

The title is framed by a large dashed rectangle. A dashed arrow points from the top-right corner of this rectangle towards the top-right corner of the slide. Another dashed arrow points from the bottom-left corner of the rectangle towards the bottom-left corner of the slide. A solid vertical arrow points upwards from the bottom-right corner of the rectangle.

Slack-stealing scheduling



1

Introduction



Contextualization and
Motivation



Contextualization

- Scheduling tasks in real time systems
- (Example) Robotics:
 - Sensor for monitoring the environment
 - Send data to server
- Run all those tasks without collision within a certain time.

Motivation

- How do we schedule hard-periodic and soft-aperiodic tasks in Fixed-priority Preemptive System?
(Concepts explained at next section)
- Is it possible to guarantee that all periodic and aperiodic tasks will be completed?



2

Review



Important concepts

LET'S REVIEW SOME CONCEPTS



Periodic Tasks

Tasks represented as a tuple $T = \langle \Phi, T, C, D \rangle$

$\Phi \rightarrow$ first time task appears

$T \rightarrow$ Period of the task

$C \rightarrow$ execution time

$D \rightarrow$ relative deadline

Aperiodic Tasks

Tasks without deadline or with soft deadline that can start in different times.



Soft Tasks

Tasks with extendable deadlines

Hard Tasks

Tasks that must be completed until the deadline

LET'S REVIEW SOME CONCEPTS



Fixed-priority system

Tasks have a fixed order of execution that must be respected

Dynamic-priority system

The order that the tasks must be executed is calculated during the system execution



Preemptive Systems

The execution of one task can be stopped by the system due to the need of processing other tasks.



3

Main Content



How to address the
scheduling problem



Older Standard Approaches

- Always Execute the Periodic Tasks first
- Await for idle time to execute aperiodics
- Or create a periodic task to handle it
- Cons:
 - Larger response time for aperiodics

Example - Using Idle Available Time

Periodic tasks:

$$T_1 = \langle \Phi_1=0, T_1=4, C_1=1, D_1=1 \rangle$$

$$T_2 = \langle \Phi_2=0, T_2=6, C_2=3, D_2=6 \rangle$$

Aperiodic tasks(arrival at 0):

$$p_1 = \langle C_1=1 \rangle \rightarrow R_1 = 6$$

$$p_2 = \langle C_2=2 \rangle \rightarrow R_2 = 12$$

task	τ_1	τ_2	τ_2	τ_2	τ_1	p_1	τ_2	τ_2	τ_1	τ_2	p_2	p_2
time	1	2	3	4	5	6	7	8	9	10	11	12

The Slack Stealing Approach

- Idea: “Steal” processing time from periodics without causing time faults.
- How: pre-compute slack time availables.
- Goal: reduce the response time of aperiodic tasks (a.k.a decrease the finish time)

Example - Using Slack-Stealing

Periodic tasks:

$$T_1 = \langle \Phi_1=0, T_1=4, C_1=1, D_1=1 \rangle$$

$$T_2 = \langle \Phi_2=0, T_2=6, C_2=3, D_2=6 \rangle$$

Aperiodic tasks(arrival at 0):

$$p_1 = \langle C_1=1 \rangle \rightarrow R_1 = 2 < 6$$

$$p_2 = \langle C_2=2 \rangle \rightarrow R_2 = 8 < 12$$

task	τ_1	p_1	τ_2	τ_2	τ_1	τ_2	p_2	p_2	τ_1	τ_2	τ_2	τ_2
time	1	2	3	4	5	6	7	8	9	10	11	12

The Slack Stealing - Basic Scenario

- All periodic tasks variables are already know.
- Tasks do not stop themselves nor sync with each other.
- Any task can be instantly preempted.
- Scheduling overhead are assumed to be zero.
- Aperiodic tasks has soft deadlines
(but can be hard with some adjustments)

The Slack Stealing - Algorithm

1. Pré-runtime configurations
 - a. Define a fixed priority algorithm for the periodic tasks
(ex: Earliest Deadline (EDS), Rate Monotonic (RTS))
 - b. Compute the Hyperperiod H (LCM of periods)
 - c. Define a priority for the aperiodic tasks
(ex: Earliest Arrival Time)

