



Iran University of Science & Technology
IUST

Digital Logic Design

Hajar Falahati

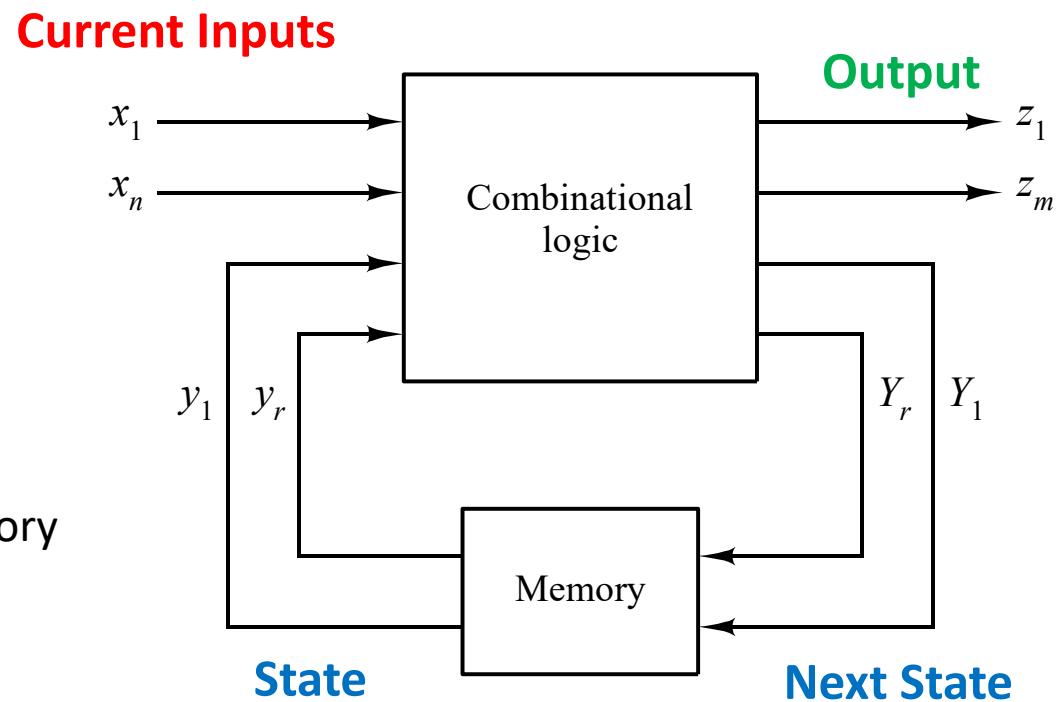
Department of Computer Engineering
IRAN University of Science and Technology

hfalahati@iust.ac.ir

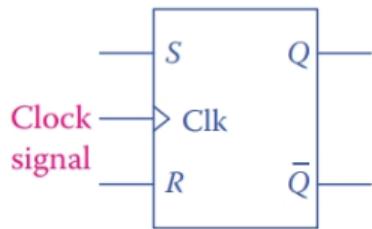
Combinational Logics + Memory

- **Output** of some logics **changes** by
 - Sequence of **inputs**
 - Time of applying inputs

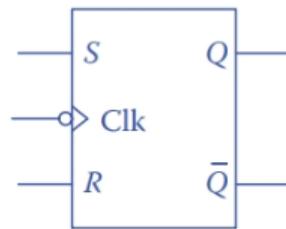
- => They **have memory**
- They **contain feedback lines**
- State
 - Produced by previous inputs
 - Information stored in a memory



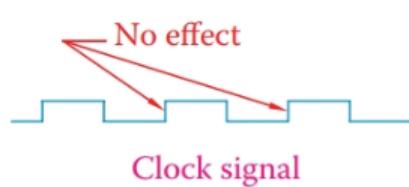
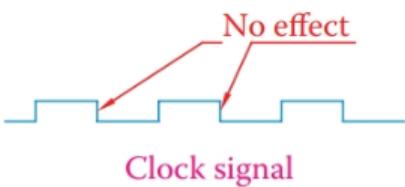
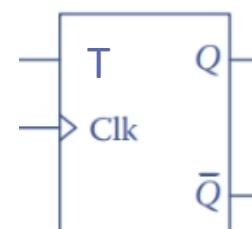
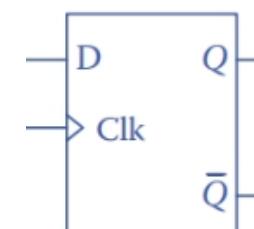
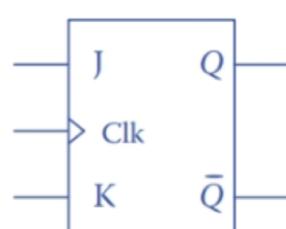
Flip Flops



Positive edge sensitive flip-flop



Negative edge sensitive flip-flop



Hold the value
Set the value
Reset the value
Invalid

	R	S	Q	\bar{Q}	J	K	Q	\bar{Q}
Hold the value	0	0	Q	\bar{Q}	0	0	Q	\bar{Q}
Set the value	0	1	1	0	0	1	0	1
Reset the value	1	0	0	1	1	0	1	0
Invalid	1	1	Invalid		1	1	\bar{Q}	Q

Hold the value
Reset the value
Set the value
Toggle the value

	D	Q	\bar{Q}	T	Q	\bar{Q}
Hold the value	0	0	1	0	Q	\bar{Q}
Toggle the value	1	1	0	1	\bar{Q}	Q

Delay FF

Outline

- Modeling of Synchronous Sequential Circuits
- Analysis of Synchronous Sequential Circuits
- Synthesis of Synchronous Sequential Circuits



Analysis

Analysis

- Determining the **functional** relation between

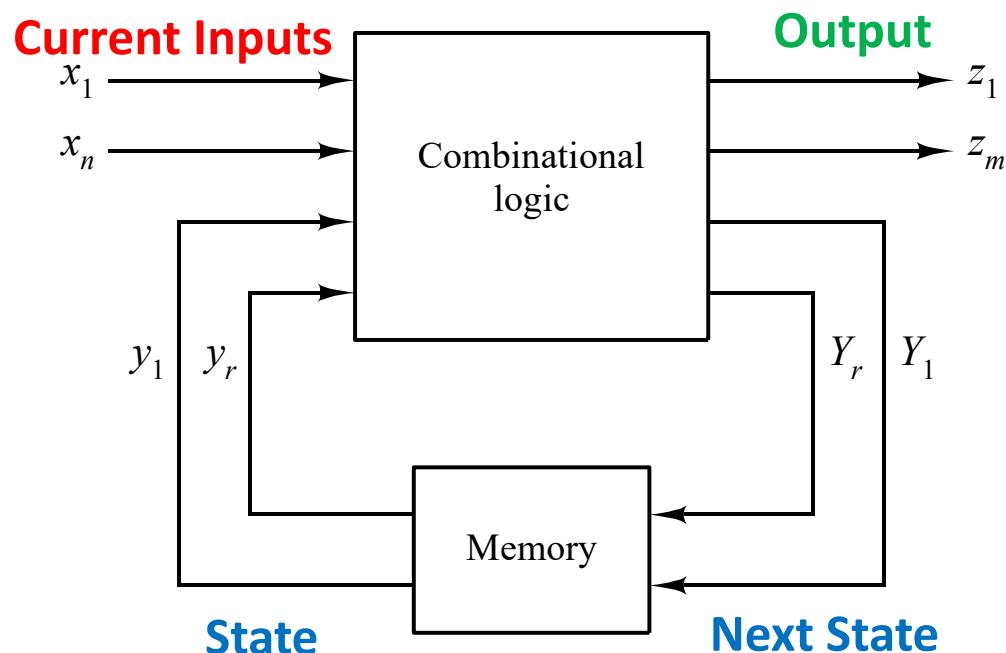
- Outputs**
 - E.g., $y(t)$
- Internal states** (content of FFs)
 - E.g., $A(t), B(t)$

- Determining the **next state**

- Internal states**
 - E.g., $A(t+1), B(t+1)$

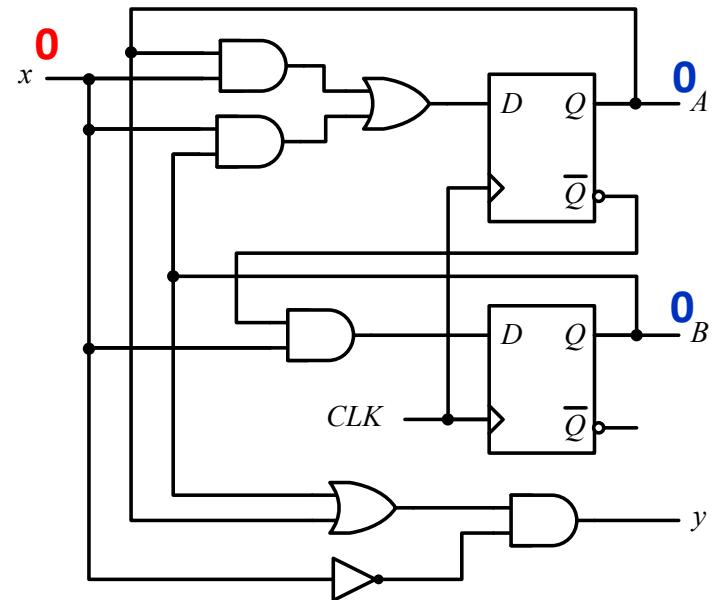
State transition

- Moving from one memory state to another one
- Under the control of a set of inputs and at a time determined by an external **clock**
- Next value of states only determine at the clock transition**



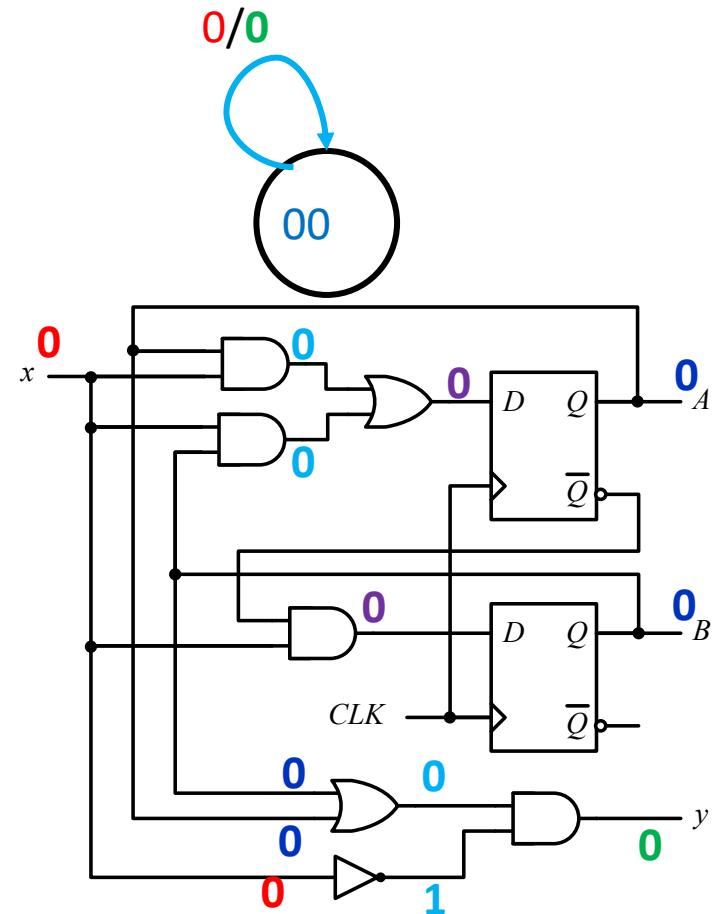
Analysis: Transition Table

Present State		Input	Next State		Output
A	B	x	A	B	y
0	0	0			
0	0	1			
0	1	0			
0	1	1			
1	0	0			
1	0	1			
1	1	0			
1	1	1			



