



## can you give a more compolete description of your proposed algorithm for PopperSAT?

PopperSAT (for your fragment: finite additivity + full product axiom on normal  $\psi$  + “abnormal  $\Rightarrow$  all conditionals = 1”, quantifier-free polynomial (in)equalities over  $P(\varphi \mid \psi)$ ) can be implemented as a **finite outer search** over which conditioning events are normal (and at what “rank”), coupled with an inner **QF\_NRA** check that looks very close to PrSAT’s stochastic truth-table encoding.<sup>[1]</sup> <sup>[2]</sup> <sup>[3]</sup>

### 1) Semantic target (what we search for)

Work over the finite space of classical valuations  $W$  for  $n$  atoms, so  $|W| = 2^n$ , exactly as in the stochastic truth-table setup used by PrSAT.<sup>[1]</sup>

Use a **finite lexicographic probability system** (LPS)  $(\mu_1, \dots, \mu_K)$ , where each  $\mu_\ell$  is an ordinary finitely additive probability measure on  $W$ , and define  $P(\varphi \mid \psi)$  by selecting the first layer where  $\mu_\ell(\psi) > 0$ .<sup>[4]</sup> <sup>[3]</sup>

In this representation, when  $\psi$  is normal at the selected layer,  $P(\neg \mid \psi)$  is literally a ratio inside  $\mu_\ell$ , so finite additivity and the Popper product axiom hold “for free”; when  $\psi$  is abnormal, Popper-style abnormality makes  $P(\neg \mid \psi)$  trivial, and under your convention this is  $P(\varphi \mid \psi) = 1$  for all  $\varphi$ .<sup>[2]</sup> <sup>[3]</sup>

### 2) Inputs and preprocessing

Input: a conjunction of polynomial equalities/inequalities over symbols  $P(\varphi \mid \psi)$ , where  $\varphi, \psi$  are propositional formulas over  $n$  atoms.

Precompute, for every propositional formula  $\theta$  that occurs in the input (and every  $\varphi \wedge \psi$ ), the set of worlds  $S_\theta = \{w \in W : w \models \theta\}$  (a bitmask is convenient).

Let  $\Psi = \{\psi_1, \dots, \psi_m\}$  be the set of distinct conditioning formulas appearing in the input.

### 3) Inner SMT encoding (division-free, conjunctive)

Choose a layer bound  $K$ . A safe, simple choice is  $K = m$  (you only need enough layers to distinguish the finitely many conditioning contexts that actually appear).

Introduce reals  $w_{\ell,i}$  for  $\ell = 1..K, i = 1..|W|$ , intended as  $\mu_\ell(\{w_i\})$ , with:

- $w_{\ell,i} \geq 0$
- $\sum_i w_{\ell,i} = 1$  for each  $\ell$

This is exactly the PrSAT “stochastic truth-table” move, just replicated across layers.<sup>[1]</sup>

For each conditional term  $P(\varphi \mid \psi_j)$  that appears, introduce a real variable  $p_{\varphi|\psi_j}$  with  $0 \leq p_{\varphi|\psi_j} \leq 1$ .<sup>[2]</sup>

Define (as abbreviations you expand into sums) for each  $j, \ell$ :

- $d_{j,\ell} := \sum_{i \in S_{\psi_j}} w_{\ell,i}$  (denominator mass of  $\psi_j$  at layer  $\ell$ )
- $n_{j,\ell}(\varphi) := \sum_{i \in S_{\varphi \wedge \psi_j}} w_{\ell,i}$  (numerator mass)

Then, given a choice of which layer is “active” for  $\psi_j$ , you add **only polynomial** constraints, no division:

- Normal at layer  $k_j$ :  $n_{j,k_j}(\varphi) = p_{\varphi|\psi_j} \cdot d_{j,k_j}$
- Abnormal:  $p_{\varphi|\psi_j} = 1$  (your convention)<sup>[2]</sup>

Finally, translate the user’s polynomial constraints by replacing each  $P(\varphi \mid \psi_j)$  with  $p_{\varphi|\psi_j}$ .

This inner problem is a conjunction of polynomial (in)equalities over reals (QF\_NRA), in the same spirit as PrSAT’s translation for Kolmogorov conditionals.<sup>[1]</sup>

#### 4) Outer search (the only “non-PrSAT” part)

You still need to connect each  $\psi_j$  to either:

- an **active layer index**  $k_j \in \{1, \dots, K\}$ , meaning “ $\psi_j$  is normal first at layer  $k_j$ ”, or
- **abnormal**, meaning it has zero mass at every layer.

Encode that by branching externally (multiple Z3 calls), so each call stays purely conjunctive:

If  $\psi_j$  is normal first at  $k_j$ , assert:

- $d_{j,\ell} = 0$  for all  $\ell < k_j$
- $d_{j,k_j} > 0$

If  $\psi_j$  is abnormal, assert:

- $d_{j,\ell} = 0$  for all  $\ell \leq K$

This “first positive layer” mechanism is exactly what makes LPSs correspond to Popper functions on finite spaces.<sup>[3]</sup>

#### Practical branching strategy

Naively you have  $(K + 1)^m$  branches, but you can prune heavily by using logical relations between conditioning events before calling Z3:

- If  $\psi \models \chi$ , then  $\text{first}(\psi)$  cannot be *later* than  $\text{first}(\chi)$  (since any layer giving  $\psi$  positive mass also gives  $\chi$  positive mass).
- If  $\psi$  is unsatisfiable, force it abnormal immediately.
- Merge syntactically different but logically equivalent  $\psi$ ’s.

Implementation-wise, do a DFS that incrementally assigns layer/abnormal statuses and uses Z3 push/pop to prune early when partial assignments already make the QF\_NRA constraints inconsistent.

## 5) Soundness/completeness intuition (for your axioms)

Soundness: any satisfying assignment to the  $w_{\ell,i}$  defines  $K$  finitely additive measures, and the “first layer” constraints make each  $\psi$  either normal (ratio-defined) or abnormal (constant-1), matching your convention.<sup>[3] [2]</sup>

On normal  $\psi$ , conditionals are computed as ratios in one underlying probability measure, so the finite additivity and Popper product axiom behavior comes along automatically, rather than being separately axiomatized in SMT.<sup>[2]</sup>

Completeness (finite case): for the finite  $W$  you care about, LPS/Popper representation results ensure you’re not losing models by searching in this form.<sup>[4] [3]</sup>

## 6) Output (model extraction)

If Z3 returns sat on some branch, you read off:

- The layers  $\mu_\ell(w_i) = w_{\ell,i}$
- For each  $\psi_j$ , whether it’s abnormal or its first layer  $k_j$
- Conditional values  $P(\varphi \mid \psi_j) = p_{\varphi|\psi_j}$

This gives you a concrete Popper model in essentially the same “stochastic truth-table” style as PrSAT, just with multiple layers.

If you want, I can take your earlier  $n = 2$  example and show the exact PopperSAT SMT-LIB generated for each of the four (A normal/abnormal)  $\times$  (B normal/abnormal) branches, so you can see the full end-to-end pipeline in the simplest nontrivial case.

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