# Data structures and algorithms Spring 2025

Lecture 8

Tree part I

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### Contents

- Tree's ADT
- Binary tree
- Tree operations
- Tree implementations part 1



# TREE

What it is (conceptual)

Why we use it (applications)

How we implement it (implementation)

### Definition [1]



**Nature View** 



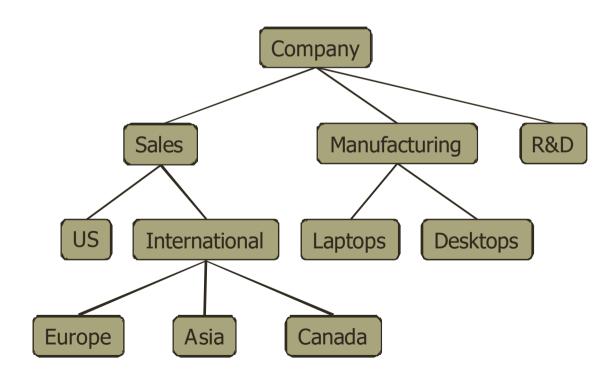
**Computer Scientists View** 

This might seem a little simplistic - but it can be quite helpfulto keep simple pictures in mind to help you workthrough the concepts and terminology.

### Definition [2]

#### Definition

 Atree (CS) is an abstract model of a hierarchical structure



### Definition [3]

- Atree consists of branches and nodes with a parent-child relation
- Example:
  - Folder structures
  - Organization structure
  - Domain name structure

```
olume serial number is 3AD5-78
   data
       ApplicationHeader
           images
       CommunityRating
       EnergyOptimizer
       Framework
       Integrator
           CommandLinks
           images
                buttons
                footer
                framework
               -layout
```

### Definition [4]

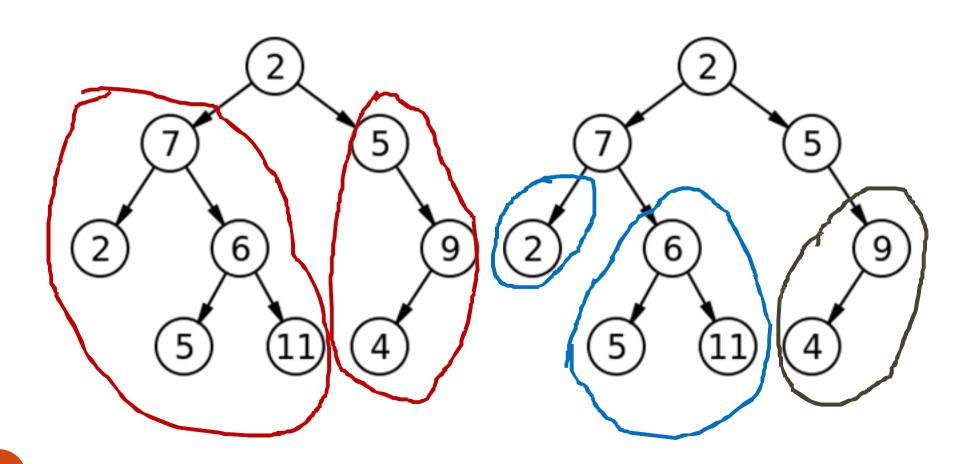
Recursive definition:

Tree (T) is a finite set of **one or more nodes** such that:

- There is one specially designated node called the root of the tree.
- The remaining nodes (excluding the root) are partitioned into m disjoint sets T1, T2, ..., Tm and each of these sets in turn is a tree.
- The trees T1, T2, ..., Tm are called the **sub-trees** of the root.

### Definition [5]

Recursive definition – Tree example



## Terminologies [1]

Aleaf node:

Anode with no children.

Abranch node (interior node):

A node which is not the root node and a leaf node.

Level of node:

The distance from the node to the root.

Depth of a tree:

The maximum level of any leaf in the tree

Degree of a node:

The number of its children

## Terminologies [2]

- Depth of a node: the number of edges from depth 0 the node to the tree's root node. height 3 - Height of a node: the number of edges from the node to the deepest possible leave. depth 1 depth 1 - Height of a tree is height of root node height 1 height 2 = Depth of a tree is depth of deepest node = maximum level of any leaf in the tree. depth 2 depth 2 height 0 height 0

root node

inner node

leaf node

depth 2

height 1

depth 3

height 0

## Terminologies [3]

#### Parent and Siblings:

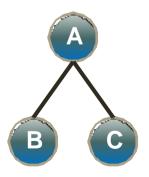
- Each root is said to be the parent of its sub-trees
- Children of the same parent are siblings.

