

# Feasibility Study of LTE Middle-Mile Networks in TV White Spaces for Rural India

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**Soubhik Deb**  
**Chaitanya Prasad Narisetty**  
**&**  
**Prof. Abhay Karandikar**

Department of Electrical Engineering  
Indian Institute of Technology Bombay

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## Introduction

- BharatNet - 638,619 villages via 250,000 Village Offices (VOs)
- Point of Presence (PoP) at VOs, wireless clusters at the villages
- Middle-mile network
- Envisioned to use the underutilized TV White Spaces
- Base Station at VOs, CPEs at associated villages

## Contribution

- Analyzed feasibility, estimates based upon Telecom Regulatory Authority of India targets rate of 512 kbps
- Proportional fair iterative algorithm for radio resource allocation

# Two-Stage System

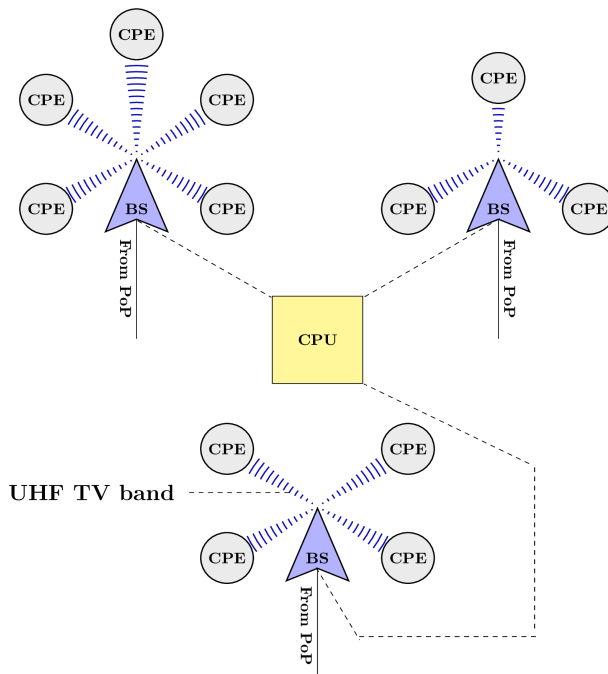


Figure: Simple overview of the two-stage system

# CoR Overview

- BaseStation - omnidirectional, 3 sectoral antennas (SAs) of  $120^\circ$
- Network partitioned into coordination regions (CoRs)
- Why truncated? **Negligible** interference from other tiers

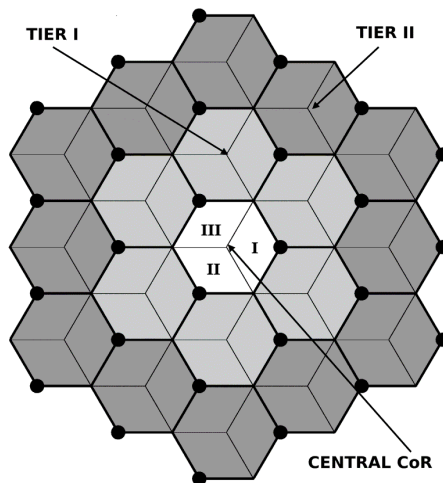


Figure: Truncated Network

# Problem Formulation

- $P_b^{\max}$  : Maximum transmission power of BaseStation  $b$
- $P_{bs}^k$  : Power of BaseStation  $b$  to a CPE in sector  $s$  on PRB  $k$
- $r_{(m,n)}^k$  : Received rate of  $CPE_{(m,n)}$  on PRB  $k$

The objective is:

$$\max \sum_{n: CPE_{(m,n)} \in \mathcal{N}_m} \sum_{k \in \mathcal{K}} r_{(m,n)}^k \quad \forall m : CoR_m \in \mathcal{M} \quad (1)$$

$$\text{subject to: } \sum_{s \in \mathcal{S}} \sum_{k \in \mathcal{K}} P_{bs}^k \leq P_b^{\max},$$

$$P_{bs}^k \geq 0.$$

- Involves optimization over continuous and discrete variables

# Interference Analysis

- No PRB is reused in a CoR  $\implies$  **zero** intra-CoR interference
- **non-zero** inter-CoR interference
- Allocating PRBs:
  - Coordinated Policy
  - Coordinated Policy with Bucketing