The Slack Stealing - Algorithm

1. Pré-runtime configurations

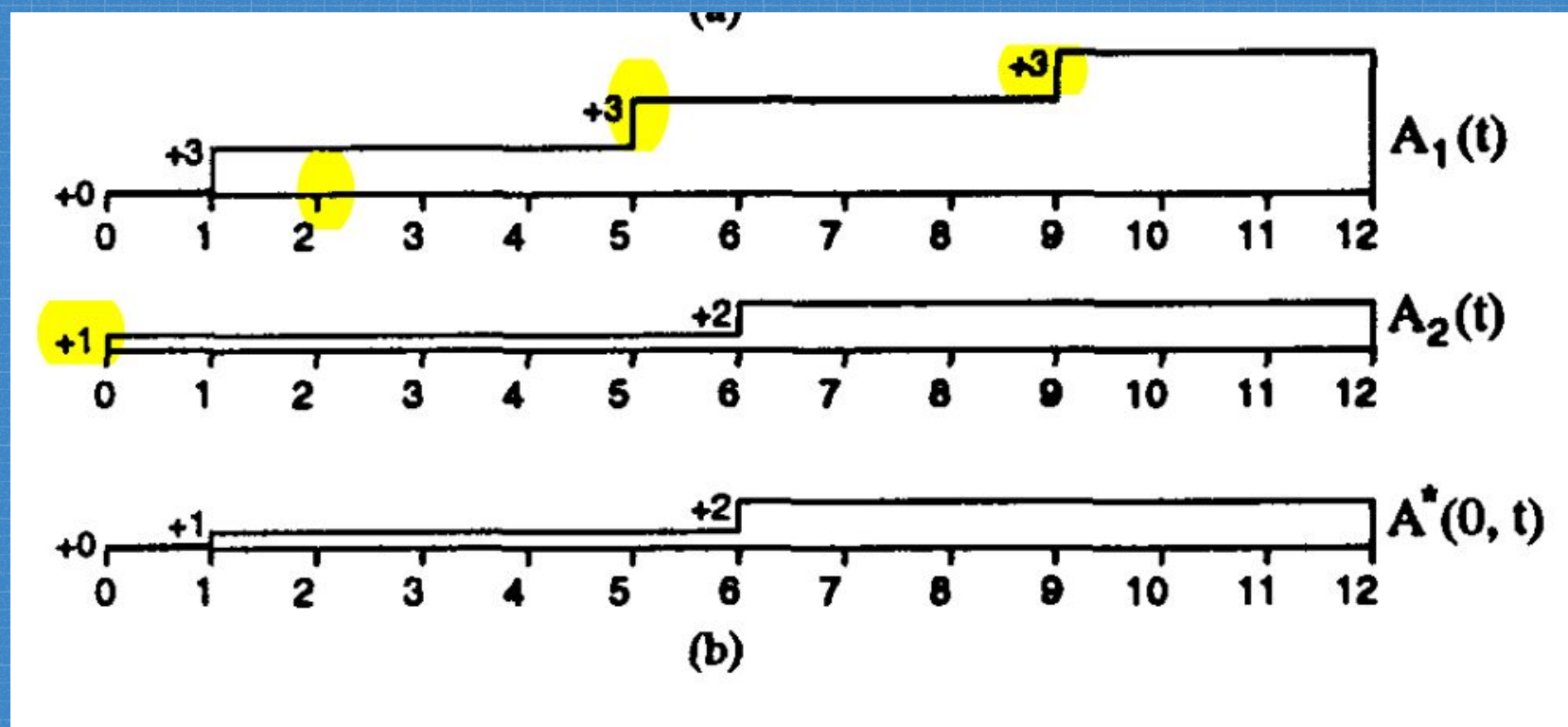
a. Compute Periodic Ready Work for all tasks

$P_i(t) \rightarrow$ qtd of processing done or waiting to be done by periodic tasks from $[0, t]$ at level i or higher.

b. Compute the Maximum Aperiodic Processing Available.

$A_i(t) \rightarrow$ max aperiodic work in $[0, t]$ at level i or higher with all τ_i deadline meet

The Slack Stealing - Algorithm



The Slack Stealing - Algorithm

1. Pré-runtime configurations

a. Initialize Accumulators

$A(s) \rightarrow$ max aperiodic work done in $[0, s]$ with all τ_i
deadlines meet (this was a horrible variable name choice)

$I_i(s) \rightarrow$ Accumulate Inactive time at level i or higher

The Slack Stealing - Algorithm

1. Scheduler Algorithm

a. The scheduler has 2 queues:

- i. The Periodic Tasks Queue
- ii. The Aperiodic Tasks Queue

b. And can be in one of these three states

- i. Idle \rightarrow Increase all $I_i(s)$ accumulators
- ii. Running $\tau_i \rightarrow$ Increase $I_j(s)$ for $j > i$
- iii. Running $p_i \rightarrow$ Increase $A(s)$

The Slack Stealing - Algorithm

1. Scheduler Algorithm - Scenarios

- a. If τ_i is running and τ_j ($j < i$) arrives, τ_j start running
- b. If a periodic ends or a aperiodic arrives we
 - i. Check if the aperiodic queue is not empty
 - ii. Check if we have slack time available to use by computing the following