Analysis: Transition Table

Present State		Input	Next State		Output
A	B	x	A	B	y
0	0	0	0	0	0
0	0	1			
0	1	0			
0	1	1			
1	0	0			
1	0	1			
1	1	0			
1	1	1			



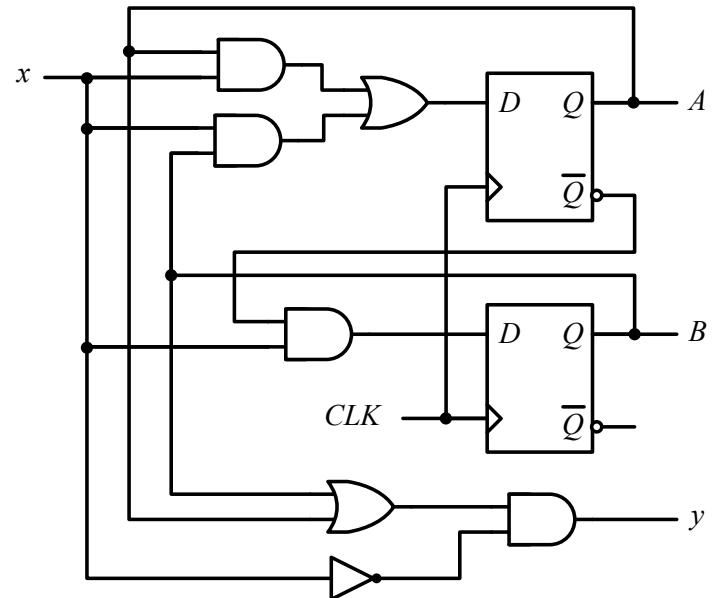
Analysis: Transition Table

Present State		Input	Next State		Output
A	B	x	A	B	y
0	0	0	0	0	0
0	0	1	0	1	0
0	1	0	0	0	1
0	1	1	1	1	0
1	0	0	0	0	1
1	0	1	1	0	0
1	1	0	0	0	1
1	1	1	1	0	0

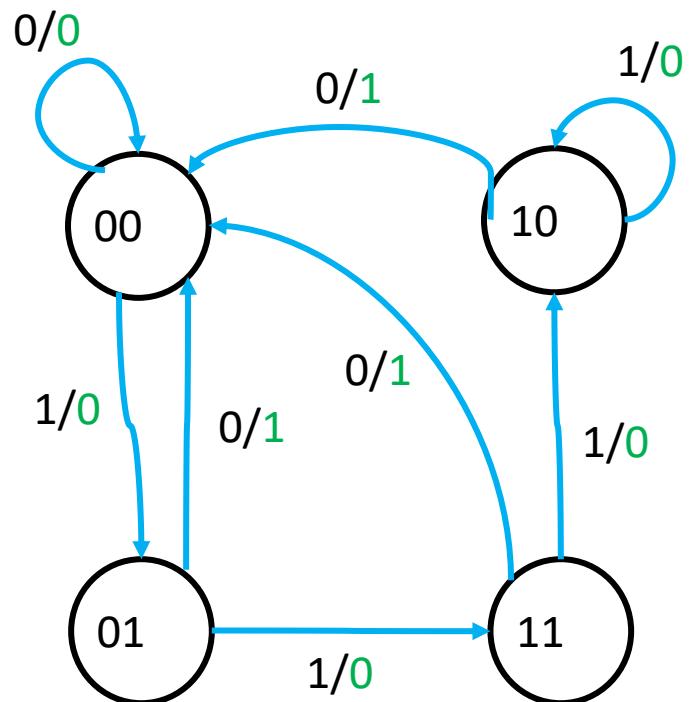
$$A(t+1) = D_A = A(t).x(t) + B(t).x(t)$$

$$B(t+1) = D_B = \overline{A(t)}.x(t)$$

$$y(t) = \overline{x(t)}.(A(t) + B(t))$$



Analysis: State Diagram

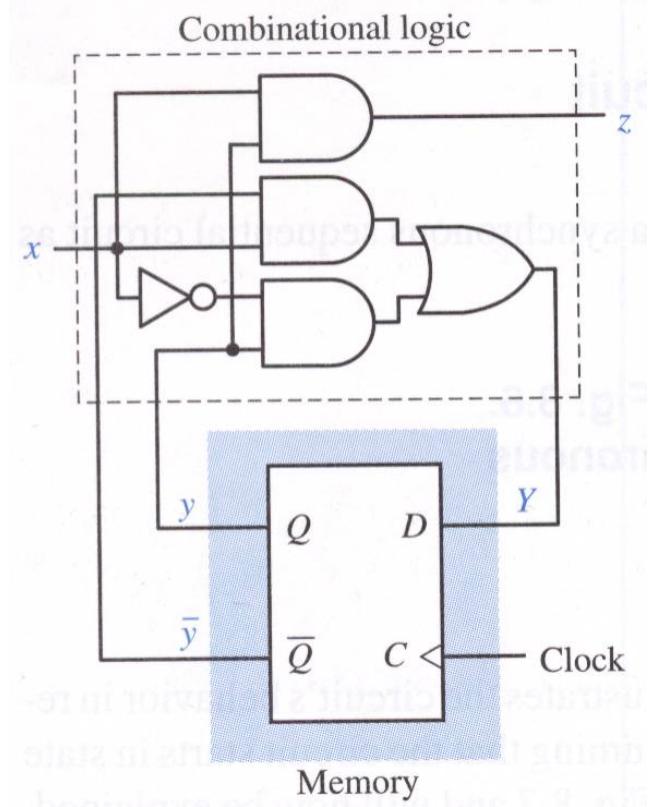


Present State		Input	Next State		Output
A	B	x	A	B	y
0	0	0	0	0	0
0	0	1	0	1	0
0	1	0	0	0	1
0	1	1	1	1	0
1	0	0	0	0	1
1	0	1	1	0	0
1	1	0	0	0	1
1	1	1	1	0	0

Analysis Sample1: Transition Table

$$Y(t + 1) = D = \overline{Y(t)}.x(t) + Y(t).\overline{x(t)}$$

$$z(t) = Y(t).x(t)$$

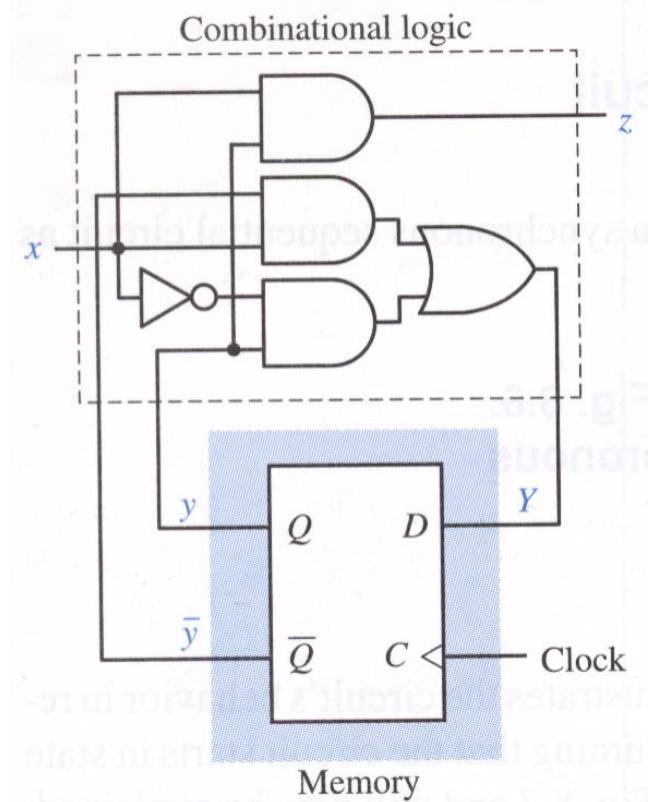


Analysis Sample1: Transition Table

$$Y(t + 1) = D = \overline{Y(t)}.x(t) + Y(t).\overline{x(t)}$$

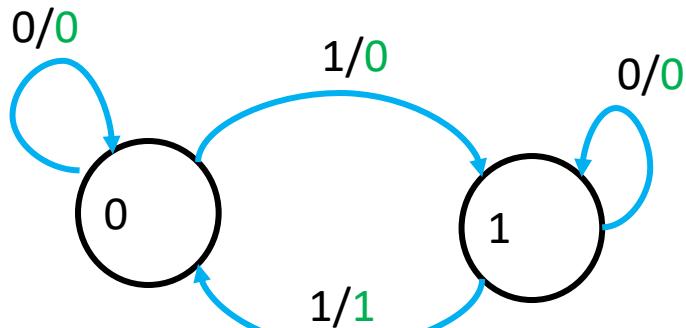
$$z(t) = Y(t).x(t)$$

Present State	Input	Next State	Output
Y	x	Y	z
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1



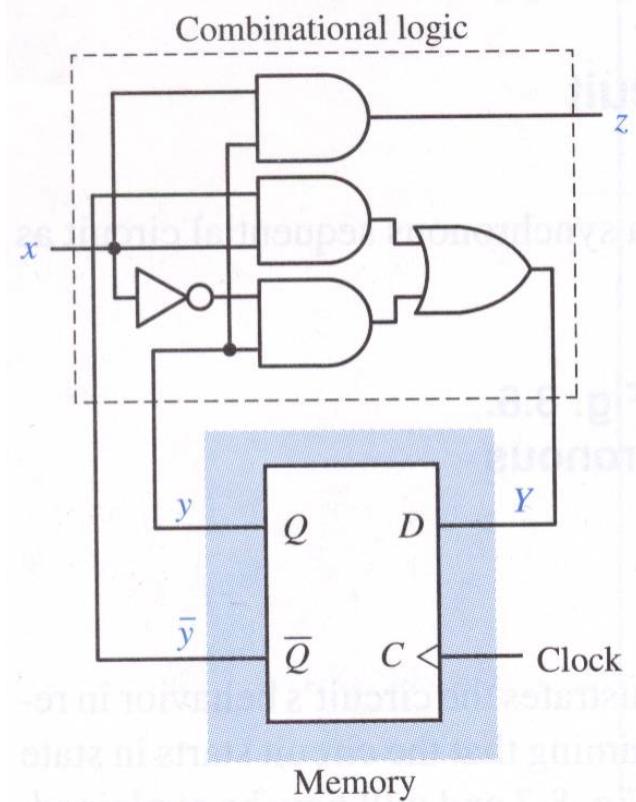
Analysis Sample1: State Diagram

Present State	Input	Next State	Output
Y	x	Y	z
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

