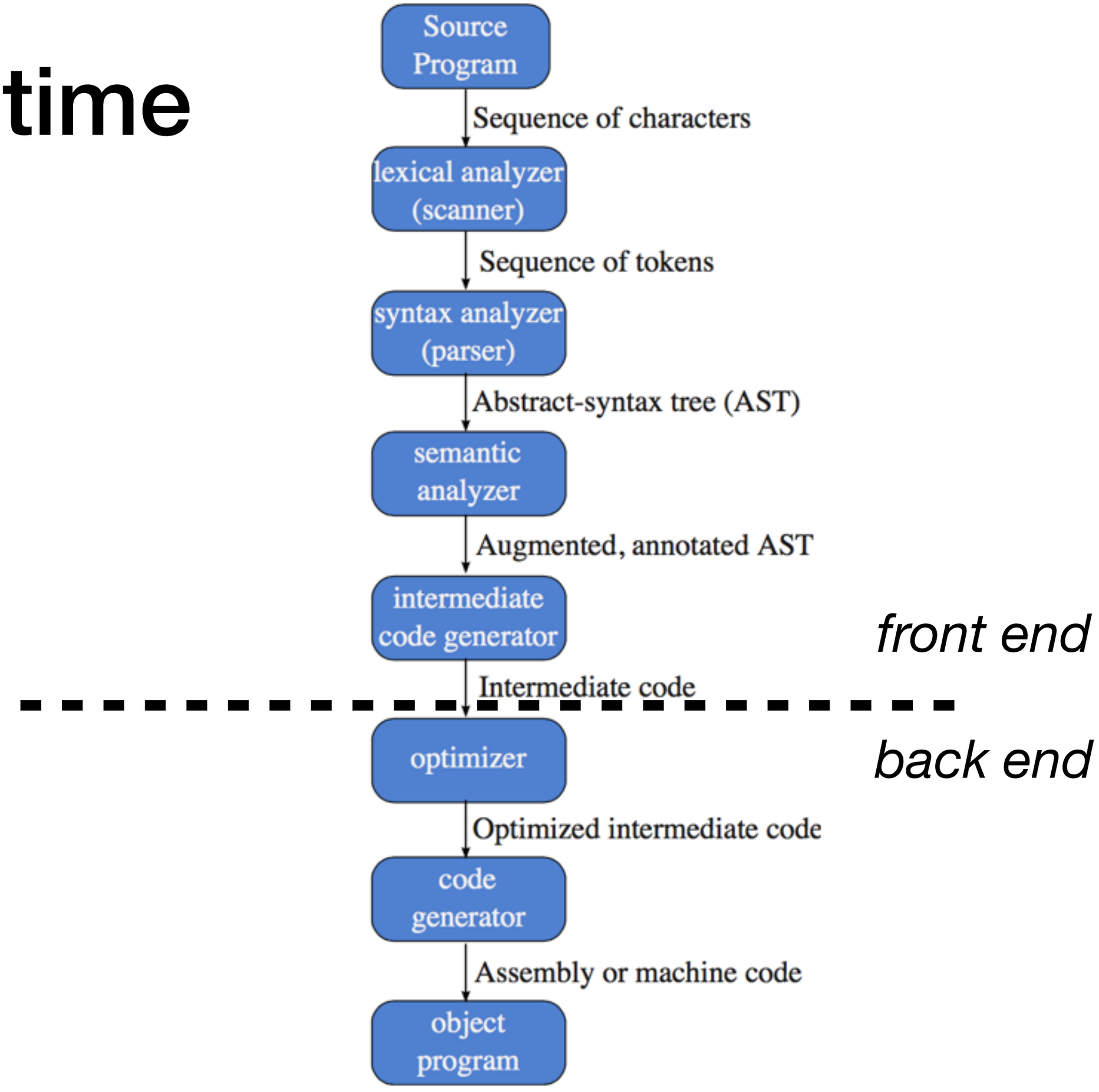
Finite-state machines

CS 536

Last time


Source Program
Sequence of characters
lexical analyzer (scanner)
Sequence of tokens
syntax analyzer (parser)
Abstract-syntax tree (AST)
semantic analyzer
Augmented, annotated AST
intermediate code generator
front end
Intermediate code
optimizer
back end
Optimized intermediate code
code generator
Assembly or machine code
object program

The scanner

Translates sequence of chars into sequence of tokens

Each time scanner is called it should:

find longest sequence of chars corresponding to a token

return that token

Scanner generator

Generates a scanner!!!

