, m\}: f_i(x) > f_i(y)$$
- $x$  is incomparable with  $y$ , if neither  $x \succsim y$  nor  $y \succsim x$

Pareto optima: 最优解

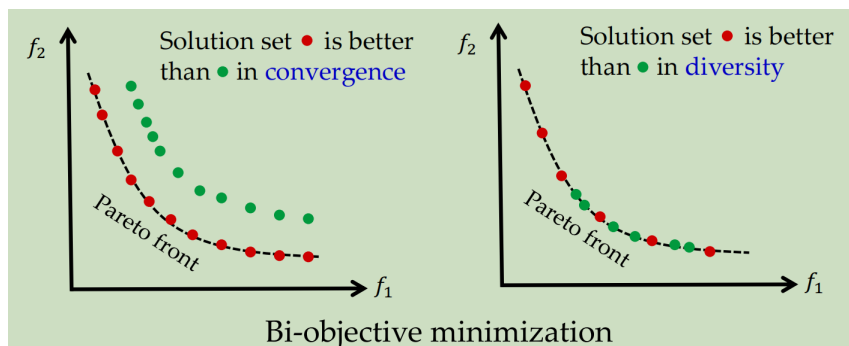
Pareto front: The collection of objective vectors of all Pareto optimal solutions 最优曲线

### Pareto front两种评价标准

- Convergence (to the Pareto front)
- Diversity (along the Pareto front)

- rank+distance

先比较rank, 后比较distance



### 解的比较方法

## Crowded comparison employed by NSGA-II

Given a set  $P$  of solutions, for any two solutions  $x, y$  in  $P$ ,  $x$  is better than  $y$ , if

- $\text{rank}(x) < \text{rank}(y)$
- or  $\text{rank}(x) = \text{rank}(y)$  but  $\text{distance}(x) > \text{distance}(y)$

Convergence

Diversity

## 二、NSGA-II

效果最差

### 1、父辈选择

Binary tournament selection

### Fast non-dominated sorting

针对rank

$n_x$ : 比当前好的解个数

$S_x$ : 比当前差的解的集合

对每个解，遍历一遍，计算上述两个参数，再根据rank，存储进 $F_i$

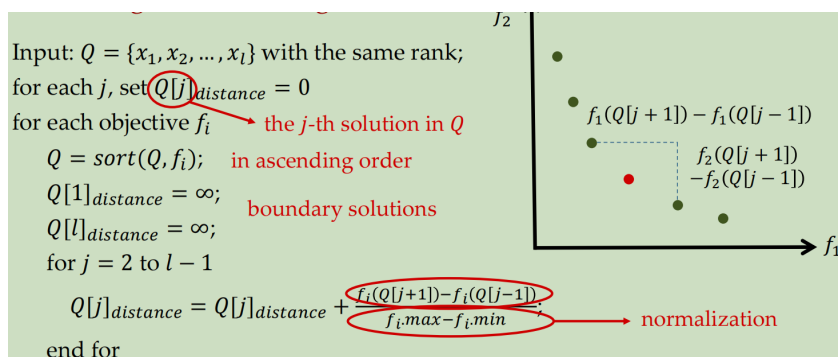
( $n_x=0$ 的所有解是 $\text{rank}=1$ )