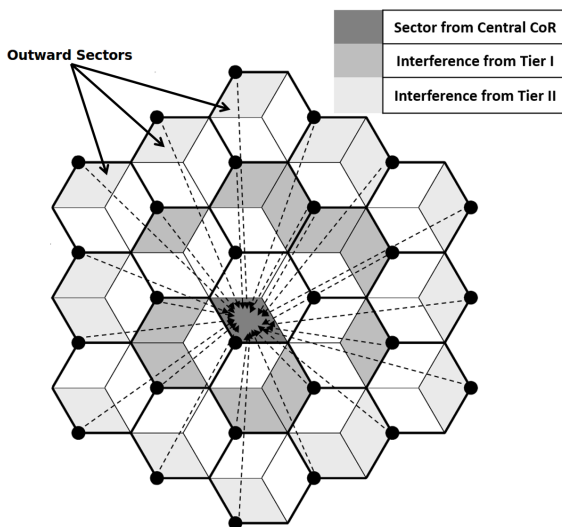


Figure: Interference Diagram for inter-CoR interference

# Topology: Modelling CPE Location

The no. of CPEs in the sector belonging to each BaseStation of a CoR is randomly chosen with  $q_s \in [0, 1]$  as follows:

$$\text{no. of CPEs in each sector} = \begin{cases} 2, & \text{w.p. } q_s \\ 1, & \text{w.p. } (1 - q_s) \end{cases}$$

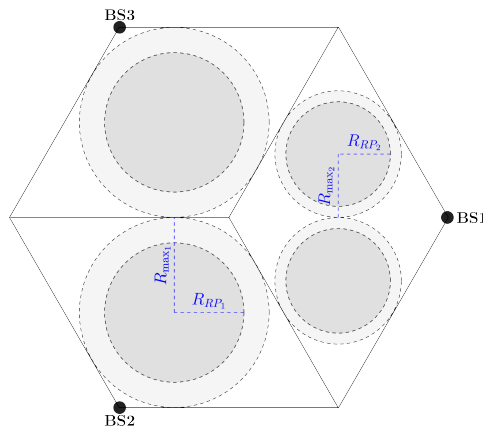


Figure:  $R_{RP1} = \gamma_1 \cdot R_{max1}$ ,  $R_{RP2} = \gamma_2 \cdot R_{max2}$ , where  $\gamma_1, \gamma_2 \in [0, 1]$

Using a standard Suboptimal Proportional Fair Allocation algorithm (SPFA), distribute the PRBs among the CPEs in each CoR



# Results

Multiple topologies with varying values of  $q_s$  and  $R_{RP}$  whose results are:

- ① Average received rate of the central CoR has little variance
- ② CoR have similar data rates, irrespective of no. of CPEs and location.

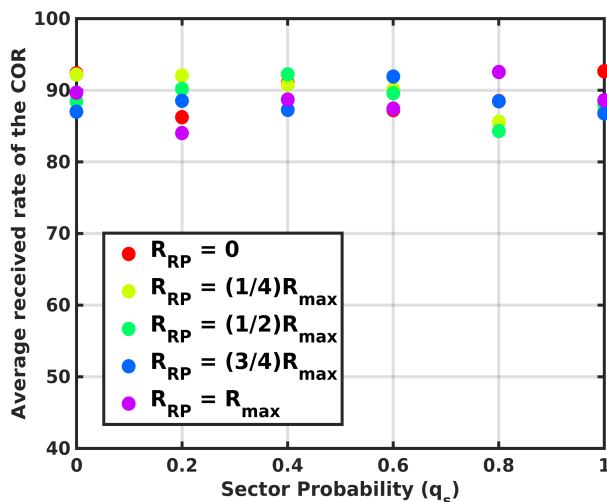


Figure: Received rate of the central CoR averaged over 10 network initializations

SPFA doesn't always allocate the PRBs fairly among all CPEs in a CoR.