Needs one regular expression for each token

Needs regular expressions for things to ignore

- comments, whitespace, etc.

To understand how it works, we need FSMs

- finite state machines

FSMs: Finite State Machines

Aka finite automata

Input: string (seq of chars)

Output: accept / reject

i.e., input is legal in language

FSMs

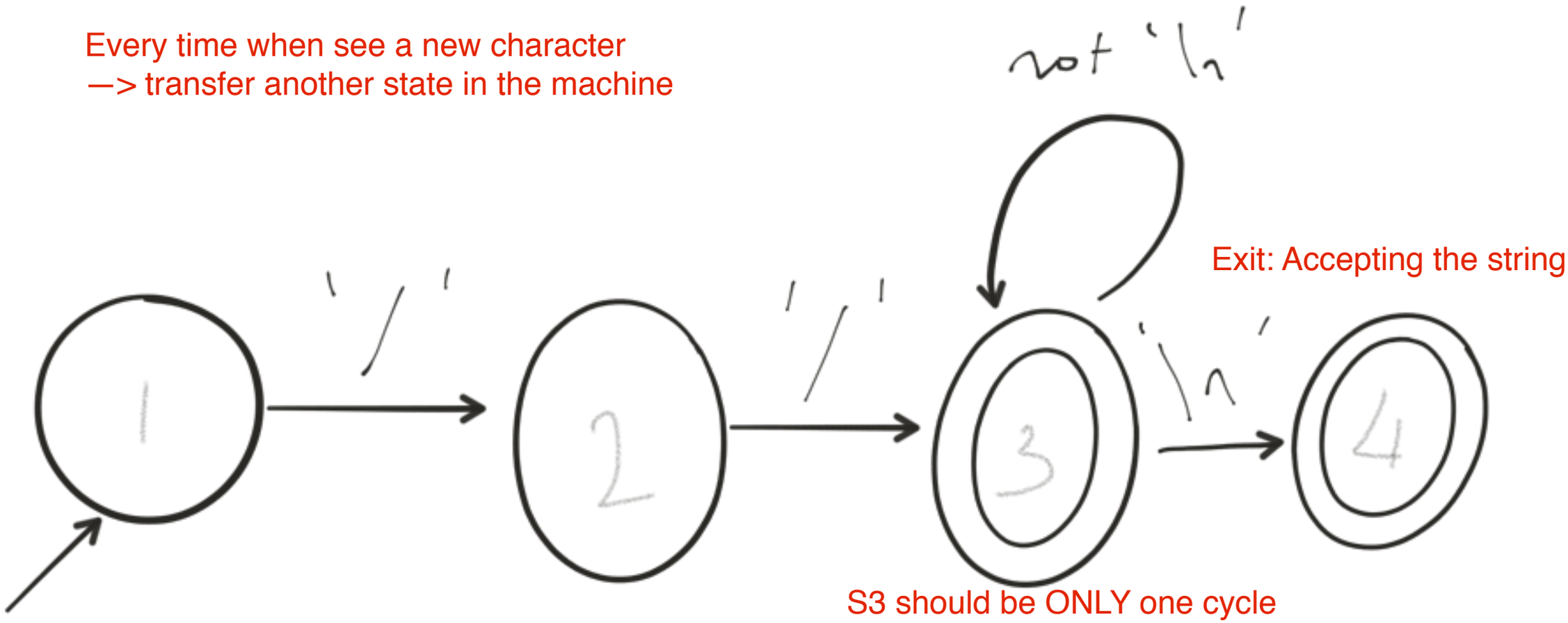
Represent regular languages

Good enough for tokens in PLs

Example 1

single line comments with //

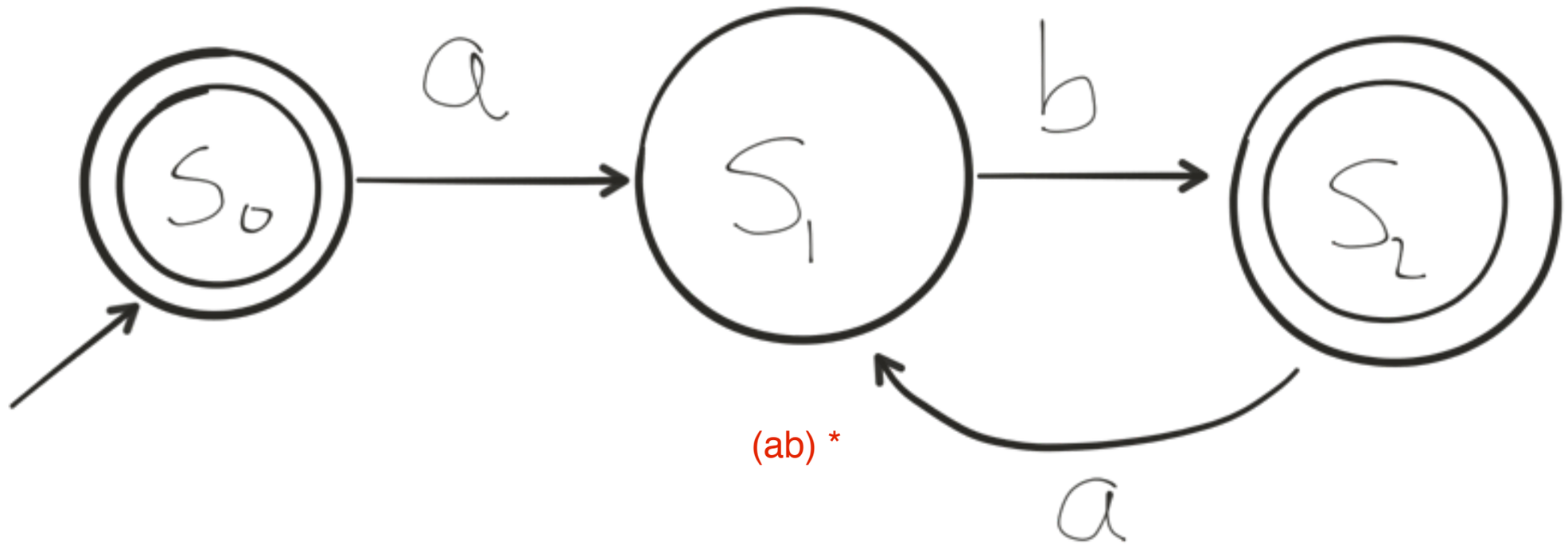
Every time when see a new character
—> transfer another state in the machine



Exit: Accepting the string

S3 should be ONLY one cycle

Example 2



What language does this accept?

—> S0 (double cycle)—>S1
I<— — — — — — — — — — I

Can you find an equivalent, but smaller, FSM?

How an FSM works

```
curr_state = start_state

let in_ch = current input char

repeat

  if there is edge out of curr_state with
  label in_ch into next_state

    cur_state = next_state

    in_ch = next char of input

  o/w stuck // error condition
```

Otherwise

```
until stuck or input string is consumed

string is accepted iff entire string is