复杂度 $O(n^2)$

### Crowding distance assignment

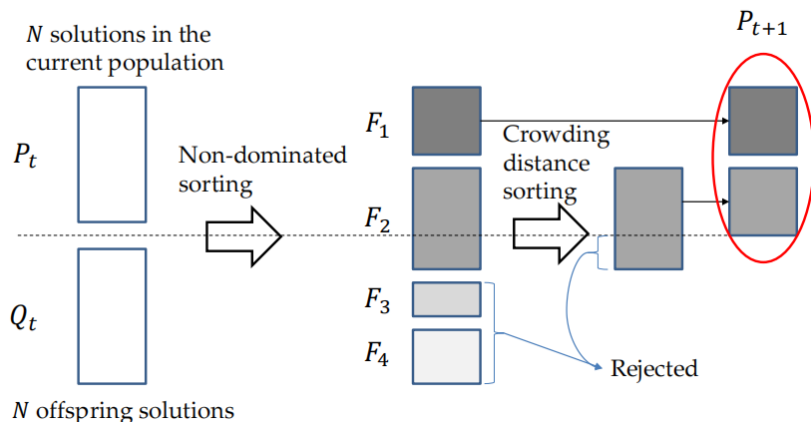
拥挤度，针对diversity

- 首尾设为无穷，因为肯定满足diversity
- $Q$ 越大越好
- 需要归一化作为评价尺度



## 2、幸存者选择

$N + N$  survivor selection

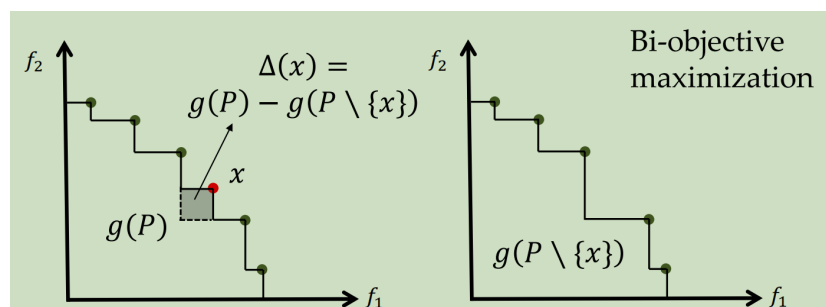


## 三、SMS-EMOA

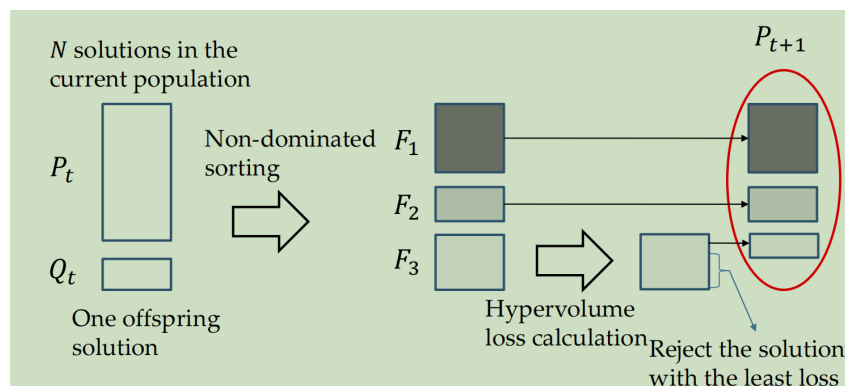
计算开销大，但收敛性好

- 增加对相同rank的解评估
- 利用quality indicators, 去掉一个解, 看损失多大
- 指标 $g(P)$

比如 the hypervolume indicator: 支配区域面积的大小 (二维好理解)



只产生一个子代offspring



## 四、MOEA/D

## 分解方法

- 加权法 (1范数)

算出每个单目标最优解，然后拟合成曲

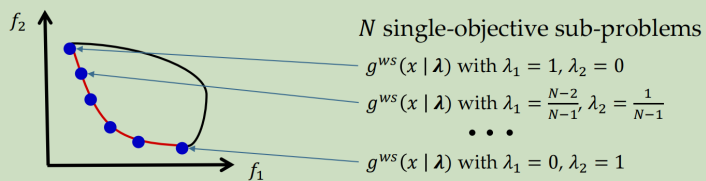
- Weighted sum approach

$$\min_{x \in \mathcal{X}} (f_1(x), f_2(x))$$



$$\min_{x \in \mathcal{X}} g^{ws}(x | \lambda) = \lambda_1 f_1(x) + \lambda_2 f_2(x)$$

$$\text{where } \lambda_1 + \lambda_2 = 1, \lambda_1, \lambda_2 \geq 0$$



- 切比雪夫方法 (无穷范数)

- Tchbycheff approach

$$\min_{x \in \mathcal{X}} (f_1(x), f_2(x))$$



$$\min_{x \in \mathcal{X}} g^t(x | \lambda, \mathbf{z}^*) = \max\{\lambda_1 |f_1(x) - z_1^*|, \lambda_2 |f_2(x) - z_2^*|\}$$

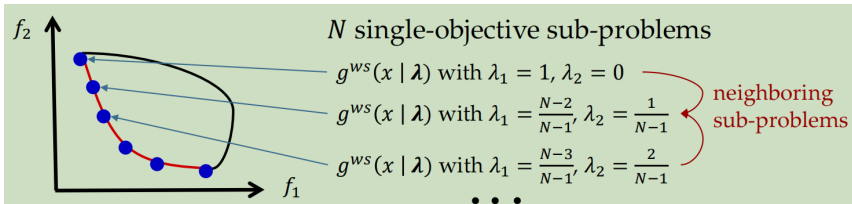
$$\text{where } \lambda_1 + \lambda_2 = 1, \lambda_1, \lambda_2 \geq 0$$

$\mathbf{z}^*$  is an Utopian point,

where  $z_1^* < \min\{f_1(x)\}$  and  $z_2^* < \min\{f_2(x)\}$

For any Pareto optimal solution  $x^*$ , there is a  $\lambda$  such that  $x^*$  is optimal to  $g^t(x | \lambda, \mathbf{z}^*)$

## 求解



For optimizing each sub-problem in each iteration

1. **Mating selection:** obtain the current solutions of some neighbours
2. **Reproduction:** generate a new solution by applying reproduction operators on its own solution and borrowed solutions
3. **Replacement:**
  - 3.1 replace its old solution by the new one if the new one is better
  - 3.2 pass the new solution on to some of its neighbours, and update its neighbor's solutions when better

串并行有区别，串行在遍历时，可能之前已经算过了