

Checkpoint/Restart in Linux

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Agenda

- What and Why Checkpoint/Restart
- Prerequisites and Requirements
- Usage Overview
- Kernel API
- Current Status (v18 posted)
- Demo
- Design/API Discussion



What is Checkpoint/Restart?

- Checkpoint: save state of a running application
- Restart: resume application from saved state
- Migration: checkpoint on one host, restart on another
 - Static migration
 - Live migration



Why Checkpoint/Restart?

- Reduced application downtime: checkpoint, reboot, restart
- Application mobility
 - User-session mobility
 - Migrate application to another server before system upgrade
- Improve system utilization
- Faster error recovery with periodic checkpoints



Why Checkpoint/Restart?

- Slow-start applications
- Debug start from last checkpoint
- General time-travel
- Other use-cases?



Pre-requisite: Freezer

- Freeze application process-tree for consistent checkpoint
- Freezer implemented as a cgroup
- Status: merged into mainline
- Usage:
 - \$ mount -t cgroup -o freezer foo /cgroups
 - \$ echo FROZEN > /cgroups/\$pid/freezer.state
 - \$ cat /cgroups/\$pid/freezer.state

FROZEN

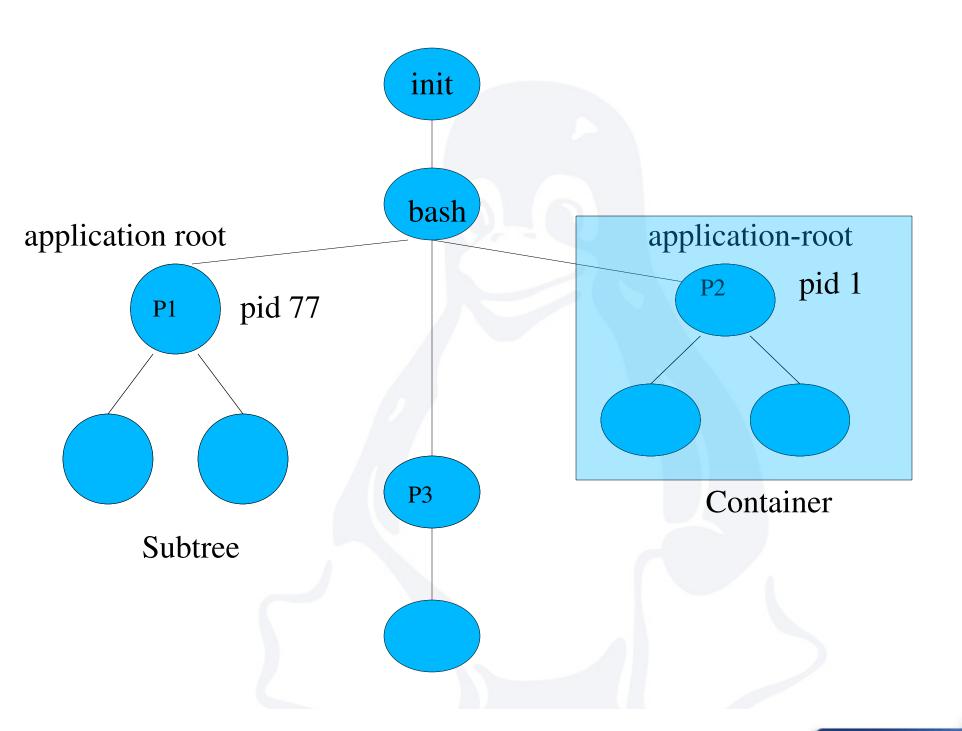
\$ echo THAWED > /cgroups/\$pid/freezer.state



Pre-requisite: Containers

- Containers isolated name spaces enable reuse of resource ids
- Needed in C/R to restore original resource ids in the application
- Create containers using clone(2) system call (or a wrapper to it)
- Status: Mount, UTS, IPC, PID, Network name spaces, devpts merged







Basic Requirements

- Transparent work with existing binaries
- Low impact on other subsystems
 - Performance
 - No duplicate code-paths
 - Maintenance overhead
- Integrated into kernel
- Generalized: Not restricted to specific applications



Basic Requirements

- Allow full-container and subtree C/R
 - Full-container: C/R of complete process tree
 - Subtree: C/R of part of process tree
- Full-container C/R needed to:
 - Restore original resource-ids
 - Prevent 'leaks' in shared resources
- Subtree C/R useful within limits for:
 - Resource-id agnostic applications
 - C/R aware applications
 - Development
- Enable self-checkpoint



Checkpoint Usage

- Create application in a container
- Freeze container (from parent)
- Checkpoint:
 - \$ checkpoint -p 1234 > checkpoint-img.1
- Snapshot file system leverage fs capability (btrfs, nilfs etc)
- Thaw application
- Terminate application (if necessary)



Restart Usage

- Restore file system state to snapshot (leverage fs capabilities)
- Restart application in new container
 \$ restart -container --wait < checkpoint-img.1



C/R Kernel API (proposed)

- sys_checkpoint(pid_t pid, int fd, ulong flags);
- sys_restart(pid_t pid, int fd, ulong flags);
 - pid: root of application process-tree to checkpoint/restart
 - fd: file descriptor or socket to/from which to write/read checkpoint image



C/R Kernel API (proposed)

```
struct clone_arg {
   u64 clone flags;
   u32 parent tid,
   u32 child tid;
   u32 nr pids;
   /* plus some reserved space */
  sys clone2() (struct clone struct *cs, pid t *pids)

    Allow additional clone-flags

    Allow ability to choose pids for child process
```