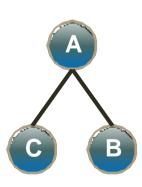
#### Ancestor and Descendant:

 Ancestor and descendant can also be used to denote the relationship that may span several level of tree.

#### Node order

- The order of the children nodes
- Default is from left to right
- Node label (data)





## Terminologies [4]

Example:

**Level of node:** State the levels of all the nodes:

A:\_\_\_\_, B:\_\_\_\_, C:\_\_\_\_,

D:\_\_\_\_\_, E:\_\_\_\_\_, F:\_\_\_\_\_,

G.\_\_\_\_, П.\_\_\_\_, I.\_\_\_\_

Root of a tree: Root of the tree is: \_\_\_\_\_

**Depth of a tree:** Depth of the tree is: \_\_\_\_\_

**Degree of a node:** State the degrees of:

A:\_\_\_\_, B:\_\_\_\_, C:\_\_\_\_,

D:\_\_\_\_, E:\_\_\_\_, F:\_\_\_\_

G:\_\_\_\_, H:\_\_\_\_, I:\_\_\_\_

**Terminal node or leaf:** State all the leaf nodes:

**Branch (interior) node:** State all the branch nodes:

# Terminologies [5]

#### **Parent and Siblings:**

State the parents of: A:\_\_\_, B:\_\_\_, C:\_\_\_,

D:\_\_\_, E:\_\_\_, F:\_\_\_,

G:\_\_\_, H:\_\_\_, I:\_\_\_

State the siblings of: A:\_\_\_\_\_, B:\_\_\_\_\_

C:\_\_\_\_\_, D:\_\_\_\_\_, E:\_\_\_\_\_,

F:\_\_\_\_\_, G:\_\_\_\_\_, H:\_\_\_\_\_, I:\_\_\_\_\_

#### **Ancestor and Descendant:**

State the ancestors of: A:\_\_\_\_\_, B:\_\_\_\_\_, C:\_\_\_\_\_, D:\_\_\_\_\_\_,

E\_\_\_\_\_\_, F:\_\_\_\_\_\_, G:\_\_\_\_\_\_\_,

H:\_\_\_\_\_\_, I:\_\_\_\_\_\_

State the descendants of: A:\_\_\_\_\_

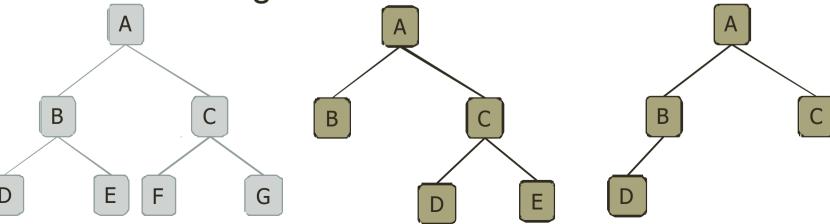
D:\_\_\_\_, E:\_\_\_\_, F:\_\_\_\_, G:\_\_\_\_, H:\_\_\_\_, I:\_\_\_\_

### Binarytree [1]

- Definition: Abinary tree is a tree that:
  - Root node and each interior node has at most two children.
  - The children of a node are an ordered pair.

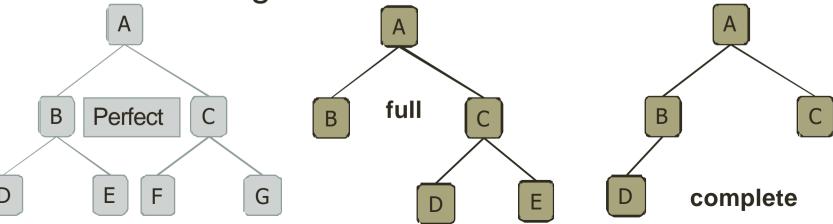
## Binarytree [2]

- Some typical types of Binary tree
  - Perfect Binary tree: every level is completely filled.
  - Full Binary tree: each node is either a leaf or has exactly two children.
  - Complete Binary tree: a perfect binary tree except perhaps for the final level which is filled from left to right.