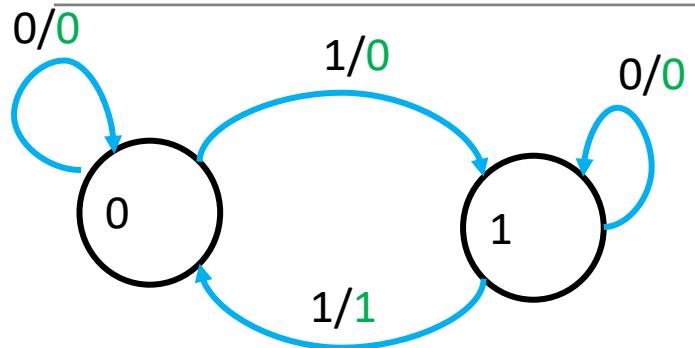


$$Y(t + 1) = D = \overline{Y(t)}.x(t) + Y(t).\overline{x(t)}$$

$$z(t) = Y(t).x(t)$$

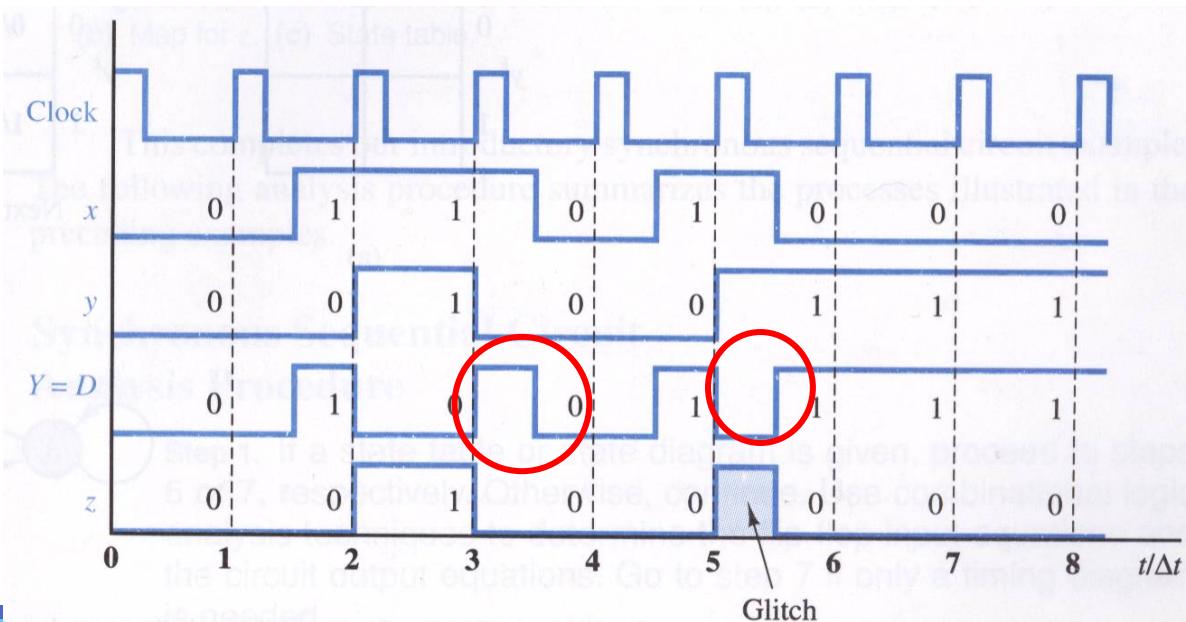
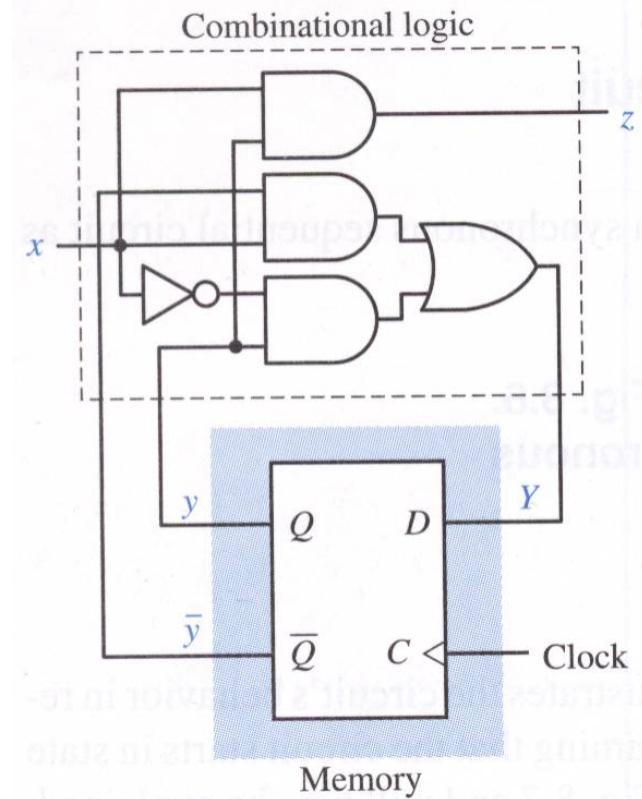


Analysis Sample1: Timing Diagram



$$Y(t+1) = D = \overline{Y(t)}.x(t) + Y(t).\overline{x(t)}$$

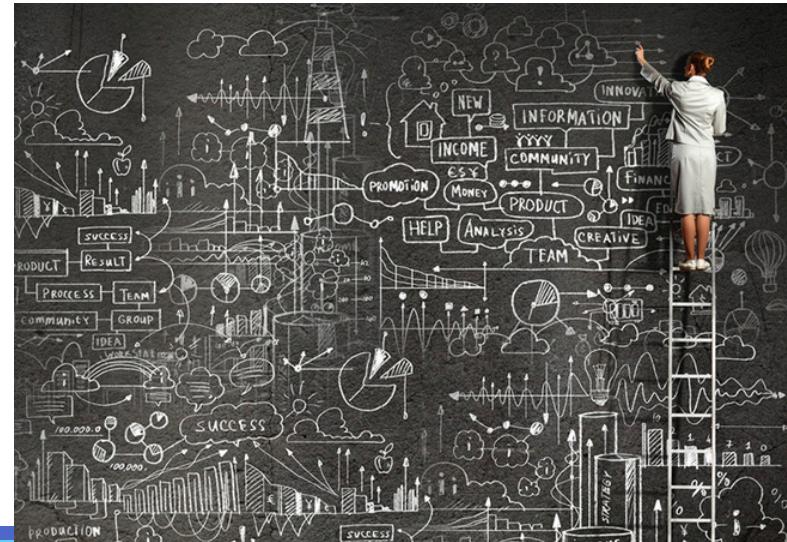
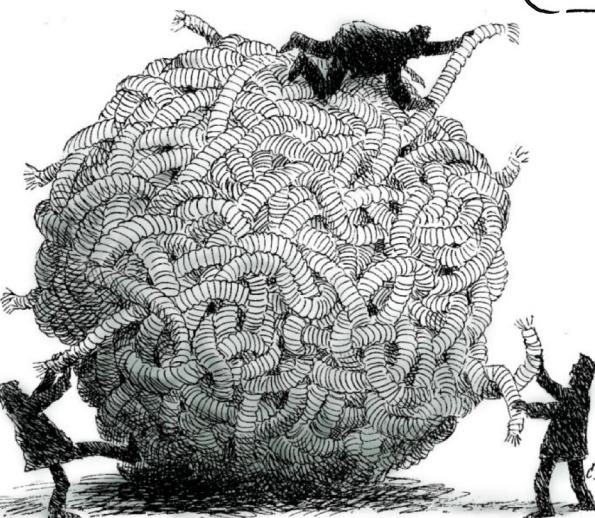
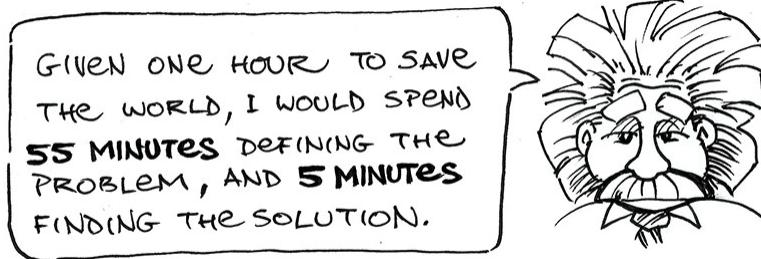
$$z(t) = Y(t).x(t)$$



Models

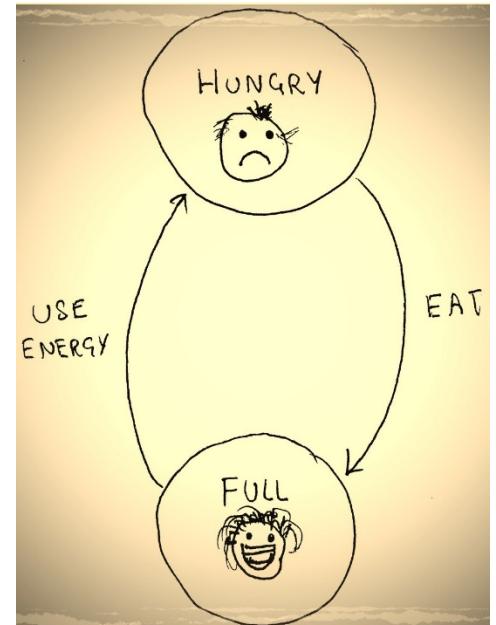
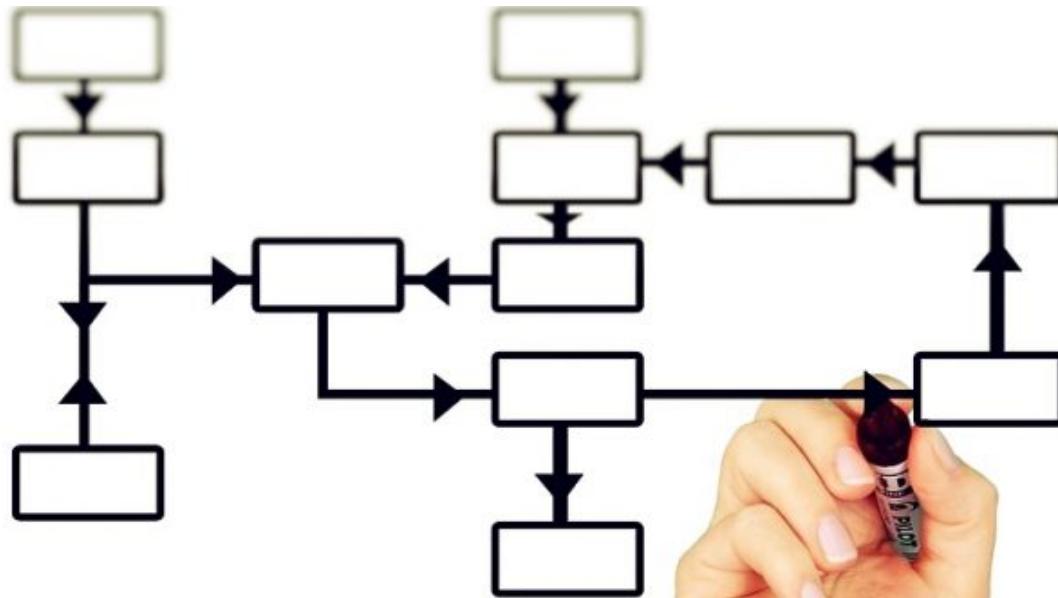
How to Implement Real Designs?