$A_i(s, t) \rightarrow$ max aperiodic work available in $[s, t]$ at level i or higher with all τ_i deadline meet

$A^*(s, t) \rightarrow$ max aperiodic work in $[s, t]$ with all τ_i deadline meet

Slack Stealing - Example

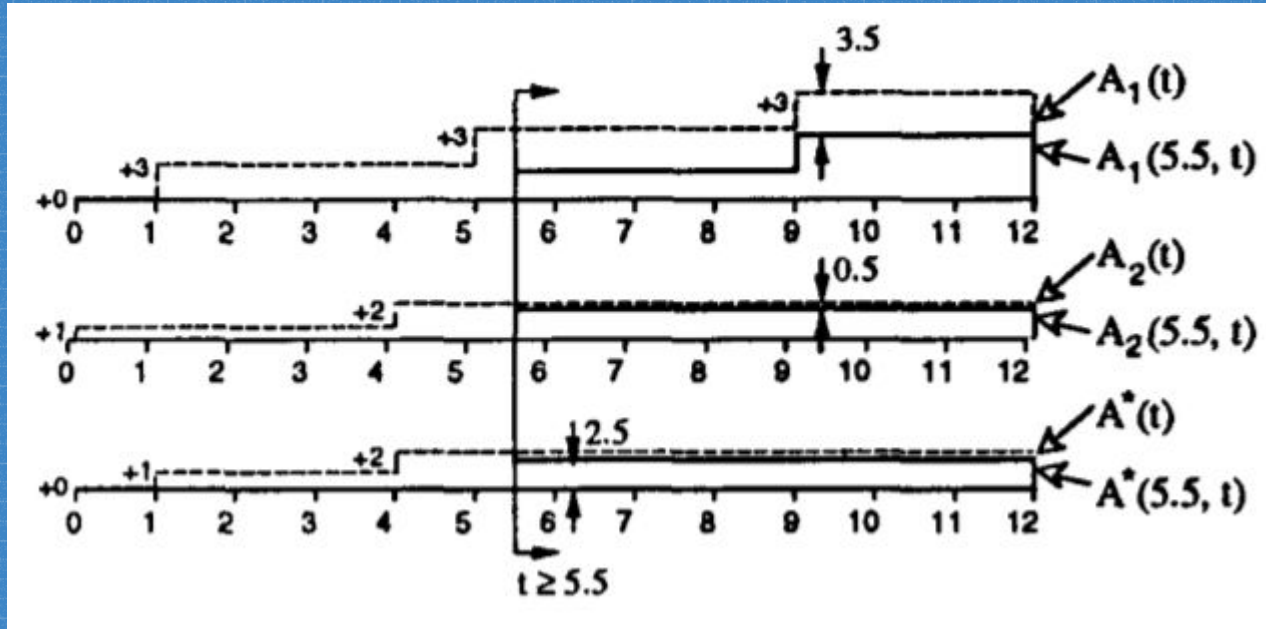
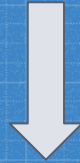


Figure: Example from slack stealer paper

Optimality of the Slack Stealer

Set of periodic tasks scheduled by fixed-priority and
aperiodic stream



Slack Stealing min response time of aperiodic and
meet deadline of periodic

(mathematical proof on original paper)



4

Illustration



Demo implementation of
slack stealing





5

References

An Optimal Algorithm for Scheduling Soft-Aperiodic Tasks in Fixed-Priority Preemptive Systems

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Abstract

This paper presents a new algorithm for servicing soft deadline aperiodic tasks in a real-time system in

solution of many practical problems which arise in actual real-time systems including task synchronization, transient overload, and simultaneous scheduling of both periodic and aperiodic tasks, among others.

- Jeon, W.; Kim, W.; Lee, H.; Lee, C.-H. Online Slack-Stealing Scheduling with Modified laEDF in Real-Time Systems. Electronics 2019, 8, 1286. doi.org/10.3390/electronics8111286
- J. P. Lehoczky and S. Ramos-Thuel, "An optimal algorithm for scheduling soft-aperiodic tasks in fixed-priority preemptive systems," [1992] Proceedings Real-Time Systems Symposium, 1992, pp. 110-123, doi: 10.1109/REAL.1992.242671.
- Real-Time Systems classes from the course
- <https://www.geeksforgeeks.org/tasks-in-real-time-systems/>

Thanks!

ANY QUESTIONS?

Tiago Marino

Gabriel Van Loon

Rodrigo Arboleda

João Vitor Ramos



Code at Github:
[/firewall1011/Slack-Stealing-Scheduling](https://github.com/firewall1011/Slack-Stealing-Scheduling)