C/R Kernel API: Image format

- Checkpoint image format:
 - Blob that may change over time
 - Has a version number
 - Stream-able
- User space tools convert image between kernel versions
- General layout:
 - Image header
 - Task hierarchy
 - Task state of each task
 - Image trailer
- Shared objects saved only once



Status: Done

- Currently restored (in ckpt-v18)
 - Process-trees, pthreads, signals, handlers
 - SYS-V IPCs, FIFOs, itimers
 - Devices: null, random, zero, pts
 - Self-checkpoint
- File systems:
 - Regular files and directories in normal fs
 - Some special fs (devpts)
- Architectures: Working on i386, ppc64, s390



C/R Network state

- AF_UNIX Sockets restored
- AF_INET Sockets:
 - Restored if both ends were checkpointed
 - If one end was checkpointed, connection restored if restarted within tcp delay



C/R: Devices

- PTY devices restored. Other virtual devices can be restored
 - Use Client/Server model and C/R server
 - Display: Use VNC
 - Audio: Pulse Audio
- If application tied to specific hardware, C/R in kernel is complex
 - Device like /dev/rtc maybe in use
 - Device may not be available
 - Restore such devices in user-space ?



Status: Pending

- Time, POSIX Timers, Timezones
- File systems
 - Pseudo FS (eg: /proc)
 - NFS ?
 - Unlinked files, directories
- Devices
- Event-poll (WIP)
- Others:
 - Inotify, mount-points, mount-ns, etc



Discussion

 We have some design/API choices that we would like some feedback on



Q&A: Time

- C/R of time presents interesting challenges
 - Restart may happen after long time
- Choose a policy on restart ?
 - Use current or original time?
 - Timer-expirations relative/absolute?
- Is policy per-process or per-restart?
 - Which policy for new children?



Q&A: Process tree

- Restore process-tree in user-space?
 - Leverage existing calls like fork(), clone()
 - Allows subtree C/R
 - Needs clone2()
 - Needs kernel synchronization of processes in tree during restart
- Or, in kernel ?
 - Avoid clone2() and synchronization in-kernel
 - Reduced flexibility ?



Q&A: Restore fs, network

- \$ ns_exec -container /bin/application -arg
- \$ echo FROZEN > /cgroups/\$pid/freezer.state
- \$ checkpoint -p 1234 > checkpoint-img.1
- \$ snapshot-filesystem
- # Thaw and terminate app
- \$ restart -container < checkpoint-img.1</pre>
- Q: Restart program creates container but how can we have it restore file system and network configuration state?



Q&A: Kernel API

- cradvise()
 - Eg: skip C/R of some portion of memory or some device and let user-space handle it
- Notify C/R-aware applications of:
 - Pending or completed checkpoint
 - Completed restart
- Restart parts of an application
 - Eg: restore regular file fds, ignore IPC
- Others?



Q&A: Kernel API

- Ability to control/optimize C/R. Eg:
 - Skip restore of part of memory
 - Skip restart of some device (let user space deal with it)
 - Use parent's UTS namespace
 - Skip C/R of IPC
- Use single system call: cradvise() with flags?
- Or separate calls: cradvise_fd(), cradvise mem() etc



Project Info

Maintainer: Oren Laadan <oren@librato.com>

Mailing list: Containers@lists.linux-foundation.org

Wiki: http://ckpt.wiki.kernel.org

Code:

- git://git.ncl.cs.columbia.edu/pub/git/linux-cr.git
- git://git.ncl.cs.columbia.edu/pub/git/user-cr.git
- git://git.sr71.net/~hallyn/cr_tests.git



Backup slides



Simple ns_exec.c

```
int child_func(void *arg)
        char **argv = arg;
        execv(argv[0], argv);
main(int argc, char *argv[])
        int pid, status, rc;
        unsigned long flags;
        void *child_stack, *stack;
        flags = CLONE_NEWNS | SIGCHLD;
        flags |= CLONE_NEWIPC | CLONE_NEWUTS | CLONE_NEWPID;
        stack = malloc(getpagesize());
        child_stack = stack + getpagesize();
        pid = clone(child_func, child_stack, flags, (void *)&argv[1]);
        rc = waitpid(pid, &status, __WALL);
```

C/R Example

```
/bin/sh
File Edit View Terminal Help
 $ mount -t cgroup -o freezer,ns /cgroups
 $ ns_exec -cpuim /bin/bash
 $ ./a.out &
 [1234]
 i is 1
 i is 2
 i is 3
 $ echo FROZEN > /cgroups/1234/freezer.state
 $ /usr/local/bin/checkpoint -p 1234 > checkpoint-1.img
 $ kill -9 1234
 $ echo THAWED > /cgroups/1234/freezer.state
 $ init 6
 $ /usr/local/bin/restart --container --wait < checkpoint-1.img</pre>
 i is 4
 i is 5
```

Resources to Checkpoint

- Process trees,
- UTS (host name)
- Memory (shared-mem, mmaps etc)
- Open-file state
- SysV IPC, FIFOs
- AF_UNIX and AF_INET Sockets
- Restart-timers
- Signal state
- Devices
- Time
- Others?
- TBD: Drop and use following two slides?



Existing implementations

- OpenVZ (Parallels)
- Zap (Columbia University)
- MCR (IBM)
- BLCR
- Others (www.checkpointing.org)