- Identify the problem
- Model the problem



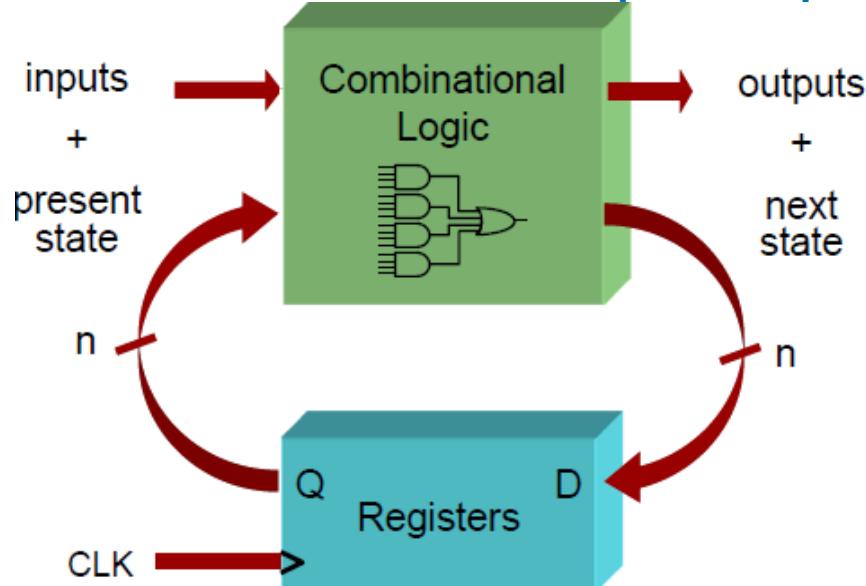
Modeling Types

- Finite State Machine
 - FSM
- Algorithmic State Machine
 - ASM



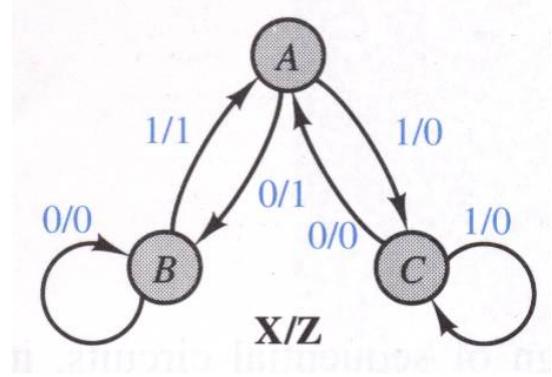
Finite State Machine (FSM)

- A system that visits a **finite** number of **logically distinct states**
- A useful **abstraction** for **sequential circuits** with **centralized states of operation**
- At each **clock edge**, combinational logic computes
 - **Outputs** and **next state** as a function of **inputs** and **present state**



FSM Representation

- State diagram & state table
 - Simple digital systems
 - A few number of inputs and outputs

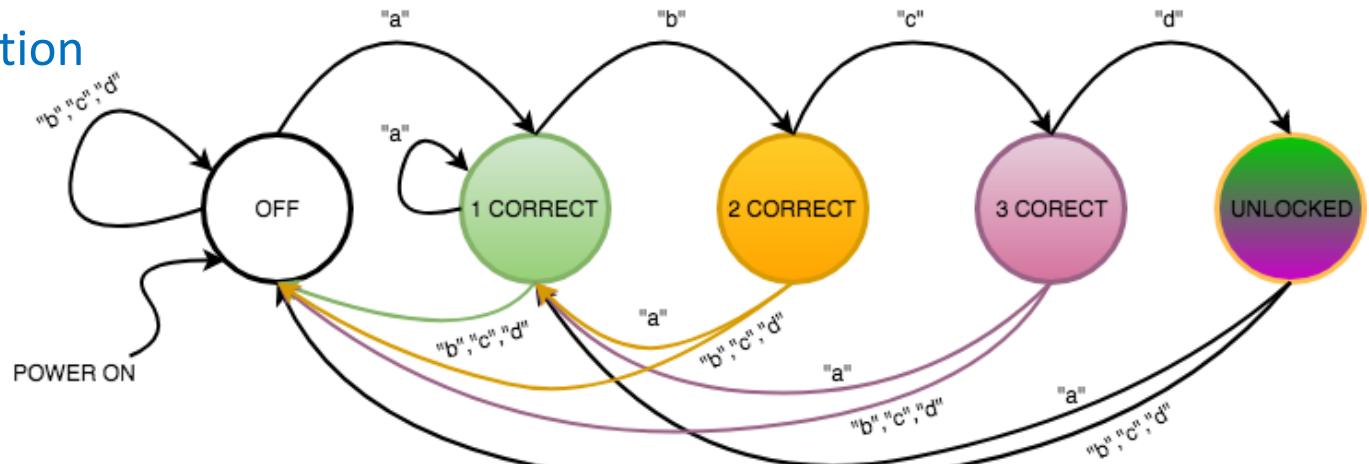


Present state	Input x	
	0	1
A	B/1	C/0
B	B/0	A/1
C	A/0	C/0

Next state/output

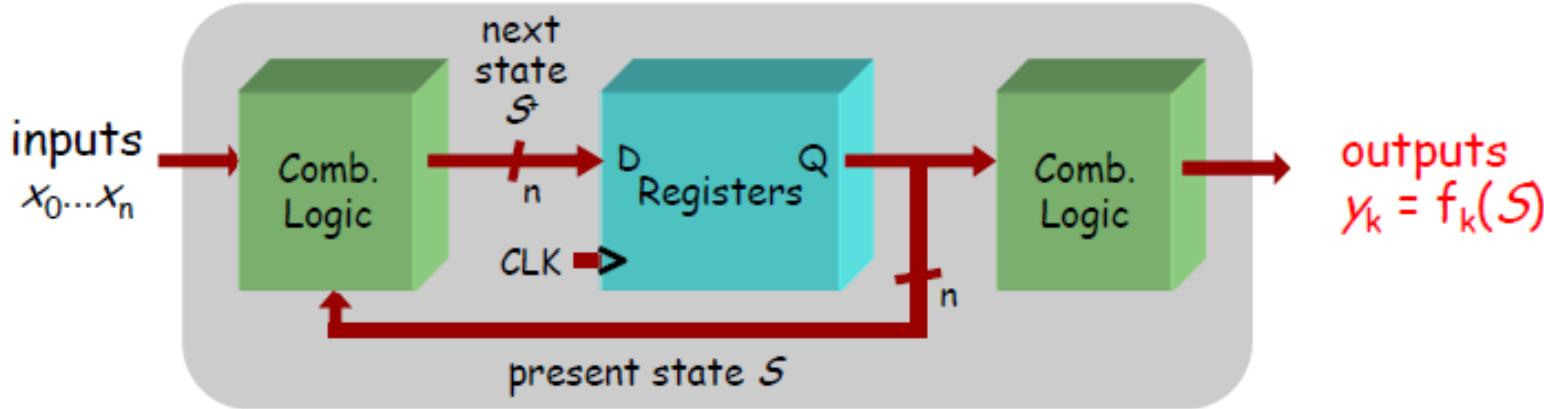
State Diagram

- A set of states
 - Nodes in a graph
- A set of inputs and outputs
 - Edges in a graph
- A set of state transition function
 - Edges in a graph
- An output function



FSM Types

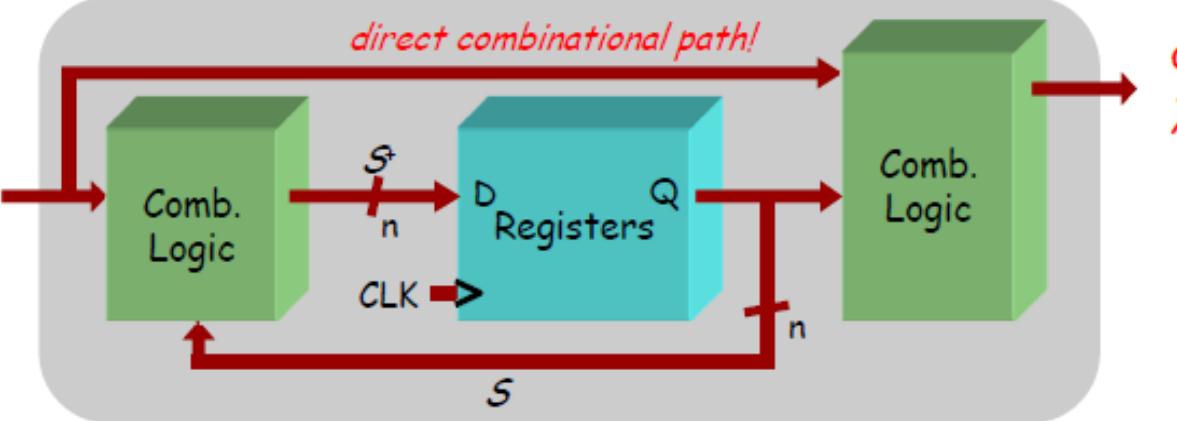
Moore FSM



$$S' = g(S, X_0 \dots X_n)$$

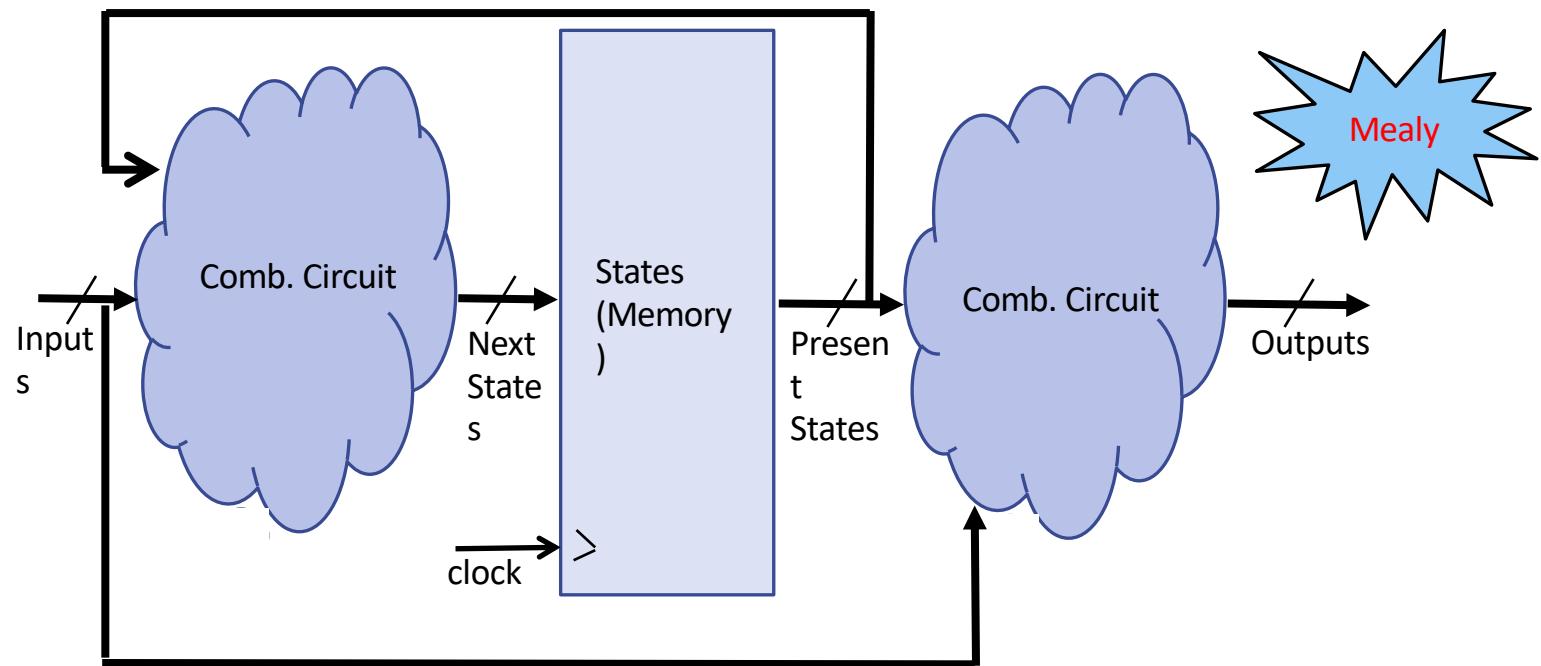
inputs
 $x_0 \dots x_n$

direct combinational path!