consumed and cur_state = final_state
```

FSMs, formally

$$(Q, \Sigma, \delta, q, F)$$

finite set of states

the alphabet (characters)

transition relation

$$\delta : Q \times \Sigma \to Q$$

start state

$$q \in Q$$

final states

$$F \subseteq Q$$

FSMs, formally

M: L(M) = {w | M(w) = true}
M = M' <=> L(M) = L(M')

$$(Q, \Sigma, \delta, q, F)$$

FSM accepts string

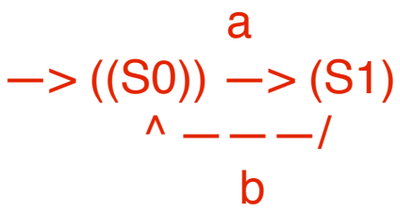
$$x_1 x_2 x_3 \ldots x_n$$

$$\Longleftrightarrow$$

$$\delta(\ldots\delta(\delta(\delta(q, x_1), x_2), x_3)\ldots, x_n) \in F$$

The language of FSM *M* is the set of all words it accepts, denoted $L(M)$

FSM example, formally



$$(Q, \Sigma, \delta, q, F)$$

$$Q = \{s_0, s_1\}$$

$$\Sigma = \{a, b, c\}$$

$$q = s_0$$

$$F = \{s_0\}$$

$$\delta = s_0, a \rightarrow s_1$$
$$s_1, b \rightarrow s_0$$

	a	b	c
s0	s1		
s1		s0	

anything else, machine is stuck

Coding an FSM

```
curr_state = start_state

done = false

while (!done)

 ch = nextChar()

 next = transition[curr state][ch]

 if (next == error || ch == EOF)

   done = true

 else

   curr_state = next

return curr_state == final_state
```