## Binarytree [2]

- Some typical types of Binary tree
  - Perfect Binary tree: every level is completely filled.
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# Binarytree [3]

- Number of nodes in a Binary tree
  - Consider a binary tree T with the depth h, n is the total nodes of T, I is the total leaves of T
  - If Tis a perfect binary tree, then:

$$l = 2^h$$
 and  $n = 2^{h+1} - 1 = 2l - 1$ 

If T is a full binary tree, then:

$$2^h \le n \le 2^{h+1} - 1$$

### Tree traversal [1]

- Definition: Atraversal of a tree:
  - Start from the root of the tree
  - Visit every node of the tree.
  - Each node is visited once.
- Traversal order
  - Order of visiting
- Visit:
  - Do something with the node
    - Print node's label
    - Perform an operation on node's data
  - Visit a node go to a node

# Tree traversal [2]

### Example

Tree traversal to printall node's label

#### Result

- A, B, C, D
  - Start from A, visit A → go to B, visit B → go to C, visit
     C → go to B → go to D, visit D.
- C, D, B, A
  - Start from A → go to B → go to C, visit C → go to B → go to D, visit D → go to B, visit B → go to A, visit A.
- C, B, D, A
  - Start from A → go to B → go to C, visit C → go to B, visit B → go to D, visit D → go to B → go to A, visit A.

### Tree traversal [3]

#### Level-order

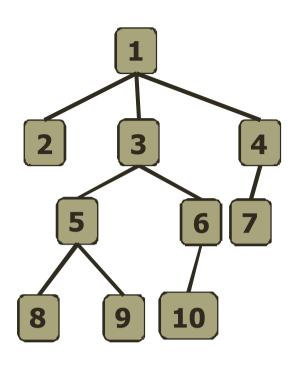
Start at root, visit root.

Visit the nodes at each level, from left to right.

$$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 7$$
$$\rightarrow 8 \rightarrow 9 \rightarrow 10$$

#### Pre-order

Visit each node, followed by its children (in pre-order) from left to right.  $1 \rightarrow 2 \rightarrow 3 \rightarrow 5 \rightarrow 8 \rightarrow 9 \rightarrow 6 \rightarrow 10 \rightarrow 4 \rightarrow 7$ 

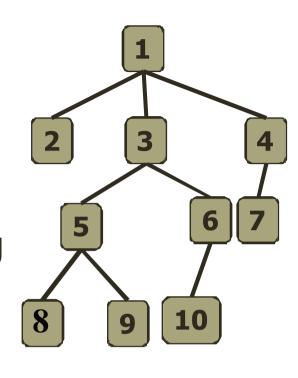


### Tree traversal [4]

#### In-order

Visit the left-most child, followed by the root, followed by the remaining children from left to right.

$$2 \rightarrow 1 \rightarrow 8 \rightarrow 5 \rightarrow 9 \rightarrow 3 \rightarrow 10 \rightarrow 6 \rightarrow 7$$
  
  $\rightarrow 4$ 



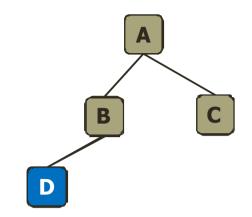
#### Post-order

Visit the left-most child, followed by the remaining children from left to right, followed by the root.

$$2 \rightarrow 8 \rightarrow 9 \rightarrow 5 \rightarrow 10 \rightarrow 6 \rightarrow 3 \rightarrow 7 \rightarrow 4$$
  
 $\rightarrow 1$ 

### Tree operations [1]

- Create an empty tree
- Add a new node to the tree



- A new node is added to the tree as a child of a parent node.
- If the parent is null, new node is the root of the tree.
- Node is added from left to right.
- Example: addNode(D,B)

### Tree operations [2]

#### Get root node of the tree

- Return the root node of the tree.
- Return null if the tree is empty.