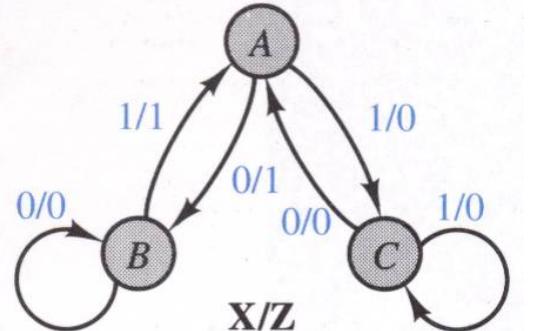


Mealy FSM

Moore and Mealy Machines

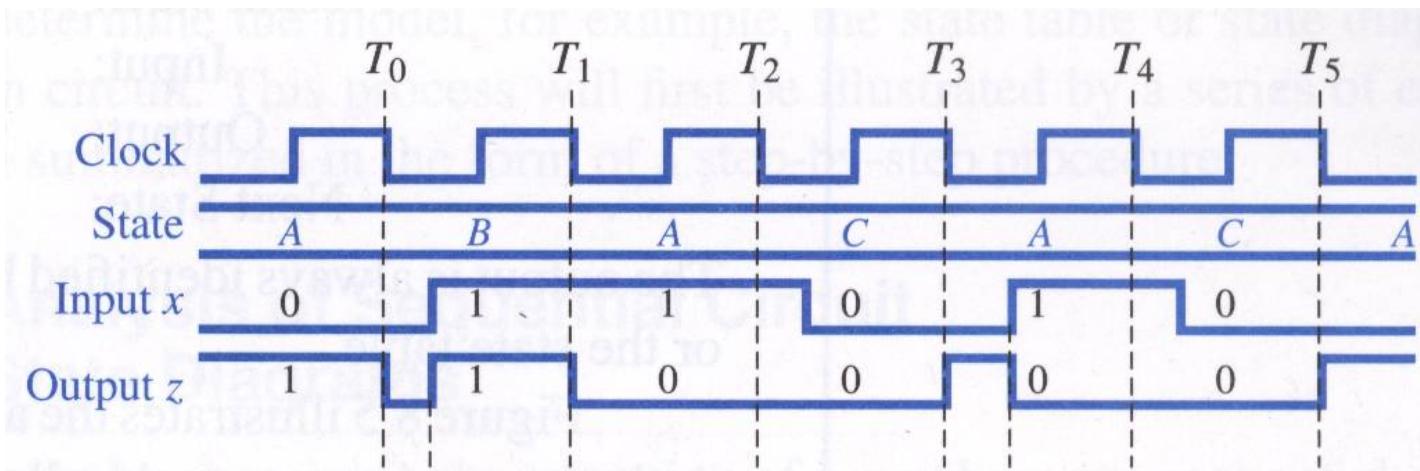


Mealy FSM: State Diagram & State Table

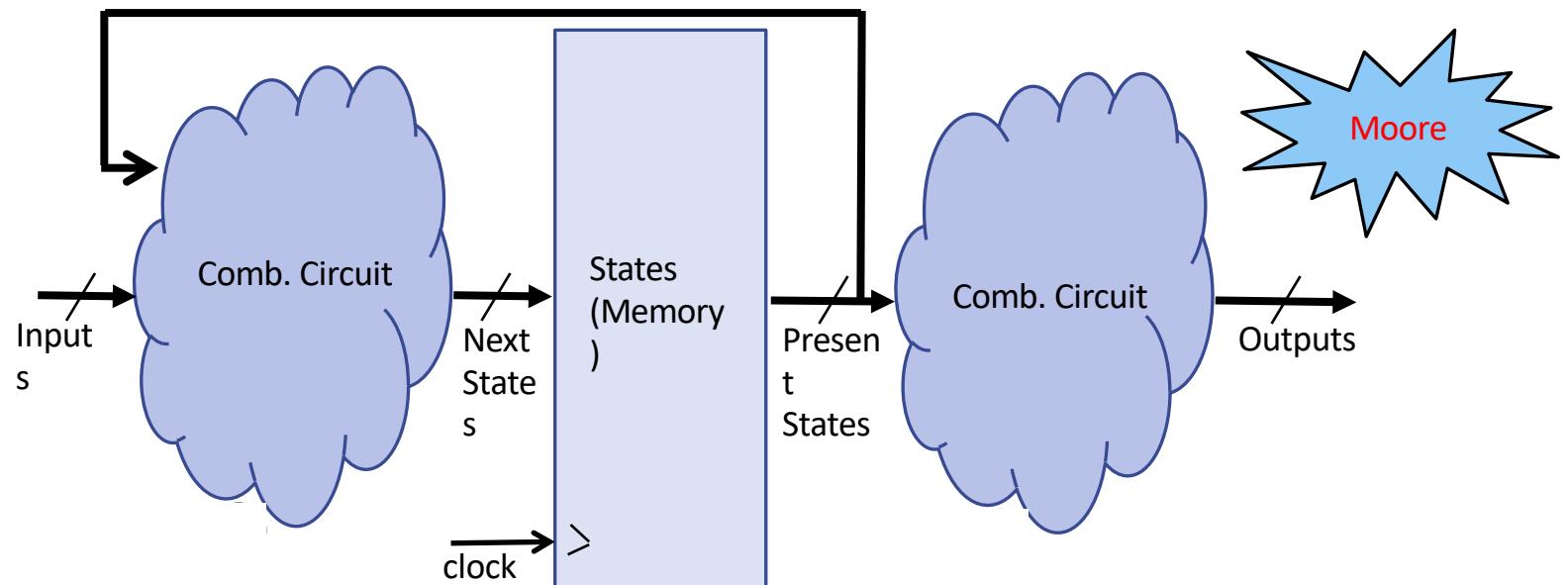


Present state	Input x		Next state/output
	0	1	
A	B/1	C/0	
B	B/0	A/1	
C	A/0	C/0	

Time:	0	1	2	3	4	5
Present state:	A	B	A	C	A	C
Input:	0	1	1	0	1	0
Output:	1	1	0	0	0	0
Next state:	B	A	C	A	C	A

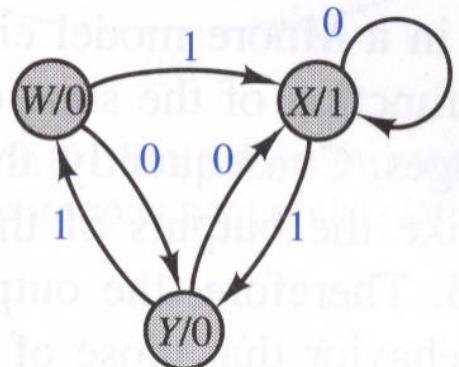


Moore and Mealy Machines



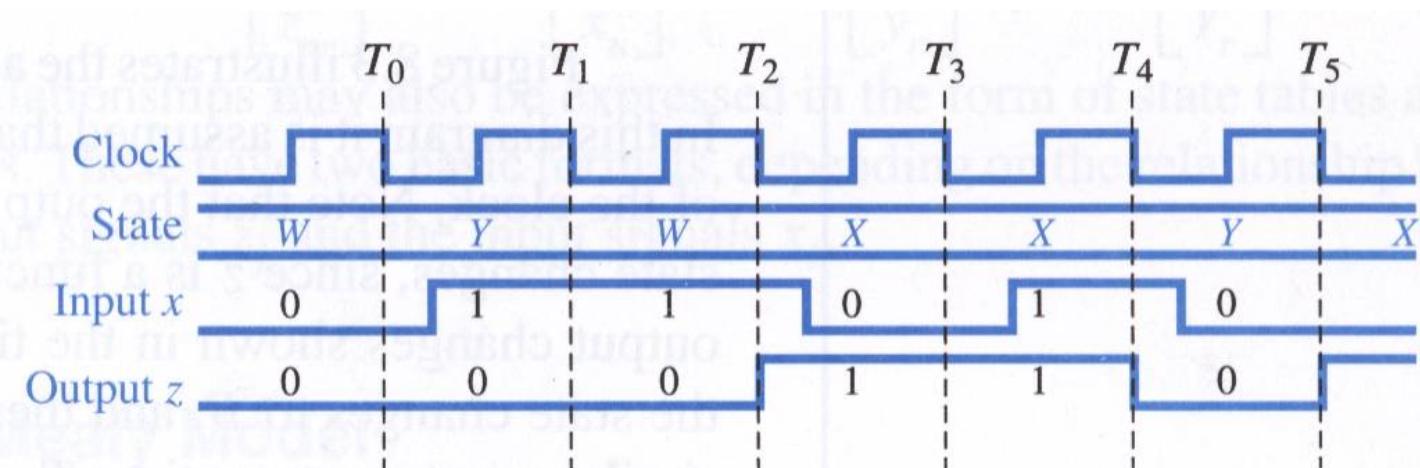
Moore FSM:

State Diagram & State Table



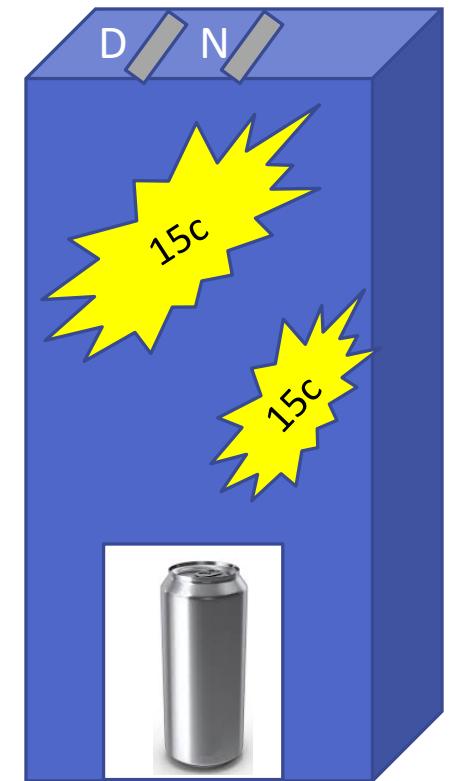
Present state	Input x		Outputs
	0	1	
W	Y	X	0
X	X	Y	1
Y	X	W	0

Time:	0	1	2	3	4	5
Present state:	W	Y	W	X	X	Y
Input:	0	1	1	0	1	0
Output:	0	0	0	1	1	0
Next State:	Y	W	X	X	Y	X



Vending Machine

- The computer department need a new soda machine
- We decided to ask computer students to design a controller for the new vending machine



Vending Machine: Characteristics

- Characteristics

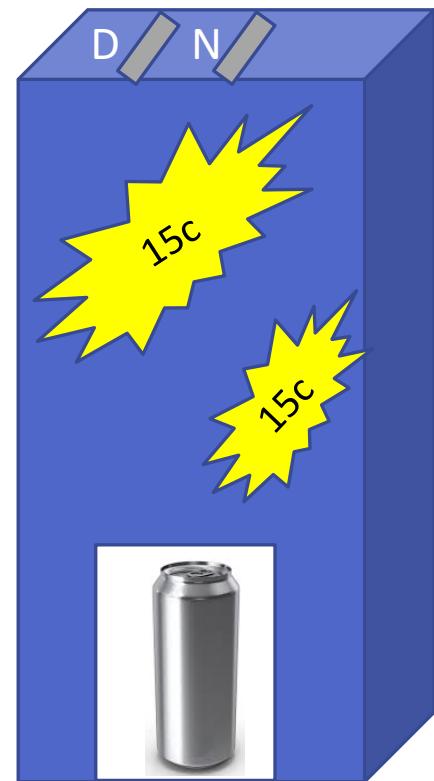
- All selections cost **15** Rial
- Machine does **not return changes!**

- Input

- **D**: dim inserted (10 Rial!)
- **F**: five inserted (5 Rial!)