FSM types: DFA & NFA

Deterministic

no state has > 1 outgoing edge with same label

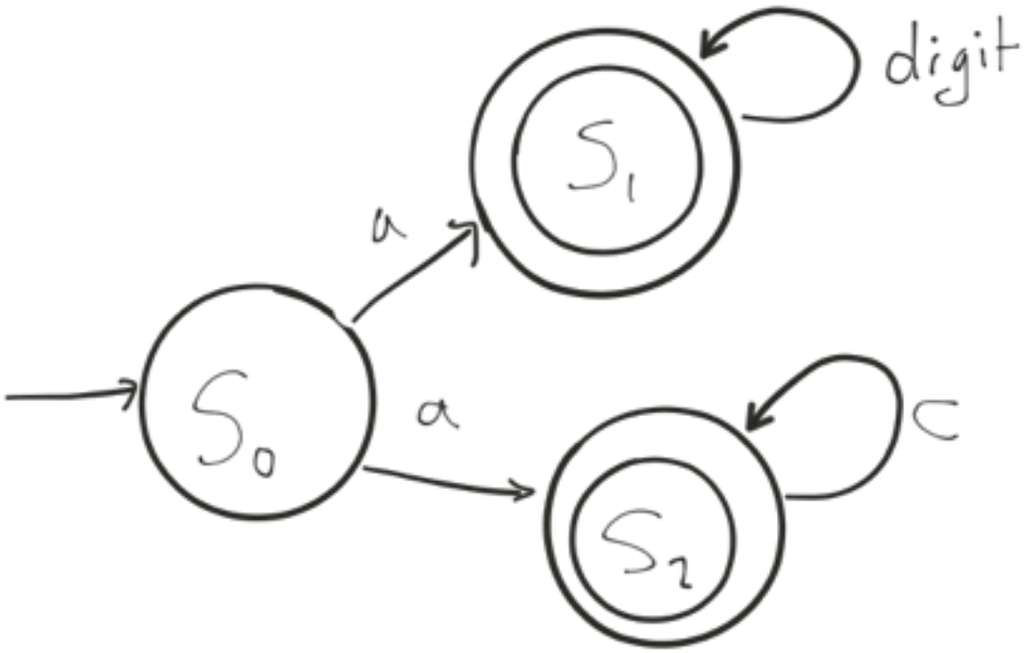
Nondeterministic

states may have multiple outgoing edges with same label

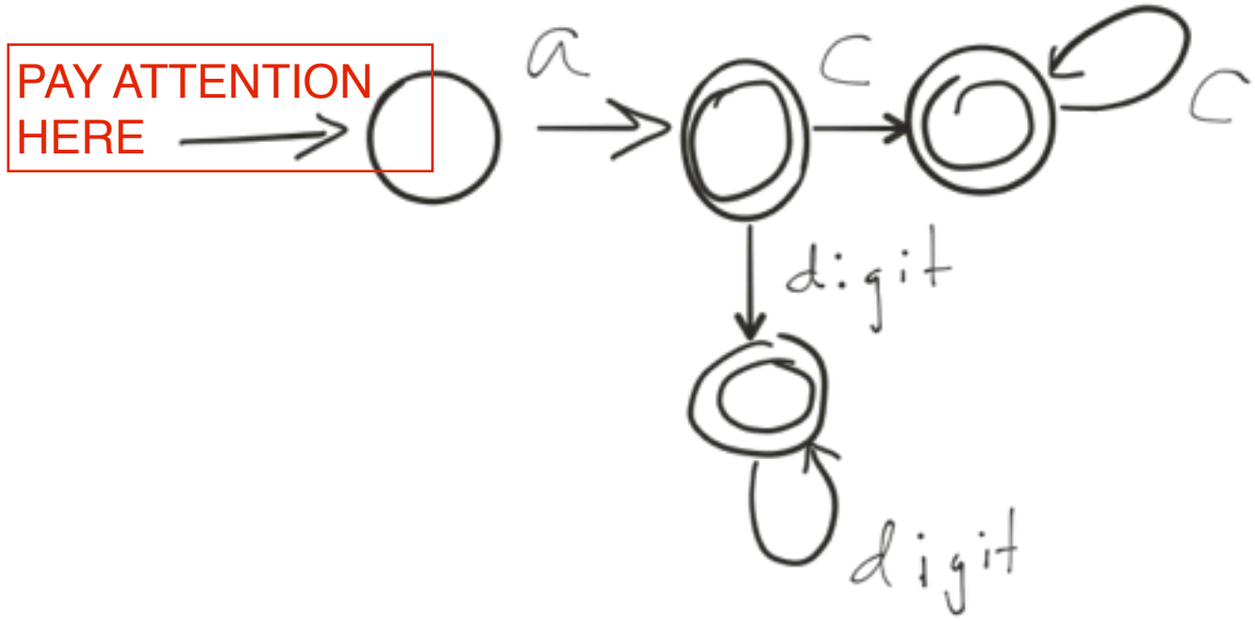
edges may be labelled with special symbol **ε** (empty string)

ε-transitions can happen without reading input

NFA example

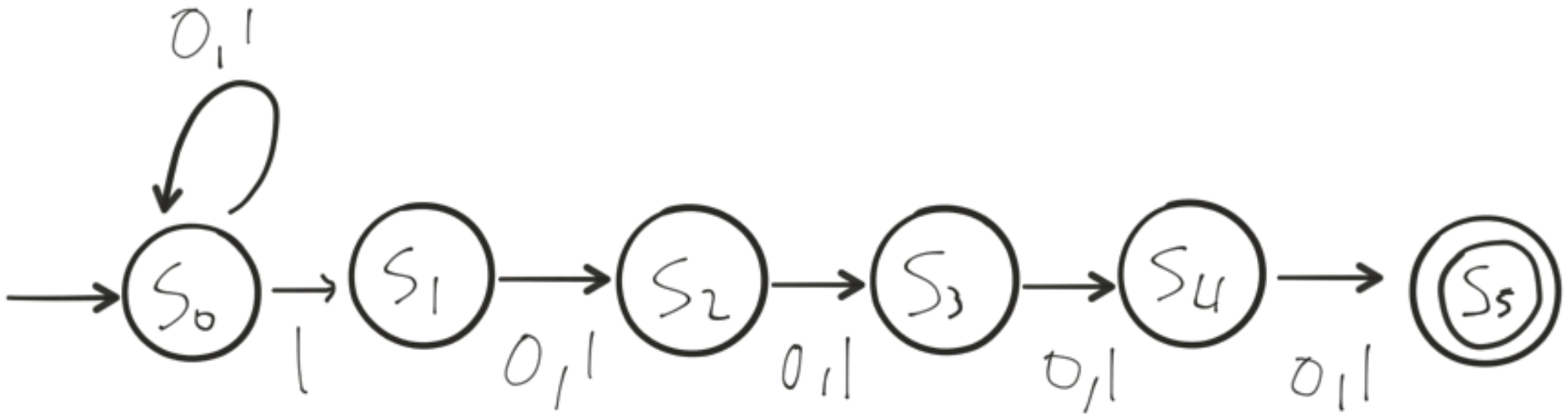


Equivalent DFA



Why NFA?

Much more compact



What does this accept?

An equivalent DFA needs 2^5 states

Extra example

Hex literals

must start with 0x or 0X

followed by at least one hex digit (0-9,a-f,A-F)

can optionally have long specifier (l,L) at the end

Extra example

A C/C++ identifier is a sequence of one or more letters, digits, or underscores. It cannot start with a digit.

What if you wanted to add the restriction that it can't end with an underscore?

automatatutor.com

Recap

The scanner reads stream of characters and finds tokens

Tokens are defined using regular expressions, which are finite-state machines

Finite-state machines can be non-deterministic

Next time: understand connection between deterministic and non-deterministic FSMs