# SPFA With Reservation (SPFAR)

- $\bar{r}_{(m,n)}$ : Latest avg. rate of the  $CPE_{(m,n)}$ ;  $\mathbf{r}_m = \{\bar{r}_{(m,n)} | CPE_{(m,n)} \in \mathcal{N}_m\}$

Equations:

$$\textcircled{1} \quad \bar{r}_m = \frac{\mathbf{r}_m \cdot \mathbf{I}}{|\mathcal{N}_m|}, \quad \chi_m : \{n \in \mathcal{N}_m | \bar{r}_{(m,n)} \leq \bar{r}_m\}$$

- $$\textcircled{2} \quad \text{For } CPE_{(m,n)} \in \chi_m, \text{ difference between } \bar{r}_{(m,n)} \text{ and } \bar{r}_m \text{ is quantified as}$$

$$K_{(m,n)} = \Gamma \cdot \frac{(\bar{r}_m - \bar{r}_{(m,n)})}{\bar{r}_m^{PRB}} \text{number of PRBs}, \quad (2)$$

where,  $K_{(m,n)}$  is rounded off to nearest integer,  $\Gamma$  is a tuning parameter and

$$\bar{r}_m^{PRB} = \frac{\sum_{n: CPE_{(m,n)} \in \mathcal{N}_m} \bar{r}_{(m,n)}}{K}.$$

- Select CPE with worst  $\bar{r}_{(m,n)}$  and assign it best  $K_{(m,n)}$  PRBs and remove them from set of PRBs that are to be assigned and so on
- Use SPFA to assign rest of PRBs
- SPFA till  $t_f$  iterations; based on average received rate over  $t_f$  iterations, SPFAR is employed in from iteration  $(t_f + 1)$  to  $(t_f + t_0)$

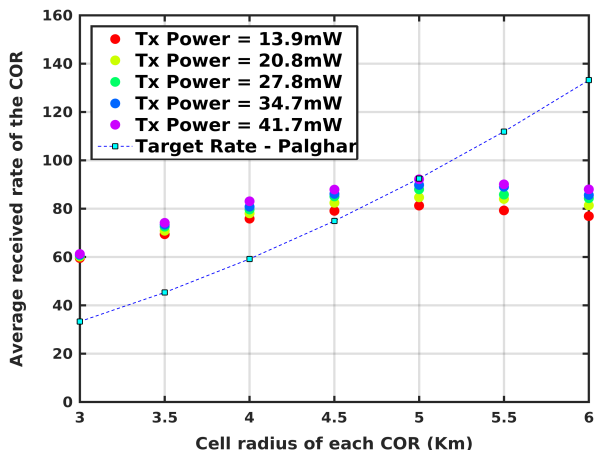
# Tradeoff

Algorithm	$\Gamma$	Fairness	CoR rate (Mbps)
SPFA	–	0.9689	89.73
SPFAR	1	0.9828	87.31
SPFAR	2	0.9862	86.26
SPFAR	3	0.9876	85.86
SPFAR	4	0.9905	85.41

**Table:** For a varying  $\Gamma$ , the tradeoff between Fairness among the CPEs vs. their total Throughput in the central CoR

- The reservation of PRBs improves fairness while introducing sub-optimality, thereby decreasing the total average rate of a CoR

# Transmission $T_x$ Power vs. Cell Radius $R_{cell}$



**Figure:** Received rate of the central CoR averaged over 10 network initializations

- Distance between the neighboring BSs decreases, the interference between them becomes more prominent
- After 5 km, path loss becomes more predominant decreasing the average rate
- Increasing the  $T_x$  power has little improvement in the total average rate due to higher inter-cell interference

# Meeting the Demands

- Palghar has a density of 285 people per sq. km
- Contention ratio is 1 : 50 as prescribed by TRAI
- For a cell radius of CoR below 5 km, the target rate of 1 Mbps per user is achievable

# Varying Demands Based on Topology

- If rates demanded CPEs  $\mathcal{N}_m$  in  $CoR_m$  are in the ratio  $\lambda_{(m,1)} : \lambda_{(m,2)} : \dots : \lambda_{(m,|\mathcal{N}_m|)}$ , then