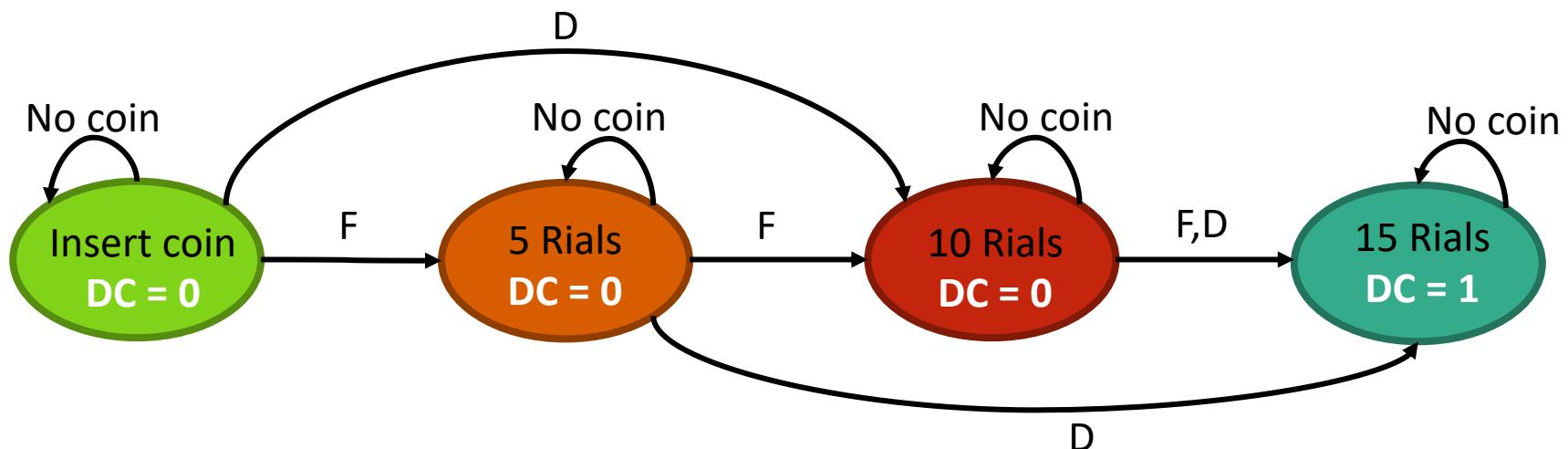
- Output

- **DC**: dispense can



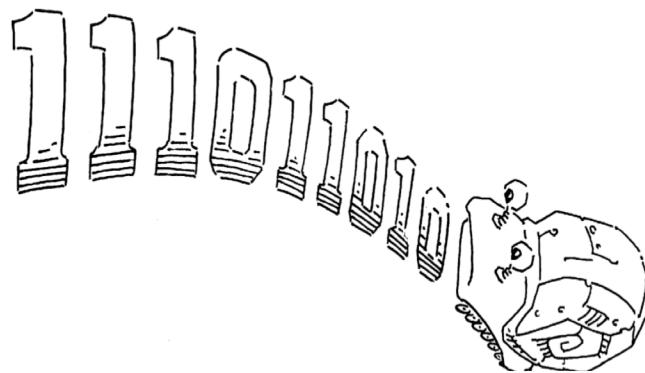
Vending Machine: FSM

- **Insert coin**
 - Default state
 - No money has been inserted
- **5 Rials**
 - A 5 Rial coin has been inserted
- **10 Rials**
 - A 10 Rial coin has been inserted
- **15 Rials**
 - Total money has been reached to 15 Rial
 - Done!



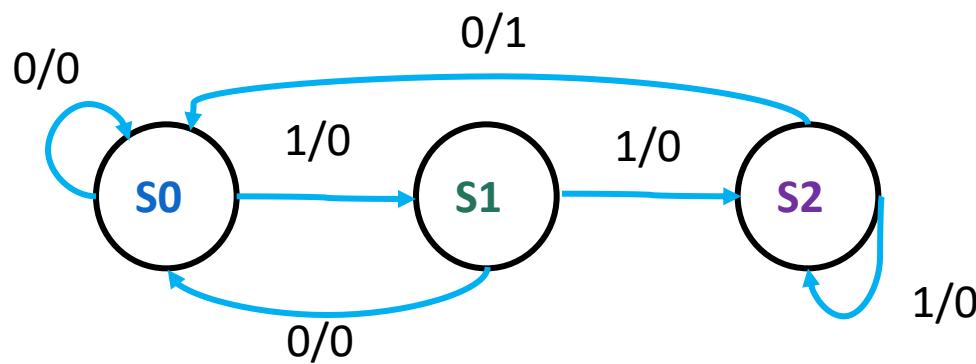
Pattern Recognition

- Recognize a specific bit pattern (110) in a bitstream
 - Mealy
 - Moore



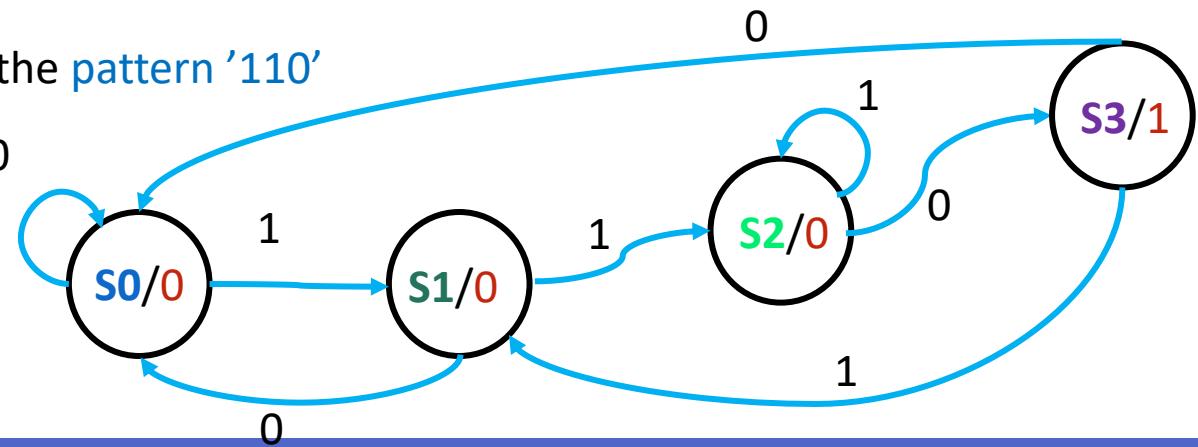
Pattern Recognition (110): Mealy

- State **S0**
 - We have **not** recognized **any** useful pattern
- State **S1**
 - We have recognized the **pattern '1'**
- State **S2:**
 - We have recognized the **pattern '11'**
 - **Output:** recognizing an input bit '0' in state S2



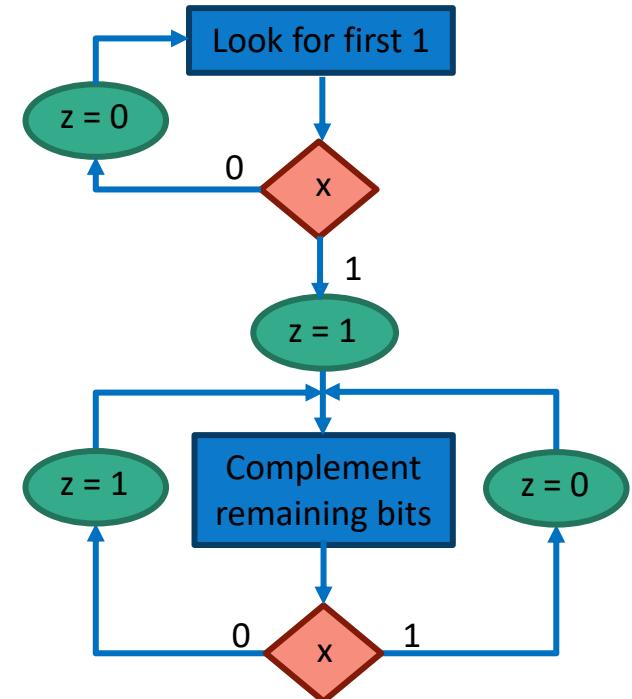
Pattern Recognition (110): Moore

- State **S0**
 - We have **not** recognized **any** useful pattern
- State **S1**
 - We have recognized the **pattern '1'**
- State **S2**:
 - We have recognized the **pattern '11'**
- State **S3**:
 - We have recognized the **pattern '110'**
 - **Output becomes 1**



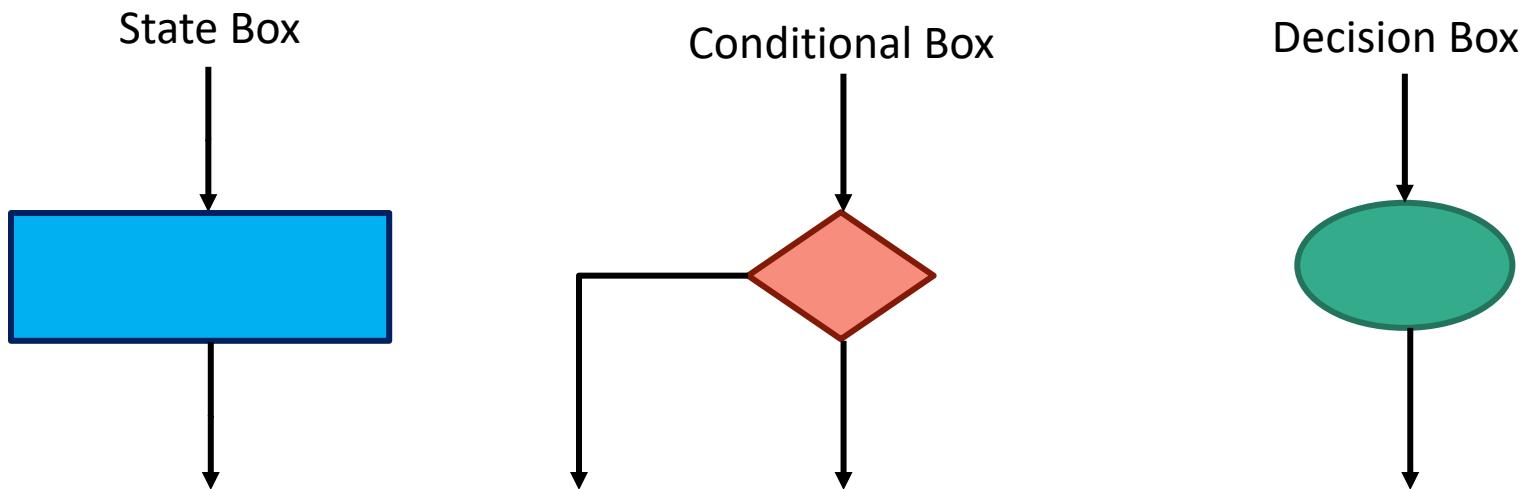
Algorithmic State Machine (ASM)