### Get the parent node of a node in the tree

- Return the parent node of a node n in the tree.
- Return null if the tree is empty or n is the root

#### Get the left most child of a node in the tree

- Return the left most child of a node n in the tree
- Return null if the tree is empty or n has no child

### Tree operations [3]

- Get the nearest right sibling of a node in the tree
  - Return the nearest right sibling of a node n in the tree.
  - Return null if the tree is empty or n has no right sibling.
- Get/Set node's label
  - Get/set the label of a node in the tree.
- Check if the tree isempty
  - Return true if the number of nodes =0, otherwise return false

### Tree ADT

```
Tree
+addNode(NodeType newNode, NodeTypeparent):
void
+getRoot(): NodeType
+getParent(NodeType n): NodeType
+leftMostChild(NodeType n): NodeType
+rightSibling(NodeType n): NodeType
+getNodeLabel(NodeType n): LabelType
+setNodeLabel(NodeType n, LabelType label): void
+isEmpty(): boolean
```



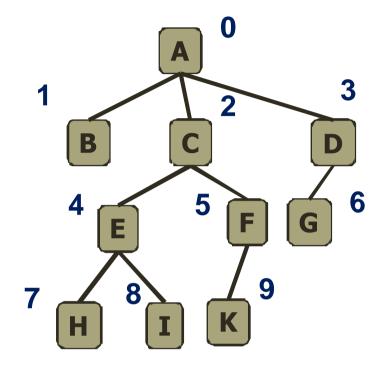
# Tree Implementation

### Array-based tree[1]

- Array is used to represent a tree
  - Atree Thas N nodes, each node is indexed by a number from 0 to N-1.
  - With node i in the tree
    - Lij is the label of node i
    - P[i] is the parent of node I
- Indexing scheme:
  - Incremental indexing started from root
  - Root is 0, parent is smaller than children
  - Children of the same parentare indexed from left to right

### Array-based tree[2]

Example



0 1 2 3 4 5 6 7 8 9

A B C D E F G H I K

-1 0 0 0 2 2 3 4 4 5

L:

P:

### Array-based tree[3]

```
Tree
-maxSize: int //maximum possible number of nodes
-n: int //current number of nodes in the tree
-I: String //Label array
-p: int[] //Parent array
+addNode(String label, int parent): void
+getParent(int node): int
+leftMostChild(int node): int
+rightSibling(int node): int
+getNodeLabel(int node): String
+setNodeLabel(int node, String label): void
```

### Array-based tree[4]

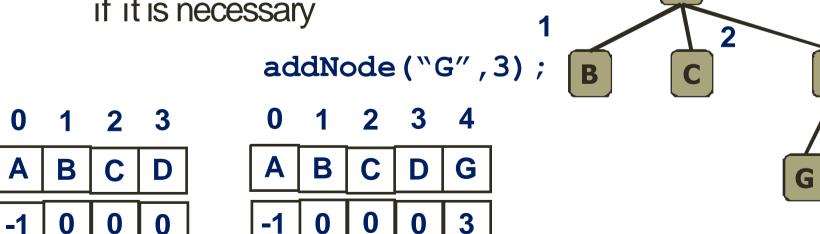
- addNode(String label, int parent)
  - If parent is -1, then we will add a root node
     1[0]=label; p[0]=-1;
  - If parent is different from -1
    - Must find a correct position for the new node
    - Shifting the array to the right if it is necessary

```
addNode("G",3);
```

```
0 1 2 3
A B C D
-1 0 0 0
```

### Array-based tree[4]

- addNode(String label, int parent)
  - If parent is -1, then we will add a root node
     1[0]=label; p[0]=-1;
  - If parent is different from -1
    - Must find a correct position for the new node
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### Array-based tree[5]

- addNode(String label, int parent)
  - If parent is different from -1
    - Must find a correct position for the new node
    - Shifting the array to the right if it is necessary

```
0 1 2 3 4

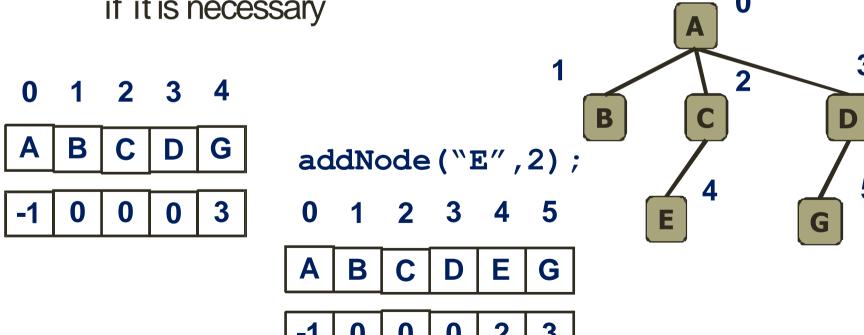
A B C D G

addNode("E",2);

1 0 0 0 3
```