## Results:

- 1 Average fairness among the CPEs with the new rate requirements is over 0.99 with an average standard deviation in CoR rate of just 3 Mbps
- 2 Average CoR rate achieved using the proposed RRA is independent of the rates demanded by its CPEs

# Future Work

- Adaptive power control can be utilized for optimality
- Framework has to be established for efficient RRA and interaction among operators in case of changes in demand and entry of a new operator
- For the coexistence among multiple operators in middle-mile network, a regulatory approach needs to be developed

# References I

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# Thank you

# Path Loss

- Hata Model appropriate for 470 – 590 MHz band

SINR of each individual PRB on a LOS link is then calculated using

$$SINR = PTx + GTx - PL - Noise \quad (3)$$

Spectral efficiency (SE) evaluated using

$$SE = \log_{10}(1 + SINR) \quad (4)$$

Using LTE CQI table, rate ( $R$ ) in that PRB of the LOS link in a SubFrame of duration  $T_{SF} = 1\text{ms}$  is obtained as follows:

$$R = \phi.(RE).(bpS).(CR).(\alpha)/T_{SF} \quad (5)$$

# Suboptimal Proportional Fair Allocation (SPFA)

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## Algorithm 1 SPFA

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1: Initialize:  $\mathcal{N}_m^* = \mathcal{N}_m$ 
2: for each iteration  $t$  do
3:   for each  $CoR_m \in \mathcal{M}$  do
4:     for each PRB  $k \in \mathcal{K}$  do
5:       if  $CoR_m \in \text{Tier II}$  then
6:         update  $N_m^*$  based on the bounds imposed by the Coordinated Policy
           with Bucketing
7:       end if
8:       for each  $CPE_{(m,n)} \in \mathcal{N}_m$  do
9:         Compute  $r_{(m,n)}^k(t)$ 
10:        Compute  $SC_{(m,n)}^k(t) = \frac{r_{(m,n)}^k(t)}{\bar{r}_{(m,n)}^k(t-1)}$ 
11:      end for
12:      Assign PRB  $k$  to
         $CPE_{(m,n^*)} = \mathcal{N}_m^* SC_{(m,n)}^k(t)$ 
13:       $\bar{r}_{(m,n^*)}^k(t) \leftarrow (1 - \frac{1}{T_c})\bar{r}_{(m,n^*)}^k(t-1) + \frac{1}{T_c}r_{(m,n^*)}^k(t)$ 
14:    end for
15:  end for
16: end for

```

## SPFA With Reservation (SPFAR)-II

- $\mathcal{K}_1$ : set of PRBs which are going to be allocated using Algorithm 2;  
 $|\mathcal{K}_1| = \sum_{n \in \mathcal{N}_m} K_{(m,n)}$  PRBs; selected uniformly at random
- Rest are assigned to  $\mathcal{K}_2$
- Allocate the PRBs in  $\mathcal{K}_2$  among all CPEs  $|\mathcal{N}_m|$  of the  $CoR_m$  using SPFA

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### Algorithm 2 Reserved PRB Allocation

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- 1: **for** each  $CoR_m \in \mathcal{M}$  **do**
  - 2:   **while**  $|\chi_m| \neq 0 \wedge |\mathcal{K}_1| \neq 0$  **do**
  - 3:     **Choose**  $CPE_{(m,n)} \in \chi_m$  such that
 
$$\bar{r}_{(m,n)} \leq \bar{r}_{(m,n')}, \forall CPE_{m,n'} \in \chi_m$$
  - 4:     **Allocate** the set of PRBs  $\mathcal{K}_3 \subset \mathcal{K}_1$  to  $CPE_{(m,n)}$  as
 
$$\mathcal{K}_3 = \left\{ y \in \mathcal{K}_1 \mid r_{(m,n)}^y \geq r_{(m,n)}^k, \forall k \in \mathcal{K}_1 - \mathcal{K}_3 \right. \\ \left. \wedge |\mathcal{K}_3| = K_{(m,n)} \right\}$$
  - 5:      $\chi_m \leftarrow \chi_m - \{CPE_{(m,n)}\}$
  - 6:      $\mathcal{K}_1 \leftarrow \mathcal{K}_1 - \mathcal{K}_3$
  - 7:   **end while**
  - 8: **end for**
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