- A **systematic** way to **design complex** digital systems
 - **Complex** digital systems
 - Large number of **inputs** and **outputs**
- Describes the **sequence of events**
- Describes timing relationship between the states
- Behavioral model
- Basic Idea
 - Flowchart



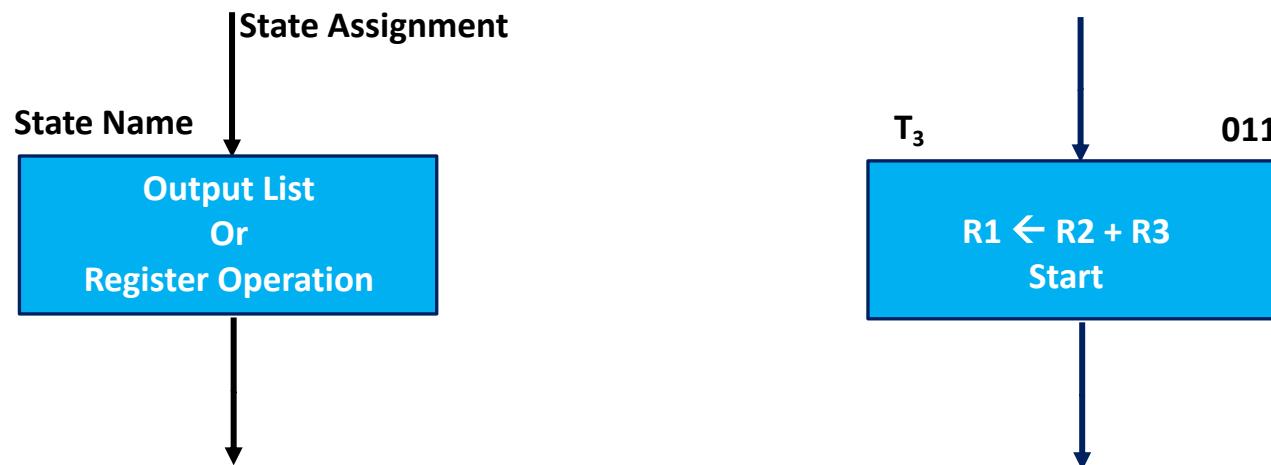
ASM chart Elements

- ASM chart elements
 - State box
 - Decision box
 - Conditional box



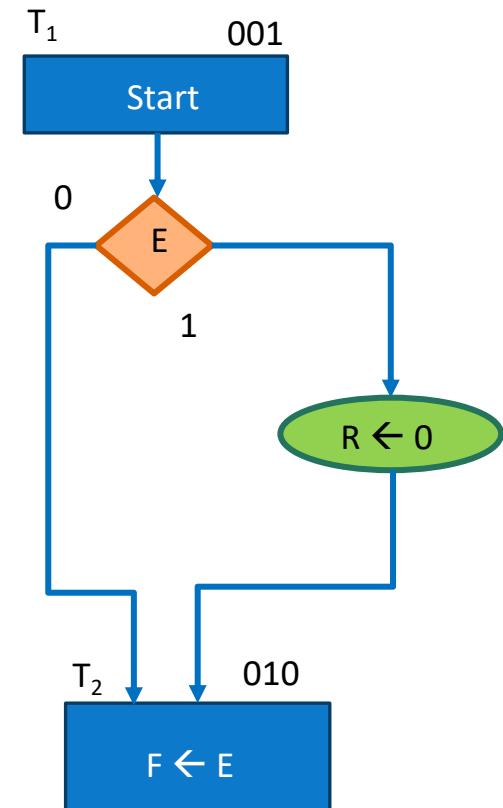
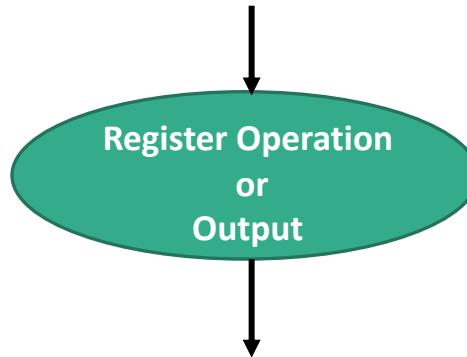
State Box

- Represents the state of the system
- Register Transfers
- Takes one cycle to be executed
 - $R1 \leftarrow R2 + R3$



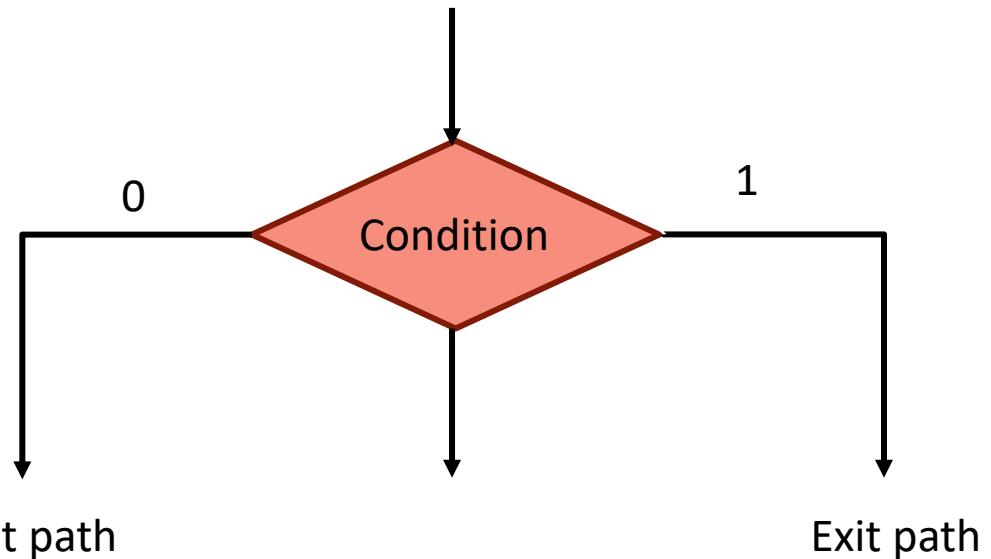
Conditional Box

- Register Transfers
- Contain conditional output list
 - Depends on both the **state** of the system and the **inputs**
 - A.k.a., **mealy** output
 - A condition output must follow a decision box
- Takes one cycle to be executed
 - $R1 \leftarrow R2 + R3$



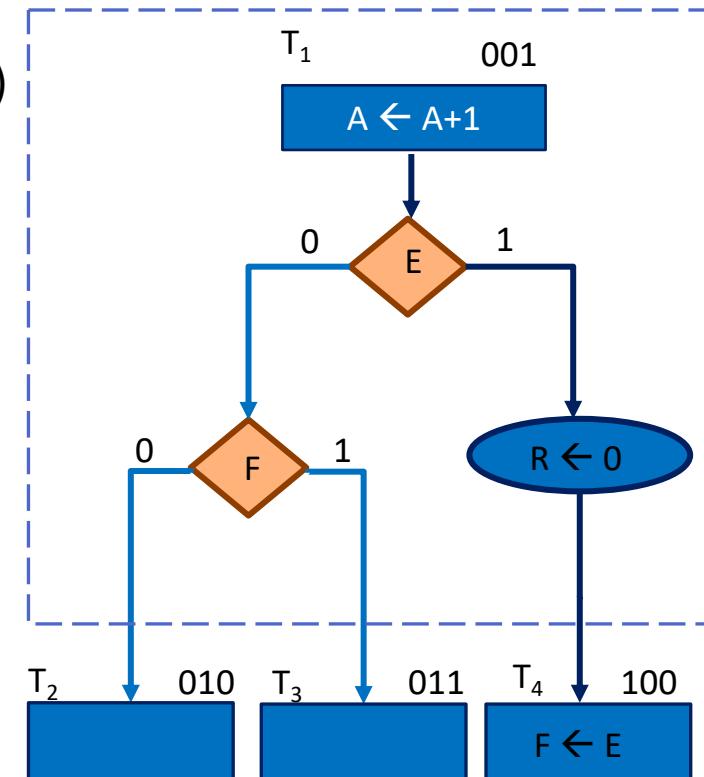
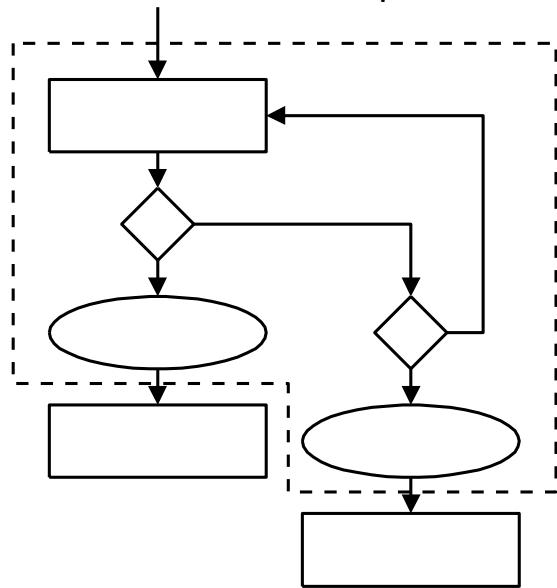
Decision Box

- Binary expressions
 - $\sim R1[7]$



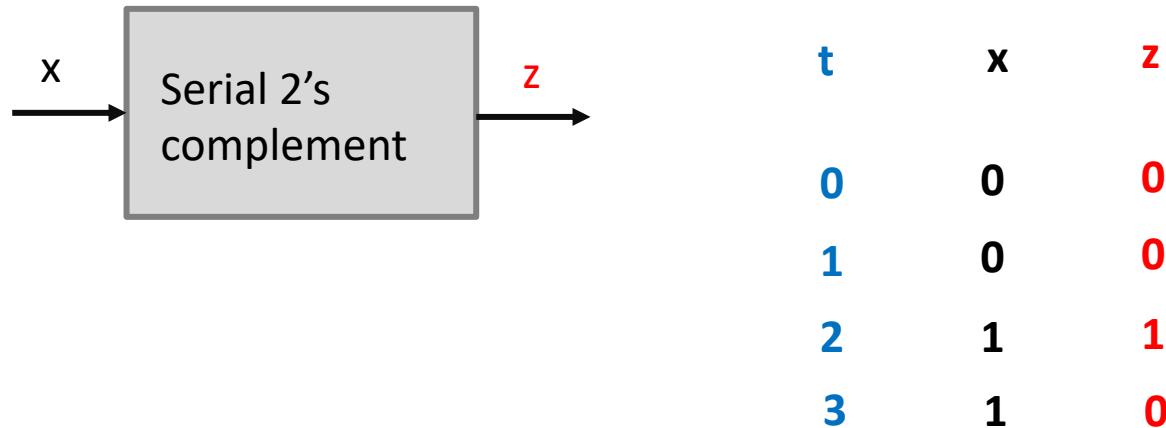
ASM Block

- Represents **what happens** in the system during **one clock cycle**
- Includes
 - Only one state box**
 - All other boxes (**decision** and **conditional** boxes)
 - Connected to the exit path of the **state box**