## Array-based tree[5]

- addNode(String label, int parent)
  - If parent is different from -1
    - Must find a correct position for the new node
    - Shifting the array to the right if it is necessary



### Array-based tree[6]

- leftMostChild(int node)
  - Example:

```
leftMostChild(0) is 1
leftMostChild(3) is 4
leftMostChild(2) is -1
```

- Start from node+1
  - Find the first node i that:

- Return -1 if could not find i

0	1	2	3	4
Α	В	С	D	G
-1	0	0	0	3

### Array-based tree[6]

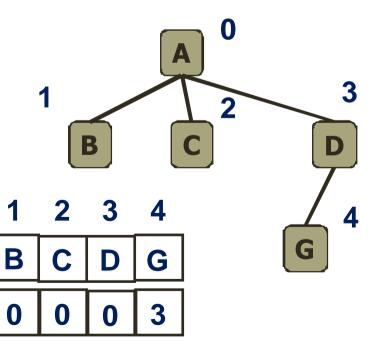
- leftMostChild(int node)
  - Example:

```
leftMostChild(0) is 1
leftMostChild(3) is 4
leftMostChild(2) is -1
```

- Start from node+1
  - Find the first node i that:

#### p[i]==node

- Return -1 if could not find i



### Array-based tree[7]

- nearestRightSibling(int node)
  - Example:

```
nearestRightSibling(1) is 2
nearestRightSibling(2) is 3
nearestRightSibling(3) is -1
```

- Check node+1
  - Return **node+1** if:

- Return -1 otherwise

0	1	2	3	4
Α	В	С	D	G
-1	0	0	0	3

### Array-based tree[7]

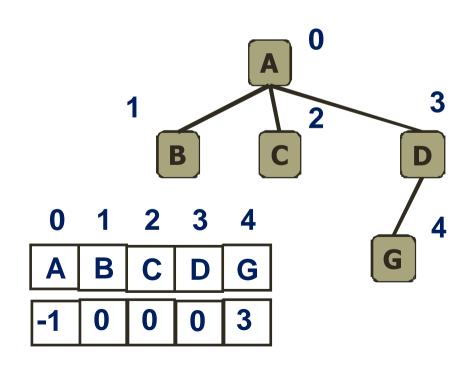
- nearestRightSibling(int node)
  - Example:

```
nearestRightSibling(1) is 2
nearestRightSibling(2) is 3
nearestRightSibling(3) is -1
```

- Check node+1
  - Return node+1 if:

#### p[node+1]==p[node]

- Return -1 otherwise



### Array-based binarytree [1]

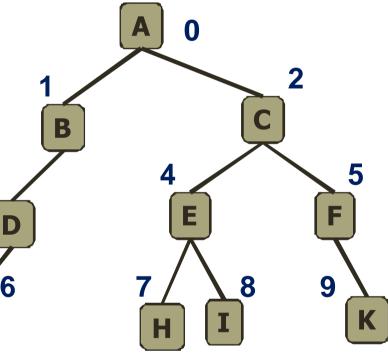
Array is used to represent a binary tree

 Using the above same indexing scheme



L[i] is the label of node i<sup>3</sup>

 P[i] is the parent of node i



0 1 2 3 4 5 6 7 8 9

L: A B C D E F G H I K

P: -1 0 0 1 2 2 3 4 4 5

## Array-based binarytree [2]

Array is used to represent a binary tree

Using perfect binary tree indexing scheme

- With node i in the tree
  - L[i] is the label of nodei
  - Left child: 2i+1
  - Right child: 2i+2



0



### Array-based binarytree [3]

```
BinaryTree
-maxSize: int //maximum possible number of nodes
-n: int // Current number of nodes
+addRoot(String label): void
+getLeftChild(int node): int
+getRightChild(int node): int
+getParent(int node): int
+addLeftChild(int node): int
+addRightChild(int node): int
+getNodeLabel(int node): String
+setNodeLabel(int node, String label): void
+preOrderTravel(int node): void
```

### Tutorial & next topic

### Preparing for the tutorial:

 Practice with examples and exercises in Tutorial 8 Example Code

### Preparing for next topic:

- Read textbook chapter 4 (4.1 4.3): Tree data structure.
- Read supplementary book chapter 10 (10.4) and chapter 12