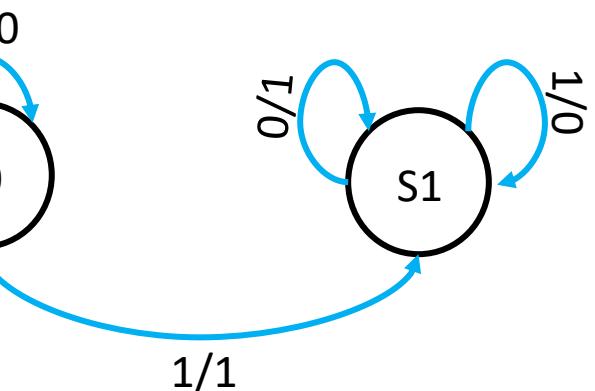
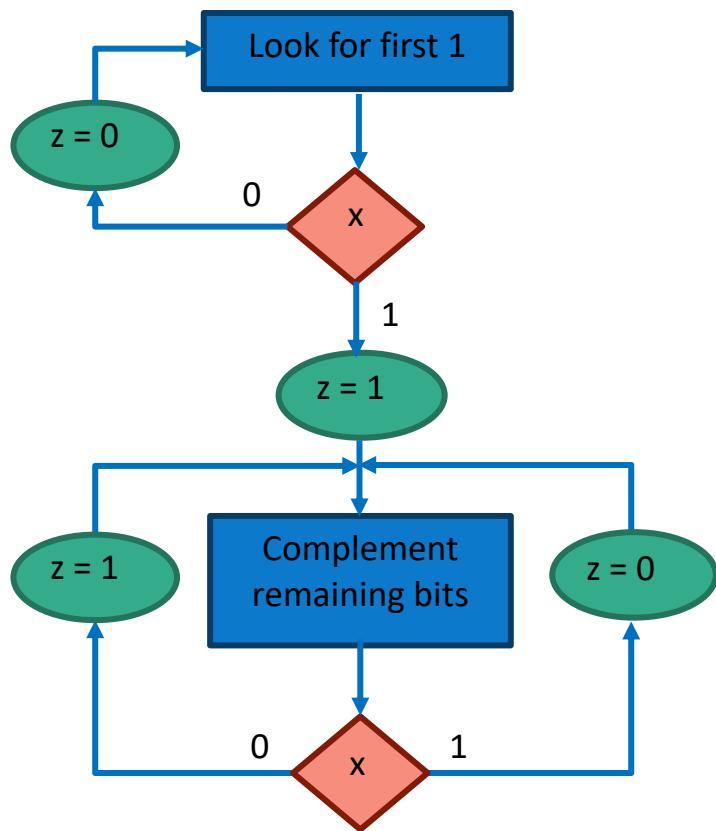


Sample Design 1

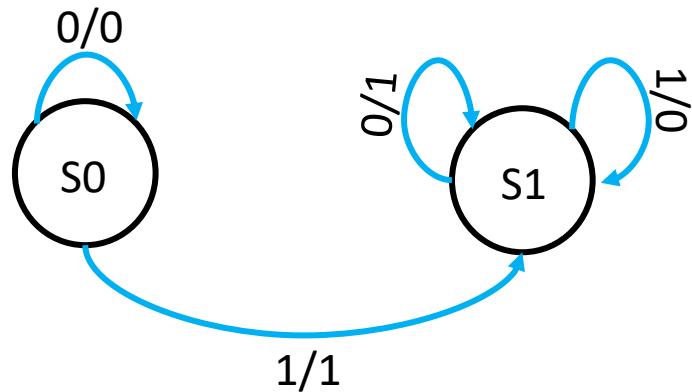
- Design a control unit for a serial 2's complement



Sample Design 1: Modeling



Sample Design 1: Excitation Table

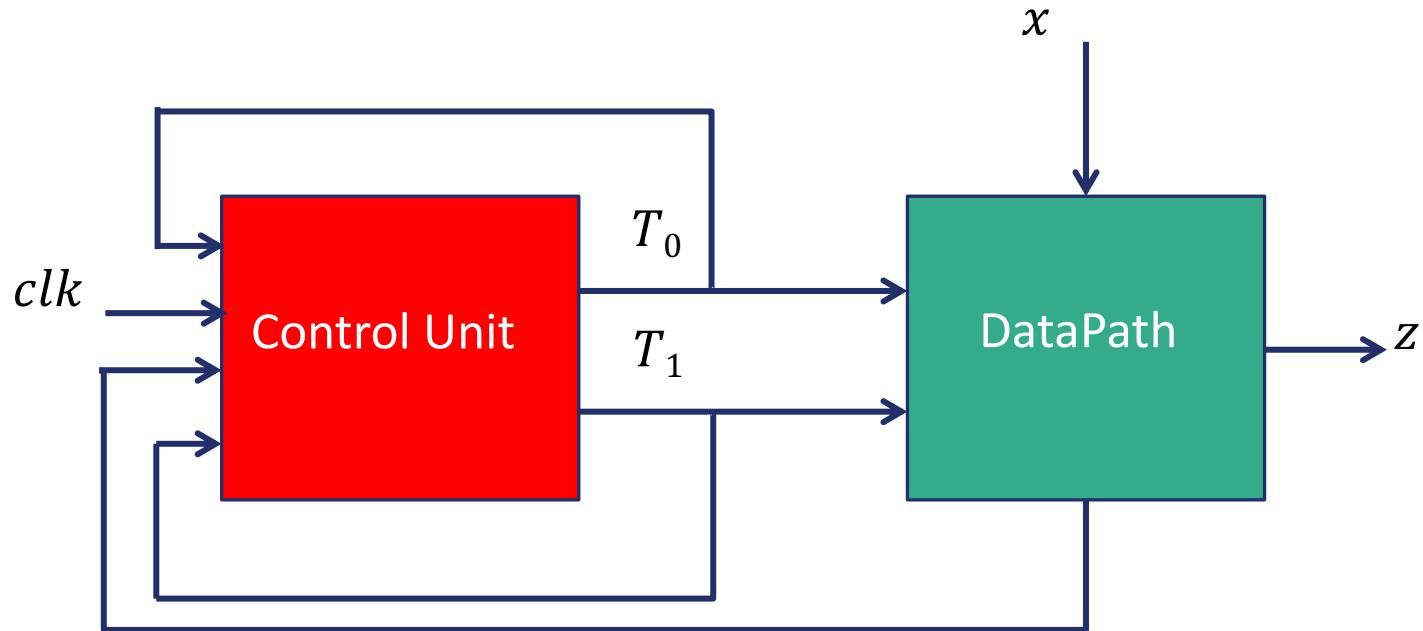


State	Input (x)	Next State	Output (z)	D
A(0)	0	A	0	0
A	1	B	1	1
B(1)	0	B	1	1
B	1	B	0	1

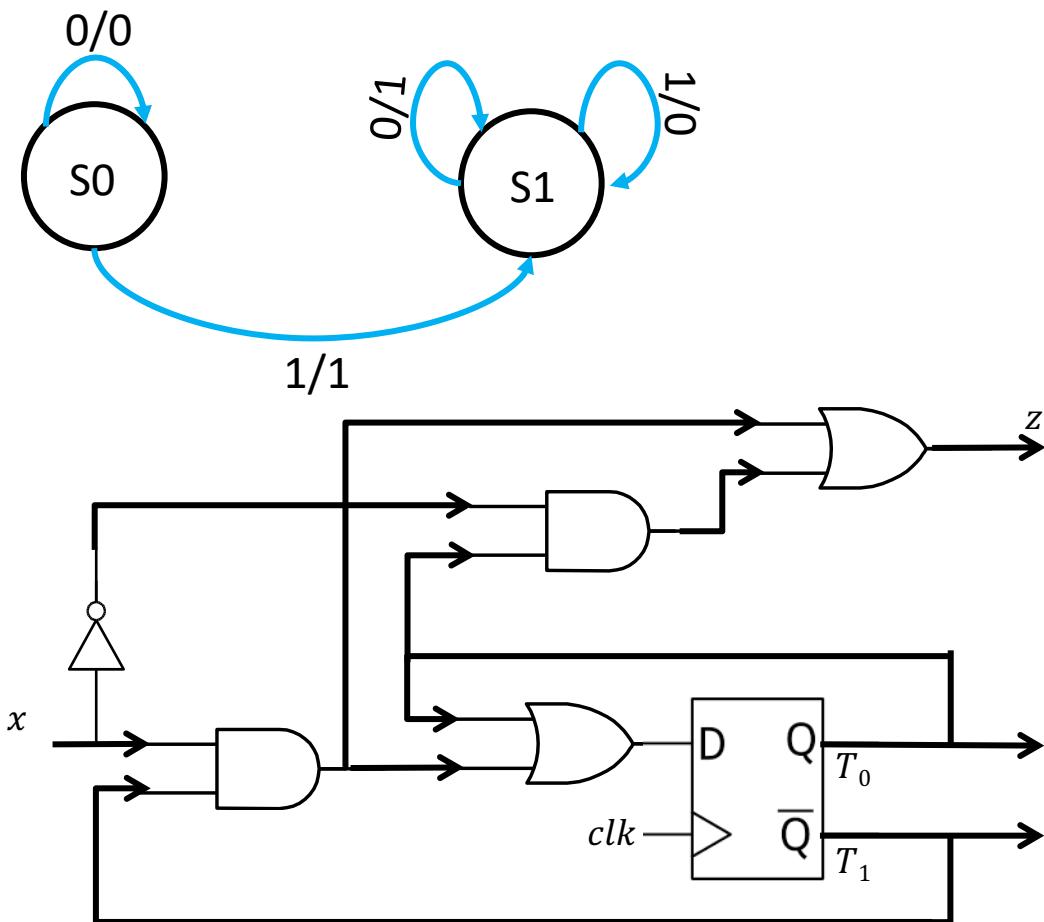
$$z = Ax + Bx'$$

$$D = Ax + B$$

Sample Design 1: Block Diagram



Sample Design 1: Design



State	Input (x)	Next State	Output (z)	D
A(0)	0	A	0	0
A	1	B	1	1
B(1)	0	B	1	1
B	1	B	0	1

$$z = Ax + Bx'$$

$$D = Ax + B$$

Moore Vs. Mealy

- Mealy
 - Has richer **description**
 - Usually requires **smaller** number of **states**
 - **Smaller** circuit **area**
 - Computes outputs as soon as inputs change
 - **Not synchronized** with the clock
 - Responds one clock cycle **sooner** than equivalent Moore FSM
 - May have **glitch**

- Moore
 - **Synchronize** with clock
 - **No glitch**, while in Mealy, you may have glitch

Thank You

